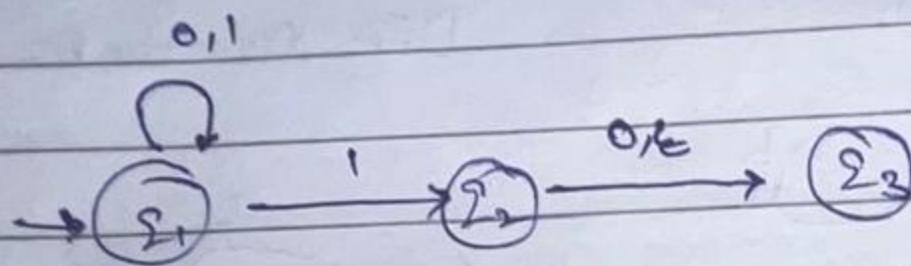


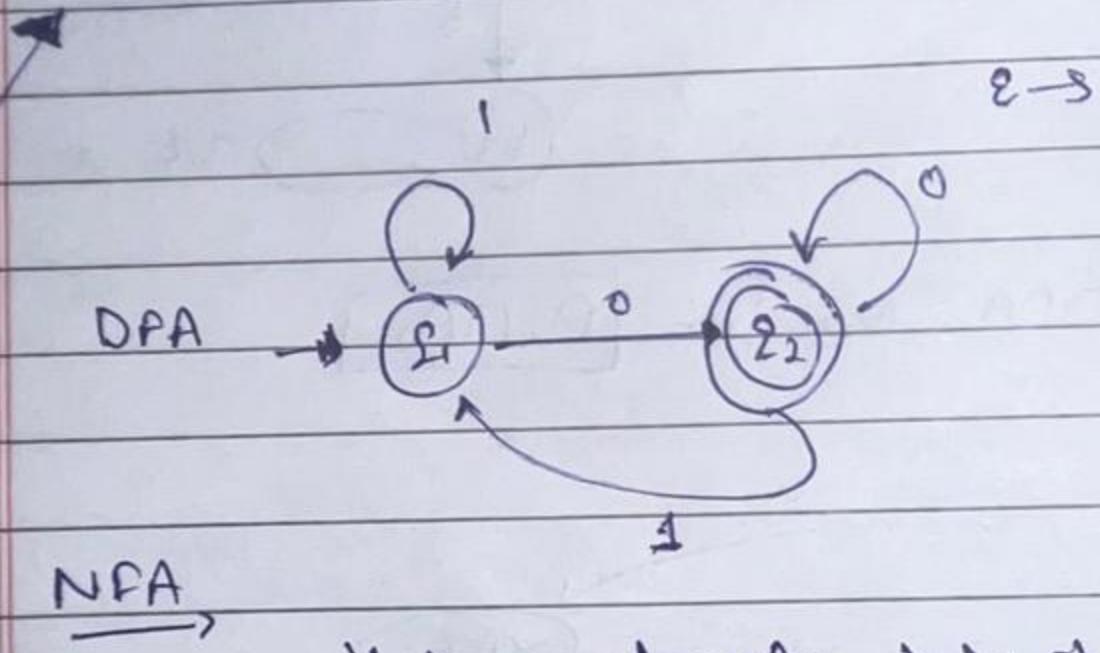
## Lecture - 4

Deterministic State is always

what is Diff b/w DFA & NDF



NFA



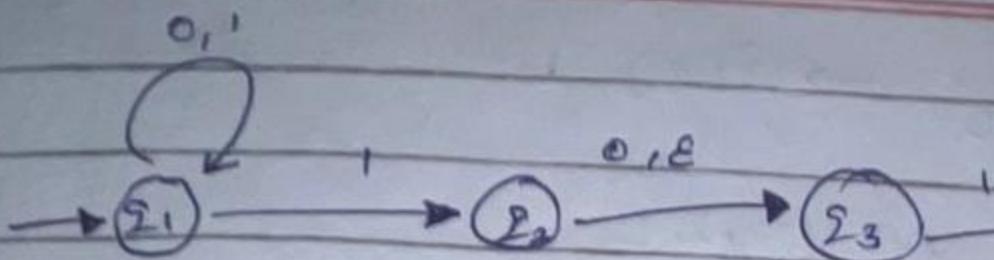
NFA

if the automaton is in st

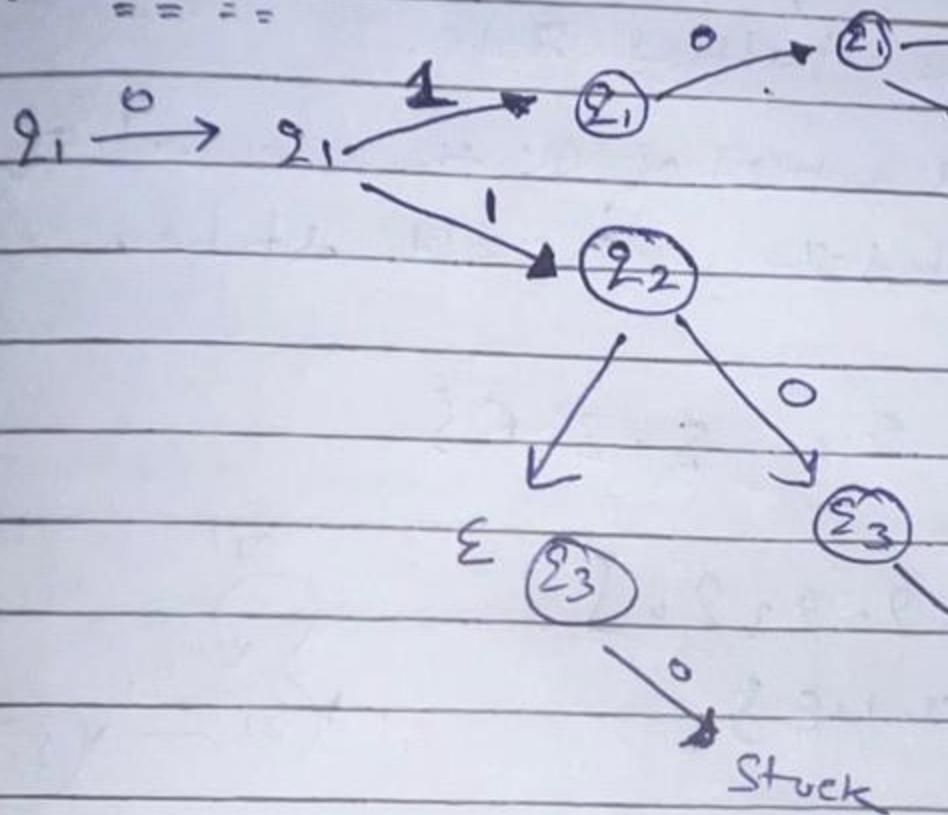
it moves into s. or other

Ex2

input String is 010110



$0 \rightarrow \underset{==}{0} \underset{==}{1} \underset{==}{0} \underset{==}{1} \underset{==}{1} \underset{==}{0}$



Computation can accept that  
input string we have 7 C  
but 2 of them are accepted

NFA:  $N = \{Q, \Sigma_\varepsilon, \delta, q_1, F\}$

$\delta: Q \times \Sigma_\varepsilon \rightarrow P(Q)$  is

$Q \rightarrow$  finite set of stat

$\Sigma_\varepsilon \rightarrow$  finite set of alphabet

$q_1 \rightarrow$  is the initial stat

$F \rightarrow$  is a subset of  $Q$ ; +  
called the acc

$N = \{Q, \Sigma_\varepsilon, \delta, q_1, F\}$

$Q = \{q_1, q_2, q_3, q_4\}$

$\Sigma_\varepsilon = \{0, 1, \varepsilon\}$

$q_1 \rightarrow \Sigma_1$

$F \rightarrow \{\varepsilon_4\}$

# COSC 3106 ) Theory of

## Lec 5# Equivalence of

By

Convert NFA to DFA.

- The key idea to convert  
keep track of all possible  
in simultaneously.

DFA  $M: \{ Q, \Sigma, \delta, q_0 \}$

NFA  $N: \{ Q, \Sigma, \delta, q_0 \}$

States : the set of states

$Q, \delta$

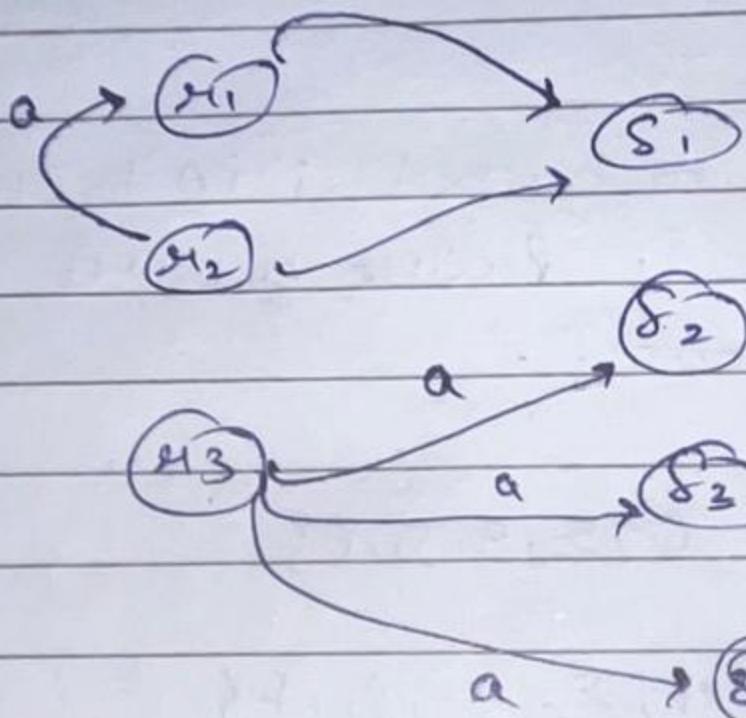
$\rightarrow$  for  $\epsilon \in R$  of set of co

■ EASY CASE :-

No  $\epsilon$  - Transition in

$\circ \underline{\epsilon x}$

$$R = \{ \circ_1 \circ_2 \circ_3$$



$$\{ s_1, s_2, s_3 \} \xrightarrow{a} \{$$

■ General Case