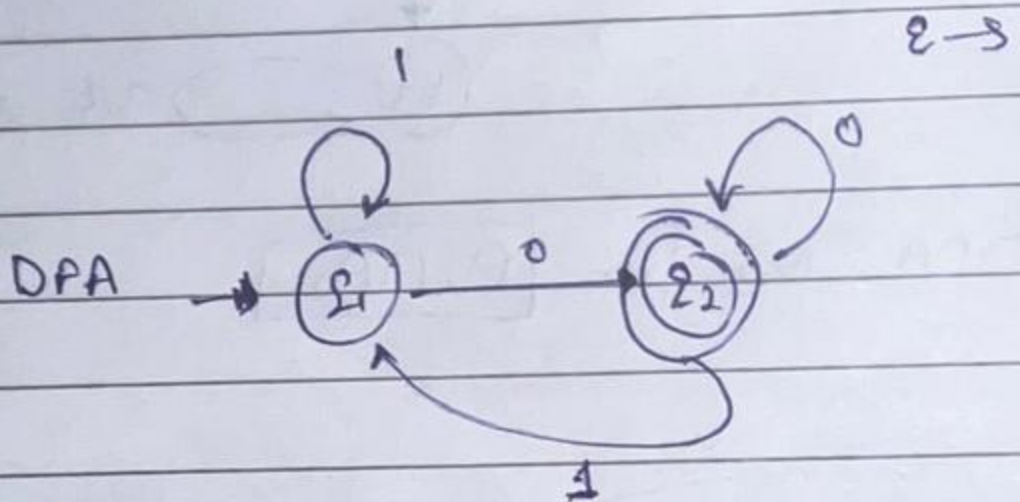
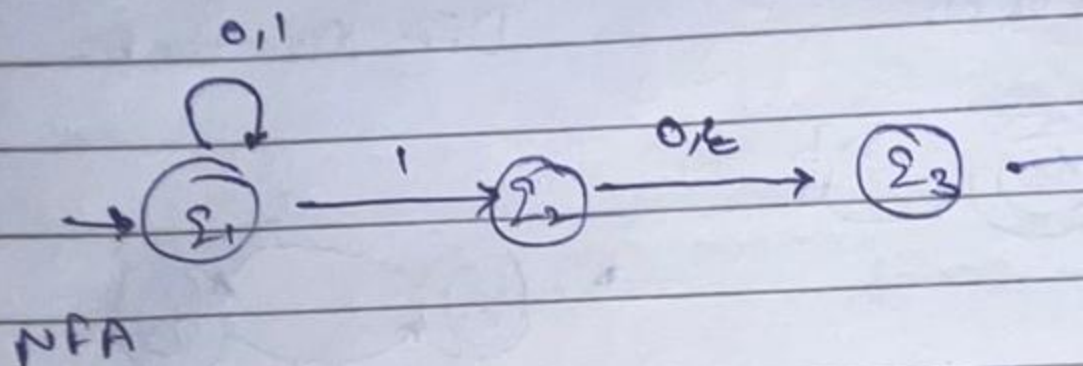


# Lecture -4

Deterministic State is always unique

what is Diff b/w DFA & NFA

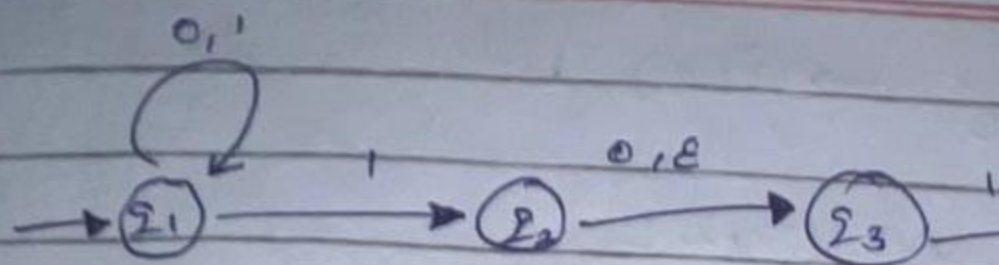


NFA

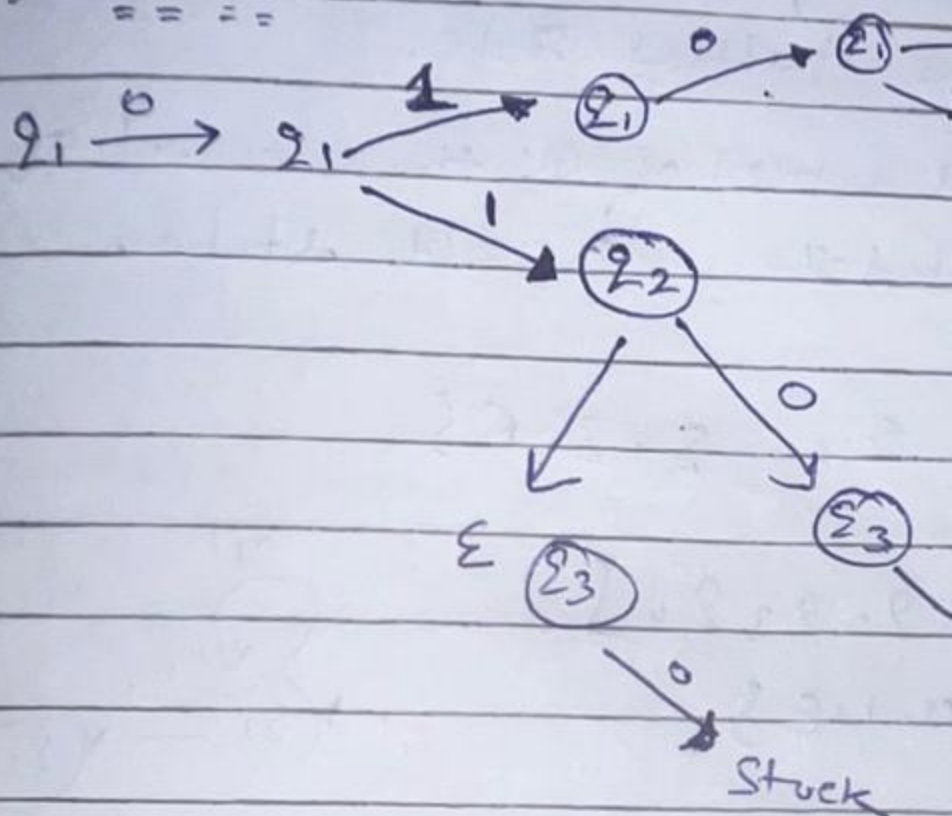
if the automaton is in state  
it may go to q1 or other

ex2

input string is 010110



0 → 010110  
=====



Computation can accept that  
input string we have 7 C  
but ② of them are accepted

NFA :  $N = \{Q, \Sigma_\epsilon, \delta, q, F\}$

$\delta: Q \times \Sigma_\epsilon \rightarrow P(Q)$  is

$Q \rightarrow$  finite set of states

$\Sigma_\epsilon \rightarrow$  finite set of alphabet

$q \rightarrow$  is the initial state

$F \rightarrow$  is a subset of  $Q$ ;  $F$   
called the accepting

$N = \{Q, \Sigma_\epsilon, \delta, q, F\}$

$Q = \{q_1, q_2, q_3, q_4\}$

$\Sigma_\epsilon = \{0, 1, \epsilon\}$

$q_1 \rightarrow \epsilon_1$

$F \rightarrow \{q_4\}$

	0	1	$\epsilon$
$q_1$			
$q_2$			
$q_3$			
$q_4$			



(COSC 3106) Theory of

Lec 5# Equivalence of

By

Convert NFA to DFA.

- The key idea to convert is to keep track of all possible states simultaneously.

DFA  $M: \{Q, \Sigma, \delta, q_0, F\}$

NFA  $N: \{Q, \Sigma, \delta, q_0, F\}$

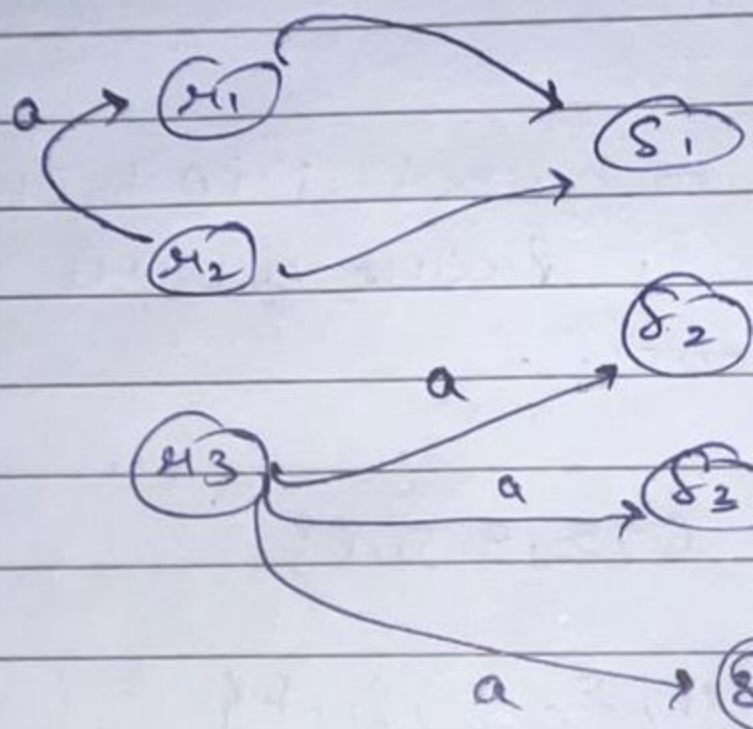
States: the set of states

$\rightarrow$  for  $\in R$  of set of  $Q$

## □ EASY CASE :-

No  $\epsilon$ -Transition in

ex  $R = \{r_1, r_2, r_3\}$



$\{r_1, r_2, r_3\} \xrightarrow{a} \{r_1, r_2, r_3\}$

## □ General Case