

A BRIEF INTRODUCTION TO DEC-POMDP

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ABSTRACT

This article introduces Dec-POMDP, the computation, and planning methods.¹

1 DEC-POMDP

Definition 1. A Dec-POMDP is a tuple $\langle \mathbb{I}, \mathcal{S}, \{\mathbb{A}_i\}, T, R, \{\mathbb{O}_i\}, O, \mathcal{H}, \gamma \rangle$:

- \mathbb{I} is a finite sets of agents, $|\mathbb{I}| = n$;
- \mathcal{S} is a set of states with designated initial state distribution b^0 ;
- \mathbb{A}_i is a set of actions for agent i with $\mathbb{A} \doteq \times_i \mathbb{A}_i$ the set of joint actions;
- T is the state transition probability function, $T: \mathcal{S} \times \mathbb{A} \times \mathcal{S} \rightarrow [0, 1]$, that specifies the probability of transitioning from state $s \in \mathcal{S}$ to $s' \in \mathcal{S}$ when the actions $\mathbf{a} \in \mathbb{A}$ are taken by agents (i.e., $T(s, \mathbf{a}, s') = P(s'|s, \mathbf{a})$);
- R is the joint reward function, where $R: \mathcal{S} \times \mathbb{A} \rightarrow \mathbb{R}$;
- \mathbb{O}_i is a set of observations for each agent i , with $\mathbb{O} \doteq \times_i \mathbb{O}_i$ the set of joint observations;
- O is an observation probability function, $O: \mathbb{A} \times \mathcal{S} \times \mathbb{O}$, that specifies the probability of seeing observation $\mathbf{o}' \in \mathbb{O}$ given the actions $\mathbf{a} \in \mathbb{A}$ are taken and state $s' \in \mathcal{S}$ is observed (i.e., $O(\mathbf{a}, s', \mathbf{o}') = P(\mathbf{o}'|\mathbf{a}, s')$);
- \mathcal{H} is the horizon (the number of steps until termination);
- γ is the discount factor for the return.

A solution to a Dec-POMDP is a joint policy $\pi: \mathbb{H}_i \rightarrow \mathbb{A}_i, \forall i \in \mathbb{I}$ over joint observation-action history $\mathbf{h} = \{\mathbf{a}^0, \mathbf{o}^1, \dots, \mathbf{o}^{\mathcal{H}-1}\}$, an optimal solution maximizes the expected return,

$$\pi^* = \operatorname{argmax}_{\pi} \mathbb{E} \left[\sum_{t=0}^{\mathcal{H}-1} R(\mathbf{h}, \pi(\mathbf{h})) \middle| b^0 \right].$$

In Dec-POMDP, the Bellman recursive formulation of the history V-function is,

$$\begin{aligned} V^\pi(\mathbf{h}) &= \sum_s P(s|b^0, \mathbf{h}) \left[R(s, \pi(\mathbf{h})) + \gamma \sum_{s'} P(s'|s, \pi(\mathbf{h})) \sum_{\mathbf{o}'} P(\mathbf{o}'|\pi(\mathbf{h}), s') V^\pi(\mathbf{h}') \right] \\ &\equiv R(\mathbf{h}, \pi) + \gamma \sum_{\mathbf{o}'} P(\mathbf{o}'|\mathbf{h}, \pi) V^\pi(\mathbf{h}'), \end{aligned} \quad (1)$$

the Bellman recursive formulation of the **history-policy** Q-function is,²

$$\begin{aligned} Q^\pi(\mathbf{h}, \pi) &= \sum_s P(s|b^0, \mathbf{h}) \left[R(s, \pi(\mathbf{h})) + \gamma \sum_{s'} P(s'|s, \pi(\mathbf{h})) \sum_{\mathbf{o}'} P(\mathbf{o}'|\pi(\mathbf{h}), s') Q^\pi(\mathbf{h}', \pi(\mathbf{h}')) \right] \\ &\equiv R(\mathbf{h}, \pi) + \gamma \sum_{\mathbf{o}'} P(\mathbf{o}'|\mathbf{h}, \pi) Q^\pi(\mathbf{h}', \pi). \end{aligned} \quad (2)$$

¹This introduction is **heavily** based on Amato et al. (2013).

²The definition of the value function is pretty flexible, which can be based on various aspects, such as the value of a state, a belief state, an observation, a state history, an observation history, an action (single-step policy), a policy (action history), observation-action history, and their combinations.

2 SUBCLASSES

Centralized Control MMDP is a fully observable version of Dec-POMDP, but it does not specify decentralized control. Dec-MDP assumes that the joint observations uniquely determine the state, while agents still act with local observations. Similarly, MPOMDP does not specify whether the control is decentralized, which could have a centralized policy $\mathbb{H} \rightarrow \mathbb{A}$.

Independent Variables A decentralized control model might be factorized with independent local variables, e.g., transition-independence (TI) $T(s, \mathbf{a}, s') = \prod_{i=1}^n T(s_i, a_i, s'_i)$ and reward-independence (RI) $R(s, \boldsymbol{\pi}) = f_{\text{mono}}(\langle R(s_i, \pi_i) \rangle_{i=1}^n)$. Network-distributed POMDP (ND-POMDP) represents the factored one with TI and block-RI, i.e., $R(s, \boldsymbol{\pi}) = f_{\text{mono}}(\langle R(s_{i, \mathcal{N}(i)}, \pi_{i, \mathcal{N}(i)}) \rangle_{i=1}^n)$, where $\mathcal{N}(i)$ are the neighbors of i .

Complexity The worst-case complexity of finite-horizon problems is: (by Amato et al. (2013))

Model	Complexity
MDP	P-complete
MMDP (Cen-MMDP)	P-complete
Dec-MDP	NEXP-complete
Dec-MDP with TI no RI	NP-complete
Dec-MDP with RI no TI	NEXP-complete
Dec-MDP with TI and RI	P-complete
POMDP	PSPACE-complete
MPOMDP (Cen-MPOMDP)	PSPACE-complete
Dec-POMDP	NEXP-complete
ND-POMDP	NEXP-complete

Theorem 1. An MDP is P-complete in finite and infinite horizons Papadimitriou & Tsitsiklis (1987).

Theorem 2. A finite POMDP is PSPACE-complete Papadimitriou & Tsitsiklis (1987).

Theorem 3. The complexity of an infinite POMDP is undecidable Madani et al. (1999), leading to the undecidability of the infinite Dec-POMDP complexity.

Theorem 4. A finite Dec-POMDP _{$n \geq 2$} is NEXP-complete, and a finite Dec-MDP _{$n \geq 3$} is also NEXP-complete Bernstein et al. (2002).

Fact 1. A Dec-MDP with TI and RI can be solved independently, resulting in P-complete.

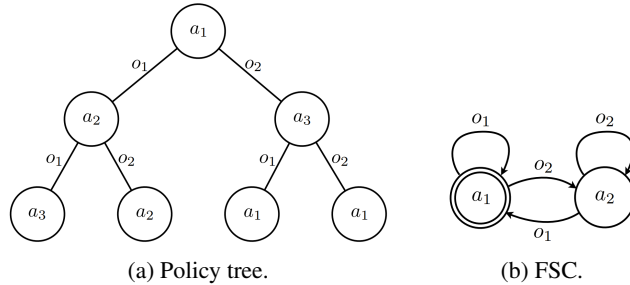
Theorem 5. A Dec-MDP with TI and joint reward is NP-complete, a Dec-MDP with RI but no TI is NEXP-complete Becker et al. (2004)

Fact 2. An ND-POMDP has the same worst-case complexity as a Dec-POMDP Nair et al. (2005).

3 PLANNING METHODS

3.1 POLICY STRUCTURE

Calculating a shared belief state in Dec-POMDP is hard, because the policy can not be recovered from the value function. The policies are normally maintained in a policy tree or FSC. Policies can be extracted by starting at the root (or initial node) and continuing to the subtree (or next node) based on observations, and can be evaluated by summing the rewards weighted by transition probability.



3.2 OPTIMAL APPROACHES

Bottom-up DP uses joint belief to find optimal solutions with policy pruning Hansen et al. (2004).

$$V^{t+1}(b^t) = \max_{a \in \mathcal{A}} \left\{ \sum_{s \in \mathcal{S}} b^t(s) \left[R(s, a) + \sum_{o \in \mathcal{O}} \mathcal{P}(o \mid s, a) V^t(b^{t+1}) \right] \right\}. \quad (3)$$

Algorithm 1 DP Backup with Policy Pruning in Dec-POMDP

Input: Depth- t policy trees Q_i^t and value vectors V_i^t for each i

- 1: Perform exhaustive backups to get Q_i^{t+1} , and compute V_i^{t+1} accordingly for each i
- 2: **repeat**
- 3: Find a policy tree $q_j \in Q_i^{t+1}$ that satisfies $\exists v_k \in \{V_i^{t+1} \setminus v_j\}, b^{t+1}v_k \geq b^{t+1}v_j, \forall b^{t+1}$
- 4: $Q_i^{t+1} \leftarrow \{Q_i^{t+1} \setminus q_j\}$, and $V_i^{t+1} \leftarrow \{V_i^{t+1} \setminus v_j\}$ accordingly
- 5: **until** no more pruning for all i

Output: Depth- $t + 1$ policy trees Q_i^{t+1} and value vectors V_i^{t+1} for each i

Top-down The policy tree can also be built using heuristic search like MAA* Szer et al. (2005).

Algorithm 2 MAA* Heuristic Search

Initialize: Joint policy tree root $\Pi \leftarrow \times_i \mathbb{A}_i$

Input: Depth- t joint policy tree Π

- 1: Select $a^* = \operatorname{argmax}_{a \in \Pi} F^{\mathcal{H}}(b^0, a)$ (heuristic and value)
- 2: Expand a^* to a^{\oplus} , and $\Pi \leftarrow \{\Pi \cup a^{\oplus}\}$
- 3: **for** Each $a \in \Pi$ **do**
- 4: **if** $F^{\mathcal{H}}(s^0, a^{\oplus}) \geq F^{\mathcal{H}}(s^0, a)$ **then**
- 5: $\Pi \leftarrow \{\Pi \setminus a\}$
- 6: **end if**
- 7: **end for**
- 8: **if** a^* is fully expanded **then**
- 9: $\Pi \leftarrow \{\Pi \setminus a^*\}$
- 10: **end if**

Input: Updated joint policy set Π .

3.3 APPROXIMATION APPROACHES

The algorithms below improve scalability to larger problems over optimal methods, but do not possess any bounds on solution quality.

MBDP Memory-bounded dynamic programming (MBDP) techniques mitigate the scalability problem of DP (which generates and evaluates all joint policy trees before pruning) by keeping a fixed number of policy trees for each agent at each step Seuken & Zilberstein (2007a). Several approaches have improved on MBDP by limiting Seuken & Zilberstein (2007b) or compressing Carlin & Zilberstein (2008) observations, replacing exhaustive backup with branch-and-bound search in the space of joint policy trees Dibangoye et al. (2009) as well as constraint optimization Kumar & Zilberstein (2010) and linear programming Wu et al. (2010) to increase the efficiency of selecting the best trees at each step.

JESP The joint equilibrium search for policies (JESP) Nair et al. (2003) uses alternating best response. Initial policies are generated for all agents, and then all but one is held fixed. The remaining agent can then calculate the best response (local optimum) to the fixed policies. The policy of this agent becomes fixed and the next agent calculates the best response. These best-response calculations to fixed other agent policies continue until no agent changes its policy.

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A COMPLEXITY CLASSES

Assuming c and k are constants, \mathcal{C} is a complexity class, the table shows complexity terminologies.

P	the set of problems solvable in polynomial time, e.g., $O(n^k)$
NP	the set of problems solvable nondeterministically in polynomial time
EXP	the set of problems solvable in exponential time, e.g., $O(c^{n^k})$
NEXP	the set of problems solvable nondeterministically in exponential time
PSPACE	the set of problems solvable in polynomial space (P and $NP \subset PSPACE$), e.g., $O(c^{n^k})$
\mathcal{C} -hard	a problem that all problems in \mathcal{C} are reducible to within polynomial time
\mathcal{C} -complete	a problem that is contained in \mathcal{C} and \mathcal{C} -hard