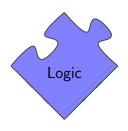


## **COMP9020**

Foundations of Computer Science

Lecture 12: Boolean Logic

# Topic 3: Logic



		[LLM]	[RW]	[Rosen]
Week 8	Boolean Logic	Ch. 3	Ch. 2, 10	Ch. 12
Week 8	Propositional Logic	Ch. 3	Ch. 2	Ch. 1

## What is logic?

### Logic is about formalizing reasoning and defining truth

- Adding rigour
- Removing ambiguity
- Mechanizing the process of reasoning

## Loose history of logic

- (Ancient times): Logic exlusive to philosophy
- Mid-19th Century: Logical foundations of Mathematics (Boole, Jevons, Schröder, etc)
- 1910: Russell and Whitehead's Principia Mathematica
- 1928: Hilbert proposes Entscheidungsproblem
- 1931: Gödel's Incompleteness Theorem
- 1935: Church's Lambda calculus
- 1936: Turing's Machine-based approach
- 1930s: Shannon develops Circuit logic
- 1960s: Formal verification; Relational databases

## Logic in Computer Science

 ${\sf Computation} \quad = \quad {\sf Calculation} \, + \, {\sf Symbolic \ manipulation}$ 

## Logic in Computer Science

Computation = Calculation + Symbolic manipulation

Logic as 2-valued computation (Boolean logic):

- Circuit design
- Code optimization
- Boolean algebra
- Nand game

## Logic in Computer Science

Computation = Calculation + Symbolic manipulation

Logic as symbolic reasoning (Propositional logic, and beyond):

- Formal verification
- Proof assistance
- Knowledge Representation and Reasoning
- Automated reasoning
- Databases

### Outline

Boolean Logic

**Boolean Functions** 

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

Boolean Algebras

### Outline

### Boolean Logic

Boolean Functions

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

Boolean Algebras

## Boolean logic

Boolean logic is about performing calculations in a "simple" mathematical structure.

- complex calculations can be built entirely from these simple ones
- can help identify simplifications that improve performance at the circuit level
- can help identify simplifications that improve presentation at the programming level

## The Boolean Algebra B

#### **Definition**

The (two-element) **Boolean algebra** is defined to be the set  $\mathbb{B}=\{0,1\}$ , together with the functions  $!:\mathbb{B}\to\mathbb{B}$ , &&:  $\mathbb{B}^2\to\mathbb{B}$ , and  $||:\mathbb{B}^2\to\mathbb{B}$ , defined as follows:

$$|x| = (1-x)$$
  $x \&\& y = \min\{x, y\}$   $x \parallel y = \max\{x, y\}$ 

## The Boolean Algebra B – Alternative definition

#### **Definition**

The (two-element) **Boolean algebra** is defined to be the set  $\mathbb{B} = \{ \text{false}, \text{true} \}$ , together with the functions  $! : \mathbb{B} \to \mathbb{B}$ , &&:  $\mathbb{B}^2 \to \mathbb{B}$ , and  $\| : \mathbb{B}^2 \to \mathbb{B}$ , defined as follows:

		X	У	x && y	X	У	$x \parallel y$
X	!x	false	false	false	false	false	false
false	true	false	true	false	false	true	true
true	false	true	false	false	true	false	true
	'	true	true	true	true	true	true
				ı			

### Alternative notation

Commonly, the following alternative notation is used:

For  $\mathbb{B}$ :  $\{F, T\}$ 

For !x:  $\overline{x}$ , x',  $\sim x$ ,  $\neg x$ 

For x && y:  $xy, x \land y$ 

For  $x \parallel y$ : x + y,  $x \vee y$ 

### Properties

We observe that !, &&, and || satisfy the following:

For all 
$$x,y,z\in\mathbb{B}$$
: 
$$x\parallel y=y\parallel x$$
 
$$x\&\&\ y=y\&\&\ x$$
 Associativity 
$$(x\parallel y)\parallel z=x\parallel (y\parallel z)$$
 
$$(x\&\&\ y)\&\&\ z=x\&\&\ (y\&\&\ z)$$
 Distribution 
$$x\parallel (y\&\&\ z)=(x\parallel y)\&\&\ (x\parallel z)$$
 
$$x\&\&\ (y\parallel z)=(x\&\&\ y)\parallel (x\&\&\ z)$$
 Identity 
$$x\parallel 0=x$$
 
$$x\&\&\ 1=x$$
 Complementation 
$$x\parallel (!x)=1$$
 
$$x\&\&\ (!x)=0$$

## Examples

### **Examples**

- Calculate x && x for all  $x \in \mathbb{B}$
- Calculate ((1 && 0) || ((!1) && (!0))

### Outline

Boolean Logic

#### **Boolean Functions**

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

Boolean Algebras

### **Boolean Functions**

#### **Definition**

An *n*-ary Boolean function is a map  $f : \mathbb{B}^n \to \mathbb{B}$ .

### Question

How many unary Boolean functions are there? How many binary functions? n-ary?

### Examples

#### **Examples**

- ! is a unary Boolean function
- &&, ∥ are binary Boolean functions
- f(x, y) = !(x && y) is a binary boolean function (NAND)
- AND $(x_0, x_1, ...) = (\cdots ((x_0 \&\& x_1) \&\& x_2) \cdots)$  is a (family) of Boolean functions
- $OR(x_0, x_1, ...) = (\cdots ((x_0 \parallel x_1) \parallel x_2) \cdots)$  is a (family) of Boolean functions

## Application: Adding two one-bit numbers

#### How can we implement:

$$\mathsf{add}: \mathbb{B}^2 \to \mathbb{B}^2$$

defined as

Χ	У	add(x, y)
0	0	00
0	1	01
1	0	01
1	1	10

Use two Boolean functions!

#### NB

Digital circuits are just sequences of Boolean functions.

### Outline

Boolean Logic

Boolean Functions

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

Boolean Algebras

## Conjunctive and Disjunctive normal form

#### **Definition**

- A literal is a unary Boolean function
- A minterm is a Boolean function of the form  $AND(I_1(x_1), I_2(x_2), \ldots, I_n(x_n))$  where the  $I_i$  are literals
- A maxterm is a Boolean function of the form  $OR(I_1(x_1), I_2(x_2), \dots, I_n(x_n))$  where the  $I_i$  are literals
- A CNF Boolean function is a function of the form  $AND(m_1, m_2, ...)$ , where the  $m_i$  are maxterms.
- A **DNF Boolean function** is a function of the form  $OR(m_1, m_2, ...)$ , where the  $m_i$  are minterms.

### Examples

#### **Examples**

- $f(x, y, z) = (x \&\& (!y) \&\& z) || (x \&\& (!y) \&\& (!z)) = x \overline{y} z + x \overline{y} \overline{z}$ : DNF, but not CNF
- $g(x, y, z) = (x \| (!y) \| z) \&\& (x \| (!y) \| (!z)) = (x + \overline{y} + z)(x + \overline{y} + \overline{z})$ : CNF function, but not DNF
- $h(x, y, z) = (x \&\& (!y) \&\& z) = x \overline{y} z$ : both CNF and DNF
- j(x, y, z) = x + y(z + x): Neither CNF nor DNF

#### NB

CNF: product of sums; DNF: sum of products

#### **Theorem**

Every Boolean function can be written as a function in DNF/CNF

Proof...

### Canonical DNF

Given an *n*-ary boolean function  $f: \mathbb{B}^n \to \mathbb{B}$  we construct an equivalent DNF boolean function as follows:

For each  $\mathbf{b} = (b_1, \dots, b_n) \in \mathbb{B}^n$  we define the minterm

$$m_{\mathbf{b}} = \text{And}(l_1(x_1), l_2(x_2), \dots, l_n(x_n))$$

where

$$I_i(x_i) = \begin{cases} x_i & \text{if } b_i = 1\\ !x_i & \text{if } b_i = 0 \end{cases}$$

We then define the DNF formula:

$$f_{\mathsf{DNF}} = \sum_{f(\mathbf{b})=1} m_{\mathbf{b}},$$

that is,  $f_{\text{DNF}}$  is the disjunction (or) over all minterms corresponding to elements  $\mathbf{b} \in \mathbb{B}$  where  $f(\mathbf{b}) = 1$ .

### Canonical DNF

#### **Theorem**

f and  $f_{DNF}$  are the same function.

### Exercise

#### **Exercises**

RW: 10.2.3 Find the canonical DNF form of each of the following expressions in variables x, y, z

- xy
- <u>Z</u>
- $xy + \overline{z}$
- f(x, y, z) = 1

### Outline

Boolean Logic

Boolean Functions

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

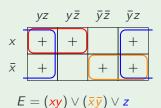
Boolean Algebras

### Karnaugh Maps

For up to four variables (propositional symbols) a diagrammatic method of simplification called **Karnaugh maps** works quite well.

For every propositional function of k=2,3,4 variables we construct a rectangular array of  $2^k$  cells. We mark the squares corresponding to the value true with eg "+" and try to cover these squares with as few rectangles with sides 1 or 2 or 4 as possible.

### Example



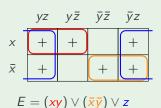
Canonical form would consist of writing all cells separately (6 clauses).

## Karnaugh Maps

For optimisation, the idea is to cover the + squares with the minimum number of rectangles. One *cannot* cover any empty cells.

- The rectangles can go 'around the corner'/the actual map should be seen as a torus.
- Rectangles must have sides of 1, 2 or 4 squares (three adjacent cells are useless).

### **Example**

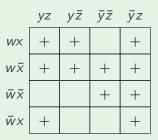


Canonical form would consist of writing all cells separately (6 clauses).

### Exercise

### **Exercise**

RW: 10.6.6(c)



### Outline

Boolean Logic

Boolean Functions

Conjunctive and Disjunctive Normal Form

Karnaugh Maps

Boolean Algebras

## Definition: Boolean Algebra

#### **Definition**

A **Boolean algebra** is a structure  $(T, \vee, \wedge, ', 0, 1)$  where

- $0, 1 \in T$
- $\vee, \wedge : T \times T \to T$  (called **join** and **meet** respectively)
- $': T \to T$  (called **complementation**)

and the following laws hold for all  $x, y, z \in T$ :

Commutativity:  $x \lor y = y \lor x, \quad x \land y = y \land x$ 

Associativity:  $(x \lor y) \lor z = x \lor (y \lor z)$ 

 $(x \wedge y) \wedge z = x \wedge (y \wedge z)$ 

Distributivity:  $x \lor (y \land z) = (x \lor y) \land (x \lor z)$ 

 $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$ 

Identity:  $x \lor 0 = x, \quad x \land 1 = x$ 

Complementation:  $x \lor x' = 1, \quad x \land x' = 0$ 

### **Example**

The set of subsets of a set X:

```
T: \operatorname{Pow}(X)
\lor (\operatorname{join}): \cup
\land (\operatorname{meet}): \cap
' (\operatorname{complement}): \cdot^{c}
\emptyset: \emptyset
1: X (\mathcal{U})
```

The Laws of Boolean algebra follow from the Laws of Set Operations.

### **Example**

The two element Boolean Algebra:

$$\mathbb{B} = (\{\mathsf{true}, \mathsf{false}\}, \|, \&\&, !, \mathsf{false}, \mathsf{true})$$

where  $!, \&\&, \parallel$  are defined as:

- !true = false; !false = true,
- true && true = true; ...
- true || true = true; ...

#### **Example**

Cartesian products of  $\mathbb{B}$ , that is *n*-tuples of 0's and 1's with Boolean operations, e.g.  $\mathbb{B}^4$ :

$$\begin{array}{ll} \textit{join:} & (1,0,0,1) \lor (1,1,0,0) = (1,1,0,1) \\ \textit{meet:} & (1,0,0,1) \land (1,1,0,0) = (1,0,0,0) \\ \textit{complement:} & (1,0,0,1)' = (0,1,1,0) \\ & \emptyset: & (0,0,0,0) \\ & \mathbb{1}: & (1,1,1,1). \end{array}$$

### **Example**

Functions from any set S to  $\mathbb{B}$ ; that is,  $\mathbb{B}^S$ 

If 
$$f, g: S \longrightarrow \mathbb{B}$$
 then

$$(f \lor g): S \to \mathbb{B}$$
 defined by  $s \mapsto f(s) \parallel g(s)$   $(f \land g): S \to \mathbb{B}$  defined by  $s \mapsto f(s) \&\& g(s)$   $f': S \to \mathbb{B}$  defined by  $s \mapsto !f(s)$   $\mathbb{C} : S \to \mathbb{B}$  is the function  $s \mapsto 0$   $\mathbb{C} : S \to \mathbb{B}$  is the function  $s \mapsto 1$ 

## Proofs in Boolean Algebras

Show an identity holds using the laws of Boolean Algebra, then that identity holds in all Boolean Algebras.

#### **Example**

In all Boolean Algebras

$$x \land x = x$$

for all  $x \in T$ .

Proof:

$$\begin{array}{lll} x &= x \wedge \mathbb{1} & & [\mathsf{Identity}] \\ &= x \wedge (x \vee x') & & [\mathsf{Complement}] \\ &= (x \wedge x) \vee (x \wedge x') & & [\mathsf{Distributivity}] \\ &= (x \wedge x) \vee \mathbb{0} & & [\mathsf{Complement}] \\ &= (x \wedge x) & & [\mathsf{Identity}] \end{array}$$

### Duality

#### **Definition**

If E is an expression defined using variables (x, y, z, etc), constants (0 and 1), and the operations of Boolean Algebra  $(\land, \lor, \text{ and }')$  then dual(E) is the expression obtained by replacing  $\land$  with  $\lor$  (and vice-versa) and 0 with 1 (and vice-versa).

#### **Definition**

If  $(T, \vee, \wedge, ', \mathbb{O}, \mathbb{I})$  is a Boolean Algebra, then  $(T, \wedge, \vee, ', \mathbb{I}, \mathbb{O})$  is also a Boolean algebra, known as the **dual** Boolean algebra.

### Theorem (Principle of duality)

If you can show  $E_1 = E_2$  using the laws of Boolean Algebra, then  $dual(E_1) = dual(E_2)$ .

## Duality

### Example

We have shown  $x \wedge x = x$ .

By duality:  $x \lor x = x$ .