

# CMPS 101

## Homework Assignment 2

### Solutions

1. p.52: 3.1-1

Let  $f(n)$  and  $g(n)$  be asymptotically non-negative functions. Using the basic definition of  $\Theta$ -notation, prove that  $f(n) + g(n) = \Theta(\max(f(n), g(n)))$ .

**Proof:**

Since  $f(n)$  and  $g(n)$  are asymptotically non-negative, there exists a positive constant  $n_0$  such that  $f(n) \geq 0$  and  $g(n) \geq 0$  for all  $n \geq n_0$ . For such  $n$  we have

$$\begin{aligned} 0 &\leq \max(f(n), g(n)) \\ &\leq \min(f(n), g(n)) + \max(f(n), g(n)) \\ &\leq 2 \cdot \max(f(n), g(n)). \end{aligned}$$

But  $f(n) + g(n) = \min(f(n), g(n)) + \max(f(n), g(n))$ , so for all  $n \geq n_0$  we have

$$0 \leq 1 \cdot \max(f(n), g(n)) \leq f(n) + g(n) \leq 2 \cdot \max(f(n), g(n)).$$

Thus  $f(n) + g(n) = \Theta(\max(f(n), g(n)))$ , as required. ■

2. p. 53: 3.1-4

Determine whether the following statements are true or false.

a.  $2^{n+1} = O(2^n)$

**Solution:** True since  $2^{n+1} = 2 \cdot 2^n = \text{const} \cdot 2^n = O(2^n)$ . ■

b.  $2^{2n} = O(2^n)$

**Solution:** False since  $2^{2n} = (2^2)^n = 4^n = \omega(2^n)$  and since  $\omega(2^n) \cap O(2^n) = \emptyset$ . ■

3. p.61: 3-2abcdef

Indicate, for each pair of expressions ( $A$ ,  $B$ ) in the table below, whether  $A$  is  $O$ ,  $o$ ,  $\Omega$ ,  $\omega$ , or  $\Theta$  of  $B$ . Assume that  $k \geq 1$ ,  $\varepsilon > 0$ , and  $c > 1$  are constants. Place 'yes' or 'no' in each of the empty cells below, and justify your answers.

|    | A           | B               | $O$ | $o$ | $\Omega$ | $\omega$ | $\Theta$ |
|----|-------------|-----------------|-----|-----|----------|----------|----------|
| a. | $\lg^k n$   | $n^\varepsilon$ | yes | yes | no       | no       | no       |
| b. | $n^k$       | $c^n$           | yes | yes | no       | no       | no       |
| c. | $\sqrt{n}$  | $n^{\sin n}$    | no  | no  | no       | no       | no       |
| d. | $2^n$       | $2^{n/2}$       | no  | no  | yes      | yes      | no       |
| e. | $n^{\lg c}$ | $c^{\lg n}$     | yes | no  | yes      | no       | yes      |
| f. | $\lg(n!)$   | $\lg(n^n)$      | yes | no  | yes      | no       | yes      |

**Justifications:**

- a. Applying l'Hopital's rule  $[k]$  times, we get  $\lim_{n \rightarrow \infty} \frac{(\lg n)^k}{n^\varepsilon} = 0$ , whence  $\lg^k n = o(n^\varepsilon)$ .
- b. Again, by  $[k]$  applications of l'Hopital's rule,  $\lim_{n \rightarrow \infty} \frac{n^k}{c^n} = 0$ , whence  $n^k = o(c^n)$ .
- c. Let  $c > 0$ , and observe that  $n^{\sin n} > c\sqrt{n}$  for  $n = 2\pi k + \frac{\pi}{2}$ , where  $k$  is a sufficiently large positive integer. Therefore the inequality  $n^{\sin n} \leq c\sqrt{n}$  is false for arbitrarily large  $n$ , and hence  $n^{\sin(n)} \neq O(\sqrt{n})$ . Also note that  $n^{\sin n} < c\sqrt{n}$  for  $n = 2\pi k + \frac{3\pi}{2}$ , where  $k$  is sufficiently large, and therefore the inequality  $n^{\sin n} \geq c\sqrt{n}$  is false for arbitrarily large  $n$ , whence  $n^{\sin(n)} \neq \Omega(\sqrt{n})$ .
- d. Observe  $\lim_{n \rightarrow \infty} (2^n / 2^{n/2}) = \lim_{n \rightarrow \infty} (\sqrt{2})^n = \infty$ , and therefore  $2^n = \omega(2^{n/2})$ .
- e. The identity  $x^{\log_b y} = y^{\log_b x}$  implies that the two functions are equal:  $n^{\lg c} = c^{\lg n}$ .
- f. Note  $\lg(n^n) = n \lg(n)$ , and Stirling's formula implies  $\lg(n!) = \Theta(n \lg(n))$ .

## 4. p.62: 3-4deh

Let  $f(n)$  and  $g(n)$  be asymptotically positive functions (i.e.  $f(n) > 0$  and  $g(n) > 0$  for sufficiently large  $n$ .) Prove or disprove the following statements.

- d.  $f(n) = O(g(n))$  implies  $2^{f(n)} = O(2^{g(n)})$ . **False**

**Counter-Example:**

Let  $f(n) = 2n$  and  $g(n) = n$ . Then  $2^{g(n)} = 2^n$  and  $2^{f(n)} = 2^{2n} = 4^n = \omega(2^n)$ , so  $2^{f(n)} = \omega(2^{g(n)})$ , and therefore  $2^{f(n)} \neq O(2^{g(n)})$ . ■

- e.  $f(n) = O((f(n))^2)$ . **False**

**Counter-Example:**

Let  $f(n) = 1/n$ . Then  $f(n) = \omega((f(n))^2)$  since  $\lim_{n \rightarrow \infty} \frac{f(n)}{f(n)^2} = \lim_{n \rightarrow \infty} \frac{1}{f(n)} = \lim_{n \rightarrow \infty} n = \infty$ , and hence  $f(n) \neq O((f(n))^2)$ . ■

- h.  $f(n) + o(f(n)) = \Theta(f(n))$ . **True**

**Proof:**

In the above formula,  $o(f(n))$  stands for some anonymous function  $h(n)$  in the class  $o(f(n))$ , whence  $\lim_{n \rightarrow \infty} \frac{h(n)}{f(n)} = 0$ . Thus  $\lim_{n \rightarrow \infty} \frac{f(n)+h(n)}{f(n)} = \lim_{n \rightarrow \infty} \left(1 + \frac{h(n)}{f(n)}\right) = 1$ , and  $f(n) + h(n) = \Theta(f(n))$ , as claimed. ■

5. Let  $g(n)$  be an asymptotically non-negative function. Prove that  $o(g(n)) \cap \Omega(g(n)) = \emptyset$ .

**Proof:**

Assume to get a contradiction that  $f(n) \in o(g(n)) \cap \Omega(g(n))$ . Then since  $f(n) = \Omega(g(n))$  we have

$$(1) \quad \exists c_1 > 0, \exists n_1 > 0, \forall n \geq n_1: \quad 0 \leq c_1 g(n) \leq f(n)$$

and since  $f(n) = o(g(n))$  we have

$$(2) \quad \forall c_2 > 0, \exists n_2 > 0, \forall n \geq n_2: \quad 0 \leq f(n) < c_2 g(n)$$

Let  $c_2 = c_1$ . Then  $c_2 > 0$ , and by (2) there exists  $n_2 > 0$  such that  $0 \leq f(n) < c_1 g(n)$  for all  $n \geq n_2$ . Pick any  $m \geq \max(n_1, n_2)$ . Then by (1) and (2) we have  $0 \leq c_1 g(m) \leq f(m) < c_1 g(m)$ , and hence  $c_1 g(m) < c_1 g(m)$ , a contradiction. Our assumption was therefore false, and no such function  $f(n)$  can exist. We conclude that  $o(g(n)) \cap \Omega(g(n)) = \emptyset$ . ■

6. Prove that  $3^{2^n} = o(2^{3^n})$ .

**Proof:**

Observe that  $\ln\left(\frac{3^{2^n}}{2^{3^n}}\right) = 2^n(\ln 3) - 3^n(\ln 2) \rightarrow -\infty$ , since  $3^n = \omega(2^n)$ . Therefore  $\frac{3^{2^n}}{2^{3^n}} \rightarrow e^{-\infty} = 0$ , whence  $3^{2^n} = o(2^{3^n})$ . ■