

Technical Document: Lab 1

Overview

Team

| Role: | Name: |
|----------------|-----------------|
| Manager | Luis Arcos |
| Technical Lead | Madylin Teel |
| Developer | Coleman Bixler |
| Developer | Garrett Snow |
| Developer | An Vu |
| Developer | Nathan Stephani |

Conventions

Naming

| | |
|--------------------------|--|
| Branch names | Example: {dev,techlead,manager}.<user>/<objective>/<task> i.e., dev.madyteel/obj-a/t-01 |
| Generated output folders | Example: task##/<operation>-versions/ i.e., task02/split-versions/ |
| C code file names | Example: var##-<operation/schedule>.x i.e., var00-ssa-schedule-blur3x3.c |
| | |
| | |

Helpful Notes

Schedules and Operations: Compatibility at a Glance

| | blur_2x2_ijqr | blur_3x3_ijqr | blur_4x4_ijqr | conv_4x3_ijqr | conv_4x3_qrji | jacobi2d_5pt_ijqr | jacobi2d_21pt_ijqr | toeplitz_4x4_ijqr |
|-------------------|---------------|---------------|---------------|---------------|---------------|-------------------|--------------------|-------------------|
| complex_0 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| inter_ijjq_qr_ri | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| inter_riqqr_jr_ji | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| split_ijqr_2 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| split_rqji_2 | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| unroll_q_r | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| unroll_r_q | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| complex_l3_v0 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| complex_l4_v0 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| complex_l5_v0 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |
| complex_l6_v0 | ✓ | ✓ | ✓ | ✓ | ✗ | ✓ | ✓ | ✓ |

Objective A: Handspun Simple Schedule Code Generator

Objective A focuses on the process of applying schedule transformation (by hand) on an operation in Static Single Assignment (SSA) form.

Static Single Assignment (SSA) – A representation used in compilers where each variable is assigned *exactly once* and defined before it is used.

Please see `LAB-SCHEDULE-DECOUPLED/lab_inital/operations/` for all operations which need to be converted to SSA format.

Please see `LAB-SCHEDULE-DECOUPLED/lab_inital/examples/basic-ssa/{ssa_var000.c, ssa_var001.c}` for examples on how to convert the operations into SSA form.

Tentative Delegation of Tasks

| | |
|-----------------|-----|
| Coleman Bixler | T03 |
| Garrett Snow | T01 |
| An Vu | T04 |
| Nathan Stephani | T02 |
| Luis Arcos | T00 |

T00: Handspun Baseline

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications in SSA form.

Implementation

Use `ssa_var000.c` and `ssa_var001.c` as references for structuring loop transformations while maintaining SSA-compliance [from lab_inital/examples/basic-ssa/].

Steps:

1. Convert each individual operation into SSA form:
 - a. Transform each operation file into SSA-compliant C code.
 - i. Rewrite each computation such that every value is assigned exactly and only once.
 - ii. Introduce new variables instead of reassigning existing ones.
 - b. Ensure there is no variable reassignment!
2. Implement the SSA-compliant C code:
 - a. Must write manual implementations of the SSA transformations.
 - b. Can verify against `baseline.c` [from lab_initial/examples/basic-ssa].
 - c. Store all ssa versions in their own folder [i.e., task00/ssa-versions/].
 - d. Ensure consistent naming across versions [i.e., `ssa-var00`].
3. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
4. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
5. Compare SSA implementations:
 - a. Compare the performance across various operations (such as blur, conv, jacobi2d, etc.).
 - b. Identify which operations benefit the most from SSA transformation.
6. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

Benchmarking

T01: Handspun + Unroll

Implement, Verify, Benchmark, and Compare for Five (5) of the Operation Specifications in SSA form.

Implementation

Steps:

1. If not already done, convert each operation into SSA form (see objective-a T00 for more details):
 - a. Each computation must be assigned to a unique variable.
 - b. No variable reassignment is allowed!
2. Apply loop unrolling:

Unrolling reduces loop overhead by manually duplicating loop iterations.

 - a. Select any (5) operations [from /lab_initial/operations/]. (For example: blur, conv, jacobi2d, etc.)
 - b. Manually unroll the inner loop:
 - i. Identify which loop needs to be unrolled (either j0, q0, or r0 for example from slides)

- ii. Example given in the slides for this lab, specifically slide #6.

Schedules: unroll <id>

The following will take the code on the left and transform it to the code on the right. Loops will be labelled by their index variable.

unroll r0

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // Q=3xR=4
    float weights[] =
    //r=0      1      2      3
    {-1.1, -1.1, 1.2, -2.1, // q=0
     1.1, 2.1, 1.2, 2.2, // q=1
     2.1, 0.1, 0.2, 1.2}; // q=2

    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0 = 0; j0 < n0; ++j0 )
            for( int q0 = 0; q0 < Q; ++q0 )
                for( int r0 = 0; r0 < R; ++r0 )
                    y[i0+n0+j0] += weights[q0*(R)+r0] *
                    x[(q0+i0)%m0*n0 + ((r0+j0)%n0)];
    END_INSTRUMENTATION;
}
```

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // BEGIN_INSTRUMENTATION; // func: COMPUTE_NAME
    for( int i0 = 0; i0 < m0; ++i0 )
        // BEGIN_INSTRUMENTATION; // loop:i0
        for( int j0 = 0; j0 < n0; ++j0 )
            // BEGIN_INSTRUMENTATION; // loop:j0
            for( int q0 = 0; q0 < Q; ++q0 )
                // BEGIN_INSTRUMENTATION; // loop:q0
                {
                    BEGIN_INSTRUMENTATION; // loop:r0
                    for( int r0 = 0; r0 < R; ++r0 )
                        y[i0+n0+j0] += weights[q0*(R)+r0] *
                        x[(q0+i0)%m0*n0 + ((r0+j0)%n0)];
                    // END_INSTRUMENTATION; // loop:r0
                }
            // END_INSTRUMENTATION; // loop:j0
        // END_INSTRUMENTATION; // loop:i0
    // END_INSTRUMENTATION; // func: COMPUTE_NAME
}
```

3. Implement the SSA-compliant unrolled C code:
 - i. Store all unrolled implementations in their own folder [i.e., /task01/unrolled-versions/].
 - ii. Ensure consistent naming across versions [i.e., `unroll-var00`].
4. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
 - b. Ensure unrolling does not introduce any unexpected behavior.
5. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Unrolled SSA vs. non-unrolled SSAx
 - ii. Unrolled SSA vs. Baseline (non-SSA)
6. Compare results:
 - a. Compare the execution time, FLOPs, and memory usage for:
 - i. Unrolled SSA vs. non-unrolled SSA
 - ii. Unrolled SSA vs. Baseline (non-SSA)
 - iii. Which operations benefit the most from loop unrolling?
 - iv. Are there any cases in which unrolling hurts performance?
7. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

T02: Handspun + Split

Implement, Verify, Benchmark, and Compare for Five (5) of the Operation Specifications in SSA form.

Implementation

Steps:

1. If not already done, convert each operation into SSA form (see objective-a T00 for more details):
 - a. Each computation must be assigned to a unique variable.
 - b. No variable reassignment is allowed!
2. Apply loop splitting:

Splitting reduces iteration overhead and improves cache locality by breaking a single loop into two nested loops.

 - a. Select any (5) operations [from /lab_initial/operations/]. (For example: blur, conv, jacobi2d, etc.)
 - b. Manually split the loops:
 - i. Identify which loops need to be split (either j0, q0, or r0 for example from slides)
 - ii. Example given in the slides for this lab, specifically slide #7.

Schedules: split <id> <out-id> <inner-id> <factor>

The following will take the code on the left and transform it to the code on the right. Loops will be labelled by their index variable.

split j0 j0_o j0_i 2

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // Q=3xR=4
    float weights[] =
    {
        //r=0      1      2      3
        -1.1, -1.1, 1.2, -2.1, // q=0
        1.1, 2.1, 1.2, 2.2, // q=1
        2.1, 0.1, 0.2, 1.2; // q=2
    };

    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0 = 0; j0 < n0; ++j0 )
            for( int q0 = 0; q0 < (Q); ++q0 )
                for( int r0 = 0; r0 < (R); ++r0 )
                    y[i0*n0+j0] += weights[q0*(R)+r0] *
                        x[ ((q0-i0)%m0)*n0 + ((r0+j0)%n0) ];
    END_INSTRUMENTATION;
}
```

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        for( int j0_o = 0; j0_o < n0; j0_o+=2 )
            for( int j0_i = 0; j0_i < 2; ++j0_i )
            {
                // BEGIN_INSTRUMENTATION; // loop:j0
                for( int q0 = 0; q0 < (Q); ++q0 )
                {
                    // BEGIN_INSTRUMENTATION; // loop:q0
                    for( int r0 = 0; r0 < (R); ++r0 )
                    {
                        BEGIN_INSTRUMENTATION; // loop:r0
                        y[i0*n0+j0_o+j0_i] += weights[q0*(R)+r0] *
                            x[ ((q0-i0)%m0)*n0 + ((r0+(j0_o+j0_i))%n0) ];
                        END_INSTRUMENTATION; // loop:r0
                    }
                }
            }
    }
}
```

3. Implement the SSA-compliant split C code:
 - i. Store all split implementations in their own folder [i.e., /task02/split-versions/].
 - ii. Ensure consistent naming across versions [i.e., `split-var00`].
4. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
 - b. Ensure loop splitting does not introduce any unexpected behavior.
5. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Split SSA vs. non-split SSA
 - ii. Split SSA vs. Baseline (non-SSA)
6. Compare results:
 - a. Compare the execution time, FLOPs, and memory usage for:
 - i. Split SSA vs. non-split SSA
 - ii. Split SSA vs. Baseline (non-SSA)
 - iii. Which operations benefit the most from loop splitting?

- iv. Are there any cases in which splitting hurts performance perhaps by introducing overhead?
- 7. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

T03: Handspun + Interchange

Implement, Verify, Benchmark, and Compare for One (1) Operation against six (6) different ways to interchange the loops in SSA form.

Implementation

Steps:

1. If not already done, convert the operation into SSA form (see objective-a T00 for more details):
 - a. Each computation must be assigned to a unique variable.
 - b. No variable reassignment is allowed!
2. Apply loop interchange in (6) different ways:

Loop interchange affects both cache locality and execution efficiency by changing the order of nested loops.

 - a. Select any (1) operation [from /lab_initial/operations/]. (For example: blur, conv, jacobi2d, etc.)
 - b. Manually interchange the loops:
 - i. Identify which loops need to be interchanged.
 - ii. Example given in the slides for this lab, specifically slide #8.

Schedules: interchange <id-a> <id-b>

The following will take the code on the left and transform it to the code on the right. Loops will be labelled by their index variable.

interchange r0 j0

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // Q=3xR=4
    float weights[] =
    {
        //r0      1      2      3
        [-1.1, -1.1, 1.2, -2.1, // q=0
         -1.1, -2.1, 1.2, 2.2, // q=1
         -2.1, 0.1, 0.2, 1.2]; // q=2
    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0 = 0; j0 < n0; ++j0 )
            for( int q0 = 0; q0 < Q; ++q0 )
                for( int r0 = 0; r0 < R; ++r0 )
                    y[(i0+n0*j0) + weights[q0*(R)r0] +
                     x[(i0+10)%m0+r0 + ((r0*j0)%n0) ];
    END_INSTRUMENTATION;
}
```

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // BEGIN_INSTRUMENTATION; // func: COMPUTE_NAME
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        // BEGIN_INSTRUMENTATION; // loop:i0
        for( int r0 = 0; r0 < R; ++r0 )
        {
            // BEGIN_INSTRUMENTATION; // loop:r0
            for( int q0 = 0; q0 < Q; ++q0 )
            {
                // BEGIN_INSTRUMENTATION; // loop:q0
                for( int j0 = 0; j0 < n0; ++j0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:j0
                    y[(i0+n0*j0) + weights[q0*(R)r0] +
                     x[(i0+10)%m0+r0 + ((r0*j0)%n0) ];
                    END_INSTRUMENTATION; // loop:j0
                }
            }
            END_INSTRUMENTATION; // loop:q0
        }
        END_INSTRUMENTATION; // loop:r0
    }
    END_INSTRUMENTATION; // loop:i0
    // END_INSTRUMENTATION; // func:COMPUTE_NAME
}
```

3. Implement the SSA-compliant split C code:
 - i. Store all interchange implementations in their own folder [i.e., /task03/interchange-versions/].
 - ii. Ensure consistent naming across versions [i.e., `ssa-interchange-var00`].
4. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
 - b. Ensure loop interchanges do not introduce any unexpected behavior.

5. Benchmark performance:
 - a. Run ``run-all-measure.sh``. (WIP)
 - i. Each interchanged version vs. baseline SSA
 - ii. Each interchanged version vs. other interchanged versions
6. Compare results:
 - a. Compare the execution time, FLOPs, and memory usage for:
 - i. Each interchanged version vs. baseline SSA
 - ii. Which interchanges give the most optimal results?
 - iii. Are there any cases in which interchanging worsens performance?
7. Clean up generated files:
 - a. Run ``run-clean-all.sh`` to clean up all files generated by a given area. (WIP)

T04: Handspun + Complex Schedule (3 Commands)

Implement, Verify, Benchmark, and Compare for three (3) of the Operation Specifications using the same schedule in SSA form.

Implementation

Steps:

1. Before applying complex scheduling, ensure the (3) desired operations are in SSA form (see objective-a T00 for more details):
 - a. Each computation must be assigned to a unique variable.
 - b. No variable reassignment is allowed!
2. Either define or utilize premade complex schedule:

A complex schedule applies (in this case, 3) transformations to the loop structure (which could be any combination of: unrolling, splitting, interchanging, etc.).

Predefined schedules can be found in `/lab_initial/schedules/complex_mix/`. *The first two options in this folder are (3) command schedules.*

Example of a (3) command complex schedule:

- a. Apply loop splitting: split `j0` into `j0_outer` and `j0_inner`
- b. Apply loop interchange: swap `q0` and `j0_outer`
- c. Apply loop unrolling: unroll `j0_inner` by a factor of 2
3. Apply the complex schedule to (3) operations:
 - a. Select any (3) operations [from `/lab_initial/operations/`]. (For example: `blur`, `conv`, `jacobi2d`, etc.)
 - b. Apply the same transformation sequence across all (3) operations.

- c. Example given in the slides for this lab, specifically slide #9 and slide #10.

Schedules: Multiple Commands

Multiple scheduling commands can be used and are executed from top to bottom.

split j0 j0_o j0_i 2
interchange r0 j0_i
unroll j0_i

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    // Q=3xR=4
    float weights[] =
    {
        // r=0
        1, 2, 3
        {-1.1, -1.1, 1.2, -2.1, // q=0
         -1.1, -2.1, 1.2, 2.2, // q=1
         -2.1, 0.1, 0.2, 1.2}; // q=2
    };

    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0 = 0; j0 < n0; ++j0 )
            for( int q0 = 0; q0 < (Q); ++q0 )
                for( int r0 = 0; r0 < (R); ++r0 )
                    y[i0*n0+j0] += weights[q0*(R)+r0] *
                        x[(q0-i0)%m0*n0 + ((r0+j0)%n0)];
    END_INSTRUMENTATION;
}
```

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        for( int j0_o = 0; j0_o < n0; j0_o+=2 )
        {
            for( int r0 = 0; r0 < (R); ++r0 )
            {
                for( int q0 = 0; q0 < (Q); ++q0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:q0
                    { // j0_i=0
                        y[i0*n0+(j0_o+0)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+0))%n0)];
                    }
                    { // j0_i=1
                        y[i0*n0+(j0_o+1)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+1))%n0)];
                    }
                }
            }
        }
        END_INSTRUMENTATION; // loop:c0
    }
}
```

Schedules: Multiple Commands

conv

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    // Q=3xR=4
    float weights[] =
    {
        // r=0
        1, 2, 3
        {-1.1, -1.1, 1.2, -2.1, // q=0
         -1.1, -2.1, 1.2, 2.2, // q=1
         -2.1, 0.1, 0.2, 1.2}; // q=2
    };

    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0_o = 0; j0_o < n0; ++j0_o )
            for( int q0 = 0; q0 < (Q); ++q0 )
                for( int r0 = 0; r0 < (R); ++r0 )
                    y[i0*n0+j0_o] += weights[q0*(R)+r0] *
                        x[(q0-i0)%m0*n0 + ((r0+j0_o)%n0)];
    END_INSTRUMENTATION;
}
```

split j0 j0_o j0_i 2

opt1 = Split(j0,j0_o,j0_i,2, conv)

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        for( int j0_o = 0; j0_o < n0; j0_o+=2 )
        {
            for( int r0 = 0; r0 < (R); ++r0 )
            {
                for( int q0 = 0; q0 < (Q); ++q0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:q0
                    { // j0_i=0
                        y[i0*n0+(j0_o+0)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+0))%n0)];
                    }
                    { // j0_i=1
                        y[i0*n0+(j0_o+1)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+1))%n0)];
                    }
                }
            }
        }
    }
}
```

opt2 = Interchange(r0,j0_i, opt1)

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        for( int j0_o = 0; j0_o < n0; j0_o+=2 )
        {
            for( int q0 = 0; q0 < (Q); ++q0 )
            {
                for( int r0 = 0; r0 < (R); ++r0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:q0
                    { // j0_i=0
                        y[i0*n0+(j0_o+0)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+0))%n0)];
                    }
                    { // j0_i=1
                        y[i0*n0+(j0_o+1)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+1))%n0)];
                    }
                }
            }
        }
    }
}
```

interchange r0 j0_i

unroll j0_i

opt3 = Unroll(j0_i, opt2)

```
void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        for( int j0_o = 0; j0_o < n0; j0_o+=2 )
        {
            for( int q0 = 0; q0 < (Q); ++q0 )
            {
                for( int r0 = 0; r0 < (R); ++r0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:q0
                    { // j0_i=0
                        y[i0*n0+(j0_o+0)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+0))%n0)];
                    }
                    { // j0_i=1
                        y[i0*n0+(j0_o+1)] += weights[q0*(R)+r0] *
                            x[(q0-i0)%m0*n0 + ((r0+(j0_o+1))%n0)];
                    }
                }
            }
        }
    }
}
```

- Implement SSA-compliant scheduled code:
 - Store all scheduled implementations in their own folder [i.e., /task04/scheduled-versions/].
 - Ensure consistent naming across versions [i.e., `ssa-schedule-blur3x3.c`].
- Verify correctness:
 - Run `run-all-verify.sh`. (WIP)
 - Ensure complex schedule does not introduce any unexpected behavior.
- Benchmark performance:
 - Run `run-all-measure.sh`. (WIP)

- i. Scheduled SSA vs. non-scheduled SSA
 - ii. Scheduled SSA vs. baseline (non-SSA)
- 7. Compare results:
 - a. Compare the execution time, FLOPs, and memory usage for:
 - i. Scheduled SSA vs. non-scheduled SSA
 - ii. Scheduled SSA vs. baseline (non-SSA)
 - iii. Which combinations of transformations performs the best?
 - iv. Are there any cases in which a specific operation benefits more from the schedule?
- 8. Clean up generated files:
 - a. Run ``run-clean-all.sh`` to clean up all files generated by a given area. (WIP)

Optional: For Additional Points

T01-Opt00: Handspun + Unroll(r) vs Unroll(q) [Optional]

Implement, Verify, Benchmark, and Compare for One (1) Operation, but two (2) different loops in SSA form.

T02-Opt00: Handspun + Split(i) vs Split(j) vs Split(q) vs Split(r) [Optional]

Implement, Verify, Benchmark, and Compare for One (1) Operation, but four (4) different loops to be split in SSA form.

T02-Opt01: Handspun + Split(2) vs Split(4) vs Split(8) [Optional]

Implement, Verify, Benchmark, and Compare for One (1) Operation, One (1) loop, but four (4) different split factors) in SSA form.

T04-Opt00: Handspun + Complex Schedule (3 Commands) [Optional]

Implement, Verify, Benchmark, and Compare for one (1) of the Operation Specification using three (3) different schedules in SSA form.

Objective B: Standalone Simple Schedule Code Generator

The goal of this objective is to focus on automating what was previously done manually in objective-a.

For this objective we will only focus on simple schedules still and additionally only those with one line, for example:

```
```
```

```
unroll j0 2
```

```
split j0 j0_outer j0_inner 4
```

```
interchange j0 q0
```

```
```
```

For this objective we do not yet want to mess with handling complex schedules.

Tentative Delegation of Tasks

| | |
|-----------------|-----|
| Coleman Bixler | T03 |
| Garrett Snow | T01 |
| An Vu | T00 |
| Nathan Stephani | T02 |

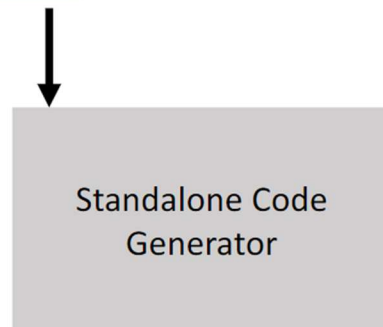
T00: Simple Standalone Baseline

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications in SSA form.

Objective B: Standalone Simple Schedule

Operation Specification

```
3 4
-1.1 -1.1 1.2 -2.1
-1.1 -2.1 -1.2 2.2
-2.1 0.1 0.2 1.2
i0 j0 q0 r0
```



```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // Q=3xR=4
    float weights[] =
    //r=0 1 2 3
    { 1.1, -1.1, 1.2, -2.1, // q=0
      1.1, -2.1, -1.2, 2.2, // q=1
      -2.1, 0.1, 0.2, 1.2 }; // q=2

    BEGIN_INSTRUMENTATION;
    for( int i0 = 0; i0 < m0; ++i0 )
        for( int j0 = 0; j0 < n0; ++j0 )
            for( int q0 = 0; q0 < (Q); ++q0 )
                for( int r0 = 0; r0 < (R); ++r0 )
                    y[i0*n0+j0] += weights[q0*(R)+r0] *
                        x[ ((q0+i0)%m0)*n0 + ((r0+j0)%n0) ];
    END_INSTRUMENTATION;
}
```

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Implementation

Steps:

1. Instead of manually writing SSA-compliant code, we will instead:
 - a. Write a generator that takes in an operation description and a simple schedule, and
 - b. Automatically generate SSA-compliant C code as the output.
2. Implement the simple standalone code generator:
 - a. Read in an operation file.
 - b. Read an optional scheduling file.
 - c. Generate an SSA-compliant C code implementation.
 - d. Ensure SSA compliance.
 - e. Handles basic scheduling, if included.
 - f. Preserves the original operation logic.
3. Verify correctness:

- a. Run `run-all-verify.sh`. (WIP)
4. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Generated SSA vs. Handwritten SSA
 - ii. Generated SSA vs. baseline (non-SSA)
5. Compare results:
 - a. Generated SSA vs. Handwritten SSA
 - b. Generated SSA vs. baseline (non-SSA)
 - c. Are there any discrepancies between manually written and generated SSA versions?
 - d. Does SSA that is generated produce unnecessary overhead?
6. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

T01: Simple Standalone + Unroll

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

Objective B: Standalone Simple (Unroll)

Operation Specification

```
3 4
-1.1 -1.1 1.2 -2.1
-1.1 -2.1 -1.2 2.2
-2.1 0.1 0.2 1.2
i0 j0 q0 r0
```

Standalone Code
Generator

unroll r0

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
    // BEGIN INSTRUMENTATION; // func: COMPUTE_NAME
    for( int i0 = 0; i0 < m0; ++i0 )
        // BEGIN INSTRUMENTATION; // loop:i0
        for( int j0 = 0; j0 < n0; ++j0 )
            // BEGIN INSTRUMENTATION; // loop:j0
            for( int q0 = 0; q0 < (Q); ++q0 )
                // BEGIN INSTRUMENTATION; // loop:q0
                {
                    // r0=0
                    y[i0+n0+j0] += weights[q0+(R)+(0)] *
                        x[(q0-16)%n0 + ((0)-j0)%n0];
                }
                // r0=1
                y[i0+n0+j0] += weights[q0+(R)+(1)] *
                    x[(q0-16)%n0 + ((1)-j0)%n0];
            }
            // r0=2
            y[i0+n0+j0] += weights[q0+(R)+(2)] *
                x[(q0-16)%n0 + ((2)-j0)%n0];
        }
        // r0=3
        y[i0+n0+j0] += weights[q0+(R)+(3)] *
            x[(q0-16)%n0 + ((3)-j0)%n0];
    }
    // END INSTRUMENTATION; // loop:j0
}
// END INSTRUMENTATION; // loop:i0
// END INSTRUMENTATION; // func:COMPUTE_NAME
}
```

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Implementation

Modify the code generator to support unrolling.

Steps:

1. Instead of manually writing SSA-compliant code, we will instead:
 - a. Write a generator that takes in an operation description and a simple unrolling schedule, and
 - b. Automatically apply unrolling to loops then generate SSA-compliant C code as the output.

2. Implement the simple standalone code generator:
 - a. Read in schedule file.
 - b. Generate an SSA-compliant C code implementation with modified loop structure according to schedule.
 - c. Ensure SSA compliance.
 - d. Preserves the original operation logic.
 - e. Store all unrolled SSA implementations in their own folder [i.e., /unrolled-versions/].
 - f. Ensure consistent naming across versions [i.e., `ssa_unroll_blur3x3.c`].
3. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
4. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Unrolled SSA vs. non-unrolled SSA
 - ii. Unrolled SSA vs. baseline (non-SSA)
5. Compare results:
 - a. Unrolled SSA vs. non-unrolled SSA
 - b. Unrolled SSA vs. baseline (non-SSA)
6. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

T02: Simple Standalone + Split

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

Objective B: Standalone Simple (split)

Operation Specification

```
3 4
-1.1 -1.1 1.2 -2.1
-1.1 -2.1 -1.2 2.2
-2.1 0.1 0.2 1.2
i0 j0 q0 r0
```

Standalone Code
Generator

split j0 j0_o j0_i 2

```
void COMPUTE_NAME( int m0,
                  int n0,
                  float *x,
                  float *y )
{
  for( int i0 = 0; i0 < m0; ++i0 )
  {
    for( int j0_o = 0; j0_o < n0; j0_o += 2 )
      for( int j0_i = 0; j0_i < 2; ++j0_i )
      {
        // BEGIN_INSTRUMENTATION; // loop:j0
        for( int q0 = 0; q0 < (0); ++q0 )
        {
          // BEGIN_INSTRUMENTATION; // loop:q0
          for( int r0 = 0; r0 < (R); ++r0 )
          {
            // BEGIN_INSTRUMENTATION; // loop:r0
            y[i0*n0+(j0_o+j0_i)] += weight0[c0*(i0)+r0] *
              x[ ((q0+i0)%m0)*n0 + ((r0+(j0_o-j0_i))%n0) ];
            // END_INSTRUMENTATION; // loop:r0
          }
        }
      }
  }
}
```

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Implementation

Modify the code generator to support splitting.

Steps:

1. Instead of manually writing SSA-compliant code, we will instead:
 - a. Write a generator that takes in an operation description and a simple splitting schedule, and
 - b. Automatically apply splitting to loops then generate SSA-compliant C code as the output.
2. Implement the simple standalone code generator:
 - a. Read in schedule file.
 - b. Generate an SSA-compliant C code implementation with modified loop structure according to schedule.
 - c. Ensure SSA compliance.
 - d. Preserves the original operation logic.
 - e. Store all split SSA implementations in their own folder [i.e., /split-versions/].
 - f. Ensure consistent naming across versions [i.e., `ssa_split_blur3x3.c`].
3. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
4. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Split SSA vs. non-split SSA
 - ii. Split SSA vs. baseline (non-SSA)
5. Compare results:
 - a. Split SSA vs. non-split SSA
 - b. Split SSA vs. baseline (non-SSA)
6. Clean up generated files:
 - a. Run `run-clean-all.sh` to clean up all files generated by a given area. (WIP)

T03: Simple Standalone + Interchange

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

Objective B: Standalone Simple (interchange)

Operation Specification

```

3 4
-1.1 -1.1 1.2 -2.1
-1.1 -2.1 -1.2 2.2
-2.1 0.1 0.2 1.2
i0 j0 q0 r0

```

Standalone Code
Generator

interchange r0 j0

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```

void COMPUTE_NAME( int m0,
                   int n0,
                   float *x,
                   float *y )
{
    // BEGIN INSTRUMENTATION; // func: COMPUTE_NAME
    for( int i0 = 0; i0 < m0; ++i0 )
    {
        // BEGIN INSTRUMENTATION; // loop:i0
        for( int r0 = 0; r0 < (R); ++r0 )
        {
            // BEGIN INSTRUMENTATION; // loop:r0
            for( int q0 = 0; q0 < (Q); ++q0 )
            {
                // BEGIN INSTRUMENTATION; // loop:q0
                for( int j0 = 0; j0 < n0; ++j0 )
                {
                    BEGIN_INSTRUMENTATION; // loop:j0
                    y[i0+r0+q0] += weights[q0+r0+r0] *
                        x[i0+(q0+0)%m0+n0 + ((r0+j0)%n0)];
                    END_INSTRUMENTATION; // loop:j0
                }
            }
        }
    }
    // END INSTRUMENTATION; // loop:r0
}
// END INSTRUMENTATION; // func:COMPUTE_NAME

```

Implementation

Modify the code generator to support loop interchanging.

Steps:

1. Instead of manually writing SSA-compliant code, we will instead:
 - a. Write a generator that takes in an operation description and a simple interchanging schedule, and
 - b. Automatically apply interchanging to loops then generate SSA-compliant C code as the output.
2. Implement the simple standalone code generator:
 - a. Read in schedule file.
 - b. Generate an SSA-compliant C code implementation with modified loop structure according to schedule.
 - c. Ensure SSA compliance.
 - d. Preserves the original operation logic.
 - e. Store all interchanged SSA implementations in their own folder [i.e., /interchanged-versions/].
 - f. Ensure consistent naming across versions [i.e., `ssa_interchange_blur3x3.c`].
3. Verify correctness:
 - a. Run `run-all-verify.sh`. (WIP)
4. Benchmark performance:
 - a. Run `run-all-measure.sh`. (WIP)
 - i. Interchanged SSA vs. non-interchanged SSA
 - ii. Interchanged SSA vs. baseline (non-SSA)
5. Compare results:
 - a. Interchanged SSA vs. non-interchanged SSA

- b. Interchanged SSA vs. baseline (non-SSA)
- 6. Clean up generated files:
 - a. Run ``run-clean-all.sh`` to clean up all files generated by a given area. (WIP)

Objective C: EoC Simple Schedule Code Generator

Tentative Delegation of Tasks

| | |
|-----------------|---------------|
| Coleman Bixler | T04,T08 |
| Garrett Snow | T02,T06 |
| An Vu | T01(?),T05(?) |
| Nathan Stephani | T03,T07 |

T00: Implement AST to C codegen

Implement an abstract syntax tree (AST) to C code generator and then test its correctness.

Implementation

T01: Implement Operation Specification to AST

Implement and test.

T02: Implement the Compiler Pass (AST to AST) for Implementing Unroll

Implement and test.

T03: Implement the Compiler Pass (AST to AST) for Implementing Split

Implement and test.

T04: Implement the Compiler Pass (AST to AST) for Implementing Interchange

Implement and test.

T05: EoC Standalone Baseline

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications in SSA form.

T06: EoC Standalone + Unroll

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

T07: EoC Standalone + Split

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

T08: EoC Standalone + Interchange

Implement, Verify, Benchmark, and Compare for All of the Operation Specifications for All relevant schedule specifications in SSA form.

Objective D: Complex Schedule Code Generator

T00: Implement Generalized Schedule (AST to AST)

T01: Complex Schedule Generator + Unroll + Unroll

T02: Complex Schedule Generator + Split + Split + Split + Split

T03: Complex Schedule Generator + Interchange + Interchange + Interchange + Interchange

T04: Complex Schedule Generator + Mixed Schedule (length 3)

T05: Complex Schedule Generator + Mixed Schedule (length 4)

T06: Complex Schedule Generator + Mixed Schedule (length 5)

Objective E: Sensitivity Analysis

T00: Analysis of Complex Schedule (length 3)

T01: Compare the Same Schedule Across Different Operations

T02: Sensitivity Study (vary the factor of a loop split)

T03: Sensitivity Study (vary architecture)

T04: Analysis of Complex Schedule (length 4)

T05: Compare the Same Schedule Across Different Operations

T06: Sensitivity Study (vary the factor of a loop split)

T07: Sensitivity Study (vary architecture)

T08: Analysis of Complex Schedule (length >4)

T09: Compare the Same Schedule Across Different Operations

T10: Sensitivity Study (vary the factor of a loop split)

T11: Sensitivity Study (vary the factor of a loop split for two split factors)

T12: Sensitivity Study (vary architecture)

T13: Post your best performance plot and schedule on Canvas

Optional: Repeats for Additional Points

T08-Opt00: Analysis of Complex Schedule (length >4)

T09-Opt00: Compare the Same Schedule Across Different Operations

T10-Opt00: Sensitivity Study (vary the factor of a loop split)

T11-Opt00: Sensitivity Study (vary the factor of a loop split for two split factors)

T12-Opt00: Sensitivity Study (vary architecture)

T13-Opt00: Post your best performance plot and schedule on Canvas

Utilities & Testing