

OS 隨堂測驗:

1. ( O ) Race conditions are prevented by requiring that critical regions be protected by locks.
2. ( X ) The value of a counting semaphore can range only between 0 and 1.
3. ( O ) Mutex locks and binary semaphores are essentially the same thing.
4. ( O ) A nonpreemptive kernel is safe from race conditions on kernel data structures.
5. A race condition **\_\_B\_\_**.
  - A) results when several threads try to access the same data concurrently
  - B) results when several threads try to access and modify the same data concurrently
  - C) will result only if the outcome of execution does not depend on the order in which instructions are executed
  - D) None of the above
6. An instruction that executes atomically **\_\_B\_\_**.
  - A) must consist of only one machine instruction
  - B) executes as a single, uninterruptible unit
  - C) cannot be used to solve the critical section problem
  - D) All of the above
7. A counting semaphore **\_\_A\_\_**.
  - A) is essentially an integer variable
  - B) is accessed through only one standard operation
  - C) can be modified simultaneously by multiple threads
  - D) cannot be used to control access to a thread's critical sections
8. A mutex lock **\_\_B\_\_**.
  - A) is exactly like a counting semaphore
  - B) is essentially a boolean variable
  - C) is not guaranteed to be atomic
  - D) can be used to eliminate busy waiting
9. ( **A** ) What is the correct order of operations for protecting a critical section using mutex locks?
  - A) release() followed by acquire()
  - B) acquire() followed by release()
  - C) wait() followed by signal()
  - D) signal() followed by wait()
10. ( **C** ) A solution to the critical section problem does not have to satisfy which of the following requirements?
  - A) mutual exclusion
  - B) progress
  - C) atomicity
  - D) bounded waiting