IGERT Data Science: Graph Algorithms

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Why study Graphs and Graph algorithms

Graphs are useful in modeling complex systems with interacting components

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Graphs are useful in modeling complex systems with interacting components

We look at, understand, and interact with graphs locally: Graph algorithms compute more global properties Graphs - Mathematics and Computer Science

Graphs have a deep Mathematical theory

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"Intrinsic" to computer science - represent hard problems, program states and execution, search structures, etc.

Network Science

Network Science :: Graphs

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Network Science

Network Science:: Graphs

Statistical Mechanics :: Mechanics

Examples of "real networks"

- ▶ The internet
- ▶ The WWW
- ► The Power Grid
- ► Collaboration network
- **▶** Citations

An indispensable tool: Visualization



 ${\bf Courtesy: The\ OPTE\ project}$

Goals of Network Science

Understand and infer "macro" properties of networks

- ▶ Network vulnerability
- ▶ How things spread : information, epidemics etc.
- ▶ Identifying "important" nodes in networks
- **..**

Graph theory and Algorithms

Page rank : Google's secret formula

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Depends on computing over the WWW network - EFFICIENTLY

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Graph algorithms: indispensable to compute statistics on large networks

Fundamental graph problems

- ► Traversals
- ► Connectivity
- ► Shortest paths
- ▶ Identifying subgraphs with specific structure
- ► Characterizing graphs
- ▶ Optimization problems
- **...**

The Shortest Path problem

Given graph and a vertex s find shortest paths from s to all other vertices.

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- ▶ Map routing, robot navigation, urban traffic planning
- ▶ Network routing protocols (OSPF, BGP, RIP)
- **...**

Traversals: Depth First Search (DFS)

Visit all vertices of the graph

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Visit all vertices of the graph

```
Procedure \mathrm{DFS}(G,v)

begin

| label v as discovered
| forall the edges from v to w in G.adjacentEdges(v) do
| if w is not labeled as discovered then
| recursively call \mathrm{DFS}(G,w)
| end
| end
| end
```

Non recursive DFS

```
Procedure DFS-iterative (G, s)
begin
    S \leftarrow \text{empty stack}
    S.\operatorname{push}(s)
    while S is not empty do
        v \leftarrow S.pop()
        {f if}\ v\ is\ not\ labeled\ as\ discovered\ {f then}
             label v as discovered
             forall the edges from v to w in G.adjacentEdges(v) do
                 S.\operatorname{push}(w)
             end
        end
    end
```

Can compute shortest paths in unweighted graphs!

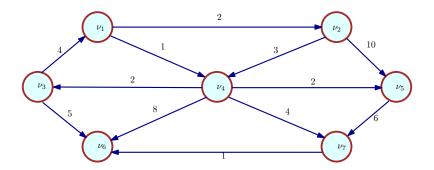
```
Procedure BFS(G, s)
begin
   s.\mathrm{dist} \leftarrow 0
   for currDist \leftarrow 0 to |V| do
       for all vertices v do
           if v not labeled as discovered and v.dist == currDist then
              Label v as discovered
              forall the edges from v to w in G.adjacentEdges(v) do
                  Handle w
              end
           end
       end
   end
end
```

Runtime : $O(|V|^2)$

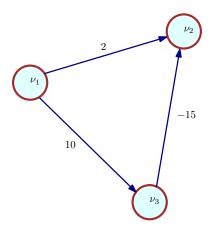
Runtime : $O(|V|^2)$

Can achieve O(|V|+|E|) using a smarter implementation

Edge weighted graphs



Negative edge weights



Main structural properties of Shortest Paths

▶ Prefixes of shortest paths are shortest paths

ightharpoonup Shortest paths from s - make up a tree

Advancing the frontier

▶ At any time we know the distance of some vertices

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▶ Relaxation of an edge $(v, w) : d(w) = \min(d(w), d(v) + c_{vw})$

Djikstra algorithm : Generalize BFS

▶ Store d_v , v.isKnown, v.path

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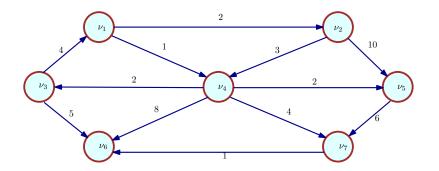
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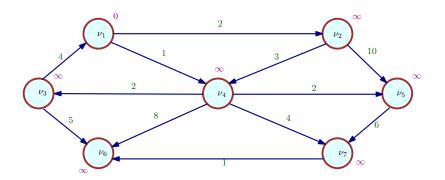
▶ Mark it known

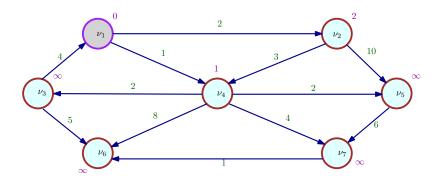
▶ Relax all edges outgoing from it

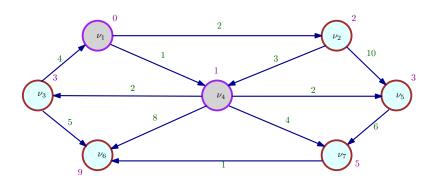
▶ Repeat until all vertices are known

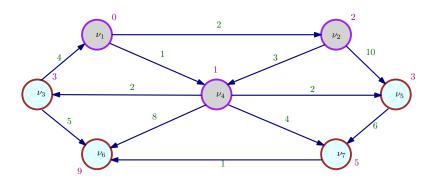
Example of Djikstra in action

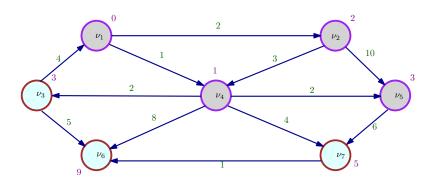


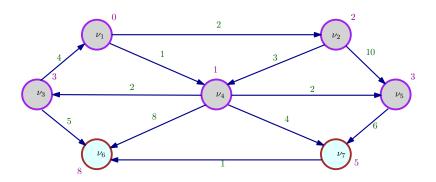


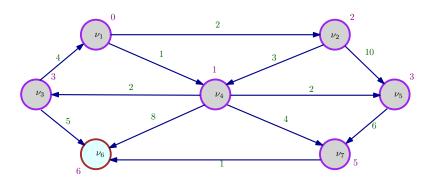


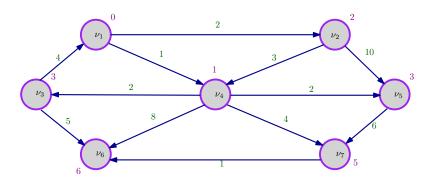












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Naive implementation : $O(|E| + |V^2|) = O(|V|^2)$

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Using a heap : $O(|E| \log |V|)$

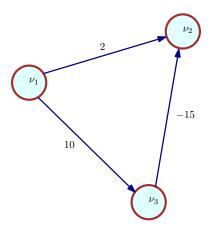
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Naive implementation :
$$O(|E| + |V^2|) = O(|V|^2)$$

Using a heap :
$$O(|E| \log |V|)$$

Using Fibonacci heap : $O(|E| + |V| \log |V|)$

Negative edge weights!



The Bellman Ford algorithm

```
Procedure Bellman-Ford(G, s) begin s.\operatorname{dist} \leftarrow 0; for i \leftarrow 1 to |V| do | Relax each edge end end
```

Runtime : O(|E||V|)

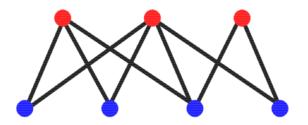
All-Pairs Shortest Path

- Can run |V| Djikstra's - $O(|E||V|\log|V|)$

▶ Floyd Warshall - $O(|V|^3)$

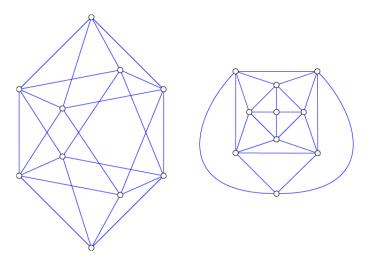
Graphs with special structure

Bipartite : $V = V_1 \cup V_2$, edges only between V_1, V_2



Graphs with special structure

Planar: Can be drawn in the plane without crossings - sparse



Back to networks

How can we model real networks?

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Doing mathematical analysis on the abstracted models helps us understand networks

Models should fit

Models should be guided by observed properties

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Test their computed properties against observed ones

Some common network models

▶ Random graph models - why?

▶ Random (n, m) graphs

► Erdös-Rényi model

▶ Albert-Barbási model - preferential attachment

Watts Strogatz model - small world phenomena

Summary statistics and Performance metrics

▶ Diameter and mean path length

▶ Clustering, average clustering, global clustering

► Degree distribution

► Centrality and its distribution

The definitions

Clustering coefficient: $C_i = 2L_i/k_i(k_i-1)$

Global clustering coefficient : $\frac{3 \times \text{Number of triangles}}{\text{Number of connected triples}}$

Betweenness Centrality

$$g(v) = \textbf{Normalizing factor} \times \sum_{s \neq v \neq t} \frac{\sigma_{st}(v)}{\sigma_{st}}$$

Why so many network models are needed?

► Erdös-Rényi model does not capture gloabl clustering

▶ Watts-Strogatz model captures small world properties

▶ Networks *grow* - how can we accommodate that?

Graphs algorithms - the message

Network Science is important

Graphs algorithms - the message

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Need approximate but very fast algorithms for basic tasks