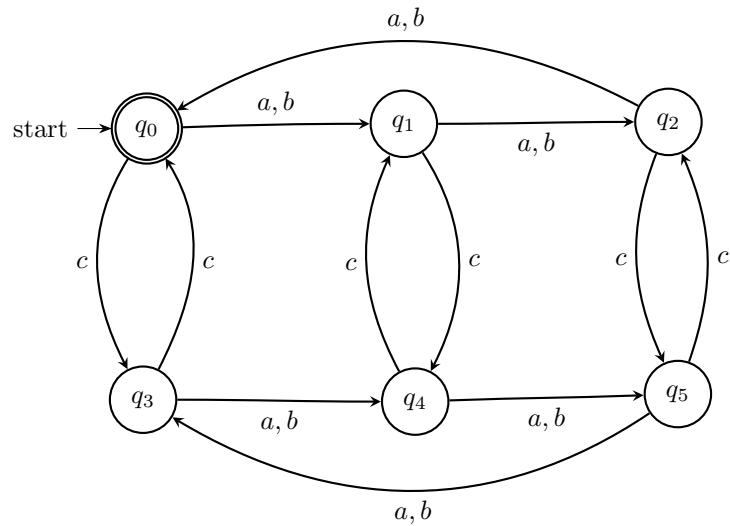


2AC3 Assignment 1

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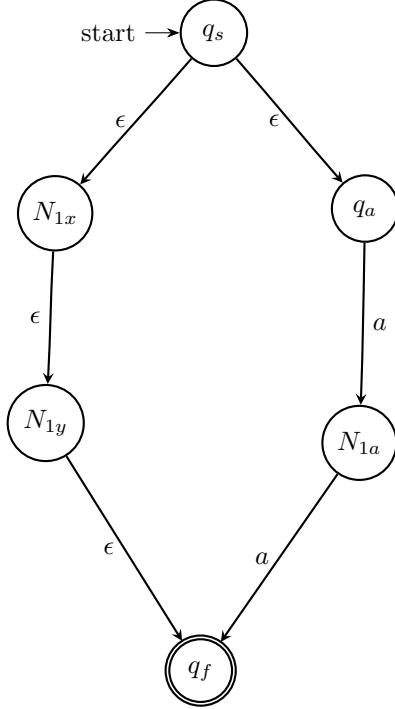
February 9, 2026

1.



q_0 acts as both the initial state and the final state. This allows ϵ to hold trivially. Furthermore, there are two rows of states that act similarly. The first row of states (q_0, q_1, q_2) represents an even number of c . The second row of states (q_3, q_4, q_5) represents an odd number of c . Both rows of states connect by cycling through a and b until the sum is divisible by 3. If the number of c is odd, the only way to get to the final state is to increase the number of c by one, causing there to be an even amount. DFA accounts for everywhere that the c may be placed.

2.



N_2 is defined as $N_2 = (Q_2, \Sigma, \Delta_2, q_s, q_f)$.

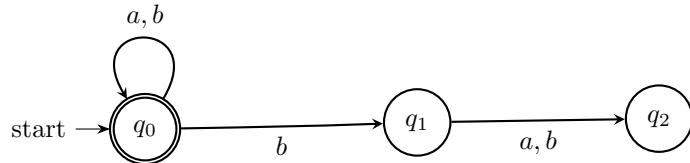
The NFA works, as it starts at state q_s , and can transition to branch xy , or branch axa on ϵ . In branch xy , $x \in L(N_1)$ is validated through N_{1x} . N_2 is able to transition to N_{1y} on ϵ , which validates $y \in L(N_1)$. Finally, N_2 reaches the final state q_f on ϵ . In branch axa , N_2 transitions to N_{1a} on a in order to fulfill the condition of axa . N_{1a} validates $x \in L(N_1)$. Then, N_2 transitions to the final state q_f on a , completing the condition of axa .

3.

N is defined as $N = (Q, \Sigma, \Delta, q_s, q_f)$.

N functions by starting at q_s . N loops through q_s for every $\delta \in \Sigma$ in order to reach v . q_s then ϵ transitions to M_1 in order to validate $v \in L(M_1)$. After validating, N ϵ transitions to q_1 , which loops for every $\delta \in \Sigma$ of w in order to reach y . q_1 then ϵ transitions to M_2 in order to validate $y \in L(M_2)$. After validating, N ϵ transitions to the final state q_f . There is no need to loop through $\delta \in \Sigma$ for z , as $z \in \Sigma^*$ holds trivially.

4.a



4.b

