

Chapter 6

Assembler

These slides support chapter 6 of the book

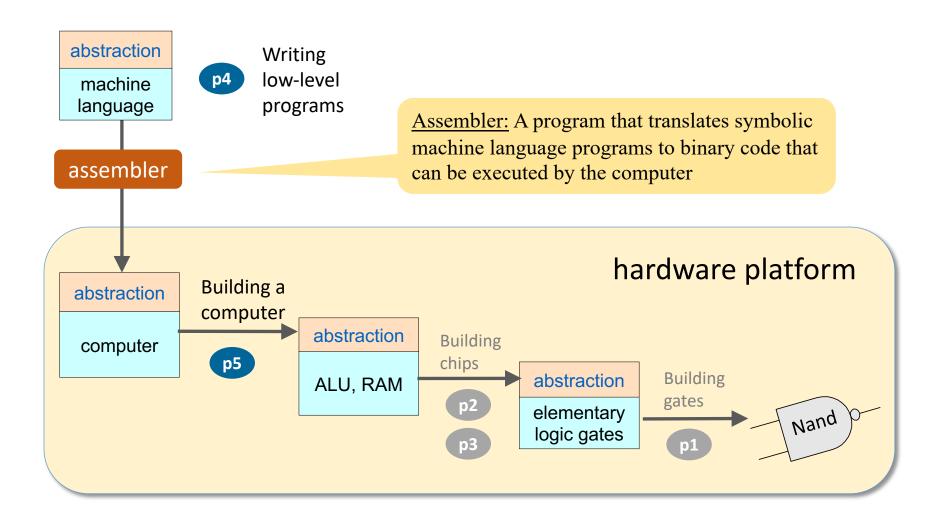
The Elements of Computing Systems

(1st and 2nd editions)

By Noam Nisan and Shimon Schocken

MIT Press

Nand to Tetris Roadmap (Part I: Hardware)



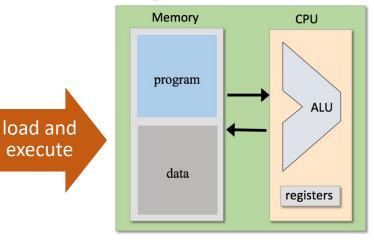
The assembler

Symbolic low-level program

```
// Program: Sum1ToN (R0 represents N)
// Computes R1 = 1 + 2 + 3 + ... + R0
// Usage: put a value >= 1 in R0
   // i = 1
   @i
   M=1
   // sum = 0
                            assembler
   @sum
   M=0
(LOOP)
   // if(i > R0) goto STOP
   @i
   D=M
   @R0
   D=D-M
   @STOP
   D; JGT
   // sum = sum + i
   @sum
   D=M
   @i
   D=D+M
   @sum
   M=D
   // i = i + 1
   @i
   M=M+1
   // goto LOOP
   @LOOP
   0;JMP
   . . .
```

Binary code

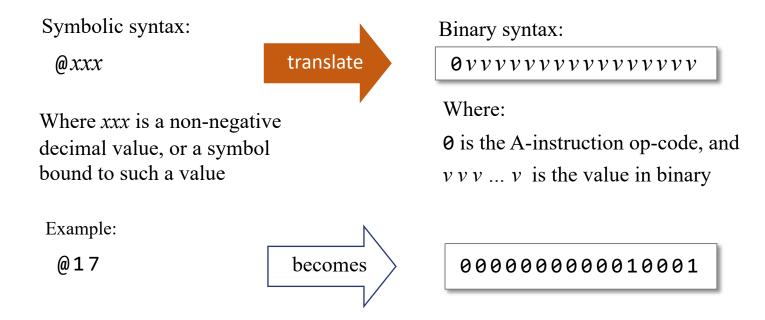
Computer



Why write an assembler?

- Because it is the "linchpin" that connects the hardware platform and the software hierarchy that sits on top of it
- Because it provides a simple example of key software engineering techniques (parsing, code generation, symbol tables, ...)

Translating A-instructions



<u>Implementation</u>

Simple: Translate the decimal value into its 16-bit representation.

Translating C-instructions

Symbolic syntax: dest = comp; jump

Binary syntax: 1 1 1 a c c c c c c d d d j j j

| CON | np | С | c | \mathcal{C} | c | c | С |
|-----|-----|---|---|---------------|---|---|---|
| 0 | | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | | 1 | 1 | 1 | 1 | 1 | 1 |
| -1 | | 1 | 1 | 1 | 0 | 1 | 0 |
| D | | 0 | 0 | 1 | 1 | 0 | 0 |
| Α | М | 1 | 1 | 0 | 0 | 0 | 0 |
| !D | | 0 | 0 | 1 | 1 | 0 | 1 |
| ! A | !M | 1 | 1 | 0 | 0 | 0 | 1 |
| -D | | 0 | 0 | 1 | 1 | 1 | 1 |
| -A | -M | 1 | 1 | 0 | 0 | 1 | 1 |
| D+1 | | 0 | 1 | 1 | 1 | 1 | 1 |
| A+1 | M+1 | 1 | 1 | 0 | 1 | 1 | 1 |
| D-1 | | 0 | 0 | 1 | 1 | 1 | 0 |
| A-1 | M-1 | 1 | 1 | 0 | 0 | 1 | 0 |
| D+A | D+M | 0 | 0 | 0 | 0 | 1 | 0 |
| D-A | D-M | 0 | 1 | 0 | 0 | 1 | 1 |
| A-D | M-D | 0 | 0 | 0 | 1 | 1 | 1 |
| D&A | D&M | 0 | 0 | 0 | 0 | 0 | 0 |
| D A | D M | 0 | 1 | 0 | 1 | 0 | 1 |

| aesi | а | а | а | effect: the value is stored in: | |
|------|---|---|---|------------------------------------|--|
| null | 0 | 0 | 0 | the value is not stored | |
| М | 0 | 0 | 1 | RAM[A] | |
| D | 0 | 1 | 0 | D register | |
| DM | 0 | 1 | 1 | D register and RAM[A] | |
| Α | 1 | 0 | 0 | A register | |
| AM | 1 | 0 | 1 | A register and RAM[A] | |
| AD | 1 | 1 | 0 | A register and D register | |
| ADM | 1 | 1 | 1 | A register, D register, and RAM[A] | |

| jump | j | j | j | effect: |
|------|---|---|---|-------------------------|
| null | 0 | 0 | 0 | no jump |
| JGT | 0 | 0 | 1 | if $comp > 0$ jump |
| JEQ | 0 | 1 | 0 | if $comp = 0$ jump |
| JGE | 0 | 1 | 1 | if $comp \ge 0$ jump |
| JLT | 1 | 0 | 0 | if <i>comp</i> < 0 jump |
| JNE | 1 | 0 | 1 | if $comp \neq 0$ jump |
| JLE | 1 | 1 | 0 | if $comp \le 0$ jump |
| JMP | 1 | 1 | 1 | Unconditional jump |

a == 0 a == 1

<u>Implementation</u>: Simple: Translate each field of the symbolic instruction (*dest*, *comp*, *jump*) into its binary code, and assemble the codes into a 16-bit instruction.

Translating C-instructions

Symbolic syntax: dest = comp; jump

Binary syntax: 1 1 1 a c c c c c c d d d j j j

| con | np | С | c | c | С | c | C |
|-----|-----|---|---|---|---|---|---|
| 0 | | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | | 1 | 1 | 1 | 1 | 1 | 1 |
| -1 | | 1 | 1 | 1 | 0 | 1 | 0 |
| D | | 0 | 0 | 1 | 1 | 0 | 0 |
| Α | М | 1 | 1 | 0 | 0 | 0 | 0 |
| !D | | 0 | 0 | 1 | 1 | 0 | 1 |
| !A | !M | 1 | 1 | 0 | 0 | 0 | 1 |
| -D | | 0 | 0 | 1 | 1 | 1 | 1 |
| -A | -M | 1 | 1 | 0 | 0 | 1 | 1 |
| D+1 | | 0 | 1 | 1 | 1 | 1 | 1 |
| A+1 | M+1 | 1 | 1 | 0 | 1 | 1 | 1 |
| D-1 | | 0 | 0 | 1 | 1 | 1 | 0 |
| A-1 | M-1 | 1 | 1 | 0 | 0 | 1 | 0 |
| D+A | D+M | 0 | 0 | 0 | 0 | 1 | 0 |
| D-A | D-M | 0 | 1 | 0 | 0 | 1 | 1 |
| A-D | M-D | 0 | 0 | 0 | 1 | 1 | 1 |
| D&A | D&M | 0 | 0 | 0 | 0 | 0 | 0 |
| D A | D M | 0 | 1 | 0 | 1 | 0 | 1 |

| dest | d | d | d | effect: the value is stored in: |
|------|---|---|---|------------------------------------|
| null | 0 | 0 | 0 | the value is not stored |
| М | 0 | 0 | 1 | RAM[A] |
| D | 0 | 1 | 0 | D register |
| DM | 0 | 1 | 1 | D register and RAM[A] |
| Α | 1 | 0 | 0 | A register |
| AM | 1 | 0 | 1 | A register and RAM[A] |
| AD | 1 | 1 | 0 | A register and D register |
| ADM | 1 | 1 | 1 | A register, D register, and RAM[A] |

| jump | j | j | j | effect: |
|------|---|---|---|-------------------------|
| null | 0 | 0 | 0 | no jump |
| JGT | 0 | 0 | 1 | if $comp > 0$ jump |
| JEQ | 0 | 1 | 0 | if $comp = 0$ jump |
| JGE | 0 | 1 | 1 | if $comp \ge 0$ jump |
| JLT | 1 | 0 | 0 | if <i>comp</i> < 0 jump |
| JNE | 1 | 0 | 1 | if $comp \neq 0$ jump |
| JLE | 1 | 1 | 0 | if $comp \le 0$ jump |
| JMP | 1 | 1 | 1 | Unconditional jump |

a == 0 a == 1

Binary:

Example: D = D+1; JLE

Translating C-instructions

Symbolic syntax: dest = comp; jump

Binary syntax: 1 1 1 a c c c c c d d d j j j

ADM

| con | np | С | c | c | С | c | С |
|-----|-----|---|---|---|---|---|---|
| 0 | | 1 | 0 | 1 | 0 | 1 | 0 |
| 1 | | 1 | 1 | 1 | 1 | 1 | 1 |
| -1 | | 1 | 1 | 1 | 0 | 1 | 0 |
| D | | 0 | 0 | 1 | 1 | 0 | 0 |
| Α | М | 1 | 1 | 0 | 0 | 0 | 0 |
| !D | | 0 | 0 | 1 | 1 | 0 | 1 |
| !A | !M | 1 | 1 | 0 | 0 | 0 | 1 |
| -D | | 0 | 0 | 1 | 1 | 1 | 1 |
| -A | -M | 1 | 1 | 0 | 0 | 1 | 1 |
| D+1 | | 0 | 1 | 1 | 1 | 1 | 1 |
| A+1 | M+1 | 1 | 1 | 0 | 1 | 1 | 1 |
| D-1 | | 0 | 0 | 1 | 1 | 1 | 0 |
| A-1 | M-1 | 1 | 1 | 0 | 0 | 1 | 0 |
| D+A | D+M | 0 | 0 | 0 | 0 | 1 | 0 |
| D-A | D-M | 0 | 1 | 0 | 0 | 1 | 1 |
| A-D | M-D | 0 | 0 | 0 | 1 | 1 | 1 |
| D&A | D&M | 0 | 0 | 0 | 0 | 0 | 0 |
| D A | D M | 0 | 1 | 0 | 1 | 0 | 1 |

| cicsi | | Ci | · · | cheet, the value is stored in. |
|-------|---|----|-----|--------------------------------|
| null | 0 | 0 | 0 | the value is not stored |
| М | 0 | 0 | 1 | RAM[A] |
| D | 0 | 1 | 0 | D register |
| DM | 0 | 1 | 1 | D register and RAM[A] |
| Α | 1 | 0 | 0 | A register |
| AM | 1 | 0 | 1 | A register and RAM[A] |

1 1 0 A register and D register

A register, D register, and RAM[A]

dest d d d effect: the value is stored in:

| jump | j | j | j | effect: |
|----------|---|---|---|---------|
| <i>J</i> | , | J | J | |

| <i>J</i> | | , | J | |
|----------|---|---|---|-----------------------|
| null | 0 | 0 | 0 | no jump |
| JGT | 0 | 0 | 1 | if $comp > 0$ jump |
| JEQ | 0 | 1 | 0 | if $comp = 0$ jump |
| JGE | 0 | 1 | 1 | if $comp \ge 0$ jump |
| JLT | 1 | 0 | 0 | if $comp < 0$ jump |
| JNE | 1 | 0 | 1 | if $comp \neq 0$ jump |
| JLE | 1 | 1 | 0 | if $comp \le 0$ jump |
| JMP | 1 | 1 | 1 | Unconditional jump |

a == 0 a == 1

Binary:

Example: A = -1



Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @i
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @i
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @i
    D=M
    @sum
    M=D+M
    // i++
    @i
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Translate

Need to handle:

- White space
- Instructions
- Symbols

We'll start with programs that have no symbols, and handle symbols later

Binary code

Binary code Symbolic code // Computes R1=1 + ... + R0 // i = 1Translate @16 M=1// sum = 0@17 Need to handle: M=0 // if i>R0 goto STOP Ignore it • White space @16 D=M @0 Instructions D=D-M @18 D; JGT • Symbols (later) // sum += i @16 D=M White space: @17 M=D+M• Empty lines, // i++ @16 • Comments, M=M+1• Indentation @4 0;JMP @17 D=Mno symbols

Symbolic code

@16 M=1@17 M=0 @16 D=M @0 D=D-M@18 D; JGT @16 D=M @17 M=D+M@16 M=M+1@4 0;JMP @17 D=M. . .



Need to handle:

- White space
- Instructions
- Symbols (later)

Binary code

Translate,

one by one

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @<u>i</u>
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D;JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @<u>i</u>
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```



Need to handle:

- White space
- Instructions
- Symbols

Original program, with symbols

Binary code

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @<u>i</u>
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D;JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @<u>i</u>
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Symbols

- Predefined symbols
- Label symbols
- Variable symbols

Original program, with symbols

Symbolic code

```
// Computes R1=1 + ... + R0
   // i = 1
    @i
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @i
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @i
    D=M
   @sum
    M=D+M
    // i++
    @i
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

The Hack language features 23 *predefined symbols*:

| <u>symbol</u> | <u>value</u> |
|---------------|--------------|
| RØ | 0 |
| R1 | 1 |
| R2 | 2 |
| • • • | |
| R15 | 15 |
| SCREEN | 16384 |
| KBD | 24576 |
| SP | 0 |
| LCL | 1 |
| ARG | 2 |
| THIS | 3 |
| THAT | 4 |
| | |

<u>Translating</u> @preDefinedSymbol:

Replace preDefinedSymbol with its value

Example: @R15 000000000001111

Symbolic code

```
// Computes R1=1 + ... + R0
        // i = 1
 0
        @i
        M=1
        // sum = 0
        @sum
        M=0
    (LOOP)
        // if i>R0 goto STOP
        @i
        D=M
        @R0
        D=D-M
        @STOP
        D; JGT
        // sum += i
        @i
10
11
        D=M
12
        @sum
13
        M=D+M
        // i++
        @i
14
15
        M=M+1
16
        @LOOP
17
        0;JMP
    (STOP)
18
        @sum
19
        D=M
```

Label symbols

- Used to label destinations of goto instructions
- Declared by the pseudo-instruction (label)
- The (*label*) directive defines the symbol *label* to refer to the memory location holding the next instruction in the program,
- Which corresponds to the instruction's *line number*

```
Example: <u>symbol</u> <u>value</u>

LOOP 4

STOP 18
```

<u>Translating</u> @labelSymbol:

Replace labelSymbol with its value

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @<u>i</u>
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @i
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Variable symbols

- Any symbol xxx which is neither predefined, nor defined elsewhere using an (xxx) label declaration, is treated as a variable
- Each variable is bound to a running memory address, starting at 16

```
Example: <u>symbol</u> <u>value</u>

i 16

sum 17
```

<u>Translating</u> @variableSymbol:

- 1. If *variableSymbol* is seen for the first time, bind to it a *value*, from 16 onward Else, it has a *value*
- 2. Replace *variableSymbol* with its *value*.

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @<u>i</u>
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @i
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Symbol table

| symbol | value |
|--------|-------|
| RØ | 0 |
| R1 | 1 |
| R2 | 2 |
| • • • | • • • |
| R15 | 15 |
| SCREEN | 16384 |
| KBD | 24576 |
| SP | 0 |
| LCL | 1 |
| ARG | 2 |
| THIS | 3 |
| THAT | 4 |
| LOOP | 4 |
| STOP | 18 |
| i | 16 |
| sum | 17 |

A data structure that the assembler creates and uses during the program translation

Contains the predefined symbols, label symbols, variable symbols, And their bindings.

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @<u>i</u>
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @<u>i</u>
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Symbol table

| symbol | value |
|--------|-------|
| RØ | 0 |
| R1 | 1 |
| R2 | 2 |
| • • • | • • • |
| R15 | 15 |
| SCREEN | 16384 |
| KBD | 24576 |
| SP | 0 |
| LCL | 1 |
| ARG | 2 |
| THIS | 3 |
| THAT | 4 |
| | |

A data structure that the assembler creates and uses during the program translation

Initialization:

Creates the table and adds all the predefined symbols

Symbolic code

```
// Computes R1=1 + ... + R0
         // i = 1
 0
         M=1
         // sum = 0
         @sum
         M=0
    (LOOP)
        // if i>R0 goto STOP
         @<u>i</u>
 4
         D=M
         @R0
         D=D-M
        @STOP
        D;JGT
         // sum += i
         @i
10
11
         D=M
12
         @sum
        M=D+M
13
         // i++
         @i
14
15
         M=M+1
16
        @LOOP
         0;JMP
17
    (STOP)
18
         @sum
19
         D=M
```

Symbol table

| symbol | value |
|--------|-------|
| RØ | 0 |
| R1 | 1 |
| R2 | 2 |
| ••• | |
| R15 | 15 |
| SCREEN | 16384 |
| KBD | 24576 |
| SP | 0 |
| LCL | 1 |
| ARG | 2 |
| THIS | 3 |
| THAT | 4 |
| LOOP | 4 |
| STOP | 18 |
| | |

A data structure that the assembler creates and uses during the program translation

Initialization:

Creates the table and adds all the predefined symbols

First pass: Counts lines and adds the label symbols

Symbolic code

```
// Computes R1=1 + ... + R0
    // i = 1
    @i
    M=1
    // sum = 0
    @sum
    M=0
(LOOP)
    // if i>R0 goto STOP
    @<u>i</u>
    D=M
    @R0
    D=D-M
    @STOP
    D; JGT
    // sum += i
    @<u>i</u>
    D=M
    @sum
    M=D+M
    // i++
    @i
    M=M+1
    @LOOP
    0;JMP
(STOP)
    @sum
    D=M
```

Symbol table

| symbol | value | |
|--------|-------|--|
| RØ | 0 | |
| R1 | 1 | |
| R2 | 2 | |
| • • • | • • • | |
| R15 | 15 | |
| SCREEN | 16384 | |
| KBD | 24576 | |
| SP | 0 | |
| LCL | 1 | |
| ARG | 2 | |
| THIS | 3 | |
| THAT | 4 | |
| LOOP | 4 | |
| STOP | 18 | |
| i | 16 | |
| sum | 17 | |
| | | |

A data structure that the assembler creates and uses during the program translation

Initialization:

Creates the table and adds all the predefined symbols

First pass: Counts lines and adds the label symbols

Second pass: Generates binary code; In the process, adds the variable symbols

(details, soon)

Assembler: Usage

<u>Input</u> (*Prog.* asm): a text file containing a sequence of lines, each being a comment, an A instruction, or a C-instruction

Output (*Prog.*hack): a text file containing a sequence of lines, each being a string of sixteen 0 and 1 characters

```
// Computes R1=1 + ... + R0
                                          0000000000010000
    // i = 1
                                          11101111111001000
                                          0000000000010001
                                          1110101010001000
    M=1
    // sum = 0
                                          0000000000010000
                       Assembler
                                          1111110000010000
    @sum
    M=0
                                          0000000000000000
(LOOP)
                                          1111010011010000
    // if i>R0 goto STOP
                                          000000000010010
                                          1110001100000001
    D=M
                                          000000000010000
                                          1111110000010000
    D=D-M
                                          000000000010001
```

<u>Usage</u>: (if the assembler is implemented in Java)

\$ java HackAssembler Prog.asm

Action: Creates a *Prog.* hack file, containing the translated Hack program

Assembler: Algorithm

Initialize:

Opens the input file (*Prog.*asm), and gets ready to process it

Constructs a symbol table,
and adds to it all the predefined symbols

First pass:

Reads the program lines, one by one, focusing only on (*label*) declarations. Adds the found labels to the symbol table

Second pass (main loop):

(starts again from the beginning of the file)

While there are more lines to process:

Gets the next instruction, and parses it

If the instruction is @symbol

If symbol is not in the symbol table, adds it

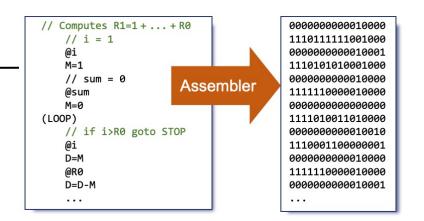
Translates the *symbol* to its binary value

If the instruction is *dest=comp*; *jump*

Translates each of the three fields into its binary value

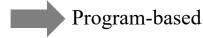
Assembles the binary values into a string of sixteen 0's and 1's

Writes the string to the output file.

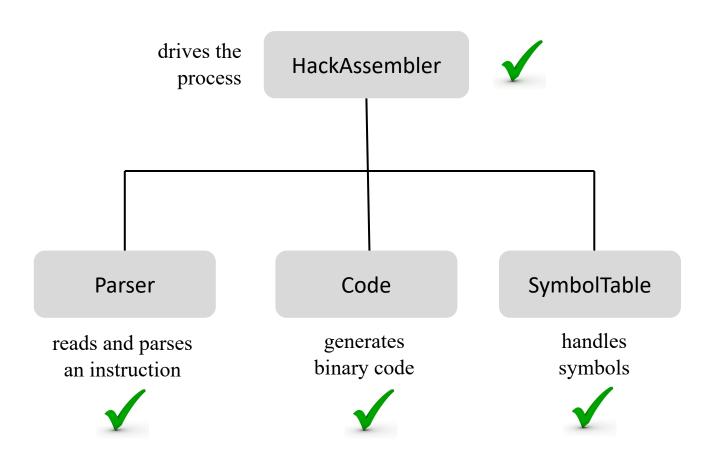


Assembler implementation options

Manual



HackAssembler: Drives the translation process



HackAssembler

Initialize:

Opens the input file (*Prog.* asm) and gets ready to process it Constructs a symbol table, and adds to it all the predefined symbols

First pass:

Reads the program lines, one by one focusing only on (*label*) declarations. Adds the found labels to the symbol table

Second pass (main loop):

(starts again from the beginning of the file)

While there are more lines to process:

Gets the next instruction, and parses it

If the instruction is @symbol

If symbol is not in the symbol table, adds it

Translates the *symbol* into its binary value

If the instruction is dest = comp; jump

Translates each of the three fields into its binary value

Assembles the binary values into a string of sixteen 0's and 1's

Writes the string to the output file.

The HackAssembler implements the assembly algorithm, using the services of:

- Parser
- Code
- SymbolTable

Parser API

Routines:

- Constructor / initializer: Creates a Parser and opens the source text file
- Getting the current instruction:

```
hasMoreLines(): Checks if there is more work to do (boolean)
```

advance(): Gets the next instruction and makes it the *current instruction* (string)

• Parsing the *current instruction*:

current instruction

Examples:

Parser API

Routines:

- Constructor / initializer: Creates a Parser and opens the source text file
- Getting the current instruction:

```
hasMoreLines(): Checks if there is more work to do
```

advance(): Gets the next instruction and makes it the current instruction

• Parsing the *current instruction*:

Parser API

Routines:

- Constructor / initializer: Creates a Parser and opens the source text file
- Getting the current instruction:

hasMoreLines(): Checks if there is more work to do

advance(): Gets the next instruction and makes it the current instruction

• Parsing the *current instruction*:

instructionType(): Returns the instruction type

symbol(): Returns the instruction's symbol (string)

dest(): Returns the instruction's *dest* field (string)

comp(): Returns the instruction's comp field (string)

jump(): Returns the instruction's jump field (string)

Used if the current instruction is dest = comp; jump

current instruction

Examples:

$$M=-1$$

Code API

Deals only with C-instructions: dest = comp; jump

Routines:

a == 0

dest(string): Returns the binary representation of the parsed dest field (string)

comp(string): Returns the binary representation of the parsed *comp* field (string)

jump(string): Returns the binary representation of the parsed jump field (string)

According to the language specification:

| cor | пр | c | c | c | c | c | c | |
|-----|-----|---|---|---|---|---|---|--|
| 0 | | 1 | 0 | 1 | 0 | 1 | 0 | |
| 1 | | 1 | 1 | 1 | 1 | 1 | 1 | |
| -1 | | 1 | 1 | 1 | 0 | 1 | 0 | |
| D | | 0 | 0 | 1 | 1 | 0 | 0 | |
| Α | М | 1 | 1 | 0 | 0 | 0 | 0 | |
| !D | | 0 | 0 | 1 | 1 | 0 | 1 | |
| !A | !M | 1 | 1 | 0 | 0 | 0 | 1 | |
| -D | | 0 | 0 | 1 | 1 | 1 | 1 | |
| -A | -M | 1 | 1 | 0 | 0 | 1 | 1 | |
| D+1 | | 0 | 1 | 1 | 1 | 1 | 1 | |
| A+1 | M+1 | 1 | 1 | 0 | 1 | 1 | 1 | |
| D-1 | | 0 | 0 | 1 | 1 | 1 | 0 | |
| A-1 | M-1 | 1 | 1 | 0 | 0 | 1 | 0 | |
| D+A | D+M | 0 | 0 | 0 | 0 | 1 | 0 | |
| D-A | D-M | 0 | 1 | 0 | 0 | 1 | 1 | |
| A-D | M-D | 0 | 0 | 0 | 1 | 1 | 1 | |
| D&A | D&M | 0 | 0 | 0 | 0 | 0 | 0 | |
| D A | D M | 0 | 1 | 0 | 1 | 0 | 1 | |

a == 1

| dest | d | d | d |
|------|---|---|---|
| null | 0 | 0 | 0 |
| М | 0 | 0 | 1 |
| D | 0 | 1 | 0 |
| DM | 0 | 1 | 1 |
| Α | 1 | 0 | 0 |
| AM | 1 | 0 | 1 |
| AD | 1 | 1 | 0 |
| ADM | 1 | 1 | 1 |
| | | | |

| jump | j | j | j |
|------|---|---|---|
| null | 0 | 0 | 0 |
| JGT | 0 | 0 | 1 |
| JEQ | 0 | 1 | 0 |
| JGE | 0 | 1 | 1 |
| JLT | 1 | 0 | 0 |
| JNE | 1 | 0 | 1 |
| JLE | 1 | 1 | 0 |
| JMP | 1 | 1 | 1 |

Examples:

dest("DM") returns "011"
comp("A+1") returns "0110111"
comp("D&M") returns "1000000"
jump("JNE") returns "101"

SymbolTable API

Routines

Constructor / initializer: Creates and initializes a SymbolTable

addEntry(symbol (string), address (int)): Adds <symbol, address> to the table (void)

contains(symbol (string)): Checks if symbol exists in the table (boolean)

getAddress(symbol (string)): Returns the address (int) associated with symbol

Symbol table: (example)

| symbol | address |
|--------|---------|
| RØ | 0 |
| R1 | 1 |
| R2 | 2 |
| | ••• |
| R15 | 15 |
| SCREEN | 16384 |
| KBD | 24576 |
| SP | 0 |
| LCL | 1 |
| ARG | 2 |
| THIS | 3 |
| THAT | 4 |
| LOOP | 4 |
| STOP | 18 |
| i | 16 |
| sum | 17 |

Assembler API (detailed)

Parser module:

| Routine | Arguments | Returns | Function |
|---------------------------|----------------------|------------------------------|--|
| Constructor / initializer | Input file or stream | _ | Opens the input file/stream and gets ready to parse it. |
| hasMoreLines | s — s | boolean | Are there more lines in the input? |
| advance | _ | _ | Skips over whitespace and comments, if necessary. |
| | | | Reads the next instruction from the input, and makes it the current instruction. |
| | | | This method should be called only if hasMoreLines is true. |
| | | | Initially there is no current instruction. |
| instructionType | _ | A_INSTRUCTION, | Returns the type of the current instruction: |
| | | C_INSTRUCTION, L_INSTRUCTION | A_INSTRUCTION for @xxx, where xxx is either a decimal number or a symbol. |
| | | (constants) | C_INSTRUCTION for dest=comp; jump |
| | | | L_INSTRUCTION for (xxx) , where xxx is a symbol. |
| symbol | _ | string | If the current instruction is (xxx) , returns the symbol xxx . If the current instruction is $@xxx$, returns the symbol or decimal xxx (as a string). |
| | | | Should be called only if instructionType is A_INSTRUCTION or L_INSTRUCTION. |
| dest | | string | Returns the symbolic <i>dest</i> part of the current <i>C</i> -instruction (8 possibilities). |
| | | | Should be called only if instructionType is C_INSTRUCTION. |
| comp | _ | string | Returns the symbolic <i>comp</i> part of the current <i>C</i> -instruction (28 possibilities). |
| | | 94500 | Should be called only if instructionType is C_INSTRUCTION. |
| jump | _ | string | Returns the symbolic <i>jump</i> part of the current <i>C</i> -instruction (8 possibilities). |
| | | | Should be called only if instructionType is C_INSTRUCTION. |

Assembler API (detailed)

Code module:

| Routine | Arguments | Returns | Function |
|---------|-----------|---------------------|--|
| dest | string | 3 bits, as a string | Returns the binary code of the <i>dest</i> mnemonic. |
| comp | string | 7 bits, as a string | Returns the binary code of the <i>comp</i> mnemonic. |
| jump | string | 3 bits, as a string | Returns the binary code of the <i>jump</i> mnemonic. |

SymbolTable module:

| Routine | Arguments | Returns | Function |
|-------------|-----------------------------------|---------|---|
| Constructor | _ | _ | Creates a new empty symbol table. |
| addEntry | symbol (string), address (int) | _ | Adds < symbol, address > to the table. |
| contains | symbol (string) | boolean | Does the symbol table contain the given symbol? |
| getAddress | symbol (string) | int | Returns the address associated with the symbol. |

HackAssembler module:

We propose no API; Implement as you see fit.

Developing a Hack Assembler

Contract

- Develop a program that translates symbolic Hack programs into binary Hack instructions
- The source program (input) is supplied as a text file named *Prog.* asm
- The generated code (output) is written into a text file named *Prog*.hack
- Assumption: *Prog.* asm is error-free

<u>Usage</u> (if the assembler is implemented in Java):

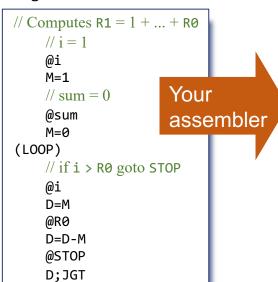
\$ java HackAssembler Prog.asm

Staged development plan

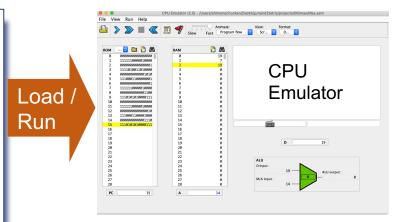
- 1. Develop a basic assembler that translates programs that have no symbols
- 2. Develop an ability to handle symbols
- 3. Morph the basic assembler into an assembler that translates any program

Testing

Prog.asm



Prog.hack



Or use a supplied test script that loads Prog. hack into the CPU emulator and tests it using pre-defined testing scenarios

Test programs

- Add.asm
- Max.asm

- MaxL.asm
- Rect.asm
- RectL.asm
- Pong.asm
- PongL.asm

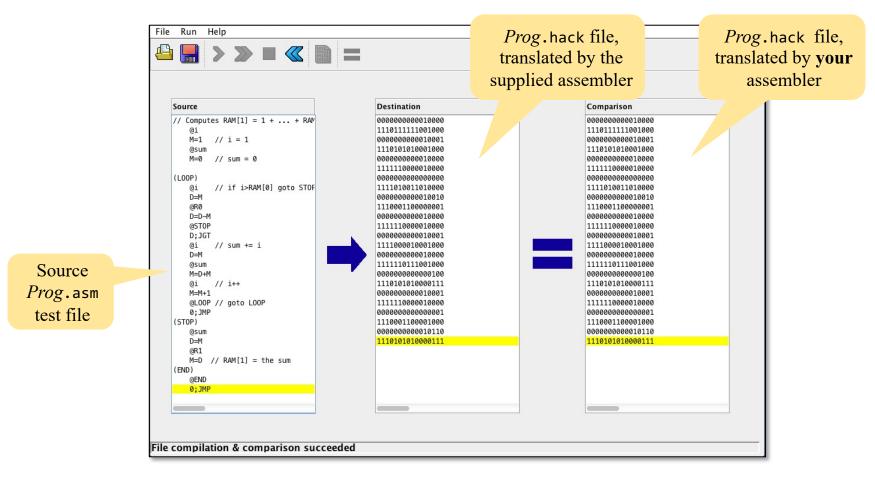
(with symbols)

(same programs, without symbols, for unit-testing the basic assembler)

Testing option II: Using the hardware simulator

- 1. Use your assembler to translate *Prog.* asm, generating the executable file *Prog.* hack
- 2. Put the *Prog.* hack file in a folder containing the chips that you developed in project 5: Computer.hdl, CPU.hdl, and Memory.hdl
- 3. Load computer.hdl into the Hardware Simulator
- 4. Load *Prog.* hack into the ROM32K chip-part
- 5. Run the clock to execute the program.

Testing option III: Using the supplied assembler



- 1. Use your assembler to translate *Prog.* asm, generating the executable file *Prog.* hack
- 2. Load *Prog.* asm into the supplied assembler, and load *Prog.* hack as a compare file
- 3. Translate *Prog.* hack, and inspect the comparison feedback messages.

Project 6

<u>Guidelines</u>: www.nand2tetris.org (projects section)

<u>Files</u>: nand2tetris/projects/06 (on your PC)

<u>Tools</u>

- The programming language in which you develop your assembler
- CPU emulator (for testing the translated programs)
- Assembler (if you plan to use it)

Guides

- <u>CPU emulator tutorial</u>
- Assembler tutorial