

**FIT5047 Second Theory Assignment**

Due date: 23rd April 2021 (23:59).

Evaluation: 100 marks = 5%.

Submission: Moodle.

# Propositional Logic - Syntax and Semantics

[ 5+5+5+5 = 20 ]

Decide whether the following propositional logic sentences are either valid, satisfiable (but not valid), or unsatisfiable. Justify your answer.

(a)  $\neg B \rightarrow (A \wedge \neg B)$  [5]

(b)  $((\neg A \vee B) \wedge (\neg C \rightarrow \neg B)) \rightarrow (A \rightarrow C)$  [5]

Convert the following propositional logic sentences into Conjunctive Normal Form (CNF):

(c)  $(C \rightarrow D) \rightarrow \neg(D \rightarrow B)$  [5]

(d)  $(A \leftrightarrow B) \rightarrow (\neg A \wedge C)$  [5]

(a)  $\neg(\neg B) \vee (A \wedge \neg B) \equiv B \vee (A \wedge \neg B) \equiv (B \vee A) \wedge (B \vee \neg B)$

A	B	$\neg B$	$(B \vee A)$	$(B \vee \neg B)$	$(B \vee A) \wedge (B \vee \neg B)$
T	T	F	T	T	T
T	F	T	T	T	T
F	T	F	T	T	T
F	F	T	F	T	F

Because  $\neg B \rightarrow (A \wedge \neg B)$  is equal to  $(B \vee A) \wedge (B \vee \neg B)$ . According to the table,  $(B \vee A) \wedge (B \vee \neg B)$  has 3 true models and 1 false model, so  $\neg B \rightarrow (A \wedge \neg B)$  is satisfiable (but not valid).

(b)  $((\neg A \vee B) \wedge (\neg C \rightarrow \neg B)) \rightarrow (A \rightarrow C) \equiv \neg((\neg A \vee B) \wedge (\neg(\neg C) \wedge \neg B)) \wedge (\neg A \wedge C) \equiv ((A \wedge \neg B) \vee (\neg C \vee B)) \wedge (\neg A \wedge C)$

A	B	C	$\neg A$	$\neg B$	$\neg C$	$(A \wedge \neg B)$	$(\neg C \vee B)$	$((A \wedge \neg B) \vee (\neg C \vee B))$	$(\neg A \wedge C)$	(a) $((A \wedge \neg B) \vee (\neg C \vee B)) \wedge (\neg A \wedge C)$
T	T	T	F	F	F	F	T	T	F	F
T	T	F	F	F	T	F	T	T	F	F
T	F	T	F	T	F	T	F	T	F	F
T	F	F	F	T	T	T	T	T	F	F
F	T	T	T	F	F	F	T	T	T	T
F	F	T	T	T	F	F	F	F	T	F
F	T	F	T	F	T	F	T	T	F	F
F	F	F	T	T	T	F	T	T	F	F

Because  $((\neg A \vee B) \wedge (\neg C \rightarrow \neg B)) \rightarrow (A \rightarrow C)$  is equal to  $((A \wedge \neg B) \vee (\neg C \vee B)) \wedge (\neg A \wedge C)$ .

According to the table, has 1 true model and 7 false models, so  $((\neg A \vee B) \wedge (\neg C \rightarrow \neg B)) \rightarrow (A \rightarrow C)$  is satisfiable (but not valid).

## 2

$$(c) (C \rightarrow D) \rightarrow \neg(D \rightarrow B) \equiv (\neg(\neg C \vee D)) \vee (\neg(\neg D \vee B)) \equiv (C \wedge \neg D) \vee (D \wedge \neg B)$$

C	D	$\neg B$	$\neg D$	$(C \wedge \neg D)$	$(D \wedge \neg B)$	$(C \wedge \neg D) \vee (D \wedge \neg B)$
T	T	T	F	F	T	T
T	T	F	F	F	F	F
T	F	T	T	T	F	T
F	T	T	F	F	T	T
F	F	T	T	F	F	F
F	T	F	F	F	F	F
F	F	F	T	F	F	F

According to the table, has 3 true models and 4 false models, so  $(C \rightarrow D) \rightarrow \neg(D \rightarrow B)$  is satisfiable (but not valid).

$$(d) (A \leftrightarrow B) \rightarrow (\neg A \wedge C) \equiv \neg((\neg A \vee B) \wedge (\neg B \vee A)) \vee (\neg A \wedge C) \equiv (A \wedge \neg B) \vee (B \wedge \neg A) \vee (\neg A \wedge C)$$

A	B	C	$\neg A$	$\neg B$	$(A \wedge \neg B)$	$(B \wedge \neg A)$	$(\neg A \wedge C)$	$(A \wedge \neg B) \vee (B \wedge \neg A) \vee (\neg A \wedge C)$
T	T	T	F	F	F	F	F	F
T	T	F	F	F	F	F	F	F
T	F	T	F	T	T	F	F	T
F	T	T	T	F	F	T	T	T
F	F	T	T	T	F	F	T	T
F	T	F	T	F	F	T	F	T
F	F	F	T	T	F	F	F	F

According to the table, has 4 true models and 3 false models, so  $(A \leftrightarrow B) \rightarrow (\neg A \wedge C)$  is satisfiable (but not valid).

### 3

## Propositional Logic - Resolution

[ 10+10 = 20 ]

Use resolution-refutation to decide whether or not the following goals can be proved, from the Knowledge Bases given below.

(a) Goal:  $C$  [10]

$$1 : (A \rightarrow B) \rightarrow B$$

$$2 : (A \rightarrow A) \rightarrow C$$

$$3 : (C \rightarrow D) \rightarrow \neg(D \rightarrow B)$$

(b) Goal:  $\neg(D \rightarrow A)$  [10]

$$1 : A \vee B \vee \neg C$$

$$2 : A \vee B \vee C$$

$$3 : D \vee \neg B$$

$$4 : \neg A \vee \neg B$$

(a) Goal:  $C$

$$1: (A \rightarrow B) \rightarrow B \equiv (A \wedge \neg B) \vee B$$

$$2: (A \rightarrow A) \rightarrow C \equiv (A \wedge \neg A) \vee C \equiv C$$

$$3: (C \rightarrow D) \rightarrow \neg(D \rightarrow B) \equiv (C \wedge \neg D) \vee (D \wedge \neg B)$$

Negated goal: 4:  $\neg C$

4 and 2: 4:  $\neg C$  2:  $C$

resolvent: NIL

This case can proved

(b) Goal:  $\neg(D \rightarrow A) \equiv D \wedge \neg A$  [10]

$$1 : A \vee B \vee \neg C$$

$$2: A \vee B \vee C$$

$$3: D \vee \neg B$$

$$4: \neg A \vee \neg B$$

Negated goal: 5:  $\neg D \vee A$

5 and 4: 5:  $\neg D \vee A$  4:  $\neg A \vee \neg B$

resolvent: 6:  $\neg D \vee \neg B$

6 and 3: 6:  $\neg D \vee \neg B$  3:  $D \vee \neg B$

**4**

resolvent: 7.  $\neg B \vee \neg B$

7 and 1: 7.  $\neg B \vee \neg B$  1.  $A \vee B \vee \neg C$

resolvent: 8:  $A \vee \neg B \vee \neg C$

8 and 2 : 8:  $A \vee \neg B \vee \neg C$  2:  $A \vee B \vee C$

resolvent: 9 :  $A$

*This case can't be proved*

## Propositional Logic - Forward Chaining

[ 5+5+5+5 = 20 ]

This question is about Horn clauses and Forward chaining.

R1:  $L \rightarrow P$

R2:  $(A \wedge B) \rightarrow L$

R3:  $P \rightarrow Q$

R4:  $L \rightarrow A$

$A$

$B$

Apply forward reasoning to the Horn clauses above in order to prove  $Q$ . After each rule application, show the values to be inserted in the columns named:

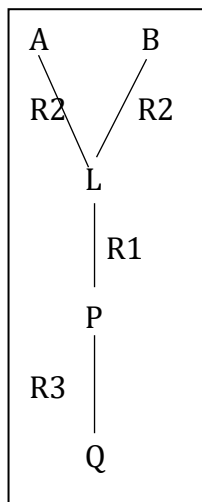
(a) AGENDA [5] (b) COUNT [5] (c) INFERRED [5]

in the table provided below, in the way seen in the lectures. Also,

(d) show the search graph resulting from the application of the rules. [5]

R1:  $\neg L \vee P$  R2:  $\neg A \vee \neg B \vee L$  R3:  $\neg P \vee Q$  R4:  $\neg L \vee A$

AGENDA	COUNT				INFERRED
	R1	R2	R3	R4	
A,B	1	2	1	1	
B	1	1	1	1	A
L	1	0	1	1	A,B
P	0	0	1	1	A,B,L
Q	0	0	0	1	A,B,L,P



## First Order Logic - Substitution and Unification

[ 5+5+5+5 = 20 ]

- (a) Calculate the composition of substitutions  $s_1s_2$  and  $s_2s_1$ , given the two substitutions  $s_1 = \{x|f(y,A), y|B, z|g(B,w)\}$  and  $s_2 = \{x|A, y|h(B,z), z|C, w|D\}$  [5]
- (b) Apply the composed substitutions  $s_1s_2$  and  $s_2s_1$  of part (a) to  $P(x, f(y), z, g(w))$ , that is, calculate predicates  $P(x, f(y), z, g(w))_{s_1s_2}$  and  $P(x, f(y), z, g(w))_{s_2s_1}$ . [5]

The following questions are about unification. Are the following literals unifiable? If so, provide the most general unifier (mgu). If not unifiable, explain why not.

(c)  $PREDICATE(x, f(7,3), g(z), z)$  and  $PREDICATE(f(3,7), y, g(w), 5)$  [5]

(d)  $PREDICATE(f(x,y))$  and  $PREDICATE(g(3,7))$  [5]

(a)  $s_1s_2 = \{x|f(y,A), y|B, z|g(B,w)\} \{x|A, y|h(B,z), z|C, w|D\} = \{x | f(B,A), y | B, z | g(B, D), w | D\}$

$s_2s_1 = \{x|A, y|h(B,z), z|C, w|D\} \{x|f(y,A), y|B, z|g(B,w)\} = \{x | A, y | h(B,C), z | C, w | D\}$

(b)  $P(x, f(y), z, g(w))_{s_1s_2} = P(f(B,A), f(B), g(B, D), g(D)) \{x | f(B,A), y | B, z | g(B, D), w | D\}$

$P(x, f(y), z, g(w))_{s_2s_1} = P(A, f(h(B,C)), C, g(D)) \{x | A, y | h(B,C), z | C, w | D\}$

(c) Unify ( $PREDICATE(x, f(7,3), g(z), z)$ ,  $PREDICATE(f(3,7), y, g(w), 5)$ )

Z1 <- Unify(P,P)

P and P are identical -> return NIL

Z2 <- Unify ((x, f(7,3), g(z), z), (f(3,7), y, g(w), 5)))

Z1 <- Unify(x, f(3,7))

Return {x| f(3,7)}

G1 <- (f(7,3), g(z), z) {x| f(3,7)} -> (f(7,3), g(z), z)

G2 <- (y, g(w), 5) {x| f(3,7)} -> (y, g(w), 5)

Z2 <- Unify ((f(7,3), g(z), z), (y, g(w), 5)))

Z1 <- Unify(f(7,3), y)

Return {y| f(7,3)}

G1 <- (g(z), z) {y| f(7,3)} -> (g(z), z)

G2 <- (g(w), 5) {y| f(7,3)} -> (g(w), 5)

Z2 <- Unify ((g(z), z), (g(w), 5)))

Z1 <- Unify((g(z), g(w)))

Z1 <- Unify(g, g)

g and g are identical return {}

Z2 <- Unify ( (z), (w))

Z1 <- Unify(z, w)

Return {z | w}

G1 <- (z) {z | w} -> (w)

G2 <- (5) {z | w} -> (5)

Z1 <- Unify ( w, 5)

Return {w|5}

Return {x| f(3,7)} {y| f(7,3)} {z | w} {w|5} = {x| f(3,7), y| f(7,3), z | 5, w|5}

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(d) Unify ( $PREDICATE(f(x, y))$ ,  $PREDICATE(g(3, 7))$ )

    Z1 <- Unify(P, P)

        P and P are identical -> return NIL

    Z2 <- Unify ( $f(x, y)$ ,  $g(3, 7)$ )

        Z1 <- Unify(f, g)

        f is a symbol and g is a symbol, but they aren't identical

        return fail

So UNIFICATION FAILS.



# 8

## First Order Logic - Resolution

[ 10+10 = 20 ]

Use resolution refutation to show the goal “G” below and show:

(a) how you construct all necessary clauses, and [10]

(b) which most general unifiers are used. [10]

$$1: P(A) \rightarrow Q(B)$$

$$2: \forall x[P(x)]$$

$$3: \forall y [\exists z [R(y,z)] \rightarrow \neg Q(y)]$$

$$G: \forall z [\neg R(B,z)]$$

$$(a) 1: P(A) \rightarrow Q(B) \equiv \neg P(A) \vee Q(B)$$

$$2: \forall x[P(x)] \equiv P(x)$$

$$3: \forall y [\exists z [R(y, z)] \rightarrow \neg Q(y)] \equiv \forall y [\neg (\exists z [R(y, z)]) \vee \neg Q(y)] \equiv \forall y [\forall z [\neg R(y, z)] \vee \neg Q(y)] \\ \equiv \neg R(y, z) \vee \neg Q(y)$$

$$G: \forall z [\neg R(B, z)] \equiv \neg R(B, z)$$

$$(b) \text{ Negated goal: } 4: R(B, z)$$

$$4 \text{ and } 3: 4: \underline{R(B, z)} \quad 3: \underline{\neg R(y, z)} \vee \neg Q(y)$$

$$\text{mgu} : \{y \mid B\}$$

$$\text{resolvent: } 5: \neg Q(B)$$

$$5 \text{ and } 1: 5: \neg Q(B) \quad 1: \neg P(A) \vee Q(B)$$

$$\text{resolvent: } 6: \neg P(A)$$

$$6 \text{ and } 2: 6: \neg P(A) \quad 2: P(x)$$

$$\text{mgu: } \{x \mid A\}$$

$$\text{resolvent: NIL}$$

$$\text{so mgu} = \{x \mid A, y \mid B\}$$