

# Genome Rearrangement

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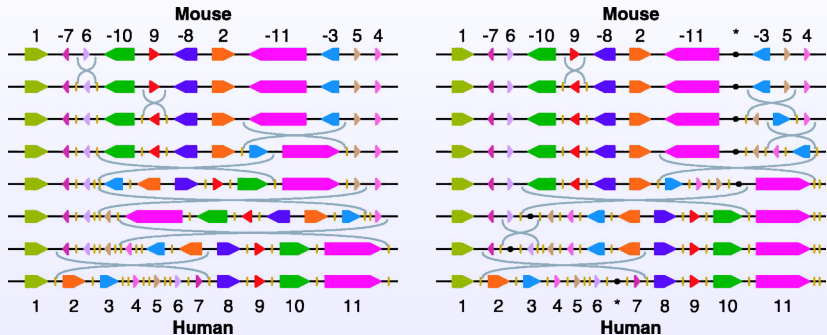
# Biological Context

Genes are sequences of nucleotides that encode the synthesis of biochemical material (be it RNA or proteins).

Biological processes can result in mutations in the genome. Some mutations are **local** and affect a single nucleotide (**substitutions**, **insertions**, **deletions**) while other can be **non-local** and affect long stretches of the sequence (**inversions**, **transpositions**, **translocations**).

Non-local mutations are rare, but accumulate over the evolutionary timescale. They thus provide a **good proxy for evolutionary distance** between species.

# Biological Context



Different species may have the same genes but in a different order (and orientation).

# Problem Statement

## Genome as a Permutation

Let the order of genes in an organism be represented as a (signed) permutation  $\pi = [\pi_1 \pi_2 \dots \pi_n]$ :

- If **gene orientation is not important**, each  $\pi_i \in \{1, \dots, n\}$  is a non-repeated **positive** integer.
- If **gene orientation is important**, each  $\pi_i$  can be **positive or negative**, and  $|\pi_i| \in \{1, \dots, n\}$  is a non-repeated positive integer.

# Problem Statement

## Reversal

### Reversal

A reversal  $\rho(i, j)$  is an operation that affects the portion of the genome between positions  $i$  and  $j$  (both inclusive), reversing this segment.

Let  $\pi = [\pi_1 \pi_2 \dots \pi_n]$ . Then, if  $\pi$  is unsigned:

$$\pi \cdot \rho(i, j) = [\pi_1 \dots \pi_{i-1} \underbrace{\pi_j \pi_{j-1} \dots \pi_{i+1} \pi_i}_{\text{reversed segment}} \pi_{j+1} \dots \pi_n]$$

If  $\pi$  is signed, the sign of each  $\pi_k$ ,  $i \leq k \leq j$ , is flipped as well.

# Problem Statement

## Minimal Reversal Distance

Let  $\pi$  and  $\pi'$  be two organisms. The **reversal distance** between them is the minimum number of reversals required to transform one of the permutations into the other.

We can assume w.l.o.g. that one of the permutations, say  $\pi'$ , is the positive identity permutation, i.e.,  $\pi'_i = i$ ,  $1 \leq i \leq n$ .

### Sorting Permutations by Reversals

Given a (signed) permutation  $\pi$ , find the shortest sequence of reversals required to transform it into the positive identity permutation.

# Unsigned Permutations

The problem of sorting an unsigned permutation by reversals is **NP-hard**.

Let us consider a naive approach to the problem, namely **PrefixSort**:

```
for each  $i \in \{1, \dots, n\}$  do  
  if  $\pi_i \neq i$  then  
     $j \leftarrow$  index for which  $\pi_j = i$   
     $\pi \leftarrow \pi \cdot \rho(i, j)$   
  endif  
endfor
```

Note that PrefixSort requires at most  $n$  reversals to produce the identity permutation.



# Unsigned Permutations

## PrefixSort in Action

Let  $\pi = [67812345]$ . PrefixSort  
requires 6 steps:

[18762345]

[12678345]

[12387645]

[12346785]

[12345876]

[12345678]

But it can be sorted in fewer  
steps:

[54321876]

[12345876]

[12345678]

# Unsigned Permutations

## A Greedy Template for Genome Rearrangement

### Greedy Genome Rearrangement

```
func GreedyRearrangement ( $\downarrow \pi$ : Permutation( $n$ )): List( $Op$ )  
variables  
     $sol, candidates$ : List( $Op$ )  
     $o, best$ :  $Op$   
begin  
     $\pi \leftarrow [0, \pi, n + 1]$  // framing the permutation  
     $sol \leftarrow \langle \rangle$   
     $candidates \leftarrow \text{getCandidates}(\pi)$  // get applicable operations  
    while  $candidates \neq \langle \rangle$  do  
         $best \leftarrow \arg \max \{ \text{getQuality}(o, \pi) \mid o \in candidates \}$  // pick best operation  
         $sol.add(best)$   
         $\pi \leftarrow best.apply(\pi)$  // apply operation  
         $candidates \leftarrow \text{getCandidates}(\pi)$  // get applicable operations  
    endwhile  
    return  $sol$   
end
```

# Unsigned Permutations

## Sorting by Breakpoints

Let  $0 \leq i \leq n$ . A pair of elements  $\pi_i, \pi_{i+1}$  are an **adjacency** if  $|\pi_i - \pi_{i+1}| = 1$ . Otherwise they are a **breakpoint**.

[0|678|12345|9]

The identity permutation is the only one that –when framed (i.e., extended with  $\pi_0 = 0$  and  $\pi_{n+1} = n + 1$ )– has no breakpoints.

It only makes sense to apply a reversal between breakpoints.

A reversal between breakpoints may remove 0, 1 or 2 of them.

**Greedy idea:** pick the reversal that maximizes the number of breakpoints removed.

# Unsigned Permutations

## Sorting by Breakpoints

$$\text{getCandidates}(\pi) = \{\rho(i+1, j) \mid i < j, \{i, j\} \subseteq \text{breakpoints}(\pi)\}$$

$$\text{getQuality}(\rho(i, j), \pi) = [|\pi_{i-1} - \pi_j| = 1] + [|\pi_{j+1} - \pi_i| = 1]$$

$\pi$	candidates	best	quality
[0 6 7 8 1 2 3 4 5 9]	$\rho(1, 3), \rho(1, 8), \rho(4, 8)$	$\rho(1, 3)$	0 (tie)
[0 8 7 6 1 2 3 4 5 9]	$\rho(1, 3), \rho(1, 8), \rho(4, 8)$	$\rho(1, 8)$	1
[0 5 4 3 2 1 6 7 8 9]	$\rho(1, 5)$	$\rho(1, 5)$	2
[0 1 2 3 4 5 6 7 8 9]	—	—	—

This algorithm is not optimal but –if an adequate tie-breaking procedure is picked– it provides a 2-approximation to the optimal.

# Complementary Bibliography



N.C. Jones, P.A. Pevzner,  
*An Introduction to Bioinformatics Algorithms*,  
MIT Press, Cambridge MA, 2004

# Image Credits

- Mouse and human genome rearrangement: Pavel Pevzner and Glenn Tesler, *PNAS*, doi: 10.1073/pnas.1330369100