

S3(7) - Exemplu

Tuesday, January 10, 2023 12:00 PM

EXEMPLU EXAMEN



```

1 class Tree inherits IO { -- empty tree
2   isEmpty() : Bool { true };
3
4   insert(k : Int) : Tree {
5     new NETree.init(k, self, self);
6   };
7
8   sum() : Int { 0 };
9 };
10
11
12 class NETree inherits Tree { -- non-empty tree
13   key : Int;
14   left : Tree;
15   right : Tree;
16
17   init(k : Int, l : Tree, r : Tree) : Tree {{
18     key <- k;
19     left <- l;
20     right <- r;
21     self;
22   }};
23
24   isEmpty() : Bool { false };
25
26   insert(k : Int) : Tree {
27     if k <= key then new SELF_TYPE.init(key,
28                                     left.insert(k),
29                                     right)
30     else new SELF_TYPE.init(key,
31                             left,
32                             right.insert(k)) fi
33     (* P1 *)
34   };
35
36   sum() : Int { key + left.sum() + right.sum() };
37 };
38
39 class Main {
40   main() : Object {
41     let tree : Tree <-
42       new Tree.insert(2).insert(1).insert(3).insert(4)
43     (* P2 *)
44     in tree.out_int(tree.sum())
45     (* P3 *)
46   };
47 };
48

```

① Organizare directe în memorie

Tree: tag 0

dim 3

Tree-diopTab → { Object, ... x3
io, ... x4
Tree.isEmpty
Tree.insert
Tree.sum

NETree: tag 1

dim 6

NETree-diopTab → { Object, ... x3
key
left
right
io, ... x4
NETree.isEmpty
NETree.insert
NETree.sum
NETree.init

Înregistrarea de activare

Conținut	Adresă
Parametru n	
\vdots	\vdots
Parametru 2	$\$fp + 16$
Parametru 1	$\$fp + 12$
$\$fp$	
$\$s0$	
$\$ra$	$\$fp$
	$\$sp$

Înregistrarea de activare cu variabile LET

Conținut	Adresă
Parametru n	
\vdots	\vdots
Parametru 2	$\$fp + 16$
Parametru 1	$\$fp + 12$
$\$fp$	
$\$s0$	
$\$ra$	$\$fp$
Variabilă let 1	$\$fp - 4$
Variabilă let 2	$\$fp - 8$
\vdots	\vdots
Variabilă let m	
	$\$sp$

2) Dimensiunea minimă a înreg. de activare pt. metoda **sum** din **NETree**?

1) regiuni: $\$fp, \$so, \$ra \Rightarrow 3$

2) parametri: 0

3) locații temporare: 1

$$\begin{aligned} key + left.sum() + right.sum() &\Rightarrow (key + left.sum()) + right.sum() \\ key + left.sum() &\Rightarrow NT = \max(NT(key), 1 + NT(left.sum())) \\ &= \max(0, 1 + 0) = 1 \\ key + left.sum() + right.sum() &\Rightarrow NT = \max(1, 1 + NT(right.sum())) \\ &= \max(1, 1 + 0) = 1 \end{aligned}$$

\Rightarrow DIMENSIUNE: 4 *curioasă*

3) Completați spațiile libere: (codul MIPS de mai jos este scris pentru **key + left.sum()** din metoda **sum()** a clasei **NETree**)

```

1 → lw    $a0 12($s0) => key
2  sw    $a0 0($sp) 2 push key
3  addiu $sp $sp -4
4  lw    $a0 16($s0) => $a0 → left
5  <verificare dispatch on void>
6  lw    $t1 8($a0) => diagTab
7  lw    $t1 36($t1) => method offset (sum)
8  jalr   $t1
9  jal    Object.copy
10 lw    $t1 4($sp)
11 addiu $sp $sp 4
12 lw    $t1 12($t1)
13 lw    $t2 12($a0)
14 add   $t1 $t1 $t2
15 sw    $t1 12($a0)

```

metoda **sum()** a clasei **NETree**

```

arithmetic(op, e1, e2) ::= <<
<e1>
  sw    $a0 0($sp) 2 push $a0
  addiu $sp $sp -4
<e2>
  jal    Object.copy → $a0 = copy $a0
  lw     $t1 4($sp) 2 pop $a0
  addiu  $sp $sp 4
  lw     $t1 12($t1) → get value from $a0
  lw     $t2 12($a0) → get value from $a0
  <op>   $t1 $t1 $t2
  sw     $t1 12($a0) # int slot
>>

```

O_{e1} = de care rezultă în urma evaluării lui $e1$

\Rightarrow poate fi \rightarrow Int(...)
Bod(...), dacă op este '='

```

Int: 0 tag 0
     4 dim 4
     8 Int - diagTab
     12 0

```

store the value at offset 12 in the new obj (\$a0)

key + left.sum()

- 1) **key** = primul atrib $\Rightarrow \$a0 + 12$
- 2) **left** = al doilea atrib $\Rightarrow \$a0 + 16$
- 3) **sum()** = offset 36 în **diagTab**

4) Semantica operațiilor

new SELF-TYPE din **insert (NETree)** \Rightarrow Tree.insert($l_1 \Rightarrow \dots$ new SELF-TYPE

$$T_0 = \begin{cases} X & \text{if } T = \text{SELF_TYPE and } so = X(\dots) \rightarrow T_0 \text{ NETree} \\ T & \text{otherwise} \end{cases}$$

$class(T_0) = (a_1 : T_1 \leftarrow e_1, \dots, a_n : T_n \leftarrow e_n) \Rightarrow$ class (NETree) = (key: Int, left: Tree, right: Tree)

$l_i = \text{newloc}(S_1)$, for $i = 1 \dots n$ and each l_i is distinct $\Rightarrow l_1, l_2, l_3 = \text{new loc}(S)$

$v_1 = T_0(a_1 = l_1, \dots, a_n = l_n) \Rightarrow v_1 = \text{NETree}(key = l_1, left = l_2, right = l_3)$

$S_2 = S_1[D_{T_1}/l_1, \dots, D_{T_n}/l_n] \Rightarrow S_2 = S_1[\text{Int}(0)/l_1, \text{void } l_2, \text{void } l_3]$

$v_1, S_2, [a_1 : l_1, \dots, a_n : l_n] \vdash \{a_1 \leftarrow e_1; \dots; a_n \leftarrow e_n\} \rightarrow v_2, S_3 \rightarrow$ nu avem expr. de inițializare pt. atribute
 $so, S_1, E \vdash \text{new } T \mapsto v_1, S_3$
 \downarrow SELF-TYPE

$so = \text{NETree}(key = lkey, left = ll, right = lr)$

$E = \{key : lkey, left : ll, right : lr, k : l_k\}$

$S = \{lkey \rightarrow \text{Int}(0), ll \rightarrow \text{void}, lr \rightarrow \text{void}, l_k \rightarrow \text{Int}(0)\}$