# Beijing Normal University School of Mathematics

# Template

app1eDog

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**28** 

math - number theory

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6 1 HPP

# 1 hpp

# 1.1 heading

```
#include <bits/stdc++.h>
         // using namespace std;
  4
        using LL = long long;
using i128 = __int128;
using PII = std::pair<int, int>;
  5
        using UI = unsigned int;
using ULL = unsigned long long;
using ULL = unsigned long long;
using PIL = std::pair<int, LL>;
using PLI = std::pair<LL, int>;
using PLI = std::pair<LL, LL>;
10
11
\overline{13}
14
        using vi = std::vector<int>;
15
        using vi = std::vector<vi>;
using vi = std::vector<vi>;
using vl = std::vector<LL>;
using vvl = std::vector<vl>;
16
17
18
        using vpi = std::vector<PII>;
19
20
21
22
23
        #define ff first
        #define ss second
#define all(v) v.begin(), v.end()
#define rall(v) v.rbegin(), v.rend()
\overline{24}
25
26
27
28
29
30
         #ifdef LOCAL
        #include "debug.h"
         #else
         #define debug(...) \
31
                do {
32
                } while (false)
33
         #endif
34
        constexpr int inf = 0x3f3f3f3f;
constexpr LL INF = 1e18;
35
36
        constexpr int lowbit(int x) { return x & -x; }
37
38
        constexpr int add(int x, int y) { return x + y < mod ? x + y : x - mod + y; }
constexpr int sub(int x, int y) { return x < y ? mod + x - y : x - y; }
constexpr int mul(LL x, int y) { return x * y % mod; }
constexpr void Add(int& x, int y) { x = add(x, y); }
constexpr void Sub(int& x, int y) { x = sub(x, y); }
constexpr void Mul(int& x, int y) { x = mul(x, y); }
constexpr void Mul(int& x, int y) { x = mul(x, y); }</pre>
39
40
41
42
43
44
        constexpr int pow(int x, int y, int z = 1) {
   for (; y; y /= 2) {
      if (y & 1) Mul(z, x);
      }
}
45
46
47
48
                        Mul(x, x);
49
50
                return z:
51
        temps constexpr int add(Ts... x) {
  int y = 0;
  (..., Add(y, x));
52 \\ 53 \\ 54 \\ 55
                return y;
56
57
        temps constexpr int mul(Ts... x) {
                int y = 1;
(..., Mul(y, x));
58
59
60
                return y;
61
62
        tandu bool Max(T& x, const U& y) { return x < y ? x = y, true : false; } tandu bool Min(T& x, const U& y) { return x > y ? x = y, true : false; }
63
64
65
66
        void solut() {
67
              ;
68
69
70
71
72
73
74
75
76
77
        int main() {
                std::ios::sync_with_stdio(false);
                std::cin.tie(0);
                int t = 1;
                std::cin >> t;
while (t--) {
                        solut();
                return 0;
```

1.2 debug.h

# 1.2 debug.h

```
template <typename T, typename U>
std::ostream& operator<<(std::ostream& os, const std::pair<T, U>& p) {
   return os << '<' << p.first << ',' << p.second << '>';';
 \frac{1}{2} \frac{1}{3} \frac{1}{4} \frac{1}{5} \frac{1}{6} \frac{1}{7} \frac{1}{8} \frac{1}{9}
          }
          template <
          typename T, typename = decltype(std::begin(std::declval<T>())),
    typename = std::enable_if_t<!std::is_same_v<T, std::string>>>
std::ostream& operator<<(std::ostream& os, const T& c) {</pre>
10
                   auto it = std::begin(c);
                  if (it == std::end(c)) return os << "{}";
for (os << '{' << *it; ++it != std::end(c); os << ',' << *it);
return os << '}';</pre>
11
13
14
          }
15
         #define debug(arg...)
     do {
16
                           std::cerr << "[" #arg "] :"; \
dbg(arg);
17
18
19
20
21
22
23
                   } while (false)
         template <typename... Ts>
void dbg(Ts... args) {
    (..., (std::cerr << ' ' << args));
    std::cerr << std::endl;</pre>
24
25
26
```

8 2 SHELL SCRIPTS

# 2 shell scripts

# 2.1 linux version

```
#!/bin/bash

cd "$1"

g++ -o main -02 -std=c++17 -DLOCAL main.cpp -ftrapv -fsanitize=address,undefined

for input in *.in; do
    output=${input%.*}.out
    answer=${input%.*}.ans

./main < $input > $ouput

echo "case ${input%.*}: "
    echo "My: "
    cat $output
    echo "Answer: "
    cat $answer

done
```

# 2.2 windows version

# 3 data structure

#### 3.1 stack

```
1  vi stk;
2  for (int i = 1; i <= n; i++){
3     while (!stk.empty() and stk.back() > a[i]) {
4         stk.pop_back();
5     }
6     stk.pop_back(a[i]);
7  }
```

## 3.2 queue

#### 3.3 DSU

```
/* DSU */
vi fa(n + 1);
std::iota(all(fa), 0);
std::function<int(int)> find = [&] (int x) -> int{
    return x == fa[x] ? x : fa[x] = find(fa[x]);
};
auto merge = [&] (int x, int y) -> void{
    x = find(x), y = find(y);
    if (x == y) return;
    // operations //
    fa[y] = x;
};
```

# 3.4 spare table

一维

```
/* spare table */
int B = 30;
       vvi f(n + 1, vi(B));
vi Log2(n + 1);
auto init = [&]() -> void {
 3
 6
7
               for (int i = 1; i <= n; i++) {
    f[i][0] = a[i];
                       if (i > 1) Log2[i] = Log2[i / 2] + 1;
 9
              int t = Log2[n];
for (int j = 1; j <= t; j++) {
   for (int i = 1; i <= n - (1 << j) + 1; i++) {
     f[i][j] = std::max(f[i][j - 1], f[i + (1 << (j - 1))][j - 1]);
}</pre>
10
11
12
13
14
15
               }
16
       };
17
       init();
       init(),
auto query = [&](int l, int r) -> int {
   int t = Log2[r - l + 1];
   return std::max(f[l][t], f[r - (1 << t) + 1][t]);</pre>
18
19
20
       };
21
```

```
/* spare table */
       intB = 30;
 3
       std::vector f(n + 1, std::vector<std::array<std::array<int, B>, B>>(m + 1));
 4
       vi Log2(n + 1);
      auto init = [&]() -> void {
   for (int i = 2; i <= std::max(n, m); i++) {
      Log2[i] = Log2[i / 2] + 1;
 \begin{array}{c} 5 \\ 6 \\ 7 \\ 8 \end{array}
 9
             for (int i = 2; i <= n; i++) {
   for (int j = 2; j <= m; j++) {
     f[i][j][0][0] = a[i][j];
}</pre>
10
11
12
13
             14
15
16
17
18
19
20 \\ 21 \\ 22 \\ 23 \\ 24
                                               f[i][j][ki][kj] =
                                                      std: max(f[i][j][ki - 1][kj], f[i + (1 << (ki - 1))][j][ki - 1][kj]);
                                         } else {
                                               f[i][j][ki][kj]
                                                      std: max(f[i][j][ki][kj-1], f[i][j+(1 << (kj-1))][ki][kj-1]);
25
                                        }
26
                                 }
\overline{27}
                          }
\frac{1}{28}
                    }
29
             }
30
31
       init();
       auto query = [&](int x1, int y1, int x2, int y2) -> int {
   int ki = Log2[x2 - x1 + 1], kj = Log2[y2 - y1 + 1];
32
33
             int k1 = Log2[x2 - x1 + 1], kj = Log2[y2 - y1 + 1];
int t1 = f[x1][y1][ki][kj];
int t2 = f[x2 - (1 << ki) + 1][y1][ki][kj];
int t3 = f[x1][y2 - (1 << kj) + 1][ki][kj];
int t4 = f[x2 - (1 << ki) + 1][y2 - (1 << kj) + 1][ki][kj];
return std::max({t1, t2, t3, t4});</pre>
34
35
36
37
38
39
      };
```

# 3.5 Cartesian tree

一种特殊的平衡树, 用元素的值作为平衡点节点的 val, 元素的下标作为 key.

```
/* cartesian tree */
vi ls(n + 1), rs(n + 1), stk(n + 1);
int top = 1;
for (int i = 1; i <= n; i++) {
    int k = top;
    while (k and a[stk[k]] > a[i]) k--;
    if (k) rs[stk[k]] = i;
    if (k < top) ls[i] = stk[k + 1];
    stk[++k] = i;
    top = k;
}</pre>
```

#### 3.6 segment tree

```
/* segment tree @ czr */
     const int N = 100010;
 3
     struct node {
 4
          int 1, r;
         ll sum, maxn, add, set;
    bool addflag, setflag;
}tr[N << 2];</pre>
    void push_up(int u) {
    tr[u].sum = tr[u << 1].sum_+ tr[u << 1 | 1].sum;</pre>
10
11
         tr[u].maxn = max(tr[u << 1].maxn, tr[u << 1 | 1].maxn);
12
    // 0 4 0 0 0
13
    // 2 6 2 2 2
// 2 6 2 4 4
14
15
16
17
    void push_down(int u) {
18
         auto& root = tr[u], &left = tr[u << 1], &right = tr[u << 1 | 1];</pre>
19
         if (root.setflag) {
```

3.6 segment tree

```
20
                 assert(!root.addflag);
 21
                 left.add = 0, left.set = root.set, left.addflag = false, left.setflag = true;
22
23
                 right.add = 0, right.set = root.set, right.addflag = false, right.setflag = true;
                 left.sum = root.set * (left.r - left.l + 1);
 24
                 right.sum = root.set * (right.r - right.l + 1);
 25
                 left.maxn = root.set;
 26
                 right.maxn = root.set;
 27
                 root.set = 0, root.setflag = false;
 28
            if (root.addflag) {
 29
                 assert(!root.setflag);
 30
                 if (left.setflag) left.set += root.add;
else left.add += root.add, left.addflag = true;
 31
 32
 33
                 if (right.setflag) right.set += root.add;
 34
                 else right.add += root.add, right.addflag = true;
 35
 36
                 left.sum += root.add * (left.r - left.l + 1);
                 right.sum += root.add * (right.r - right.l + 1);
left.maxn += root.add;
 37
 38
 39
                 right.maxn += root.add
 40
                 root.add = 0, root.addflag = false;
 41
 42
            assert(root.add == 0);
 43
 44
      void build(int u, int 1, int r, vector<11>& a) {
   if (1 == r) {
      tr[u].1 = tr[u].r = 1, tr[u].sum = tr[u].maxn = a[1];
      tr[u].add = 0, tr[u].set = 0;
      tr[u].addflag = tr[u].setflag = false;
} also {
 45
 46
 47
 48
 49
 50
 51
                 tr[u].1 = 1, tr[u].r = r, tr[u].add = 0, tr[u].set = 0;
 52
                 tr[u].addflag = tr[u].setflag = false;
 53
                 int mid = l + r >> 1;
                 build(u << 1, 1, mid, a);
build(u << 1 | 1, mid + 1, r, a);
 54
 55
 56
                 push_up(u);
 57
           }
      }
 58
 59
 60
      // 区间加
 61
      void modify(int u, int 1, int r, 11 d) {
            if (1 > r) return;
if (tr[u].1 >= 1 && tr[u].r <= r) {
 62
 63
                 if (tr[u].setflag) tr[u].set += d;
else tr[u].add += d, tr[u].addflag = true;
tr[u].sum += d * (tr[u].r - tr[u].l + 1);
 64
 65
 66
                 tr[u].maxn += d;
 67
 68
            } else {
 69
                 push_down(u);
                 int mid = tr[u].l + tr[u].r >> 1;
if (1 <= mid) modify(u << 1, 1, r, d);
if (r > mid) modify(u << 1 | 1, 1, r, d);</pre>
 70
 71
 72
 73
                 push_up(u);
 74
      }
 76
      // 区间赋值
 77
 78
      void update(int u, int l, int r, ll x) {
 79
            if (1 > r) return;
 80
            if (tr[u].1 >= 1 && tr[u].r <= r) {</pre>
                 tr[u].set = x, tr[u].setflag = true;
tr[u].add = 0, tr[u].addflag = false;
 81
 83
                 tr[u].sum = x * (tr[u].r - tr[u].l + 1);
 84
                 tr[u].maxn = x;
 85
            } else {
 86
                 push_down(u)
                 int mid = tr[u].l + tr[u].r >> 1;
 87
                 if (1 <= mid) update(u << 1, 1, r, x);
if (r > mid) update(u << 1 | 1, 1, r, x);</pre>
 88
 89
 90
                 push_up(u);
 91
      }
 92
 93
      11 query_sum(int u, int 1, int r) {
    if (1 > r) return 0;
    if (tr[u].1 >= 1 && tr[u].r <= r) return tr[u].sum;</pre>
 94
 95
 96
            else {
    ll res = 0;
 97
 98
 99
                 push_down(u);
                 int mid = tr[u].l + tr[u].r >> 1;
100
101
                 if (1 <= mid) res += query_sum(u << 1, 1, r);</pre>
                 if (r > mid) res += query_sum(u << 1 | 1, 1, r);</pre>
102
103
                 return res;
104
      }
105
106
```

 $3-DATA\ STRUCTURE$ 

```
107
     |ll query_maxn(int u, int l, int r) {
            if (1 > r) return -1e18;
if (tr[u].1 >= 1 && tr[u].r <= r) return tr[u].maxn;
108
109
110
            else {
111
                 11 \text{ res} = -1e18;
                push_down(u);
112
113
                 int mid = tr[u].l + tr[u].r >> 1;
                 if (1 <= mid) res = max(res, query_maxn(u << 1, 1, r));
114
                 if (r > mid) res = max(res, query_maxn(u << 1 | 1, 1, r));</pre>
115
116
                 return res;
           }
117
118
      }
119
      // 找到最小 i 使得 sum(1, i) >= k
ll find_presum_idx(int_u, int 1, int r, int x) {
120
121
122
            if (tr[u].1 == tr[u].r) return tr[u].1;
123
            else {
124
                push_down(u);
                 int mid = tr[u].l + tr[u].r >> 1;
125
126
                 if (r <= mid) {</pre>
127
                      return find_presum_idx(u << 1, 1, r, x);</pre>
128
                 } else if (1 > mid) {
129
                      return find_presum_idx(u << 1 | 1, 1, r, x);</pre>
130
                 } else {
                      il lsum = query_sum(u << 1, 1, r);
if (lsum >= x) return find_presum_idx(u << 1, 1, mid, x);
else return find_presum_idx(u << 1 | 1, mid + 1, r, x - lsum);</pre>
131
132
133
134
                }
           }
135
      }
136
```

# 3.7 segment tree split

12

```
/* segment tree split @ wrb */
     #include<bits/stdc++.h>
 \bar{3}
     using namespace std;
     namespace Acc{
 \begin{array}{c} 4\\5\\6\\7 \end{array}
          using i64=int64_t;
enum{N=200009,M=10000000};
          i64 v[M];
 8
          int lc[M],rc[M],tot,a[N],r[N];
 9
          auto up=[](int o){
10
               v[o]=v[lc[o]]+v[rc[o]];
11
12
          void bd(int&o,int 1,int r){
13
               if(o=++tot,l==r)return cin>>v[o],void();
14
               int md=l+r>>1;
               bd(lc[o],1,md),bd(rc[o],md+1,r),up(o);
15
16
17
          void spl(int&o,int&x,int l,int r,int L,int R){
18
19
               if(l<=L&&R<=r)return o=x,x=0,void();</pre>
               int md=L+R>>1;
20
21
22
23
24
25
               o=++tot;
               if(l<=md)spl(lc[o],lc[x],l,r,L,md);</pre>
               if(r>md)spl(rc[o],rc[x],l,r,md+1,R);
               up(o),up(\bar{x});
          void mg(int&o,int x,int l,int r){
   if(!o||!x)return o|=x,void();
\frac{1}{26}
27
               if(l==r)return v[o]+=v[x],void();
28
               int md=l+r>>1;
               mg(lc[o],lc[x],l,md);
mg(rc[o],rc[x],md+1,r);
29
30
31
32
               up(o);
\frac{33}{34}
          void ins(int&o,int l,int r,int x,int k){
               if(!o)o=++tot;
35
               if(v[o]+=k,l==r)return;
36
               int md=l+r>>1;
37
               x<=md?ins(lc[o],1,md,x,k):ins(rc[o],md+1,r,x,k);</pre>
38
          i64 qry(int o,int 1,int r,int L,int R){
   if(!o)return 0;
39
40
41
               if(1<=L&&R<=r)return v[o];</pre>
               int md=L+R>>1;i64 z=0;
if(l<=md)z=qry(lc[o],l,r,L,md);
42
43
               if(r>md)z+=qry(rc[o],1,r,md+1,R);
44
45
               return z;
46
47
          int kth(int o,int l,int r,int k){
48
               if(l==r)return 1;
49
               if(k>v[o])return -1;
```

```
50
                     int md=l+r>>1;
51
                     if(k<=v[lc[o]])return kth(lc[o],1,md,k);</pre>
52
                     else return kth(rc[o],md+1,r,k-v[lc[o]]);
53
54
              auto work=[](){
55
                     int n,m,i,x,y,o=1;
                     for(cin>>n>m,bd(r[1],1,n);m--;)switch(cin>>i,i){
    case 0:cin>>i>x>>y,spl(r[++o],r[i],x,y,1,n);break;
56
57
                            case 0.cim>>i>>x>>y,spi([++0],f[1],x,y,1,m),break;
case 1:cim>>x>>y,mg(r[x],r[y],1,n);break;
case 2:cim>>i>>x>>y,ins(r[i],1,n,y,x);break;
case 3:cim>>i>>x>>y,cout<<qry(r[i],x,y,1,n)<<'\n';break;
case 4:cim>>i>>x,cout<<kth(r[i],1,n,x)<<'\n';break;</pre>
58
59
60
61
62
                     }
63
              };
64
65
       int main(){
66
              ios::sync_with_stdio(0);
              cin.tie(0),Acc::work();
67
68
       }
```

# 3.8 persistent segment tree

### 单点修改, 版本拷贝

n 个数, m 次操作, 操作分别为

- 1.  $v_i$  1  $loc_i$   $value_i$ : 将第  $v_i$  个版本的  $a[loc_i]$  修改为  $value_i$ ,
- 2.  $v_i$  2  $loc_i$ : 拷贝第  $v_i$  个版本, 并查询该版本的  $a[loc_i]$ .

```
// 洛谷 P3919 【模板】可持久化线段树 1 (可持久化数组)
 3
     struct node {
 4
         int 1, r, key;
 5
    };
 6
 7
     int main() {
 8 9
         std::ios::sync_with_stdio(false);
std::cin.tie(0);
10
         std::cout.tie(0);
11
         int n, m;
std::cin >> n >> m;
12
13
         vi a(n + 1);
for (int i = 1; i <= n; i++) {</pre>
14
15
16
              std::cin >> a[i];
18
         /* hjt segment tree */
int idx = 0;
19
20
21
         vi root(m + 1);
22
         std::vector<node> tr(n * 25);
23
\overline{24}
         std::function<int(int, int)> build = [&](int 1, int r) -> int {
              int p = ++idx;
if (1 == r) {
\overline{25}
26
27
                  tr[p].key = a[1];
\frac{1}{28}
                  return p;
29
30
              int mid = (1 + r) >> 1;
31
              tr[p].l = build(1, mid);
32
              tr[p].r = build(mid + 1, r);
33
              return p;
34
35
         36
37
              int q = ++idx;
tr[q].1 = tr[p].1, tr[q].r = tr[p].r;
if (tr[q].1 == tr[q].r) {
    tr[q].key = x;
38
39
40
41
42
                  return q;
43
              int mid = (1 + r) >> 1;
44
45
              if (k <= mid) {</pre>
46
                  tr[q].l = modify(tr[q].l, l, mid, k, x);
              } else {
```

3 DATA STRUCTURE

```
48
                   tr[q].r = modify(tr[q].r, mid + 1, r, k, x);
49
              }
50
              return q;
51
52
         };
53
         std::function<int(int, int, int, int)> query = [&](int p, int l, int r, int k) -> int {
54
55
56
57
              if (tr[p].1 == tr[p].r) {
                   return tr[p].key;
              int mid = (1 + r) >> 1;
if (k <= mid) {</pre>
58
59
                   return query(tr[p].1, 1, mid, k);
              } else {
60
61
                   return query(tr[p].r, mid + 1, r, k);
62
              }
63
64
65
         root[0] = build(1, n);
66
67
         for (int i = 1; i <= m; i++) {</pre>
              int op, ver, k, x;
std::cin >> ver >> op;
68
69
              if (op == 1) {
70
71
72
73
74
75
76
77
78
79
                   std::cin >> k >> x;
                   root[i] = modify(root[ver], 1, n, k, x);
              } else {
                   std::cin >> k;
                   root[i] = root[ver];
                   std::cout << query(root[ver], 1, n, k) << ' \n';
              }
80
         return 0;
    }
```

#### 区间第 k 小

长度为 n 的序列 a, m 次查询, 每次查询 [l,r] 中的第 k 小值.

```
// 洛谷P3834 【模板】可持久化线段树 2

    \begin{array}{r}
      123456789
    \end{array}

     struct node {
          int 1, r, cnt;
     };
     int main() {
          std::ios::sync_with_stdio(false);
          std::cin.tie(0)
10
          std::cout.tie(0);
11
12
          int n, m;
13
          std::cin >> n >> m;
          vi a(n + 1), v;
for (int i = 1; i <= n; i++) {
    std::cin >> a[i];
14
15
16
17
                v.push_back(a[i]);
18
19
          std::sort(all(v));
          v.erase(unique(all(v)), v.end());
auto find = [&](int x) -> int { return std::lower_bound(all(v), x) - v.begin() + 1; };
20
21
22
23
24
25
26
27
28
29
           /* hjt segment tree */
          std::vector<node>(n * 25);
           vi root(n + 1);
          int idx = 0;
          std::function<int(int, int)> build = [&](int 1, int r) -> int {
                int p = ++idx;
if (l == r) return p;
30
31
                int mid = (1 + r) > 1;
\begin{array}{c} 32 \\ 33 \\ 34 \\ 35 \end{array}
                tr[p].1 = build(1, mid), tr[p].r = build(mid + 1, r);
                return p;
          };
36
37
          std::function<int(int, int, int, int)> modify = [&](int p, int 1, int r, int x) -> int {
                int q = ++idx;
tr[q] = tr[p];
38
39
                if (tr[q].l == tr[q].r) {
40
                     tr[\bar{q}].cnt++;
41
                     return q;
42
43
                int mid = (1 + r) >> 1;
```

3.9 sweep line

```
44
             if (x <= mid) {</pre>
             tr[q].1 = modify(tr[q].1, 1, mid, x);
} else {
45
46
47
                 tr[q].r = modify(tr[q].r, mid + 1, r, x);
48
49
             tr[q].cnt = tr[tr[q].1].cnt + tr[tr[q].r].cnt;
50
51
        };
52
53
         std::function<int(int, int, int, int, int)> query = [&](int p, int q, int l, int r,
                                                                      int x) -> int {
54
55
             if (1 == r) return 1;
56
             int cnt = tr[tr[p].1].cnt - tr[tr[q].1].cnt;
             int mid = (1 + r) >> 1;
57
             if (x <= cnt) {</pre>
59
                 return query(tr[p].1, tr[q].1, 1, mid, x);
60
             } else {
61
                 return query(tr[p].r, tr[q].r, mid + 1, r, x - cnt);
             }
62
63
        };
64
65
        root[0] = build(1, v.size());
66
67
68
         for (int i = 1; i <= n; i++) {</pre>
69
             root[i] = modify(root[i - 1], 1, v.size(), find(a[i]));
70
71
72
73
74
75
76
77
         for (int i = 1; i <= m; i++) {</pre>
             int 1, r, k;
             std::cin >> 1 >> r >> k;
             std::cout << v[query(root[r], root[l - 1], 1, v.size(), k) - 1] << '\n';
        return 0;
78
    }
```

# 3.9 sweep line

```
/* sweep line @ czr */
 3
     struct Node {
          int 1, r;
 4
          ll sum, length, res;
 5
     }tr[N << 2];</pre>
 6
7
     void push_up(int u) {
 8 9
          tr[u].res = (tr[u << 1].res + tr[u << 1 | 1].res) % Mod;
10
     void update_length(int u) {
   if (tr[u].sum) {
11
12
13
               tr[u].length = tr[u].res;
14
            else
15
               if (tr[u].1 == tr[u].r) tr[u].length = 0;
16
               else tr[u].length = (tr[u << 1].length + tr[u << 1 | 1].length) % Mod;</pre>
17
          }
     }
19
     void build(int u, int 1, int r) {
   if (1 == r) tr[u] = {1, r, 0, 0, 0};
20
21
22
          else {
23
               tr[u] = {1, r, 0, 0, 0};
               int mid = 1 + r >> 1;
build(u << 1, 1, mid);
2\overline{4}
\overline{25}
26
               build(u << 1 | 1, mid + 1, r);
27
              push_up(u);
28
          }
29
     }
30
31
     void modify(int u, int l, int r, int op) {
32
          if (tr[u].1 >= 1 && tr[u].r <= r) {
33
               tr[u].sum += op;
34
               update_length(u);
35
          } else {
36
               int mid = tr[u].1 + tr[u].r >> 1;
               if (1 <= mid) modify(u << 1, 1, r, op);
if (r > mid) modify(u << 1 | 1, 1, r, op);</pre>
37
38
39
               push_up(u);
40
               update_length(u);
41
          }
42
43
     void change(int u, int x, ll d) {
```

```
45
            if (tr[u].l == tr[u].r) {
46
                  tr[u].res = (tr[u].res + d) % Mod;
47
                  update_length(u);
48
            } else {
                  int mid = tr[u].l + tr[u].r >> 1;
if (x <= mid) change(u << 1, x, d);
else change(u << 1 | 1, x, d);</pre>
49
50
51
52
53
54
                  push_up(u);
                  update_length(u);
            }
55
```

```
/* sweep line @ wrb */
 23
     #define int long long
     const int N = 2e5+10;
 4
    int b[N<<1],n,len,ans;</pre>
    struct node{
    int y1,y2,x,k;
}a[N<<1];</pre>
 6
7
     struct Seg{
 9
    #define lc (o<<1)</pre>
    #define rc (o<<1|1)
    static const int N = 5e6+10;</pre>
10
11
12
         int sum[N],cnt[N],tag[N];
13
         void push_up(int o,int 1,int r){
14
              if(sum[o])cnt[o]=b[r+1]-b[1];
15
              else cnt[o]=cnt[lc]+cnt[rc];
16
17
          void add(int o,int l,int r,int L,int R,int k){
18
              if(r<L || 1>R)return;
19
              if(l==L && r==R)return (void)(sum[o]+=k,push_up(o,1,r));
20
              int mid=L+R>>1;
              if(r<=mid)add(lc,l,r,L,mid,k);
else if(l>mid)add(rc,l,r,mid+1,R,k);
\overline{21}
\overline{22}
23
              else add(lc,1,mid,L,mid,k),add(rc,mid+1,r,mid+1,R,k);
24
              push_up(o,L,R);
\overline{25}
         }
\frac{26}{27}
    #undef lc
     #undef rc
28
29
30
    void work(){
         cin>>n;
31
32
         for(int i=1,x1,y1,x2,y2;i<=n;i++){</pre>
              cin>>x1>>y1>>x2>>y2;
              b[i*2-1]=y1,b[i*2]=y2,a[i*2-1]={y1,y2,x1,1},a[i*2]={y1,y2,x2,-1};
33
34
35
         n << =1;
36
         sort(b+1,b+n+1),len=unique(b+1,b+n+1)-b-1;
37
         for(int i=1;i<=n;i++)a[i].y1=lower_bound(b+1,b+len+1,a[i].y1)-b,a[i].y2=lower_bound(b+1,b+len+1,a[i].</pre>
         y2)-b;
sort(a+1,a+n+1,[](node a,node b)->bool{return a.x<b.x;});
38
39
40
              t.add(1,a[i].y1,a[i].y2-1,1,len-1,a[i].k);
41
              ans+=t.cnt[1]*(a[i+1].x-a[i].x);
         }
42
43
          cout<<ans;
44
    #undef int
```

#### 3.10 treap

# fhq treap

n 次操作, 操作分为如下 6 种:

- 1. 插入数 x;
- 2. 删除数 x (若有多个相同的数,只删除一个);
- 3. 查询数 x 的排名 (排名定义为小于 x 的数的个数 + 1);
- 4. 查询排名为 x 的数;
- 5. 求 x 的前驱 (前驱定义为小于 x 的最大数);

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6. 求 x 的后继 (后继定义为大于 x 的最小数).

```
struct node {
 3
            node *ch[2];
             int key, val;
 4
             int cnt, size;
 5
            node(int _key) : key(_key), cnt(1), size(1) {
    ch[0] = ch[1] = nullptr;
 6
7
 8
                   val = rand();
10
             // node(node *_node) {
11
            // key = _node->key, val = _node->val, cnt = _node->cnt, size = _node->size;
// }
12
13
14
15
             inline void push_up() {
                   size = cnt;
if (ch[0] != nullptr) size += ch[0]->size;
16
17
18
                   if (ch[1] != nullptr) size += ch[1]->size;
19
20
      };
21
22
      struct treap {
      #define _2 second.first
#define _3 second.second
23
24
25
26
            node *root:
27
            pair<node *, node *> split(node *p, int key) {
   if (p == nullptr) return {nullptr, nullptr};
   if (p->key <= key) {</pre>
28
29
30
                         auto temp = split(p->ch[1], key);
p->ch[1] = temp.first;
31
32
33
                         p->push_up();
34
                         return {p, temp.second};
35
                         auto temp = split(p->ch[0], key);
p->ch[0] = temp.second;
36
37
                        p->push_up();
return {temp.first, p};
38
39
40
                   }
            }
41
42
            pair<node *, pair<node *, node *> > split_by_rank(node *p, int rank) {
   if (p == nullptr) return {nullptr, {nullptr, nullptr}};
   int ls_size = p->ch[0] == nullptr ? 0 : p->ch[0]->size;
   if (rank <= ls_size) {</pre>
43
44
45
46
                         auto temp = split_by_rank(p->ch[0], rank);
p->ch[0] = temp._3;
47
48
                  p > comp._o,
p -> push_up();
return {temp.first, {temp._2, p}};
} else if (rank <= ls_size + p->cnt) {
   node *lt = p->ch[0];
49
50
51
52
                         node *rt = p->ch[1];
p->ch[0] = p->ch[1] = nullptr;
53
54
                         return {lt, {p, rt}};
55
56
                         auto temp = split_by_rank(p->ch[1], rank - ls_size - p->cnt);
p->ch[1] = temp.first;
57
58
59
                         p->push_up();
60
                         return {p, {temp._2, temp._3}};
61
                   }
62
            }
63
            node *merge(node *u, node *v) {
   if (u == nullptr && v == nullptr) return nullptr;
   if (u != nullptr && v == nullptr) return u;
64
65
66
                   if (v != nullptr && u == nullptr) return v;
67
                   if (u->val < v->val) {
    u->ch[1] = merge(u->ch[1], v);
68
69
70
71
72
73
74
75
76
77
78
                         u->push_up();
                         return u;
                   } else {
                         v\rightarrow ch[0] = merge(u, v\rightarrow ch[0]);
                         v->push_up();
                         return v;
                   }
79
             void insert(int key) {
                   auto temp = split(root, key);
auto l_tr = split(temp.first, key - 1);
80
81
82
                   node *new_node;
83
                   if (l_tr.second == nullptr) {
                         new_node = new node(key);
```

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```
85
                } else {
 86
                     1_tr.second->cnt++;
 87
                     1_tr.second->push_up();
 88
 89
                node *l_tr_combined = merge(l_tr.first, l_tr.second == nullptr ? new_node : l_tr.second);
 90
                root = merge(l_tr_combined, temp.second);
 91
 92
           void remove(int key) {
   auto temp = split(root, key);
   auto l_tr = split(temp.first, key - 1);
   if (l_tr.second->cnt > 1) {
 93
 94
 95
 96
                     1_tr.second->cnt--
 97
 98
                     1_tr.second->push_up();
 99
                     l_tr.first = merge(l_tr.first, l_tr.second);
100
                } else {
                     if (temp.first == l_tr.second) temp.first = nullptr;
delete l_tr.second;
101
102
103
                     1_tr.second = nullptr;
104
105
                root = merge(l_tr.first, temp.second);
106
107
108
           int get_rank_by_key(node *p, int key) {
                auto temp = split(p, key - 1);
int ret = (temp.first == nullptr ? 0 : temp.first->size) + 1;
109
110
111
                root = merge(temp.first, temp.second);
112
                return ret;
113
114
           int get_key_by_rank(node *p, int rank) {
   auto temp = split_by_rank(p, rank);
115
116
                int ret = temp._2->key;
root = merge(temp.first, merge(temp._2, temp._3));
117
118
119
                return ret;
120
121
122
           int get_prev(int key) {
                auto temp = split(root, key - 1);
123
124
                int ret = get_key_by_rank(temp.first, temp.first->size);
125
                root = merge(temp.first, temp.second);
126
                return ret;
127
           }
128
129
           int get_nex(int key) {
130
                auto temp = split(root, key);
                int ret = get_key_by_rank(temp.second, 1);
root = merge(temp.first, temp.second);
131
132
133
                return ret;
134
     };
135
136
137
      treap tr;
138
139
      int main() {
           ios::sync_with_stdio(false);
140
141
           cin.tie(0)
142
           cout.tie(0);
143
144
           srand(time(0));
145
           int n;
146
147
           cin >> n;
148
           while (n--) {
                int op, x;
cin >> op >> x;
if (op == 1) {
149
150
151
                tr.insert(x);
} else if (op == 2) {
152
153
154
                tr.remove(x);
} else if (op == 3) {
   cout << tr.get_rank_by_key(tr.root, x) << '\n';</pre>
155
156
157
                } else if (op == 4) {
158
                     cout << tr.get_key_by_rank(tr.root, x) << '\n';</pre>
159
                } else if (op == 5) {
160
                     cout << tr.get_prev(x) << '\n';</pre>
161
                  else {
162
                     cout << tr.get_nex(x) << '\n';</pre>
                }
163
164
           return 0;
165
      }
166
```

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#### 用 01 trie 实现的一种方式

同样的题目, 注意使用 01 trie 只能存在非负数. 速度能快不少, 但只能单点操作, 而且有点费空间.

```
// 洛谷 P3369 【模板】普通平衡树
 \overline{2}
      struct Treap {
   int id = 1, maxlog = 25;
   int ch[N * 25][2], siz[N * 25];
 3
 4
 5
 6
           int newnode() {
 8
 9
                 ch[id][0] = ch[id][1] = siz[id] = 0;
10
                 return id;
12
13
           void merge(int key, int cnt) {
14
                 int \ddot{u} = 1;
                 for (int i = maxlog - 1; i >= 0; i--) {
   int v = (key >> i) & 1;
   if (!ch[u][v]) ch[u][v] = newnode();
15
16
17
18
                      u = ch[u][v];
19
                      siz[u] += cnt;
                 }
20
\overline{21}
           }
22
           int get_key_by_rank(int rank) {
   int u = 1, key = 0;
   for (int i = maxlog - 1; i >= 0;
      if (siz[ch[u][0]] >= rank) {
23
24
25
                                             - 1; i >= 0; i--) {
26
                      u = ch[u][0];
} else {
27
                            key |= (1 << i);
29
                            rank -= siz[ch[u][0]];
30
31
                            u = ch[u][1];
32
                      }
33
34
                 return key;
35
36
37
           int get_rank_by_key(int rank) {
38
                 int key = 0;
39
                 int u = 1;
                 for (int i = maxlog - 1; i >= 0; i--) {
   if ((rank >> i) & 1) {
40
41
                            key += siz[ch[u][0]];
u = ch[u][1];
42
43
44
                      } else {
                            u = ch[u][0];
45
46
47
                       if (!u) break;
48
49
                 return key;
50
51
           int get_prev(int x) { return get_key_by_rank(get_rank_by_key(x)); }
int get_next(int x) { return get_key_by_rank(get_rank_by_key(x + 1) + 1); }
53
54
      } treap;
55
56
      const int num = 1e7;
57
      int n, op, x;
58
59
      int main() {
60
           std::ios::sync_with_stdio(false);
61
           std::cin.tie(0);
62
           std::cout.tie(0);
63
64
           std::cin >> n;
           for (int i = 1; i <= n; i++) {
   std::cin >> op >> x;
   if (op == 1) {
65
66
67
68
                      treap.merge(x + num, 1);
69
                 } else if (op == 2) {
                 treap.merge(x + num, -1);
} else if (op == 3) {
    std::cout << treap.get_rank_by_key(x + num) + 1 << '\n';
}</pre>
70
71
72
73
74
75
                 } else if (op == 4) {
                      std::cout << treap.get_key_by_rank(x) - num << '\n';</pre>
                 } else if (op == 5) {
                      std::cout << treap.get_prev(x + num) - num << '\n';</pre>
76
77
                   else if (op == 6) {
78
                       std::cout << treap.get_next(x + num) - num << '\n';</pre>
79
80
           return 0;
```

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82 |}

20

# 3.11 splay

#### 文艺平衡树

初始为 1 到 n 的序列, m 次操作, 每次将序列下标为  $[l \sim r]$  的区间翻转.

```
// 洛谷 P3391 【模板】文艺平衡树
 1
2
3
     struct node {
 \begin{array}{c} 4\\5\\6\\7\end{array}
          int ch[2], fa, key;
          int siz, flag;
          void init(int _fa, int _key) { fa = _fa, key = _key, siz = 1; }
     };
 9
     struct splay {
   node tr[N];
10
11
12
          int n, root, idx;
13
14
15
          bool get(int u) { return u == tr[tr[u].fa].ch[1]; }
          void pushup(int u) { tr[u].siz = tr[tr[u].ch[0]].siz + tr[tr[u].ch[1]].siz + 1; }
16
17
18
19
          void pushdown(int u) {
                if (tr[u].flag) {
20
21
22
                    std::swap(tr[u].ch[0], tr[u].ch[1]);
tr[tr[u].ch[0]].flag ^= 1, tr[tr[u].ch[1]].flag ^= 1;
                    tr[u].flag = 0;

    \begin{array}{r}
      23 \\
      24 \\
      25
    \end{array}

               }
          }
\frac{1}{26}
          void rotate(int x) {
                int y = tr[x].fa, z = tr[y].fa;
\overline{27}
28
                int op = get(x);
29
30
31
               tr[y].ch[op] = tr[x].ch[op ^ 1];
if (tr[x].ch[op ^ 1]) tr[tr[x].ch[op ^ 1]].fa = y;
tr[x].ch[op ^ 1] = y;
tr[y].fa = x, tr[x].fa = z;
32
33
34
                if (z) tr[z].ch[y == tr[z].ch[1]] = x;
               pushup(y), pushup(x);
35
36
          37
38
39
40
41
                if (k == 0) root = u;
          }
42
43
44
          void output(int u) {
45
               pushdown(u);
46
                if (tr[u].ch[0]) output(tr[u].ch[0]);
               if (tr[u].key >= 1 && tr[u].key <= n) {
   std::cout << tr[u].key << ' ';</pre>
47
48
49
50
51
52
53
54
55
56
                if (tr[u].ch[1]) output(tr[u].ch[1]);
          }
          void insert(int key) {
                idx++:
                tr[idx].ch[0] = root;
                tr[idx].init(0, key);
57
58
               tr[root].fa = idx;
               root = idx;
59
               pushup(idx);
60
61
62
          int kth(int k) {
                int u = root;
63
                while (1) {
64
65
                    pushdown(u);
                     if (tr[u].ch[0] && k <= tr[tr[u].ch[0]].siz) {
66
67
                          u = tr[u].ch[0];
                    } else {
    k -= tr[tr[u].ch[0]].siz + 1;
68
69
70 \\ 71 \\ 72
                          if (k <= 0) {</pre>
                               opt(u, 0);
                               return u;
73
                          } else {
```

3.11 splay 21

```
74
75
76
77
78
79
                              u = tr[u].ch[1];
                         }
                    }
               }
 80
      } splay;
 81
 82
      int n, m, 1, r;
 83
 84
      int main() {
 85
           std::ios::sync_with_stdio(false);
 86
           std::cin.tie(0);
 87
           std::cout.tie(0);
 88
 89
           std::cin >> n >> m;
 90
           splay.n = n;
 91
           splay.insert(-inf);
 92
           rep(i, 1, n) splay.insert(i);
 93
           splay.insert(inf);
 94
           rep(i, 1, m) {
 95
               std::cin >> 1 >> r;
 96
               1 = \text{splay.kth}(1), r = \text{splay.kth}(r + 2);
               splay.opt(1, 0), splay.opt(r, 1);
splay.tr[splay.tr[r].ch[0]].flag ^= 1;
 97
 98
 99
100
           splay.output(splay.root);
101
102
           return 0;
103
      }
```

# 普通平衡树

```
// 洛谷 P3369 【模板】普通平衡树
 2
 3
     struct node {
 4
          int ch[2], fa, key, siz, cnt;
 5
 6
7
          void init(int _fa, int _key) { fa = _fa, key = _key, siz = cnt = 1; }
 8
9
          void clear() { ch[0] = ch[1] = fa = key = siz = cnt = 0; }
    };
10
11
     struct splay
12
         node tr[N];
13
          int n, root, idx;
14
          bool get(int u) { return u == tr[tr[u].fa].ch[1]; }
15
16
          void pushup(int u) { tr[u].siz = tr[tr[u].ch[0]].siz + tr[tr[u].ch[1]].siz + tr[u].cnt; }
17
18
19
          void rotate(int x) {
              int y = tr[x].fa, z = tr[y].fa;
int op = get(x);
tr[y].ch[op] = tr[x].ch[op ^ 1];
20
\overline{21}
\overline{22}
              if (tr[x].ch[op ^ 1]) tr[tr[x].ch[op ^ 1]].fa = y;
tr[x].ch[op ^ 1] = y;
tr[y].fa = x, tr[x].fa = z;
if (z) tr[z].ch[y == tr[z].ch[1]] = x;
23
24
25
26
27
              pushup(y), pushup(x);
28
29
30
          void opt(int u, int k) {
    for (int f = tr[u].fa; f = tr[u].fa, f != k; rotate(u)) {
31
                   if (tr[f].fa != k) {
32
33
                        rotate(get(u) == get(f) ? f : u);
34
35
36
               if (k == 0) root = u;
37
38
          void insert(int key) {
39
40
               if (!root) {
41
                   idx++;
42
                   tr[idx].init(0, key);
43
                   root = idx;
44
                   return;
45
               int u = root, f = 0;
46
               while (1) {
47
                   if (tr[u].key == key) {
48
49
                        tr[u].cnt++;
                        pushup(u), pushup(f);
50
```

3 DATA STRUCTURE

```
opt(u, 0);
 52
                        break;
 53
 54
                   f = u, u = tr[u].ch[tr[u].key < key];
 55
                   if (!u) {
 56
                        idx++;
                       tr[idx].init(f, key);
tr[f].ch[tr[f].key < key] = idx;
pushup(idx), pushup(f);</pre>
 57
 58
 59
 60
                        opt(idx, 0);
 61
                        break;
 62
                   }
 63
              }
          }
 64
 65
 66
          // 返回节点编号 //
 67
          int kth(int rank) {
 68
               int u = root;
69
               while (1) {
                   if (tr[u].ch[0] && rank <= tr[tr[u].ch[0]].siz) {</pre>
70
71
72
73
74
75
76
77
78
79
                       u = tr[u].ch[0];
                   } else {
                        rank -= tr[tr[u].ch[0]].siz + tr[u].cnt;
                        if (rank <= 0) {</pre>
                            opt(u, 0);
                            return u;
                        } else {
                            u = tr[u].ch[1];
                        }
 80
                   }
 81
              }
 82
          }
 83
 84
          // 返回排名 //
 85
          int nlt(int key) {
 86
               int rank = 0, u = root;
               while (1) {
 87
                   if (tr[u].key > key) {
    u = tr[u].ch[0];
 88
 89
 90
                   } else {
                        rank += tr[tr[u].ch[0]].siz;
 91
                        if (tr[u].key == key) {
    opt(u, 0);
 92
 93
 94
                            return rank + 1;
 95
                        }
 96
                        rank += tr[u].cnt;
 97
                        if (tr[u].ch[1]) {
                            u = tr[u].ch[1];
 98
 99
                        } else {
100
                            return rank + 1;
                        }
101
102
                   }
103
              }
104
105
106
          int get_prev(int key) { return kth(nlt(key) - 1); }
107
108
          int get_next(int key) { return kth(nlt(key + 1)); }
109
110
          void remove(int key) {
111
              nlt(key);
112
               if (tr[root].cnt > 1) {
                   tr[root].cnt--;
113
114
                   pushup(root);
115
                   return;
116
               int u = root, l = get_prev(key);
tr[tr[u].ch[1]].fa = 1;
117
118
              tr[1].ch[1] = tr[u].ch[1];
tr[u].clear();
119
120
121
              pushup(root);
122
123
          124
125
126
               if (tr[u].ch[1]) output(tr[u].ch[1]);
127
128
129
130
     } splay;
131
132
     int n, op, x;
133
134
     int main() {
135
          std::ios::sync_with_stdio(false);
136
          std::cin.tie(0);
          std::cout.tie(0);
```

3.12 link cut tree 23

```
138
139
          splay.insert(-inf), splay.insert(inf);
140
141
          std::cin >> n;
142
          for (int i = 1; i <= n; i++) {
              std::cin >> op >> x;
if (op == 1) {
143
144
145
                   splay.insert(x)
146
              } else if (op == 2)
147
                   splay.remove(x);
              } else if (op == 3) {
148
                   std::cout << splay.nlt(x) - 1 << endl;</pre>
149
150
              } else if (op == 4) {
151
                   std::cout << splay.tr[splay.kth(x + 1)].key << endl;</pre>
              } else if (op == 5) {
152
153
                   std::cout << splay.tr[splay.get_prev(x)].key << endl;</pre>
              } else if (op == 6) {
154
155
                   std::cout << splay.tr[splay.get_next(x)].key << endl;</pre>
156
          }
157
158
159
          return 0;
     }
160
```

#### 3.12 link cut tree

```
/* link cut tree @ wrb */
 2
      struct LCT{
            int v[N],r[N],f[N],s[N][2],st[N],tp;
void pu(int x){v[x]=a[x]^v[s[x][0]]^v[s[x][1]];}
void flp(int x){r[x]^=1,std::swap(s[x][0],s[x][1]);}
void pd(int x){if(r[x])flp(s[x][0]),flp(s[x][1]),r[x]=0;}
 3
 4
 5
 6
 7
            bool isrt(int x){return s[f[x]][0]!=x&&s[f[x]][1]!=x;}
 8
                 int y=f[x],z=f[y],k=(s[y][1]==x);if(!isrt(y)) s[z][y==s[z][1]]=x;
f[x]=z,f[y]=x,f[s[x][k^1]]=y,s[y][k]=s[x][k^1],s[x][k^1]=y,pu(y),pu(x);}
 9
10
11
            void spl(int x){
                 st[tp++]=x;for(int i=x;!isrt(i);i=f[i])st[tp++]=f[i];
12
                 while(tp)pd(st[--tp]);
13
                  while(!isrt(x)){
14
                       if(!isrt(f[x]))rtt((s[f[x]][0]==x)^(s[f[f[x]]][0]==f[x])?x:f[x]);
15
16
                       rtt(x);
17
                 }pu(x);
18
19
            void acc(int x){for(int y=0;x;y=x,x=f[x]) spl(x),s[x][1]=y,pu(x);}
            void mkrt(int x){acc(x),spl(x),flp(x);}
int fdrt(int x){acc(x),spl(x);while(s[x][0])x=s[x][0];spl(x);return x;}
20
21
            void cut(int x,int y){mkrt(x);if(x==fdrt(y)&&f[y]==x&&!s[y][0])s[x][1]=f[y]=0,pu(x);}
void lk(int x,int y){mkrt(x);if(x!=fdrt(y))f[x]=y;}
22
23
      }t;
```

#### 3.13 Lichao tree

```
/* Lichao tree @ wrb */
     #include<bits/stdc++.h>
 3
     using namespace std;
     namespace Acc{
     #define lc (o<<1)
#define rc (o<<1|1)
 5
 6
          const int N = 4e5+10;
          int v[N],n,1,r,z;
 8
 9
          double k[N],b[N];
          inline void r1(int&x){x=(x+z-1)%39989+1;}
10
11
          inline void r2(int&x){x=(x+z-1)%1000000000+1;}
12
          double f(int o,int x){
13
              return k[o]*x+b[o];
14
15
          int beat(int x,int a,int b){
              double u=f(a,x),v=f(b,x);
return fabs(u-v)<=1e-8?a<b:u>v;
16
17
18
19
          void add(int o,int L,int R,int x){
20
              int md=L+R>>1;
\overline{21}
               if(1<=L&&R<=r){</pre>
22
                    if(!v[o])return (void)(v[o]=x);
                   if(beat(L,v[o],x) && beat(R,v[o],x))return;
if(beat(L,x,v[o]) && beat(R,x,v[o]))return (void)(v[o]=x);
23
24
25
                    if(beat(md,x,v[o]))swap(x,v[o]);
```

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```
\frac{26}{27}
                  if(beat(L,x,v[o]))add(lc,L,md,x);
                  else add(rc,md+1,R,x);
28
                  return;
29
30
              if(r>md)add(rc,md+1,R,x);
31
              if(l<=md)add(lc,L,md,x);</pre>
32
33
         int ask(int o,int L,int R){
34
35
36
37
              if(L==R)return v[o];
              int md=L+R>>1,h=1<=md?ask(lc,L,md):ask(rc,md+1,R);</pre>
             return beat(1,h,v[o])?h:v[o];
38
         void work(){
39
              cin>>n;
40
              for(int op,y1,y2,c=0;n--;){
41
                  cin>>op;
42
                  if(op){
43
                       cin>>l>>y1>>r>>y2,++c,r1(l),r2(y1),r1(r),r2(y2);
44
                       if(l==r)k[c]=0,b[c]=max(y1,y2);
45
                       else {
                           if(l>r)swap(l,r),swap(y1,y2);
k[c]=(y2-y1+0.)/(r-1),b[c]=y1-k[c]*1;
46
47
48
49
                       add(1,1,4e4+10,c);
50
                  }else cin>>1,r1(1),cout<<(z=ask(1,1,4e4+10))<<'\n';
51
             }
52
53
54
    int main(){
55
         return Acc::work(),0;
56
    }
```

# 3.14 ODT

```
/* ODT @ wrb */
 23
     struct T{
           int 1,r,v;
 4
          T(int a, int b=-1, int c=-1):l(a),r(b),v(c){}
 5
          bool operator<(const T&_)const{return 1<_.1;}</pre>
 \frac{\tilde{6}}{7}
     };
     set<T>s;
 8 9
     auto spl(int p){
          auto it=s.lower_bound(p);
10
          if(it!=end(s) && it->l==p)return it;
11
           --it;
12
          int l=it->1,r=it->r,v=it->v;
13
          s.erase(it),s.insert(T(1,p-1,v));
14
          return s.insert(T(p,r,v)).first;
15
16
     void asgn(int 1,int r,int v){
          auto ed=spl(r+1),bg=spl(1);
17
18
          s.erase(bg,ed);
          auto i=s.insert(T(1,r,v)).first,j=prev(i);
if(i!=begin(s)&&j->v==v)l=j->1,s.erase(j);
if((j=next(i))!=end(s)&&j->v==v)r=j->r,s.erase(j);
19
20
\overline{21}
22
           s.erase(i),s.insert(T(1,r,v));
\overline{23}
     }
```

# 4 string

# 4.1 kmp

```
/* kmp */
auto kmp = [&](const std::string& s) -> vi {
   int n = s.length();
   vi next(n);
   for (int i = 1; i < n; i++) {
        int j = next[i - 1];
        while (j > 0 and s[i] != s[j]) j = next[j - 1];
        if (s[i] == s[j]) j++;
        next[i] = j;
   }
   return next;
};
```

4.2 z function 25

# 4.2 z function

```
/* exkmp */
      auto exkmp = [&](const std::string& s) -> vi {
 3
           int n = s.size();
 4
           vi z(n);
           for (int i = 1, l = 0, r = 0; i < n; i++) {
    if (i <= r and z[i - 1] < r - i + 1) {
        z[i] = z[i - 1];
 5
 6
 8
                 } else {
                      z[i] = std::max(0, r - i + 1);
while (z[i] + i < n \text{ and } s[z[i]] == s[z[i] + i]) z[i]++;
 9
10
11
12
                 if (z[i] + i - 1 > r) {
13
                       1 = i;
                      r = z[i] + i - 1;
14
15
16
17
           return z;
      };
```

#### 4.3 manacher

```
/* manacher @ wrb */
auto Manacher = [&](const std::string& t) {
    std::string s = "#";
    for (char c : t) s += c, s += '#';
    int i, o = 0, r = 0, n = s.size();
    std::vector<int> p(n, 1), q(n);
    for (i = 0; i < n; ++i) {
        if (i <= r) p[i] = std::min(r - i + 1, p[2 * o - i]);
        for (; p[i] <= i && s[i + p[i]] == s[i - p[i]]; ++p[i]);
        if (i + p[i] - 1 > r) r = i + p[i] - 1, o = i;
    }
    return p;
}
```

# 4.4 AC automaton

```
/* AC auto */
      int cnt = 0;
const int N = 2e5 + 10;
     static std::array<std::array<int, 26>, N> tr;
static std::array<int, N> exist, fail, ans, point;
 6
      vi order;
 8
      auto insert = [&](const auto& s) {
            int p = 0;
           for (const auto& ch : s) {
  int c = ch - 'a';
  if (!tr[p][c]) tr[p][c] = ++cnt;
10
11
12
13
                 p = tr[p][c];
14
15
            exist[p]++;
16
           return p;
     };
17
18
19
      auto build = [&]() {
           std::queue<int> q;
for (int i = 0; i < 26; i++) {
    if (tr[0][i]) q.push(tr[0][i]);</pre>
20
21
22
23
24
            while (!q.empty()) {
25
                 auto u = q.front();
26
                 q.pop();
                 forder.push_back(u);
for (int i = 0; i < 26; i++) {
    if (tr[u][i]) {</pre>
27
28
29
                             fail[tr[u][i]] = tr[fail[u]][i];
30
31
                             q.push(tr[u][i]);
                       } else {
32
                             tr[u][i] = tr[fail[u]][i];
33
34
                 }
35
36
           }
    };
```

26 4 STRING

```
39
     auto query = [&](const auto& s) {
         int p = 0;
for (const auto ch : s) {
40
41
42
              p = tr[p][ch - 'a'];
43
              ans[p]++;
44
45
         return;
46
    };
47
     void solve (){
48
         for (int i = 0; i < n; i++) {
49
             point[i] = insert(t);
50
51
         build();
         query(s);
/* fail 树上子树求和 */
53
54
         reverse(all(order));
for (const auto& i : order) ans[fail[i]] += ans[i];
55
56
```

# 4.5 PAM

```
/* PAM @ ddl */
     std::vector<node> tr;
 3
     std::vector<int> stk;
 4
     auto newnode = [&](int len) {
          tr.emplace_back();
 6
7
          tr.back().len = len;
         return (int) tr.size() - 1;
 8
    auto PAMinit = [&]() {
   newnode(0), tr.back().fail = 1;
   newnode(-1), tr.back().fail = 0;
 9
10
11
12
          stk.push_back(-1);
13
14
    PAMinit();
15
     auto getfail = [&](int v) {
16
          while (stk.end()[-2 - tr[v].len] != stk.back()) {
              v = tr[v].fail;
18
19
          return v;
20
    };
\overline{21}
     auto insert = [&](int last, int c, int cnt) {
22
          stk.emplace_back(c);
\overline{23}
          int x = getfail(last);
\frac{23}{24} \frac{25}{25}
          if (!tr[x].ch[c]) {
               int u = newnode(tr[x].len + 2);
\frac{26}{27}
               tr[u].fail = tr[getfail(tr[x].fail)].ch[c];
               tr[x].ch[c] = u;
28
               /* tr[u].size = tr[tr[u].fail].size + 1; */
29
               /* Can be used to count the number of types of palindromic strings ending at the current
30
                * position */
31
32
          tr[tr[x].ch[c]].size += cnt;
\begin{array}{c} 33 \\ 34 \\ 35 \\ 36 \\ 37 \end{array}
          return tr[x].ch[c];
     auto build = [&]() { /* DP on fail tree */
          int ans = 0;
for (int i = (int) tr.size() - 1; i > 1; i--) {
38
               tr[tr[i].fail].size += tr[i].size;
39
               /* options */
40
41
          return ans;
42
    j};
     /* PAM */
43
    int ans = 0, last = 0;
for (int i = 0; i < n; i++) {
44
45
          last = insert(last, s[i] - 'a', 1);
46
```

4.6 Suffix Array 27

```
11
          vector<int> bl;
12
          size_t count() const {
13
               return t.size() - 2;
14
15
          const T& operator[](const size_t& p) const {
16
               return t[p];
17
18
          const T& ask(const size_t& p) const {
19
               return t[bl[p]];
20
21
          int gf(int o, int p) {
\frac{22}{23}
               while (p - t[o].len - 1 < 0 \mid | s[p - t[o].len - 1] != s[p]) o = t[o].fa;
24
25
          void append(int c) {
               int p = s.size(), o;
s += c, o = gf(las, p);
if (t[o].ch[c] == 0) {
26
27
28
\frac{1}{29}
                     t.emplace_back();
                     t.back().len = t[o].len + 2;
t.back().fa = t[gf(t[o].fa, p)].ch[c];
30
31
                     t.back().d = t[t.back().fa].d + 1;
32
33
                     t[o].ch[c] = t.size() - 1;
34
35
               bl.emplace_back(las = t[o].ch[c]);
36
          PAM(): las(), s(), t(2) {
   t[0].fa = t[1].fa = 1, t[1].len = -1;
37
38
39
          PAM(const string& str, int h) : las(), s(), t(2) {
   t[0].fa = t[1].fa = 1, t[1].len = -1;
40
41
42
               for (char c : str) append(c - h);
43
44
     };
```

### 4.6 Suffix Array

```
/* suffix array and ST table @ jiangly */
auto suffixArray = [&](const std::string& s) {
 3
           int n = s.length();
           vi sa(n), rk(n);
 4
           std::iota(all(sa), 0);
std::sort(all(sa), [&](int a, int b) { return s[a] < s[b]; });</pre>
 5
 6
 7
           rk[sa[0]] = 0;
 8
           for (int i = 1; i < n; ++i) {</pre>
 9
                rk[sa[i]] = rk[sa[i - 1]] + (s[sa[i]] != s[sa[i - 1]]);
10
11
           int k = 1;
12
           vi tmp(n), cnt(n);
13
           tmp.reserve(n);
           while (rk[sa[n-1]] < n-1) {
14
15
                 tmp.clear();
                 for (int i = 0; i < k; ++i) tmp.push_back(n - k + i);
for (const auto& i : sa) {
    if (i >= k) tmp.push_back(i - k);
16
17
18
19
20
                 std::fill(all(cnt), 0);
21
                 for (int i = 0; i < n; i++) cnt[rk[i]]++;
for (int i = 1; i < n; i++) cnt[i] += cnt[i - 1];
for (int i = n - 1; i >= 0; i--) sa[--cnt[rk[tmp[i]]]] = tmp[i];
22
23
24
                 std::swap(rk, tmp);
25
                 rk[sa[0]] = 0;
26
                 for (int i = 1; i < n; i++) {
                      rk[sa[i]] = rk[sa[i - 1]] + (tmp[sa[i - 1]] < tmp[sa[i]] or sa[i - 1] + k == n or tmp[sa[i - 1] + k] < tmp[sa[i] + k]);
27
28
29
30
                k *= 2;
31
           vi height(n);
for (int i = 0, j = 0; i < n; ++i) {
    if (rk[i] == 0) continue;</pre>
32
33
34
                if (j) --j;
while (s[i + j] == s[sa[rk[i] - 1] + j]) ++j;
height[rk[i]] = j;
35
36
37
38
39
           return std::make_tuple(sa, rk, height);
40
     auto [sa, rk, height] = suffixArray(s);
41
42
     vvi f(n, vi(30, inf));
43
      vi Log2(n);
44
      auto init = [&]() -> void {
           for (int i = 0; i < n; i++) {</pre>
```

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```
46
                   f[i][0] = height[i];
47
                   if (i > 1) Log2[i] = Log2[i / 2] + 1;
48
49
             int t = Log2.back();
            for (int j = 1; j <= t; j++) {
    for (int i = 0; i <= n - (1 << j); i++) {
        f[i][j] = std::min(f[i][j - 1], f[i + (1 << (j - 1))][j - 1]);
50
51
52
53
54
            }
55
      };
56
      init();
      auto query = [&](int 1, int r) -> int {
   int t = Log2[r - 1 + 1];
   return std::min(f[1][t], f[r - (1 << t) + 1][t]);</pre>
57
58
59
60
61
      auto lcp = [&](int i,
                                        <u>int</u> j) {
            i = rk[i], j = rk[j];
if (i > j) std::swap(i, j);
62
63
64
            return query(i + 1, j);
      };
65
```

```
/* suffix array @ wrb */
                    auto SA = [](std::string s) {
                                  int n = s.size(), m = 128, i, j, l;
std::vector<int> ct(m), sa(n), rk(n), h(n), a(n);
for (i = 0; i < n; ++i) ++ct[rk[i] = s[i]];
for (i = 1; i < m; ++i) ct[i] += ct[i - 1];
for (i = n - 1; ~i; --i) sa[--ct[rk[i]]] = i;
for (l = 1; l < n; l *= 2) {
   for (j = 0, i = n - 1; i >= n - 1; --i) a[j++] = i;
   for (i = 0; i < n; ++i) if (sa[i] >= 1) a[j++] = sa[i] - 1;
   ct = std::vector<int>(m):
    \overline{3}
    4
    5
    6
    8
    9
10
                                                       ct = std::vector<int>(m);
11
                                                     for (i = 0; i < n; ++i) ++ct[rk[a[i]]];
for (i = 1; i < m; ++i) ct[i] += ct[i - 1];</pre>
12
13
14
                                                       for (i = n - 1; ~i; --i) sa[--ct[rk[a[i]]]] = a[i];
15
                                                       std::swap(rk, a), rk[sa[0]] = 0;
                                                     for (i = 1; i < n; ++i) {
    rk[sa[i]] = rk[sa[i - 1]] + (a[sa[i]] != a[sa[i - 1]] || a[(sa[i] + 1) % n] != a[(sa[i - 1] + 1) % n] || a[(sa[i] + 1) % n] 
16
17
                                                                                         1) % n]);
18
                                                     if ((m = rk[sa[n - 1]] + 1) == n) break;
19
20
\overline{21}
                                    for (i = j = 0; i + 1 < n; h[rk[i++]] = j) {
\frac{1}{22}
                                                     for (j? --j: 0; s[i+j] == s[sa[rk[i] - 1] + j]; ++j);
23
24
25
26
                                    sa.erase(sa.begin());
                                     rk.erase(rk.begin());
                                    h.erase(h.begin());
27
                                    return make_tuple(sa, rk, h);
28
                  //h[i] : LCP(rk[i], rk[i - 1])
```

#### 4.7 Cantor expansion

```
/* Cantor expresion @ wrb */
    std::cin >> n, fac[0] = 1;
    for (int i = 1; i <= n; ++i) {
    std::cin >> a[i];
 3
 4
 5
         fac[i] = 111 * fac[i - 1] * i % P;
 \frac{\tilde{6}}{7}
    auto ins = [&](int x) {
         for (; x <= n; x += x & -x) ++t[x];</pre>
 9
10
    auto ask = [\&] (int x) {
         int z = 0;
11
         for (; x; x ^= x & -x) z += t[x];
12
13
         return z;
14
15
    int z = 0;
16
    for (int i = n; i; --i) {
17
         z = (z + 111 * fac[n - i] * ask(a[i])) % P;
18
         ins(a[i]);
19
    std::cout << ++z << '\n';
20
```

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#### 4.8 trie

# 普通字典树 (单词匹配)

```
/* trie */
    int cnt;
 3
     std::vector<std::array<int, 26>> trie(n + 1);
     vi exist(n + 1);
     auto insert = [&](const std::string& s) -> void {
         int p = 0;
 6
         for (const auto ch : s) {
   int c = ch - 'a';
   if (!trie[p][c]) trie[p][c] = ++cnt;
 8
 9
10
              p = trie[p][c];
11
12
          exist[p] = true;
13
    };
14
     auto find = [&](const string& s) -> bool {
         int p = 0;
15
         for (const auto ch : s) {
   int c = ch - 'a';
16
17
              if (!trie[p][c]) return false;
18
19
              p = trie[p][c];
20
21
         return exist[p];
     };
```

## 01 字典树 (求最大异或值)

给定 n 个数, 取两个数进行异或运算, 求最大异或值.

```
/* trie */
int cnt = 0;
 3
        std::vector\langlestd::array\langleint, 2>> trie(N);
auto insert = [&](int x) -> void {
 4
                int p = 0;

for (int i = 30; i >= 0; i--) {

   int c = (x >> i) & 1;

   if (!trie[p][c]) trie[p][c] = ++cnt;
 5
 6
 7
  8
 9
                        p = trie[p][c];
10
11
        };
        auto find = [&](int x) -> int {
12
                int sum = 0, p = 0;

for (int i = 30; i >= 0; i--) {

   int c = (x >> i) & 1;

   if (trie[p][c ^ 1]) {

      p = trie[p][c ^ 1];

      cum += (1 << i);
13
14
15
16
17
18
                                sum += (1 << i);
19
                        } else {
                               p = trie[p][c];
20
21
22
\overline{23}
                return sum;
        };
```

# 字典树合并

来自浙大城市学院 2023 校赛 E 题.

给定一棵根为 1 的树,每个点的点权为  $w_i$ . 一共 q 次询问,每次给出一对 u,v,询问以 v 为根的子树上的点与 u 的权值最大异或值.

```
int main() {
 2
         std::ios::sync_with_stdio(false);
3
         std::cin.tie(0);
 4
         int n, m;
std::cin >> n;
5
6
7
         vi w(n + 1);
8
         for (int i = 1; i <= n; i++) std::cin >> w[i];
9
         vvi e(n + 1);
10
         for (int i = 1, u, v; i < n; i++) {</pre>
```

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```
11
                 std::cin >> u >> v;
12
                 e[u].push_back(v);
13
                 e[v].push_back(u);
14
15
           // 离线询问 //
16
            std::cin >> m;
17
18
           std::vector<vpi> q(n + 1);
19
            vi ans(m + 1);
20
21
           for (int i = 1; i <= m; i++) {</pre>
                 int u, v;
22 \\ 23 \\ 24 \\ 25
                 std::cin >> u >> v;
                 q[v].emplace_back(u, i);
26
           // 01 trie //
27
           std::vector<std::array<int, 2>> tr(1);
28
           auto new_node = [&]() -> int {
   tr.emplace_back();
\overline{29}
30
31
32
                 return tr.size() - 1;
            vi id(n + 1);
33
34
           auto insert = [&](int root, int x) {
                 int p = root;
for (int i = 29; i >= 0; i--) {
   int c = x >> i & 1;
35
36
                       if (!tr[p][c]) tr[p][c] = new_node();
p = tr[p][c];
37
38
39
                 }
           };
40
41
           auto query = [&](int root, int x) -> int {
                 int ans = 0, p = root;
for (int i = 29; i >= 0; i--) {
   int c = x >> i & 1;
   if (tr[p][c ^ 1]) {
      p = tr[p][c ^ 1];
      ans += (1 << i);
}</pre>
42
43
44
45
46
47
48
                       } else {
49
                            p = tr[p][c];
50
51
                 }
52
53
54
                 return ans;
           std::function<int(int, int)> merge = [&](int a, int b) -> int {
55
                 // b 的信息挪到 a 上 //
                 if (!a) return b;
if (!b) return a;
tr[a][0] = merge(tr[a][0], tr[b][0]);
tr[a][1] = merge(tr[a][1], tr[b][1]);
56
57
58
59
60
                 return a;
61
62
            std::function<void(int, int)> dfs = [&](int u, int fa) {
63
                 id[u] = new_node();
                 insert(id[u], w[u]);
for (auto v : e[u]) {
   if (v == fa) continue;
64
65
66
67
                       dfs(v, u);
68
                       id[u] = merge(id[u], id[v]);
69
                 for (auto [v, i] : q[u]) {
   ans[i] = query(id[u], w[v]);
70 \\ 71 \\ 72 \\ 73 \\ 74 \\ 75 \\ 76 \\ 77
                 }
           dfs(1, 0);
           for (int i = 1; i <= m; i++) std::cout << ans[i] << endl;</pre>
           return 0;
     }
```

# 5 math - number theory

#### $5.1 \mod int$

```
template <int P>
      struct Mint {
 3
            int v = 0;
 4
 5
            // reflection //
 6
            template <typet = int>
            constexpr operator T() const {
 7
 8
                 return v;
 9
10
            // constructor //
constexpr Mint() = default;
template <typet>
11
12
13
            constexpr Mint(T x) : v(x % P) {}
constexpr int val() const { return v; }
14
15
16
            constexpr int mod() { return P; }
17
18
19
            friend std::istream& operator>>(std::istream& is, Mint& x) {
20
                 LL y;
                 is >> y;
21
22
                 x = y;
\overline{23}
                 return is;
24
25
            friend std::ostream& operator<<(std::ostream& os, Mint x) { return os << x.v; }</pre>
26
27
            // comparision //
28
            friend constexpr bool operator == (const Mint& lhs, const Mint& rhs) { return lhs.v == rhs.v; }
           friend constexpr bool operator!=(const Mint& Ins, const Mint& Ins) { return Ins.v != rhs.v; } friend constexpr bool operator!=(const Mint& Ihs, const Mint& rhs) { return Ihs.v != rhs.v; } friend constexpr bool operator<=(const Mint& Ihs, const Mint& rhs) { return Ihs.v <= rhs.v; } friend constexpr bool operator>=(const Mint& Ihs, const Mint& rhs) { return Ihs.v > rhs.v; } friend constexpr bool operator>(const Mint& Ihs, const Mint& rhs) { return Ihs.v > rhs.v; }
29
30
31
32
33
            friend constexpr bool operator>=(const Mint& lhs, const Mint& rhs) { return lhs.v >= rhs.v; }
34
35
            // arithmetic //
36
            template <typet>
37
            friend constexpr Mint power(Mint a, T n) {
38
                 Mint ans = 1;
39
                 while (n) {
40
                       if (n & 1) ans *= a;
41
                       a *= a;
42
                       n >>= 1:
43
                 return ans;
44
45
46
            friend constexpr Mint inv(const Mint& rhs) { return power(rhs, P - 2); }
            friend constexpr Mint operator+(const Mint& lhs, const Mint& rhs) {
   return lhs.val() + rhs.val() < P ? lhs.val() + rhs.val() : lhs.val() - P + rhs.val();</pre>
47
48
49
50
            friend constexpr Mint operator-(const Mint& lhs, const Mint& rhs) {
                 return lhs.val() < rhs.val() ? lhs.val() + P - rhs.val() : lhs.val() - rhs.val();</pre>
51
52
            friend constexpr Mint operator*(const Mint& lhs, const Mint& rhs) {
   return static_cast<LL>(lhs.val()) * rhs.val() % P;
53
54
55
           friend constexpr Mint operator/(const Mint& lhs, const Mint& rhs) { return lhs * inv(rhs); }
Mint operator+() const { return *this; }
Mint operator-() const { return Mint() - *this; }
56
57
58
59
            constexpr Mint& operator++() {
60
                 if (v == P) v = 0;
61
62
                 return *this;
63
64
            constexpr Mint& operator--() {
65
                 if (v == 0) v = P;
66
67
                 return *this;
68
69
70
71
            constexpr Mint& operator++(int) {
   Mint ans = *this;
                 ++*this;
72
73
                 return ans;
74
75
            constexpr Mint operator--(int) {
                 Mint ans = *this;
                  --*this;
76
77
                 return ans;
78
79
            constexpr Mint& operator+=(const Mint& rhs) {
```

```
80
             v = v + rhs;
81
            return *this;
82
83
        constexpr Mint& operator-=(const Mint& rhs) {
84
             v = v - rhs;
            return *this;
86
87
        constexpr Mint& operator*=(const Mint& rhs) {
88
             v = v * rhs;
89
             return *this;
90
91
        constexpr Mint& operator/=(const Mint& rhs) {
92
             v = v / rhs;
return *this;
93
94
    };
95
    using Z = Mint<998244353>;
96
```

#### 5.2 Eculid

#### 欧几里得算法

```
1 std::gcd(a, b)
```

# 扩展欧几里得算法

```
/* exgcd */
 3
      auto exgcd = [&](LL a, LL b, LL& x, LL& y) {
   LL x1 = 1, x2 = 0, x3 = 0, x4 = 1;
   while (b != 0) {
 4
                 LL c = a / b;
 5
 6
7
                 std::tie(x1, x2, x3, x4, a, b) = std::make_tuple(x3, x4, x1 - x3 * c, x2 - x4 * c, b, a - b * c);
 9
           x = x1, y = x2;
10
     auto exgcd = [&](auto&& self, LL a, LL b, LL& x, LL& y) {
    if (!b) {
        x = 1, y = 0;
    }
11
12
13
14
                 return a;
15
16
           LL d = self(self, b, a % b, y, x);
17
           y -= a / b * x;
18
           return d;
19
     };
```

```
1 auto exgcd = [&](auto&& self, LL a, LL b, LL& x, LL& y) {
2     if (!b) {
        x = 1, y = 0;
        return;
     }
6     self(self, b, a % b, y, x);
     y -= a / b * x;
};
```

```
1 auto exgcd = [&](auto&& self, LL a, LL b, LL& x, LL& y) {
2     if (!b) {
3         x = 1, y = 0;
4         return a;
5     }
6     LL d = self(self, b, a % b, y, x);
7     y -= a / b * x;
8     return d;
9 };
```

# 类欧几里得算法

```
一般形式: 求 f(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor
```

5.3 inverse 33

```
f(a, b, c, n) = nm - f(c, c - b - 1, a, m - 1)
```

```
LL f(LL a, LL b, LL c, LL n) {
    if (a == 0) return ((b / c) * (n + 1));
    if (a >= c || b >= c)
        return f(a % c, b % c, c, n) + (a / c) * n * (n + 1) / 2 + (b / c) * (n + 1);
    LL m = (a * n + b) / c;
    LL v = f(c, c - b - 1, a, m - 1);
    return n * m - v;
}
```

```
更进一步, 求: g(a,b,c,n) = \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor 以及 h(a,b,c,n) = \sum_{i=0}^{n} \lfloor \frac{ai+b}{c} \rfloor^2 g(a,b,c,n) = \lfloor \frac{mn(n+1)-f(c,c-b-1,a,m-1)-h(c,c-b-1,a,m-1)}{2} \rfloor h(a,b,c,n) = nm(m+1) - 2f(c,c-b-1,a,m-1) - 2g(c,c-b-1,a,m-1) - f(a,b,c,n)
```

```
const int inv2 = 499122177;
const int inv6 = 166374059;
 3
 4
        LL f(LL a, LL b, LL c, LL n);
       LL g(LL a, LL b, LL c, LL n);
LL h(LL a, LL b, LL c, LL n);
 5
 6
        struct data {
 9
             LL f, g, h;
10
11
        data calc(LL a, LL b, LL c, LL n) {
    LL ac = a / c, bc = b / c, m = (a * n + b) / c, n1 = n + 1, n21 = n * 2 + 1;
13
14
               data d;
15
               if (a == 0) {
                     d.f = bc * n1 \% mod;
16
                     d.g = bc * n % mod * n1 % mod * inv2 % mod;
d.h = bc * bc % mod * n1 % mod;
17
18
19
                     return d;
20
\frac{1}{21}
               if (a >= c || b >= c) {
22
                     d.f = n * n1 \% mod * inv2 \% mod * ac % mod + bc * n1 % mod;
23
                     ac * n % mod * n1 % mod * n21 % mod * inv6 % mod + bc * n % mod * n1 % mod * inv2 % mod;
d.h = ac * ac % mod * n % mod * n1 % mod * n21 % mod * inv6 % mod +
bc * bc % mod * n1 % mod + ac * bc % mod * n % mod * n1 % mod;
24
25
26
                     d.f %= mod, d.g %= mod, d.h %= mod;
data e = calc(a % c, b % c, c, n);
d.h += e.h + 2 * bc % mod * e.f % mod + 2 * ac % mod * e.g % mod;
d.g += e.g, d.f += e.f;
27
29
30
31
32
                     d.f %= mod, d.g %= mod, d.h %= mod;
                     return d;
\begin{array}{c} 33 \\ 34 \end{array}
              data e = calc(c, c - b - 1, a, m - 1);

d.f = n * m % mod - e.f, d.f = (d.f % mod + mod) % mod;

d.g = m * n % mod * n1 % mod - e.h - e.f, d.g = (d.g * inv2 % mod + mod) % mod;

d.h = n * m % mod * (m + 1) % mod - 2 * e.g - 2 * e.f - d.f;
35
36
37
38
               d.h = (d.h \% mod + mod) \% mod;
39
```

#### 5.3 inverse

#### 线性递推

$$a^{-1} \equiv -\lfloor \frac{p}{a} \rfloor \times (p\%a)^{-1}$$

#### 求 n 个数的逆元

```
/* inverse */
 2
3
      auto inverse =[&](const vi& a) {
            int n = a.size();
 4
            vi b(n), f(n), ivf(n);
            f[0] = a[0];

for (int i = 1; i < n; i++) {

f[i] = 111 * f[i - 1] * a[i] % p;
 5
 6
7
 8
 9
            ivf.back() = quick_power(f.back(), p - 2, p);
for (int i = n - 1; i; i--) {
   ivf[i - 1] = 111 * ivf[i] * a[i] % p;
10
11
12
            b[0] = ivf[0];
for (int i = 1; i < n; i++) {
13
14
15
                  b[i] = 111 * ivf[i] * f[i - 1] % p;
16
17
            return b;
18
     };
```

## 5.4 sieve

#### 素数

```
vi prime, is_prime(n + 1, 1);
auto Euler_sieve = [&] (int n) {
    for (int i = 2; i <= n; i++) {
        if (is_prime[i]) prime.push_back(i);
        for (auto p : prime) {
            if (i * p > n) break;
            is_prime[i * p] = 0;
            if (i % p == 0) break;
        }
    }
}
```

#### 欧拉函数

```
\begin{array}{c}2\\3\\4\\5\\6\\7\end{array}
 8
                        prime.push_back(i);
                        phi[i] = i - 1;
10
                  for (auto p : prime) {
   if (i * p > n) break;
   is_prime[i * p] = 0;
   if (i % p) {
11
12
13
14
15
                              phi[i * p] = phi[i] * phi[p];
16
                        } else {
                              phi[i * p] = phi[i] * p;
break;
17
18
19
20
                        }
                  }
\overline{21}
            }
22
      };
```

#### 约数和

```
vi g(n + 1), d(n + 1), prime;
vi is_prime(n + 1, 1);
auto get_d = [&](int n) {
   int tot = 0;
   g[1] = d[1] = 1;
   for (int i = 2; i <= n; i++) {
      if (is_prime[i]) {</pre>
```

5.4 sieve 35

```
8
9
                                 prime.push_back(i);
                                 d[i] = g[i] = i + 1;
10
                        for (auto p : prime) {
    if (i * p > n) break;
    is_prime[i * p] = 0;
    if (i % p == 0) {
        g[i * p] = g[i] * p + 1;
        d[i * p] = d[i] / g[i] * g[i * p];
        break;
11
12
13
14
15
16
17
                                         break;
                                 } else {
18
                                         d[i * p] = d[i] * d[p];
g[i * p] = 1 + p;
19
20
21
22
23
                }
        };
```

# 莫比乌斯函数

```
vi mu(n + 1), prime;
       vi is_prime(n + 1, 1);
       auto get_mu = [&](int n) {
 4
              mu[1] = 1;
              for (int i = 2; i <= n; i++) {
    if (is_prime[i]) {</pre>
 5
 6
7
                            prime.push_back(i);
mu[i] = -1;
 9
                     for (auto p : prime) {
    if (i * p > n) break;
    is_prime[i * p] = 0;
    if (i % p == 0) {
        mu[i * p] = 0;
        break;
    }
}
10
11
12
13
14
15
                                    break;
16
17
                            mu[i * p] = -mu[i];
                     }
19
20
       };
```

# 杜教筛

```
const int N = 1e7;
         vi mu(N + 1), phi(N + 1), prime;
         vl sum_phi(N + 1), sum_mu(N + 1);
vi is_prime(N + 1, 1);
std::map<LL, LL> mp_mu;
 6
         /* 计算 1 ~ 10<sup>7</sup> 的 mu */
auto get_mu = [&](int n) {
    phi[1] = mu[1] = 1;
    for (int i = 2; i <= n; i++) {
        if (is_prime[i]) {
 9
10
11
                                     prime.push_back(i);
phi[i] = i - 1;
mu[i] = -1;
12
13
14
15
                           for (auto p : prime) {
    if (i * p > n) break;
    is_prime[i * p] = 0;
    if (i % p == 0) {
        phi[i * p] = phi[i] * p;
        mu[i * p] = 0;
        break;
}
16
18
19
20
\overline{21}
22
                                               break;
23
\frac{1}{24}
                                     phi[i * p] = phi[i] * phi[p];
mu[i * p] = -mu[i];
\frac{1}{25}
26
27
                   }
28
         };
        get_mu(N);
for (int i = 1; i <= N; i++) {
    sum_phi[i] = sum_phi[i - 1] + phi[i];
    sum_mu[i] = sum_mu[i - 1] + mu[i];</pre>
29
30
31
32
33
34
       /* 杜教筛: 求 mu 的前缀和 */
```

```
std::function<LL(LL)> S_mu = [&](LL x) -> LL {
37
            if (x <= N) return sum_mu[x];</pre>
38
            auto it = mp_mu.find(x);
39
            if (it != mp_mu.end()) return mp_mu[x];
           LL ans = 1;

for (LL i = 2, j; i <= x; i = j + 1) {

   j = x / (x / i);

   ans -= S_mu(x / i) * (j - i + 1);
40
41
42
43
44
45
            return mp_mu[x] = ans;
46
      };
47
      /* 杜教筛: 求 phi 的前缀和 */
auto S_phi = [&] (LL x) -> LL {
48
49
            if (x <= N) return sum_phi[x];
LL ans = 0;
50
51
52
53
54
55
            for (LL i = 1, j; i <= x; i = j + 1) {
    j = x / (x / i);
    ans += 1ll * (S_mu(j) - S_mu(i - 1)) * (x / i) * (x / i);
56
            return (ans - 1) / 2 + 1;
      };
```

# 5.5 powerful number

目标: 求积性函数 f(n) 的前缀和. 做法如下:

- 1. 构造积性函数 g(n),满足其易求前缀和且素数处函数值等于 f 的函数值.
- 2. 构造 h = f/g, 即 f = h\*g (狄利克雷卷积), 容易知道 h(1) = 1. 容易计算出 h 在非 powerful number 处函数值均为 0.
- 3. 根据

$$F(n) = \sum_{i=1}^{n} f(n)$$

$$= \sum_{i=1}^{n} \sum_{d|i} g(i)h(i/d)$$

$$= \sum_{d=1}^{n} \sum_{i=1}^{\left\lfloor \frac{n}{d} \right\rfloor} h(d)g(i)$$

$$= \sum_{d=1}^{n} h(d)G\left(\left\lfloor \frac{n}{d} \right\rfloor\right)$$

$$= \sum_{d=1,2,\cdots,n,d \text{ is powerful number}} h(d)G\left(\left\lfloor \frac{n}{d} \right\rfloor\right)$$

发现只需要计算 h 在 powerful number 处的函数值即可, 可以边搜索边计算.

给定  $f(p^k) = p^k(p^k - 1)$  为积性函数, 计算其前缀和.

```
auto powerfulNumber = [&](LL n) {

  \begin{array}{c}
    2 \\
    3 \\
    4 \\
    5 \\
    6 \\
    7 \\
    8 \\
    9
  \end{array}

              /* 1. construct g, and compute G */
             /* int m = sqrt(n); // maybe TLE // */
int m = 2e6;
             std::vector<Z> Gs(m + 1);
              vi prime, is_prime(m + 1, 1);
             for (int i = 2; i <= m; i++) {
   if (is_prime[i]) {
     prime.push_back(i);
     Gs[i] = i - 1;
}</pre>
10
11
12
                    13
14
15
16
                                 Gs[i * p] = Gs[i] * Gs[p];
17
                          } else {
```

5.6 block 37

```
19
                               Gs[i * p] = Gs[i] * p;
20
21
22
                  }
23
            for (int i = 2; i <= m; i++) {
   Gs[i] = Gs[i - 1] + Z(i) * Gs[i];
24
25
26
            std::map<LL, Z> mp;
auto G = [&](auto&& self, LL n) {
27
28
29
                  if (n <= m) return Gs[n];</pre>
30
                   if (mp.find(n) != mp.end()) return mp[n];
                  Z ans = Z(n) * Z(n + 1) * Z(n * 2 + 1) * inv6;

for (LL 1 = 2, r, k; 1 <= n; 1 = r + 1) {

k = n / 1, r = n / (n / 1);
31
32
33
34
                        ans -= (Z(r) * Z(r + 1) - Z(1 - 1) * Z(1)) * inv2 * self(self, k);
35
36
                  return mp[n] = ans;
37
            };
/* 2. compute h(p^c) */
38
39
            vvl ps(prime.size());
            std::vector<std::vector<Z>> hs(prime.size());
40
            int len = 0;
41
            for (int i = 0; i < prime.size(); i++) {
   LL p = prime[i], now = p * p, c = 2;
   ps[i] = {1, p}, hs[i] = {1, 0};
   while (now <= n) {
42
43
44
45
                        ps[i].push_back(now);
46
                        Z ans = Z(ps[i][c]) * (Z(ps[i][c]) - 1);
for (int j = 1; j <= c; j++) {
    ans -= Z(ps[i][j]) * Z(ps[i][j - 1]) * Z(p - 1) * hs[i][c - j];</pre>
47
48
49
50
                        hs[i].push_back(ans);
51
52
                        now *= p, c += 1;
53
54
                  len += ps[i].size();
55
56
            debug(len);
57
            /* 3. search powerful number */
58
            Z ans = 0;
59
            auto dfs = [&](auto&& self, int id, LL now, Z hd) -> void {
                  ans += hd * G(G, n / now);
for (int i = id; i < prime.size(); i++) {</pre>
60
61
                        int p = prime[i], c = 2;
if (now > n / p / p) break;
for (LL x = now * p * p; x <= n; x *= p, c++) {
    if (hs[i][c]) self(self, i + 1, x, hd * hs[i][c]);
}</pre>
62
63
64
65
66
67
                  }
68
69
            dfs(dfs, 0, 1, 1);
70
            return ans;
      };
```

### 5.6 block

### 分块的逻辑

下取整  $\lfloor \frac{n}{q} \rfloor = k$  的分块  $()g \leq n)$ 

```
1  for(int l = 1, r, k; l <= n; l = r + 1){
2     k = n / 1;
3     r = n / (n / 1);
4     debug(l, r, k);
}</pre>
```

 $k = \lfloor \frac{n}{q} \rfloor$  从大到小遍历  $\lfloor \frac{n}{q} \rfloor$  的所有取值, [l, r] 对应的是 g 取值的区间.

```
for(int l = 1, r, k; l < n; l = r + 1){
    k = (n + 1 - 1) / l;
    r = (n + k - 2) / (k - 1) - 1;
    debug(l, r, k);
}</pre>
```

 $k = \lceil \frac{n}{q} \rceil$  从大到小遍历  $\lceil \frac{n}{g} \rceil$  的所有取值, [l, r] 对应的是 g 取值的区间.

## 一般形式

计算  $\sum_{i=1}^{n} f(i) \lfloor \frac{n}{i} \rfloor$ , 设 s(i) 为 f(i) 的前缀和。

```
1  for (int l = 1, r; l <= n; l = r + 1) {
2     r = n / (n / 1);
3     ans += (s[r] - s[l - 1]) * (n / 1);
}</pre>
```

 $\sum_{i=1}^{n} f(i) \lfloor \frac{a}{i} \rfloor \lfloor \frac{b}{i} \rfloor$ 

```
for (int 1 = 1, r, r1, r2; 1 <= n; 1 = r + 1) {
    if (a / 1) {
        r1 = a / (a / 1);
    } else {
        r1 = n;
    }
    if (b / 1) {
        r2 = b / (b / 1);
    } else {
        r2 = n;
    }
    remin(min(r1, r2), n);
    ans += (s[r] - s[1 - 1]) * (a / 1) * (b / 1);
}</pre>
```

### 5.7 CRT & exCRT

求解

$$\begin{cases}
N \equiv a_1 \bmod m_1 \\
N \equiv a_2 \bmod m_2 \\
\dots \\
N \equiv a_n \bmod m_n
\end{cases}$$

```
有 N \equiv \sum_{i=1}^{k} a_i \times \operatorname{inv}\left(\frac{M}{m_i}, m_i\right) \times \left(\frac{M}{m_i}\right) \operatorname{mod} M
```

```
1    /* CRT */
2    auto crt = [&](int n, const vi& a, const vi& m) -> LL{
3        LL ans = 0, M = 1;
4        for(int i = 1; i <= n; i++) M *= m[i];
5        for(int i = 1; i <= n; i++){
            ans = (ans + a[i] * inv(M / m[i]) * (M / m[i])) % M;
7        }
8        return (ans % M + M) % M;
9    };</pre>
```

扩展中国剩余定理

```
1    /* exCRT */
2    auto excrt = [&](int n, const vi& a, const vi& m) -> LL{
3        LL A = a[1], M = m[1];
```

 $5.8 \quad BSGS \& exBSGS$  39

```
for (int i = 2; i <= n; i++) {
   LL x, y, d = std::gcd(M, m[i]);
   exgcd(M, m[i], x, y);
   LL mod = M / d * m[i];
   x = x * (a[i] - A) / d % (m[i] / d);
   A = ((M * x + A) % mod + mod) % mod;
   M = mod;
}
return A;
};</pre>
```

### 5.8 BSGS & exBSGS

求解满足  $a^x \equiv b \mod p$  的 x

```
/* BSGS + [&] (LL a, LL b, LL p) {
    if (1 % p == b % p) return 011;
    LL k = std::sqrt(p) + 1;
    std::unordered_map<LL, LL> hash
    for (LL i = 0, j = b % p; i < k; i++) {
        hash[j] = i;
        i = i * a % p;
    }
}</pre>
  2
 3
  4
  5
 6
 8
  9
                         j = j * a % p;
10
                LL ak = 1;
11
                for (int i = 1; i <= k; i++) ak = ak * a % p;
12
                for (int i = 1, j = ak; i <= k; i++) {
    if (hash.count(j)) return 111 * i * k - hash[j];</pre>
13
14
15
                         j = 111 * j * ak % p;
16
17
                 return -INF;
18
        };
```

## $(a,p) \neq 1$ 的情形

```
/* exBSGS */
         /* return value < 0 means no solution */
auto exBSGS = [&](auto&& self, LL a, LL b, LL p) {</pre>
 3
                 b exbsds = [&](altowx self, LL a, b
b = (b % p + p) % p;
if (111 % p == b % p) return 011;
LL x, y, d = std::gcd(a, p);
exgcd(exgcd, a, p, x, y);
if (d > 1) {
  4
 5
  6
  8
                          if (b % d != 0) return -INF;
exgcd(exgcd, a / d, p / d, x, y);
return self(self, a, b / d * x % (p / d), p / d) + 1;
  9
10
11
12
                  return BSGS(a, b, p);
13
        };
14
```

### 5.9 Miller Rabin

```
/* Miller Rabin */
    v1 vv = \{2, 325, 9375,
                         28178, 450775, 9780504, 1795265022};
3
    auto quick_power = [&](LL a, LL n, LL mod) {
4
       LL ans = 1;
5
       while (n) {
6
           if (n & 1) ans = (i128) ans * a % mod;
           a = (i128) a * a % mod;
8
           n >>= 1;
9
10
       return ans;
11
   };
12
   13
14
15
16
17
18
           LL p = quick_power(a % n, d, n);
           int i = s;
while (p != 1 and p != n - 1 and a % n and i--) p = (i128) p * p % n;
19
20
21
           if (p != n - 1 and i != s) return false;
22
23
       return true;
```

24 |};

### 5.10 Pollard Rho

能在  $O(n^{\frac{1}{4}})$  的时间复杂度随机出一个 n 的非平凡因数.

```
/* pollard rho */
    \begin{array}{c}2\\3\\4\\5\\6\\7\end{array}
 8 9
10
                    LL d = std::gcd(val, x);
11
                    if(d > 1) return d;
12
13
14
            LL d = std::gcd(val, x);
15
            if(d > 1) return d;
16
        }
17
    };
```

利用 Miller Rabin 和 Pollard Rho 进行素因数分解

```
auto factorize = [&](LL a) -> v1{
              vl ans, stk;

  \begin{array}{c}
    2 \\
    3 \\
    4 \\
    5 \\
    6 \\
    7 \\
    8 \\
    9
  \end{array}

              for (auto p : prime) {
   if (p > 1000) break;
   while (a % p == 0) {
      ans.push_back(p);
   }
}
                            a /= p;
                     if (a == 1) return ans;
              }
10
11
              stk.push_back(a);
12
              while (!stk.empty()) {
                     LL b = stk.back();
13
                     stk.pop_back();
if (miller_rabin(b)) {
14
15
16
17
                            ans.push_back(b);
                            continue;
18
19
                     LL c = b;
20
21
22
23
24
                     while (c >= b) c = pollard_rho(b);
                     stk.push_back(c);
stk.push_back(b / c);
              return ans;
25
       };
```

### 5.11 quadratic residu

```
/* cipolla */
        auto cipolla = [&](int x) {
               std::srand(time(0));
 \frac{3}{4} \\ \frac{5}{6}
               auto check = [\&] (int x) -> bool { return pow(x, (mod - 1) / 2) == 1; };
               if (!x) return 0;
if (!check(x)) return -1;
               int a, b;
while (1) {
 7
 8
 9
                      a = rand() % mod;
                      b = sub(mul(a, a), x);
if (!check(b)) break;
10
11
13
               PII t = {a, 1};
              PII t = {a, 1};
PII ans = {1, 0};
auto mulp = [&](PII x, PII y) -> PII {
    auto [x1, x2] = x;
    auto [y1, y2] = y;
    int c = add(mul(x1, y1), mul(x2, y2, b));
    int d = add(mul(x1, y2), mul(x2, y1));
    return [a, d];
14
15
16
17
18
19
20
                       return {c, d};
21
               };
```

5.12 Lucas 41

### **5.12** Lucas

## 卢卡斯定理

用于求大组合数,并且模数是一个不大的素数.

$$\left(\begin{array}{c} n \\ m \end{array}\right) \bmod p = \left(\begin{array}{c} \lfloor n/p \rfloor \\ \lfloor m/p \rfloor \end{array}\right) \cdot \left(\begin{array}{c} n \bmod p \\ m \bmod p \end{array}\right) \bmod p$$

$$\left(\begin{array}{c} n \bmod p \\ m \bmod p \end{array}\right)$$
 可以直接计算, $\left(\begin{array}{c} \lfloor n/p \rfloor \\ \lfloor m/p \rfloor \end{array}\right)$  可以继续使用卢卡斯计算.

递归至 m=0 的时候, 返回 1.

p 不太大, 一般在  $10^5$  左右.

```
auto C = [&](LL n, LL m, LL p) -> LL {
    if (n < m) return 0;
    if (m == 0) return 1;
    return fac[n] * inv_fac[m] % p * inv_fac[n - m] % p;
};

/* lucas */
auto lucas = [&](auto&& self, LL n, LL m, LL p) -> LL {
    if (n < m) return 0;
    if (m == 0) return 1;
    return C(n % p, m % p, p) * self(self, n / p, m / p, p) % p;
}</pre>
```

### 素数在组合数中的次数

Legengre 给出一种 n! 中素数 p 的幂次的计算方式为:

$$\sum_{1 \leqslant j} \left\lfloor \frac{n}{p^j} \right\rfloor.$$

另一种计算方式利用 p 进制下各位数字和:

$$v_p(n!) = \frac{n - S_p(n)}{p - 1}.$$

则有

$$v_p(C_m^n) = \frac{S_p(n) + S_p(m-n) - S_p(m)}{p-1}.$$

### 扩展卢卡斯定理

计算

$$\binom{n}{m} \mod p$$
,

p 可能为合数.

第一部分: CRT.

原问题变成求

$$\begin{cases} \begin{pmatrix} n \\ m \end{pmatrix} \equiv a_1 \bmod p_1^{\alpha_1} \\ \begin{pmatrix} n \\ m \end{pmatrix} \equiv a_2 \bmod p_2^{\alpha_2} \\ \dots \\ \begin{pmatrix} n \\ m \end{pmatrix} \equiv a_k \bmod p_k^{\alpha_1} \end{cases}$$

在求出 a<sub>i</sub> 之后就可以利用 CRT 求出答案.

第二部分: 移除分子分母中的素数

问题转换成求解

$$\binom{n}{m} \mod q^k$$

等价于

$$\frac{\frac{n!}{q^x}}{\frac{m!}{q^y}\frac{(n-m)!}{q^z}}q^{x-y-z} \bmod q^k,$$

其中 x 表示 n! 中 q 的次数, y, z 同理.

第三部分: 威尔逊定理的推论

问题转换为求

$$\frac{n!}{q^x} \mod q^k$$
.

可以利用威尔逊定理的推论.

```
/* exlucas */
                     auto exLucas = [&](LL n, LL m, LL p) {

\begin{array}{c}
2 \\
3 \\
4 \\
5 \\
6 \\
7 \\
8 \\
9 \\
10 \\
11 \\
12
\end{array}

                                        auto inv = [&](LL a, LL p) {
                                                         LL x, y;
exgcd(a, p, x, y);
return (x % p + p) % p;
                                        auto func = [&](auto&& self, LL n, LL pi, LL pk) {
                                                            if (!n) return 111;
                                                          LL ans = 1;

for (LL i = 2; i <= pk; i++) {

   if (i % pi) ans = ans * i % p;
13
14
                                                          ans = quick_power(ans, n / pk, pk);
for (LL i = 2; i <= n % pk; i++) {
   if (i % pi) ans = ans * i % pk;</pre>
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
32
33
34
35
36
37
38
                                                           ans = ans * self(self, n / pi, pi, pk) % pk;
                                                          return ans;
                                        };
                                        auto multiLucas = [&](LL n, LL m, LL pi, LL pk) {
                                                          b Multilucas = [@](LL ii, LL ii, LL iii, lt iii, 
                                                          return ans;
                                        auto crt = [&](const vl& a, const vl& m, int k) {
   LL ans = 0;
                                                           for (int i = 0; i < k; i++) {</pre>
                                                                               ans = (ans + a[i] * inv(p / m[i], m[i]) * (p / m[i])) % p;
39
```

5.13 Wilson 43

```
40
              return (ans % p + p) % p;
41
42
         43
44
45
46
             prime.push_back(1);
while (pp % i == 0) {
    prime.back() *= i;
47
48
49
50
51
                  pp /= i;
52
              a.push_back(multiLucas(n, m, i, prime.back()));
53
54
         if (pp > 1) {
             prime.push_back(pp);
a.push_back(multiLucas(n, m, pp, pp));
55
56
57
58
         return crt(a, prime, a.size());
59
    };
```

### 5.13 Wilson

## 简单结论

对于素数 p 有

$$(p-1)! \equiv -1 \mod p$$
.

### 推论

令  $(n!)_p$  表示不大于 n 且不被 p 整除的正整数的乘积.

特殊情形: n 为素数 p 时即为上述结论.

一般结论: 对素数 p 和正整数 q 有

$$((p^q)!)_p \equiv \pm 1 \bmod p^q.$$

详细定义:

$$((p^q)!)_p = \begin{cases} 1 & \text{if } p = 2 \text{ and } q \geqslant 3, \\ -1 & \text{other wise.} \end{cases}$$

### 5.14 LTE

将素数 p 在整数 n 中的个数记为  $v_p(n)$ .

(n,p)=1

对所有素数 p 和满足 (n,p)=1 的整数 n, 有

1. 若  $p \mid x - y$ , 则有

$$v_p(x^n - y^n) = v_p(x - y).$$

2. 若  $p \mid x - y$ , 则对奇数 n 有

$$v_p(x^n + y^n) = v_p(x + y).$$

## p 是奇素数

对所有奇素数 p 有

1. 若  $p \mid x - y$ , 则有

$$v_p(x^n - y^n) = v_p(x - y) + v_p(n).$$

2. 若  $p \mid x - y$ , 则对奇数 n 有

$$v_p(x^n + y^n) = v_p(x + y) + v_p(n).$$

p = 2

对 
$$p=2$$
 且  $p \mid x-y$  有

1. 对奇数 n 有

$$v_2(x^n - y^n) = v_2(x - y).$$

2. 对偶数 n 有

$$v_2(x^n - y^n) = v_2(x - y) + v_2(x + y) + v_2(n) - 1.$$

除此之外, 对上述 x, y, n, 若  $4 \mid x - y$ , 有

- 1.  $v_2(x+y) = 1$ .
- 2.  $v_2(x^n y^n) = v_2(x y) + v_2(n)$ .

### 5.15 Mobius inversion

### 莫比乌斯函数

$$\mu(n) = \begin{cases}
1 & n = 1, \\
0 & n 含有平方因子, \\
(-1)^k & k 为 n 的本质不同素因子个数.
\end{cases}$$

性质

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & n = 1 \\ 0 & n \neq 1 \end{cases}.$$
$$\varphi(n) = \sum_{d|n} d \cdot \mu(\frac{n}{d}).$$

反演结论

$$[\gcd(i,j)=1] = \sum_{d|\gcd(i,j)} \mu(d).$$

 $O(n \log n)$  求莫比乌斯函数

```
mu[1] = 1;
for (int i = 1; i <= n; i++){
   for (int j = i + i; j <= n; j += i){</pre>
```

5.15 Mobius inversion 45

# 莫比乌斯变换

设 
$$f(n), F(n)$$
.

1. 
$$F(n) = \sum_{d|n} f(d)$$
,  $\mathbb{M} f(n) = \sum_{d|n} \mu(d) F\left(\frac{n}{d}\right)$ .

2. 
$$F(n) = \sum_{n|d} f(d)$$
, 则  $f(n) = \sum_{n|d} \mu\left(\frac{d}{n}\right) F(d)$ .

# 6 math - polynomial

### 6.1 FTT

## FFT 与拆系数 FFT

```
const int sz = 1 \ll 23;
 2 3
    int rev[sz];
    int rev_n;
 4
     void set_rev(int n) {
         if (n == rev_n) return;
 6
7
         for (int i = 0; i < n; i++) rev[i] = (rev[i / 2] | (i & 1) * n) / 2;
         rev_n = n;
 8
 9
    tempt void butterfly(T* a, int n) {
10
         set_rev(n);
for (int i = 0; i < n; i++) {</pre>
11
12
              if (i < rev[i]) std::swap(a[i], a[rev[i]]);</pre>
13
14
15
16
    namespace Comp {
18
    long double pi = 3.141592653589793238;
19
\frac{20}{21}
    tempt struct complex {
         T x, y;
complex(T x = 0, T y = 0) : x(x), y(y) {}
complex operator+(const complex& b) const { return complex(x + b.x, y + b.y); }
22
23
\overline{24}
25
26
27
         complex operator*(const complex& b) const {
\frac{1}{28}
              return complex<T>(x * b.x - y * b.y, x * b.y + y * b.x);
29
30
          complex operator~() const { return complex(x, -y); }
31
         static complex unit(long double rad) { return complex(std::cos(rad), std::sin(rad)); }
32
33
34
          // namespace Comp
35
36
37
    struct fft_t {
    typedef Comp::complex<double> complex;
38
         complex wn[sz];
39
40
         fft_t() {
41
              for (int i = 0; i < sz / 2; i++) {
42
                   wn[sz / 2 + i] = complex::unit(2 * Comp::pi * i / sz);
43
44
              for (int i = sz / 2 - 1; i; i--) wn[i] = wn[i * 2];
45
46
47
         void operator()(complex* a, int n, int type) {
48
              if (type == -1) std::reverse(a + 1, a + n);
              butterfly(a, n);
49
50
              for (int i = 1; i < n; i *= 2) {
                   const complex* w = wn + i;
51
52
                   for (complex *b = a, t; b != a + n; b += i + 1) {
53
54
55
56
57
                        t = \bar{b}[i];
                       b[i] = *b - t;
*b = *b + t;
for (int j = 1; j < i; j++) {
    t = (++b)[i] * w[j];</pre>
58
59
                            b[i] = *b - t;
                             *b = *b + t;
60
                        }
                   }
61
62
63
              if (type == 1) return;
              for (int i = 0; i < n * 2; i++) ((double*) a)[i] /= n;
64
65
66
    } FFT;
67
68
     typedef decltype(FFT)::complex complex;
69
\frac{70}{71}
     vi fft(const vi& f, const vi& g) {
    static complex ff[sz];
72
         int n = f.size(), m = g.size();
73
          vi h(n + m - 1);
74
          if (std::min(n, m) <= 50) {</pre>
              for (int i = 0; i < n; i++) {</pre>
```

 $6.2 ext{ } FWT$ 

```
for (int j = 0; j < m; ++j) {
   h[i + j] += f[i] * g[j];</pre>
 76
 77
78
 79
                  }
 80
                  return h;
 81
 82
             int c = 1:
             while (c + 1 < n + m) c *= 2;
 83
            std::memset(ff, 0, sizeof(decltype(*(ff))) * (c));
for (int i = 0; i < n; i++) ff[i].x = f[i];</pre>
 84
 85
 86
             for (int i = 0; i < m; i++) ff[i].y = g[i];</pre>
 87
            FFT(ff, c, 1);
             for (int i = 0; i < c; i++) ff[i] = ff[i] * ff[i];</pre>
 88
 89
 90
            for (int i = 0; i + 1 < n + m; i++) h[i] = std::llround(ff[i].y / 2);</pre>
 91
 92
       }
 93
      94
 95
 96
 97
             int n = f.size(), m = g.size();
 98
             vi h(n + m - 1);
             if (std::min(n, m) <= 50) {</pre>
 99
                  for (int i = 0; i < n; ++i) {
   for (int j = 0; j < m; ++j) {</pre>
100
101
102
                             Add(h[i + j], mul(f[i], g[j]));
103
104
105
                  return h;
106
107
             int c = 1;
108
             while (c + 1 < n + m) c *= 2;
            for (int i = 0; i < 2; ++i) {
109
                  std::memset(ff[i], 0, sizeof(decltype(*(ff[i]))) * (c));
std::memset(gg[i], 0, sizeof(decltype(*(ff[i]))) * (c));
for (int j = 0; j < n; ++j) ff[i][j].x = f[j] / s[i] % s[1];
for (int j = 0; j < m; ++j) gg[i][j].x = g[j] / s[i] % s[1];</pre>
110
111
112
113
114
                  FFT(ff[i], c, 1);
115
                  FFT(gg[i], c, 1);
116
117
             for (int i = 0; i < c; ++i) {
                  ff[2][i] = ff[1][i] * gg[1][i];
ff[1][i] = ff[1][i] * gg[0][i];
gg[1][i] = ff[0][i] * gg[1][i];
118
119
120
                  ff[0][i] = ff[0][i] * gg[0][i];
121
122
123
            for (int i = 0; i < 3; ++i) {</pre>
                  FFT(ff[i], c, -1);
for (int j = 0; j + 1 < n + m; ++j) {
124
125
                       Add(h[j], mul(std::llround(ff[i][j].x) % mod, s[i]));
126
127
128
            FFT(gg[1], c, -1);
for (int i = 0; i + 1 < n + m; ++i) {
129
130
                  Add(h[i], \; mul(std::llround(gg[1][i].x) \; \% \; mod, \; s[1])); \\
131
132
133
             return h;
134
```

### 6.2 FWT

各种分治过程: and:

```
\begin{split} & FWT[A] = merge(FWT[A_0] + FWT[A_1], FWT[A_1]), \\ & UFWT[A'] = merge(UFWT[A'_0] - UFWT[A'_1], UFWT[A'_1]). \end{split}
```

or:

$$\begin{aligned} & \text{FWT}[\mathbf{A}] = merge(\text{FWT}[\mathbf{A}_0], \text{FWT}[\mathbf{A}_0] + \text{FWT}[\mathbf{A}_1]), \\ & \text{UFWT}[\mathbf{A}'] = merge(\text{UFWT}[\mathbf{A}'_0], -\text{UFWT}[\mathbf{A}'_0] + \text{UFWT}[\mathbf{A}'_1]). \end{aligned}$$

xor:

```
FWT[A] = merge(FWT[A_0] + FWT[A_1], FWT[A_0] - FWT[A_1]),
```

```
\text{UFWT[A']} = merge\left(\frac{\text{UFWT[A'_0]} + \text{UFWT[A'_1]}}{2}, \frac{\text{UFWT[A'_0]} - \text{UFWT[A'_1]}}{2}\right).
```

```
/* FWT */
     auto FWT_and = [&](vi v, int type) -> vi {
 2
          \begin{array}{c} 3 \\ 4 \\ 5 \\ 6 \\ 7 \\ 8 \end{array}
 9
                                v[i] = add(x, y);
10
                          } else {
                                v[i] = sub(x, y);
11
12
13
               }
14
15
          }
16
          return v;
17
     };
18
19
     auto FWT_or = [&](vi v, int type) -> vi {
\frac{20}{21}
          int n = v.size();
for (int mid = 1; mid < n; mid <<= 1) {</pre>
               for (int block = mid << 1, j = 0; j < n; j += block) {
  for (int i = j; i < j + mid; i++) {
    LL x = v[i], y = v[i + mid];
    if (type == 1) {
        v[i + mid] = add(x, y);
    }
}</pre>
22
23
24
25
26
\overline{27}
                          } else {
   v[i + mid] = sub(y, x);
28
29
30
                     }
31
               }
32
          }
33
          return v;
34
35
36
     auto FWT_xor = [&](vi v, int type) -> vi {
37
          int n = v.size();
          38
39
40
41
                          v[i] = add(x, y);
42
43
                          v[i + mid] = sub(x, y);
if (type == -1) {
44
                                Mul(v[i], inv2);
Mul(v[i + mid], inv2);
45
46
47
48
                     }
49
               }
50
51
          return v;
52
53
     };
     a = FWT(a, 1), b = FWT(b, 1);
for (int i = 0; i < (1 << n); i++) {
54
55
56
           c[i] = mul(a[i], b[i]);
57
58
     c = FWT(c, -1);
```

```
/* FWT @ wrb */
 2
       void FMTor(int f[]) {
            for (int i = 0; i < n; ++i)
for (int j = 0; j < m; ++j)
if (j >> i & 1) f[j] = (f[j] + f[j ^ 1 << i]) % P;
 3
 4
 5
 \frac{6}{7}
      void FMToriv(int f[]) {
            for (int i = 0; i < n; ++i)
    for (int j = 0; j < m; ++j)
        if (j >> i & 1) f[j] = (f[j] - f[j ^ 1 << i] + P) % P;</pre>
 9
10
11
      void FMTand(int f[]) {
            for (int i = 0; i < n; ++i)
for (int j = 0; j < m; ++j)
if (~j >> i & 1) f[j] = (f[j] + f[j ^ 1 << i]) % P;
13
14
15
16
17
      void FMTandiv(int f[]) {
            for (int i = 0; i < n; ++i)
    for (int j = 0; j < m; ++j)
        if (~j >> i & 1) f[j] = (f[j] - f[j ^ 1 << i] + P) % P;
18
19
20
21
    void FWT(int f[]) {
```

```
23
                   for (int len = 1; len < m; len *= 2) {</pre>
24
                            for (int i = 0; i < m; i += len * 2) {</pre>
                                     for (int j = i; j < i + len; ++j) {
    int x = f[j], y = f[j + len];
    f[j] = (x + y) % P;
    f[j + len] = (x - y + P) % P;
25
26
27
28
29
30
                            }
31
                   }
         }
32
          void FWTiv(int f[]) {
33
                  for (int len = 1; len < m; len *= 2) {
    for (int i = 0; i < m; i += len * 2) {
        for (int j = i; j < i + len; ++j) {
            int x = f[j], y = f[j + len];
            f[j] = 111 * (x + y) * iv2 % P;
            f[j + len] = 111 * (x - y + P) * iv2 % P;
}</pre>
34
35
36
37
38
39
40
                                     }
                            }
41
                  }
42
         }
43
```

# 6.3 class polynomial

```
class polynomial : public vi {
   public:
 2
 3
         polynomial() = default;
         polynomial(const vi& v) : vi(v) {}
 4
 5
         polynomial(vi&& v) : vi(std::move(v)) {}
 6
          int degree() { return size() - 1; }
 8
 9
          void clearzero() {
10
              while (size() && !back()) pop_back();
11
12
    };
13
14
15
     polynomial& operator+=(polynomial& a, const polynomial& b) {
         a.resize(std::max(a.size(), b.size()), 0);
for (int i = 0; i < b.size(); i++) {
16
17
18
              Add(a[i], b[i]);
19
20
          a.clearzero();
21
         return a;
22
    }
23
24
    polynomial operator+(const polynomial& a, const polynomial& b) {
25
         polynomial ans = a;
26
          return ans += b;
27
     }
28
29
     polynomial& operator-=(polynomial& a, const polynomial& b) {
         a.resize(std::max(a.size(), b.size()), 0);
for (int i = 0; i < b.size(); i++) {
30
31
32
              Sub(a[i], b[i]);
33
\frac{34}{35}
          a.clearzero();
         return a;
36
     }
37
38
    polynomial operator-(const polynomial& a, const polynomial& b) {
39
         polynomial ans = a;
40
          return ans -= b;
41
42
     class ntt_t {
43
44
45
         static const int maxbit = 22;
         static const int sz = 1 << maxbit;
static const int mod = 998244353;
46
47
48
          static const int g = 3;
49
         std::array<int, sz + 10> w;
std::array<int, maxbit + 10> len_inv;
50
51
52
53
         ntt_t() {
              int wn = pow(g, (mod - 1) >> maxbit);
54
              w[0] = 1;
55
              for (int i = 1; i <= sz; i++) {</pre>
56
57
                   w[i] = mul(w[i - 1], wn);
58
59
              len_inv[maxbit] = pow(sz, mod - 2);
```

```
60
                  for (int i = maxbit - 1; ~i; i--) {
 61
                       len_inv[i] = add(len_inv[i + 1], len_inv[i + 1]);
 62
 63
 64
 65
            void operator()(vi& v, int& n, int type) {
                  int bit = 0;
while ((1 << bit) < n) bit++;
int tot = (1 << bit);</pre>
 66
 67
 68
 69
                  v.resize(tot, 0);
 70
71
72
73
74
75
76
77
78
                  vi rev(tot);
                  n = tot;
                  for (int i = 0; i < tot; i++) {
    rev[i] = rev[i >> 1] >> 1;
                       if (i & 1) {
    rev[i] |= tot >> 1;
                  for (int i = 0; i < tot; i++) {</pre>
                       if (i < rev[i]) {</pre>
 80
                             std::swap(v[i], v[rev[i]]);
 81
82
 83
                  for (int midd = 0; (1 << midd) < tot; midd++) {</pre>
 84
                        int mid = 1 << midd;</pre>
 85
                        int len = mid << 1;</pre>
 86
                        for (int i = 0; i < tot; i += len) {</pre>
                             for (int j = 0; j < mid; j++) {
   int w0 = v[i + j];</pre>
 87
 88
 89
                                   int w1 = mul(
                                  w[-max(
   w[type == 1 ? (j << maxbit - midd - 1) : (len - j << maxbit - midd - 1)],
   v[i + j + mid]);
v[i + j] = add(w0, w1);
v[i + j + mid] = sub(w0, w1);</pre>
 90
 91
 92
 93
 94
                             }
                       }
 95
 96
                  if (type == -1) {
 97
 98
                       for (int i = 0; i < tot; i++) {</pre>
 99
                             v[i] = mul(v[i], len_inv[bit]);
100
101
                  }
102
      } NTT;
103
```

### 乘法

```
polynomial& operator*=(polynomial& a, const polynomial& b) {
   if (!a.size() || !b.size()) {
 2
 3
                  a.resize(0);
 4
5
                  return a;
 6
7
            polynomial tmp = b;
int deg = a.size() + b.size() - 1;
            int temp = deg;
 8
 9
10
            // 项数较小直接硬算
11
12
            if ((LL) a.size() * (LL) b.size() <= (LL) deg * 50LL) {</pre>
13
                   tmp.resize(0);
                  tmp.resize(deg, 0);
for (int i = 0; i < a.size(); i++) {
   for (int j = 0; j < b.size(); j++) {
      tmp[i + j] = add(tmp[i + j], mul(a[i], b[j]));
}</pre>
14
15
16
17
18
19
20
                  }
                   a = tmp;
\overline{21}
                  return a;
22
23
\frac{24}{25}
            // 项数较多跑 NTT
26
            NTT(a, deg, 1);
\overline{27}
            NTT(tmp, deg, 1);
for (int i = 0; i < deg; i++) {
\frac{1}{28}
\overline{29}
                  Mul(a[i], tmp[i]);
\frac{29}{30} 31
            NTT(a, deg, -1);
32
             a.resize(temp);
33
            return a;
34
      }
35
```

```
36 | polynomial operator*(const polynomial& a, const polynomial& b) {
37 | polynomial ans = a;
38 | return ans *= b;
39 | }
```

逆

```
polynomial inverse(const polynomial& a) {
            polynomial ans({pow(a[0], mod - 2)});
polynomial temp;
 \begin{array}{c} 3 \\ 4 \\ 5 \end{array}
            polynomial tempa;
            int deg = a.size();
for (int i = 0; (1 << i) < deg; i++) {
 6
 7
                  tempa.resize(0);
 8
                  tempa.resize(1 << i << 1, 0);
for (int j = 0; j != tempa.size() and j != deg; j++) {
   tempa[j] = a[j];</pre>
10
11
12
                  temp = ans * (polynomial({2}) - tempa * ans);
                  if (temp.size() > (1 << i << 1)) {
    temp.resize(1 << i << 1, 0);</pre>
13
14
15
16
                  temp.clearzero();
17
                  std::swap(temp, ans);
18
19
            ans.resize(deg);
20
            return ans;
21
```

## 对数

```
polynomial diffrential(const polynomial& a) {
 3
          if (!a.size()) {
              return a:
 4
          polynomial ans(vi(a.size() - 1));
 5
6
7
          for (int i = 1; i < a.size(); i++) {
   ans[i - 1] = mul(a[i], i);</pre>
 8 9
          return ans;
     }
10
11
12
     polynomial integral(const polynomial& a) {
          polynomial ans(vi(a.size() + 1));
for (int i = 0; i < a.size(); i++) {</pre>
13
14
               ans[i + 1] = mul(a[i], pow(i + 1, mod - 2));
15
16
17
          return ans;
18
     }
19
20
     polynomial ln(const polynomial& a) {
\overline{21}
          int deg = a.size();
22
          polynomial da = diffrential(a);
23
24
25
          polynomial inva = inverse(a);
          polynomial ans = integral(da * inva);
          ans.resize(deg);
26
          return ans;
```

## 指数

```
polynomial exp(const polynomial& a) {
    polynomial ans({1});
    polynomial temp;
    polynomial tempa;
    polynomial tempaa;
    int deg = a.size();
    for (int i = 0; (1 << i) < deg; i++) {
        tempa.resize(0);
        tempa.resize(1 << i << 1, 0);
        for (int j = 0; j != tempa.size() and j != deg; j++) {
            tempa[j] = a[j];
        }
        tempaa = ans;
    }
}</pre>
```

```
14
                    tempaa.resize(1 << i << 1);</pre>
                    temp = ans * (tempa + polynomial({1}) - ln(tempaa));
if (temp.size() > (1 << i << 1)) {
    temp.resize(1 << i << 1, 0);</pre>
15
16
17
18
19
                    temp.clearzero();
20
                   std::swap(temp, ans);
\overline{21}
\overline{22}
             ans.resize(deg);
23
             return ans;
24
```

### 根号

```
polynomial sqrt(polynomial& a) {
 2
           polynomial ans({cipolla(a[0])});
if (ans[0] == -1) return ans;
 \frac{1}{3}
           polynomial temp;
 5
6
7
           polynomial tempa;
           polynomial tempaa;
           int deg = a.size();
for (int i = 0; (1 << i) < deg; i++) {</pre>
 8
                 tempa.resize(0);
10
                 tempa.resize(1 << i << 1, 0);
                 for (int j = 0; j != tempa.size() and j != deg; j++) {
   tempa[j] = a[j];
11
12
13
14
                 tempaa = ans;
                tempaa - ams,
tempaa.resize(1 << i << 1);
temp = (tempa * inverse(tempaa) + ans) * inv2;
if (temp.size() > (1 << i << 1)) {</pre>
15
16
17
18
                      temp.resize(1 << i << 1, 0);
19
\frac{1}{20}
                 temp.clearzero();
                 std::swap(temp, ans);
22
23
           ans.resize(deg);
24
           return ans;
25
26
     // 特判 //
27
28
29
     int cnt = 0;
for (int i = 0; i < a.size(); i++) {</pre>
30
           if (a[i] == 0) {
31
32
                cnt++;
33
           } else {
34
                break;
35
36
37
     if (cnt) {
38
           if (cnt == n) {
                for (int i = 0; i < n; i++) {
    std::cout << "0";
39
40
41
42
                 std::cout << endl;</pre>
43
                return 0;
44
45
           if (cnt & 1) {
46
                std::cout << "-1" << endl;
47
                return 0;
48
           polynomial b(vi(a.size() - cnt));
for (int i = cnt; i < a.size(); i++) {
   b[i - cnt] = a[i];</pre>
49
50
51
52
53
           a = b;
54
     a.resize(n - cnt / 2);
55
     a = sqrt(a);
if (a[0] == -1) {
    std::cout << "-1" << endl;</pre>
56
57
58
59
           return 0;
60
61
     reverse(all(a));
62
     a.resize(n):
     reverse(all(a));
```

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## 6.4 wsy poly

```
#include <bits/stdc++.h>
     using ul = std::uint32_t;
 3
     using li = std::int32_t;
     using ll = std::int64_t;
     using ull = std::uint64_t;
using llf = long double;
using lf = double;
     using vul = std::vector;
using vvul = std::vector<vul>;
using pulb = std::pair<ul, bool>;
11
     using vpulb = std::vector<pulb>;
using vvpulb = std::vector<vpulb>
13
     using vb = std::vector<bool>;
14
15
     const ul base = 998244353;
17
18
     std::mt19937 rnd;
19
20
     ul plus(ul a, ul b) { return a + b < base ? a + b : a + b - base; }
21
     ul minus(ul a, ul b) { return a < b ? a + base - b : a - b; }
23
\overline{24}
     ul mul(ul a, ul b) { return ull(a) * ull(b) % base; }
\overline{25}
26
     void exgcd(li a, li b, li& x, li& y) {
\frac{1}{27}
          if (b) {
28
               exgcd(b, a % b, y, x);
y -= x * (a / b);
\frac{1}{29}
          } else {
30
               x = 1;

y = 0;
31
32
33
34
     }
35
36
     ul inverse(ul a) {
          li x, y;
exgcd(a, base, x, y);
return x < 0 ? x + li(base) : x;
37
38
39
40
     }
41
     ul pow(ul a, ul b) {
42
          ul ret = 1;
ul temp = a;
43
44
          while (b) {
   if (b & 1) {
45
46
47
                   ret = mul(ret, temp);
48
49
               temp = mul(temp, temp);
50
               b > = 1;
51
52
          return ret;
53
     }
54
55
     ul sqrt(ul x) {
56
57
          ul a;
          ul w2;
58
          while (true) {
    a = rnd() % base;
59
60
61
               w2 = minus(mul(a, a), x);
62
               if (pow(w2, base - 1 >> 1) == base - 1) {
63
                    break;
64
65
66
          ul b = base + 1 >> 1;
67
          ul rs = 1, rt = 0;
68
          ul as = a, at = 1;
          ul qs, qt;
while (b) {
69
70
71
               if (b & 1) {
72
73
74
75
76
                    qs = plus(mul(rs, as), mul(mul(rt, at), w2));
                    qt = plus(mul(rs, at), mul(rt, as));
                    rs = qs;
                    rt = qt;
77
               b >>= 1;
78
               qs = plus(mul(as, as), mul(mul(at, at), w2));
79
               qt = plus(mul(as, at), mul(as, at));
               as = qs;
80
               at = qt;
81
82
83
          return rs + rs < base ? rs : base - rs;
```

```
ul log(ul x, ul y, bool inited = false) {
 87
          static std::map<ul, ul> bs;
 88
           const ul d = std::round(std::sqrt(lf(base - 1)));
 89
           if (!inited) {
               bs.clear();
 90
               for (ul i = 0, j = 1; i != d; ++i, j = mul(j, x)) {
   bs[j] = i;
 91
 92
 93
               }
 94
          ul temp = inverse(pow(x, d));
for (ul i = 0, j = 1;; i += d, j = auto it = bs.find(mul(y, j));
 95
 96
                                               j = mul(j, temp)) {
 97
 98
               if (it != bs.end()) {
 99
                    return it->second + i;
100
101
          }
     }
102
103
     ul powroot(ul x, ul y, bool inited = false) {
   const ul g = 3;
104
105
106
           ul lgx = log(g, x, inited);
107
          li s, t;
exgcd(y, base - 1, s, t);
if (s < 0) {</pre>
108
109
110
               s += base - 1;
111
112
          return pow(g, ull(s) * ull(lgx) % (base - 1));
113
114
115
      class polynomial : public vul {
116
          public:
          void clearzero() {
   while (size() && !back()) {
117
118
                   pop_back();
119
120
121
          polynomial() = default;
122
123
          polynomial(const vul& a) : vul(a) {}
          polynomial(vul&& a) : vul(std::move(a)) {}
ul degree() const { return size() - 1; }
124
125
126
          ul operator()(ul x) const {
127
               ul ret = 0;
128
               for (ul i = size() - 1; ~i; --i) {
129
                   ret = mul(ret, x);
130
                    ret = plus(ret, vul::operator[](i));
131
132
               return ret;
133
          }
134
     };
135
136
     polynomial& operator+=(polynomial& a, const polynomial& b) {
          a.resize(std::max(a.size(), b.size()), 0);
for (ul_i = 0; i != b.size(); ++i) {
137
138
139
               a[i] = plus(a[i], b[i]);
140
141
          a.clearzero();
142
          return a;
143
144
145
      polynomial operator+(const polynomial& a, const polynomial& b) {
146
          polynomial ret = a;
147
           return ret += b;
148
149
150
     polynomial& operator = (polynomial& a, const polynomial& b) {
          a.resize(std::max(a.size(), b.size()), 0);
for (ul i = 0; i != b.size(); ++i) {
151
152
153
               a[i] = minus(a[i], b[i]);
154
155
          a.clearzero();
156
          return a;
157
158
159
     polynomial operator-(const polynomial& a, const polynomial& b) {
160
          polynomial ret = a;
161
          return ret -= b;
162
163
164
      class ntt_t {
165
          public:
1.66
           static const ul lgsz = 20;
167
          static const ul sz = 1 << lgsz;</pre>
168
           static const ul g = 3;
169
          ul w[sz + 1];
          ul leninv[lgsz + 1];
170
171
          ntt_t() {
```

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```
172
                ul wn = pow(g, (base - 1) >> lgsz);
173
                 w[0] = 1;
174
                for (ul i = 1; i <= sz; ++i) {</pre>
175
                      w[i] = mul(w[i - 1], wn);
176
177
                leninv[lgsz] = inverse(sz);
178
179
                for (ul i = lgsz - 1; ~i; --i) {
    leninv[i] = plus(leninv[i + 1], leninv[i + 1]);
180
181
182
            void operator()(vul& v, ul& n, bool inv) {
                ul lgn = 0;
while ((1 << lgn) < n) {
183
184
185
                      ++1gn;
186
187
                n = 1 \ll lgn;
                v.resize(n, 0);
for (ul i = 0, j = 0; i != n; ++i) {
   if (i < j) {</pre>
188
189
190
191
                          std::swap(v[i], v[j]);
192
                     193
194
195
196
                           k >>= 1;
                      }
197
198
                      j |= k;
199
200
                for (ul lgmid = 0; (1 << lgmid) != n; ++lgmid) {</pre>
201
                      ul mid = 1 << lgmid;
202
                      ul len = mid << 1;
203
                      for (ul i = 0; i != n; i += len) {
                          for (ul j = 0; j != mid; ++j) {
    ul t0 = v[i + j];
204
205
                                ul t1 =
206
                                    mul(w[inv ? (len - j << lgsz - lgmid - 1) : (j << lgsz - lgmid - 1)],
    v[i + j + mid]);</pre>
207
208
                                v[i + j] = plus(t0, t1);
209
                                v[i + j + mid] = minus(t0, t1);
210
211
                          }
212
                     }
213
214
                 if (inv) {
                     for (ul i = 0; i != n; ++i) {
215
216
                           v[i] = mul(v[i], leninv[lgn]);
217
218
                }
219
           }
220
      } ntt;
221
222
      polynomial& operator*=(polynomial& a, const polynomial& b) {
223
            if (!b.size() || !a.size()) {
224
                a.resize(0);
225
                return a;
226
227
           polynomial temp = b;
ul npmp1 = a.size() + b.size() - 1;
228
229
            if (ull(a.size()) * ull(b.size()) <= ull(npmp1) * ull(50)) {</pre>
230
                temp.resize(0);
                temp.resize(npmp1, 0);
for (ul i = 0; i != a.size(); ++i) {
    for (ul j = 0; j != b.size(); ++j) {
        temp[i + j] = plus(temp[i + j], mul(a[i], b[j]));
}
231
232
233
234
235
                      }
236
\bar{2}37
                a = temp;
238
                a.clearzero();
239
                return a;
240
241
           ntt(a, npmp1, false);
           ntt(temp, npmp1, false);
for (ul i = 0; i != npmp1; ++i) {
    a[i] = mul(a[i], temp[i]);
242
243
244
245
246
           ntt(a, npmp1, true);
247
            a.clearzero();
248
           return a;
249
      }
250
251
      polynomial operator*(const polynomial& a, const polynomial& b) {
252
           polynomial ret = a;
253
            return ret *= b;
\begin{array}{c} 254 \\ 255 \end{array}
256
      polynomial& operator*=(polynomial& a, ul b) {
257
           if (!b) {
258
                a.resize(0);
```

```
259
                 return a;
260
261
            for (ul i = 0; i != a.size(); ++i) {
                 a[i] = mul(a[i], b);
262
263
264
            return a;
265
       }
266
       polynomial operator*(const polynomial& a, ul b) {
    polynomial ret = a;
267
268
269
            return ret *= b;
270
271
272
       polynomial inverse(const polynomial& a, ul lgdeg) {
272
273
274
275
276
277
            polynomial ret({inverse(a[0])});
            polynomial temp;
            polynomial tempa;
             for (ul i = 0; i != lgdeg; ++i) {
                 tempa.resize(0);
                 tempa.resize(0',
tempa.resize(1 << i << 1, 0);
for (ul j = 0; j != tempa.size() && j != a.size(); ++j) {
   tempa[j] = a[j];</pre>
278
279
280
281
                 temp = ret * (polynomial({2}) - tempa * ret);
if (temp.size() > (1 << i << 1)) {</pre>
282
283
284
                       temp.resize(1 << i << 1, 0);
285
286
                  temp.clearzero();
287
                  std::swap(temp, ret);
288
289
            return ret;
290 }
291
       void quotientremain(const polynomial& a, polynomial b, polynomial& q, polynomial& r) {
   if (a.size() < b.size()) {</pre>
292
293
294
                 q = polynomial();
r = std::move(a);
295
296
                 return;
297
298
            std::reverse(b.begin(), b.end());
299
            auto ta = a;
300
            std::reverse(ta.begin(), ta.end());
            ul n = a.size() - 1;
ul m = b.size() - 1;
301
302
303
            ta.resize(n - m + 1);
            ul lgnmmp1 = 0;
while ((1 << lgnmmp1) < n - m + 1) {
304
305
306
                  ++lgnmmp1;
307
            q = ta * inverse(b, lgnmmp1);
q.resize(n - m + 1);
308
309
310
311
            std::reverse(b.begin(), b.end());
            std::reverse(q.begin(), q.end());
312
            r = a - b * q;
313
314
315
       polynomial mod(const polynomial& a, const polynomial& b) {
316
            polynomial q, r;
317
             quotientremain(a, b, q, r);
318
            return r;
319
320
321
       polynomial quotient(const polynomial& a, const polynomial& b) {
322
            polynomial q, r; quotientremain(a, b, q, r);
323
324
            return q;
325
326
327
       polynomial sqrt(const polynomial& a, ul lgdeg) {
    polynomial ret({sqrt(a[0])});
328
329
            polynomial temp;
330
            polynomial tempa;
            for (ul i = 0; i != lgdeg; ++i) {
    tempa.resize(0);
331
332
                 tempa.resize(1 << i << 1, 0);
for (ul j = 0; j != tempa.size() && j != a.size(); ++j) {
    tempa[j] = a[j];</pre>
333
334
335
336
                  temp = (tempa * inverse(ret, i + 1) + ret) * (base + 1 >> 1);
if (temp.size() > (1 << i << 1)) {
   temp.resize(1 << i << 1, 0);</pre>
337
338
339
340
341
                  temp.clearzero();
342
                  std::swap(temp, ret);
343
344
            return ret;
345 }
```

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```
346
347
      polynomial diffrential(const polynomial& a) {
348
           if (!a.size()) {
349
                return a;
350
           polynomial ret(vul(a.size() - 1, 0));
351
           for (ul i = 1; i != a.size(); ++i) {
   ret[i - 1] = mul(a[i], i);
352
353
354
355
           return ret;
      }
356
357
358
      polynomial integral(const polynomial& a) {
359
           polynomial ret(vul(a.size() + 1, 0));
360
           for (ul i = 0; i != a.size(); ++i) {
361
                ret[i + 1] = mul(a[i], inverse(i + 1));
362
363
           return ret;
364
      }
365
      polynomial ln(const polynomial& a, ul lgdeg) {
   polynomial da = diffrential(a);
366
367
368
           polynomial inva = inverse(a, lgdeg);
           polynomial ret = integral(da * inva);
if (ret.size() > (1 << lgdeg)) {</pre>
369
370
                ret.resize(1 << lgdeg);</pre>
371
372
                ret.clearzero();
373
374
           return ret;
375
      }
376
      polynomial exp(const polynomial& a, ul lgdeg) {
    polynomial ret({1});
    polynomial temp;
377
378
379
           polynomial tempa;
for (ul i = 0; i != lgdeg; ++i) {
380
381
382
                tempa.resize(0);
                tempa.resize(1 << i << 1, 0);
383
384
                for (ul j = 0; j != tempa.size() && j != a.size(); ++j) {
385
                     tempa[j] = a[j];
386
387
                temp = ret * (polynomial({1}) - ln(ret, i + 1) + tempa);
388
                if (temp.size() > (1 << i << 1)) {</pre>
389
                     temp.resize(1 << i << 1, 0);
390
391
                temp.clearzero();
392
                std::swap(temp, ret);
393
394
           return ret;
395
396
397
      polynomial pow(const polynomial& a, ul k, ul lgdeg) { return exp(ln(a, lgdeg) * k, lgdeg); }
398
399
      polynomial alpi[1 << 16][17];</pre>
400
401
      polynomial getalpi(const ul x[], ul 1, ul lgrml) {
           if (lgrml == 0)
402
403
                return alpi[1][lgrm1] = vul({minus(0, x[1]), 1});
404
           return alpi[1][lgrm1] = getalpi(x, 1, lgrm1 - 1) * getalpi(x, 1 + (1 << lgrm1 - 1), lgrm1 - 1);</pre>
405
      }
406
407
      void multians(const polynomial& f, const ul x[], ul y[], ul 1, ul lgrml) {
    if (f.size() <= 700) {
        for (ul_i = 1; i_! = 1 + (1 << lgrml); ++i) {</pre>
408
409
410
411
                     y[i] = f(x[i]);
412
413
                return;
414
           if (lgrml == 0) {
415
                y[1] = f(x[1]);
416
417
                return:
418
           multians(mod(f, alpi[1][lgrml - 1]), x, y, 1, lgrml - 1);
multians(mod(f, alpi[1 + (1 << lgrml - 1)][lgrml - 1]), x, y, 1 + (1 << lgrml - 1), lgrml - 1);</pre>
419
420
421
      }
422
423
      ul sqrt(ul x) {
424
           ūl a;
425
           ul w2;
426
           while (true) {
427
                a = rnd() % base;
                w2 = minus(mul(a, a), x);
if (pow(w2, base - 1 >> 1) == base - 1) {
428
429
430
                     break:
                }
431
           }
432
```

```
433
           ul b = base + 1 >> 1;
434
           ul rs = 1, rt = 0;
435
           ul as = a, at = 1;
           ul qs, qt;
while (b) {
436
437
438
               if (b & 1) {
                    qs = plus(mul(rs, as), mul(mul(rt, at), w2));
qt = plus(mul(rs, at), mul(rt, as));
439
440
441
442
                    rt = qt;
443
444
                b >>= 1:
445
                qs = plus(mul(as, as), mul(mul(at, at), w2));
                qt = plus(mul(as, at), mul(as, at));
446
                as = qs;
447
448
                at = qt;
449
450
           return rs + rs < base ? rs : base - rs;</pre>
451
452
453
      ul log(ul x, ul y, bool inited = false) {
454
           static std::map<ul, ul> bs;
455
           const ul d = std::round(std::sqrt(lf(base - 1)));
456
           if (!inited) {
457
                bs.clear();
               for (ul_i = 0, j = 1; i != d; ++i, j = mul(j, x)) {
458
459
                    bs[j] = i;
460
461
           lul temp = inverse(pow(x, d));
for (ul i = 0, j = 1;; i += d, j = mul(j, temp)) {
    auto it = bs.find(mul(y, j));
    if (it l= bs.gad()) {
462
463
464
465
                if (it != bs.end()) {
466
                    return it->second + i;
467
468
           }
469
470
      ul powroot(ul x, ul y, bool inited = false) {
   const ul g = 3;
   ul lgx = log(g, x, inited);
471
472
473
474
           li s, t;
           exgcd(y, base - 1, s, t);
if (s < 0) {
475
476
477
               s += base - 1;
478
479
           return pow(g, ull(s) * ull(lgx) % (base - 1));
480
481
482
      ul n;
483
484
      int main() {
           std::scanf("%u", &n);
485
486
           polynomial f;
487
           for (ul i = 0; i <= n; ++i) {
488
               ul t;
                std::scanf("%u", &t);
489
490
               f.push_back(t % base);
491
492
           polynomial g = \exp(\ln(f * inverse(f[0]), 17) * inverse(3), 17) * powroot(f[0], 3);
493
           while (g.size() <= n) {</pre>
                g.push_back(0);
494
495
496
           for (ul i = 0; i <= n; ++i) {</pre>
                if (i) {
497
                    std::putchar(' ');
498
499
500
                std::printf("%u", g[i]);
501
502
           std::putchar(' \ n');
503
           return 0;
504
```

## 6.5 Lagrange interpolation

### 一般的插值

插值的结果是

$$f(x) = \sum_{i=1}^{n} y_i \prod_{j \neq i} \frac{x - x_j}{x_i - x_j},$$

直接带入 k 并且取模即可, 时间复杂度  $O(n^2)$ .

## 坐标连续的插值

给出的点是  $(0, y_0), \dots, (n, y_n)$ .

```
/* interpolation */
auto lagrange(int n, std::vector<Z> y, int x) -> Z {
    if (x <= n) return y[x];
    Z s1 = 1, s2 = 0;
    for (int i = 1; i <= n; i++) s1 *= (x - i);
    for (int i = 1; i <= n; i++) {
        Z res = ((n - i) & 1 ? -y[i] : y[i]);
        res /= (x - i);
        res *= fiv[i - 1];
        res *= fiv[n - i];
        s2 += res;
    }
    return s1 * s2;
}</pre>
```

# 7 math - game theory

# 7.1 nim game

若 nim 和为 0, 则先手必败.

暴力打表.

```
| vi SG(21, -1); /* 记忆化 */
| std::function<int(int, int)> sg = [&](int x) -> int {
| if (/* 为最终态 */) return SG[x] = 0; |
| if (SG[x] != -1) return SG[x]; |
| vi st; |
| for (/* 枚举所有可到达的状态 y */) {
| st.push_back(sg(y)); |
| } |
| std::sort(all(st)); |
| for (int i = 0; i < st.size(); i++) {
| if (st[i] != i) return SG[x] = i; |
| } |
| return SG[x] = st.size(); |
```

# 7.2 anti - nim game

若

- 1. 所有堆的石子均为一个, 且 nim 和不为 0,
- 2. 至少有一堆石子超过一个, 且 nim 和为 0,

则先手必败.

# 8 math - linear algebra

### 8.1 matrix

#### determinant mod 998244353

```
/* determinant */
       auto det = [&](int n, vvi e) -> int {
 \bar{3}
              int ans = 1;
 4
5
              for (int i = 1; i <= n; i++) {</pre>
                    if (a[i][i] == 0) {
    for (int j = i + 1; j <= n; j++) {
        if (a[j][i] != 0) {
 6
7
                                         for (int k = i; k <= n; k++) {</pre>
 9
                                               std::swap(a[i][k], a[j][k]);
10
11
                                         ans = sub(mod, ans);
12
                                         break;
                                  }
13
                           }
14
15
                    if (a[i][i] == 0) return 0;
Mul(ans, a[i][i]);
16
17
                    int x = pow(a[i][i], mod - 2);
for (int k = i; k <= n; k++) {
    Mul(a[i][k], x);</pre>
18
19
20
21
                    for (int j = i + 1; j <= n; j++) {
   int x = a[j][i];
   for (int k = i; k <= n; k++) {
      Sub(a[j][k], mul(a[i][k], x));
}</pre>
\overline{22}
23
\overline{24}
25
26
27
                    }
28
29
              return ans;
30
       };
```

### determinant mod non-prime

```
/* determinant @ wrb */
       int a[609][609];
      int a[cos];cos;
int main() {
   int n, P, z = 1;
   std::cin >> n >> P;
   for (int i = 1; i <= n; ++i) {
      for (int j = 1; j <= n; ++j) std::cin >> a[i][j];
}
 \frac{1}{3}
 5
 6
 9
              for (int i = 1; i <= n; ++i) {</pre>
                    for (int j = i + 1; j <= n; ++j) {
    while (a[j][i]) {
10
                                 z = -z;
int d = a[i][i] / a[j][i];
for (int k = i; k <= n; ++k) {
   int x = (a[i][k] - 111 * d * a[j][k]) % P;</pre>
13
14
15
                                         a[i][k] = a[j][k], a[j][k] = x;
16
17
                           }
18
19
20
                    z = 111 * z * a[i][i] % P;
21
22
              std::cout << (z + P) % P;
23
```

## matrix multiplication

 $A_{n \times m}$  与  $B_{m \times k}$  相乘并模 998244353.

```
1    /* matrix multiplication */
2    auto matmul = [&](int n, int m, int k, const vvi& a, const vvi& b) -> vvi {
3       vvi c(n + 1, vi(k + 1));
4       for (int i = 1; i <= n; i++) {
5          for (int l = 1; 1 <= m; l++) {
</pre>
```

### 8.2 linear basis

```
/* linear basis */
vl p(63), s(63);
                                            /* basis and case */
       vl p(63), s(63);  /* basis and case *;
auto insert = [&](LL x, int id) {
    LL ans = 0;
    for (int i = 62; i >= 0; i--) {
        if (~(x >> i) & 1) continue;
        if (!p[i]) {
            p[i] = x;
            s[i] = ans ^ (111 << id);
            break;</pre>
 3
4
5
6
  7
8
  9
10
                                break;
11
12
                        x ^= p[i], ans ^= s[i];
13
14
                return x;
15
        auto_query = [&](LL x) {
16
17
               LL ans = 0;
for (int i = 62; i >= 0; i--) {
    if (~(x >> i) & 1) continue;
    x ^= p[i], ans ^= s[i];
18
19
\frac{10}{20}
22
                return (x ? -1 : ans);
23
\overline{24}
        auto queryMax = [&]() {
25
                LL ans = 0;
                for (int i = 62; i >= 0; i--)
26
                        if ((ans ^ p[i]) > ans) ans ^= p[i];
27
\overline{28}
                return ans:
29
        };
```

```
/* linear basis @ wrb */
template<typename T, const int M = sizeof(T) * 8>
struct Liner_Basis {
 \bar{2}
 \overline{3}
 \begin{array}{c} 4 \\ 5 \\ 6 \end{array}
            T a[M];
            size_t sz;
Liner_Basis() : a(), sz() {}
size_t size() const {
 7
8
                  return sz;
 9
10
            void clear() {
11
                  memset(a, 0, sizeof a);
12
            bool ins(T x) {
13
                  for (size_t i = M - 1; ~i && x; --i)
    if (x >> i & 1) {
        if (a[i]) x ^= a[i];
14
15
16
17
                               else return a[i] = x, true;
                         }
18
19
                  return false;
\begin{array}{c} 20 \\ 21 \\ 22 \end{array}
            Liner_Basis& operator+=(const Liner_Basis&_) {
                  for (T x : _.a) if (x) this->ins(x);
return *this;
\frac{23}{24}
\begin{array}{c} 25 \\ 26 \end{array}
            Liner_Basis operator+(const Liner_Basis&_) {
   Liner_Basis z = *this;
                  return z += _;
27
\frac{1}{28}
29
            T qry(T x = 0) {
                  for (size_t i = M - 1; ~i; --i)
if ((x ^ a[i]) > x) x ^= a[i];
30
31
32
33
                  return x;
            }
34
      template<typename T>
using LB = Liner_Basis<T>;
35
36
37
      38
39
40
```

```
41
     | struct Liner_Basis {
            using u64 = unsigned long long;
static const size_t M = 60;
42
43
44
            u64 a[M + 1];
            size_t sz;
size_t size() {
45
46
47
            return sz;
48
            Liner_Basis& operator+=(u64 x) {
for (size_t i = M; ~i && x; --i)
49
50
            if (x >> i & 1)
if (a[i]) x ^= a[i];
51
52
            else return a[i] = x, ++sz, *this;
53
            return *this;
54
55
56
            Liner_Basis& operator+=(const Liner_Basis&_) {
            for (u64 x : _.a) if (x) *this += x; return *this;
57
58
59
            Liner_Basis operator+(u64 x) {
Liner_Basis z = *this;
60
61
62
            return z += x;
63
            Liner_Basis operator+(const Liner_Basis&_) {
Liner_Basis z = *this;
64
65
66
            return z += _;
}
67
            u64 qry(u64 x = 0) {
for (size_t i = M; ~i; --i)
if ((x ^ a[i]) > x) x ^= a[i];
68
69
70
71
            return x;
72
73
74
            u64 rank(u64 x) {
            uo4 falk(uo4 x) {
uo4 h = 1, z = 0;
for (size_t i = 0; i <= M; ++i)
if (a[i]) {
if (x >> i & 1) z += h;
75
76
77
78
            h <<= 1;
79
80
            return z;
81
            u64 kth(u64 x) {
82
83
            u64 z = 0;
84
            for (size_t i = M; ~i; --i)
if (x >> i & 1) z ^= a[i];
85
86
            return z;
87
88
      }v;
      using LB = Liner_Basis;
```

### 8.3 linear programming

## 8.4 bm

```
/* bm @ wrb */
     const int p = 998244353;
     auto power = [](int a, int b = p - 2) {
 3
          int z = 1;
 4
          while (b) {
 5
 6
              if (b & 1) z = 111 * z * a % p;
 7
              a = 111 * a * a % p, b >>= 1;
 8
 9
          return z;
10
     };
11
     vector<int> berlekamp_massey(const vector<int> &a) {
       vector<int> v, last; // v is the answer, 0-based, p is the module
int k = -1, delta = 0;
12
13
14
       for (int i = 0; i < (int)a.size(); i++) {</pre>
15
          int tmp = 0;
for (int j = 0; j < (int)v.size(); j++)
   tmp = (tmp + (long long)a[i - j - 1] * v[j]) % p;</pre>
16
17
18
19
\frac{20}{21}
          if (a[i] == tmp) continue;
22
23
          if (k < 0) {
            k = i;
24
            delta = (a[i] - tmp + p) % p;
25
            v = vector < int > (i + 1);
26
27
            continue;
```

# 9 complex number

```
tandu struct Comp {

  \begin{array}{c}
    2 \\
    3 \\
    4 \\
    5 \\
    6 \\
    7 \\
    8 \\
    9
  \end{array}

                     T a, b;
                     Comp(T _a = 0, T _b = 0) { a = _a, b = _b; }
Comp operator+(const Comp& x) const { return Comp(a + x.a, b + x.b); }
Comp operator-(const Comp& x) const { return Comp(a - x.a, b - x.b); }
Comp operator*(const Comp& x) const { return Comp(a * x.a - b * x.b, a * x.b + b * x.a); }
bool operator==(const Comp& x) const { return a == x.a and b == x.b; }
                    bool operator==(const Comp& x) const { return T real() { return a; }
T imag() { return b; }
U norm() { return (U) a * a + (U) b * b; }
Comp conj() { return Comp(a, -b); }
Comp operator/(const Comp& x) const {
    Comp y = x;
    Comp c = Comp(a, b) * y.conj();
    T d = y.norm();
10
11
12
\overline{13}
14
                               T d = y.norm();
return Comp(c.a / d, c.b / d);
15
16
17
18
19
           typedef Comp<LL, LL> complex;
           complex gcd(complex a, complex b) {
  LL d = b.norm();
  if (d == 0) return a;
21
22
                    if (d == 0) return a;
std::vector<complex> v(4);
complex c = a * b.conj();
auto fdiv = [&](LL a, LL b) -> LL { return a / b - ((a ^ b) < 0 && (a % b)); };
v[0] = complex(fdiv(c.real(), d), fdiv(c.imag(), d));
v[1] = v[0] + complex(1, 0);
v[2] = v[0] + complex(0, 1);</pre>
23
\overline{24}
25
26
27
\frac{1}{28}
                     v[3] = v[0] + complex(1, 1);
29
30
                     for (auto& x : v) {
31
                              x = a - x * b;
32
33
                     std::sort(all(v), [&](complex a, complex b) { return a.norm() < b.norm(); });</pre>
34
                     return gcd(b, v[0]);
35
          };
```

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# 10 graph

## 10.1 topology sort

```
/* topology sort */
 \begin{array}{c} 1\\2\\3\\4\\5 \end{array}
      vi top;
      auto topsort = [&]() -> bool {
            vi d(n + 1);
            std::queue<int> q;
for (int i = 1; i <= n; i++) {
    d[i] = e[i].size();
    if (!d[i]) a push(i);</pre>
 6
7
 8
                  if (!d[i]) q.push(i);
 9
10
            while (!q.empty()) {
                  int u = q.front();
11
12
                  q.pop();
13
                  top.push_back(u);
                  for (auto v : e[u]) {
    d[v]--;
14
15
16
                        if (!d[v]) q.push(v);
17
18
19
            if (top.size() != n) return false;
            return true;
20
\overline{21}
      };
```

## 10.2 shortest path

Floyd

```
/* floyd */
  \begin{array}{c} 2 \\ 3 \\ 4 \\ 5 \\ 6 \\ 7 \end{array}
         auto floyd = [&]() -> vvi {
                 vvi dist(n + 1, vi(n + 1, inf));
for (int i = 1; i <= n; i++) {
    for (int j = 1; j <= n; j++) {
        Min(dist[i][j], e[i][j]);
    }</pre>
  8 9
                          dist[i][i] = 0;
10
                 for (int k = 1; k <= n; k++) {</pre>
                          for (int i = 1; i <= n; i++) {
   for (int j = 1; j <= n; j++) {
      Min(dist[i][j], dist[i][k] + dist[k][j]);
}</pre>
11
12
13
14
15
                          }
16
                 return dist;
17
        };
18
```

### Dijkstra

```
/* dijkstra */
 \begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
        auto dijkstra = [&](int s) -> vl {
                v1 dist(n + 1, INF);
vi vis(n + 1, 0);
                dist[s] = 0;
 6
7
                std::priority_queue<PLI, std::vector<PLI>, std::greater<PLI>> q;
q.emplace(OLL, s);
                q.empidecold, sy,
while (!q.empty()) {
    auto [dis, u] = q.top();
    q.pop();
    if (vis[u]) continue;
 8 9
10
11
12
                       vis[u] = 1;
                       for (const auto& [v, w] : e[u]) {
    if (dist[v] > dis + w) {
        dist[v] = dis + w;
        q.emplace(dist[v], v);
}
13
14
15
16
17
18
                       }
19
20
                return dist;
       };
```

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#### **SPFA**

```
/* SPFA */
      int n, m, s;
vl dist(n + 1, INF);
      std::vector<bool> vis(n + 1);
std::vector<PLI > e(n + 1);
      void spfa(int s){
            for (int i = 1; i <= n; i++) dist[i] = INF;
dist[s] = 0;</pre>
 8 9
            std::queue<int> q;
10
            q.push(s);
11
            vis[s] = true;
12
            while(q.size()){
                  auto u = q.front();
q.pop();
vis[u] = false;
13
14
15
                  for(const auto& [v, w] : e[u]){
    if(dist[v] > dist[u] + w){
        dist[v] = dist[u] + w;
16
18
19
                               if(!vis[v]){
20
                                     q.push(v);
21
                                     vis[v] = true;
22
                               }
23
                         }
\overline{24}
                  }
25
            }
26
      }
```

### Johnson

```
/* johnson */
 \bar{2}
     auto johnson = [&]() -> vvl {
          /* 负环 */
 3
 4
          vl dist1(n + 1);
         vi vis(n + 1), cnt(n + 1);
auto spfa = [&]() -> bool {
 5
 6
              std::queue<int> q;
 7
              for (int u = 1; u <= n; u++) {
 8
 9
                   q.push(u);
10
                   vis[u] = false;
11
12
               while (!q.empty()) {
13
                   auto u = q.front();
14
                   q.pop();
15
                   vis[u] = false;
                   for (auto [v, w] : e[u]) {
    if (dist1[v] > dist1[u] + w) {
        dist1[v] = dist1[u] + w;
    }
16
18
                             Max(cnt[v], cnt[u] + 1);
if (cnt[v] >= n) return true;
19
20
21
                             if (!vis[v]) {
\frac{1}{22}
                                  q.push(v);
23
24
25
                                  vis[v] = true;
                        }
26
                   }
27
28
              return false;
29
30
          /* dijkstra */
31
          vl dist2(n + 1);
         32
33
34
35
                   vis[u] = false;
36
37
              dist2[s] = 0;
              std::priority_queue<PLI, std::vector<PLI>, std::greater<PLI>> q;
q.emplace(0, s);
38
39
              while (!q.empty()) {
    auto [d, u] = q.top();
40
41
42
                   q.pop();
43
                   if (vis[u]) continue;
                   44
45
46
47
48
                             q.emplace(dist2[v], v);
49
                   }
50
```

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```
51
                   }
             };
if (spfa()) return vvl{};
52
53
             for (int u = 1; u <= n; u++) {
   for (auto& [v, w] : e[u]) {
      w += dist1[u] - dist1[v];
}</pre>
54
55
56
57
58
59
             vvl dist(n + 1, vl(n + 1));
for (int u; u <= n; u++) {</pre>
60
                   dijkstra(u);
61
                   for (int v = 1; v <= n; v++) {
   if (dist2[v] == 1e9) {</pre>
62
63
64
                                dist[u][v] = INF;
65
                            else {
66
                                dist[u][v] = dist2[v] + dist1[v] - dist1[u];
67
68
                   }
69
70
71
             return dist;
      };
```

# 最短路计数 - Dijkstra

```
/* dijkstra */
 3
      auto dijkstra = [&](int s) -> std::pair<vl, vi> {
            vl dist(n + 1, INF);
 4
            vi cnt(n + 1), vis(n + 1);
            dist[s] = 0;
 5
 6
7
            cnt[s] = 1;
           std::priority_queue<PLI, std::vector<PLI>, std::greater<PLI>> q;
q.emplace(OLL, s);
           while (!q.empty()) {
    auto [dis, u] = q.top();
 9
10
11
                  q.pop();
12
                  if (vis[u]) continue;
13
                  vis[u] = 1;
14
                  for (const auto& [v, w] : e[u]) {
                       if (dist[v] > dis + w) {
    dist[v] = dis + w;
15
16
17
                             cnt[v] = cnt[u];
                       q.push({dist[v], v});
} else if (dist[v] == dis + w) {
   // cnt[v] += cnt[u];
   cnt[v] += cnt[u];
18
19
\frac{10}{20} 21
22
23
                             cnt[v] %= mod;
\begin{array}{c} 23 \\ 24 \\ 25 \end{array}
                 }
26
           return {dist, cnt};
27
      };
```

# 最短路计数 - Floyd

```
/* floyd */
         auto floyd() = [&] -> std::pair<vvi, vvi> {
  \begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
                  6
  78
  9
                            dist[i][i] = 0;
10
11
                  for (int k = 1; k <= n; k++) {</pre>
                           (int k = 1; k <= n; k++) {
  for (int i = 1; i <= n; i++) {
    for (int j = 1; j <= n; j++) {
      if (dist[i][j] == dist[i][k] + dist[k][j]) {
         cnt[i][j] += cnt[i][k] * cnt[k][j];
      } else if (dist[i][j] > dist[i][k] + dist[k][j]) {
         cnt[i][j] = cnt[i][k] * cnt[k][j];
         dist[i][j] = dist[i][k] + dist[k][j];
    }
}
12
13
14
15
16
17
18
19
\frac{10}{20}
                                    }
                           }
22
23
                  return {dist, cnt};
         };
```

### 负环

判断的是最短路长度.

```
/* SPFA */
 3
     auto spfa = [&]() -> bool {
         std::queue<int> q;
vi vis(n + 1), cnt(n + 1);
 4
 5
6
7
8
9
         for (int i = 1; i <= n; i++) {</pre>
              q.push(i);
              vis[i] = true;
          while (!q.empty()) {
10
              auto u = q.front();
11
              q.pop();
              12
13
14
                        dist[v] = dist[u] + w;
15
                        cnt[v] = cnt[u] + 1;
if (cnt[v] >= n) return true;
if (!vis[v]) {
16
17
18
19
                             q.push(v);
\frac{20}{21}
                             vis[v] = true;
22
                   }
23
24
          }
25
         return false;
26
     }
```

# 10.3 minimum spanning tree

## Kruskal

```
std::vector<std::tuple<int, int, int>> edge;
     auto kruskal = [&]() -> int {
    std::sort(all(edge), [&](std::tuple<int, int, int> a, std::tuple<int, int, int> b) {
        auto [x1, y1, w1] = a;
        auto [x2, y2, w2] = b;
        return w1 < w2;
    }
}</pre>
 3
 5
 6
7
 8
           10
11
                 a = find(a), b = find(b);
if (a != b) {
12
13
14
                       fa[a] = b;
15
                       res += w;
16
                       /* res = std::max(res, w); */
                       cnt++;
18
                 }
19
20
           if (cnt < n - 1) return -1;</pre>
21
           return res;
22
     }
```

# 10.4 SCC

### Tarjan

```
/* tarjan */
vi dfn(n + 1), low(n + 1), stk(n + 1), belong(n + 1);
int timestamp = 0, top = 0, scc_cnt = 0;
std::vector<bool> in_stk(n + 1);
auto tarjan = [&](auto&& self, int u) -> void {
    dfn[u] = low[u] = ++timestamp;
    stk[++top] = u;
    in_stk[u] = true;
    for (const auto& v : e[u]) {
        if (!dfn[v]) {
            self(self, v);
        }
```

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```
12
                     Min(low[u], low[v]);
13
                } else if (in_stk[v]) {
14
                     Min(low[u], dfn[v]);
15
16
17
           if (dfn[u] == low[u]) {
                scc_cnt++;
int v;
18
19
                do {
20
                     v = stk[top--];
in_stk[v] = false;
belong[v] = scc_cnt;
21
22
23
24
25
                } while (v != u);
           }
26
     };
```

### 10.5 DCC

### 点双连通分量

求点双连通分量.

```
vi dfn(n + 1), low(n + 1), is_bcc(n + 1), stk;
int timestamp = 0, bcc_cnt = 0, root = 0;
vvi bcc(2 * n + 1);
std::function<void(int, int)> tarjan = [&](int u, int fa) {
 3
 4
 \begin{array}{c} 5 \\ 6 \\ 7 \end{array}
             dfn[u] = low[u] = ++timestamp;
             int child = 0;
             stk.push_back(u);
             if (u == root and e[u].empty()) {
 9
                   bcc_cnt++;
10
                   bcc[bcc_cnt].push_back(u);
11
                   return;
12
13
             for (auto v : e[u]) {
                   if (!dfn[v]) {
14
15
                         tarjan(v, u);
low[u] = std::min(low[u], low[v]);
if (low[v] >= dfn[u]) {
16
17
18
                                child++;
19
                                if (u != root or child > 1) {
20 \\ 21 \\ 22 \\ 23 \\ 24
                                      is_bcc[u] = 1;
                                bcc_cnt++;
                                int z;
                                do {
25
26
                                      z = stk.back();
                                      stk.pop_back();
27
                                bcc[bcc_cnt].push_back(z);
} while (z != v);
\frac{1}{28}

    \begin{array}{c}
      29 \\
      30 \\
      31 \\
      32
    \end{array}

                                bcc[bcc_cnt].push_back(u);
                         }
                   } else if (v != fa) {
                         low[u] = std::min(low[u], dfn[v]);
                   }
33
34
             }
35
      for (int i = 1; i <= n; i++) {
   if (!dfn[i]) {</pre>
36
37
38
                   root = i;
39
                   tarjan(i, i);
40
             }
41
      }
```

求割点.

 $10.5 \quad DCC \qquad \qquad 71$ 

```
14
                                  is_bcc[u] = 1;
15
                            }
16
17
                 } else if (v != fa) {
                      low[u] = std::min(low[u], dfn[v]);
18
19
           }
20
\overline{21}
     };
\overline{22}
     for (int i = 1; i <= n; i++) {
   if (!dfn[i]) {</pre>
\frac{1}{23}
24
25
                 root = i;
                 tarjan(i, i);
26
27
```

### 边双连通分量

求边双连通分量.

```
std::vector<vpi> e(n + 1);
 3
     for (int i = 1; i <= m; i++) {</pre>
          int u, v;
std::cin >> u >> v;
 4
 5
          e[u].emplace_back(v, i);
 6
          e[v].emplace_back(u, i);
     vi dfn(n + 1), low(n + 1), is_ecc(n + 1), fa(n + 1), stk;
int timestamp = 0, ecc_cnt = 0;
vvi ecc(2 * n + 1);
 8
10
11
     std::function<void(int, int)> tarjan = [&](int u, int id) {
          low[u] = dfn[u] = ++timestamp;
12
          stk.push_back(u);
for (auto [v, idx] : e[u]) {
13
14
               if (!dfn[v]) {
15
16
                     tarjan(v, idx);
               low[u] = std::min(low[u], low[v]);
} else if (idx != id) {
17
18
19
                     low[u] = std::min(low[u], dfn[v]);
20
21
22
23
          if (dfn[u] == low[u]) {
               ecc_cnt++;
24
                int v;
25
                do {
26
                     v = stk.back();
27
                     stk.pop_back();
               ecc[ecc_cnt].push_back(v);
} while (v != u);
\frac{1}{28}
29
30
          }
31
     };
    for (int i = 1; i <= n; i++) {
   if (!dfn[i]) {</pre>
32
33
34
               tarjan(i, 0);
35
     }
```

## 另一个版本

```
/* DCC @ wrb */
    |// 割点
     std::vector<int> G[N];
     int dfn[N], low[N], is_cut[N], tm, rt;
 5
     void tar(int u) {
 6
           int c = 0;
          dfn[u] = low[u] = ++tm;
for (int v : G[u]) {
 8 9
                if (!dfn[v]) {
                     ++c, tar(v);
10
               low[u] = std::min(low[u], low[v]);

if (low[v] == dfn[u]) is_cut[u] = 1;

} else low[u] = std::min(low[u], dfn[v]);
11
12
13
14
15
           if (u == rt) is_cut[u] = c > 1;
     }
16
17
     int main() {
           int n, m;
19
           std::cin >> n >> m;
20
          for (int x, y; m--; ) {
```

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```
21
                 std::cin >> x >> y;
 22
                 G[x].emplace_back(y);
 23
                 G[y].emplace_back(x);
 24
 25
           for (rt = 1; rt <= n; ++rt) {</pre>
 \frac{1}{26}
                 if (!dfn[rt]) tar(rt);
 27
 \overline{28}
           std::vector<int> cut_v;
 29
30
           for (int i = 1; i <= n; ++i) {
    if (is_cut[i]) cut_v.emplace_back(i);</pre>
 31
 32
            std::cout << cut_v.size() << '\n';
 33
           for (int x : cut_v) std::cout << x << ' ';</pre>
 34
 35
 36
 37
      // 桥
      std::vector<std::pair<int, int>> G[N], brg;
int dfn[N], low[N], rt, tm;
 38
 39
      void tar(int u, int fa, int fr) {
    dfn[u] = low[u] = ++tm;
 40
 41
 42
           for (auto[v, i] : G[u]) {
 43
                 if (!dfn[v]) {
                tar(v, u, i), low[u] = std::min(low[u], low[v]);
} else if (i != fr) {
 44
 45
                      low[u] = std::min(low[u], dfn[v]);
 46
 47
 48
 49
50
           if (u != rt && dfn[u] == low[u]) {
                 brg.emplace_back(std::minmax(u, fa));
 51
 52
 5\overline{3}
54
      int main() {
            int n, m;
 55
            std::cin >> n >> m;
           for (int i = 1, x, y; i <= m; ++i) {
    std::cin >> x >> y;
 56
 57
                G[x].emplace_back(y, i);
G[y].emplace_back(x, i);
 58
 59
 60
           for (rt = 1; rt <= n; ++rt) {
    if (!dfn[rt]) tar(rt, -1, -1);</pre>
 61
 62
 63
 64
           std::sort(begin(brg), end(brg));
           for (auto[u, v] : brg) {
    std::cout << u << ' ' << v << '\n';</pre>
 65
 66
 67
 68
 69
 70
 \begin{array}{c} 71 \\ 72 \end{array}
       // 点双
      std::vector<int> G[N];
 73\\74
      std::vector<std::vector<int>> bcc;
      int dfn[N], low[N], tm, st[N], tp, rt;
 75
76
77
      void tar(int u) {
           dfn[u] = low[u] = ++tm, st[++tp] = u;
           if (G[u].empty()) bcc.push_back({u});
for (int v : G[u]) {
 78
79
                 if (!dfn[v]) {
                      tar(v), low[u] = std::min(low[u], low[v]);
if (low[v] == dfn[u]) {
 80
 81
82
83
                           bcc.push_back({u});
 84
                                bcc.back().emplace_back(st[tp]);
 85
                           } while(st[tp--] != v);
 86
 87
                 } else low[u] = std::min(low[u], dfn[v]);
 89
 90
      int main() {
 91
           std::ios::sync_with_stdio(0);
 92
            std::cin.tie(0);
 93
           int n, m;
 94
            std::cin >> n >> m;
           for (int x, y; m--; ) {
    std::cin >> x >> y;
 95
96
 97
                 G[x].emplace_back(y);
 98
                 G[y].emplace_back(x);
 99
100
           for (int i = 0; i < n; ++i) tar(i);</pre>
           std::cout << bcc.size() << '\n';</pre>
101
           for (auto v : bcc) {
102
                 std::cout << v.size() << ' ';
103
                 for (int x : v) std::cout << x << ' ';
std::cout << '\n';</pre>
104
105
           }
106
    | }
107
```

10.6 2-sat 73

```
108
109
110
      std::vector<std::pair<int, int>> G[N];
std::vector<std::vector<int>> becc;
111
112
       int dfn[N], low[N], tm, st[N], tp;
113
       void tar(int u, int fr) {
    dfn[u] = low[u] = ++tm, st[++tp] = u;
    for (auto[v, i] : G[u]) {
        if (!dfn[v]) {
114
115
116
117
                 tar(v, i), low[u] = std::min(low[u], low[v]);
} else if (i != fr) {
118
119
120
                       low[u] = std::min(low[u], dfn[v]);
121
122
123
            if (dfn[u] == low[u]) {
124
                 becc.emplace_back();
125
                      becc.back().emplace_back(st[tp]);
126
127
                  } while (st[tp--] != u);
128
129
130
       int main() {
131
            int n, m;
132
            std::cin >> n >> m;
            for (int i = 1, x, y; i <= m; ++i) {
    std::cin >> x >> y;
    G[x].emplace_back(y, i);
133
134
135
136
                 G[y].emplace_back(x, i);
137
138
            for (int i = 0; i < n; ++i) {</pre>
139
                  if (!dfn[i]) tar(i, -1);
140
141
            std::cout << becc.size() << '\n';</pre>
142
            for (auto& v : becc) {
143
                 std::cout << v.size() << ' ';
                 for (int x : v) std::cout << x << ' ';
std::cout << '\n';</pre>
144
145
            }
146
       }
147
```

#### 10.6 2-sat

给出 n 个集合,每个集合有 2 个元素,已知若干个数对 (a,b),表示 a 与 b 矛盾.要从每个集合各选择一个元素,判断能否一共选 n 个两两不矛盾的元素.

```
/* two sat */
      auto 2sat = [&](int n, const vpi& v) -> vi {
 3
            /* tarjan */
           vvi e(2 * n);
vi dfn(2 * n), low(2 * n), stk(2 * n), belong(2 * n);
int timestamp = 0, top = 0, scc_cnt = 0;
std::vector<bool> in_stk(2 * n);
 4
5
 67
            auto tarjan = [&](auto&& self, int u) -> void {
 8
                  dfn[u] = low[u] = ++timestamp;
 9
                  stk[++top] = u;
in_stk[u] = true;
10
11
                  for (const auto& v : e[u]) {
12
                       if (!dfn[v]) {
13
                             self(self, v);
14
                       Min(low[u], low[v]);
} else if (in_stk[v]) {
15
16
                             Min(low[u], dfn[v]);
17
18
19
20
21
                  if (dfn[u] == low[u]) {
                       scc_cnt++;
22
                        int v;
23
                             v = stk[top--];
in_stk[v] = false;
belong[v] = scc_cnt;
24
25
26
27
                       } while (v != u);
28
                  }
\frac{1}{29}
30
            for (const auto& [a, b] : v) {
    e[a].push_back(b ^ 1);
    e[b].push_back(a ^ 1);
31
32
33
34
            for (int i = 0; i < 2 * n; i++) {</pre>
35
                  if (!dfn[i]) tarjan(tarjan, i);
```

```
36 | }
37 | vi ans;
38 | for (int i = 0; i < 2 * n; i += 2) {
39 | if (belong[i] == belong[i + 1]) return vi{};
40 | ans.push_back(belong[i] > belong[i + 1] ? i + 1 : i);
41 | }
42 | return ans;
43 |};
```

上述将 i 与 i+1 作为一个集合里的元素, 编号为 0 至 2n-1.

# 10.7 minimum ring

Floyd

```
/* minimum ring */
  \begin{smallmatrix}2&3&4&5\\5&6&7&8\end{smallmatrix}
         auto min_circle = [&]() -> int {
                 vvi dist(n + 1, vi(n + 1, inf));
for (int i = 1; i <= n; i++) {
    for (int j = 1; j <= n; j++) {
        Min(dist[i][j], g[i][j]);
                          dist[i][i] = 0;
  9
                 for (int k = 1; k <= n; k++) {
   for (int i = 1; i < k; i++) {
     for (int j = 1; j < i; j++) {
        Min(ans, dist[i][j] + g[i][k] + g[k][j]);
     }</pre>
10
11
12
13
14
15
16
                          for (int i = 1; i <= n; i++) {</pre>
                                  for (int j = 1; j <= n; j++) {
    Min(dist[i][j], dist[i][k] + dist[k][j]);</pre>
17
18
19
20
                          }
\overline{21}
                 }
22
                 return ans;
\frac{-}{23}
         };
```

tree - diameter

#### 10.8 tree - center of gravity

```
/* center of gravity */
int sum; /* 点权和 */
 1
      vi size(n + 1), weight(n + 1), w(n + 1), depth(n + 1);
std::array<int, 2> centroid = {0, 0};
 3
 4
 5
       auto get_centroid = [&](auto&& self, int u, int fa) -> void {
 \frac{\tilde{6}}{7}
             size[u] = w[u];
weight[u] = 0;
             for (auto v : e[u]) {
   if (v == fa) continue;
 8
 9
                    self(self, v, u);
size[u] += size[v];
10
11
12
                    Max(weight[u], size[v]);
13
             Max(weight[u], sum - size[u]);
if (weight[u] <= sum / 2) {
    centroid[centroid[0] != 0] = u;</pre>
14
15
16
17
             }
      };
18
```

#### 10.9 tree - DSU on tree

给出一课 n 个节点以 1 为根的树, 每个节点染上一种颜色, 询问以 u 为节点的子树中有多少种颜色.

```
1 // Problem: U41492 树上数颜色 int main() {
```

 $10.10 \quad tree - AHU$ 

```
std::ios::sync_with_stdio(false);
5
         std::cin.tie(0);
6
         std::cout.tie(0);
 8
         int n, m, dfn = 0, cnttot = 0;
9
         std::cin >> n;
         vvi e(n + 1);
vi siz(n + 1), col(n + 1), son(n + 1), dfnl(n + 1), dfnr(n + 1), rank(n + 1);
10
11
12
         vi ans(n + 1), cnt(n + 1);
13
14
         for (int i = 1; i < n; i++) {</pre>
              int u, v;
std::cin >> u >> v;
15
16
17
              e[u].push_back(v);
18
              e[v].push_back(u);
19
20
         for (int i = 1; i <= n; i++) {
21
              std::cin >> col[i];
22
23
         auto add = [&](int u) -> void {
24
              if (cnt[col[u]] == 0) cnttot++;
25
              cnt[col[u]]++;
26
27
         auto del = [&](int u) -> void {
28
              cnt[col[u]]--
29
              if (cnt[col[u]] == 0) cnttot--;
30
         auto dfs1 = [%](auto&% self, int u, int fa) -> void {
    dfnl[u] = ++dfn;
31
32
33
              rank[dfn] = u;
34
              siz[u] = 1;
              for (auto v : e[u]) {
   if (v == fa) continue;
35
36
37
                   self(self, v, u);
38
                   siz[u] += siz[v];
39
                   if (!son[u] or siz[son[u]] < siz[v]) son[u] = v;</pre>
40
41
              dfnr[u] = dfn;
42
43
         auto dfs2 = [&](auto&& self, int u, int fa, bool op) -> void {
              for (auto v : e[u]) {
   if (v == fa or v == son[u]) continue;
   self(self, v, u, false);
44
45
46
47
48
              if (son[u]) self(self, son[u], u, true);
49
              for (auto v : e[u]) {
   if (v == fa or v == son[u]) continue;
50
51
                   rep(i, dfnl[v], dfnr[v]) { add(rank[i]); }
52
              add(u);
ans[u] = cnttot;
if (op == false)
53
54
55
56
                   rep(i, dfnl[u], dfnr[u]) { del(rank[i]); }
57
58
         dfs1(dfs1, 1, 0);
59
         dfs2(dfs2, 1, 0, false);
std::cin >> m;
60
61
         for (int i = 1; i <= m; i++) {</pre>
62
63
              int u;
64
              std::cin >> u;
              std::cout << ans[u] << endl;
65
66
67
         return 0;
68
    }
```

#### 10.10 tree - AHU

```
/* AHU */
    std::map<vi, int> mapple;
    std::function<int(vvik, int, int)> tree_hash = [&](vvik e, int u, int fa) -> int {
         vi code;
         if (u == 0) code.push_back(-1);
6
        for (auto v : e[u]) {
             if (v == fa) continue;
7
8
             code.push_back(tree_hash(e, v, u));
9
10
         std::sort(all(code));
11
         int id = mapple.size();
         auto it = mapple.find(code);
if (it == mapple.end()) {
12
13
             mapple[code] = id;
```

```
15 | } else {
16 | id = it->ss;
17 | }
18 | return id;
19 |};
```

#### 10.11 tree - LCA

```
/* LCA */
 3
     int B = 30;
     vvi e(n + 1), fa(n + 1, vi(B));
 4
5
     vi dep(n + 1);
     auto dfs = [&] (auto&& self, int u) -> void {
 6
7
           for (auto v : e[u]) {
                if (v == fa[u][0]) continue;
 8 9
                dep[v] = dep[u] + 1;
fa[v][0] = u;
10
                self(self, v);
           }
11
12
13
     auto init = [&]() -> void {
14
           dep[root] = 1;
          15
16
17
18
19
20
           }
\overline{21}
     };
     init();
auto LCA = [&](int a, int b) -> int {
    if (dep[a] > dep[b]) std::swap(a, b);
    if dep[a] - dep[a];
22 \\ 23 \\ 24 \\ 25
          int d = dep[b] - dep[a];
for (int i = 0; (1 << i) <= d; i++) {
   if (d & (1 << i)) b = fa[b][i];</pre>
26
27
28
29
           if (a == b) return a;
           for (int i = B - 1; i >= 0 and a != b; i--) {
    if (fa[a][i] == fa[b][i]) continue;
30
31
32
33
34
                a = fa[a][i];
                b = fa[b][i];
           }
35
           return fa[a][0];
36
     auto dist = [&](int a, int b) -> int { return dep[a] + dep[b] - dep[LCA(a, b)] * 2; };
```

#### 10.12 tree - heavy light decomposion

对一棵有根树进行如下 4 种操作:

- 1.1 x y z: 将节点 x 到节点 y 的最短路径上所有节点的值加上 z.
- 2.  $2 \times y$ : 查询节点 x 到节点 y 的最短路径上所有节点的值的和.
- 3. 3xz: 将以节点 x 为根的子树上所有节点的值加上 z.
- 4. 4 x: 查询以节点 x 为根的子树上所有节点的值的和.

```
/* heavy light decomposion */
 3
     int cnt = 0;
     vi son(n + 1), fa(n + 1), siz(n + 1), depth(n + 1);
     vi dfn(n + 1), rank(n + 1), top(n + 1), botton(n + 1);
auto dfs1 = [&] (auto&& self, int u) -> void {
 \begin{array}{c} 4\\5\\6\\7 \end{array}
           son[u] = -1, siz[u] = 1;
for (auto v : e[u]) {
 8 9
                if (depth[v] != 0) continue;
                depth[v] = depth[u] + 1;
10
                fa[v] = u;
11
                self(self, v)
12
                siz[u] += siz[v];
13
                if (son[u] == -1 \text{ or } siz[v] > siz[son[u]]) son[u] = v;
14
           }
    };
15
```

10.13 tree - virtual tree 77

```
16
     auto dfs2 = [&](auto&& self, int u, int t) -> void {
          top[u] = t;
dfn[u] = ++cnt;
17
19
          rank[cnt] = u;
          botton[u] = dfn[u];
20
          if (son[u] == -1) return;
\overline{21}
22
          self(self, son[u], t);
\frac{-}{23}
          Max(botton[u], botton[son[u]]);
24
          for (auto v : e[u]) {
   if (v != son[u] and v != fa[u]) {
\overline{25}
26
                    self(self, v, v);
27
                    Max(botton[u], botton[v]);
28
29
30
31
     depth[root] = 1;
32
     dfs1(dfs1, root);
33
     dfs2(dfs2, root, root);
34
35
     /* 求 LCA */
36
     auto LCA = [&](int a, int b) -> int {
          while (top[a] != top[b]) {
37
38
               if (depth[top[a]] < depth[top[b]]) std::swap(a, b);</pre>
               a = fa[top[a]];
39
40
          return (depth[a] > depth[b] ? b : a);
41
42
     };
43
     /* 维护 u 到 v 的路径 */
44
     while (top[u] != top[v]) {
   if (depth[top[u]] < depth[top[v]]) std::swap(u, v);</pre>
45
46
47
          opt(dfn[top[u]], dfn[u]);
48
          u = fa[top[u]];
49
50
     if (dfn[u] > dfn[v]) std::swap(u, v);
51
     opt(dfn[u], dfn[v]);
52
53
     /* 维护 u 为根的子树 */
54
     opt(dfn[u], botton[u]);
55
56
     线段树的 build() 函数中
57
     if(1 == r) tree[u] = {1, 1, w[rank[1]], 0};
58
59
60
     build(1, 1, n);
for (int i = 1; i <= m; i++) {</pre>
61
62
63
          int op, u, v;
          LL k;
64
          std::cin >> op;
65
          if (op == 1) {
66
               std::cin >> u >> v >> k;
while (top[u] != top[v]) {
67
68
                   if (depth[top[u]] < depth[top[v]]) std::swap(u, v);</pre>
69
70
                   modify(1, dfn[top[u]], dfn[u], k);
71
                   u = fa[top[u]];
73
               if (dfn[u] > dfn[v]) std::swap(u, v);
74
75
          modify(1, dfn[u], dfn[v], k);
} else if (op == 2) {
76
77
               std::cin >> u >> v;
               LL ans = 0;
               while (top[u] != top[v]) {
   if (depth[top[u]] < depth[top[v]]) std::swap(u, v);</pre>
78
79
                    ans = (ans + query(1, dfn[top[u]], dfn[u])) % p;
80
81
                    u = fa[top[u]];
82
               if (dfn[u] > dfn[v]) std::swap(u, v);
ans = (ans + query(1, dfn[u], dfn[v])) % p;
std::cout << ans << endl;</pre>
83
84
85
          } else if (op == 3) {
   std::cin >> u >> k;
86
87
88
              modify(1, dfn[u], botton[u], k);
89
          } else {
90
               std::cin >> u;
91
               std::cout << query(1, dfn[u], botton[u]) % p << endl;</pre>
92
     }
93
```

#### 10.13 tree - virtual tree

```
2 3
     auto build_vtree = [&](vi ver) -> void {
          std::sort(all(ver), [&](int x, int y) { return dfn[x] < dfn[y]; });</pre>
          vi stk = {1};
for (auto v : ver) {
 4
 5
               int u = stk.back();
int lca = LCA(v, u);
 6
7
               if (lca != u) {
 8
                    while (dfn[lca] < dfn[stk.end()[-2]]) {
   g[stk.end()[-2]].push_back(stk.back());</pre>
 9
10
11
                          stk.pop_back();
12
                    }
13
                    u = stk.back();
                    if (dfn[lca] != dfn[stk.end()[-2]]) {
14
                         g[lca].push_back(u);
stk.pop_back();
15
16
                          stk.push_back(lca);
17
18
                    } else {
                          g[lca].push_back(u);
19
20
                          stk.pop_back();
\overline{21}
                    }
22
23
24
25
               }
               stk.push_back(v);
          while (stk.size() > 1) {
26
               int u = stk.end()[-2];
27
               int v = stk.back();
               g[u].push_back(v);
28
29
               stk.pop_back();
30
          }
     };
31
```

# 10.14 tree - pseudo tree

```
/* ring detection (directed) */
     vi vis(n + 1), fa(n + 1), ring;
auto dfs = [&](auto&& self, int u) -> bool {
 \begin{array}{c} 3 \\ 4 \\ 5 \end{array}
          vis[u] = 1;
          for (const auto& v : e[u]) {
    if (!vis[v]) {
 6
7
                    fa[v] = u;
 8 9
                    if (self(self, v)) {
                         return true;
                    }
10
11
               } else if (vis[v] == 1) {
                    ring.push_back(v);
for (auto x = u; x != v; x = fa[x]) {
12
13
                         ring.push_back(x);
14
15
16
                    reverse(all(ring));
17
                    return true;
18
               }
19
          }
20
21
          vis[u] = 2;
          return false;
22
23
24
25
     for (int i = 1; i <= n; i++) {
    if (!vis[i]) {</pre>
               if (dfs(dfs, i)) {
\frac{1}{26}
                    // operations //
27
               }
28
          }
29
     }
30
31
32
     /* cycle detection (undirected) */
     vi vis(n + 1), ring;
33
34
     vpi fa(n + 1)
     auto dfs = [&](auto&& self, int u, int from) -> bool {
35
          vis[u] = 1;
36
          for (const auto& [v, id] : e[u]) {
37
               if (id == from) continue;
               if (!vis[v]) {
   fa[v] = {u, id};
38
39
                    if (self(self, v, id)) {
40
41
                         return true;
42
               } else if (vis[v] == 1) {
43
                    ring.push_back(v);
44
                    for (auto x = u; x != v; x = fa[x].ff) {
45
46
                         ring.push_back(x);
47
48
                    return true;
49
               }
```

```
50
51
          vis[u] = 2;
52
          return false;
53
    };
    for (int i = 1; i <= n; i++) {
   if (!vis[i]) {</pre>
54
55
56
               if (dfs(dfs, i, 0)) {
                    // operations //
57
58
59
     }
60
```

### 10.15 tree - divide and conquer on tree

#### 点分治

# 第一个题

一棵  $n \leq 10^4$  个点的树, 边权  $w \leq 10^4$ .  $m \leq 100$  次询问树上是否存在长度为  $k \leq 10^7$  的路径.

```
// 洛谷 P3806 【模板】点分治1

  \begin{array}{c}
    2 \\
    3 \\
    4 \\
    5 \\
    6 \\
    7
  \end{array}

      int main() {
           std::ios::sync_with_stdio(false);
std::cin.tie(0);
            std::cout.tie(0);
           int n, m, k;
std::cin >> n >> m;
 8
10
11
            std::vector<vpi> e(n + 1);
12
            std::map<int, PII> mp;
13
14
            for (int i = 1; i < n; i++) {</pre>
15
                 int u, v, w;
                 std::cin >> u >> v >> w;
16
                 e[u].emplace_back(v, w);
e[v].emplace_back(u, w);
17
18
19
           for (int i = 1; i <= m; i++) {
    std::cin >> k;
    mp[i] = {k, 0};
20
21
22
\frac{-}{23}
\overline{24}
\overline{25}
            /* centroid decomposition */
26
           int top1 = 0, top2 = 0, root;
vi len1(n + 1), len2(n + 1), vis(n + 1);
static std::array<int, 20000010> cnt;
\frac{1}{27}
\frac{1}{28}
29
30
            std::function<int(int, int)> get_size = [&](int u, int fa) -> int {
31
                 if (vis[u]) return 0;
32
                 int ans = 1;
                 for (auto [v, w] : e[u]) {
   if (v == fa) continue;
33
34
35
                       ans += get_size(v, u);
36
37
                 return ans;
38
           };
39
40
            std::function<int(int, int, int, int&)> get_root = [&](int u, int fa, int tot,
41
                                                                                           int& root) -> int {
42
                 if (vis[u]) return 0;
43
                 int sum = 1, maxx = 0;
                 for (auto [v, w] : e[u]) {
44
45
                       if (v == fa) continue;
                       int tmp = get_root(v, u, tot, root);
Max(maxx, tmp);
46
47
48
                       sum += tmp;
49
50
                 Max(maxx, tot - sum);
                 if (2 * maxx <= tot) root = u;</pre>
51
52
                 return sum;
53
54
           std::function<void(int, int, int)> get_dist = [&](int u, int fa, int dist) -> void {
   if (dist <= 10000000) len1[++top1] = dist;</pre>
55
56
                 for (auto [v, w] : e[u]) {
   if (v == fa or vis[v]) continue;
57
58
59
                       get_dist(v, u, dist + w);
60
```

```
61
           };
 62
 63
            auto solve = [&](int u, int dist) -> void {
 64
                 top2 = 0;
 65
                 for (auto [v, w] : e[u]) {
                      if (vis[v]) continue;
 66
                      top1 = 0;
 67
                      get_dist(v, u, w);
for (int i = 1; i <= top1; i++) {
   for (int tt = 1; tt <= m; tt++) {</pre>
 68
 69
70
71
72
73
74
75
76
77
78
                                 int k = mp[tt].ff;
                                 if (k >= len1[i]) mp[tt].ss |= cnt[k - len1[i]];
                      for (int i = 1; i <= top1; i++) {
    len2[++top2] = len1[i];</pre>
                           cnt[len1[i]] = 1;
 80
81
82
83
84
85
                for (int i = 1; i <= top2; i++) cnt[len2[i]] = 0;</pre>
           };
            std::function<void(int)> divide = [&](int u) -> void {
                 vis[u] = cnt[0] = 1;
                 solve(u, 0);
 86
                 for (auto [v, w] : e[u]) {
 87
                      if (vis[v]) continue;
 88
                      get_root(v, u, get_size(v, u), root);
 89
                      divide(root);
 90
                }
 91
           };
 92
            get_root(1, 0, get_size(1, 0), root);
 93
 94
           divide(root);
 95
 96
           for (int i = 1; i <= m; i++) {
   if (mp[i].ss == 0) {
     std::cout << "NAY" << endl;</pre>
 97
 98
99
                 } else {
100
                      std::cout << "AYE" << endl;
101
                 }
102
           }
103
104
           return 0;
105
```

#### 第二个题

一棵  $n \le 4 \times 10^4$  个点的树, 边权  $w \le 10^3$ . 询问树上长度不超过  $k \le 2 \times 10^4$  的路径的数量.

```
12
     // 洛谷 P4178 Tree
 \frac{3}{4} \\ \frac{4}{5} \\ \frac{6}{7}
     int main() {
          std::ios::sync_with_stdio(false);
std::cin.tie(0);
          std::cout.tie(0);
 8
          int n, k;
std::cin >> n;
 9
10
          std::vector<vpi> e(n + 1);
          for (int i = 1; i < n; i++) {</pre>
11
12
               int u, v, w;
std::cin >> u >> v >> w;
13
14
               e[u].emplace_back(v, w);
15
               e[v].emplace_back(u, w);
16
17
          std::cin >> k;
18
19
           /* centroid decomposition */
20 \\ 21 \\ 22 \\ 23 \\ 24
           int root;
          vi len, vis(n + 1);
          std::function<int(int, int)> get_size = [&](int u, int fa) -> int {
               if (vis[u]) return 0;
25
               int ans = 1:
               for (auto [v, w] : e[u]) {
   if (v == fa) continue;
\frac{1}{26}
27
28
29
30
31
32
                    ans += get_size(v, u);
               }
               return ans;
          };
33
          std::function<int(int, int, int, int&)> get_root = [&](int u, int fa, int tot,
34
                                                                                  int& root) -> int {
35
               if (vis[u]) return 0;
```

10.16 tree - matrix tree

```
36
               int sum = 1, maxx = 0;
37
               for (auto [v, w] : e[u]) {
38
                   if (v == fa) continue;
39
                   int tmp = get_root(v, u, tot, root);
                   maxx = std::max(maxx, tmp);
40
                   sum += tmp;
41
42
              }
43
              maxx = std::max(maxx, tot - sum);
              if (2 * maxx <= tot) root = u;
44
45
              return sum;
46
47
48
          std::function<void(int, int, int)> get_dist = [&](int u, int fa, int dist) -> void {
49
               len.push_back(dist);
              for (auto [v, w] : e[u]) {
    if (v == fa || vis[v]) continue;
50
51
52
                   get_dist(v, u, dist + w);
53
54
         };
55
          auto solve = [&](int u, int dist) -> int {
56
              len.clear();
get_dist(u, 0, dist);
57
58
59
               std::sort(all(len));
60
               int ans = 0;
              for (int 1 = 0, r = len.size() - 1; 1 < r;) {
   if (len[1] + len[r] <= k) {</pre>
61
62
63
                        ans += r - 1++;
                   } else {
64
65
                        r--;
66
67
68
              return ans;
69
70
71
          std::function<int(int)> divide = [&](int u) -> int {
72
73
              vis[u] = true
               int ans = solve(u, 0);
74
               for (auto [v, w] : e[u]) {
75
                   if (vis[v]) continue;
76
77
                   ans -= solve(v, w);
                   get_root(v, u, get_size(v, u), root);
ans += divide(root);
78
79
80
              return ans;
81
          };
82
          get_root(1, 0, get_size(1, 0), root);
std::cout << divide(root) << endl;</pre>
83
84
85
86
          return 0:
87
     }
```

81

#### 10.16 tree - matrix tree

```
const int N=33,M=152599,P=998244353;
     int qpow(int a,int b=P-2){
 3
           int r=1;for(;b;b>>=1,a=111*a*a%P)if(b&1)r=111*r*a%P;return r;
     struct T{int x,y,z;T(int a=0,int b=0,int c=0):x(a),y(b),z(c){}}e[N*N];
 6
     struct F{
           int a,b;
 8
          F():a(),b(){}
          F().a(),b();

F(int x,int y):a(x),b(y){}

F operator+(const F&_)const{return F((a+_.a)\%P,(b+_.b)\%P);}

F operator+=(const F&_)freturn *this=*this+_;}

F operator-(const F&_)const{return F((a-_.a+P)\%P,(b-_.b+P)\%P);}
 9
10
11
12
          F operator-=(const F&_){return *this=*this-_;}
F operator*(const F&_)const{return F((111*a*_.b+111*b*_.a)%P,111*b*_.b%P);}
13
14
15
           F operator*=(const F&_){return *this=*this*_;}
           int operator&()const{return b?2:(a?1:0);}
16
17
           bool operator!()const{return !a&&!b;}
18
           F operator~()const{
19
                int d=qpow(b);
                return F((P-111*a*d%P*d%P)%P,d);
20
21
22
\frac{-}{23}
     int fa[N],phi[M],n,m;
24
     int gf(int x){return x==fa[x]?x:fa[x]=gf(fa[x]);}
25
          cal(int p){
F a[N][N],d,iv,z=F(0,1);
26
27
           int i,j,k,l,x=0;iota(fa,fa+n+1,0);
```

```
28
           for(i=1;i<=m;++i)if(e[i].z%p==0){</pre>
29
                if((j=gf(e[i].x))!=(k=gf(e[i].y)))fa[j]=k;
                j=e[i].x,k=e[i].y,l=e[i].z,++x;
a[j][k]-=F(1,1),a[k][j]-=F(1,1);
30
31
32
                a[j][j]+=F(1,1),a[k][k]+=F(1,1);
33
          for(j=0,i=1;i<=n;++i)if(fa[i]==i)++j;
if(j>1 || x<n-1)return 0;</pre>
34
35
36
37
           for(i=1;i<n;++i){</pre>
               for(k=i,j=i+1;j<n;++j)if(&a[j][i]>&a[k][i])k=j;
if(k!=i)swap(a[i],a[k]),z*=F(0,P-1);
38
39
                if(!a[i][i])return 0;
                for(z*=a[i][i],iv=~a[i][i],j=i;j<n;++j)a[i][j]*=iv;
for(j=i+1;j<n;++j)for(d=a[j][i],k=i;k<n;++k)a[j][k]-=a[i][k]*d;
40
41
42
43
          return z.a;
44
45
     void work(){
46
          int h=0,i,j,x,y,z;
47
          for(cin>>n>m, i=1; i<=m; ++i) cin>>x>>y>>z, e[i]=T(x,y,z), h=max(h,z);
48
          iota(phi+1,phi+h+1,1);
49
           for(i=1;i<=h;++i)for(j=i<<1;j<=h;j+=i)phi[j]=(phi[j]-phi[i]+P)%P;</pre>
           for(z=0,i=1;i<=h;++i)z=(z+1ll*phi[i]*cal(i)%P)%P;</pre>
50
51
           cout<<z<'\n';
52
```

# 10.17 Prefür sequence

```
\* prefur @ wrb *\
2
   for(int i=1;i<n;i++)cin>>fa[i],d[fa[i]]++;
   3
4
5
       p[i]=fa[j];
6
       while(i<n-1&&!--d[p[i]]&&j>p[i])p[i+1]=fa[p[i]],i++;
   }
9
    10
11
   for(int i=1;i<n-1;i++)cin>>p[i],d[p[i]]++;
12
   p[n-1]=n;
   for(int i=1,j=1;i<n;i++,j++){
    while(d[j])j++;</pre>
13
14
15
       fa[j]=p[i];
16
       while(i < n\&\&! - d[p[i]]\&\&j > p[i])fa[p[i]]=p[i+1],i++;
   }
```

#### 10.18 network flow - maximal flow

Dinic

```
/* dinic */
      struct edge {

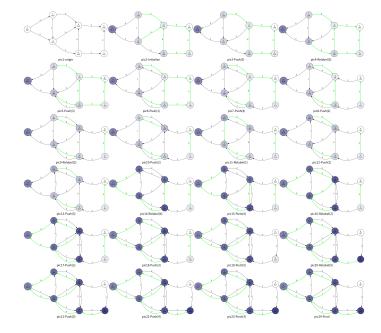
    \begin{array}{r}
      23 \\
      45 \\
      67 \\
      89
    \end{array}

            int from, to;
            LL cap, flow;
            edge(int u, int v, LL c, LL f) : from(u), to(v), cap(c), flow(f) {}
     };
      struct Dinic {
10
            int n, m = 0, s, t;
11
            std::vector<edge> e;
12
            vi g[N];
13
            int d[N], cur[N], vis[N];
14
15
            void init(int n) {
                  for (int i = 0; i < n; i++) g[i].clear();</pre>
16
17
                  e.clear();
18
                 m = 0;
19
\overline{20}
21
22
           void add(int from, int to, LL cap) {
   e.push_back(edge(from, to, cap, 0));
   e.push_back(edge(to, from, 0, 0));
23
24
25
                  g[from].push_back(m++);
                  g[to].push_back(m++);
26
```

```
27
28
29
             bool bfs() {
                   for (int i = 1; i <= n; i++) {
    vis[i] = 0;
30
31
32
                   std::queue<int> q;
q.push(s), d[s] = 0, vis[s] = 1;
while (!q.empty()) {
    int = p fount();
}
33
34
35
36
                         int u = q.front();
                         q.pop();
                         for (int i = 0; i < g[u].size(); i++) {
    edge& ee = e[g[u][i]];</pre>
37
38
                                if (!vis[ee.to] and ee.cap > ee.flow) {
   vis[ee.to] = 1;
39
40
41
                                      d[ee.to] = d[u] + 1;
42
                                      q.push(ee.to);
43
44
                         }
45
46
                   return vis[t];
47
48
            LL dfs(int u, LL now) {
    if (u == t || now == 0) return now;
49
50
                   LL flow = 0, f;

for (int& i = cur[u]; i < g[u].size(); i++) {
    edge& ee = e[g[u][i]];
    edge& er = e[g[u][i] ^ 1];
    if (d[u] + 1 == d[ee.to] and (f = dfs(ee.to, std::min(now, ee.cap - ee.flow))) > 0) {
51
52
53
54
55
56
                                ee.flow += f, er.flow -= f;
                               flow += f, now -= f;
if (now == 0) break;
57
58
59
                         }
60
                   }
61
                   return flow;
62
63
64
            LL dinic() {
                   LL ans = 0;
65
66
                   while (bfs()) {
                         for (int i = 1; i <= n; i++) cur[i] = 0;
ans += dfs(s, INF);</pre>
67
68
69
70
                   return ans;
71
       } maxf;
```

#### **HLPP**



```
/* 点的连边编号 */
            std::vector<int> g[N];
 6
7
8
            std::priority_queue<node> q;
            std::queue<int> qq;
            bool vis[N];
 9
10
           int cnt[N];
11
12
           void init() {
13
                 e.clear();
                 nd.clear();
for (int i = 0; i <= n + 1; i++) {</pre>
14
15
                       nd.pushback(node(inf, i, 0));
16
17
                       g[i].clear();
18
                       vis[i] = false;
19
                 }
           }
20 \\ 21 \\ 22 \\ 23 \\ 24
           void add(int u, int v, LL w) {
                 e.pushback(edge(u, v, w));
e.pushback(edge(v, u, 0));
g[u].pushback(m++);
25
                 g[v].pushback(m++);
26
27
28
\overline{29}
           void bfs() {
30
                 nd[t].hight = 0;
31
                 qq.push(t);
                 while (!qq.empty()) {
   int u = qq.front();
32
33
                       qq.pop();
vis[u] = false;
34
35
                       for (auto j : g[u]) {
   int v = e[j].to;
   if (e[j].cap == 0 && nd[v].hight > nd[u].hight + 1) {
      nd[v].hight = nd[u].hight + 1;
      if (vis[v] == false) {
36
37
38
39
40
41
                                        qq.push(v);
42
                                        vis[v] = true;
                                  }
43
44
                             }
45
                       }
46
                 }
47
                 return;
48
\overline{49}
50
           void _push(int u) {
51
52
53
54
                 for (auto j : g[u]) {
    edge &ee = e[j], &er = e[j ^ 1];
                       int v = ee.to;
                       node &nu = nd[u], &nv = nd[v];
55
                       if (ee.cap && nv.hight + 1 == nu.hight) {
                             LL flow = std::min(ee.cap, nu.flow);
ee.cap -= flow, er.cap += flow;
56
57
58
                             nu.flow -= flow, nv.flow += flow;
                             if (vis[v] == false && v != t && v != s) {
59
60
                                   q.push(nv);
                                   vis[v] = true;
61
62
63
                             if (nu.flow == 0) break;
64
                       }
                 }
65
66
           }
67
68
           void relabel(int u) {
69
                 nd[u].hight = inf;
                 for (auto j : g[u]) {
   int v = e[j].to;
   if (e[j].cap && nd[v].hight + 1 < nd[u].hight) {</pre>
70
71
72
73
74
75
76
77
78
79
80
                             nd[u].hight = nd[v].hight + 1;
                 }
           }
           LL hlpp() {
                 bfs();
                 if (nd[s].hight == inf) return 0;
81
                 nd[s].hight = n;
                 for (int i = 1; i <= n; i++) {
82
83
                       if (nd[i].hight < inf) cnt[nd[i].hight]++;</pre>
84
                 for (auto j : g[s]) {
   int v = e[j].to;
   int flow = e[j].cap;
85
86
87
88
                       if (flow) {
                            e[j].cap -= flow, e[j ^ 1].cap += flow;
nd[s].flow -= flow, nd[v].flow += flow;
if (vis[v] == false && v != s && v != t) {
89
90
91
92
                                  q.push(nd[v]);
```

```
93
                             vis[v] = true;
 94
                        }
 95
                    }
 96
 97
               while (!q.empty()) {
98
                    int u = q.top().id;
                    q.pop();
vis[u] = false;
99
100
101
                    _push(u);
                    if (nd[u].flow) {
102
103
                         cnt[nd[u].hight]--;
104
                         if (cnt[nd[u].hight] == 0) {
                             for (int i = 1; i <= n; i++) {
    if (i != s && i != t && nd[i].hight > nd[u].hight && nd[i].hight < n + 1) {</pre>
105
106
107
                                      nd[i].hight = n + 1;
108
109
                             }
110
                        }
                        relabel(u);
111
112
                         cnt[nd[u].hight]++;
113
                         q.push(nd[u]);
114
                         vis[u] = true;
115
116
117
               return nd[t].flow;
118
119
      } maxf;
```

## 10.19 network flow - minimum cost flow

#### Dinic + SPFA

```
/* Dinic + SPFA */
      struct edge {
 3
           int from, to;
           LL cap, cost;
 5
 6
           edge(int u, int v, LL c, LL w) : from(u), to(v), cap(c), cost(w) {}
     };
 8 9
      const int N = 2000;
10
11
      struct MCMF {
12
           int n, m = 0, s, t;
13
           std::vector<edge> e;
           vi g[N];
int cur[N], vis[N];
14
15
           LL dist[N], minc;
16
17
18
           void init(int n) {
19
                for (int i = 0; i < n; i++) g[i].clear();</pre>
                e.clear();
20
\frac{20}{21}
                minc = m = 0;
22
23
24
           void add(int from, int to, LL cap, LL cost) {
                e.push_back(edge(from, to, cap, cost));
e.push_back(edge(to, from, 0, -cost));
25
26
27
                g[from] .push_back(m++);
                g[to].push_back(m++);
29
30
31
32
           bool spfa() {
                for (int i = 1; i <= n; i++) {
    dist[i] = INF, cur[i] = 0;</pre>
33
34
                std::queue<int> q;
q.push(s), dist[s] = 0, vis[s] = 1;
35
36
37
                while (!q.empty()) {
                      int u = q.front();
q.pop();
38
39
                     q.pop();
vis[u] = 0;
for (int j = cur[u]; j < g[u].size(); j++) {
    edge& ee = e[g[u][j]];
    int u = ce to:</pre>
40
41
42
43
                           if (ee.cap && dist[v] > dist[u] + ee.cost) {
    dist[v] = dist[u] + ee.cost;
44
45
46
                                 if (!vis[v]) {
47
                                      q.push(v);
48
                                      vis[v] = 1;
49
```

```
50
                         }
51
                    }
52
               }
53
               return dist[t] != INF;
54
55
          LL dfs(int u, LL now) {
   if (u == t) return now;
56
57
58
59
               vis[u] = 1;
               LL ans = 0;
               for (int& i = cur[u]; i < g[u].size() && ans < now; i++) {
  edge &ee = e[g[u][i]], &er = e[g[u][i] ^ 1];</pre>
60
61
62
                    int v = ee.to;
                    if (!vis[v] && ee.cap && dist[v] == dist[u] + ee.cost) {
63
64
                         LL f = dfs(v, std::min(ee.cap, now - ans));
65
                         if (f) {
66
                              minc += f * ee.cost, ans += f;
67
                              ee.cap -= f;
68
                              er.cap += f;
69
70
71
72
73
74
75
76
77
78
79
                    }
               }
               vis[u] = 0;
               return ans;
          PLL mcmf() {
               LL \max f = 0;
               while (spfa()) {
                    LL tmp;
80
                    while ((tmp = dfs(s, INF))) maxf += tmp;
81
82
               return std::make_pair(maxf, minc);
83
84
     } minc_maxf;
```

#### Primal-Dual 原始对偶算法

```
/* primal dual */
 \begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
     struct edge {
           int from, to;
           LL cap, cost;
 \frac{6}{7}
           edge(int u, int v, LL c, LL w) : from(u), to(v), cap(c), cost(w) {}
 9
     struct node {
10
           int v, e;
\overline{11}
12
           node(int _v = 0, int _e = 0) : v(_v), e(_e) {}
13
     };
14
15
     const int maxn = 5000 + 10;
16
17
      struct MCMF {
           int n, m = 0, s, t;
18
19
           std::vector<edge> e;
\frac{1}{20}
           vi g[maxn];
           int vis[maxn];
          LL dis[maxn], h[maxn]; node p[maxn * 2];
22
\frac{1}{23} 24
\frac{25}{26}
           void add(int from, int to, LL cap, LL cost) {
    e.push_back(edge(from, to, cap, cost));
\overline{27}
                e.push_back(edge(to, from, 0, -cost));
28
                g[from].push_back(m++);
29
30
31
                g[to].push_back(m++);
32
33
34
           bool dijkstra() {
                std::priority_queue<PIL, std::vector<PIL>, std::greater<PIL>> q;
for (int i = 1; i <= n; i++) {
    dis[i] = INF;</pre>
35
36
                     vis[i] = 0;
37
                dis[s] = 0;
38
                q.push({0, s});
39
                 while (!q.empty()) {
40
41
                     auto u = q.top().ss;
42
                      q.pop();
                      if (vis[u]) continue;
43
44
                      vis[u] = 1;
45
                     for (auto i : g[u]) {
```

```
46
                         edge ee = e[i];
47
                         int v = ee.to;
48
                         LL nc = ee.cost + h[u] - h[v];
                        if (ee.cap and dis[v] > dis[u] + nc) {
    dis[v] = dis[u] + nc;
49
50
                             p[v] = node(u, i);
51
                             if (!vis[v]) q.push({dis[v], v});
52
                        }
53
                   }
54
               }
55
               return dis[t] != INF;
56
57
          }
58
59
          void spfa() {
               std::queue<int> q;
60
61
               for (int i = 1; i <= n; i++) h[i] = INF;</pre>
               h[s] = 0, vis[s] = 1;
62
               q.push(s);
63
               while (!q.empty()) {
   int u = q.front();
64
65
66
                    q.pop();
67
                    vis[u] = 0;
                    for (auto i : g[u]) {
68
69
                         edge ee = e[i];
                         int v = ee.to;
70
71
72
73
74
75
                         if (ee.cap and h[v] > h[u] + ee.cost) {
                             h[v] = h[u] + ee.cost;
                             if (!vis[v]) {
                                  vis[v] = 1;
                                  q.push(v);
76
77
78
79
                             }
                        }
                   }
               }
80
81
         }
82
         PLL mcmf() {
83
               LL \max f = 0, \min c = 0;
84
               spfa();
85
               while (dijkstra()) {
86
                    LL minf = INF;
                   for (int i = 1; i <= n; i++) h[i] += dis[i];
for (int i = t; i != s; i = p[i].v) minf = std::min(minf, e[p[i].e].cap);
87
88
                    for (int i = t; i != s; i = p[i].v) {
89
                        e[p[i].e].cap -= minf;
e[p[i].e ^ 1].cap += minf;
90
91
92
93
                   maxf += minf;
94
                   minc += minf * h[t];
95
               }
96
               return std::make_pair(maxf, minc);
97
98
     } minc_maxf;
```

#### 存在负环的网络

流满后推流, 转化为上下界网络流.

## 10.20 network flow - minimal cut

最小割解决的问题是将图中的点集 V 划分成 S 与 T, 使得 S 与 T 之间的连边的容量总和最小.

#### 最大流最小割定理

网络中s到t的最大流流量的值等于所要求的最小割的值, 所以求最小割只需要跑 Dinic 即可.

#### 获得 S 中的所有点

在 Dinic 的 bfs 函数中,每次将所有点的 d 数组值改为无穷大,最后跑完最大流之后 d 数组不为无穷大的就是和源点一起在 S 集合中的点.

## 例子

最小割的本质是对图中点集进行 2-划分, 网络流只是求解答案的手段.

- 1. 在图中花费最小的代价断开一些边使得源点 s 无法流到汇点 t. 直接跑最大流就得到了答案.
- 2. 在图中删除最少的点使得源点 s 无法流到汇点 t. 对每个点进行拆点, 在 i 与 i' 之间建立容量为 1 的有向边.

# 10.21 network flow - upper / lower bound

## 无源汇上下界可行流

每条有向边有流量的上下界限制,但整张图并未确定源点与汇点.如果存在满足每个点的流入量等于流出量,且每条边的流量满足其上下界限制的流,称之为可行流.

- 1. 将每条边先给予大小为下界的流量,
- 2. 对每个点计算总流入量  $in_u$  与总流出量  $out_u$  的值,
- 3. 建立超级源点到每个点, 容量大小为  $\max\{0, \text{in}_u \text{out}_u\}$  的边; 建立每个点到超级汇点, 容量大小为  $\max\{0, \text{out}_u \text{in}_u\}$ ,
- 4. 跑从超级源点到超级汇点的最大流,如果超级源点每条边都流满意味着存在可行流. 将每条边的流量加上预先给每条边设置的下界流量即为可行流方案.

# 有源汇上下界可行流

1. 建立汇点 t 到源点 s 的, 容量为 ∞ 的有向边, 将其转化为无源汇的问题.

## 有源汇上下界最大流

- 1. 建立汇点 t 到源点 s 的, 容量为 ∞ 的有向边, 将其转化为无源汇的问题,
- 2. 跑上下界可行流, 可行流流量为边  $t \xrightarrow{\infty} s$  的流量.
- 3. 删除  $t \xrightarrow{\infty} s$  的边, 再残量网络上跑 s 到 t 的最大流,
- 4. 答案等于可行流流量 + 最大流流量.

#### 有源汇上下界最小流

- 1. 建立汇点 t 到源点 s 的, 容量为  $\infty$  的有向边, 将其转化为无源汇的问题,
- 2. 跑上下界可行流, 可行流流量为边  $t \xrightarrow{\infty} s$  的流量.
- 3. 删除  $t \xrightarrow{\infty} s$  的边, 再残量网络上跑 t 到 s 的最大流,
- 4. 答案等于可行流流量 最大流流量.

#### 有源汇上下界最小费用可行流

- 1. 按下界流满并计算费用,
- 2. 类似有源汇上下界最大流建图, 跑超级源点到超级汇点的费用流,
- 3. 答案等于按下界的费用加上后续残量网络.

## 10.22 network flow - other versions

```
/* dinic @ wrb */
       template<typename T, T inf = numeric_limits<T>::max()>
struct Max_Flow {
 2
 3
 4
               vector<int> he, cur, d, ne, to;
 5
               vector<T> c;
 6
               int s, t;
              Max_Flow(int m) : he(m, -1), s(-1), t(-1) {}
void add(int x, int y, T z = inf, T w = 0) {
    // cerr << x << ' ' << y << ' ';
    // if (z == inf) cerr << "inf\n";
    // else cerr << z << '\n';</pre>
 7
 8
 9
10
11
                     ne.emplace_back(he[x]);
he[x] = ne.size() - 1;
to.emplace_back(y);
13
14
                      c.emplace_back(z);
15
16
                      ne.emplace_back(he[y]);
                      he[y] = ne.size() - 1;
17
18
                      to.emplace_back(x);
19
                      c.emplace_back(w);
20
21
               int bfs() {
22
                      queue<int> q;
                     queue<int> q;
d.assign(he.size(), -1);
q.emplace(s), d[s] = 0;
for (; q.size(); q.pop()) {
   int u = q.front(), v;
   for (int i = he[u]; ~i; i = ne[i]) {
      if (c[i] && d[v = to[i]] == -1) {
        d[v] = d[u] + 1;
      if (v == t) return 1.
23
24
25
26
\overline{27}
28
29
30
                                            if (v == t) return 1;
31
                                            q.emplace(v);
32
                                     }
33
                             }
34
35
                      return 0;
36
              T dfs(int u, T fl) {
   if (u == t) return fl;
37
38
                     if (u -- c, zc...
T z = 0, r;
for (int& i = cur[u], v; ~i; i = ne[i]) {
    if (c[i] && d[v = to[i]] == d[u] + 1) {
        -- dfe(v min(fl, c[i]));
}
39
40
41
                                    r = dfs(v, min(fl, c[i]));
if (r == 0) d[v] = -1;
else {
42
43
44
45
                                            fl = r, z += r, c[i] = r, c[i ^ 1] += r;
46
                                            if (fl == 0) return z;
47
                             }
48
49
                      }
50
                      return z;
51
              T dinic(int _s, int _t) {
52
                      Tz = 0;
53
                      for (s = _s, t = _t; bfs();) {
   cur = he, z += dfs(s, inf);
54
55
                      }
56
57
                      return z;
58
59
       };
```

```
/* bounded flow @ lys */
#include <iostream>
#include <cstdio>
#include <algorithm>
#include <queue>
#include <vector>
#include <vector>
#define int long long
using namespace std;
```

```
10
    const int maxn = 50020;
11
     const int inf = 1e18;
13
     struct Dinic_limit
14
          int st, sgn; // st = 1 表示有源汇; sgn 表示最大(1)最小(-1)流
15
16
          struct edge
17
18
               int x, y, cap, flow, cost;
19
20 \\ 21 \\ 22 \\ 23 \\ 24
          int deg[maxn]; // rd - cd
          vector<int> e[maxn];
          vector<edge> edges;
          int mx;
          int mcmf;
25
          void add(int x, int y, int cap, int cost)
26
27
               edges.push_back({x, y, cap, 0, cost});
edges.push_back({y, x, 0, 0, -cost});
mx = max({mx, x, y});
28
\overline{29}
30
               int m = edges.size();
31
               e[x].push_back(m - 2), e[y].push_back(m - 1);
32
33
          void add(int x, int y, int 1, int r, int cost)
34
35
                if (cost >= 0)
                    add(x, y, r - 1, cost), deg[y] += 1, deg[x] -= 1, mcmf += 1 * cost;
36
37
38
                    add(y, x, r - 1, -cost), deg[y] += r, deg[x] -= r, mcmf += r * cost;
39
          int s, t;
int vis[maxn], dis[maxn];
40
41
42
          bool spfa()
43
44
               queue<int> q;
               fill(vis, vis + mx + 1, 0), fill(dis, dis + mx + 1, inf);
dis[s] = 0, q.push(s), vis[s] = 1;
45
46
47
               while (!q.empty())
48
                    int x = q.front();
q.pop(), vis[x] = 0;
for (int i : e[x])
49
50
51
52
                         auto k = edges[i];
if (k.cap - k.flow > 0 && k.cost + dis[x] < dis[k.y])</pre>
53
54
55
                               dis[k.y] = dis[x] + k.cost;
56
57
                               if (!vis[k.y])
58
                                    q.push(k.y), vis[k.y] = 1;
59
                         }
60
                    }
               }
61
               return dis[t] != inf;
62
63
64
          int cur[maxn];
          int dfs(int x, int lim)
65
66
               if (x == t || lim == 0)
67
68
                    return lim;
69
               vis[x] = 1;
70
71
72
73
74
75
76
77
78
79
               int res = 0, f;
               for (int &i = cur[x]; i < (int)e[x].size(); i++)</pre>
                     auto &k = edges[e[x][i]];
                    if (!vis[k.y] && k.cost + dis[x] == dis[k.y] && (f = dfs(k.y, min(lim, k.cap - k.flow))))
    res += f, lim -= f, k.flow += f, edges[e[x][i] ^ 1].flow -= f, mcmf += f * k.cost;
if (lim == 0)
                         break:
               }
               vis[x] = 0;
80
               return res;
81
82
          int dinic(int s_, int t_)
83
84
               int ss = mx + 1, tt = ss + 1;
85
               int tot = 0;
86
               for (int i = 1; i <= mx; i++)</pre>
87
                    if (deg[i] > 0)
                    add(ss, i, deg[i], 0), tot += deg[i];
else if (deg[i] < 0)
89
                         add(i, tt, -deg[i], 0);
90
91
               if (st)
               add(t_, s_, 0, inf, 0);
s = ss, t = tt;
92
93
               int res = 0;
94
95
               while (spfa())
96
                    fill(cur, cur + mx + 1, 0), res += dfs(s, inf);
```

```
// cerr << res << " " << tot << endl;
 97
 98
                if (res != tot)
 99
                     return -1;
100
                 if (st == 0)
101
                     return 1;
102
                res = -edges.back().flow;
                edges.back().cap = edges.back().flow = 0;
edges[edges.size() - 2].cap = edges[edges.size() - 2].flow = 0;
s = s_, t = t_;
if (sgn == -1)
103
104
105
106
                     swap(s, t);
107
108
                while (spfa())
109
                     fill(cur, cur + mx + 1, 0), res += sgn * dfs(s, inf);
110
                return res;
111
112
           void clear()
113
                for (int i = 0; i <= mx; i++)
    e[i].clear(), deg[i] = 0;</pre>
114
115
116
                edges.clear();
                mx = 0, mcmf = 0;
117
118
           Dinic_limit(int st_ = 1, int sgn_ = 1) { st = st_, sgn = sgn_; } // st = 1 表示有源汇; sgn 表示最大(1)最小(-1)流
119
120
121
           // 使用时调用 dinic 函数,返回-1表示无解,否则返回最大/最小流
      };
122
123
124
      Dinic_limit G(1, 1);
125
126
      signed main()
127
           ios::sync_with_stdio(false), cin.tie(0);
int n, m, S, T;
cin >> n >> m >> S >> T;
128
129
130
131
           for (int i = 0; i < m; i++)</pre>
132
                int s, t, 1, r, c;
cin >> s >> t >> 1 >> r >> c;
133
134
                G.add(s, t, 1, r, c);
135
           }
136
137
           int res = G.dinic(S, T);
           if (res == -1)
138
                cout << -1 << endl;
139
140
           {
141
                cout << res << " " << G.mcmf << endl;</pre>
142
143
           }
144
      }
```

## 10.23 matching - matching on bipartite graph

# 二分图最大匹配

#### Kuhn-Munkres

时间复杂度:  $O(n^3)$ .

```
/* Kuhn-Munkres */
     auto KM = [&](int n1, int n2, vvi e) -> std::pair<vi, vi> {
         vi vis(n2 + 1);
vi l(n1 + 1, -1), r(n2 + 1, -1);
std::function<bool(int)> dfs = [&](int u) -> bool {
 3
 4
 5
 6
7
               for (auto v : e[u]) {
                    if (!vis[v]) {
                        vis[v] = 1;
 8
 9
                        if (r[v] == -1 \text{ or } dfs(r[v])) {
10
                             r[v] = u;
11
                             return true;
                        }
                   }
13
              }
14
15
              return false;
16
          for (int i = 1; i <= n1; i++) {</pre>
17
              std::fill(all(vis), 0);
18
19
              dfs(i);
20
21
          for (int i = 1; i <= n2; i++) {</pre>
22
               if (r[i] == -1) continue;
```

```
23 | 1[r[i]] = i;

24 | }

25 | return {1, r};

26 |};

27 | auto [mchl, mchr] = KM(n1, n2, e);

28 | std::cout << mchl.size() - std::count(all(mchl), -1) << endl;
```

#### Hopcroft-Karp

据说时间复杂度是  $O(m\sqrt{n})$  的, 但是快的飞起.

```
/* Hopcroft-Karp */

  \begin{array}{c}
    2 \\
    3 \\
    4 \\
    5 \\
    6 \\
    7
  \end{array}

       vpi e(m);
       auto hopcroft_karp = [&] (int n, int m, vpi& e) -> std::pair<vi, vi> {
   vi g(e.size()), l(n + 1, -1), r(m + 1, -1), d(n + 2);
   for (auto [u, v] : e) d[u]++;
   std.:pair<[-2,1](), d(n + 2);</pre>
              std::partial_sum(all(d), d.begin());
for (auto [u, v] : e) g[--d[u]] = v;
 8
              for (vi a, p, q(n + 1);;) {
    a.assign(n + 1, -1);
                    p.assign(n + 1, -1);
int t = 1;
10
11
12
                     for (int i = 1; i <= n; i++) {
13
                            if (1[i] == -1) {
14
                                  q[t++] = a[i] = p[i] = i;
15
16
                    bool match = false;
for (int i = 1; i < t; i++) {
   int u = q[i];</pre>
17
18
19
20
21
22
23
24
25
                            if (l[a[\bar{u}]] != -1) continue;
                            for (int j = d[u]; j < d[u + 1]; j++) {
                                  int v = g[j];
if (r[v] == -1) {
                                         while (v != -1) {
    r[v] = u;
\frac{26}{27}
                                                 std::swap(1[u], v);
                                                u = p[u];
28
29
30
31
32
33
                                          }
                                         match = true;
                                          break;
                                   if (p[r[v]] == -1) {
                                         q[t++] = v = r[v];
34
                                         p[v] = u;
35
                                          a[v] = a[u];
36
37
38
39
                     if (!match) break;
40
41
              return {1, r};
      };
42
```

### 二分图最大权匹配

## **Kuhn-Munkres**

注意是否为完美匹配,非完美选0,完美选-INF. (存疑)

```
\begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
      auto KM = [&](int n, vvl e) -> std::tuple<LL, vi, vi> {
            vl la(n + 1), lb(n + 1), pp(n + 1), vx(n + 1);
vi l(n + 1, -1), r(n + 1, -1);
             vi va(n + 1), vb(n + 1);
            LL delta;
 6
7
8
            auto bfs = [&](int x) -> void {
  int a, y = 0, y1 = 0;
  std::fill(all(pp), 0);
  std::fill(all(vx), INF);
 9
10
11
                   r[y] = x;
12
                   do {
13
                          a = r[y], delta = INF, vb[y] = 1;
                         for (int b = 1; b <= n; b++) {
    if (!vb[b]) {</pre>
14
15
16
                                      if (vx[b] > la[a] + lb[b] - e[a][b]) {
```

```
vx[b] = la[a] + lb[b] - e[a][b];
17
18
19
                                                 pp[b] = y;
\begin{array}{c} 20 \\ 21 \\ 22 \\ 23 \\ 24 \end{array}
                                          if (vx[b] < delta) {</pre>
                                                 delta = vx[b];
                                                 y1 = b;
                                   }
25
26
27
28
29
30
                            for (int b = 0; b <= n; b++) {
   if (vb[b]) {
      la[r[b]] -= delta;
}</pre>
                                          lb[b] += delta;
                                   } else
                                          vx[b] -= delta;
31
32
33
34
35
                            }
                     y = y1;
} while (r[y] != -1);
while (y) {
   r[y] = r[pp[y]];
   r = r [r]
36
37
38
39
                            y = pp[y];
              for (int i = 1; i <= n; i++) {
    std::fill(all(vb), 0);</pre>
40
41
                     bfs(i);
42
43
              LL ans = 0;
for (int i = 1; i <= n; i++) {
    if (r[i] == -1) continue;</pre>
44
45
46
                     l[r[i]] = i;
ans += e[r[i]][i];
47
48
49
50
              return {ans, 1, r};
       };
51
52
       auto [ans, mchl, mchr] = KM(n, e);
```

# 10.24 matching - matching on general graph

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# 11 geometry

#### 11.1 two demention

点与向量

```
struct Point {
  \frac{1}{2}
                      LL x = 0, y = 0;
Point() = default;
                     Point() = default,
Point(long long x, long long y) : x(x), y(y) {}
operator bool() { return *this != Point{}; }
friend bool operator==(Point p, Point q) { return p.x == q.x and p.y == q.y; }
friend bool operator!=(Point p, Point q) { return !(p == q); }
friend Point operator+(Point p, Point q) { return {p.x + q.x, p.y + q.y}; }
friend Point operator-(Point p, Point q) { return {p.x - q.x, p.y - q.y}; }
  4
  5
  6
7
  8
  9
                     friend LL dot(Point p, Point q) { return p.x * q.x + p.y * q.y; }
friend LL det(Point p, Point q) { return p.x * q.y - q.x * p.y; }
friend bool operator<(Point p, Point q) {
    return std::pair{p.quad(), det(q, p)} < std::pair{q.quad(), Oll};
    return (p.x == q.x ? p.y < q.y : p.x < q.x);
}</pre>
10
11
12
13
14
15
16
                      int quad() const {
                               if (x > 0 & x y >= 0) return 1;
if (x <= 0 \text{ and } y > 0) return 2;
if (x <= 0 \text{ and } y <= 0) return 3;
if (x >= 0 \text{ and } y <= 0) return 4;
17
18
19
20
21
22
23
24
25
26
27
                                return 0;
                      friend LL dist(Point p, Point q) {
   return (p.x - q.x) * (p.x - q.x) + (p.y - q.y) * (p.y - q.y);
           std::istream& operator>>(std::istream& is, Point& p) { return is >> p.x >> p.y; }
          std::ostream& operator<<(std::ostream& os, Point p) {
    return os << '(' << p.x << ',' << p.y << ')';
28
30
```

线段

```
struct line {
 3
          point a, b;
 4
5
          line(point _a = {}, point _b = {}) { a = _a, b = _b; }
           /* 交点类型为 double */
 \begin{matrix} 6\\7\\8\\9\end{matrix}
          friend point iPoint(line p, line q) {
               point v1 = p.b - p.a;
               point v2 = q.b - q.a;
               point u = q.a - p.a;
return q.a + (q.b - q.a) * ((u ^ v1) * 1. / (v1 ^ v2));
10
11
13
           /* 极角排序 */
14
          bool operator<(const line& p) const {
   double t1 = std::atan2((b - a).y, (b - a).x);</pre>
15
16
                double t2 = std::atan2((p.b - p.a).y, (p.b - p.a).x);
17
18
19
               if (fabs(t1 - t2) > eps) {
    return t1 < t2;</pre>
20
21
               return ((p.a - a) ^ (p.b - a)) > eps;
22
          }
     };
```

#### 11.2 convex

2D

```
1    /* andrew */
2    auto andrew = [&](std::vector<point>& v) -> std::vector<point> {
3        std::sort(all(v));
```

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```
std::vector<point> stk;
5
        for (int i = 0; i < n; i++) {</pre>
 6
            point x = v[i];
             while (stk.size() > 1 \text{ and } ((stk.end()[-1] - stk.end()[-2]) ^ (x - stk.end()[-2])) <= 0) 
8
                 stk.pop_back();
 9
10
             stk.push_back(x);
11
12
        int tmp = stk.size();
        for (int i = n - 2; i \ge 0; i - -) {
13
             point x = v[i];
14
15
             while (stk.size() > tmp and ((stk.end()[-1] - stk.end()[-2]) ^ (x - stk.end()[-2])) <= 0) {
16
                 stk.pop_back();
             stk.push_back(x);
18
19
20
        return stk;
21
    };
```

## 求凸包 and 判断点与凸包关系

```
#include<bits/stdc++.h>
     #ifdef LOCAL
 3
     #include "debug.h"
     #else
     #define debug(...) 0
 6
     #endif
     #define all(v) begin(v), end(v)
 8
     using namespace std;
     using pii = pair<int, int>;
10
     template<typename T, typename P, T inf = numeric_limits<T>::max()>
struct Convex_Hull {
    ____
11
12
          using ptt = pair<T, T>;
13
         ., using 1128 = __int128; vector<ptt> a, b; T low har-
14
15
16
          T lox, hix;
         P crs(const ptt& a, const ptt& b) {
17
18
              return (P)a.first * b.second - (P)a.second * b.first;
19
20
         ptt mns(const ptt% a, const ptt% b) {
              return ptt{a.first - b.first, a.second - b.second};
21
22
23
          Convex_Hull(vector<ptt> c) {
24
              assert(c.size() > 0);
25
              sort(begin(c), end(c));
26
              vector<int> st = {0};
27
              int n = c.size(), tp = 0;
for (int i = 1; i < n; ++i) {</pre>
28
29
                   while (tp > 0 &&
30
                        crs(mns(c[st[tp]], c[st[tp-1]]), mns(c[i], c[st[tp]])) \leftarrow 0) 
31
                             --tp, st.pop_back();
32
33
                   st.emplace_back(i), ++tp;
34
              }
35
              int tmp = tp;
              for (int i = n - 1; ~i; --i) {
   while (tp > tmp &&
36
37
38
                        crs(mns(c[st[tp]], c[st[tp - 1]]), mns(c[i], c[st[tp]])) <= 0) {</pre>
39
                             --tp, st.pop_back();
40
41
                   st.emplace_back(i), ++tp;
42
43
              for (int i = 0; i <= tmp; ++i) {</pre>
44
                   a.emplace_back(c[st[i]]);
45
46
              for (int i = tmp; i <= tp; ++i) {
47
                   b.emplace_back(c[st[i]]);
              }
48
49
50
          // n >= 3
51
         pair<int, vector<ptt>> insd(T x, T y) { // 0: outside, 1: invertex, 2, {u,v}: inedge, 3: inside
              ptt o = {x, y};
if (x < a[0].first || x > b[0].first) return {0, {}};
52
53
54
              int li = lower_bound(begin(a), end(a), ptt{x, -inf}) - begin(a);
              if (o == a[li]) return {1, {}};
int hi = lower_bound(begin(b), end(b), ptt{x, inf}, greater{}) - begin(b);
if (o == b[hi]) return {1, {}};
55
56
57
              if (li == 0) {
58
                   if (hi + 1 == b.size()) return {0, {}};
assert(b.end()[-1].first == b.end()[-2].first);
59
60
61
                   if (y < b.end()[-1].second || y > b.end()[-2].second) return {0, {}};
                   return {2, {b.end()[-1], b.end()[-2]}};
62
```

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```
63
                  if (hi == 0) {
 64
                        if (li + 1 == a.size()) return {0, {}};
assert(a.end()[-2].first == a.end()[-1].first);
 65
                       if (y < a.end()[-2].second || y > a.end()[-1].second) return {0, {}};
return {2, {a.end()[-2], a.end()[-1]}};
 67
 68
 69
70
71
72
73
74
75
76
77
78
                  P v1 = crs(mns(o, a[li - 1]), mns(a[li], a[li - 1]));
                  if (v1 == 0) return {2, {a[ii - 1], a[ii]}};
P v2 = crs(mns(o, b[hi - 1]), mns(b[hi], b[hi - 1]));
if (v2 == 0) return {2, {b[hi - 1], b[hi]}};
                  debug(v1, v2);
                  return {v1 > 0 && v2 > 0 || v1 < 0 && v2 < 0 ? 3 : 0, {}};
       };
       namespace Acc {
            auto work = []() {
    string ans[] = {"OUT", "ON", "ON", "IN"};
 80
81
82
83
                  int n, q;
cin >> n;
 84
85
                  vector<pair<int, int>> a(n);
for (auto& [x, y] : a) cin >> x >> y;
 86
                  Convex_Hull<int, long long> ch(a);
 87
                  debug(ch.a, ch.b);
 88
                  cin >> q;
 89
                  for (int x, y; q--; ) {
 90
                        cin >> x >>
 91
                        cout << ans[ch.insd(x, y).first] << '\n';</pre>
 92
 93
            };
 94
 95
       int main() {
 96
            std::ios::sync_with_stdio(0);
 97
            std::cin.tie(0);
int T = 1;
 98
 99
             // std::cin >> T;
100
             while (T--) Acc::work();
101
      }
```

# 11.3 half plane union

```
/* half plane union */
        auto half_plane = [&](std::vector<line>& ln) -> std::vector<point> {
 \begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
               std::sort(all(ln));
               ln.erase(
                      unique(
 6
7
                             all(ln),
                              [](line& p, line& q) {
                                    double t1 = std::atan2((p.b - p.a).y, (p.b - p.a).x);
double t2 = std::atan2((q.b - q.a).y, (q.b - q.a).x);
return fabs((t1 - t2)) < eps;</pre>
 8 9
10
                             })
11
                      ln.end());
12
               auto check = [&](line p, line q, line r) -> bool {
   point a = iPoint(p, q);
   return ((r.b - r.a) ^ (a - r.a)) < -eps;</pre>
13
14
15
16
17
               line q[ln.size() + 2];
               int hh = 1, tt = 0;
q[++tt] = ln[0];
18
19
20
               q[++tt] = ln[1];
              qi'fttj - In[i];
for (int i = 2; i < (int) ln.size(); i++) {
    while (hh < tt and check(q[tt - 1], q[tt], ln[i])) tt--;
    while (hh < tt and check(q[hh + 1], q[hh], ln[i])) hh++;
    q[++tt] = ln[i];</pre>
21
22
23
24
25
26
27
               while (hh < tt and check(q[tt - 1], q[tt], q[hh])) tt--;
while (hh < tt and check(q[hh + 1], q[hh], q[tt])) hh++;</pre>
28
29
30
               q[tt + 1] = q[hh];
               std::vector<point> ans;
for (int i = hh; i <= tt; i++) {</pre>
31
                      ans.push_back(iPoint(q[i], q[i + 1]));
32
33
               return ans;
34
       };
```

11.5 Simpson 97

```
/* rotate @ wrb */
    #include<cstdio>
    #include<algorithm>
    #define db double
 5
    namespace Acc{
         const int N = 5e4+10;
 6
         struct node{
         int x,y;
}a[N],stk[N];
 8
 9
10
         11
         int dis(node a,node b){return ((a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y));}
12
         int n,tp,ans;
13
         void work(){
14
             scanf("%d",&n);
             for(int i=1;i<=n;i++)scanf("%d%d",&a[i].x,&a[i].y);</pre>
15
16
             std::sort(a+1,a+n+1,[=](node a,node b)->bool{return a.x<b.x || (a.x==b.x && a.y<b.y);});
             stk[1]=a[1],tp=1;
18
             for(int i=2;i<=n;i++){</pre>
                 while(tp>1 && cmp(stk[tp-1],stk[tp],a[i])<=0)tp--;</pre>
19
20
                 stk[++tp]=a[i];
\overline{21}
22
             int tmp=tp;
for(int i=n-1;i>=1;i--){
23
                 while(tp>tmp && cmp(stk[tp-1],stk[tp],a[i])<=0)tp--;</pre>
24
\frac{1}{25}
                 stk[++tp]=a[i];
26
27
             for(int i=1,j=3;i<tp;i++){</pre>
28
                 \label{lem:while} \begin{tabular}{ll} while (cmp(stk[i],stk[i+1],stk[j]) \le cmp(stk[i],stk[i+1],stk[j+1])) = j\%(tp-1)+1; \\ \end{tabular}
29
                 ans=std::max(ans,std::max(dis(stk[i],stk[j]),dis(stk[i+1],stk[j])));
30
31
             printf("%d",ans);
32
         }
33
    }
34
    int main(){
35
         return Acc::work(),0;
    }
36
```

# 11.5 Simpson

```
#include <bits/stdc++.h>
      using namespace std;
 3
      double a;
 4
      double f(double x) {
 5
            return pow(x, a / x - x);
 6
      double simpson(double 1, double r) {
   double mid = (1 + r) / 2;
   return (r - 1) * (f(1) + 4 * f(mid) + f(r)) / 6;
 9
10
      }
      double asr(double l, double r, double eps, double ans, int d) {
   double mid = (1 + r) / 2;
   double Fl = simpson(l, mid), Fr = simpson(mid, r);
   if (abs(Fl + Fr - ans) <= 15 * eps && d < 0) {
      return Fl + Fr + (Fl + Fr - ans) / 15;
   }</pre>
11
12
13
14
15
16
17
            return asr(1, mid, eps / 2, Fl, d - 1) + asr(mid, r, eps / 2, Fr, d - 1);
18
19
      double calc(double 1, double r, double eps) {
20
            return asr(l, r, eps, simpson(l, r), 12);
21
22
      int main() {
23
            cin >> a:
            if (a < 0) {
\overline{24}
                  cout << "orz\n";
25
26
            } else {
27
                  cout << fixed << setprecision(5) << calc(1e-8, 15, 1e-8) << '\n';
28
29
      }
```

# 12 offline algorithm

# 12.1 discretization

```
std::sort(all(a));
a.erase(unique(all(a)), a.end());
auto get_id = [&](const int& x) -> int { return lower_bound(all(a), x) - a.begin() + 1; };
```

# 12.2 Mo algorithm

#### 普通莫队

```
int block = n / sqrt(2 * m / 3);
     std::sort(all(q), [&](node a, node b) {
    return a.l / block == b.l / block ? (a.r == b.r ? 0 : ((a.l / block) & 1) ^ (a.r < b.r))</pre>
 \frac{2}{3}
                                                        : a.1 < b.1;
 5
     auto move = [&](int x, int op) -> void {
          if (op == 1) {
              /* operations */
 9
          } else {
10
              /* operations */
11
12
    for (int k = 1, 1 = 1, r = 0; k <= m; k++) {
  node Q = q[k];</pre>
13
14
          while (1 > Q.1) {
15
16
               move(--1, 1);
17
18
          while (r < Q.r) {</pre>
19
              move(++r, 1);
20
21
          while (1 < Q.1) {</pre>
22
23
               move(l++, -1);
24
25
26
27
          while (r > Q.r) {
               move(r--, -1);
```

# 12.3 回滚莫队

```
/* rollback Mo */
      #include<bits/stdc++.h>
 \frac{1}{3}
      namespace Acc {
    const int N = 200009;
 5
            int a[N], b[N], id[N], f[N], g[N], p[N], z[N];
 6
7
            std::pair<int, int> st[N];
            struct T {
                  int 1, r, o;
 9
            }q[N];
10
            auto work = []() {
11
                  int n, m;
std::cin >> n;
12
13
                  for (int i = 1; i <= n; ++i) {</pre>
14
                         std::cin >> a[i], b[i] = a[i];
15
                  std::sort(b + 1, b + n + 1);
int ct = std::unique(b + 1, b + n + 1) - b - 1;
for (int i = 1; i <= n; ++i) {</pre>
16
18
19
                         a[i] = std::lower_bound(b + 1, b + ct + 1, a[i]) - b;
\frac{20}{21}
                  std::cin >> m;
                  for (int i = 1; i <= m; ++i) {
    auto&[1, r, o] = q[i];
    std::cin >> 1 >> r, o = i;
23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28
                  int B = ceil(n / sqrt(m));
for (int i = 1; i <= n; ++i) {
   id[i] = (i - 1) / B + 1;</pre>
29
30
                   std::sort(q + 1, q + m + 1, [](T a, T b) {
```

12.4 CDQ 99

```
31
                         return id[a.1] == id[b.1] ? a.r < b.r : a.l < b.1;</pre>
32
33
                   int ans = 0, L = 1, R = 0;
                  for (int i = 1; i <= m; ++i) {
   auto[1, r, o] = q[i];
   if (id[1] != id[q[i - 1].1]) {</pre>
34
35
36
37
                               ans = 0;
                              mems td:min(n, id[1] * B), L = R + 1;
memset(f + 1, 0, ct << 2);
memset(g + 1, 0, ct << 2);</pre>
38
39
40
                         }
41
42
                         if (id[1] == id[r]) {
                               for (int j = 1; j <= r; ++j) {
    if (p[a[j]] == 0) p[a[j]] = j;
    else ans = std::max(ans, j - p[a[j]]);</pre>
43
44
45
46
                              for (int j = 1; j <= r; ++j) p[a[j]] = 0; z[o] = ans, ans = 0;
47
48
49
                         } else {
                               while (R < r) {
50
                                     ++R, g[a[R]] = R;
if (f[a[R]] == 0) f[a[R]] = R;
51
52
                                     else ans = std::max(ans, R - f[a[R]]);
53
54
                               int las = ans, t = L;
while (1 < L) {</pre>
55
56
57
                                     --L;
                                     int x = f[a[L]], y = g[a[L]];
st[L] = std::make_pair(x, y);
58
59
                                     f[a[L]] = L;
60
61
                                     if (g[a[L]] == 0) g[a[L]] = L;
                                     else ans = std::max(ans, g[a[L]] - L);
62
                               }
63
                               z[o] = ans;
64
                              for (int j = 1; j < t; ++j) {
    auto[x, y] = st[j];
    f[a[j]] = x, g[a[j]] = y;
65
66
67
68
69
                               ans = las, L = t;
70
                         }
71
                  for (int i = 1; i <= m; ++i) {
    std::cout << z[i] << '\n';</pre>
73
74
                  }
75
            };
76
77
      }
      int main() {
78
             std::ios::sync_with_stdio(0);
79
             std::cin.tie(0), Acc::work();
80
      }
```

## 12.4 CDQ

n 个三维数对  $(a_i, b_i, c_i)$ , 设 f(i) 表示  $a_j \leq a_i, b_j \leq b_i, c_j \leq c_i (i \neq j)$  的个数. 输出 f(i)  $(0 \leq i \leq n-1)$  的值.

```
// 洛谷 P3810 【模板】三维偏序(陌上花开)
 3
     struct data {
 4
          int a, b, c, cnt, ans;
 5
          data(int _a = 0, int _b = 0, int _c = 0, int _cnt = 0, int _ans = 0) {
   a = _a, b = _b, c = _c, cnt = _cnt, ans = _ans;
 6
 9
10
          bool operator!=(data x) {
               if (a != x.a) return true;
if (b != x.b) return true;
11
12
13
               if (c != x.c) return true;
               return false;
14
15
          }
     };
16
17
18
     int main() {
19
          std::ios::sync_with_stdio(false);
20
21
          std::cin.tie(0);
22
          int n, k;
23
          std::cin >> n >> k;
          static data v1[N], v2[N];
for (int i = 1; i <= n; i++) {</pre>
24
25
```

```
26
                std::cin >> v1[i].a >> v1[i].b >> v1[i].c;
27
28
           std::sort(v1 + 1, v1 + n + 1, [\&](data x, data y) {
                if (x.a != y.a) return x.a < y.a;
if (x.b != y.b) return x.b < y.b;
return x.c < y.c;</pre>
29
30
31
32
           }):
33
34
35
           int t = 0, top = 0;
           for (int i = 1; i <= n; i++) {
                t++:
36
37
                 if (v1[i] != v1[i + 1]) {
                      v2[++top] = v1[i];
38
39
                      v2[top].cnt = t;
                      t = 0;
                }
40
41
42
           vi tr(N);
           auto add = [&](int pos, int val) -> void {
   while (pos <= k) {</pre>
43
44
45
                      tr[pos] += val;
46
                      pos += lowbit(pos);
47
                }
48
49
50
           auto query = [&](int pos) -> int {
                 int ans = 0;
51
                while (pos > 0) {
52
                      ans += tr[pos];
                      pos -= lowbit(pos);
53
54
55
                return ans;
56
57
           std::function<void(int, int)> CDQ = [&](int 1, int r) -> void {
                if (1 == r) return;
int mid = (1 + r) >> 1;
58
59
                CDQ(1, mid), CDQ(mid + 1, r);

std::sort(v2 + 1, v2 + mid + 1, [&](data x, data y) {

    if (x.b != y.b) return x.b < y.b;

    return x.c < y.c;
60
61
62
63
64
                });
65
                 std::sort(v2 + mid + 1, v2 + r + 1, [&](data x, data y) {
                     if (x.b != y.b) return x.b < y.b;
return x.c < y.c;</pre>
66
67
                int i = 1, j = mid + 1;
while (j <= r) {</pre>
69
70
71
72
73
74
75
76
77
78
79
                      while (i <= mid && v2[i].b <= v2[j].b) {</pre>
                           add(v2[i].c, v2[i].cnt);
                           i++;
                      v2[j].ans += query(v2[j].c);
                      j++;
                for (int ii = 1; ii < i; ii++) {
                      add(v2[ii].c, -v2[ii].cnt);
80
                }
81
                return;
82
83
           CDQ(1, top);
           vi ans(n + 1);
for (int i = 1; i <= top; i++) {
   ans[v2[i].ans + v2[i].cnt] += v2[i].cnt;</pre>
84
85
86
87
88
           for (int i = 1; i <= n; i++) {
    std::cout << ans[i] << endl;</pre>
89
90
91
           return 0;
92
```

#### 12.5 segment tree devide and conquer

```
/* seg div @ wrb */
#include<bits/stdc++.h>
using namespace std;
namespace Acc {
    const int N = 1e5;
    pair<int, int> q[N * 2];
    vector<int> v[N * 4];
    int n, l, r, p;
    int fa[N * 2], sz[N * 2];
    pair<int, int> st[N * 2];
    int tp;
    void ins(int o, int L, int R) {
```

```
13
                      if (r < L || 1 > R) return ;
14
                      if (1 <= L && R <= r) {</pre>
15
                            v[o].emplace_back(p);
16
                            return ;
18
                     int md = L + R >> 1;
                     ins(o << 1, L, md);
ins(o << 1 | 1, md + 1, R);
19
20
\overline{21}
              auto gf = [](int x) {
    while (x != fa[x]) x = fa[x];
\overline{22}
23
24
25
                     return x;
26
               auto mg = [](int x, int y) {
                     x = gf(x), y = gf(y);
if (x != y) {
27
28
                            if (sz[x] < sz[y]) swap(x, y);
fa[y] = x, sz[x] += sz[y];
st[++tp] = {x, y};</pre>
29
30
31
32
                     }
33
34
35
              f;
void dfs(int o, int L, int R) {
   int lastp = tp;
   for (int i : v[o]) {
      auto[x, y] = q[i];
      mg(x, y + n), mg(x + n, y);
      if (gf(x) == gf(x + n)) {
        for (int i = L; i <= R; ++i) {
            cout << "Mo\n":</pre>
36
37
38
39
40
                                          cout << "No\n";
41
42
43
                                   goto _;
                            }
44
45
                     if (L == R) {
    cout << "Yes\n";</pre>
46
47
                     } else {
48
49
                            int md = L + R >> 1;
                            dfs(o << 1, L, md);
dfs(o << 1 | 1, md + 1, R);
50
51
52
                     }
53
                     for (; tp > lastp; --tp) {
   auto[x, y] = st[tp];
   fa[y] = y, sz[x] -= sz[y];
54
55
56
57
                     }
58
59
               auto work = []() {
                     int m, k;
cin >> n >> m >> k;
for (int i = 1; i <= m; ++i) {</pre>
60
61
62
\overline{63}
                            int x, y;
cin >> x >> y >> 1 >> r;
if (++1 <= r) {</pre>
64
65
66
                                   q[p = i] = \{x, y\}, ins(1, 1, k);
67
68
                     iota(fa + 1, fa + n * 2 + 1, 1);
fill(sz + 1, sz + n * 2 + 1, 1);
dfs(1, 1, k);
69
70
71
72
73
74
              };
       }
       int main() {
75
              ios::sync_with_stdio(0);
76
               cin.tie(0), Acc::work();
       }
```

# 13 Print All Cases

# 13.1 print all trees with n nodes

构造所有 n 个节点的树.

## 13.1.1 有根树

```
表示其数量的数列在 oeis 上编号为 A000081. n=1,2,3\cdots,20 的项分别为: 1,1,2,4,9, 20,48,115,286,719, 1842,4766,12486,32973,87811, 235381,634847,1721159,4688676,12826228.
```

构造所有  $n \le 20$  的有根树的 (平均) 运行时间为 15.7054s.

```
/* integer partition */
      int n = 5;
 \begin{array}{c} 3 \\ 4 \\ 5 \\ 6 \\ 7 \\ 8 \end{array}
      std::vector<vvi> part(n + 1);
      auto integerPartition = [&](int n) {
   // part[1] = {{1}};
   for (int i = 1; i <= n; i++) {</pre>
                 part[i].push_back({i});
for (int j = 1; j < i; j++) {
    for (const auto& v : part[i - j]) {</pre>
 9
10
                             vi tmp = v;
                             tmp.push_back(j);
std::sort(all(tmp));
11
12
13
14
15
                             part[i].push_back(tmp);
16
                 std::sort(all(part[i]));
17
                 part[i].erase(unique(all(part[i])), part[i].end());
18
19
20
     integerPartition(n);
\overline{21}
      /* find all trees */
22
      std::vector<std::string>> trees(n + 1);
\begin{array}{c} 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 30 \\ 31 \\ 32 \\ 33 \\ 34 \\ 35 \\ 36 \\ 37 \end{array}
      auto allTrees = [&](int n) {
           std::string s;
for (int i = 1; i < n; i++) s += '(';
for (int i = 1; i < n; i++) s += ')';
            trees[n].push_back(s);
           for (const auto& v : part[n - 1]) {
    std::vector<std::string> now;
                  auto dfs = [&](auto&& self, int i) {
                       if (i == v.size()) {
                             std::string s = "";
                             auto tmp = now;
                             std::sort(all(tmp));
                             for (const auto& ss : tmp) s += '(' + ss + ')';
                             trees[n].push_back(s);
                             return;
38
39
                       for (const auto& s : trees[v[i]]) {

40

41

42

43

                             now.push_back(s);
self(self, i + 1);
                             now.pop_back();
                       }
44
45
                  dfs(dfs, 0);
46
            std::sort(all(trees[n]));
48
            trees[n].erase(unique(all(trees[n])), trees[n].end());
49
     };
for (int i = 1; i <= n; i++) {
50
51
52
53
54
           allTrees(i);
           debug(i, trees[i].size());
            std::cout << '\n';
55
     for (const auto& s : trees[n]) {
56
            vvi e(n + 1);
            vi fa(n + 1);
            int cnt = 1, now = 1;
```

```
for (const auto& c : s) {
    if (c == '(') {
        cnt += 1;
        e [now].push_back(cnt);
        e[cnt].push_back(now);
    fa[cnt] = now;
        now = cnt;
    } else {
        now = fa[now];
    }
}
debug(e);
/* do the things you need */
}
```

104 14 MAGIC

# 14 Magic

# 14.1 magic heap

对顶堆维护中位数.

```
/* magic heap */

    \begin{array}{r}
      2 \\
      3 \\
      4 \\
      5 \\
      6 \\
      7 \\
      8 \\
      9
    \end{array}

       struct MagicHeap {
             LL sum1 = 0, sumr = 0;
             std::priority_queue<int> ql;
             std::priority_queue<int, std::vector<int>, std::greater<int>> qr;
void le2ri() {
                    auto x = ql.top();
                   suml -= x, ql.pop();
sumr += x, qr.push(x);
10
11
             void ri2le() {
                   auto x = qr.top();
12
13
                   sumr -= x, qr.pop();
                   suml += x, ql.push(x);
14
15
             void pushL(int x) { suml += x, ql.push(x); }
void pushR(int x) { sumr += x, qr.push(x); }
void push(int x) {
16
17
18
19
                    if (ql.empty()) {
20
21
22
23
                         pushL(x);
                   } else if (qr.empty()) {
   (x <= q1.top() ? le2ri(), pushL(x) : pushR(x));</pre>
                    } else {
                         int le = ql.top(), ri = qr.top();
if (le <= x and x <= ri) {
    (ql.size() == qr.size() ? pushL(x) : pushR(x));
} else if (x < le) {</pre>
24
25
26
27
28
                                if (ql.size() != qr.size()) le2ri();
29
30
                                pushL(x);
                          } else {
31
                                if (ql.size() <= qr.size()) ri2le();</pre>
32
                                pushR(x);
33
34
                   }
35
36
             int size() { return ql.size() + qr.size(); }
bool empty() { return ql.empty() and qr.empty(); }
37
38
             LL val() { return suml + sumr; }
39
             LL mid() { return ql.top(); }

LL dist() { return sumr - suml + ql.top() * (ql.size() - qr.size()); }
40
41
      };
```

## 14.2 operator queue

双栈维护队列半群.

```
template <typename T, typename Op>
     struct OpQueue {
 \begin{array}{c} 2\\ 3\\ 4\\ 5 \end{array}
          static_assert(std::is_convertible_v<std::invoke_result_t<0p, T, T>, T>);
           const T e;
          const 1 e,
const 0p op;
std::vector<T> 1, r, a;
OpQueue(T e, 0p op) : e(e), op(op), l{e}, r{e} {}
T val() const { return op(l.back(), r.back()); }
 6
7
 8 9
           void push(T x) {
10
               r.push_back(op(r.back(), x));
11
                a.push_back(x);
12
13
           void pop() {
14
                if (1.size() == 1) {
15
                     for (; !a.empty(); a.pop_back()) {
16
                          1.push_back(op(a.back(), 1.back()));
17
18
                     r.resize(1);
19
20
21
                assert(l.size() > 1);
                1.pop_back();
22
23
           int size() const { return 1.size() + r.size() - 2; }
\overline{24}
           bool empty() const { return 1.size() + r.size() == 2; }
25
```

14.3 Fast GCD 105

```
26 | /* When using this, remember to replace "T" with correct type and "e" with identity in the half group. */
28 | auto op = [](T a, T b) -> T {
29 | /* You operations */
30 | };
31 | OpQueue<T, decltype(op)> a(e, op);
```

## 14.3 Fast GCD

O(V) 预处理, O(1) 查询 GCD.

```
/* fast GCD @ luogu shit */
       const int N = 5005, M = 1e6 + 5, S = 1000, P = 998244353;
 3
       int _a[M], _b[M], _c[M];
       int v[M], p[M], r;
int f[S + 1][S + 1];
 6
       int a[N], b[N];
 9
       void Init() {
              a[1] = _b[1] = _c[1] = 1;
for (int i = 2; i <= M - 5; ++i) {
    if (!v[i]) {
        p[++r] = i;
        c[i] = b[i] = 1        c[i] =</pre>
10
11
12
13
                           a[i] = b[i] = 1, c[i] = i;
14
                    }
15

int tp;
for (int j = 1; j <= r && (tp = i * p[j]) <= M - 5; ++j) {
    v[tp] = 1;
        _a[tp] = _a[i] * p[j];
        _b[tp] = _b[i];
        _c[tp] = _c[i];
    if (_a[tp] > _b[tp]) {
        swap(_a[tp], _b[tp]);
        if (_b[tp] > _c[tp]) {
            swap(_b[tp], _c[tp]);
        }
}

16
17
18
19
20
21
22
23
24
25
26
27
28
                           if (!(i % p[j])) {
\overline{29}
                                 break;
30
31
                    }
             }
32
33
34
             for (int i = 1; i <= S; ++i) {</pre>
                    f[0][i] = f[i][0] = i;
for (int j = 1; j <= S; ++j) {
f[i][j] = f[j % i][i];
35
36
37
38
39
40
             return;
       }
41
42
       int gcd(int x, int y) {
43
              int A = 1, tp = f[_a[x]][y % _a[x]];
44
45
              A *= tp;
             y /= tp;
46
              tp = f[_b[x]][y \% _b[x]];
47
              A *= tp;
48
49
              tp = (v[_c[x]] ? f[_c[x]][y % _c[x]] : (y % _c[x] ? 1 : _c[x]));
A *= tp;
50
51
              y /= tp;
52
53
              return A;
54
       }
```

# **14.4** $q \equiv \frac{a}{b} \mod mod$

 $14 \quad MAGIC$ 

10 [}