

2 Keys Keyboard

Initially on a notepad only one character 'A' is present. You can perform two operations on this notepad for each step:

1. **Copy All** : You can copy all the characters present on the notepad (partial copy is not allowed).
2. **Paste** : You can paste the characters which are copied **last time**.

Given a number n . You have to get **exactly** n 'A' on the notepad by performing the minimum number of steps permitted. Output the minimum number of steps to get n 'A'.

Example 1:

Input: 3

Output: 3

Explanation:

Initially, we have one character 'A'.

In step 1, we use **Copy All** operation.

In step 2, we use **Paste** operation to get 'AA'.

In step 3, we use **Paste** operation to get 'AAA'.

Note:

1. The n will be in the range $[1, 1000]$.

Solution 1

```
public int minSteps(int n) {  
    int[] dp = new int[n+1];  
  
    for (int i = 2; i <= n; i++) {  
        dp[i] = i;  
        for (int j = i-1; j > 1; j--) {  
            if (i % j == 0) {  
                dp[i] = dp[j] + (i/j);  
                break;  
            }  
        }  
    }  
    return dp[n];  
}
```

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Solution 2

We look for a divisor d so that we can make d copies of (n / d) to get n

The process of making d copies takes d steps (1 step of **Copy All** and $d - 1$ steps of **Paste**)

We keep reducing the problem to a smaller one in a loop.

The **best** cases occur when n is decreasing fast, and method is almost $O(\log(n))$

For example, when $n = 1024$ then n will be divided by 2 for only 10 iterations, which is much faster than $O(n)$ DP method.

The **worst** cases occur when n is some multiple of large prime, e.g. $n = 997$ but such cases are rare.

```
public int minSteps(int n) {  
    int s = 0;  
    for (int d = 2; d <= n; d++) {  
        while (n % d == 0) {  
            s += d;  
            n /= d;  
        }  
    }  
    return s;  
}
```

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Solution 3

```
/**
 * It take 2 op to double, 3 ops to triple, ...
 * if  $n \% 2 == 0$ , then  $f(n) = f(n/2) + 2$ 
 * if  $n \% 3 == 0$ , then  $f(n) = f(n/3) + 3$ 
 *  $2 * 2 = 2 + 2$ ,  $2 * 3 > 2 + 3$ ,  $4 * 4 > 4 + 4$ , so it is always better to divide whenever possible.
 * now it became a problem for finding all possible factors;
 */
class Solution {
public:
    int minSteps(int n) {
        if (n == 1) return 0;
        for (int i = 2; i < n; i++)
            if (n % i == 0) return i + minSteps(n / i);
        return n;
    }
};
```

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