

应用密码学作业 #4

XXX:202XX80XXXXXXXX

2023年4月14日

# 1 DES 第一轮的计算

第 1 轮的输出为 FFFFFFFF 61545920。 附录 A,B 为原始代码。若要得到结果,需运行  $\$ruby\ DES.rb$ 。

## 2 AES 初始化轮和第一轮的计算

附录 C, D 为原始代码。若要得到结果,需运行  $$ruby\ AES.rb$ 。

- (1)State 最初的内容
- 00 04 08 0C
- 01 05 09 0D
- 02 06 0A 0E
- 03 07 0B 0F
- (2) 初始化轮密钥加密后的值
- 01 05 09 0D
- 00 04 08 0C
- 03 07 0B 0F
- $02\quad 06\quad 0A\quad 0E$
- (3) 字节替换后的结果
- 7C 6B 01 D7
- 63 F2 30 FE
- 7B C5 2B 76
- 77 6F 67 AB
- (4) 行移位后的结果
  - 7C 6B 01 D7
  - F2 30 FE 63
  - 2B 76 7B C5
  - AB 77 6F 67
- (5) 列混淆之后的结果
- 75 87 0F B2
- 55 E6 04 22
- 3E 2E B8 8C
- $10 \quad 15 \quad 58 \quad 0 \mathrm{A}$
- (6) 第一轮使用的轮密钥

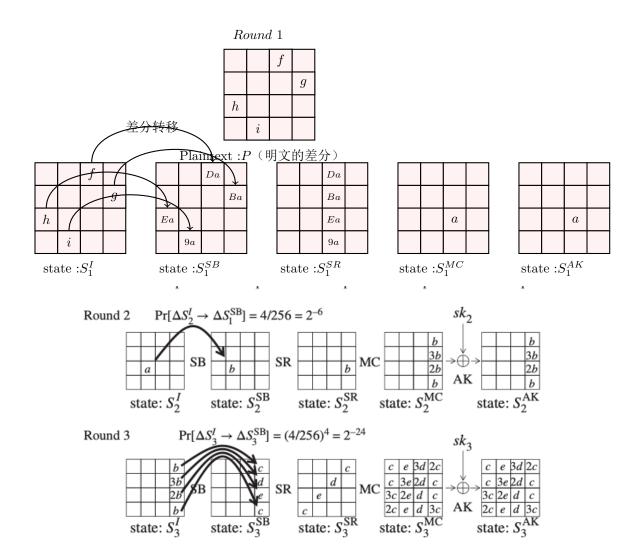
```
7F 7E 7F 7E
7D 7C 7D 7C
7D 7C 7D 7C
7D 7C 7D 7C
```

# 3 DES 中与 Sbox 有关偏差的计算

根据课件, $X_2$  应当是 S 盒的 6 位输入的次高位比特,而  $Y_1,Y_2,Y_3,Y_4$  为输出的 4 个比特。

$Sbox_i$	该随机变量等于 0 的次数	偏差
0	14	-0.281250
1	20	-0.187500
2	42	0.156250
3	40	0.125000
4	12	-0.312500
5	42	0.156250
6	18	-0.218750
7	16	-0.250000
附代码:		

```
require "./init.rb"
puts " 第 i 个 S 盒 & 该随机变量等于 0 的次数 & 偏差 \\\\"
total = 64.0
8.times do |i|
    count = 0.0
    4.times do |j|
        16.times do |k|
            count_ones = (Sbox[i][j][k]*2 + k/8).to_s(2).scan("1").length
        if count_ones % 2 == 0
            count += 1
        end
    end
    end
    puts "%d & %d & %f \\\\" %[i, count, count/total - 0.5]
end
```



### 4 AES 的差分特征

如上图所示,前三轮的一个可能的查分特征寻找方法如上图所示,明文一 开始的差分如第一个矩阵所示,我们不指出其具体的差分值是多少,我们仅以 字母表示该处的差分。为了构造一个查分,我们首先确定第一轮末的一个差分, 假定其是 a,设 a 能差分转移到 b,可以向前和向后推到这个差分转移向第一轮 传播和向后面两轮传播的情况。如图所示,具体不再阐述。其中二三轮的查分 推导截图来自《Security of Block Ciphers From Algorithm Design to Hardware Implementation-Wiley 2015

### 5 偏差

记  $X_1 \oplus X_2$  为随机变量  $X_1 \oplus X_2$ ,  $X_2 \oplus X_3$  为随机变量  $X_2 \oplus X_3$ , 则  $X_1 \oplus X_2$ 与  $X_2 \oplus X_3$  互相独立  $\iff$   $P(X_1 \oplus X_2 = i, X_2 \oplus X_3 = j) = P(X_1 \oplus X_2 = i, X_2 \oplus X_3 = j)$  $i)P(X_2 \oplus X_3 = j), i, j \in \{0, 1\} \iff$ 

$$\begin{cases} P(X_1 \oplus X_2 = 0, X_2 \oplus X_3 = 0) = P(X_1 \oplus X_2 = 0)P(X_2 \oplus X_3 = 0) \\ P(X_1 \oplus X_2 = 0, X_2 \oplus X_3 = 1) = P(X_1 \oplus X_2 = 0)P(X_2 \oplus X_3 = 1) \\ P(X_1 \oplus X_2 = 1, X_2 \oplus X_3 = 0) = P(X_1 \oplus X_2 = 1)P(X_2 \oplus X_3 = 0) \\ P(X_1 \oplus X_2 = 1, X_2 \oplus X_3 = 1) = P(X_1 \oplus X_2 = 1)P(X_2 \oplus X_3 = 1) \end{cases}$$

$$\begin{cases} (\frac{1}{2} + \epsilon_1)(\frac{1}{2} + \epsilon_2)(\frac{1}{2} + \epsilon_3) + (\frac{1}{2} - \epsilon_1)(\frac{1}{2} - \epsilon_2)(\frac{1}{2} - \epsilon_3) = (\frac{1}{2} + 2\epsilon_1\epsilon_2)(\frac{1}{2} + 2\epsilon_2\epsilon_3) \\ (\frac{1}{2} + \epsilon_1)(\frac{1}{2} + \epsilon_2)(\frac{1}{2} - \epsilon_3) + (\frac{1}{2} - \epsilon_1)(\frac{1}{2} - \epsilon_2)(\frac{1}{2} + \epsilon_3) = (\frac{1}{2} + 2\epsilon_1\epsilon_2)(\frac{1}{2} - 2\epsilon_2\epsilon_3) \\ (\frac{1}{2} + \epsilon_1)(\frac{1}{2} - \epsilon_2)(\frac{1}{2} - \epsilon_3) + (\frac{1}{2} - \epsilon_1)(\frac{1}{2} + \epsilon_2)(\frac{1}{2} + \epsilon_3) = (\frac{1}{2} - 2\epsilon_1\epsilon_2)(\frac{1}{2} + 2\epsilon_2\epsilon_3) \\ (\frac{1}{2} - \epsilon_1)(\frac{1}{2} + \epsilon_2)(\frac{1}{2} - \epsilon_3) + (\frac{1}{2} + \epsilon_1)(\frac{1}{2} - \epsilon_2)(\frac{1}{2} + \epsilon_3) = (\frac{1}{2} - 2\epsilon_1\epsilon_2)(\frac{1}{2} - 2\epsilon_2\epsilon_3) \end{cases}$$

上述四式化简后均得到  $4\epsilon_1\epsilon_2^2\epsilon_3=\epsilon_1\epsilon_3$ 。

四式化简后均得到 
$$4\epsilon_1\epsilon_2^2\epsilon_3 = \epsilon_1\epsilon_3$$
。  
 $4\epsilon_1\epsilon_2^2\epsilon_3 = \epsilon_1\epsilon_3 \iff (4\epsilon_2^2 - 1)\epsilon_1\epsilon_2 = 0 \iff \epsilon_1 = 0 \text{ or } \epsilon_3 = 0 \text{ or } \epsilon_2 = \pm \frac{1}{2}$ 

## 6 RSA 加密

- (1) 因为公钥 e=17,  $n=p \times q=77$ ,  $\phi(n)=6 \times 10=60$ , 私钥  $d=e^{-1}$  $\mod \phi(n) = 53$ ,故密文  $c = m^d \mod n = 8^{53} \mod 77 = 50$
- (2) 若采用蒙哥马利算法,则由 53 的二进制展开 (110101)2,根据蒙哥马 利算法,从次高位开始至最低位,若该位为1,则进行一次模平方运算和

一次模乘运算,即两次模乘运算;若该位为 0,则进行一次模平方运算(一次模乘)。次高位至最低位共有 3 个 1 和 2 个 0,故共需  $3\times2+2\times1=8$  次模乘运算。

### 7 RSA 解密

公开密钥 (e, n) = (5, 35), 密文 c 为 10, 求明文 m, 意即已知  $m^5 \mod 35 = 10$ , 求 m。而只要求出加密所用的私钥 d 即可。

因为 d, e 满足  $de \mod \phi(n) = 1$ ,  $n = 5 \times 7$ ,  $\phi(n) = 4 \times 6 = 24$ , 根据欧几里得算法,可计算 d = 5, 故明文  $m = c^5 \mod 35 = 10^5 \mod 35 = 5$ 。

#### $\mathbf{A}$ init.rb

```
# in all of the matrixs, the index starts with 1
# the initial permutaion of plaintext
IP = [
    58,50,42,34,26,18,10,2,60,52,44,36,28,20,12,4,
    62,54,46,38,30,22,14,6,64,56,48,40,32,24,16,8,
    57,49,41,33,25,17,9,1,59,51,43,35,27,19,11,3,
    61,53,45,37,29,21,13,5,63,55,47,39,31,23,15,7
]
# the permutation matrix of key extension
PC_1 = [
    57, 49, 41, 33, 25, 17, 9,
    1, 58, 50, 42, 34, 26, 18,
    10, 2, 59, 51, 43, 35, 27,
    19, 11, 3, 60, 52, 44, 36,
    63, 55, 47, 39, 31, 23, 15,
    7, 62, 54, 46, 38, 30, 22,
    14, 6, 61, 53, 45, 37, 29,
    21, 13, 5, 28, 20, 12, 4
]
PC 2 = [
    14, 17, 11, 24, 1, 5,
```

3, 28, 15, 6, 21, 10,

```
23, 19, 12, 4, 26, 8,
    16, 7, 27, 20, 13, 2,
    41, 52, 31, 37, 47, 55,
    30, 40, 51, 45, 33, 48,
    44, 49, 39, 56, 34, 53,
    46, 42, 50, 36, 29, 32
]
# the eight sbox used between AddDey and Pbox permutation
Sbox = [
    [14,4,13,1,2,15,11,8,3,10,6,12,5,9,0,7],
        [0,15,7,4,14,2,13,1,10,6,12,11,9,5,3,8],
        [4,1,14,8,13,6,2,11,15,12,9,7,3,10,5,0],
        [15,12,8,2,4,9,1,7,5,11,3,14,10,0,6,13],
    ],
    [15,1,8,14,6,11,3,4,9,7,2,13,12,0,5,10],
        [3,13,4,7,15,2,8,14,12,0,1,10,6,9,11,5],
        [0,14,7,11,10,4,13,1,5,8,12,6,9,3,2,15],
        [13,8,10,1,3,15,4,2,11,6,7,12,0,5,14,9],
    ],
    [10,0,9,14,6,3,15,5,1,13,12,7,11,4,2,8],
        [13,7,0,9,3,4,6,10,2,8,5,14,12,11,15,1],
        [13,6,4,9,8,15,3,0,11,1,2,12,5,10,14,7],
        [1,10,13,0,6,9,8,7,4,15,14,3,11,5,2,12],
    ],
    [7,13,14,3,0,6,9,10,1,2,8,5,11,12,4,15],
        [13,8,11,5,6,15,0,3,4,7,2,12,1,10,14,9],
        [10,6,9,0,12,11,7,13,15,1,3,14,5,2,8,4],
        [3,15,0,6,10,1,13,8,9,4,5,11,12,7,2,14],
    ],
    [2,12,4,1,7,10,11,6,8,5,3,15,13,0,14,9],
        [14,11,2,12,4,7,13,1,5,0,15,10,3,9,8,6],
        [4,2,1,11,10,13,7,8,15,9,12,5,6,3,0,14],
```

```
[11,8,12,7,1,14,2,13,6,15,0,9,10,4,5,3],
    ],
    [12,1,10,15,9,2,6,8,0,13,3,4,14,7,5,11],
        [10,15,4,2,7,12,9,5,6,1,13,14,0,11,3,8],
        [9,14,15,5,2,8,12,3,7,0,4,10,1,13,11,6],
        [4,3,2,12,9,5,15,10,11,14,1,7,6,0,8,13],
    ],
    [4,11,2,14,15,0,8,13,3,12,9,7,5,10,6,1],
        [13,0,11,7,4,9,1,10,14,3,5,12,2,15,8,6],
        [1,4,11,13,12,3,7,14,10,15,6,8,0,5,9,2],
        [6,11,13,8,1,4,10,7,9,5,0,15,14,2,3,12],
    ],
    [13,2,8,4,6,15,11,1,10,9,3,14,5,0,12,7],
        [1,15,13,8,10,3,7,4,12,5,6,11,0,14,9,2],
        [7,11,4,1,9,12,14,2,0,6,10,13,15,3,5,8],
        [2,1,14,7,4,10,8,13,15,12,9,0,3,5,6,11],
    ]
]
# the permutation that extends the right 32 bits of last round to 48 bits
E = [
    32,1,2,3,4,5,4,5,6,7,8,9,
    8,9,10,11,12,13,12,13,14,15,16,17,
    16,17,18,19,20,21,20,21,22,23,24,25,
    24,25,26,27,28,29,28,29,30,31,32,1
]
# the permutation after the Sbox
P = [
    16,7,20,21,29,12,28,17,
    1,15,23,26,5,18,31,10,
    2,8,24,14,32,27,3,9,
    19,13,30,6,22,11,4,25
]
```

### $\mathbf{B}$ DES.rb

```
require "./init.rb"
class ConvertArray
    #convert array to binary string, delimited by a space
    def self.toBinary(input)
        input.each_slice(4) . map . inject("") {|sum, s| sum << (s . join . to_s + " ")}</pre>
    end
    #convert array to hex string, delimited by a space
    def self.toHex(input)
        x = input.each\_slice(4) . map . inject("") {|sum, s| sum << s . join . to\_s}
        "%08X" % x.to_i(2)
    end
end
class Bignum
    def hex_to_binary_arr
        bin_len = self.to_s(16) . length * 4
        format = "%#{bin_len}d"
        sprintf(format, self.to_s(2)) . split("") . map(&:to_i)
    end
end
class Fixnum
    def hex_to_binary_arr
        bin_len = self.to_s(16) . length * 4
        format = "%#{bin_len}d"
        sprintf(format, self.to_s(2)) . split("") . map(&:to_i)
    end
end
def printHex(input)
    puts "#{ConvertArray.toHex(input)}"
end
def printBin(input)
    puts "#{ConvertArray.toBinary(input)}"
end
```

```
class DES
    def AddCheckBits(key)
        key_bits = key.hex_to_binary_arr
        # key_bits = @key.hex_to_binary_arr
        i = 8
        # extend key to 64 bits
        until i < 1 do
            sum = 0
            j = 7
            until j < 1 do
                sum += key_bits[7*i-j]
                j = j - 1
            end
            key_bits[7*i, 0] = sum % 2
            i = i - 1
        end
        key_bits
    end
    def permutation(count, origin_arr, permute_arr)
        out_arr = []
        count.times do |i|
            out_arr[i] = origin_arr[permute_arr[i] - 1]
        end
        out_arr
    end
    def shift(c_in, d_in)
        if @round == 1 || @round == 2 || @round == 9 || @round == 16
            c_out = @c_in[1..27] + [@c_in[0]]
            d_out = @d_in[1..27] + [@d_in[0]]
        else
            c_out = @c_in[2..27] + @c_in[0..1]
            d_out = @d_in[2..27] + @d_in[0..1]
        end
        return c_out, d_out
    end
```

```
def round_key(c_in, d_in)
    #@c_in, @d_in = shift()
    tmp = @c_in + @d_in
   roundkey = permutation(48, tmp, PC_2)
end
def getPlaintex_IP
    @plaintext
end
def getKey
    @key
end
def round_function
    @c_in, @d_in = shift(@c_in, @d_in)
   roundkey = round_key(@c_in, @d_in)
   1 = 0r
   r_e = permutation(48, @r, E)
    sbox_in_array = r_e.zip(roundkey).map { |x,y| x^y }
    sbox_8_in = sbox_in_array . each_slice(6) . map . inject([]) { |sum, s| sum << s.j</pre>
    sbox_outputs = []
    8.times do |i|
        tmp = sbox_8_in[i]
        j = tmp / 32 * 2 +
                            (tmp % 2)
        k = (tmp \% 32) / 2
        sbox_outputs[i] = Sbox[i][j][k]
    end
    sb_out = sbox_outputs.each.map.inject([]) { |sum, s| sum << sprintf("%04d", s.to_s
    sb_out = sb_out.flatten.each.map(&:to_i)
   sbox_outputs_P = permutation(32, sb_out, P)
   r = sbox_outputs_P.zip(@1).map { |x,y| x^y }
```

```
printHex(1)
        printHex(r)
        @1, @r = 1, r
        @round = @round + 1
    end
    def sixteen_rounds
        15.times do |i|
            puts " 第%d 轮的输出" % (i+2)
            round_function
        end
    end
    def initialize(key, plaintext)
        if key.to_s(16) .length == 14
            @key = AddCheckBits(key)
        else
            @key = key.hex_to_binary_arr
        @key = permutation(56, @key, PC_1)
        @plaintext = permutation(64, plaintext.hex_to_binary_arr, IP)
        @round = 1
        @c_in = @key[0..27]
        0d_{in} = 0key[28..55]
        @1 = @plaintext[0..31]
        @r = @plaintext[32..63]
    end
end
key = 0x5555555555555
plaintext = OxFEFEFEFEFEFEFE
des = DES.new(key, plaintext)
puts " 第 1 轮的输出"
des.round_function()
des.sixteen_rounds()
```

#### C AESSbox.rb

```
AesBox=[
    [0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0
    [0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0
    [0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0
    [0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0
    [0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0
    [0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0
    [0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0
    [0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0
    [0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0
    [0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0
    [0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0
    [0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0
    [0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0xe8, 0xdd, 0x74, 0x1f, 0x4b, 0xbd, 0
    [0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0
    [0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0
    [0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0
]
```

#### $\mathbf{D}$ AES.rb

```
puts "%02X & %02X & %02X & %02X \\\" % [self[0], self[4], self[8], self[12]]
        puts "%02X & %02X & %02X & %02X \\\\" % [self[1], self[5], self[9], self[13]]
        puts "%02X & %02X & %02X & %02X \\\\" % [self[2], self[6], self[10], self[14]]
        puts "%02X & %02X & %02X & %02X \\\\" % [self[3], self[7], self[11], self[15]]
    end
end
class Fixnum
    def add(addtion)
        self ^ addtion
    end
    def mult(multiplier)
        res=0
        if multiplier == 1
            res = self
        end
        if multiplier == 2
            if self >= 128
                res = self * 2 - 256
                res = res ^ OX1B
            else
                res = self * 2
            end
        end
        if multiplier==3
            res = self.mult(2).add(self)
        end
        res
    end
end
class XOR
```

block\_one.zip(block\_two).map {|x,y| x^y}

def self.sum(block\_one, block\_two)

```
end
end
class AES
   def init
        @plaintext_matrix, @key_matrix = _init_()
        #@key_matrix.print
        @round = 0
        @rc = [0x01, 0x02, 0x04, 0x08, 0x10, 0x20, 0x40, 0x80, 0x1b, 0x36]
        @round_enc_result = @plaintext_matrix
   end
   def subByte(byte)
        i = byte / 16
        j = byte % 16
       AesBox[i][j]
   end
   def computeRoundKey
        i = @round * 16
        # cycle shift left 1 byte
        shift1B = [@key_matrix[i-3], @key_matrix[i-2], @key_matrix[i-1], @key_matrix[i-4]]
        #puts shift1B
        4.times do |i|
            puts "%02X" % ( shift1B[i] = subByte(shift1B[i]) )
        end
        #puts shift1B
        constant = [@rc[@round], 0, 0, 0]
        #puts constant
        # compute key_matrix 的第四列, 即 key_matrix[16-19]
        \# w[i] = w[i-4] + constant + subbytes(shiftibyte(w[i-1])) ; w_{i-4} = key_matrix[i+1]
        # the i in last line is the i of aes algorithm,
        # not the index i of key_matrix array in the program
        @key_matrix += XOR.sum( XOR.sum(shift1B, constant), @key_matrix[i-16,4] )
        # compute w[i+1], w[i+2], w[i+3]
        3.times do
            i += 4
            @key_matrix += XOR.sum(@key_matrix[i-4, 4], @key_matrix[i-16, 4])
```

end

```
@key_matrix[16,16].print
end
def SubBytes
    # @round = 1
    16.times do |i|
        @round_enc_result[i] = subByte(@round_enc_result[i])
    end
end
def AddKey
    # key is in the @key_matrix
    @round_enc_result = XOR.sum(@key_matrix[@round * 16, 16], @round_enc_result)
    @round += 1
end
def ShiftRows
    round_enc_result = Array.new(@round_enc_result)
    \#@round\_enc\_result.print
    # 循环移位写成列向量的形式
    \# PC = [0, 4, 8, 12,
           5, 9,13, 1,
           10,14, 2, 6,
           15, 3, 7,11
    # ]
    pc = [0, 5, 10, 15,
          4, 9, 14, 3,
          8, 13, 2, 7,
          12, 1, 6, 11
    ٦
    16.times do |i|
        @round_enc_result[i] = round_enc_result[pc[i]]
    end
end
def MixColumn
    #@round_enc_result =
```

```
round_enc_result_matrix_2d = [
        [@round_enc_result[0],@round_enc_result[4],@round_enc_result[8],@round_enc_res
        [@round_enc_result[1],@round_enc_result[5],@round_enc_result[9],@round_enc_res
        [@round_enc_result[2],@round_enc_result[6],@round_enc_result[10],@round_enc_result[10]
        [@round_enc_result[3],@round_enc_result[7],@round_enc_result[11],@round_enc_re
   ]
   mix_matrix = [
        [2,3,1,1],
        [1,2,3,1],
        [1,1,2,3],
        [3,1,1,2]
    ٦
   mix_round_result_2dmatrx = Array.new(4) { Array.new(4) }
    4.times do |i|
        4.times do |j|
            mix_round_result_2dmatrx[i][j] = 0
            4.times do |k|
                mix_round_result_2dmatrx[i][j] ^= round_enc_result_matrix_2d[k][j].mul
            end
        end
    end
    pc_1 = [
        [0,0],[1,0],[2,0],[3,0],
        [0,1],[1,1],[2,1],[3,1],
        [0,2], [1,2], [2,2], [3,2],
        [0,3],[1,3],[2,3],[3,3]
   1
    16.times do |i|
        @round_enc_result[i] = mix_round_result_2dmatrx[pc_1[i][0]][pc_1[i][1]]
    end
end
def RoundFunction
    init
    self.AddKey
    puts "初始化轮密钥后的结果"
    @round_enc_result.print
    self.SubBytes
```

```
puts "字节替换后的结果"
       @round_enc_result.print
       self.ShiftRows
       puts " 行移位后的结果"
       @round_enc_result.print
       self.MixColumn
       puts " 列混淆后的结果"
       @round_enc_result.print
       self.computeRoundKey
       puts " 第一轮密钥: "
       @key_matrix[16, 16].print
       self.AddKey
       #puts "第一轮加密的结果"
       \#@round\_enc\_result.print
   end
key_matrix.print()
aes = AES.new
aes.init
aes.RoundFunction
```

end