

PRACTICAL No. 3

Topic: Parser Construction

Platform: Windows or Linux

Language to be used: Python or Java (based on the companies targeted for placement)

Aim: To find FIRST and Follow of a grammar.

(A) Write a program to find FIRST for any grammar. All the following rules of FIRST must be implemented.

For a generalized grammar: $A \rightarrow \alpha XY$

$FIRST(A) = FIRST(\alpha XY)$

$= \alpha$ if α is the terminal symbol (Rule-1)

$= FIRST(\alpha)$ if α is a non-terminal and $FIRST(\alpha)$ does not contain ϵ (Rule-2)

$= FIRST(\alpha) - \epsilon \cup FIRST(XY)$ if α is a non-terminal and $FIRST(\alpha)$ contains ϵ (Rule-3)

Input: Grammar rules from a file or from console entered by user.

Code⇒

```
import sys
sys.setrecursionlimit(60)

def first(string):
    #print("first({})".format(string))
    first_ = set()
    if string in non_terminals:
        alternatives = productions_dict[string]

        for alternative in alternatives:
            first_2 = first(alternative)
            first_ = first_ | first_2

    elif string in terminals:
        first_ = {string}

    elif string==' ' or string=='@':
        first_ = {'@'}

    else:
```

```

first_2 = first(string[0])
if '@' in first_2:
    i = 1
    while '@' in first_2:
        #print("inside while")

        first_ = first_ | (first_2 - {'@'})
        #print('string[i:]=', string[i:])
        if string[i:] in terminals:
            first_ = first_ | {string[i:]}
            break
        elif string[i:] == '':
            first_ = first_ | {'@'}
            break
        first_2 = first(string[i:])
        first_ = first_ | first_2 - {'@'}
        i += 1
    else:
        first_ = first_ | first_2

#print("returning for first({})".format(string), first_)
return first_

def follow(nT):
    #print("inside follow({})".format(nT))
    follow_ = set()
    #print("FOLLOW", FOLLOW)
    prods = productions_dict.items()
    if nT==starting_symbol:
        follow_ = follow_ | {'$'}
    for nt,rhs in prods:
        #print("nt to rhs", nt,rhs)
        for alt in rhs:
            for char in alt:
                if char==nT:
                    following_str = alt[alt.index(char) + 1:]
                    if following_str=='':
                        if nt==nT:
                            continue
                        else:
                            follow_ = follow_ | follow(nt)
                    else:
                        follow_2 = first(following_str)
                        if '@' in follow_2:

```

```

        follow_ = follow_ | follow_2-{'@'}
        follow_ = follow_ | follow(nt)
    else:
        follow_ = follow_ | follow_2
    #print("returning for follow({})".format(nT),follow_)
    return follow_

no_of_terminals=4

terminals = ['a','b','c','p']

no_of_non_terminals=4

non_terminals = ['S','A','B','C']

starting_symbol = 'S'

no_of_productions = 4

productions = ['S->AB/C',
               'A->a/BC',
               'B->p/@',
               'C->c'
               ]

# no_of_terminals=int(input("Enter no. of terminals: "))

# terminals = []

# print("Enter the terminals :")
# for _ in range(no_of_terminals):
#     terminals.append(input())

# no_of_non_terminals=int(input("Enter no. of non terminals: "))

# non_terminals = []

# print("Enter the non terminals :")
# for _ in range(no_of_non_terminals):
#     non_terminals.append(input())

# starting_symbol = input("Enter the starting symbol: ")

```

```

# no_of Productions = int(input("Enter no of productions: "))

# productions = []

# print("Enter the productions:")
# for _ in range(no_of Productions):
#     productions.append(input())

#print("terminals", terminals)

#print("non terminals", non_terminals)

#print("productions", productions)

productions_dict = {}

for nT in non_terminals:
    productions_dict[nT] = []

#print("productions_dict", productions_dict)

for production in productions:
    nonterm_to_prod = production.split(">")
    alternatives = nonterm_to_prod[1].split("/")
    for alternative in alternatives:
        productions_dict[nonterm_to_prod[0]].append(alternative)

#print("productions_dict", productions_dict)

#print("nonterm_to_prod", nonterm_to_prod)
#print("alternatives", alternatives)

FIRST = {}
FOLLOW = {}

for non_terminal in non_terminals:
    FIRST[non_terminal] = set()

for non_terminal in non_terminals:
    FOLLOW[non_terminal] = set()

#print("FIRST", FIRST)

```

```

for non_terminal in non_terminals:
    FIRST[non_terminal] = FIRST[non_terminal] | first(non_terminal)

#print("FIRST",FIRST)

FOLLOW[starting_symbol] = FOLLOW[starting_symbol] | {'$'}
for non_terminal in non_terminals:
    FOLLOW[non_terminal] = FOLLOW[non_terminal] | follow(non_terminal)

#print("FOLLOW", FOLLOW)

print("{: ^20}{: ^20}{: ^20}".format('Non Terminals','First','Follow'))
for non_terminal in non_terminals:
    print("{: ^20}{: ^20}{: ^20}".format(non_terminal,str(FIRST[non_terminal]),str(FOLLOW[non_terminal])))

```

Implementation: FIRST rules

Output: FIRST information for each non-terminal

```

PROBLEMS OUTPUT TERMINAL JUPYTER DEBUG CONSOLE
Windows PowerShell
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Install the latest PowerShell for new features and improvements! https://aka.ms/PSWindows

PS C:\Users\asus\OneDrive\Desktop\college_sem\5th sem\cd_lab\prac3> python -u "c:\Users\asus\OneDrive\Desktop\college_sem\5th sem\cd_lab\prac3\LL1_PARS.py"
Non Terminals      First      Follow
S      {'a', 'b', 'c', 'p', '@'}  {'$'}
A      {'a', 'b', '@'}        {'$', 'p'}
B      {'@', 'p'}              {'$'}
C      {'c'}                   {'$'}
PS C:\Users\asus\OneDrive\Desktop\college_sem\5th sem\cd_lab\prac3>

```

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(B) Calculate Follow for the given grammar manually, input the follow information and Construct the LL (1) parsing table using the FIRST and FOLLOW values computed above.

⇒

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a) LL(1) parser

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$S \rightarrow ABC$

$A \rightarrow a|b|e$

$B \rightarrow p|e$

$C \rightarrow c$

First

Follow

S $\{a, b, c, p, e\}$

$\{ \$ \}$

A $\{a, b, e\}$

$\{ \$, p \}$

B $\{e, p\}$

$\{ \$ \}$

C $\{c\}$

$\{ \$ \}$

Parsing table

	a	b	c	p	\$
S	$S \rightarrow AB$	$S \rightarrow AB$	$S \rightarrow C$	$S \rightarrow AB$	
A	$A \rightarrow a$	$A \rightarrow b$		$A \rightarrow e$	$A \rightarrow e$
B				$B \rightarrow p$	
C			$C \rightarrow c$		