

Converting a CFG into CNF

$S \rightarrow aSb \mid \epsilon$ (getting rid of ϵ -prod?)

$S \rightarrow aSb \mid ab$

$S \rightarrow aAb \mid A$
 $A \rightarrow S \mid ab$

Give CFG G :

- ① Add "some productions" $\rightarrow \hat{G}$
- ② Remove unit and ϵ -productions $\rightarrow \hat{\hat{G}}$
- ③ Simplify to CNF $\rightarrow G'$

step 1: \hat{P} is initialized with P
(productions), $\gamma \in (V \cup \Sigma)^*$, $A, B \in V$

Repeat the following two steps
until no further updates are made:

- * if $A \rightarrow B$ and $B \rightarrow \gamma$ are in \hat{P} , then add $A \rightarrow \gamma$ to \hat{P} .
- * if $A \rightarrow aB\gamma$ and $B \rightarrow \epsilon$ are in \hat{P} , then add $A \rightarrow a\gamma$ to \hat{P} .

Step 2: Remove all unit and ϵ -productions

Also, remove unreachable non-terminals.

$S \rightarrow aTb \mid \epsilon \mid ab \mid$

$T \rightarrow S \mid ab \mid \epsilon \mid aTb$



step 1

step 2: $S \rightarrow aTb \mid ab$

$T \rightarrow ab \mid aTb$

step 3: For each terminal a ,
create a new non terminal (say A)
and update the rules.

.. step 3.1:

and update ...
step 3.1:

$$S \rightarrow ATB \mid AB$$

$$T \rightarrow AB \mid ATB$$

$$A \rightarrow a$$

$$B \rightarrow b$$

Then break the long sequences
by introducing new non-terminals.

$$S \rightarrow LB \mid AB$$

$$T \rightarrow AB \mid LB$$

$$L \rightarrow AT$$

$$A \rightarrow a$$

$$B \rightarrow b$$

note: we could
have merged S
and T but not
necessary

$$S \rightarrow SS \mid [S] \mid \epsilon \quad \underbrace{[] \mid S}_{\text{step 1}}$$

↓

$$\text{step 2: } S \rightarrow SS \mid [S] \mid []$$

↓

$$\text{step 3: } S \rightarrow SS \mid AR \mid LR$$

$$A \rightarrow LS$$

$$L \rightarrow [$$

$R \rightarrow]$

Closure properties of CFLs

* Union: if $A = L(G_1)$ and $B = L(G_2)$ for CFGs G_1 and G_2 , then $A \cup B$ is CFG? yes

$S \rightarrow S_1 \mid S_2 \rightarrow$ start nonterminal of G_2 .
start symbol of G_1

* Concatenation: And then $S_1 \rightarrow \dots$
 $S_2 \rightarrow \dots$
 $S \rightarrow S_1 S_2$ yes!

* Intersection:

the intersection of two CFLs is not necessarily CF.

$C = \{a^n b^n c^n : n \geq 0\} \rightarrow$ not CFG

$A = \{a^n b^n c^k : n \geq 0, k \geq 0\}$

$B = \{a^k b^n c^n : n \geq 0, k \geq 0\}$

$A \cap B = C$
↓ ↓ ↘ not CFL
CFL CFL