

# COMPSCI 2AC3 Assignment 1

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10 February 2024

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# 1 Answer 1

The given DFA can be defined as follows:

- $Q = q1, q2, q3, q4, q5$
- $\Sigma = 0, 1$
- $S = q1$
- $F = q3$

This DFA has 5 states, state  $q1$  is the starting state, state  $q3$  is the only final state and state  $q4$  is considered as the garbage state for rejected inputs.

The DFA starts at state  $q1$ , it then check the left most bit. If the left most bit is '0' then the input is rejected and the DFA sends it to state  $q4$  which is considered as the garbage state. After entering the garbage state the input cannot exit it and has been rejected as asked in the question.

If the left most bit is '1' then the input is sent to state  $q2$ , if the input ends in this state, the remainder is 1 and the input is not divisible by 3. If the input ends at state  $q5$  then, the remainder is 2 and the input is not divisible by 3 either.

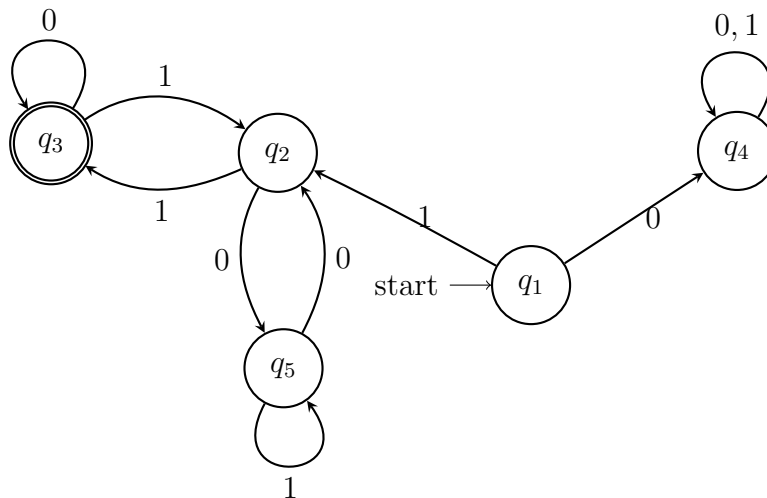


Figure 1: DFA for A

## 2 Answer 2

Let  $N_3 = (Q_3, \Sigma_3, \Delta_3, S_3, F_3)$  then, the NFA of  $N_3$  can be defined as follows:

- States of  $N_3$  will be a union of states from  $N_1$  and  $N_2$   
 $Q_3 = Q_1 \cup Q_2$
- All the  $\Sigma$  are the same for  $N_3$   
 $\Sigma = \Sigma_1 = \Sigma_2$
- The starting state of  $N_3$  is the starting state of  $N_1$   
 $S_3 = S_1$
- The final state of  $N_3$  is the final state of  $N_2$   
 $F_3 = F_2$
- The transitions  $\Delta$  are as follows:
  - For each transition in  $\Delta_1(p, a, q)$  where,  $p, q \in Q_1$  and  $a \in \Sigma$ , we include that in  $\Delta_3$ .
  - For each transition in  $\Delta_2(p, a, q)$  where,  $p, q \in Q_2$  and  $a \in \Sigma$ , we include that in  $\Delta_3$ .
  - For each transition in  $\Delta_1(p, a, q)$  where,  $p \in F_1$ ,  $q \in S_2$  and  $a \in \Sigma$ , we include that in  $\Delta_3$ .

### Informal explanation:

The basic idea of  $N_3$  is that it starts at the starting states of  $S_1$  and once is reached and once it reached any  $F_1$ , it transitions into  $S_2$  then, it carries out transitions of  $N_2$  until it reaches a state in  $F_2$ . This ensures that strings accepted by  $N_2$  are a concatenation of string accepted by  $N_1$  and string accepted by  $N_2$ .

### 3 Answer 3

The statement " $L(N_1) = \sim L(N_2)$ " is incorrect. I used the NFA I made for Question 4 to prove that the statement is incorrect using a counterexample.

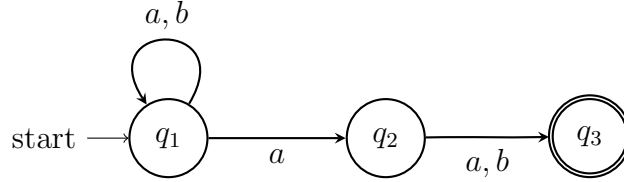


Figure 2: NFA 1

The NFA below accepts strings whose second letter from the last is an 'a'. From that we can see that the example 'abaa' is an accepted input. Therefore,  $abaa \in L = N_1$

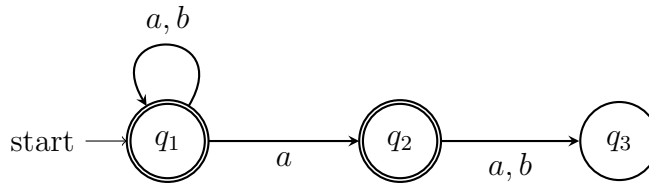


Figure 3: NFA 2 with switched final states.

According to the given statement " $L(N_1) = \sim L(N_2)$ ", the above NFA should not accept any string whose second last letter is 'a'. However, after simulating the example 'abaa' we see that the NFA does in fact accept it since, state q1 has a self loop and it is an accepted final state. this proves that  $abaa \in L = N_2$

Hence, statement is proven to be incorrect.

## 4 Question 4

### 4.1 Part A

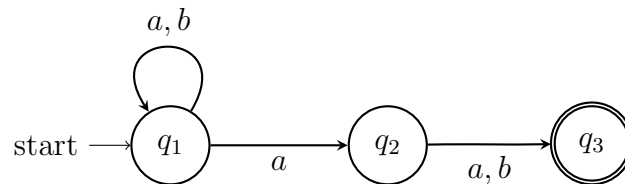


Figure 4: NFA for A

### 4.2 Part B

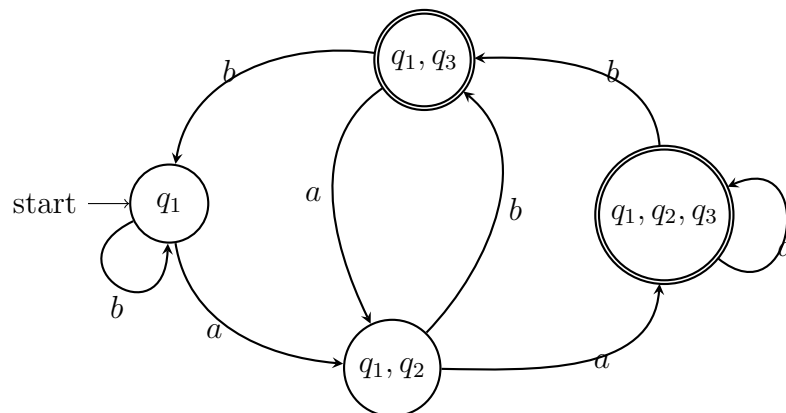


Figure 5: DFA for A

The steps to convert the NFA onto a DFA are:

- After transition  $b$  and a state ' $q_1$ ' goes to ' $q_1$ ' and state ' $q_1, q_2$ ' respectively.
- After transition  $b$  and a state ' $q_1, q_3$ ' goes to ' $q_1$ ' and state ' $q_1, q_2$ ' respectively.

- After transition b and a state 'q1,q2' goes to 'q1,q3' and state 'q1,q2,q3' respectfully.
- After transition b and a state 'q1,q2,q3' goes to 'q1,q3' and state 'q1,q2,q3' respectfully.