The Entity-Relationship Model COMPSCI 2DB3: Databases

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The relational data model

Data model

The rules by which real-world data can be represented and structured.

The relational data model All data is modeled as a collection of *tables*.

Schema Describe how stored data is structured.

- Conceptual schema: in terms of the data model.
- Physical schema: details in terms of layout on disk: files, file structure, auxiliary data structures (indices),

Instance The actual data (adhering to some schema).

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Example: The relational data model

 $\textbf{cgrades} \quad (\underline{id}, \qquad \text{name}, \qquad \text{program}, \quad \text{age}, \quad \text{score})$

The relational schema describing a table

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Instance The actual data (adhering to some schema).

Example: The relational data model

 cgrades
 (id, name, program, age, score)

 Table name
 Columns (attributes) & their domains

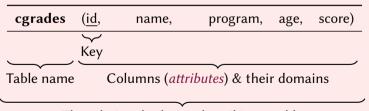
The relational schema describing a table

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The relational schema describing a table

Schema Describe how stored data is structured.

- Conceptual schema: in terms of the data model.
- ► Physical schema: details in terms of *layout on disk*: files, file structure, auxiliary data structures (indices),

Instance The actual data (adhering to some schema).

Example: The relational data model

| cgrades | (<u>id</u> , | name, | program, | age, | score) | |
|---------|---------------|---------|----------|------|--------|------|
| | 53666 | Jones | cs | 18 | 3.4 | |
| | 53688 | Smith | ee | 18 | 3.2 | |
| | 53650 | Smith | math | 19 | 3.8 | |
| | 53831 | Madayan | music | 11 | 1.8 | |
| | 53832 | Guldu | music | 12 | 2.0 | |

How to create a schema?

Requirements Analysis

- ▶ What data can this application collect?
- ► What data does this application need?
- What does the application do with the data?

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Is a database management system the right fit? E.g., big data analytics, machine learning, visualization, ...

How to create a schema?

Requirements Analysis

- What data can this application collect?
- ► What data does this application need?
- What does the application do with the data?

Is a database management system the right fit?
E.g., big data analytics, machine learning, visualization, ...

What about agile development?

It is complex, costly, and sometimes impossible to refactor away wrong data choices. E.g.,

- ► How do you add missing data after the fact?
- How do you restructure data to improve performance?

Agile development *requires* expertise to make the right choices at the right time.

Requirements Analysis: the entity-relationship data model

Relational data: Tables represent everything Faculty(fid, name, location, mail), Course(cid, name, year, credits).

Teaches(fid, cid).

Requirements Analysis: the entity-relationship data model

```
Relational data: Tables represent everything

Faculty(\underline{\text{fid}}, name, location, mail),

Course(\underline{\text{cid}}, name, year, credits).

Factorises

Relationship
```

Requirements Analysis: the entity-relationship data model

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Relational data: Tables represent everything

Faculty(fid, name, location, mail),

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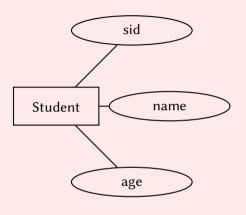
Relationship
```

Entity-Relationship data model

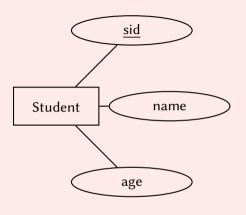
- ► A semantic data model: Specify meaning and not representation.
- ► Can model high-level *concepts*: entities, relationships, attributes,
- ► Can easily be translated to well-designed tables (*representation*).
- ► Resulting ER-diagrams are easier to document and discuss with *non-specialists*.

Student

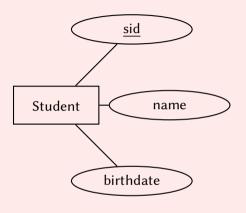
The *entity set* student.



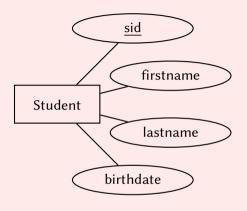
Students have attributes: student id, name, and age.



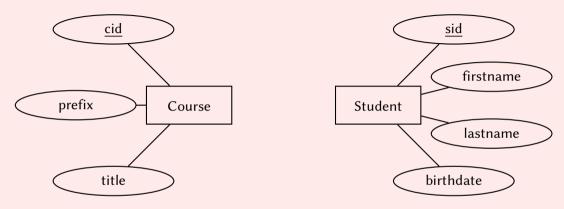
Student id is a key: a unique identifier.



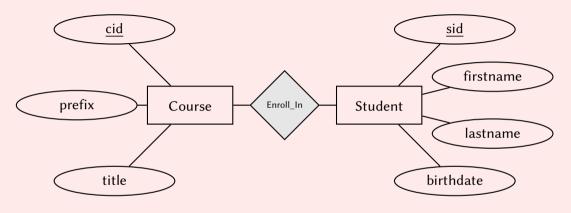
Age is a *derived* attribute: birthdate is easier to maintain!



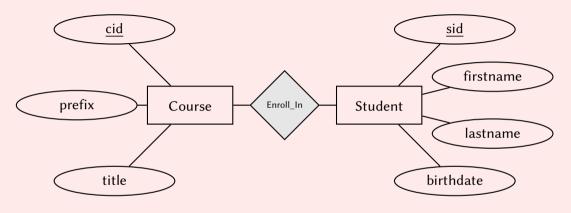
Names are *complex*: does the application need first and last names?



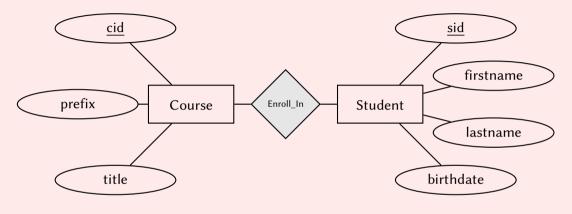
Another *entity*: a course.



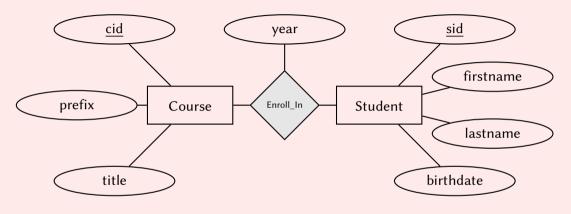
Each student can enroll in classes: a relationship set.



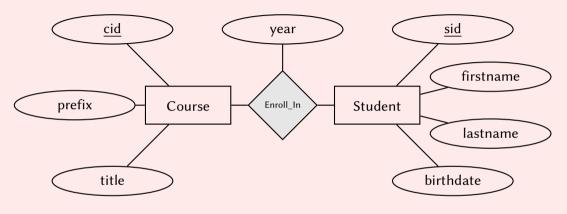
This is a many-to-many relationship set.



Relationships are *unique*: A student cannot take a single course multiple times.



Relationship sets can have attributes: the *year* in which the course was taken.



Even with relationship attributes: Students can take a course only once!

Intermission: Names...

Question: Did you ever have problems with your name?

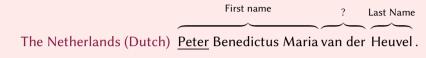
Vote at https://strawpoll.com/cp6fysdxr.

Or: go to https://strawpoll.live and use the code 112128.

Intermission: Names—an example

The Netherlands (Dutch) Peter Benedictus Maria van der Heuvel .

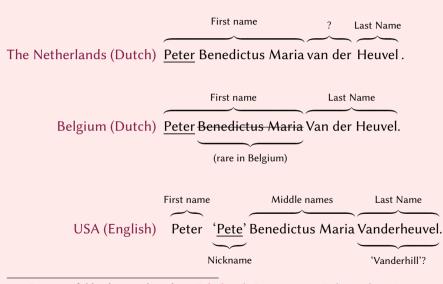
Intermission: Names—an example





Some useful background reading: Falsehoods Programmers Believe About Names.

Intermission: Names—an example



Some useful background reading: Falsehoods Programmers Believe About Names.

Intermission: What to do with names?

Ask yourself: Why do you need the name?

- 1. Do not put up limitations, unless you have a specific goal.
- 2. Be clear to users so they can provide correct info. E.g.,

Informal Ask what they want to be called (e.g., Hi!).

Formal Ask what they want to be called formally (e.g., letters).

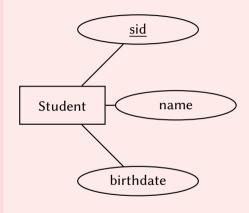
Billing Ask the name as used by their bank.

Shipping Ask the name as used by the mail carrier.

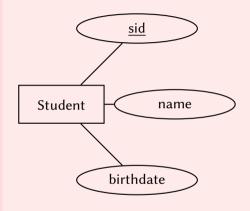
Travel Ask the name as printed in their travel documents.

3. Have sensible defaults: e.g., prefill billing names with account information.

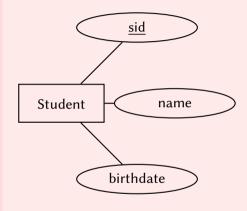
There is no standard solution that works in all cases!



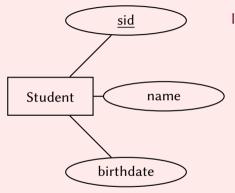
- ► An *entity* is an object in the real world.
- Entities are described via a set of attributes.
- Attributes have a *domain* (e.g., text, number, date).



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- Attributes have a *domain* (e.g., text, number, date).
- ► *Key*: set of attributes that identify an entity.
- ► All keys are *candidate keys*.
- One is chosen as the primary key.



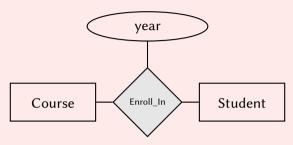
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- ► *Key*: set of attributes that identify an entity.
- ► All keys are *candidate keys*.
- ► One is chosen as the *primary key*.
- ► An *entity set*: collection of similar entities.
- An *instance*: snapshot of the entities ("data").



Instance

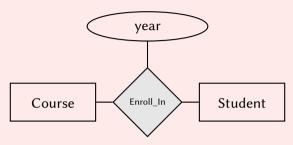
| <u>sid</u> ı | | name | birthdate | | |
|--------------|-------|---------|-------------------|--|--|
| | 53666 | Jones | May 18, 1999 | | |
| | 53688 | Smith | June 2, 1994 | | |
| | 53650 | Smith | December 21, 1995 | | |
| | 53831 | Madayan | February 7, 2000 | | |
| | 53832 | Guldu | April 1, 1991 | | |
| | | | | | |

The entity-relationship model: Relationships



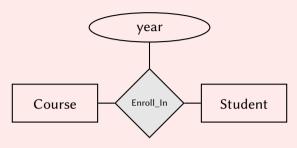
- ► A *relationship* relates two or more entities.
- Standard relationship: many-to-many: Students can enroll in many courses, courses can have many enrolled students.
- ► Relationships can have additional *attributes*.

The entity-relationship model: Relationships



- ► A *relationship* relates two or more entities.
- Standard relationship: many-to-many: Students can enroll in many courses, courses can have many enrolled students.
- ► Relationships can have additional *attributes*.
- ► A *relationship set*: collection of similar relationships between entities.
- ► A *instance*: snapshot of the relations ("data").

The entity-relationship model: Relationships



Instance

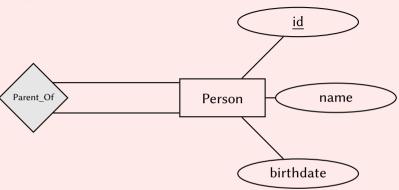
| Course | Student | year | Course | Student | year |
|----------|-----------|------|----------|-----------|------|
| Course_1 | Student_1 | 2019 | Course_2 | Student_1 | 2019 |
| Course_1 | Student_2 | 2020 | Course_2 | Student_3 | 2021 |

More on relationships

Relationships are at the core of the entity-relationship model.

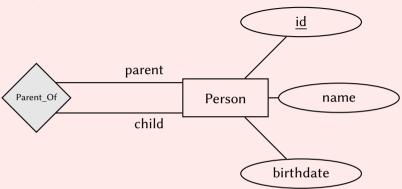
- ► Self-referential relationship sets.
- ► *n*-ary relationship sets.
- ► Key constraints.

Self-relationships



► Entities in an entity set can be related to each others. E.g., *Persons* can be *parents*, *friends*, *partners*, *colleagues*.

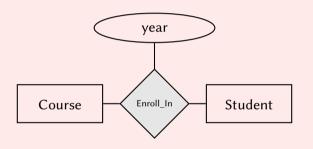
Self-relationships



- Entities in an entity set can be related to each others.
 E.g., Persons can be parents, friends, partners, colleagues.
- ► Typically, one names the *roles* in such relationships sets.

Ternary relationships

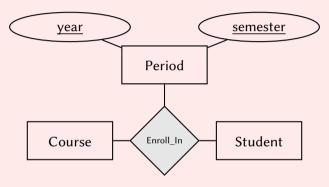
What if a student can enroll several times?



Pairs (student, course) must be unique.

Ternary relationships

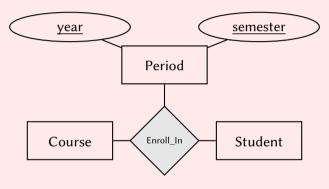
What if a student can enroll several times?



Pairs (student, course, period) must be unique.

Ternary relationships

What if a student can enroll several times?



Ternary relationships are rare: this can be modeled in many other ways.

Types of binary relations (2DM3)

A binary relation $R \subseteq A \times B$ is

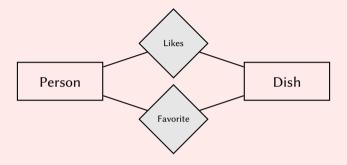
Many-to-Many if there are no restrictions on the relation.

One-to-Many if each $a \in A$ is related to *at-most-one* $b \in B$.

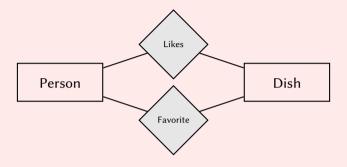
► If A has total participation: R is a function of A, each $a \in A$ is related to exactly-one $b \in B$.

One-to-One each A is related to at-most-one B and vice-versa.

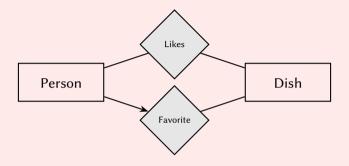
- ► If A has total participation: R is an injection, each $a \in A$ is related to exactly-one-and-unique $b \in B$.
- ▶ If *B* has total participation: *R* is an surjection, each $b \in B$ is related to exactly-one-and-unique $a \in A$.
- ▶ If *A*, *B* have *total participation*: *R* is a *bijection*, there is a one-to-one mapping between all $a \in As$ and $b \in Bs$.



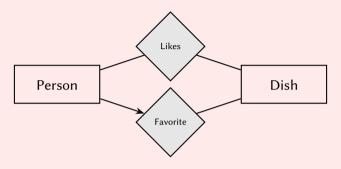
A person *likes* many dishes, but has *one* favorite dish.



This says: a person has many favorite dishes!



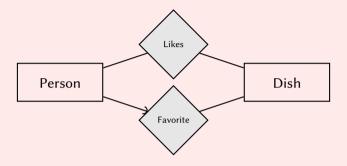
Key constraint: A person has *at-most-one* favorite dish.



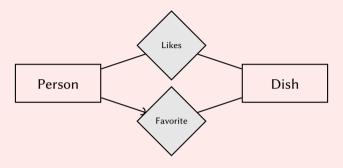
Key constraint: A person has *at-most-one* favorite dish.

→ R: entity E partakes at-most-once in a R-relationship (partial).

6/2

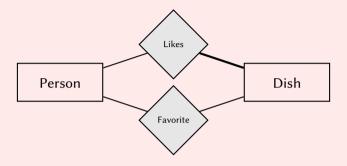


Key constraint: A person has *exactly-one* favorite dish.

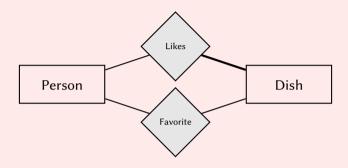


Key constraint: A person has *exactly-one* favorite dish. $E \longrightarrow R$: entity E partakes exactly-once in a R-relationship (total).

6/2

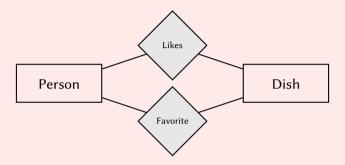


Participation constraint: Every dish is *liked*.



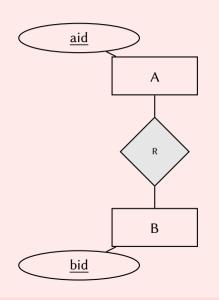
Participation constraint: Every dish is *liked*.

 $E \longrightarrow R$: entity E partakes at-least-once in a R-relationship (total).



Remarks on notation

- ▶ We follow the notation of the textbook.
- ▶ Many other sources use different notations (e.g., arrows reversed).
- ► We use dedicated notation for exactly-once participation (———).

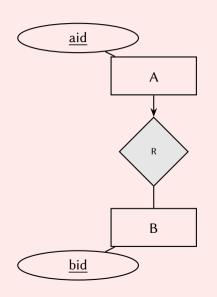


| Entities A |
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| δ |
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| β γ |

| Relationships R | | | |
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| | β | 1 | |
| | β | 2 | |
| | γ | 3 | |

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many-to-many

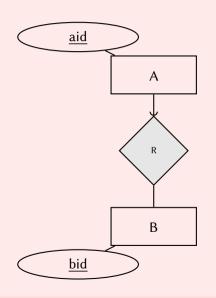


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one-to-many (partial)

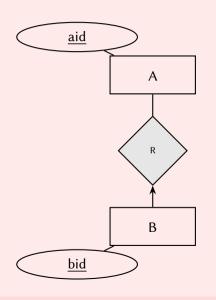


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one-to-many (total) (function)



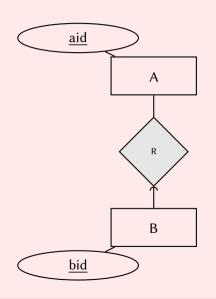
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Entities B

many-to-one (partial)

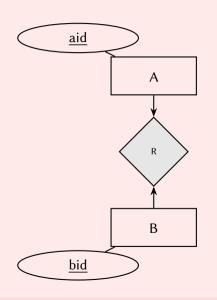


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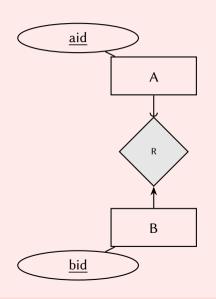
many-to-one (total)



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one-to-one (partial, partial)



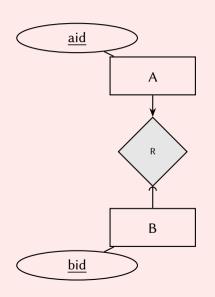
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Entities B

one-to-one (total, partial) (injection)

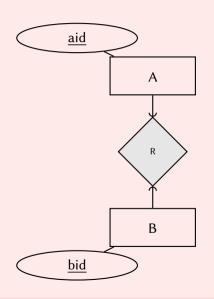


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| Entities B |
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one-to-one (partial, total) (surjection)



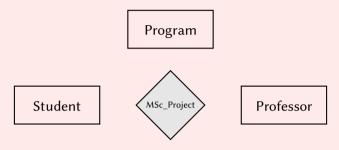
| Entities A |
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| Relationships R | | |
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| α | 1 | |
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Entities B

one-to-one (total, total) (bijection)

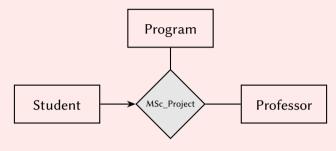


- A student can do *one* Master Project.
- Each project has a supervising professor and is performed in a degree program.

Question: Where should the key constraints go?

Vote at https://strawpoll.com/3f6eovfz8.

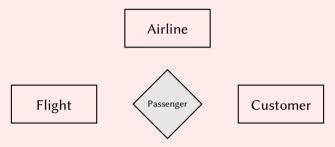
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- A student can do *one* Master Project.
- Each project has a supervising professor and is performed in a degree program.

Answer: Constraint on Student

- ► Key constraint for the participation of students.
- ► Professors can supervise *many* Master Projects.
- Many students can perform a Master Project within a program.

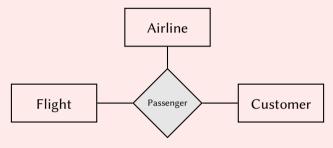


Each passenger (customer) of a flight bought a single ticket via a single airline.

Question: Where should the key constraints go?

Vote at https://strawpoll.com/gwbwsh9hy.

Or: go to https://strawpoll.live and use the code 241367.



Each passenger (customer) of a flight bought a single ticket via a single airline.

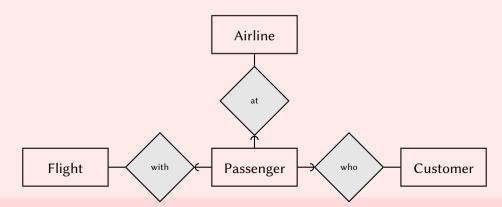
Answer: This is not a key constraint

- ► Customers can be on *many* flights.
- ► Airlines can sell tickets for *many* flights.
- ► Flights can have *many* customers, each from different airlines.

Each passenger (customer) of a flight bought a single ticket via a single airline.

Answer: This is not a key constraint

We can make Passenger a entity!



The basics: entities, relationships, and constraints

We covered enough to model *most typical* situations. We have already seen situations we could not model, however!

Advanced modeling features

- Weak entities.
- ► ISA Hierarchies.

The need for weak entities

Definition

A weak entity is 'owned by' another identity:

A weak entity can only be uniquely identified in conjunction with owning entities.

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Examples

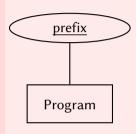
Assignments belong to courses:

- Many courses have a 1-st assignment.
- ► The 1-st assignment of COMPSCI 2DB3: *ER data model*.

Courses belong to a degree program (at McMaster):

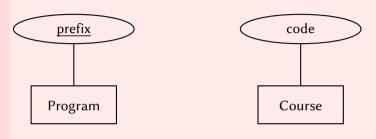
- ► There are several courses with code 2DB3 (e.g., COMPSCI 2DB3 versus SFWRENG 2DB3).
- ► There is only one COMPSCI 2DB3 (this course).

Weak entities in practice



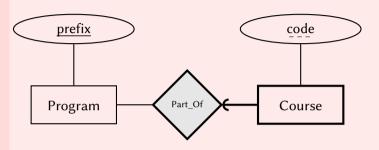
A degree program is an entity with a *unique prefix* (COMPSCI, SFWRENG, ...).

Weak entities in practice



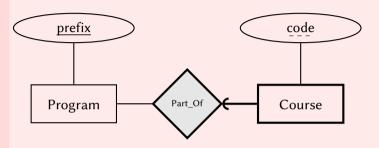
A course has a a *unique code* within a degree program (2DB3 within COMPSCI).

Weak entities in practice

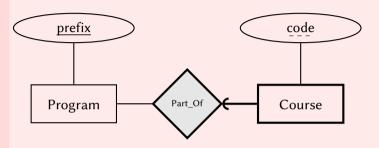


Course is a weak entity and Part_Of is an identifying relationship (bold).

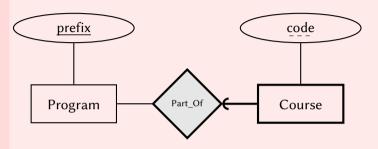
2/2



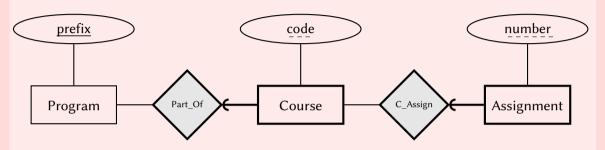
Weak entities must participate in their identifying relationship(s):



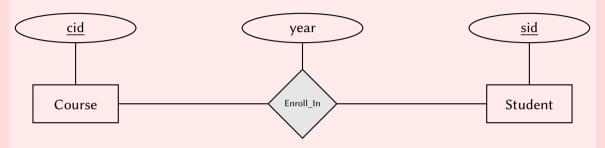
Attribute code is a *partial key* for the weak entity Course.



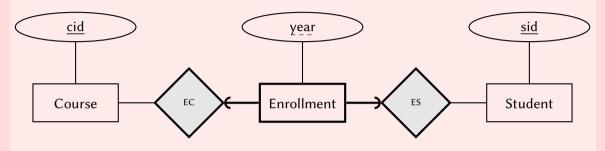
The primary key for weak entity Course is the pair (<u>prefix</u>, <u>code</u>)!



The primary key for weak entity Assignment is the triple (<u>prefix</u>, <u>code</u>, <u>number</u>).



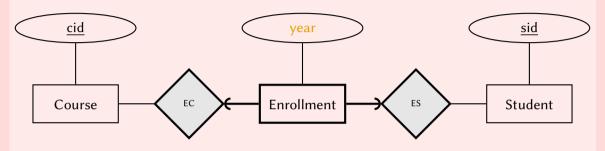
Problem: Student could only enroll once



Problem: Student could only enroll once

Turn relationship Enroll_In into a weak entity Enrollment.

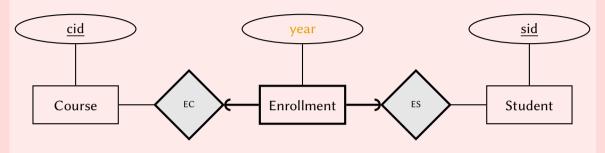
- Enrollment has partial key year.
- Enrollment is owned by both a Student and a Course.
- ▶ The primary key of Enrollment is the triple ($\underline{\text{sid}}$, $\underline{\text{cid}}$, $\underline{\text{year}}$).



A word of caution

That something *can* be written, does not mean it *should* be written.

Do *not* use weak entities when simpler constructs suffice!



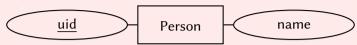
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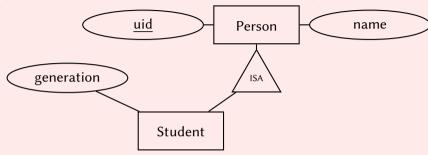
Do *not* use weak entities when simpler constructs suffice!

E.g., do not use weak entities to express normal relationships.

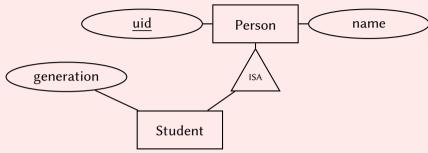
This is a classic example of a *class hierarchy*.



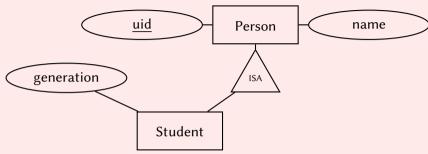
Students, Faculty, and Staff are all *peoples* with names and University IDs.



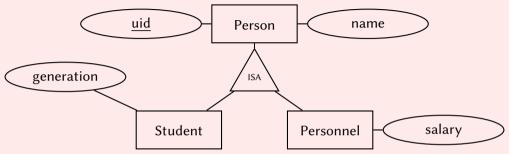
Each student is part of a *generation* (e.g., Class of 2021).



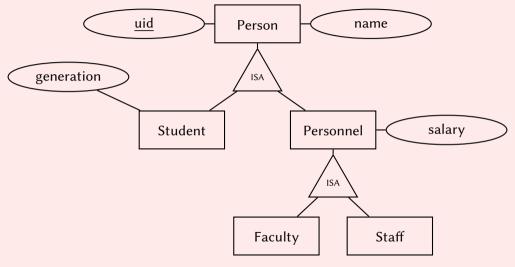
We say that entity Student is a *specialization* or *subclass* of entity Person.



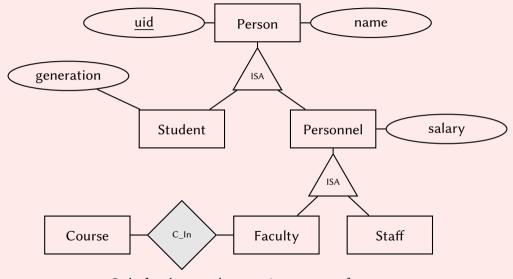
We say that entity Person is a *generalization* or *superclass* of Student.



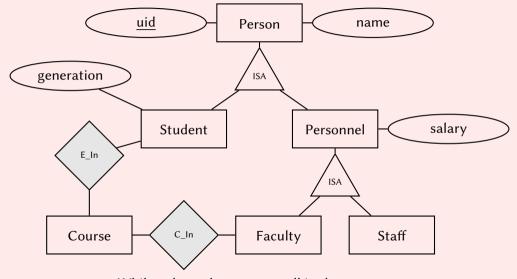
Each faculty and staff member is personnel with a *salary*.



We can further specialize faculty and staff.



Only faculty members are instructors of courses.



While only students can enroll in these courses.

ISA hierarchies require documentation

Every combination of entities in the hierarchy can exists. E.g.,

- there can be Person entities that are not students, faculty, or staff;
- there can be Faculty entities that are also Staff and Students,

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Terminology

Overlap constraints Can a single entity belong to *multiple* subclasses? E.g., can an entity be both Faculty and Student?

Covering constraints Must every superclass entity *also belong* to one of its subclasses? E.g., must every Person entity be a Student, Faculty, and/or Staff?

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Document all constraints.

On good design: Attributes versus Entities

Attributes should represent a *single atomic* value.

Atomic here means a base value without internal structure: No lists, complex objects,

Entities represent *complex objects*: multiple attributes.

Example: Phone numbers

If you want to model a Person entity with

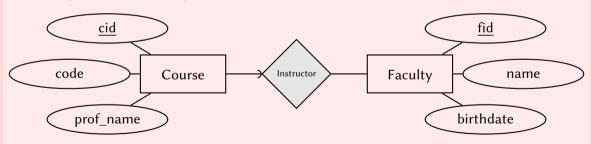
- ▶ a *single phone number*: use an attribute.
- ► multiple phone numbers: use a phone number (weak) entity.

On good design: Redundancy

What *redundancy*: storing the same information in multiple ways.

Why Redundancies cause *inconsistencies* during updates. Redundancies *waste* storage space.

Example of Redundancy



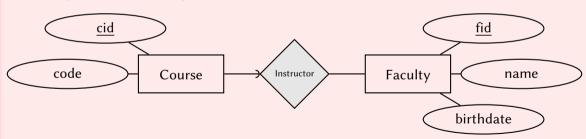
The instructor name is stored twice.

On good design: Redundancy

What redundancy: storing the same information in multiple ways.

Why Redundancies cause *inconsistencies* during updates. Redundancies *waste* storage space.

Example of Redundancy (removed)



On good design: Keys & weak entities

- Real-world entities might not have a clear key.
 E.g., not everyone has a government-provided SIN.
- We often create unique internal identifiers for entity sets.
 E.g., Mac ID as used within McMaster University.
 E.g., sequential counters in many systems.

Weak entities and normal entities

- Every weak entity can be turned into a normal entity.
- Use weak entities only if there there is a clear identifier in the context of an owner.
 E.g., the course code '2DB3' is identifying a course if you know the program.
 E.g., the shirt number '5' is identifying a player if you know the sport team.

Summary

Entity-relationship modeling

- ▶ yields a *conceptual schema*: high-level description of the data in terms of:
 - entities, attributes, and relationships,
 - weak entities, ISA hierarchies, and aggregation,
 - constraints: keys, relationship participation, ISA hierarchies;
- depends on a complete requirements analysis;
- ▶ is highly subjective: many ways to model the same data.

Later lectures: Translating an ER-diagram into tables.