

$$\begin{aligned} & \# \cap \sim a \quad \Sigma = \{a, b\} \\ L(\# \cap \sim a) &= L(\#) \cap L(\sim a) \\ &= \Sigma \cap (\Sigma^* - L(a)) \\ &= \{a, b\} \cap (\Sigma^* - \{a\}) = \{b\} \end{aligned}$$


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\*  $\{x \in \{a, b, c\}^*, x \text{ does not have any } a \text{ in it}\}$

\*  $\alpha = \sim a$  X  $aa \in L(\alpha)$

\*  $\beta = \sim(@a@)$  ✓

\*  $\gamma = \sim(a^+)$  X  $bac \in L(\gamma)$

\*  $\delta = (c + b)^*$  ✓

$$L(\beta) = L(\delta) \\ \text{but } \beta \neq \delta$$


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Patterns are redundant. We actually don't need all of the atomic patterns or compound patterns.

\* We can replace  $\#$  with  $a + b + c$  (for  $\Sigma = \{a, b, c\}$ )

\*  $\varepsilon$  with  $\emptyset^*$

\*  $\alpha \cap \beta$  with  $\sim(\sim\alpha + \sim\beta)$

\*  $\sim\alpha$  with ..? (more difficult)

... ✓ ... ← c . 1 . 2

\*  $\sim \alpha$  with ... (more ...)

\*  $\alpha = \sim (@ b @ b @)$  where  $\Sigma = \{a, b\}$   
 find "equivalent"  $\beta$  without  $\sim$ ?

$$a^* b a^* + a^*$$

## Regular Expressions

A regular expression is a pattern that only uses the following:

\* Atomic:  $\emptyset, a, \epsilon$   $a \in \Sigma$

\* Compound:  $+, \cdot, *$   
 $\downarrow$  concat

\* The order of precedence for compound operations:  $* > \cdot > +$

$$a c^* b + a^* b^* \equiv (a(c^*)b) + ((a^*)(b^*))$$

\* Regular expression define equivalency classes.

$$\alpha \equiv \beta \text{ iff } L(\alpha) = L(\beta)$$

\* Reflexive  $\alpha \equiv \alpha$

\* Symmetric:  $\alpha \equiv \beta \iff \beta \equiv \alpha$

\* Transitive:  $\alpha \equiv \beta, \beta \equiv \gamma \implies \alpha \equiv \gamma$

\*  $\alpha \leq \beta$  if  $L(\alpha) \subseteq L(\beta)$   
 partial order

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Theorem. The following statements are equivalent for every set  $A \subseteq \Sigma^*$ :

(i)  $A$  is regular (has a DFA)

(ii)  $\exists$  NFA  $N$ , s.t.  $L(N) = A$

(iii)  $\approx \approx N$  with  $\epsilon$ -transitions s.t.  
 $L(N) = A$

(iv)  $\exists$  pattern  $\alpha$  s.t.  $L(\alpha) = A$

(v)  $\exists$  regular expression  $\beta$  s.t.  $L(\beta) = A$