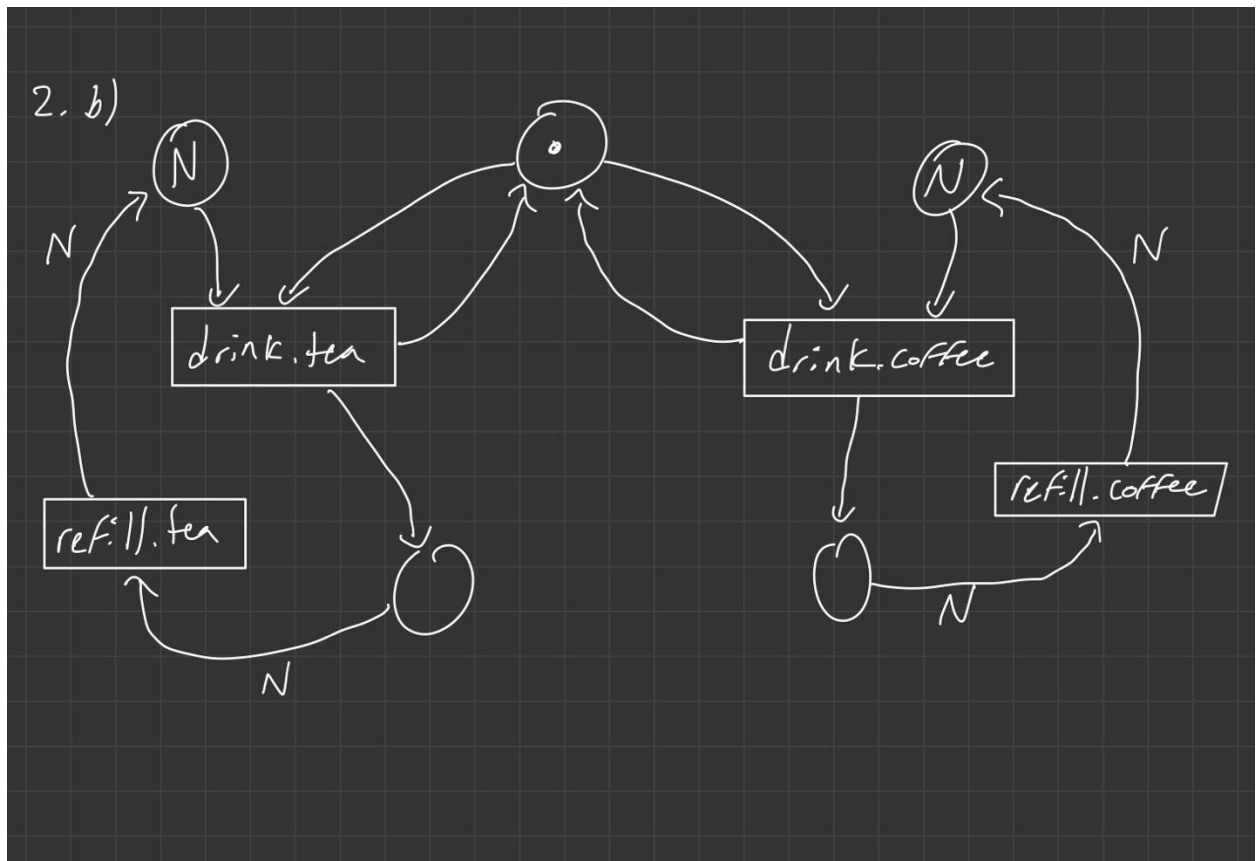


2SD3 Midterm

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1. DOOR_CONTROLLER = CLOSED,
OPEN = (rear -> CLOSED | front -> OPEN | both -> OPEN | neither -> CLOSED),
CLOSED = (rear -> CLOSED | front -> OPEN | both -> OPEN | neither -> CLOSED).
2. A) const N = 2
CUSTOMER = (drink.tea -> CUSTOMER | drink.coffee -> CUSTOMER).
STAFF_MEMBER = (refill.tea -> STAFF_MEMBER | refill.coffee -> STAFF_MEMBER).
TEA_MACHINE = TEA_MACHINE[N],
TEA_MACHINE[i:0..N] = (when (i>0) drink.tea -> TEA_MACHINE[i-1]
| when (i==0) refill.tea -> TEA_MACHINE[N]).
COFFEE_MACHINE = COFFEE_MACHINE[N],
COFFEE_MACHINE[i:0..N] = (when (i>0) drink.coffee -> COFFEE_MACHINE[i-1]
| when (i==0) refill.coffee -> COFFEE_MACHINE[N]).
const C = 2
|| CUSTOMERS = (forall [c:0..C] customer [c] : CUSTOMER).
|| HOTEL_CAF = (CUSTOMERS || STAFF_MEMBER || TEA_MACHINE || COFFEE_MACHINE).

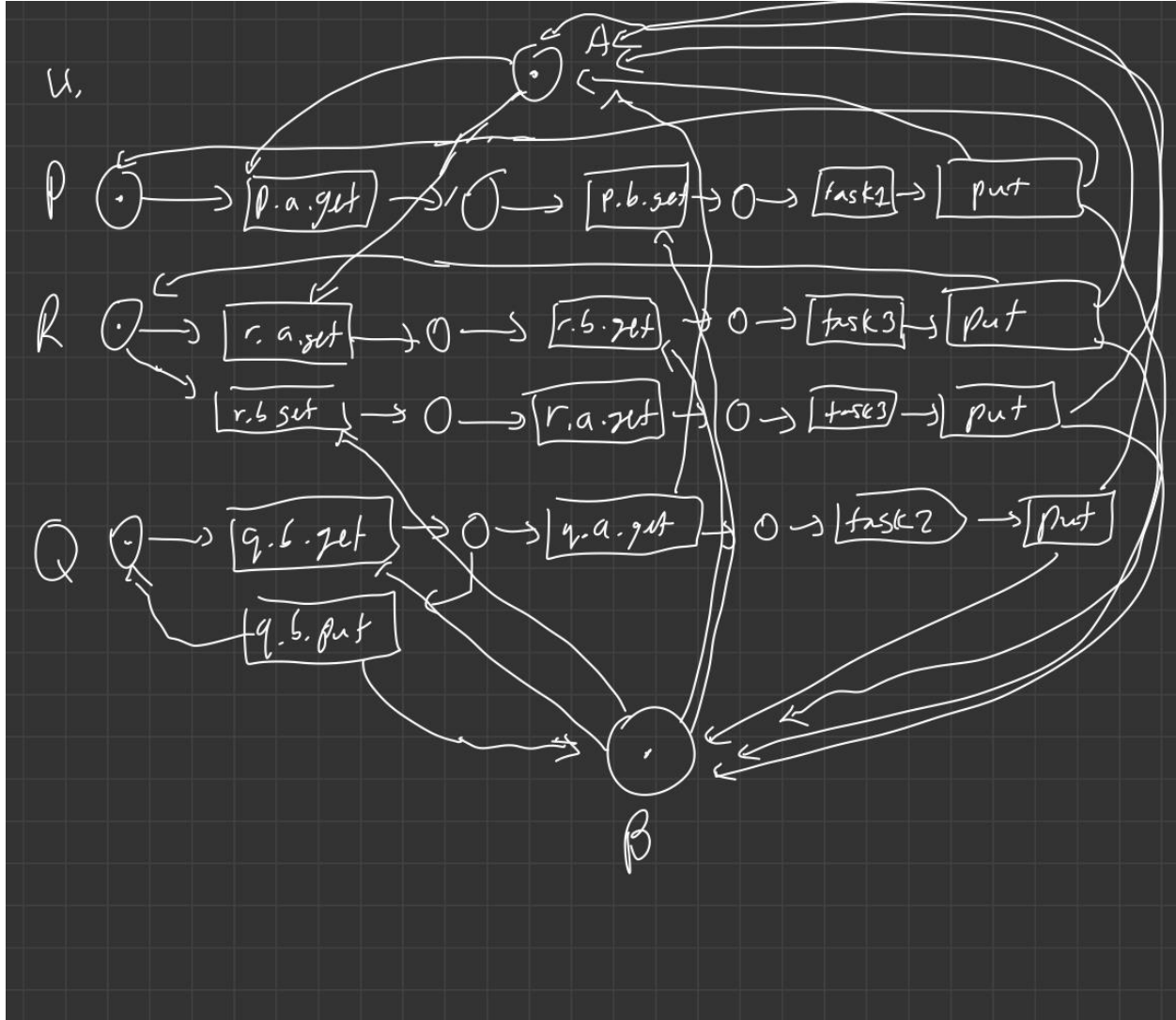


3. $A = (b \rightarrow C \mid c \rightarrow b \rightarrow B)$,

$B = (b \rightarrow B \mid a \rightarrow A),$
 $C = (a \rightarrow b \rightarrow B \mid a \rightarrow B \mid b \rightarrow A).$

4. A) RESOURCE = (get \rightarrow put \rightarrow RESOURCE).
 P = (p.a.get \rightarrow p.b.get \rightarrow task1 \rightarrow p.a.put \rightarrow p.b.put \rightarrow P).
 Q = (q.b.get \rightarrow (q.b.put \rightarrow Q \mid q.a.get \rightarrow task2 \rightarrow q.a.put \rightarrow q.b.put \rightarrow Q)).
 R = (r.a.get \rightarrow r.b.get \rightarrow task3 \rightarrow r.a.put \rightarrow r.b.put \rightarrow R \mid r.b.get \rightarrow r.a.get \rightarrow task3 \rightarrow r.a.put \rightarrow r.b.put \rightarrow R).
 \parallel SYSTEM = (P \parallel Q \parallel R \parallel {p,q,r} :: {a,b} : RESOURCE).

4b)



5. const N = 5
 CLERK = (show_seats \rightarrow CLERK \mid select_seat \rightarrow print_ticket \rightarrow CLERK).
 SEATS = SEATS[N],
 SEATS[i:0..N] = (when (i>0) select_seat \rightarrow SEATS[i-1] \mid when (i==0) select_seat \rightarrow ERROR).
 CLIENT = (select_seat \rightarrow CLIENT).
 \parallel COMPUTER = (CLERK \parallel SEATS \parallel CLIENT).

We set the value N to be the number of seats that are available in the system. The SEATS_AVAILABLE process keeps track of the seats that are still available that the client can select. The CLERK process will show the available seats and the customers can select an available seat. When the seat is selected, the seat will be removed from the SEATS_AVAILABLE process and lower the number of available seats by 1. This will not allow another customer from booking that seat. If there are no more seats available, the SEATS process will hit the ERROR state. Double booking is prevented by removing the seat from the available seats and not showing any taken seats in the show_seats action.

6. A) P is not bisimilar to Q as if we give in the trace "ba", the LTS P would be in state p_2 , while LTS Q would be in either state q_2 or q_3 . The state p_2 allows for the actions b and c while state q_2 allows for only the action b . State q_3 similarly only allows the action c . Thus p_2 is not bisimilar to q_2 and p_2 is not bisimilar to q_3 and thus, P and Q are not bisimilar.

B) To prove that P is bisimilar to S , we have to show that there is a pair of bisimilar states for all states in P and S .

State p_0 and s_0 both only have the action b , and are therefore bisimilar.

State p_1 and s_1 both only have the action a when we follow the trace "b" and are therefore bisimilar.

State p_2 and s_2 both have the actions b and c when we follow the trace "ba" and are therefore bisimilar.

When we enter the trace "bab", this case is the same as p_0 and s_0 which is already bisimilar.

When we enter the trace "bac" the states p_0 and s_3 only have the action b and are therefore bisimilar.

Since it has been shown all cases have bisimilar pairs, P is therefore bisimilar to S .

7. If we trace both the processes P_1 and P_2 , we see that they have the same traces as $\text{Traces}(P_1) = \text{Traces}(P_2) = \text{Prefix}((a(bc \cup cb))^*)$.

If we draw the Petri Nets of $|P_1Q$ and $|P_2Q$, we will see that taking the right transition of "a" for $|P_1Q$, we will enter a deadlock state as there is no token to run the transition "c". Thus, $|P_1Q$ will

7. $P1$

$P2$

$|| P1 Q$

deadlocked if this path taken

$|| P2 Q$

