

# Supplementary Material for “TechHNet: Scalable Temporal Hypergraph Neural Network”

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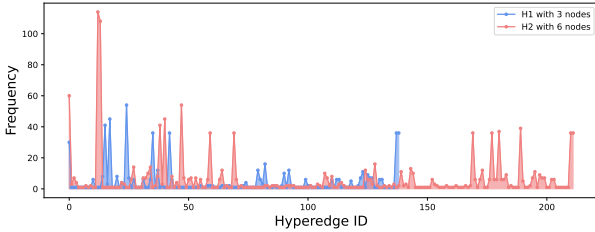
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## 1 PRELIMINARIES AND BACKGROUNDS

### 1.1 One of the Motivation

As discussed in Section 3.3, the existing solutions suffer from the *over-smoothing* problem due to the multi-hop hyperedge-centric neighbors and hyperedge-centric message passing. To illustrate, we take the CATWalk as an example, visualizing the sampled hyperedge neighbors for three randomly selected hyperedges, where we set  $L = 3$  and  $k = 6$ . For clarity, we renumber the IDs of these hyperedge-centric neighbors from 0 to  $K^L \cdot |H_e| - 1$ . We observe that (1) For both  $H_1$  and  $H_2$ , it is common to find repeated neighbors among multi-hop hyperedge neighbors, with the maximum frequency reaching 114. (2) The frequency distribution is long-tailed: a few hyperedges are seen more than 100 times, but most appear fewer than 5 times. These frequently occurring hyperedges may overwhelm the hyperedge embeddings’ expressiveness, resulting in sub-optimal embedding quality and poor generalization capability. The above observation motivates us to develop more effective techniques for hyperedge-centric neighbor extraction and the message passing mechanism.



**Figure 1: The frequency of hyperedge sampled by CATWalk for two incoming hyperedges ( $H_1$  and  $H_2$ ) on the Congress Bills dataset, where  $L$  and  $K$  are set to 3 and 6, respectively. We re-number sampled hyperedges and count them.**

## 2 THEORETICAL RESULTS

### 2.1 Proof of Theorem 1

**THEOREM 1.** *Given a new hyperedge  $(e_i, t)$  with  $m$  nodes at time  $t$ , denoted by  $\mathcal{V}_i = \{v_1, \dots, v_m\}$ , and suppose the hash tables corresponding to these nodes are already full at time  $t$ . For any node  $v_j$  within the hyperedge  $(e_i, t)$ , its newest hyperedge-centric neighbors can be inserted in its hash table  $\mathcal{N}_{v_j}^{t^-}(K)$  with probability  $\Pr \propto$*

*$\exp(\frac{\theta \cdot \lambda}{K}(t - t^-))$ , where  $\lambda$  is a constant,  $\theta$  is the hyperparameter that controls this retention probability, and  $t^-$  is the previous timestamp.*

**PROOF.** Here, we assume that hyperedges come in for any node (e.g.,  $v$ ) by following a Poisson process with a constant intensity (e.g.,  $\lambda$ ). Let  $v_m$  be one of the historical temporal neighbors of  $v_j$ , where  $v_j, v_m \in (e_i, t)$ . According to Section 4.2.2,  $v_m$  is hashed to the position  $\text{hash}(v_m, t)$ , and the replacement probability that any other temporal neighbors are inserted to  $\mathcal{N}_{v_j}^{t^-}(K)$  is  $\theta$ . Since we suppose  $v_m$  is already inserted to  $\mathcal{N}_{v_j}^{t^-}(K)$ , we only need to evaluate the probability that it does not get replaced by another temporal neighbor. We define the event of hyperedges with  $v_j$  arriving since  $t^-$  as  $\Theta_{v_j}$  and we know from the Poisson process that

$$\mathbb{E}[\Theta_{v_j}] = \lambda \cdot (t - t^-). \quad (1)$$

On average, each hyperedge has probability  $\frac{\theta}{K}$  of replacing  $(v_m, t)$ . Then, we define the event of replacement for  $\text{hash}(v_m, t)$  as  $\Theta_{\text{hash}(v_m, t)}$  and we have

$$\mathbb{E}[\Theta_{\text{hash}(v_m, t)}] = \frac{\theta \cdot \lambda \cdot (t - t^-)}{K}. \quad (2)$$

By the property of a Poisson process, the probability that none of the hyperedge-centric neighbors get inserted into the same position from  $t^-$  to  $t$  is  $\exp(-\frac{\theta \cdot \lambda}{K}(t - t^-))$ . Last, we have

$$\Pr((v_m, t) \in \mathcal{N}_{v_j}^{t^-}(K)) \propto \exp(-\frac{\theta \cdot \lambda}{K}(t - t^-)) \quad (3)$$

Overall, the proof is completed.  $\square$

### 2.2 Proof of Lemma 1

**LEMMA 1.** *Hyper2Token is permutation invariant and is a universal approximation of invariant multi-set functions.*

**PROOF.** Let  $\pi(S)$  be a given permutation of set  $S$ , we aim to show that  $\Psi(S) = \Psi(\pi(S))$ . We first recall the Hyper2Token: Let  $S = \{z_1(t), \dots, z_{|H_e|}(t)\}$ , where  $z_i(t) \in \mathbb{R}^{d_n}$ , be the input set and  $Z_e(t)$  be its matrix representation:

$$\Psi(Z_e(t)) = \frac{1}{|H_e|} \cdot \sum_{i=1}^{|H_e|} \text{MLP}(z_i(t)), \quad (4)$$

$$= \text{MEAN}(\text{MLP}(z_1(t)), \dots, \text{MLP}(z_{|H_e|}(t))). \quad (5)$$

Let  $\Psi(\pi(S)) = [z_{\pi(1)}(t), \dots, z_{\pi(|H_e|)}(t)]$  be a permutation of the input matrix  $Z_e(t)$ . Accordingly, we can see the output of the

**Table 1: Dataset Statistic.**

Alias	Dataset	$ \mathcal{V} $	$ \mathcal{E} $	#Timestamps	$\max  H_{\mathcal{E}} $	average length $\overline{ H _e}$	Domain
EE	Email-Enron	143	10,883	10,788	18	2.5	Social network
NC	NDC-Classes	1,161	49,724	5,891	24	3.2	Drug network
NS	NDC-Substances	5,311	112,405	7,734	25	2.6	Drug network
UT	Users-Threads	125,602	192,947	189,917	14	2.2	Social network
CB	Congress-Bills	1,718	260,851	5,936	25	4.3	Bill network
TMS	Threads-Math-Sx	176,445	719,792	718,340	21	2.5	Social network
CD	Coauth-DBLP	1,924,991	3,700,067	83	25	3.0	Scholar network
TSO	Threads-Stack-Overflow	2,675,955	11,305,343	11,260,218	25	3.5	Social network

Hyper2Token is permuted by  $\pi(\cdot)$ :

$$\Psi(\pi(Z_e(t))) = \text{MEAN}(\text{MLP}(z_{\pi(1)}(t)), \dots, \text{MLP}(z_{\pi(|H_e|)}(t))), \quad (6)$$

$$= \text{MEAN}(\pi(\text{MLP}(z_1(t)), \dots, \text{MLP}(z_{|H_e|}(t)))), \quad (7)$$

$$= \Psi(Z_e(t)). \quad (8)$$

In the last step, we leverage the property that  $\text{MEAN}(\cdot)$  is permutation invariant. Based on Eqs. (6)-(8), we can see that Hyper2Token possesses permutation invariance. Given that Hyper2Token employs a two-layer MLP, it functions as a universal approximator capable of modeling any function. Therefore, Hyper2Token serves as a universal approximator.  $\square$

### 3 MODEL TRAINING

During the training phase, for each hyperedge in the training set, we adopt the commonly used negative sample generation method [9] to generate a negative sample. Next, for each hyperedge in the training set such as  $(e, t) = \{v_1, v_2, \dots, v_m\}$ , including both positive and negative samples. During the batch processing, we do the node-wise and hyperedge-wise message passing to generate the hyperedge embeddings. For the hyperedge prediction, we use a 2-layer perception over hyperedge embeddings to generate a predicted score, which can be plugged into the cross-entropy loss for training or compared with a threshold to make the final prediction.

## 4 EXPERIMENTAL SETTING

### 4.1 Datasets

We employ 8 real-world datasets for model evaluation. We provide a more detailed introduction in Table 1, including the averaged hyperedge length  $\overline{|H_e|}$  and application domains.

### 4.2 Baselines

- NHP [13] is a classical neural hyperlink prediction method, which is an enhanced version of the self-attention-based graph convolutional network for hypergraphs (HyperSAGC N) [14]. We use the time encoding function to encode timestamps and regard them as the hyperedge features. The source code is provided at [1].
- CHESHIRE [8] treats a hyperedge as a fully connected graph (clique) and uses a Chebyshev spectral graph convolution neural network to refine the embeddings of the

nodes within the hyperedge, which is a hyperedge prediction method. We also use the time encoding function to encode timestamps and regard them as the hyperedge features. The source code is provided at [4].

- CE-CAW: We apply CAW [12] on the CE of the temporal hypergraph. CAW is a temporal edge prediction method that uses causal anonymous random walks to capture the dynamic laws of the network in an inductive manner. The source code is provided at [2].
- CE-Zebra: We apply Zebra [11] on the CE of the temporal hypergraph. Zebra accelerates the computation of T-GNN by directly aggregating the features of its neighbors returned by the top- $k$  temporal Personalized PageRank (T-PPR). The source code is provided at [6].
- CE-Orca: We apply Orca [10] on the CE of the temporal hypergraph. Orca proposes a caching-based framework to address the neighborhood explosion problem by non-trivially caching and reusing intermediate embeddings. The source code is provided at [5].
- CATWalk [7] directly models the temporal and high-order structures to generate hyperedge embeddings. Concretely, it extracts hyperedge-centric neighbors by designing a set-based sampling method and encodes sampled hyperedges and nodes using the permutation invariant message passing. The source code is provided at [3].

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