

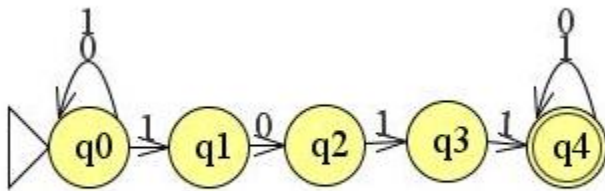
Lucy Wilcox

Assignment for Day 6

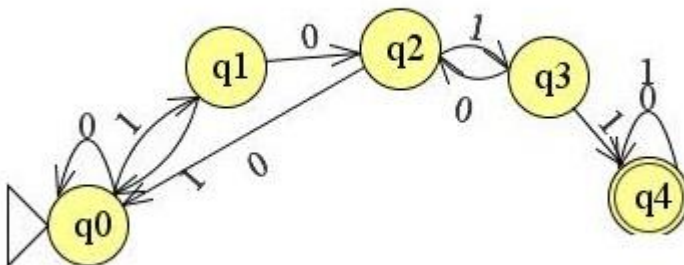
I used Sipser's book to assist me with the homework and the JFLAP program.

## Part 1

### Problem 1



### Problem 2



2) Transition table

state	0	1
$q_0$	$q_0$	$q_1$
$q_1$	$q_2$	$q_0$
$q_2$	$q_0$	$q_3$
$q_3$	$q_2$	$q_4$
$q_4$	$q_4$	$q_4$

### Problem 3

I do not think that a large alphabet is more powerful than a short one.

An alphabet  $L = \{a, b, c, d\}$  contains no more information than an alphabet  $N = \{0, 1\}$ .

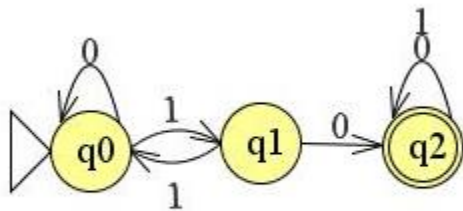
From  $N$ , we can form the alphabet  $N^* = \{\text{empty}, 0, 1, 00, 01, 10, 11, 000, \dots\}$  each element of  $N^*$ .

Regular languages are closed under the star operation so elements in  $N^*$  are part of  $N$ .  $N^*$  also contains infinite entries, therefore it contains enough information that it can compute at least the same amount as the four entries  $L$  which has. They could for example be represented as  $\{00, 01, 10, 11\}$ . This concept is used in the binary alphabet which encodes letters as numbers.

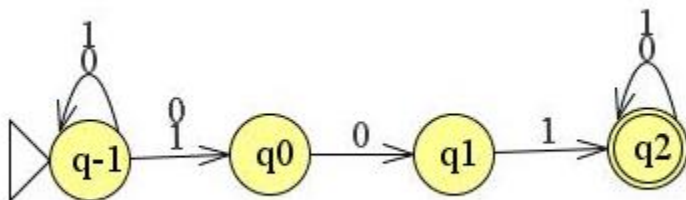
### Problem 4

Here's my attempt:

This is A which looks for "10". Any DFA looking a string can be represented as a forward chain of transitions which are the values we are looking for, in this case  $q_0 \rightarrow q_1$  if 1,  $q_1 \rightarrow q_2$  if 0, so 1 then 0. It will also have other arrows going to previous states,  $q$ 's with a smaller index than themselves. And on the final state, has 0 and 1 returning to itself.



A NFA can be made which looks for the reverse by adding a new starting node which loops 0, 1 back to itself, and to the  $q_0$  of the DFA, A. Remove all arrows which go to  $q$ 's with a smaller index than themselves, and perform a not operation on all other transition values (the 0s and 1s).



### Part 2

#### Problem 5

Done

**Problem 6**

$0 \cdot 10^*$