

# Sorting Basics

Unsorted Array

9	1	3	2	7	4
---	---	---	---	---	---



sorting algorithm

Sorted Array

1	2	3	4	7	9
---	---	---	---	---	---

# Lecture Flow

- 1) Pre-requisites
- 2) Problem Definitions and Applications
- 3) Different approaches
- 4) Bubble sort
- 5) Selection sort
- 6) Insertion sort
- 7) Checkpoint
- 8) Counting sort
- 9) Practice questions
- 10) Resources

# Pre-requisites

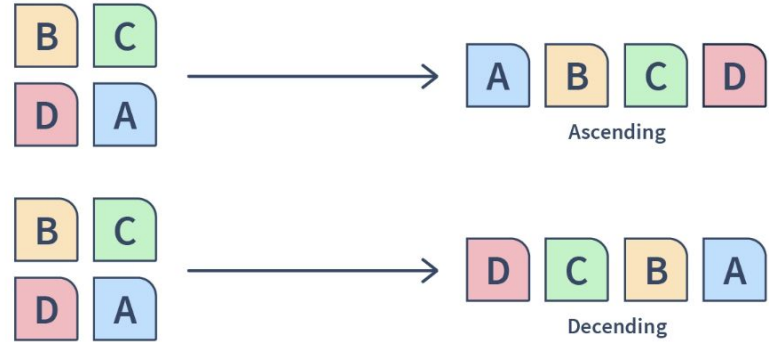
- Basic python
- Asymptotic Analysis
- Arrays



# Problem Definition

# What is Sorting?

- Arranging a set of data in some logical order.
- **Increasing** or **Decreasing** manner.
- Helps finding **largest**, **smallest**, **median** and **nth** value, or **group** items by quality.

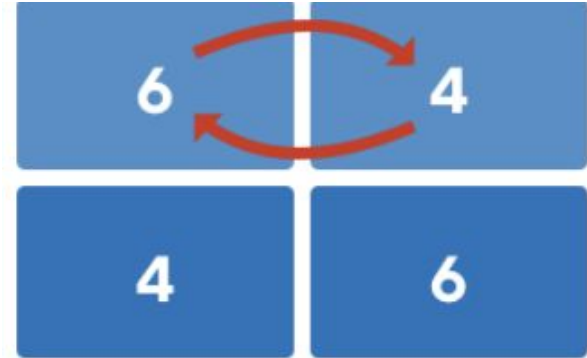


# Sorting by Element Comparison

## Sorting by Element Comparison

A data item is **compared** with other items in the list of items in order to find its place in the sorted list.

*Ex: Bubble, Selection, Insertion ...*



Q: How else can we sort?

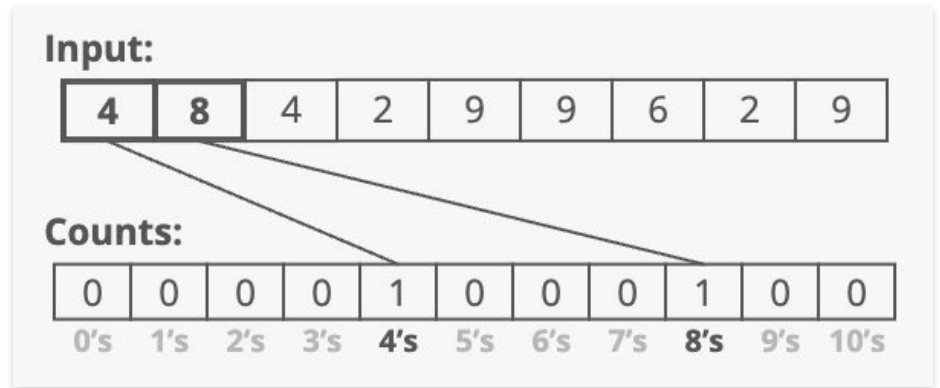


# Sorting by Distribution

## Sorting by Distribution

All items under sorting are distributed over a storage space and then grouped together to get the sorted list.

*Ex: Counting Sort, Bucket Sort ...*

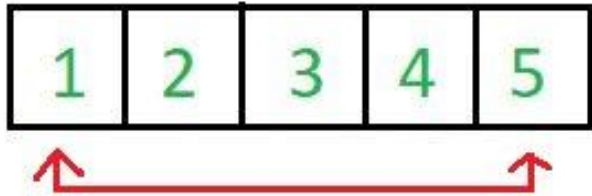




# In-Place vs Out-of-Place

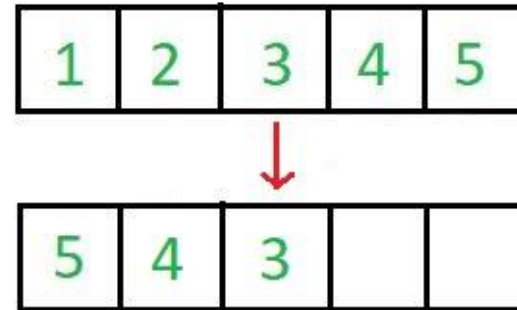
## In-Place

Uses constant space by modifying the order of the elements within the list.



## Out-of-Place

Uses extra space to modify order of elements.



# In-Place vs Out-of-Place

**Input:**

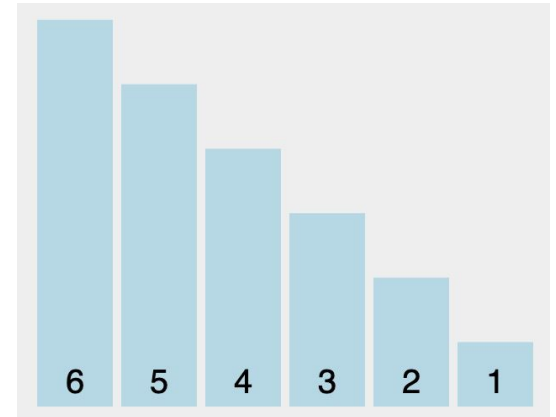
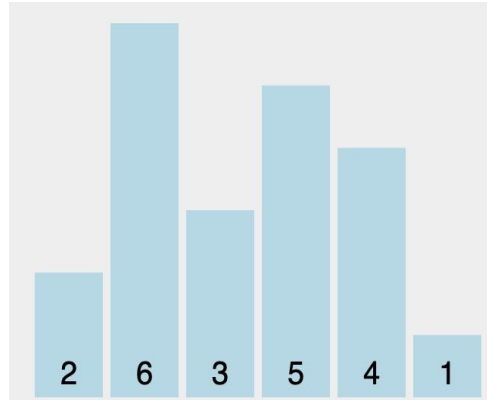
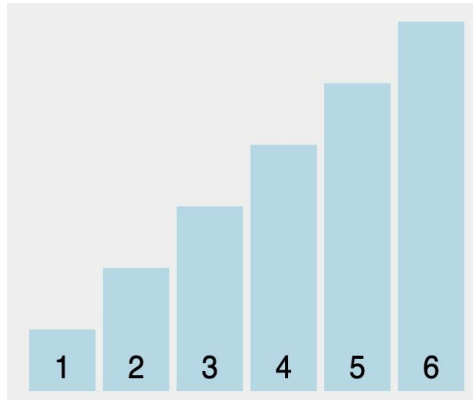
4	8	4	2	9	9	6	2	9
---	---	---	---	---	---	---	---	---

**Counts:**

0	0	0	0	1	0	0	0	1	0	0
0's	1's	2's	3's	4's	5's	6's	7's	8's	9's	10's

# Efficiency of Sorting Algorithm

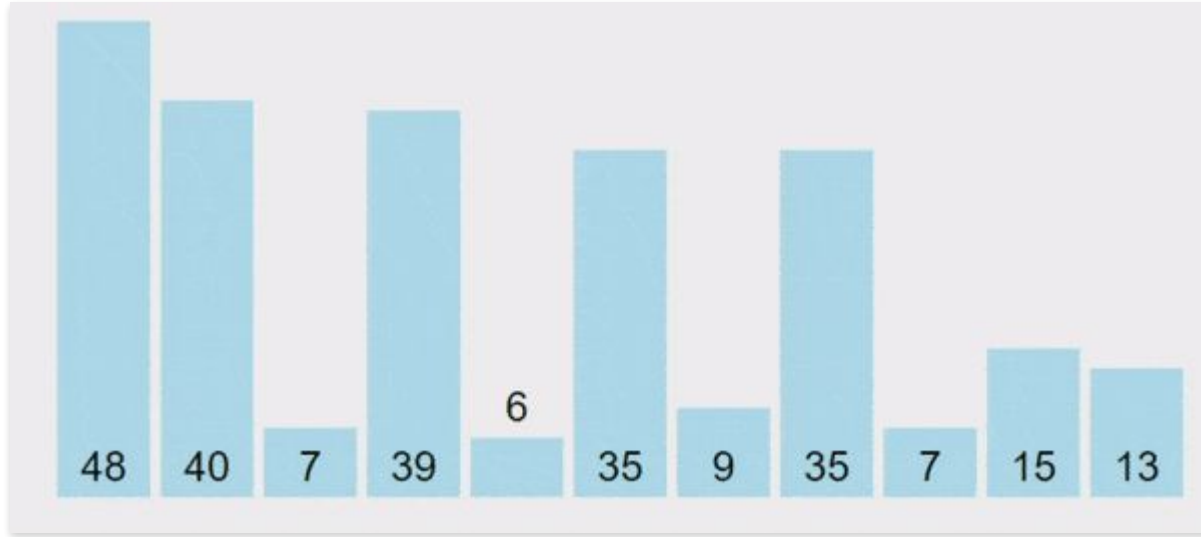
- The complexity of a sorting algorithm measures the running time of a function in which **n** number of items are to be sorted.
- Various sorting algorithms are analyzed in the cases like - **Best case, Worst case** or **Average case**.



# Comparison Sorting

# 1. Bubble Sort

# Bubble Sort

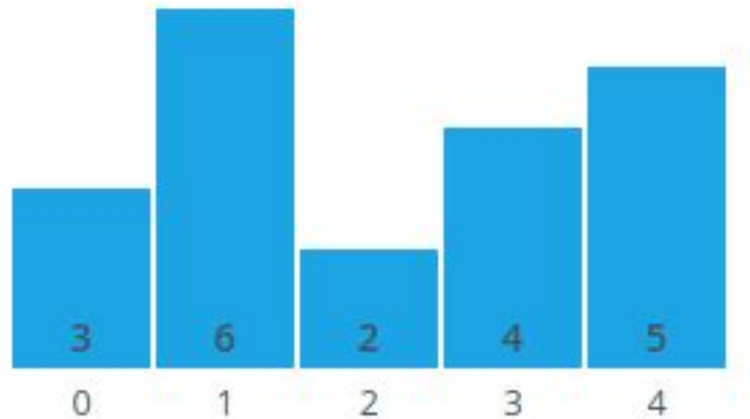


Bubble Sort works by repeatedly **swapping** the adjacent elements if they are in the **wrong order**. Keep doing this until sorted.

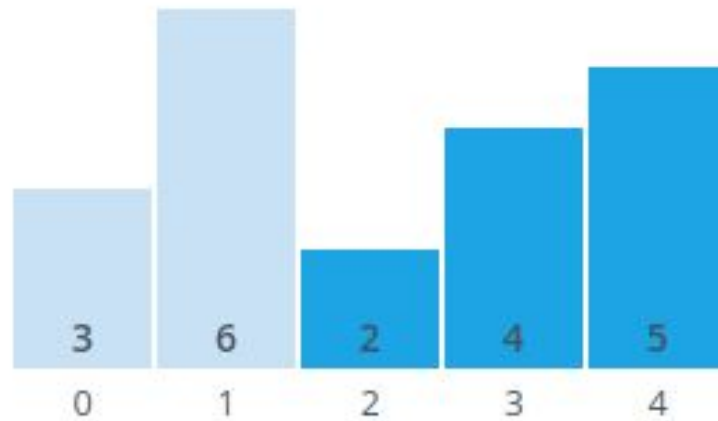
visualization from: [VisuAlgo](#)

# Bubble Sort Simulation

For each pass, we will move left to right swapping adjacent elements as needed. Each pass moves the next largest element into its final position (these will be shown in green).



# Compare The Elements





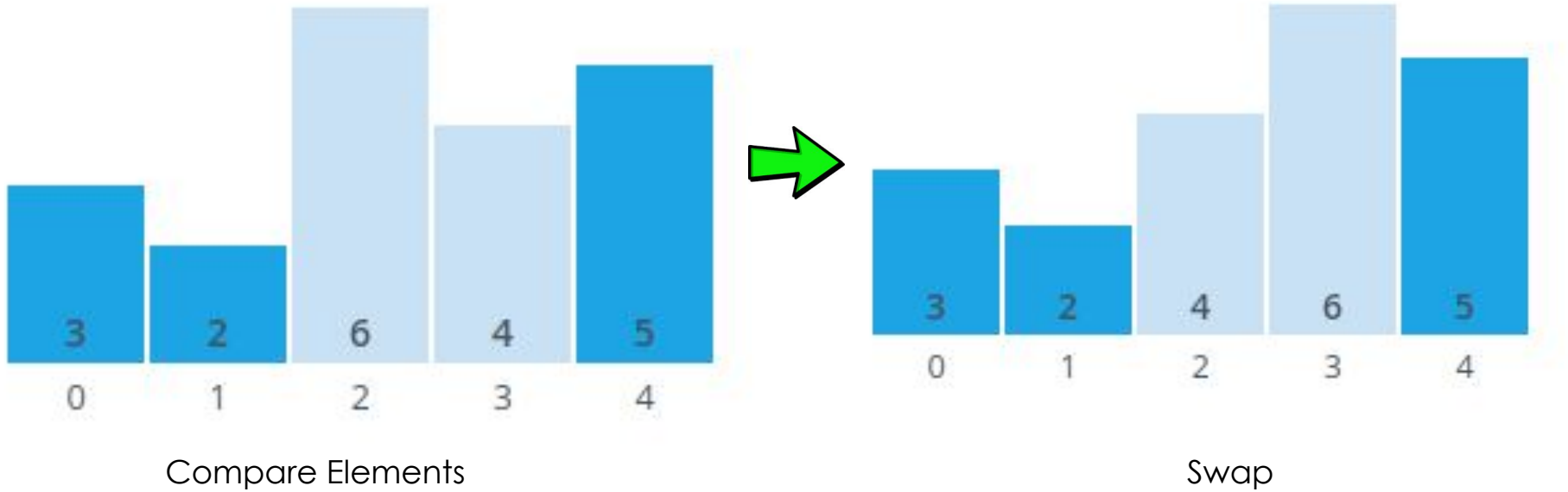
# Compare The Elements and Swap



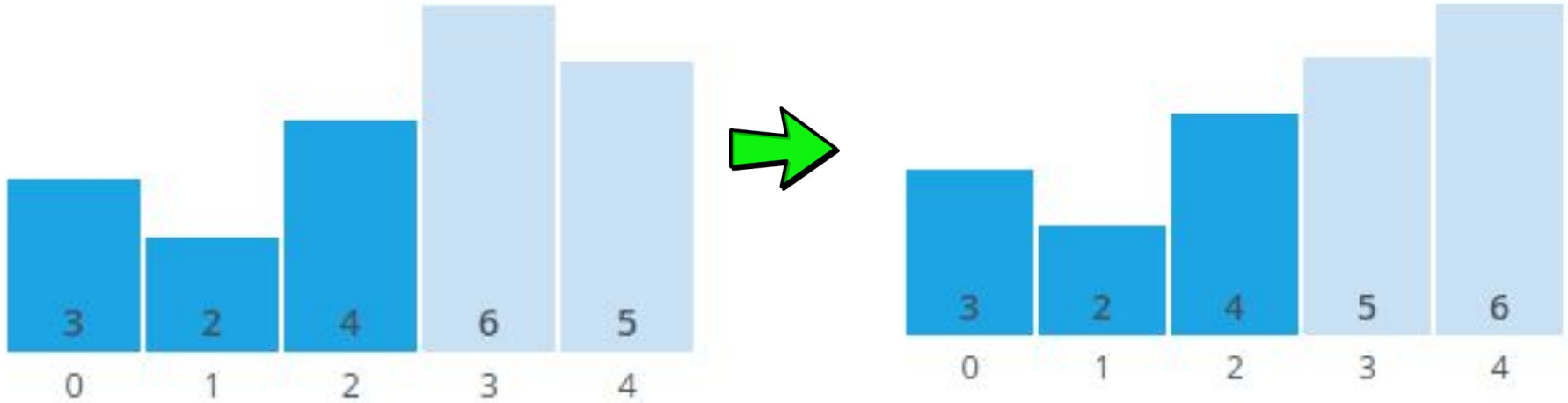
Compare Elements

Swap

# Compare The Elements and Swap



# Compare The Elements and Swap

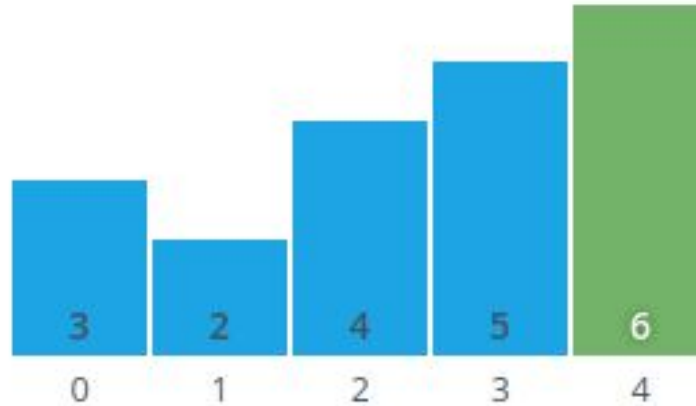


Compare Elements

Swap

# First Pass Done

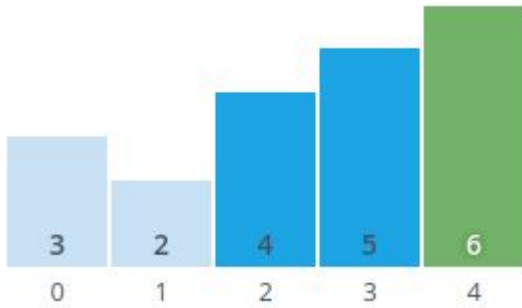
The last element processed is now in its final position.  
Now we will start the next pass for each element moving through the list



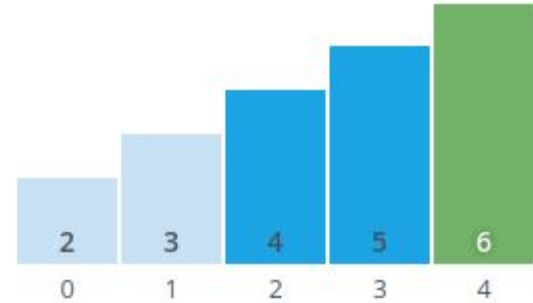
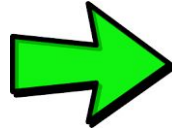
Q: What happens with the second pass?



# Compare The Elements and Swap

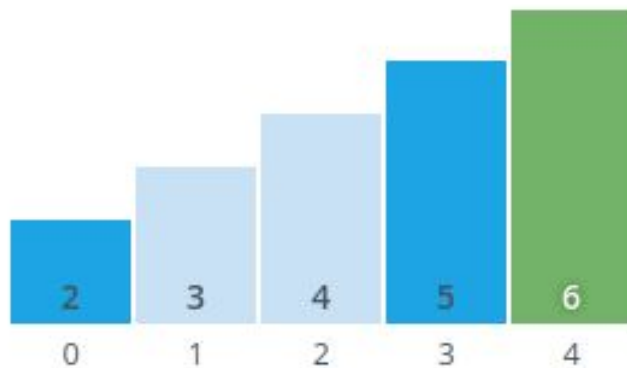


Compare the elements



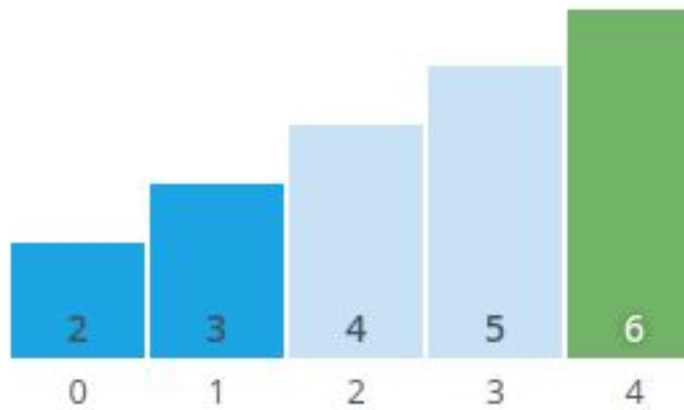
Swap

# Compare The Elements



Compare elements

# Compare The Elements

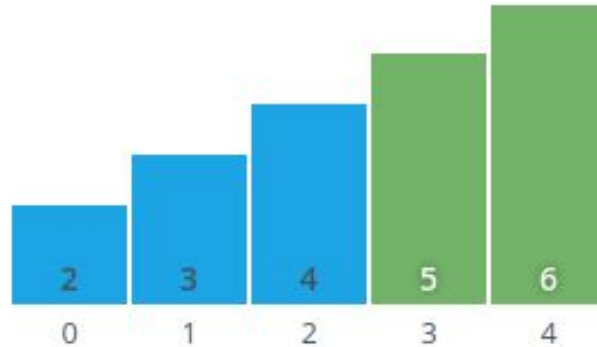


Compare elements

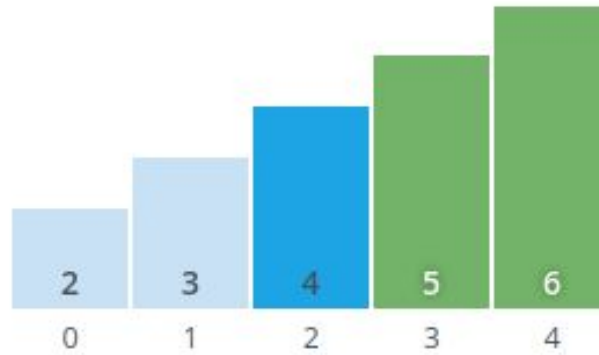


# Done!

The last element processed is now in its final position. Now we will start the next move For each element moving through the list

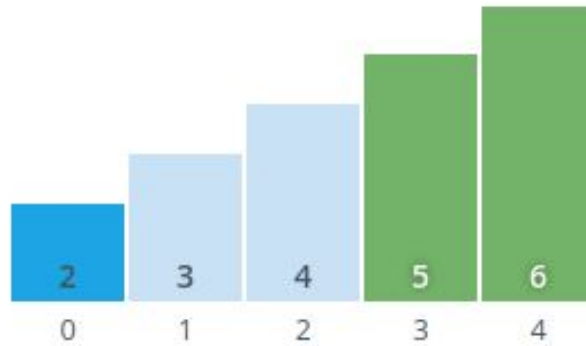


# Compare The Elements



Compare elements

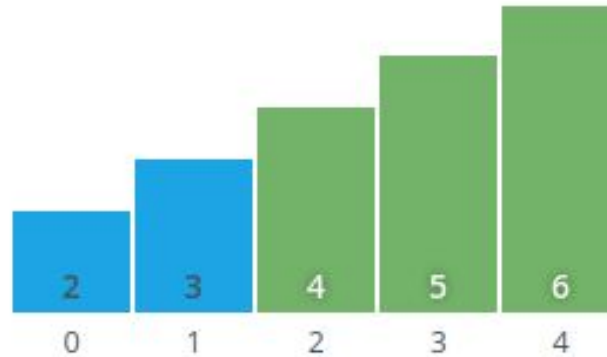
# Compare The Elements



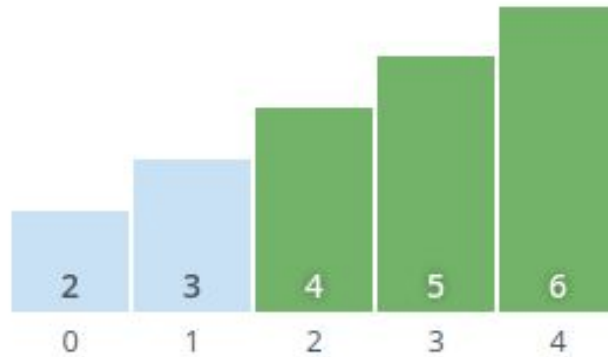
Compare elements

# Done!

The last element processed is now in its final position. We will start the next path For each element moving through the list



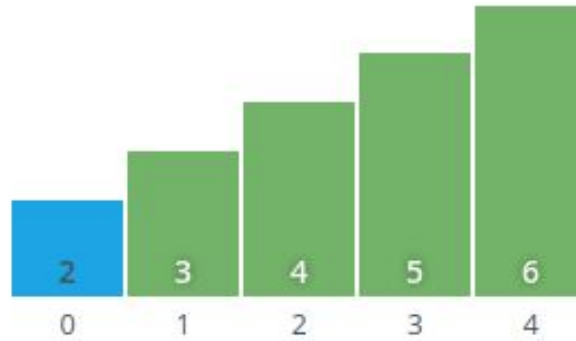
# Compare The Elements



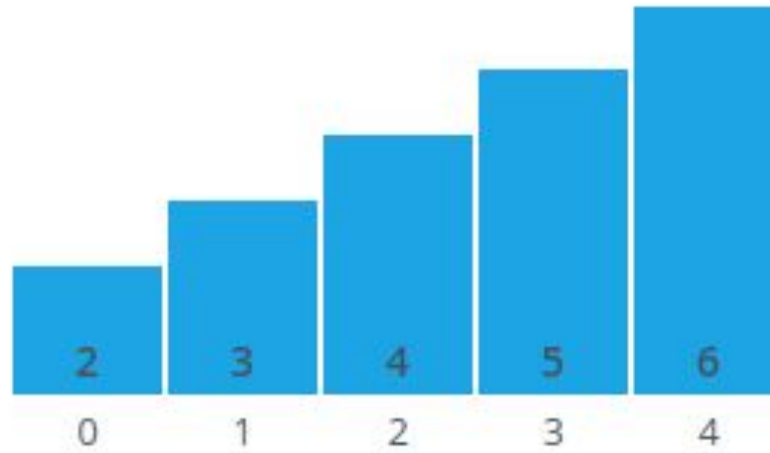
Compare elements

# Done!

The last element processed is now in its final position.



## Final Sorted Output



# Time & Space Complexity

Worst case ? \_\_\_\_\_

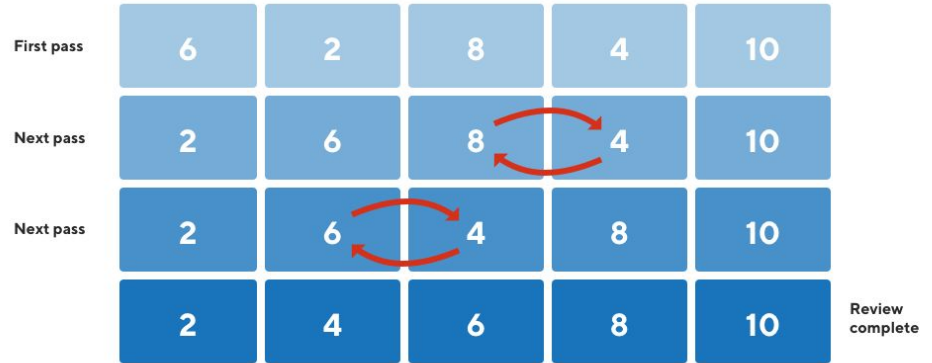
Best case ? \_\_\_\_\_

Average case ? \_\_\_\_\_

Stable ? \_\_\_\_\_

In Place? \_\_\_\_\_

## Bubble Sort





# Time & Space Complexity

Time complexity:  $O(n^2)$

Space complexity:  $O(1)$

Worst case  $O(n^2)$

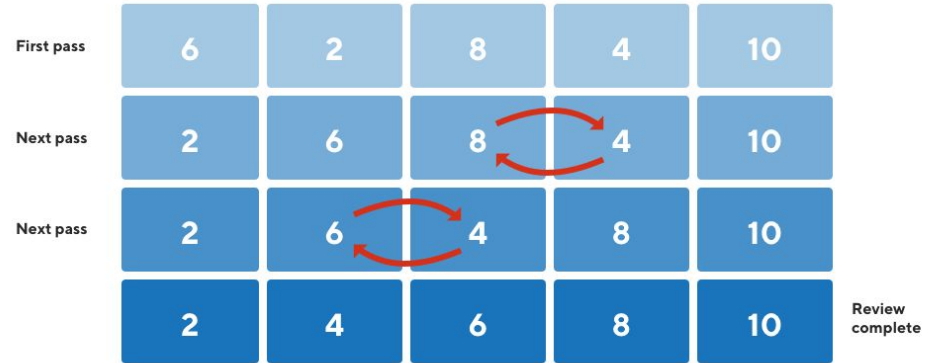
Best case  $O(n^2)$

Average case  $O(n^2)$

Stable ? Yes

In Place? Yes

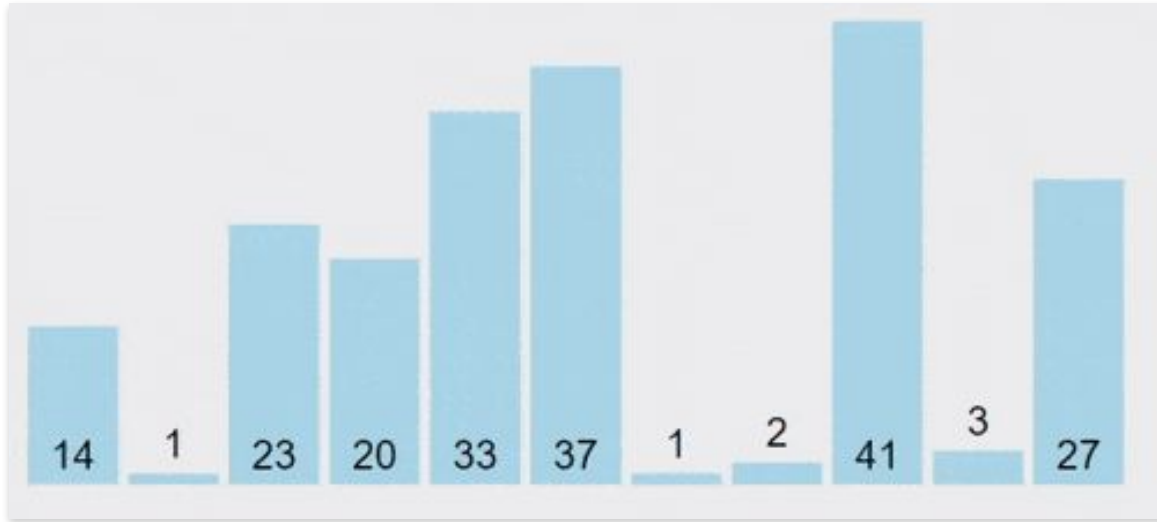
## Bubble Sort



Solve by Using Bubble Sort

## 2. Selection Sort

# Selection Sort



Selection sort selects the **smallest element** from an unsorted list in each iteration and places that element **at the beginning** of the unsorted list.

visualization from: [VisuAlgo](#)

**Q:** What if we selected the largest one first and place it at the end?



Q: What happens when we put the smallest at the end?



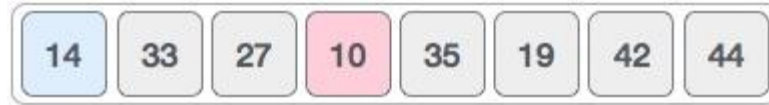
# Selection Sort Simulation

Consider the following array as an example.



# Selection Sort

**For the first position** in the sorted list, **the whole list is scanned sequentially**. The first position where 14 is stored presently, we search the whole list and find that 10 is the lowest value.





# Selection Sort

We then replace 14 with 10. After one iteration 10, which happens to be the minimum value in the list, appears in the first position of the sorted list.

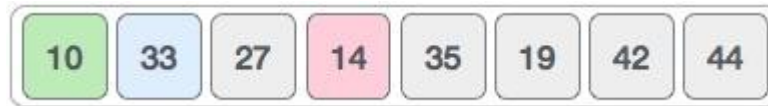


# Selection Sort

For the second position, where 33 is residing, we start scanning the rest of the list in a linear manner.



We find that 14 is the second lowest value in the list and it should appear at the second place. We swap these values.



# Selection Sort

After two iterations, two least values are positioned at the beginning in a sorted manner.

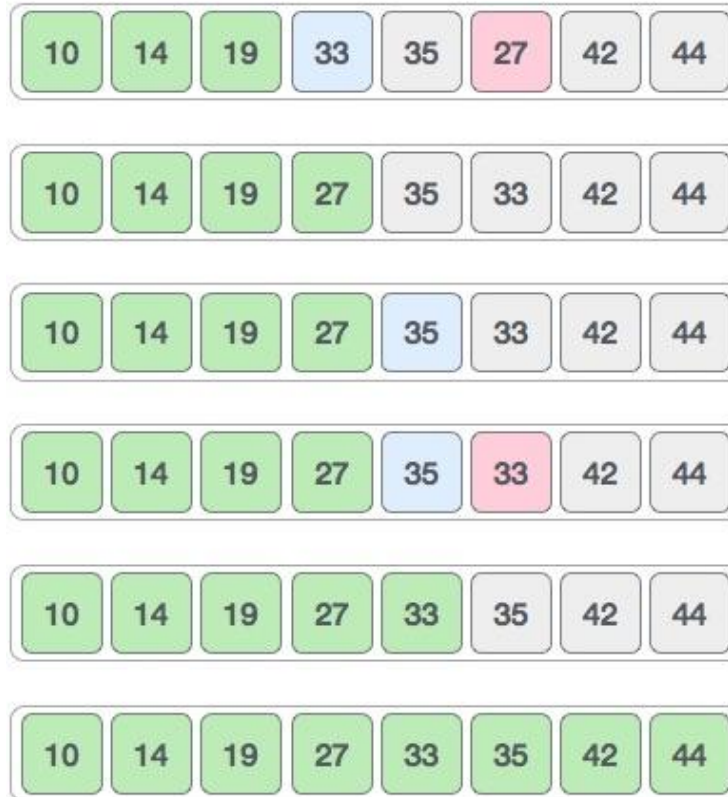


The same process is applied to the rest of the items in the array.



# Final Sorted Output

The same process is applied to the rest of the items in the array.



# Algorithm

- Set MIN\_INDEX to location 0
- Repeat until list is sorted
  - Search the minimum element in the list
  - Swap with value at location MIN\_INDEX
  - Increment MIN\_INDEX to point to next element

**Note:** Every selection requires a search through the input list.

# Time & Space Complexity

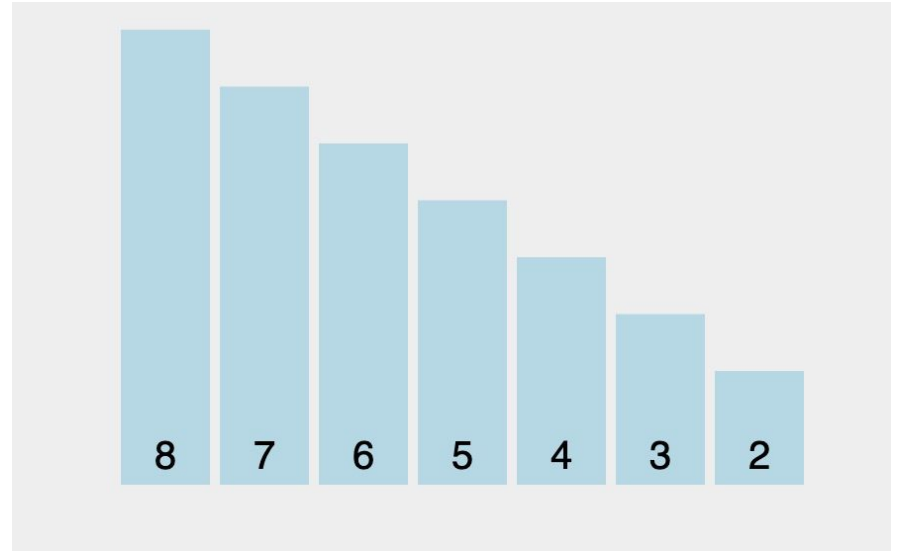
Worst case ? \_\_\_\_\_

Best case ? \_\_\_\_\_

Average case ? \_\_\_\_\_

Stable ? \_\_\_\_\_

In Place? \_\_\_\_\_



# Time & Space Complexity

Time complexity:  $O(n^2)$

Space complexity:  $O(1)$

Worst case  $O(n^2)$

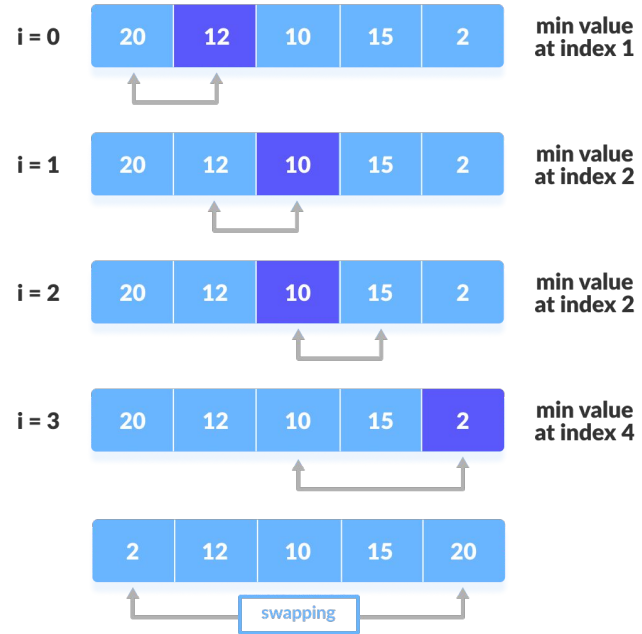
Best case  $O(n^2)$

Average case  $O(n^2)$

Stable ? No

In Place? Yes

step = 0

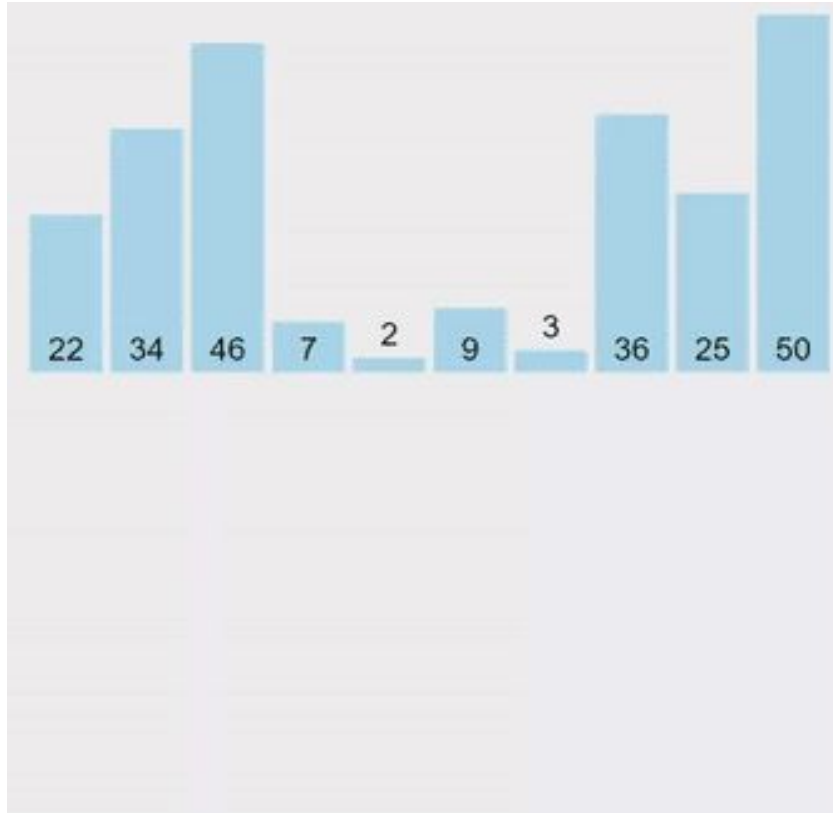


Solve by Using Selection Sort



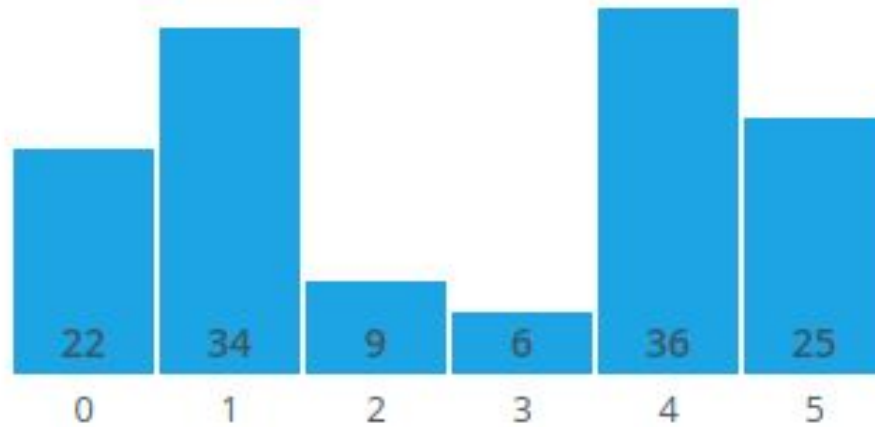
### 3. Insertion Sort

# Insertion Sort



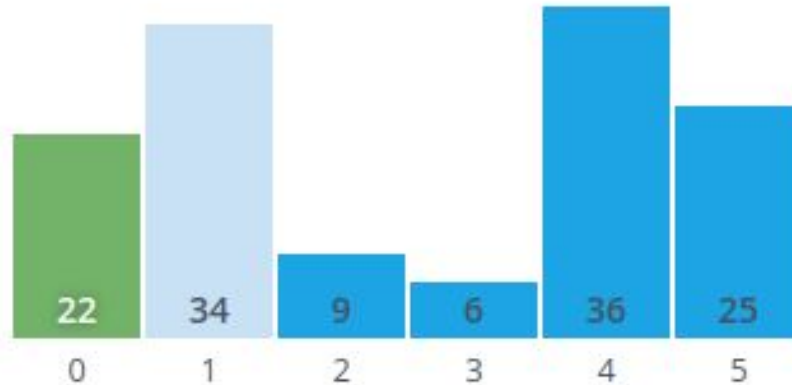
Let's sort the beginning of the list, and insert new elements to the sorted part one by one.

# Insertion Sort Simulation



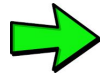
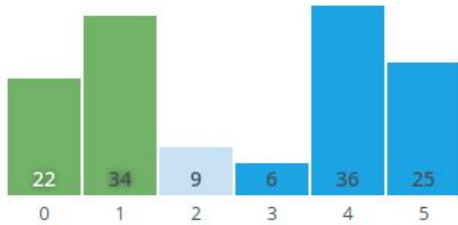
# Insertion Sort Simulation

Green records to the left are always sorted.  
We begin with the record in position 0 in the sorted portion, and we will be moving the record in position 1 (in blue) to the left until it is sorted.

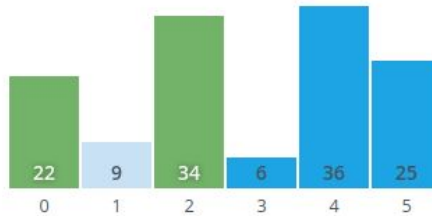


# Insertion Sort Simulation

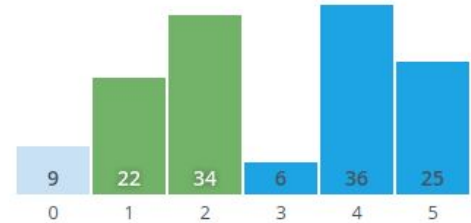
Move the blue record to the left until it reaches the correct position.



Swap

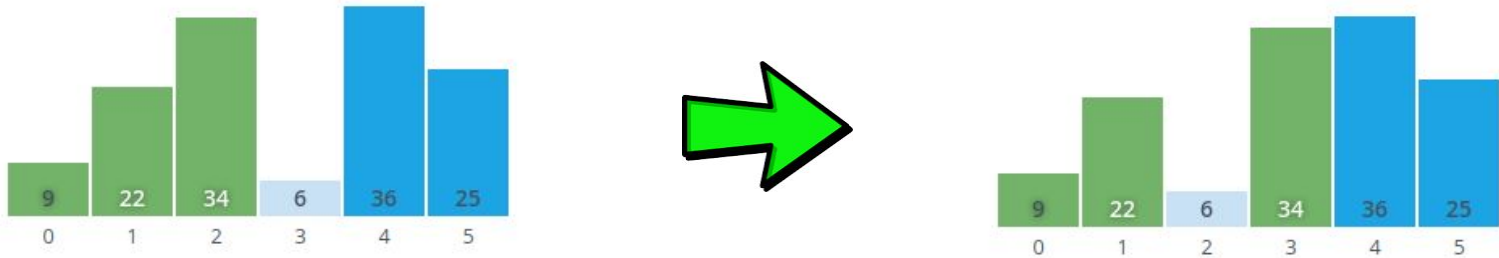


Swap



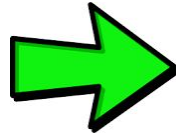
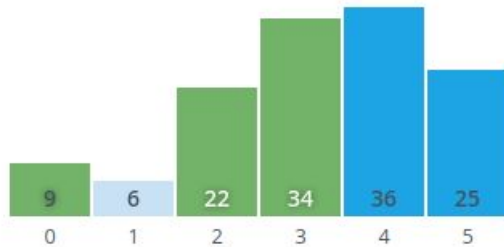
# Insertion Sort Simulation

Move the blue record to the left until it reaches the correct position.

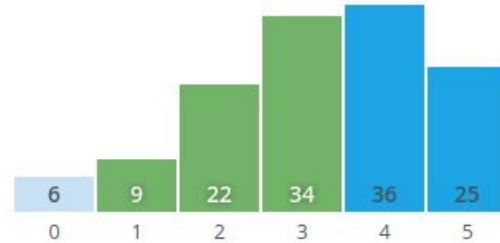


# Insertion Sort Simulation

Move the blue record to the left until it reaches the correct position.

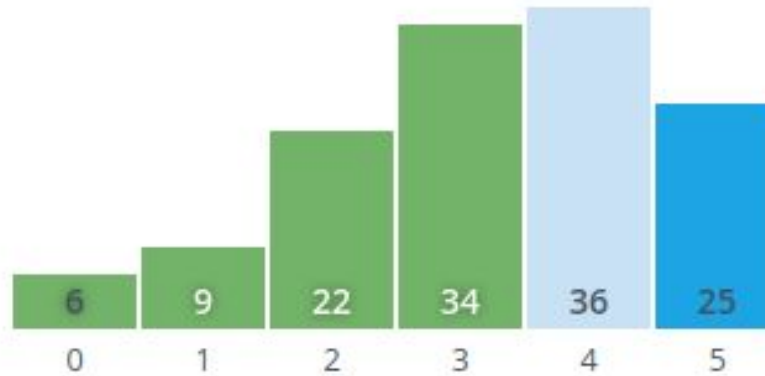


Swap



# Insertion Sort Simulation

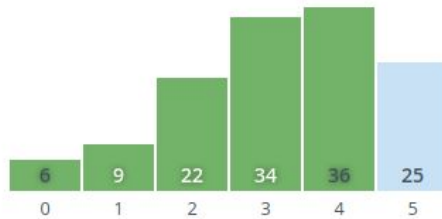
Move the blue record to the left until it reaches the correct position.



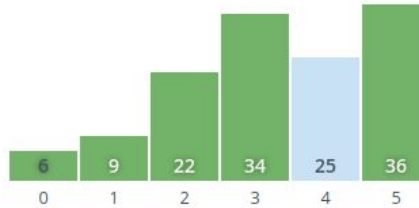


# Insertion Sort Simulation

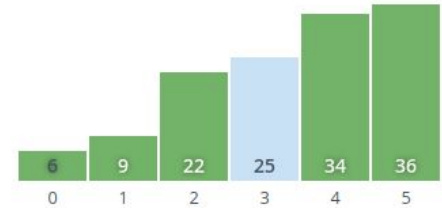
Move the blue record to the left until it reaches the correct position.



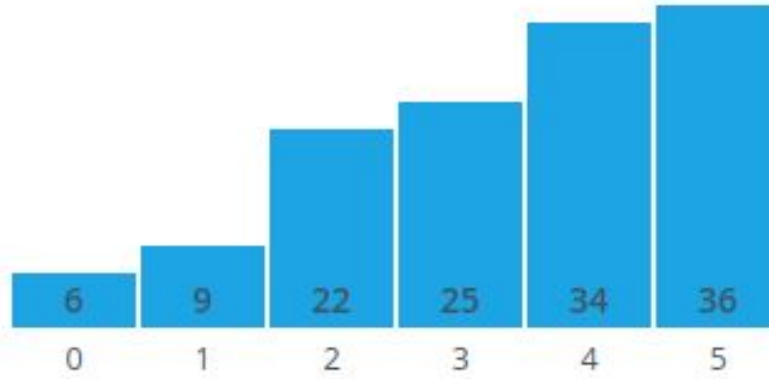
Swap



Swap



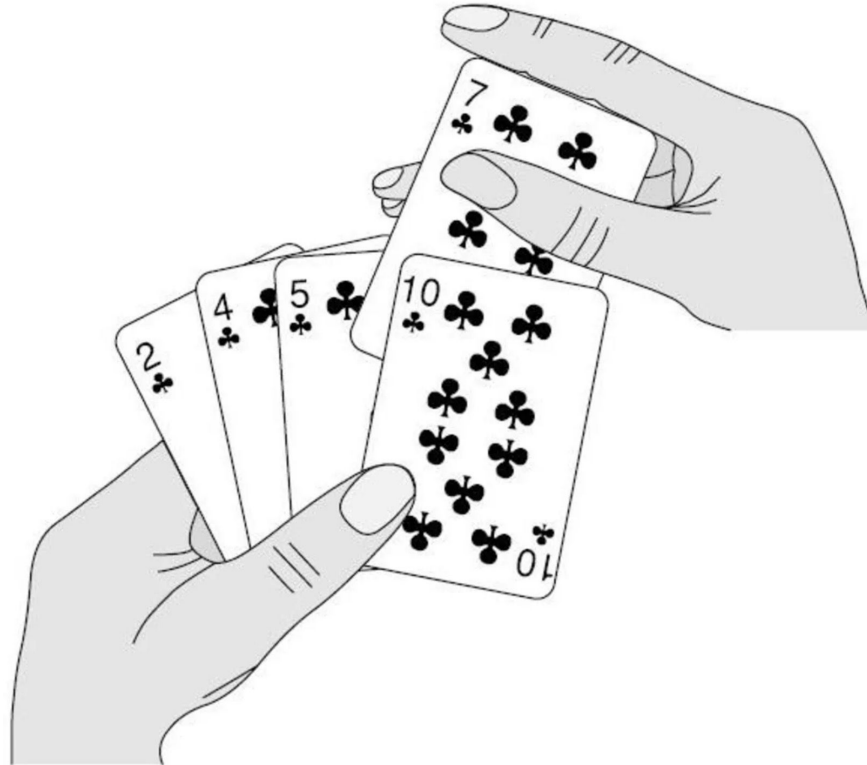
## Final Sorted Output



**Q:** What would be a real life example of Insertion sort?



# Insertion Sort



# Algorithm

- If the element is the first one, it is already sorted.
- Move to next element
- Compare the current element with all elements in the sorted array
- Shift all the the elements in sorted sub-list that is greater than the value to be sorted.
- Insert the value at the correct position
- Repeat until the complete list is sorted

# Time & Space Complexity

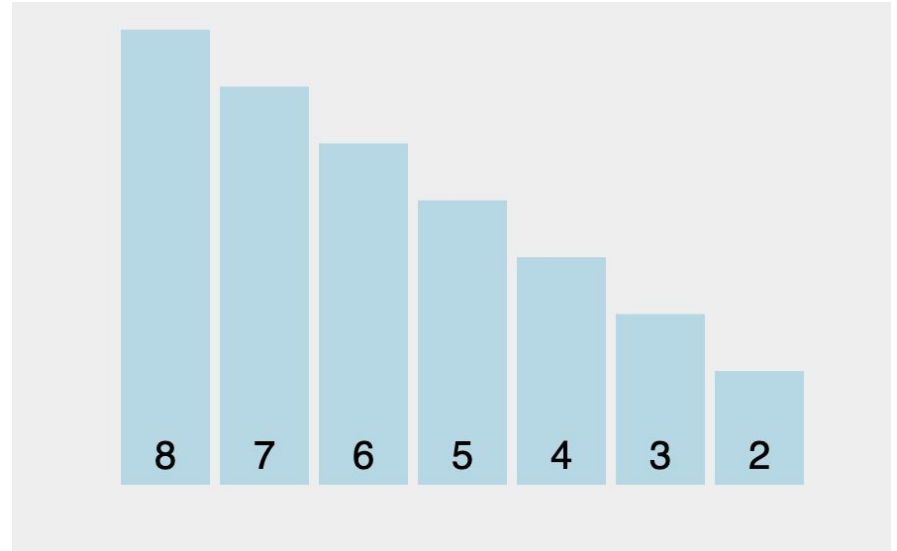
Worst case ? \_\_\_\_\_

Best case ? \_\_\_\_\_

Average case ? \_\_\_\_\_

Stable ? \_\_\_\_\_

In Place? \_\_\_\_\_



# Time & Space Complexity

Time complexity:  **$O(n^2)$**

Space complexity:  **$O(1)$**

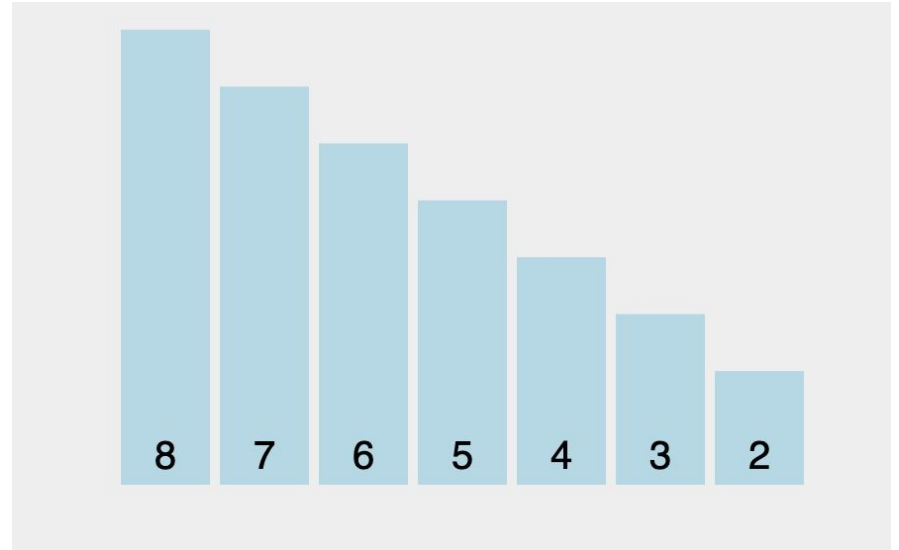
Worst case       **$O(n^2)$**

Best case       **$O(n)$**

Average case       **$O(n^2)$**

Stable ?      Yes

In Place?      Yes



Solve by Using Insertion Sort



# Distribution Sorting



# 1. Counting Sort



# Definition



If the range of the numbers is small enough that can fit in memory;

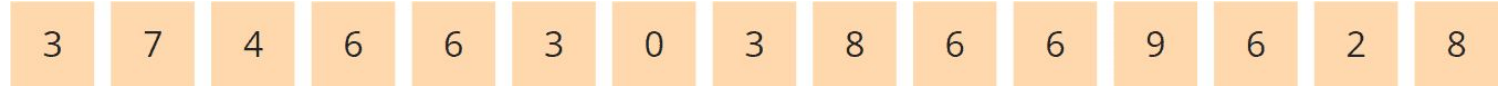
Count the occurrence of items then generate output based on counts in counter array.

# Illustration

## Phase 1: **Counting**

First, a **storage array** is created whose length corresponds to the number range. Then you iterate once over the elements to be sorted, and, for each element, you **increment the value** in the array at the position corresponding to the element.

Consider the following array as an example.

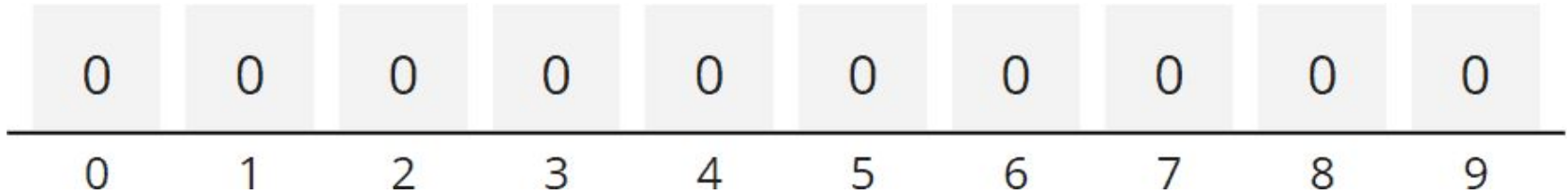


Find the maximum number

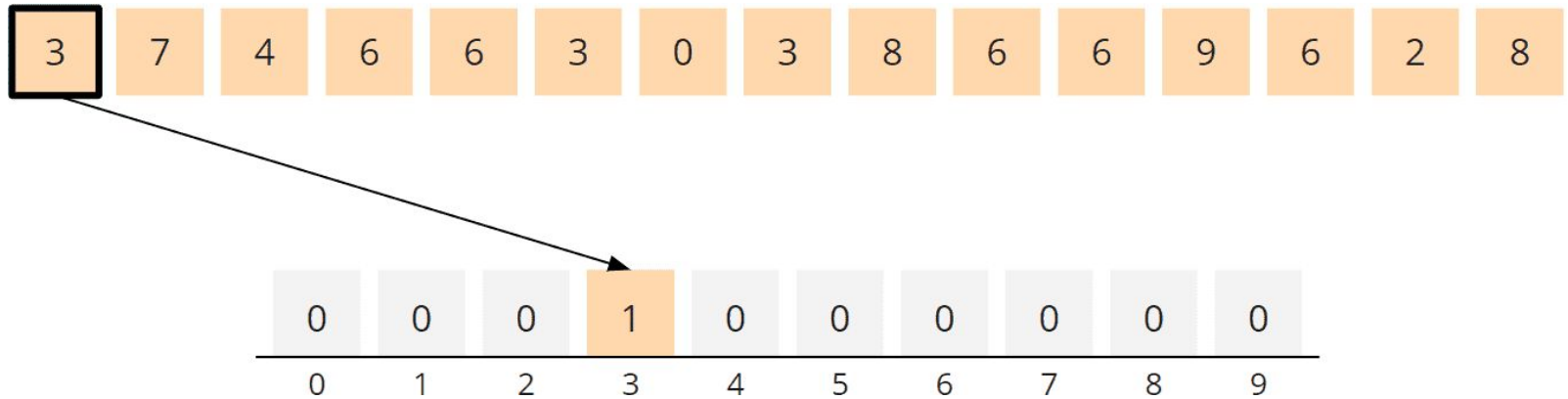


max

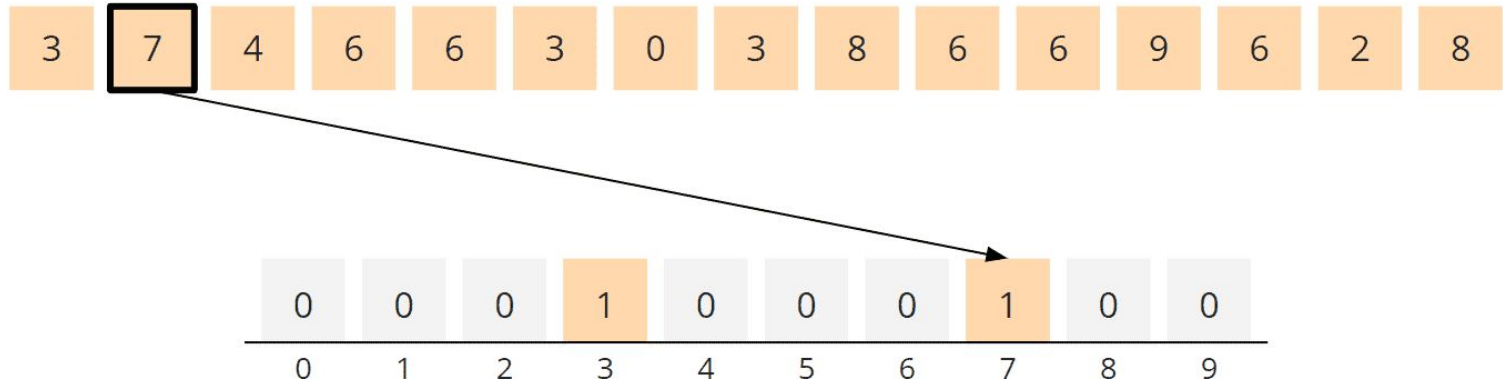
We create an additional array with the length of the maximum number +1. In our case we create an array with length of 10, initialized with zeros.



Now we iterate over the array to be sorted. The first element is a 3 – accordingly, we increase the value in the auxiliary array at position 3 by one:

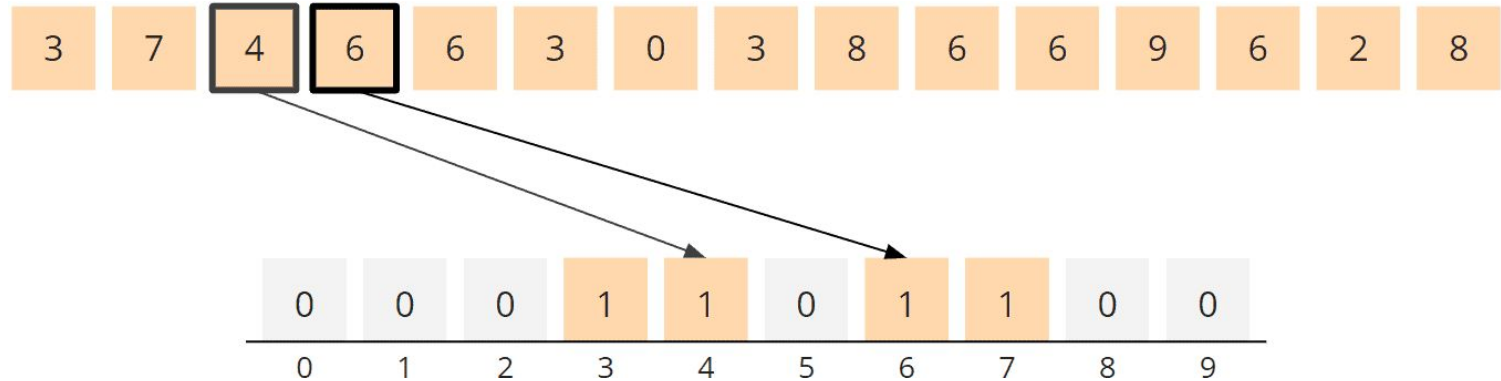


The second element is a 7. We increment the field at position 7 in the helper array

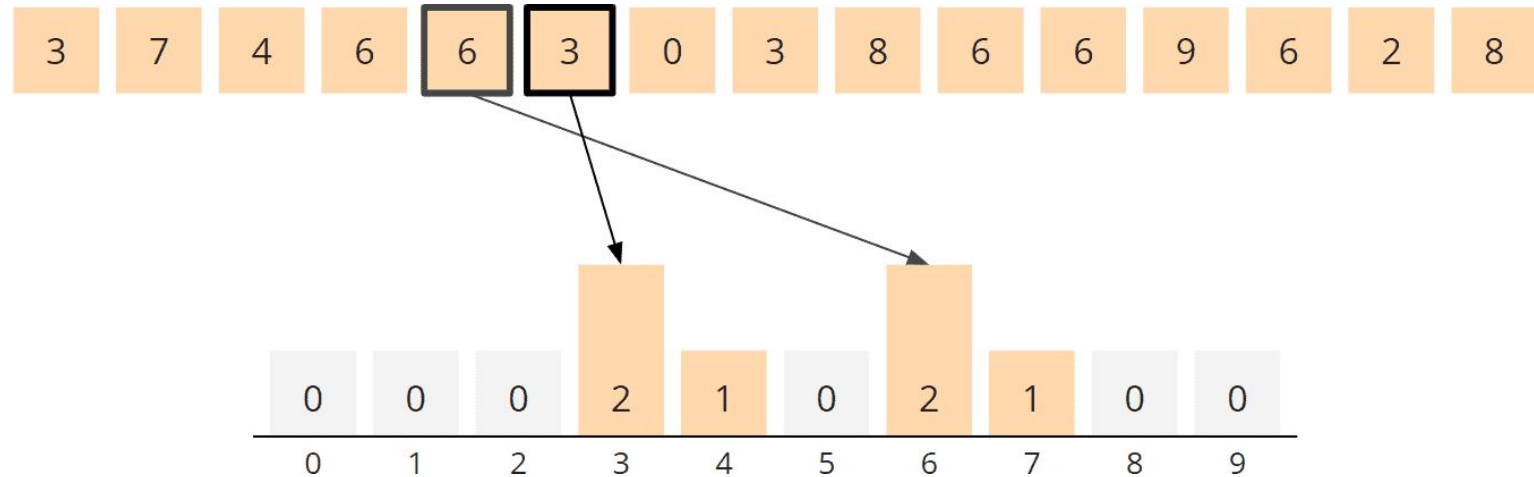




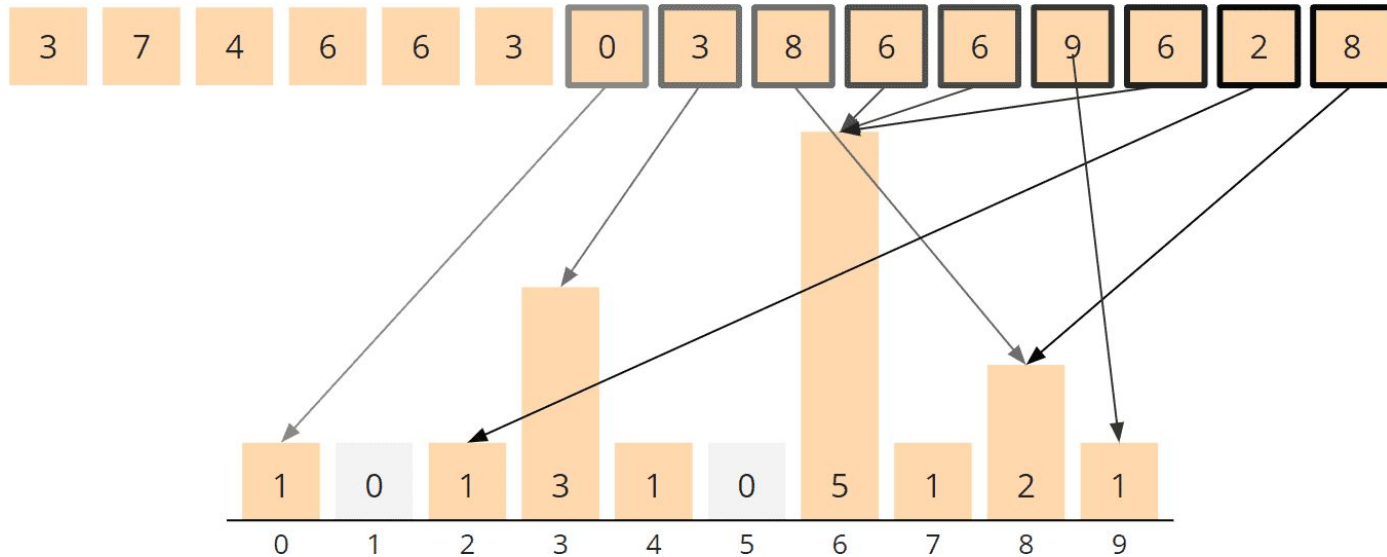
Elements 4 and 6 follow – thus, we increase the values at positions 4 and 6 by one each



The next two elements – the 6 and the 3 – are two elements that have already occurred before. Accordingly, the corresponding fields in the auxiliary array are increased from 1 to 2



The principle should be clear now. After also increasing the auxiliary array values for the remaining elements, the auxiliary array finally looks like this

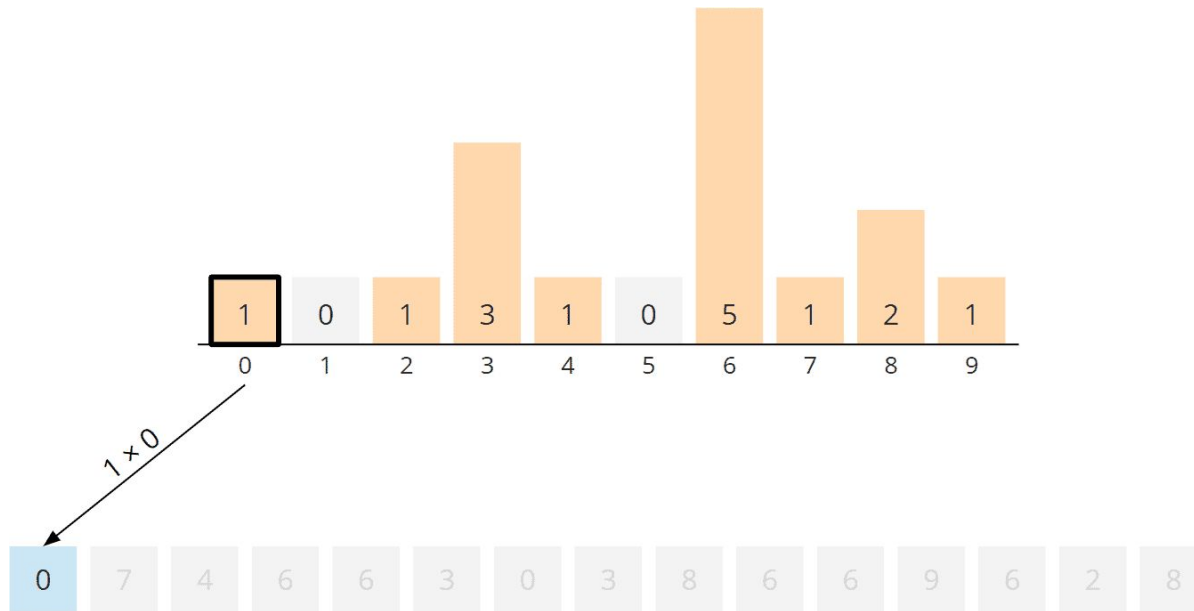


# Illustration

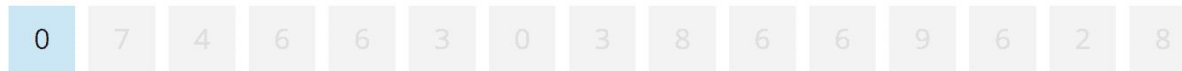
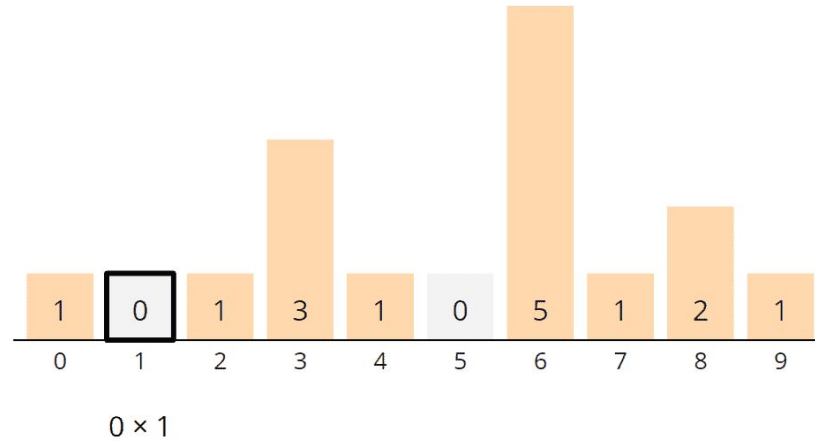
## Phase 2: Rearranging

We iterate once over the histogram array. We write the respective array index into the array to be sorted as often as the histogram indicates at the corresponding position.

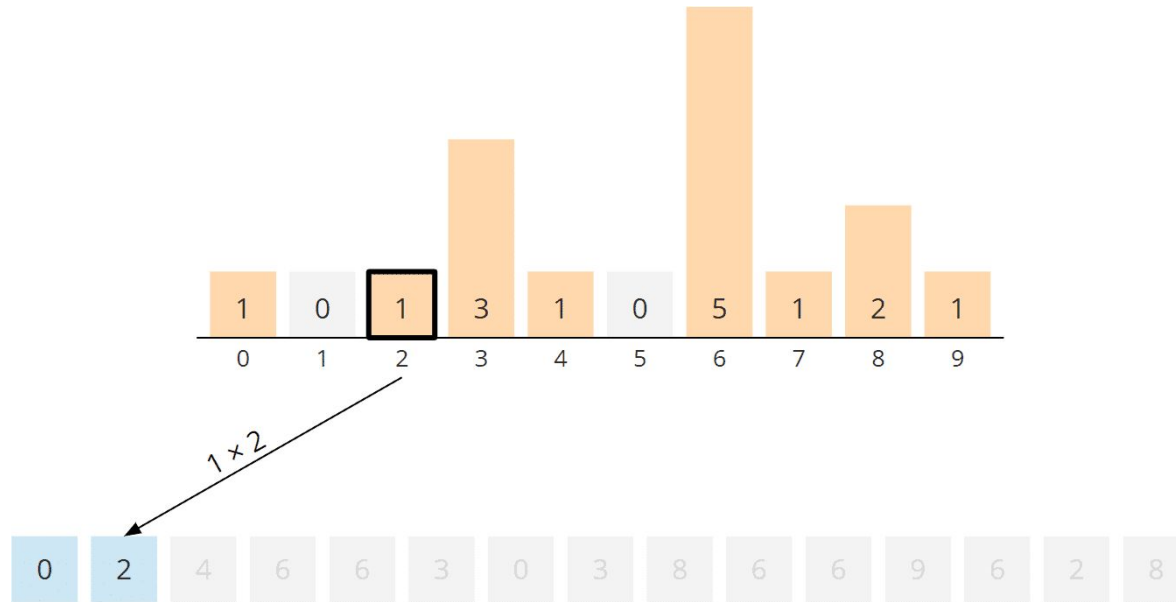
In the example, we start at position 0 in the auxiliary array. That field contains a 1, so we write the 0 exactly once into the array to be sorted.



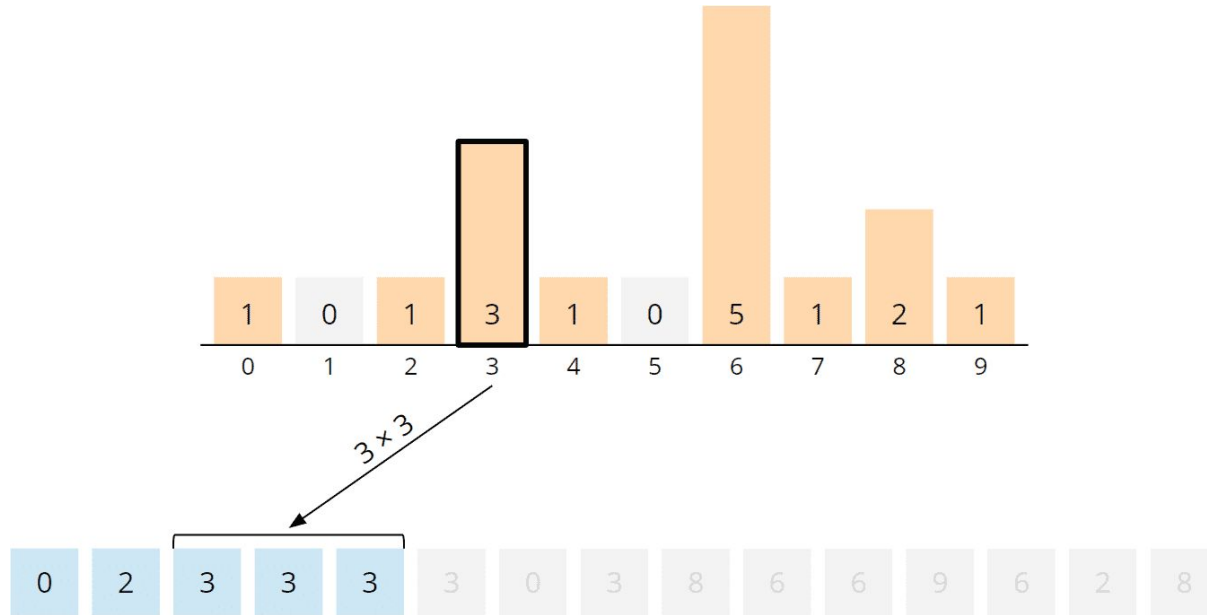
At position 1 in the histogram, there is a 0, meaning we skip this field – no 1 is written to the array to be sorted.



Position 2 of the histogram again contains a 1, so we write one 2 into the array to be sorted

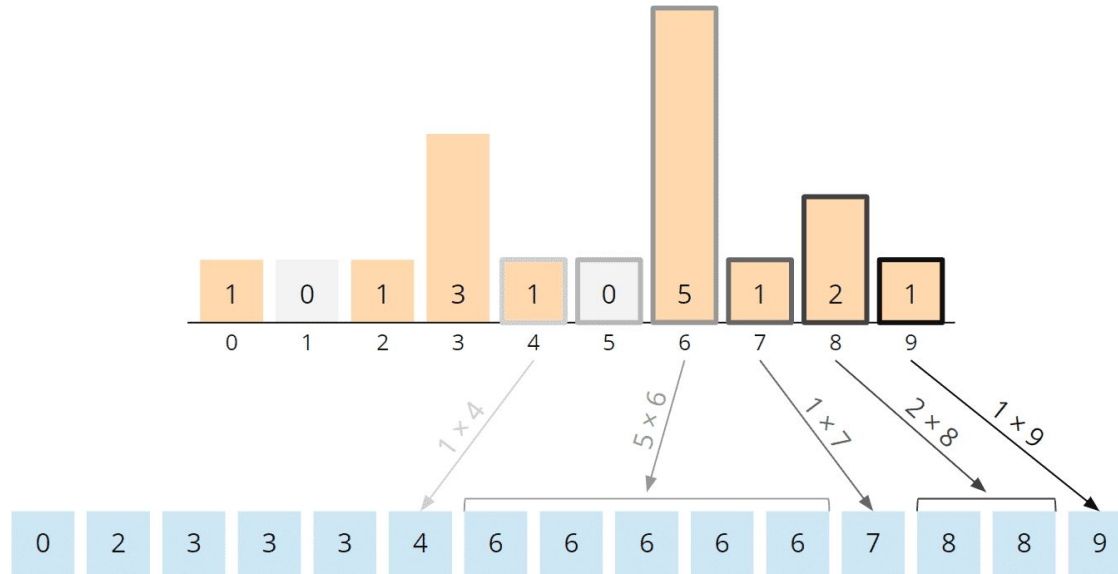


We come to position 3, which contains a 3; so we write three times a 3 into the array





And so it goes on. We write once the 4, five times the 6, once the 7, twice the 8 and finally once the 9 into the array to be sorted



Solve by Using Counting Sort

## Count Sorting with Negative Numbers

**Q:** Will the above implementation work when we include negative numbers?

**Q:** If NO, what can we do to make it work?

## Count Sorting with Negative Numbers

The problem with the previous counting sort was that we could not sort the elements if we have negative numbers in them. Because there are **no negative array indices**.

So what we do is, find the minimum element and we will store the count of that **minimum element at the zero** index. Then include **an offset** with the absolute value of the minimum element.

## Time complexity



Worst case ? \_\_\_\_\_

Best case ? \_\_\_\_\_

Average case ? \_\_\_\_\_

Stable ? \_\_\_\_\_

In Place? \_\_\_\_\_

**Note:** Counting sort is most efficient if the range of input values is not greater than the number of values to be sorted.

## Time complexity

The time complexity of counting sort algorithm is  $O(n+k)$  where  $n$  is the number of elements in the array and  $k$  is the range of the elements.

Worst case  $O(n+k)$

Best case  $O(n+k)$

Average case  $O(n+k)$

Stable ? \_\_\_\_\_No\_\_\_\_\_

In Place? \_\_\_\_\_No\_\_\_\_\_

**Note:** Counting sort is most efficient if the range of input values is not greater than the number of values to be sorted.



- Any **improvement in sorting time** significantly affect the **overall efficiency** and saves a great deal of computer time.
- Space constraints are usually less important than time complexity, for most sorting algorithms the amount of space needed is closer to  **$O(n)$** .

# Summary

	Time Complexity			Space	Stable	In Place
	Best Case	Average Case	Worst Case			
<b>Bubble</b>	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$	Yes	Yes
<b>Insertion</b>	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$	Yes	Yes
<b>Selection</b>	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$	Yes	Yes
<b>Counting</b>	$O(n + k)$	$O(n + k)$	$O(n + k)$	$O(n)$	No	No



# Using built-in functions to sort

# Python libraries: `sorted()` and `.sort()`

```
array = [1,2,5,4,3,6]
```

```
array.sort()
```

```
print(array)
```

```
# 1 2 3 4 5 6
```

```
array = [1,2,5,4,3,6]
```

```
array.sort(reverse=True)
```

```
print(array)
```

```
# 6 5 4 3 2 1
```

```
array = [1,2,5,4,3,6]
```

```
sorted_array = sorted(array)
```

```
print(sorted_array)
```

```
# 1 2 3 4 5 6
```

# Writing custom comparator

Default sorting sorts **based on the values**

What if you wanted to sort based on:

- In reverse
- Sorting based on some other criteria besides the values

# Example: sort an array based on a cost

```
costs = [1,3,2,5,1,3]
students = [1,2,3,4,5,6]

studentToCost = {}
for idx in range(len(students)):
    studentToCost[students[idx]] = costs[idx]

def customComparator(item):
    return studentToCost[item]

students.sort(key = customComparator)
print(students)
# [1, 5, 3, 2, 6, 4]
```

# Practice time

Relative Sort Array



# Helpful Resources

- [Big O Cheat Sheet](#)
- [Insertion Sort](#)
- [Selection Sort](#)
- [Counting Sort](#)
- [Stable vs Unstable Sort](#)

# Selecting a sorting algorithm

Language libraries often have sorting algorithms so we won't be coding our own sorting algorithms every time.

**Q:** When the range of numbers is small enough which algorithm should we use?

**Q:** What about When half of the array is already sorted?

# Practice Questions

[Bubble Sort](#)

[Insertion Sort](#)

[Counting Sort](#)

[Selection Sort](#)

[Sort Colors](#)

[Pancake Sorting](#)

[Find Target Indices After Sorting Array](#)

[Maximum Number Of Coins You Can Get](#)

[How Many Numbers Are Smaller Than The Current Number](#)