

# Toy Compiler

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## **Abstract**

A compiler for the Toy language in pure Java. It compiles to C as Intermediate code representation, then the *Clang* compiler is used to generate machine code

# Toy language Compiler

A simple Parser for the Toy language, made with Java CUP and JFlex. This implementation results in the creation of an executable file given a compliant toy file.

## Build

### Requirements

- Java 11
- Maven

### IntelliJ Configuration

This project is provided with a set of *IntelliJ Configurations*:

To run it just import it as a *Maven project* and click *Run*

### Maven Build

Alternatively you can compile manually typing:

```
mvn package
```

The *jar* file will be placed under the *target/* directory.

## Assignment overview

This compiler translates a toy program into a Clang-compliant C program. The generated *.c* file is then compiled and, after that, it's ready to run. The stages of the execution are the following:

- Lexical analysis
- Syntactic analysis
- (Optional) AST visualization
- Semantic analysis
- Toy2C translation

### Differences with the assignment

This implementation doesn't go that far from the assignment. Although, some variations have been made by the authors:

- The token MAIN has been introduced;
- The productions "*Main ::= PROC MAIN ...*" have been added, slightly changing the syntax of the language. These productions, along with the *ProcList*, make every Toy program syntactically compliant when the procedure *Main*:
  - appears one single time,
  - is the last one in the file,
  - returns an INT.

## Lexical analysis

This step is carried out by a Lexer written in **flex** and compiled with **Jflex** (an open source tool for generating Lexer in Java).

It is composed of a single **Lexer.flex** source file containing all the logic to generate a Lexer class, which is crucial for the next stage. The Lexer resulting by the compiling does the whole Lexical analysis.

## Lexical specification

This section describes the set of tokens and their corresponding pattern.

```
1  //Procedures
2  PROC "proc"
3  CORP "corp"
4  MAIN "main"
5
6  //Type
7  INT "int"
8  FLOAT "float"
9  BOOL "bool"
10 STRING "string"
11 VOID "void"
12
13 //Statements
14 WHILE "while"
15 DO "do"
16 OD "od"
17 READ "readln"
18 WRITE "write"
19 ASSIGN "assign"
20 IF "if"
21 THEN "then"
22 FI "fi"
23 ELSE "else"
24 ELIF "elif"
25
26 //Separators
27 LPAR "("
28 RPAR ")"
29 COLON ":"
30 COMMA ","
31 SEMI ";"
32
33
34 //Operators
35 ASSIGN ":="
36 PLUS "+"
37 MINUS "-"
38 TIMES "*"
39 DIV "/"
40 EQ "="
41 NE "<>"
42 LT "<"
43 LE "<="
44 GT ">"
```

```
45     GE ">="
46     AND "&&"
47     OR "||"
48     NOT "!"
49     NULL "null"
50     TRUE "true"
51     FALSE "false"
52     RETURN "->"
53
54
55     ALPHA=[A-Za-z]
56     DIGIT=[0-9]
57     NONZERO_DIGIT=[1-9]
58     NEWLINE=\\r\\n|\\r\\n\\n
59     WHITESPACE = | [ \\t\\f]
60     ID = (|_)*
61     INT = ((*)|0)
62     FLOAT = +)
63     STRING_TEXT = [^\\"]*
64     COMMENT_TEXT = [\\w\\.\\@]*
```

## Syntactic analysis

Just like the lexical analysis, the syntactic analysis is done by a generated Java class. This has been done with CUP (Construction of Useful Parsers), which implements *LALR(1)* parsing.

## Syntactic specification

This section describes the whole syntactic specification of the Toy language that was implemented. The grammar as-is doesn't allow *LALR(1)* parsing. In order to generate the parser, the *PEM-DAS* (Parenthesis, Exponents, Multiplications/Divisions, Additions/Subtractions) rule has been introduced.

```
1 Program ::= VarDeclList ProcList
2         ;
3
4     VarDeclList ::= /* empty */
5                 | VarDecl VarDeclList
6                 ;
7
8     VarDecl ::= Type IdListInit SEMI
9             ;
10
11     ProcList ::= Main
12              | Proc ProcList
13              ;
14
15     Type ::= INT
16           | BOOL
17           | FLOAT
18           | STRING
19           ;
20
21     IdListInit ::= ID
22                | IdListInit COMMA ID
23                | ID ASSIGN Expr
24                | IdListInit COMMA ID ASSIGN Expr
25                ;
26
27     Proc ::= PROC ID LPAR ParamDeclList RPAR ResultTypeList COLON
28            VarDeclList StatList RETURN ReturnExprs CORP SEMI
29            | PROC ID LPAR RPAR ResultTypeList COLON
30              VarDeclList StatList RETURN ReturnExprs CORP SEMI
31            |   PROC ID LPAR ParamDeclList RPAR ResultTypeList COLON
32              VarDeclList RETURN ReturnExprs CORP SEMI
33            | PROC ID LPAR RPAR ResultTypeList COLON
34              VarDeclList RETURN ReturnExprs CORP SEMI
35            ;
36
37     Main ::= PROC MAIN LPAR ParamDeclList RPAR INT COLON
38            VarDeclList StatList RETURN ReturnExprs CORP SEMI
39            | PROC MAIN LPAR RPAR INT COLON
40              VarDeclList StatList RETURN ReturnExprs CORP SEMI
41            |   PROC MAIN LPAR ParamDeclList RPAR INT COLON
42              VarDeclList RETURN ReturnExprs CORP SEMI
43            | PROC MAIN LPAR RPAR INT COLON
44              VarDeclList RETURN ReturnExprs CORP SEMI
45            ;
```

```

46
47 ResultTypeList ::= ResultType
48     | ResultType COMMA ResultTypeList
49     ;
50
51 ResultType ::= Type
52     | VOID
53     ;
54
55 ReturnExprs ::= ExprList
56     | /* empty */
57     ;
58
59 ParamDeclList ::= ParDecl
60     | ParamDeclList SEMI ParDecl
61     ;
62
63 ParDecl ::= Type IdList
64     ;
65
66 IdList ::= ID
67     | IdList COMMA ID
68     ;
69
70 StatList ::= Stat
71     | Stat StatList
72     ;
73
74 Stat ::= IfStat SEMI
75     | WhileStat SEMI
76     | ReadlnStat SEMI
77     | WriteStat SEMI
78     | AssignStat SEMI
79     | CallProc SEMI
80     ;
81
82 WhileStat ::= WHILE StatList RETURN Expr DO StatList OD
83     | WHILE Expr DO StatList OD
84     ;
85
86 IfStat ::= IF Expr THEN StatList ElifList Else FI
87     ;
88
89 ElifList ::= /* empty */
90     | Elif ElifList
91     ;
92
93 Elif ::= ELIF Expr THEN StatList
94     ;
95
96 Else ::= /* empty */
97     | ELSE StatList
98     ;
99
100 ReadlnStat ::= READ LPAR IdList RPAR
101     ;

```

```

102
103 WriteStat ::= WRITE LPAR ExprList RPAR
104         ;
105
106 AssignStat ::= IdList ASSIGN ExprList
107         ;
108
109 CallProc ::= ID LPAR ExprList RPAR
110           | ID LPAR RPAR
111           ;
112
113 ExprList ::= Expr
114           | Expr COMMA ExprList
115           ;
116
117 Expr ::= NULL
118       | TRUE
119       | FALSE
120       | INT_CONST
121       | FLOAT_CONST
122       | STRING_CONST
123       | ID
124       | ID LPAR ExprList RPAR
125       | ID LPAR RPAR
126       | Expr1 PLUS Expr2
127       | Expr1 MINUS Expr2
128       | Expr1 TIMES Expr2
129       | Expr1 DIV Expr2
130       | Expr1 AND Expr2
131       | Expr1 OR Expr2
132       | Expr1 GT Expr2
133       | Expr1 GE Expr2
134       | Expr1 LT Expr2
135       | Expr1 LE Expr2
136       | Expr1 EQ Expr2
137       | Expr1 NE Expr2
138       | MINUS Expr
139       | NOT Expr
140       ;

```

## The Abstract Syntax Tree

The implemented Grammar has been enhanced with a series of *actions*, one for each production.

Generally, these actions instantiate a *Node* object related to each of the *non-terminals* appearing in the right-hand side of the production itself. There are several different actions in this parser.

For example:

```
1 VarDeclList ::= /* empty */  {: RESULT = new
    LinkedList<VariableDeclarationNode>(); :}
2 | VarDecl VarDeclList      {: v1.add(vd); RESULT = v1; :}
3 ;
```

The *empty* production creates a new list containing the variables declared, while the second one appends a given variable to the list. In a successful situation, one would expect the process to eventually resolve in the *empty* statement, thus instantiating the list and adding the items found.

As the *Parser* elaborates a given source file, the corresponding (and unique) *Syntax Tree* is generated recursively. With the completion of the parsing process, a pointer to the root of the *Syntax tree* is returned, which comes in handy for the next step of the compiler.

### Tree visualization

Moreover, this implementation provides with a visualization of the *Syntactic Tree*, constructed in the previous step, via XML. The *Visitor* pattern fits this role perfectly. Once the parser has finished its job, the user can call the *ASTVisitor* on the root of the tree, which will generate a .xml file based on the instance of the *Syntactic tree*. The user can open the generated file using any Web Browser.



## Semantic Analysis

In order to check whether the input program is semantically compliant, a *SemanticVisitor* object is created and invoked on the root of the AST generated by the parser. Given the set of rules, a single visit of such tree is sufficient towards the semantic analysis.

### Inference table

The following tables describe each and every inference rule required to the type checking. The implementation of these rules can be found in the *TypeCheck* class.

Binary_op	First operand type	Second operand type	Resulting type
:=	int	int	int
:=	float	float	float
:=	string	string	string
:=	boolean	boolean	boolean
/ + -	int	int	int
/ + -	int	float	float
/ + -	float	int	float
/ + -	float	float	float
<= < == <> > >=	int	int	boolean
<= < == <> > >=	int	float	boolean
<= < == <> > >=	float	int	boolean
<= < == <> > >=	float	float	boolean
&&	boolean	boolean	boolean

Unary_op	Operand type	Resulting type
-	int	int
-	float	float
!	boolean	boolean

### Inference rules

These are the inference rules implemented in the *SemanticVisitor* class.

- TypeCheck rules

$$\frac{(expr : \tau) \in \Gamma}{\Gamma \vdash expr : \tau}$$

- Binary operation

$$\frac{\Gamma \vdash expr_1 : \tau_1 \quad \Gamma \vdash expr_2 : \tau_2 \quad \Gamma \vdash binary\_op(binOp, \tau_1, \tau_2) = \tau}{\Gamma \vdash (expr_1 binOp expr_2) : \tau}$$

- Unary operation

$$\frac{\Gamma \vdash expr : \tau_1 \quad \Gamma \vdash unary\_op(unOp, \tau_1) = \tau}{\Gamma \vdash (unOp expr) : \tau}$$

- Call procedure

$$\frac{\Gamma \vdash proc : \tau_i^{i \in \mathbb{N}} \rightarrow \tau_j^{j \in \mathbb{N}} \quad \Gamma \vdash par_i^{i \in \mathbb{N}} : \tau_i}{\Gamma \vdash proc(par_i^{i \in \mathbb{N}}) : \tau_j^{j \in \mathbb{N}}}$$

- While statement

$$\frac{\Gamma \vdash \text{cnd\_expr} : \mathbf{boolean} \quad \Gamma \vdash \text{while\_stmt}_i^{i \in \mathbb{N} \setminus \{0\}} \quad \Gamma \vdash \text{do\_stmt}_j^{j \in \mathbb{N}}}{\Gamma \vdash \mathbf{while} (\text{while\_stmt}_i^{i \in \mathbb{N} \setminus \{0\}}) \mathbf{do} (\text{do\_stmt}_j^{j \in \mathbb{N}}) \mathbf{od}}$$

- Assign statement

$$\frac{(x_i^{i \in \mathbb{N} \setminus \{0\}} : \tau_i) \in \Gamma \quad \Gamma \vdash (\text{expr}_j^{j \in \mathbb{N} \setminus \{0\}}) \rightarrow \tau_i^{i \in \mathbb{N} \setminus \{0\}}}{\Gamma \vdash x_i^{i \in \mathbb{N} \setminus \{0\}} = \text{expr}_j^{j \in \mathbb{N} \setminus \{0\}}}$$

- If statement

$$\frac{\Gamma \vdash \text{cnd\_expr} : \mathbf{boolean} \quad \Gamma \vdash \text{then\_stmt}_x^{x \in \mathbb{N}} \quad \Gamma \vdash \text{elif\_stmt}_y^{y \in \mathbb{N}} \quad \Gamma \vdash \text{else\_stmt}_z^{z \in \mathbb{N}}}{\Gamma \vdash \mathbf{if} \text{cnd\_expr} \mathbf{then} (\text{then\_stmt}_x^{x \in \mathbb{N}}) \mathbf{elif} (\text{elif\_smt}_y^{y \in \mathbb{N}}) \mathbf{else} (\text{else\_stmt}_z^{z \in \mathbb{N}}) \mathbf{fi}}$$

## Translating to the C language