Corso di Compilatori A.A. 2023/24

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1 Introduzione

Il seguente documento contiene le specifiche del linguaggio Toy2.

- La sezione delle specifiche lessicali contiene la lista dei token con i rispettivi pattern
- La sezione delle specifiche sintattiche contiene la grammatica utilizzata con all'interno tutti i non terminali (scritti in minuscolo e maiuscolo) ed i terminali (scritti interamente in maiuscolo), questa sezione contiene inoltre la tabella delle precedenze e i nodi dell'Abstract Syntax Tree
- La sezione delle specifiche semantiche riporta le regole di Type Checking e le tabelle per gli operatori
- La sezione di testing mostra tutti i test effettuati

2 Specifiche Lessicali

Token	Pattern		
VAR	var		
COLON	:		
ASSIGN	· ^ =		
SEMI			
COMMA	;		
TRUE	true		
FALSE	false		
REAL	real		
INTEGER	integer		
STRING	_		
BOOLEAN	string boolean		
RETURN	return		
	func		
FUNCTION			
TYPERETURN	->		
LPAR			
RPAR)		
PROCEDURE	proc		
WHILE	while		
ENDPROCEDURE	$_{10}^{\circ}$		
ENDFUNCTION	endfunc		
OUT	out		
WRITE	>		
WRITERETURN	>!		
DOLLARSIGN	\$		
READ	<		
IF	if		
THEN	then		
ELSE	else		
ENDIF	endif		
ELIF	elseif		
DO	do		
ENDWHILE	endwhile		
PLUS	+		
MINUS	- *		
TIMES	,		
DIV	/		
EQ	=		
NE	<>		
LT	<		
LE	<=		
GT	>		
GE	>=		
AND	&&		
OR	ļ.		
NOT	!		
ENDVAR	\\		
REF	(i) (i) (ii) (ii) (iii)		
ID	[a-zA-Z] ([a-zA-Z] [0-9] _)*		
STRING_CONST	\" ~\" [0-9]+		
INTEGER_CONST	[0-9]+		
REAL_CONST	[0-9]+ ("." [0-9]+)?		

3 Specifiche Sintattiche

3.1 Grammatica

```
{\bf Program} ::= {\bf Iter} \ {\bf Procedure} \ {\bf Iter}
{\tt IterNoProcedure} ::= {\tt VarDecls} \ {\tt IterNoProcedure}
    Function IterNoProcedure
   | /* empty */
{\rm Iter} ::= {\rm VarDecl} \ {\rm Iter}
    Function Iter
    Procedure Iter
   | /* empty */
VarDecl ::= VAR Decls
{\it Decls}::={\it Ids}~{\it COLON}~{\it Type}~{\it SEMI}~{\it Decls}
    Ids ASSIGN Consts SEMI Decls
    Ids COLON Type SEMI ENDVAR
    Ids ASSIGN Consts SEMI ENDVAR
Ids ::= ID COMMA Ids
   | ID
Consts ::= Const COMMA Consts
   Const
Const ::= REAL\_CONST
    INTEGER_CONST
    STRING_CONST
    TRUE
    FALSE
Type ::= REAL
    INTEGER
    STRING
    BOOLEAN
Function ::= FUNCTION ID LPAR FuncParams RPAR TYPERETURN Types
COLON Body ENDFUNCTION
```

 $Func Params ::= ID \ COLON \ Type \ Other Func Params$

```
| /* empty */
OtherFuncParams ::= COMMA ID COLON Type OtherFuncParams
  | /* empty */
Types ::= Type COMMA Types
  | Type
Procedure ::= PROCEDURE ID LPAR ProcParams RPAR COLON Body
ENDPROCEDURE
ProcParams::= ProcParamId COLON Type OtherProcParams
  | /* empty */
OtherProcParams ::= COMMA ProcParamId COLON Type OtherProcParams
  | /* empty */
ProcParamId ::= ID
  | OUT ID
Body ::= VarDecl Body
   Stat Body
   | /* empty */
Stat ::= Ids ASSIGN Exprs SEMI
   ProcCall SEMI
    RETURN Exprs SEMI
    WRITE IOArgs SEMI
    WRITERETURN IOArgs SEMI
    READ IOArgs SEMI
    IfStat SEMI
   WhileStat SEMI
FunCall ::= ID LPAR Exprs RPAR
   | ID LPAR RPAR
ProcCall ::= ID LPAR ProcExprs RPAR
   | ID LPAR RPAR
IfStat ::= IF Expr THEN Body Elifs Else ENDIF
{\bf Elifs}::={\bf Elif}\;{\bf Elifs}
  | /* empty */
Elif ::= ELIF Expr THEN Body
```

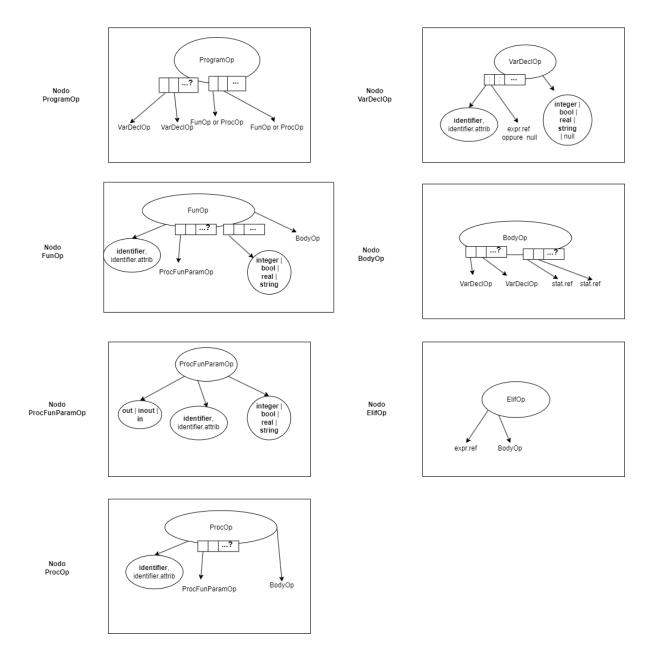
```
Else ::= ELSE \ Body
   | /* empty */
While Stat ::= WHILE Expr DO Body ENDWHILE
IOArgs ::= OtherIOArgs IOArgs
   | DOLLARSIGN LPAR Expr RPAR IOArgs
   | /* empty */
OtherIOArgs ::= OtherIOArgs PLUS OtherIOArgs
   | STRING_CONST
\label{eq:procExprs} \operatorname{ProcExprs} ::= \operatorname{Expr} \, \operatorname{COMMA} \, \operatorname{ProcExprs}
    REF ID COMMA ProcExprs
    Expr
    REF ID
Exprs ::= Expr COMMA Exprs
   Expr
Expr ::= FunCall
    REAL_CONST
    INTEGER_CONST
    STRING_CONST
    ID
    TRUE
    FALSE
    Expr PLUS Expr
    Expr MINUS Expr
    Expr TIMES Expr
    Expr DIV Expr
    Expr AND Expr
    Expr OR Expr
    Expr GT Expr
    Expr GE Expr
    Expr LT Expr
    Expr LE Expr
    Expr EQ Expr
    Expr NE Expr
    LPAR Expr RPAR %PAR
    MINUS Expr\% UMINUS
    NOT Expr
```

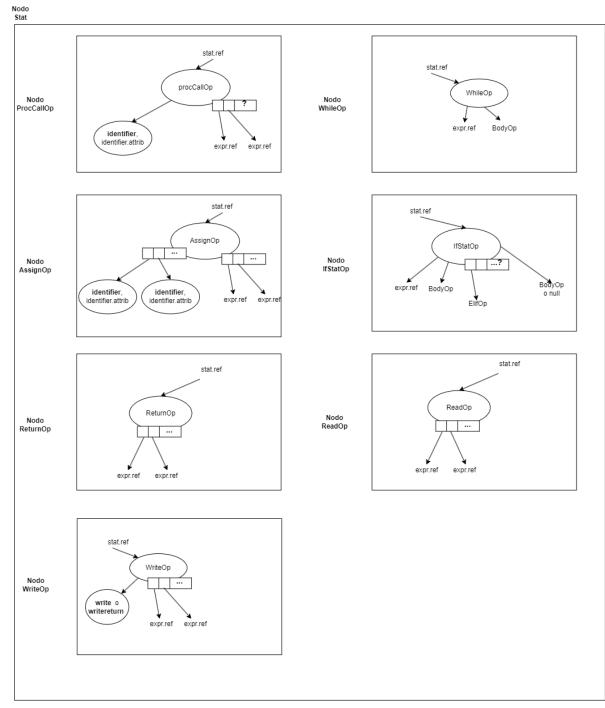
3.2 Tabella delle precedenze

La priorità della seguente tabella viene specificata come nell'ordine fornito da Java CUP, riga più in basso equivale a priorità più alta.

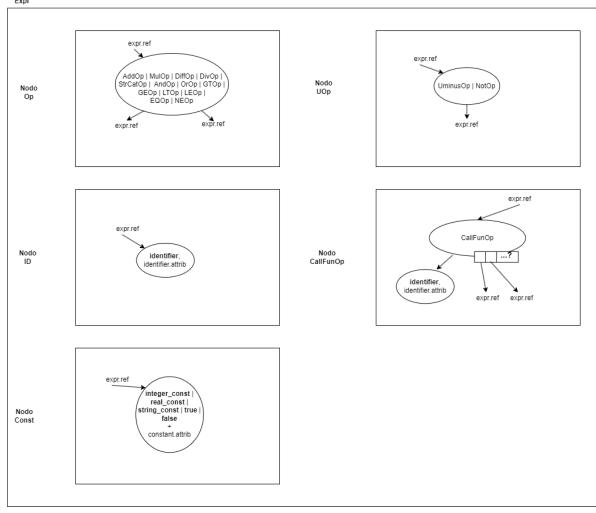
Token	Associatività	
ID VAR	SINISTRA	
COMMA	SINISTRA	
ASSIGN	DESTRA	
OR	SINISTRA	
AND	SINISTRA	
EQ NE	SINISTRA	
LE GE GT LT	SINISTRA	
PLUS MIUS	SINISTRA	
TIMES DIV	SINISTRA	
NOT UMINUS	DESTRA	
PAR	SINISTRA	

3.3 Abstract Syntax Tree









4 Specifiche Semantiche

Identificatore

 $\frac{\Gamma(id){=}\tau}{\Gamma{\vdash}id{:}\tau}$

Costanti

 $\Gamma \vdash \text{real_const} : \text{real}$

 $\Gamma \vdash \text{integer_const} : \text{integer}$ $\Gamma \vdash \text{string_const} : \text{string}$ $\Gamma \vdash \text{true} : \text{boolean}$ $\Gamma \vdash \text{false} : \text{boolean}$

Lista di istruzioni

 $\frac{\Gamma \vdash stmt_1 : notype \ \Gamma \vdash stmt_2 : notype}{\Gamma \vdash stmt_1, \ stmt_2 : notype}$

Chiamata a funzione

$$\frac{\Gamma \vdash f: \tau_1 \times \ldots \times \tau_n -> \sigma_1 \ldots \sigma_m \ \Gamma \vdash e_i: \tau_i^{i \in 1 \ldots n}}{\Gamma \vdash f(e_1, \ldots, e_n): \sigma_1 \ldots \sigma_m}$$

Chiamata a procedura

$$\frac{\Gamma \vdash f: \tau_1 \times \ldots \times \tau_n -> notype \ \Gamma \vdash e_i: \tau_i^{i \in 1 \ldots n}}{\Gamma \vdash f(e_1, \ldots, e_n): notype}$$

Assegnazione

$$\frac{\Gamma(id_i):\tau_i^{i\in 1...n}\ \Gamma\vdash e_i:\tau_i^{i\in 1...n}}{\Gamma\vdash id_1,...,id_n\wedge=e_1,...,e_n:notype}$$

Dichiarazione-Istruzione

 $\frac{\Gamma[id {\rightarrow} \tau] {\vdash} stmt : notype}{\Gamma {\vdash} \tau \ id; \ stmt : notype}$

While

 $\frac{\Gamma \vdash e:boolean \; \Gamma \vdash body:notype}{\Gamma \vdash \mathbf{while} \; e \; \mathbf{do} \; body \; \mathbf{endwhile}:notype}$

If-elseif-else

 $\frac{\Gamma \vdash e_1 : boolean \ \Gamma \vdash body_1 : notype \ \Gamma \vdash e_2, \ ..., \ e_n : boolean \ \Gamma \vdash body_2, \ ..., \ body_n : notype \ body_{n+1} : notype \ \Gamma \vdash \mathbf{if} \ e_1 \ \mathbf{then} \ body_1 \ \mathbf{elseif} \ e_2 \ \mathbf{then} \ body_2 \ ... \ \mathbf{elseif} \ e_n \ \mathbf{then} \ body_n \ \mathbf{else} \ body_{n+1} \ \mathbf{endif} : notype \ body_{n+1} \ \mathbf{endif} : notype \ \mathbf{else} \ body_{n+1} \ \mathbf{endif} : notype \ \mathbf{else} \$

Write

 $\frac{\Gamma {\vdash} e : \tau}{\Gamma {\vdash} {-->} e : notype}$

Read

$$\frac{\Gamma(id{=}\tau)\ \Gamma{\vdash}e:\tau}{\Gamma{\vdash}<{--}\ id\ e:notype}$$

\mathbf{Return}

$$\frac{\Gamma \vdash e_i : \tau_i^{i \in 1...n}}{\Gamma \vdash return \; e_1, \, ..., \, e_n : notype}$$

Operatori Unari

$$\frac{\Gamma \vdash e : \tau_1 \ optype1(op_1, \, \tau_1) = \tau}{\Gamma \vdash op_1 \ e : \tau}$$

op1	operando	risultato
MINUS	integer	integer
MINUS	real	real
NOT	boolean	boolean

Operatori Binari

$$\frac{\Gamma \vdash e_1 : \tau_1 \ \Gamma \vdash e_2 : \tau_2 \ optype2(op_2, \, \tau_1, \, \tau_2) = \tau}{\Gamma \vdash e_1 \ op_2 \ e_2 : \tau}$$

op2	operando1	operando2	risultato
DIV	integer	integer	real
ADD	string	integer	string
ADD	integer	string	string
ADD	string	real	string
ADD	real	string	string
ADD	string	boolean	string
ADD	boolean	string	string
ADD, MINUS, TIMES	integer	integer	integer
ADD, MINUS, TIMES, DIV	integer	real	real
ADD, MINUS, TIMES, DIV	real	integer	real
AND, OR	boolean	boolean	boolean
EQ, NE	integer	integer	boolean
EQ, NE	real	integer	boolean
EQ, NE	integer	real	boolean
EQ, NE	real	real	boolean
EQ, NE	string	string	boolean
EQ, NE	boolean	boolean	boolean
LT, LE, GT, GE	integer	integer	integer
LT, LE, GT, GE	real	integer	integer
LT, LE, GT, GE	integer	real	integer
LT, LE, GT, GE	real	real	integer

5 Test

Test	Descrizione
TestCubo	Pass
TestCuboError	Errore: endvar mancante
TestFibonacci	Pass
TestFibonacciError	Errore: manca il return
TestTabelline	Pass
TestTabellineError	Errore: il while non viene chiuso
TestConvertitore	Pass
TestConvertitoreError	Errore: gradi è di tipo integer, ma viene utilizzato come se fosse real
TestContaspazi	Pass
TestContaspaziError	Errore: c non è stato dichiarato
Sample	Pass
ProgramEs5	Pass