

Architetture dei Sistemi di Elaborazione

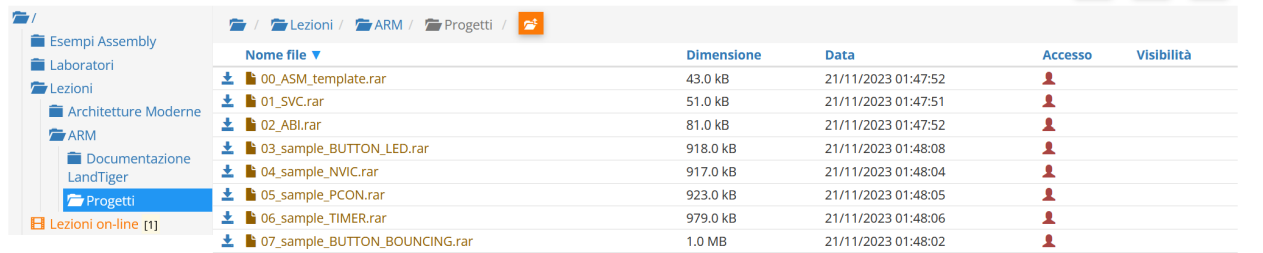
Delivery date:
30th November 2023

Laboratory 6

Expected delivery of **lab_06.zip** must include:

- Solutions of the exercises 1, 2 and 3
- this document compiled possibly in pdf format.

Starting from the ASM_template project (available on Portale della Didattica), solve the following exercises.



Nome file	Dimensione	Data	Accesso	Visibilità
00_ASM_template.rar	43.0 kB	21/11/2023 01:47:52		
01_SVC.rar	51.0 kB	21/11/2023 01:47:51		
02_ABI.rar	81.0 kB	21/11/2023 01:47:52		
03_sample_BUTTON_LED.rar	918.0 kB	21/11/2023 01:48:08		
04_sample_NVIC.rar	917.0 kB	21/11/2023 01:48:04		
05_sample_PCON.rar	923.0 kB	21/11/2023 01:48:05		
06_sample_TIMER.rar	979.0 kB	21/11/2023 01:48:06		
07_sample_BUTTON_BOUNCING.rar	1.0 MB	21/11/2023 01:48:02		

- 1) Write a program using the ARM assembly that performs the following operations:
- Initialize registers R1, R3 and R4 to random signed values
 - Sum R1 to R3 ($R1+R3$) and store the result in R2
 - Subtract R4 to R2 ($R4-R2$) and store the result in R5
 - Force, using the debug register window, a set of specific values to be used in the program to provoke the following flag to be updated **once at a time** (whenever possible) to 1:
 - carry
 - overflow
 - negative
 - zero
 - Report the selected values in the table below.

	Please, report the hexadecimal representation of the values			
Updated flag	R1 + R3		R4 - R2	
	R1	R3	R4	R2
Carry = 1	0xB1000000	0x50000000	0x00000002	0x00000001
Carry = 0	0x01000001	0x50000000	0x00000005	0x00000004
Overflow	0x50000000	0x50000000	0x0F000000	0x80000000
Negative	0x80000000	0x40000000	0xE0000001	0x00000001
Zero	0x00000000	0x00000000	0x00000001	0x00000001

Please explain the cases when it is **not** possible to force a **single** FLAG condition:
Per l'addizione, l'unico caso in cui non è possibile forzare un solo FLAG è quello dell'overflow, perché se sommo 2 numeri positivi molto grandi per avere overflow avrò un risultato negativo (quindi anche il flag N sarà attivo) mentre se sommo 2 numeri negativi grandi per avere overflow avrò un risultato positivo ma avrò anche carry. Per quanto riguarda la sottrazione invece, i casi sono 2: sempre quello dell'overflow, per lo stesso motivo di prima; quello del flag Z, perché sottraendo due numeri uguali (per ottenere zero) non posso avere borrow e il carry nella sottrazione è definito come $C=1-B$ quindi avrò anche $C=1$.

- 2) Write two versions of a program that performs the following operations:
- Initialize registers R2 and R3 to random signed values
 - Compare the two registers:
 - If they differ, store in the register R5 the minimum among R2 and R3

- Otherwise, perform on R3 a logical left shift of 1 (is it equivalent to what?), sum R2 and store the result in R4 (i.e. $r4=(r3<<1)+r2$).

First, solve it by resorting to 1) a traditional assembly programming approach using conditional branches and then compare the execution time with a 2) conditional instructions execution approach.

Report the execution time in the two cases in the table that follows.

NOTE, report the number of clock cycles (cc), as well as the simulation time in milliseconds (ms) considering a cpu clock (clk) frequency of 16 MHz.

Refer to the guide “howto_setup_keil” to change the clock frequency in Keil.

	R2==R3 [cc]	R2==R3 [ms]	R2!=R3 [cc]	R2!=R3 [ms]
1) Traditional	12	0.00075	14	0.00087
2) Conditional Execution	13	0.00081	11	0.00063

- 3) Write a program that calculates the trailing zeros of a variable. The trailing zeros are computed by counting the number of zeros starting from the least significant bit and stopping at the first 1 encountered: e.g., the trailing zeros of 0b10100000 are 5. The variable to check is in R1. After the count, if the number of trailing zeros is odd, perform the sum between R2 and R3. If the number of trailing zeros is even, perform the difference between R2 and R3. In both cases the result is placed in R4.

Implement the ASM code that performs the following operations:

- Determines whether the number of trailing zeros of R1 is odd or even.
- As a result, the value of R4 is computed as follows:
 - If the trailing zeros are even, R4 is the difference between R2 and R3
 - Else, R4 is the sum of R2 and R3
- Report code size and execution time (with 15MHz clk) in the following table.

Code size [Bytes]	Execution time [s]	
	If R1 is even	Otherwise
72	0.00000093	0.00000100

ANY USEFUL COMMENT YOU WOULD LIKE TO ADD ABOUT YOUR SOLUTION:

Nel caso dell'esercizio 2 ho approssimato i cicli di clock al numero intero più grande.

Per l'esercizio 3 ho considerato gli zeri del valore in binario, shiftando il valore a sinistra fin quando attivo il Flag Z. Il numero di iterazioni che mi sono servite per portare il numero al valore 0 sono i bit che non fanno parte dei trailing zeros. Quindi sottraggo a 32 il numero di cifre.