Team notebook

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1 Data structures

1.1 Centroid decomposition

```
namespace decomposition {
  int cnt[MAX], depth[MAX], f[MAX];
  int dfs (int u, int p = -1) {
    cnt[u] = 1;
    for (int v : g[u])
      if (!depth[v] && v != p)
        cnt[u] += dfs(v, u);
```

```
return cnt[u];
int get_centroid (int u, int r, int p = -1) {
 for (int v : g[u])
   if (!depth[v] && v != p && cnt[v] > r)
     return get_centroid(v, r, u);
 return u;
int decompose (int u, int d = 1) {
 int centroid = get_centroid(u, dfs(u) >> 1);
 depth[centroid] = d;
 /// magic function
 for (int v : g[centroid])
   if (!depth[v])
     f[decompose(v, d + 1)] = centroid;
 return centroid:
int lca (int u, int v) {
 for (; u != v; u = f[u])
   if (depth[v] > depth[u])
     swap(u, v);
 return u;
```

1.2 Fenwick tree

```
/// Complexity: log(|N|)

/// Tested: https://tinyurl.com/y88y7ws7

int lower_find(int val) { /// last value < or <= to val

int idx = 0;

for(int i = 31-__builtin_clz(n); i >= 0; --i) {

int nidx = idx | (1 << i);

if(nidx <= n && bit[nidx] <= val) { /// change <= to <

val -= bit[nidx];

idx = nidx;

}

return idx;

}
```

1.3 Heavy light decomposition

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/ybdbmbw7(problem L)
int idx;
vector<int> len, hld_child, hld_index, hld_root, up;
void dfs( int u, int p = 0 ) {
```

```
len[u] = 1;
 up[u] = p;
 for( auto& v : g[u] ) {
   if( v == p ) continue;
   depth[v] = depth[u]+1;
   dfs(v, u);
   len[u] += len[v];
   if( hld_child[u] == -1 || len[hld_child[u]] < len[v] )</pre>
    hld child[u] = v:
void build_hld( int u, int p = 0 ) {
 hld index\lceil u \rceil = idx++:
 narr[hld_index[u]] = arr[u]; /// to initialize the segment tree
 if( hld_root[u] == -1 ) hld_root[u] = u;
 if( hld child[u] != -1 ) {
   hld root[hld child[u]] = hld root[u]:
   build hld(hld child[u], u):
 for( auto& v : g[u] ) {
   if( v == p || v == hld child[u] ) continue;
   build_hld(v, u);
void update_hld( int u, int val ) {
 update_tree(hld_index[u], val);
data query_hld( int u, int v ) {
 data val = NULL_DATA;
 while( hld_root[u] != hld_root[v] ) {
   if( depth[hld_root[u]] < depth[hld_root[v]] ) swap(u, v);</pre>
   val = val+query_tree(hld_index[hld_root[u]], hld_index[u]);
   u = up[hld_root[u]];
 if( depth[u] > depth[v] ) swap(u, v);
 val = val+querv tree(hld index[u], hld index[v]);
 return val:
/// when updates are on edges use:
/// if (depth[u] == depth[v]) return val;
/// val = val+query_tree(hld_index[u] + 1, hld_index[v]);
void build(int n, int root) {
 len = hld index = up = depth = vector<int>(n+1):
 hld child = hld root = vector<int>(n+1, -1):
 idx = 1; /// segtree index [1, n]
 dfs(root, root): build hld(root, root):
 /// initialize data structure
```

1.4 Mo's

```
/// Complexity: O(|N+Q|*sart(|N|)*|ADD/DEL|)
/// Tested: Not vet
// Requires add(), delete() and get_ans()
struct query {
 int l. r. idx:
 query (int 1, int r, int idx) : 1(1), r(r), idx(idx) {}
int S; // s = sqrt(n)
bool cmp (query a, query b) {
 if (a.1/S != b.1/S) return a.1/S < b.1/S;
 return a.r > b.r:
S = sqrt(n); // n = size of array
sort(q.begin(), q.end(), cmp);
int 1 = 0, r = -1;
for (int i = 0; i < q.size(); ++i) {</pre>
 while (r < q[i].r) add(++r);
 while (1 > q[i].1) add(--1);
 while (r > q[i].r) del(r--);
 while (1 < q[i].1) del(1++);</pre>
 ans[q[i].idx] = get_ans();
```

3

1.5 Order statistics

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
//methods
tree.find_by_order(k) //returns pointer to the k-th smallest element
tree.order_of_key(x) //returns how many elements are smaller than x
//if element does not exist
tree.end() == tree.find_by_order(k) //true
```

1.6 Persistent segment tree

```
/// Complexity: O(|N|*log|N|)
/// Tested: Not yet
struct node {
  node *left, *right;
  int v;
  node() : left(this), right(this), v(0) {}
  node(node *left, node *right, int v) :
   left(left), right(right), v(v) {}
  node* update(int 1, int r, int x, int value) {
```

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```
if (1 == r) return new node(nullptr, nullptr, v + value);
int m = (1 + r) / 2;
if (x <= m)
    return new node(left->update(1, m, x, value), right, v + value);
    return new node(left, right->update(m + 1, r, x, value), v + value);
};
```

1.7 Rmq

```
/// Complexity: 0(|N|*log|N|)

/// Tested: https://tinyurl.com/y739tcsj

struct rmq {

    vector<vector<int> > table;
    rmq(vector<int> &v) : table(v.size() + 1, vector<int>(20)) {
        int n = v.size()+1;
        for (int i = 0; i < n; i++) table[i][0] = v[i];
        for (int j = 1; (1<<j) <= n; j++)
            for (int i = 0; i + (1<<j-1) < n; i++)
                table[i][j] = max(table[i][j-1], table[i + (1<<j-1)][j-1]);
    }

int query(int a, int b) {
    int j = 31 - __builtin_clz(b-a+1);
    return max(table[a][j], table[b-(1<<j)+1][j]);
    }
};
```

1.8 Sack

```
/// Complexity: |N|*log(|N|)
/// Tested: https://tinyurl.com/y9fz8vdt
int dfs(int u, int p = -1) {
    who[t] = u; fr[u] = t++;
    pii best = {0, -1};
    int sz = 1;
    for(auto v : g[u])
        if(v != p) {
        int cur_sz = dfs(v, u);
        sz += cur_sz;
        best = max(best, {cur_sz, v});
    }
    to[u] = t-1;
    big[u] = best.second;
    return sz;
}
void add(int u, int x) { /// x == 1 add, x == -1 delete
        cnt[u] += x;
}
```

```
void go(int u, int p = -1, bool keep = true) {
 for(auto v : g[u])
  if(v != p && v != big[u])
     go(v, u, 0);
 if(big[u] != -1) go(big[u], u, 1);
 /// add all small
 for(auto v : g[u])
   if(v != p && v != big[u])
     for(int i = fr[v]; i <= to[v]; i++)</pre>
       add(who[i]. 1):
 add(u, 1):
 ans[u] = get(u);
 if(!keep)
   for(int i = fr[u]; i <= to[u]; i++)</pre>
     add(who[i], -1):
void solve(int root) {
 t = 0:
 dfs(root):
 go(root);
```

1.9 Sqrt Decomposition

```
struct bucket {
 int 1, r, lazy;
 bucket(int 1, int r) : 1(1), r(r), lazy(0) {}
 void build() {
   for(int i = 1; i <= r; i++) a[i] += lazy;</pre>
   for(int i = 1: i \le r: i++) {} /// build DS from scratch
   lazv = 0:
 void update(int L, int R, 11 v) {
   if(L == 1 && R == r) lazv += v:
   else { /// handle by hand
     for(int i = L: i <= R: i++) a[i] += v:
     build():
   }
 int query(int L, int R) {
   int ans = INT MIN:
   if(L == 1 && R == r) ans = ds.get_max(x);
   else { /// handle by hand
     for(int i = L; i <= R; i++) ans = max(ans, abs(a[i] + x) * b[i]);
   return ans;
};
{ /// at main(), update from a to b, len is the size of bucket
 int 1 = a / len, r = b / len;
 for(int i = 1 i <= r; i++) { /// in theory, all are complete</pre>
```

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```
int x = max(l, i*len);
int y = min(r, (i+1)*len-1);
bucket[i].operation(x, y);
}
}
```

1.10 Treap

```
/// Complexity: O(|N|*log|N|)
/// Tested: Not vet
mt19937_64 rng64(chrono::steady_clock::now().time_since_epoch().count());
uniform int distribution<ll> dis64(0, 111<<60):
template <typename T>
class treap {
private:
       struct item:
        typedef struct item * pitem:
       pitem root = NULL;
       struct item {
               11 prior; int cnt, rev;
               T key, add, fsum;
               pitem 1, r;
               item(T x, ll p) {
                       add = 0*x; key = fsum = x;
                       cnt = 1; rev = 0;
                       1 = r = NULL; prior = p;
               }
       }:
       int cnt(pitem it) { return it ? it->cnt : 0; }
       void upd cnt(pitem it) {
               if(it) it->cnt = cnt(it->1) + cnt(it->r) + 1:
        void upd_sum(pitem it) {
               if(it) {
                       it->fsum = it->key;
                       if(it->1) it->fsum += it->1->fsum:
                       if(it->r) it->fsum += it->r->fsum:
               }
       void update(pitem t, T add, int rev) {
               if(!t) return:
               t->add = t->add + add;
               t->rev = t->rev ^ rev;
               t \rightarrow key = t \rightarrow key + add;
               t \rightarrow fsum = t \rightarrow fsum + cnt(t) * add;
       }
       void push(pitem t) {
               if(!t || (t->add == 0*T() && t->rev == 0)) return;
               update(t->1, t->add, t->rev);
               update(t->r, t->add, t->rev);
               if(t\rightarrow rev) swap(t\rightarrow 1,t\rightarrow r);
```

```
t->add = 0*T(); t->rev = 0;
       void merge(pitem & t, pitem 1, pitem r) {
              push(1); push(r);
              if(!1 || !r) t = 1 ? 1 : r;
              else if(1->prior > r->prior) merge(1->r, 1->r, r), t = 1;
              else merge(r->1, 1, r->1), t = r;
              upd_cnt(t); upd_sum(t);
       void split(pitem t, pitem & 1, pitem & r, int index) { // split index = how many
           elements
              if(!t) return void(1 = r = 0);
              push(t):
              if(index \le cnt(t->1)) split(t->1, 1, t->1, index), r = t;
              else split(t->r, t->r, r, index - 1 - cnt(t->1)), 1 = t;
              upd cnt(t): upd sum(t):
       void insert(pitem & t. pitem it. int index) { // insert at position
              push(t):
              if(!t) t = it:
              else if(it->prior > t->prior) split(t, it->l, it->r, index), t = it;
              else if(index <= cnt(t->1)) insert(t->1, it, index);
              else insert(t->r, it, index-cnt(t->l)-1);
              upd_cnt(t); upd_sum(t);
       void erase(pitem & t, int index) {
              push(t):
              if(cnt(t->1) == index) merge(t, t->1, t->r);
              else if(index < cnt(t->1)) erase(t->1, index);
              else erase(t->r, index - cnt(t->1) - 1);
              upd_cnt(t); upd_sum(t);
       T get(pitem t, int index) {
              push(t);
              if(index < cnt(t->1)) return get(t->1, index);
              else if(index > cnt(t->1)) return get(t->r, index - cnt(t->1) - 1):
              return t->kev:
 T query_sum (pitem &t, int 1, int r) {
              pitem 11. r1:
              split (t, l1, r1, r + 1);
              pitem 12, r2:
              split (11, 12, r2, 1):
              T ret = r2 -> fsum:
              pitem t2:
              merge (t2, 12, r2);
              merge (t, t2, r1);
              return ret:
public:
       int size() { return cnt(root); }
       void insert(int pos, T x) {
              if(pos > size()) return;
```

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2 Dp optimization

2.1 Convex hull trick dynamic

```
/// Complexity: O(|N|*log(|N|))
/// Tested: Not yet
typedef 11 T;
const T is_query = -(1LL<<62); // special value for query</pre>
struct line {
 T m, b;
 mutable multiset<line>::iterator it, end;
  const line* succ(multiset<line>::iterator it) const {
   return (++it == end ? nullptr : &*it);
  bool operator < (const line& rhs) const {</pre>
   if(rhs.b != is query) return m < rhs.m;</pre>
    const line *s = succ(it):
    if(!s) return 0:
    return b-s->b < (s->m-m)*rhs.m:
};
struct hull dynamic : public multiset<line> { // for maximum
 bool bad(iterator v) {
    iterator z = next(v):
    if(v == begin()){
     if(z == end()) return false;
     return v->m == z->m && v->b <= z->b:
    iterator x = prev(y);
    if(z == end()) return y->m == x->m && y->b <= x->b;
    return (x->b - y->b)*(z->m - y->m) >=
          (y->b - z->b)*(y->m - x->m);
  iterator next(iterator y){ return ++y; }
  iterator prev(iterator y){ return --y; }
  void add(T m, T b){
   iterator y = insert((line){m, b});
```

```
y->it = y; y->end = end();
if(bad(y)){ erase(y); return; }
while(next(y) != end() && bad(next(y))) erase(next(y));
while(y != begin() && bad(prev(y))) erase(prev(y));
}
T eval(T x){
line 1 = *lower_bound((line){x, is_query});
return l.m*x+l.b;
}
};
```

2.2 Convex hull trick

```
struct line {
 11 m. b:
 ll eval (ll x) { return m*x + b; }
struct cht {
 vector<line> lines:
 vector<lf> inter:
 lf get_inter(line &a, line &b) { return lf(b.b - a.b) / (a.m - b.m); }
 inline bool ok(line &a. line &b. line &c) {
   return lf(a,b-c,b) / (c,m-a,m) > lf(a,b-b,b) / (b,m-a,m);
 void add(line 1) {
   n = lines.size():
   if(n && lines.back().m == l.m && lines.back().b >= l.b) return;
   if(n == 1 && lines.back().m == l.m && lines.back().b < l.b) lines.pop back(), n--:
   while(n >= 2 && !ok(lines[n-2], lines[n-1], 1)) {
     lines.pop_back(); inter.pop_back();
   lines.push_back(1); n++;
   if(n >= 2) inter.push_back(get_inter(lines[n-1], lines[n-2]));
 ll get_max(lf x) {
   if(lines.size() == 0) return LLONG_MIN;
   if(lines.size() == 1) return lines[0].eval(x);
   int pos = lower_bound(inter.begin(), inter.end(), x) - inter.begin();
   return lines[pos].eval(x);
};
```

2.3 Divide and conquer

```
/// Complexity: O(|N|*|K|*log|N|))
/// ****** Theory ******
```

```
/// dp[k][i]=min(dp[k1][i]+C[i][i]), i < i
/// opt[k][i] opt[k][i+1].
/// A sufficient (but not necessary) condition for above is
/// C[a][c] + C[b][d] C[a][d] + C[b][c] where a < b < c < d .
void go(int k, int l, int r, int opl, int opr) {
 if(1 > r) return;
 int mid = (1 + r) / 2, op = -1;
 ll &best = dp[mid][k];
 best = INF:
 for(int i = min(opr, mid); i >= opl; i--) {
   ll cur = dp[i][k-1] + cost(i+1, mid):
   if(best > cur) {
    best = cur:
     op = i;
 }
 go(k, 1, mid-1, opl, op);
 go(k, mid+1, r, op, opr);
```

2.4 Knuth

```
/// Complexity: O(|N|^2))
/// Tested: https://tinyurl.com/y6ofp8wb
/// ****** Theory ******
/// dp[i][j] = min(dp[i][k]+dp[k][j])+C[i][j], i<k<j
/// where opt[i][j1]
                       opt[i][j] opt[i+1][j].
/// sufficient (but not necessary) condition for above is
/// C[a][c] + C [b][d] C[a][d] + C [b][c] and C[b][c] C[a][d] where
for(int i = 1; i <= n; i++) {</pre>
 opt[i][i] = i;
 dp[i][i] = sum[i] - sum[i-1];
for(int len = 2; len <= n; len++)</pre>
 for(int 1 = 1; l+len-1 <= n; l++) {</pre>
   int r = 1+len-1;
   dp[1][r] = oo;
   for(int i = opt[l][r-1]; i <= opt[l+1][r]; i++) {</pre>
     11 \text{ cur} = dp[1][i-1] + dp[i+1][r] + sum[r] - sum[1-1];
     if(cur < dp[1][r]) {</pre>
       dp[1][r] = cur;
       opt[1][r] = i;
```

3 Formulas

3.1 2-SAT rules

- $p \to q \equiv \neg p \lor q$
- $p \to q \equiv \neg q \to \neg p$
- $p \lor q \equiv \neg p \to q$
- $p \land q \equiv \neg(p \to \neg q)$
- $\neg(p \to q) \equiv p \land \neg q$
- $(p \to q) \land (p \to r) \equiv p \to (q \land r)$
- $(p \to q) \lor (p \to r) \equiv p \to (q \lor r)$
- $(p \to r) \land (q \to r) \equiv (p \land q) \to r$
- $(p \to r) \lor (q \to r) \equiv (p \lor q) \to r$
- $(p \land q) \lor (r \land s) \equiv (p \lor r) \land (p \lor s) \land (q \lor r) \land (q \lor s)$

3.2 Burnside's lemma

$$\#orbitas = \frac{1}{|G|} \sum_{g \in G} |fix(g)|$$

- 1. **G**: Las acciones que se pueden aplicar sobre un elemento, incluyendo la identidad, eg. Shift 0 veces, Shift 1 veces...
- 2. Fix(g): Es el número de elementos que al aplicar g vuelven a ser ser ellos mismos
- 3. Órbita: El conjunto de elementos que pueden ser iguales entre si al aplicar alguna de las acciones de G

3.3 Catalan Numbers

- $C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{(n+1)!n!} \text{ con } n \ge 0, C_0 = 1 \text{ y } C_{n+1} = \frac{2(2n+1)!}{n+2} C_n$
- 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012, 742900, 2674440, 9694845, 35357670

3.4 Combinatorics

- Distribute N objects among K people $\binom{n}{k} = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Hockey-stick identity $\sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1}$

3.5 Compound Interest

• N is the initial population, it grows at a rate of R. So, after X years the popularion will be $N \times (1+R)^X$

3.6 DP optimization theory

Name	Original Recurrence	Sufficient Condition		
CH 1	$dp[i] = min_{j < i} \{dp[j] + b[j] * a[i]\}$	$b[j] \ge b[j+1]$ Optionally	$O(n^2)$	O(n)
		$a[i] \le a[i+1]$		
CH 2	$dp[i][j] = min_{k < j} \{ dp[i-1][k] +$	$b[k] \ge b[k+1]$ Option-	$O(kn^2)$	O(kn)
	$b[k]*a[j]$ }	ally $a[j] \le a[j+1]$		
D&Q	$dp[i][j] = min_{k < j} \{ dp[i-1][k] +$	$A[i][j] \le A[i][j+1]$	$O(kn^2)$	$O(kn\log n)$
	$C[k][j]\}$			
Knuth	$dp[i][j] = min_{i < k < j} \{dp[i][k] +$	$A[i,j-1] \leq A[i,j] \leq$	$O(n^3)$	$O(n^2)$
	$dp[k][j]\} + C[i][j]$	A[i+1,j]		

Notes:

- A[i][j] the smallest k that gives the optimal answer, for example in dp[i][j] = dp[i-1][k] + C[k][j]
- C[i][j] some given cost function
- We can generalize a bit in the following way $dp[i] = \min_{j < i} \{F[j] + b[j] * a[i]\}$, where F[j] is computed from dp[j] in constant time

3.7 Euler Totient properties

- $\phi(p) = p 1$
- $\phi(n*m) = \phi(n)*\phi(m)$ si $\gcd(n,m) = 1$
- $\phi(n) = n(1-\frac{1}{p_1})(1-\frac{1}{p_2})...(1-\frac{1}{p_k})$ donde p_i es primo y divide a n

3.8 Fermat's theorem

Let m be a prime and x and m coprimes, then:

- $\bullet \ x^{m-1} \mod m = 1$
- $x^k \mod m = x^{k \mod (m-1)} \mod m$
- $x^{\phi(m)} \mod m = 1$

3.9 Great circle distance or geographical distance

Great circle distance or geographical distance

- $d = \text{great distance}, \phi = \text{latitude}, \lambda = \text{longitude}, \Delta = \text{difference}$ (all the values in radians)
- σ = central angle, angle form for the two vector
- $d = r * \sigma$, $\sigma = 2 * \arcsin(\sqrt{\sin^2(\frac{\Delta\phi}{2}) + \cos(\phi_1)\cos(\phi_2)\sin^2(\frac{\Delta\lambda}{2})})$

3.10 Heron's Formula

- $s = \frac{a+b+c}{2}$
- $Area = \sqrt{s(s-a)(s-b)(s-c)}$
- a, b, c there are the lengths of the sides

3.11 Interesting theorems

- $a^d \equiv a^{d \mod \phi(n)} \mod n$ if $a \in Z^{n_*}$ or $a \notin Z^{n_*}$ and $d \mod \phi(n) \neq 0$
- $a^d \equiv a^{\phi(n)} \mod n$ if $a \notin Z^{n_*}$ and $d \mod \phi(n) = 0$
- thus, for all a, n and d (with $d \ge \log_2(n)$) $a^d \equiv a^{\phi(n)+d \mod \phi(n)} \mod n$

3.12 Law of sines and cosines

- a, b, c: lengths, A, B, C: opposite angles, d: circumcircle
- $\frac{a}{\sin(A)} = \frac{b}{\sin(B)} = \frac{c}{\sin(C)} = d$
- $c^2 = a^2 + b^2 2ab\cos(C)$

3.13 Number of divisors

•
$$\tau(n) = \prod_{i=1}^{k} (\alpha_i + 1)$$

3.14 Product of divisors of a number

$$u(n) = n^{\frac{\tau(n)}{2}}$$

- if p is a prime, then: $\mu(p^k) = p^{\frac{k(k+2)}{2}}$
- if a and b are coprimes, then: $\mu(ab) = \mu(a)^{\tau(b)} \mu(b)^{\tau(a)}$

3.15 Pythagorean triples ($a^2 + b^2 = c^2$)

- Given an arbitrary pair of integers m and n with m > n > 0: $a = m^2 n^2$, b = 2mn, $c = m^2 + n^2$
- The triple generated by Euclid's formula is primitive if and only if m and n are coprime and not both odd.
- To generate all Pythagorean triples uniquely: $a = k(m^2 n^2), b = k(2mn), c = k(m^2 + n^2)$
- If m and n are two odd integer such that m > n, then: a = mn, $b = \frac{m^2 n^2}{2}$, $c = \frac{m^2 + n^2}{2}$
- If n = 1 or 2 there are no solutions. Otherwise n is even: $\left(\left(\frac{n^2}{4} 1\right)^2 + n^2 = \left(\frac{n^2}{4} + 1\right)^2\right)$ n is odd: $\left(\left(\frac{n^2 1}{2}\right)^2 + n^2 = \left(\frac{n^2 + 1}{2}\right)^2\right)$

3.16 Simplex Rules

The simplex algorithm operated on linear programs in standard form:

 $\mathbf{Maximixe}: c^T \cdot x$

Subject to: $Ax \leq b, x_i \geq 0$

- $x = (x_1, ..., x_n)$ the variables of the problem
- $c = (c_1, ..., c_n)$ are the coefficients of the objective function
- A is a $p \times n$ matrix and $b = (b_1, ..., b_p)$ constants with $b_j \ge 0$

3.17 Sum of divisors of a number

•
$$\sigma(n) = \prod_{i=1}^{k} (1 + p_i + \dots + p_i^{\alpha_i}) = \prod_{i=1}^{k} \frac{p_i^{\alpha_i + 1} - 1}{p_i - 1}$$

3.18 Summations

- $\bullet \ \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$
- $\sum_{i=1}^{n} i^3 = (\frac{n(n+1)}{2})^2$
- $\sum_{i=1}^{n} i^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$
- $\sum_{i=0}^n x^i = \frac{x^{n+1}-1}{x-1}$ para $x \neq 1$

3.19 Theorems

- There is always a prime between numbers n^2 and $(n+1)^2$, where n is any positive integer
- There is an infinite number of pairs of the from $\{p, p+2\}$ where both p and p+2 are primes.
- Every even integer greater than 2 can be expressed as the sum of two primes.
- Every integer greater than 2 can be written as the sum of three primes.

4 Geometry

4.1 3D

```
typedef double T;
struct p3 {
 T x, v, z;
 // Basic vector operations
 p3 operator + (p3 p) { return {x+p.x, y+p.y, z+p.z }; }
 p3 operator - (p3 p) { return {x - p.x, y - p.y, z - p.z}; }
 p3 operator * (T d) { return {x*d, y*d, z*d}; }
 p3 operator / (T d) { return {x / d, y / d, z / d}; } // only for floating point
 bool operator == (p3 p) { return tie(x, y, z) == tie(p.x, p.y, p.z); }
 bool operator != (p3 p) { return !operator == (p); }
p3 zero {0, 0, 0 };
T operator | (p3 v, p3 w) \{ /// dot \}
 return v.x*w.x + v.y*w.y + v.z*w.z;
p3 operator * (p3 v, p3 w) { /// cross
 return { v.y*w.z - v.z*w.y, v.z*w.x - v.x*w.z, v.x*w.y - v.y*w.x };
T sq(p3 v) { return v | v; }
double abs(p3 v) { return sqrt(sq(v)); }
p3 unit(p3 v) { return v / abs(v): }
double angle(p3 v, p3 w) {
 double cos theta = (v \mid w) / abs(v) / abs(w):
 return acos(max(-1.0, min(1.0, cos_theta)));
T orient(p3 p, p3 q, p3 r, p3 s) { /// orient s, pqr form a triangle
 return (q - p) * (r - p) | (s - p);
T orient_by_normal(p3 p, p3 q, p3 r, p3 n) { /// same as 2D but in n-normal direction
 return (q - p) * (r - p) | n;
struct plane {
 p3 n; T d;
 /// From normal n and offset d
  plane(p3 n, T d): n(n), d(d) {}
```

```
/// From normal n and point P
 plane(p3 n, p3 p): n(n), d(n | p) {}
 /// From three non-collinear points P,Q,R
 plane(p3 p, p3 q, p3 r): plane((q - p) * (r - p), p) \{\}
 /// - these work with T = int
 T side(p3 p) { return (n | p) - d; }
 double dist(p3 p) { return abs(side(p)) / abs(n); }
 plane translate(p3 t) {return \{n, d + (n \mid t)\}; }
 /// - these require T = double
 plane shift_up(double dist) { return {n, d + dist * abs(n)}; }
 p3 proj(p3 p) { return p - n * side(p) / sq(n); }
 p3 refl(p3 p) { return p - n * 2 * side(p) / sq(n); }
struct line3d {
 p3 d. o:
 /// From two points P. Q
 line3d(p3 p, p3 q): d(q - p), o(p) {}
 /// From two planes p1, p2 (requires T = double)
 line3d(plane p1, plane p2) {
   d = p1.n * p2.n;
   o = (p2.n * p1.d - p1.n * p2.d) * d / sq(d);
 /// - these work with T = int
 double sq_dist(p3 p) { return sq(d * (p - o)) / sq(d); }
 double dist(p3 p) { return sqrt(sq_dist(p)); }
 bool cmp_proj(p3 p, p3 q) { return (d | p) < (d | q); }</pre>
 /// - these require T = double
 p3 proj(p3 p) { return o + d * (d | (p - o)) / sq(d); }
 p3 refl(p3 p) { return proj(p) * 2 - p; }
 p3 inter(plane p) { return o - d * p.side(o) / (p.n | d); }
double dist(line3d 11, line3d 12) {
 p3 n = 11.d * 12.d:
 if(n == zero) // parallel
   return 11.dist(12.o):
 return abs((12.o - 11.o) | n) / abs(n):
p3 closest on line1(line3d 11, line3d 12) { /// closest point on 11 to 12
 p3 n2 = 12.d * (11.d * 12.d);
 return 11.0 + 11.d * ((12.0 - 11.0) | n2) / (11.d | n2):
double small angle(p3 v, p3 w) { return acos(min(abs(v | w) / abs(v) / abs(w), 1.0)); }
double angle(plane p1, plane p2) { return small_angle(p1.n, p2.n); }
bool is parallel(plane p1, plane p2) { return p1.n * p2.n == zero; }
bool is_perpendicular(plane p1, plane p2) { return (p1.n | p2.n) == 0; }
double angle(line3d 11, line3d 12) { return small_angle(11.d, 12.d); }
bool is_parallel(line3d 11, line3d 12) { return 11.d * 12.d == zero; }
bool is_perpendicular(line3d 11, line3d 12) { return (11.d | 12.d) == 0; }
double angle(plane p, line3d 1) { return _pI / 2 - small_angle(p.n, 1.d); }
bool is_parallel(plane p, line3d 1) { return (p.n | 1.d) == 0; }
bool is_perpendicular(plane p, line3d 1) { return p.n * 1.d == zero; }
```

```
line3d perp_through(plane p, p3 o) { return line(o, o + p.n); }
plane perp_through(line3d 1, p3 o) { return plane(1.d, o); }
```

4.2 General

```
const lf eps = 1e-9:
typedef double T:
struct pt {
 T x, v:
 pt operator + (pt p) { return {x+p.x. v+p.v}: }
 pt operator - (pt p) { return {x-p.x, y-p.y}; }
 pt operator * (pt p) { return {x*p.x-y*p.y, x*p.y+y*p.x}; }
 pt operator * (T d) { return {x*d, y*d}; }
 pt operator / (lf d) { return {x/d, y/d}; } /// only for floating point
 bool operator == (pt b) { return x == b.x && y == b.y; }
 bool operator != (pt b) { return !(*this == b); }
 bool operator < (const pt &o) const { return v < o.v || (v == o.v && x < o.x); }
 bool operator > (const pt &o) const { return y > o.y || (y == o.y && x > o.x); }
int cmp (lf a, lf b) { return (a + eps < b ? -1 : (b + eps < a ? 1 : 0)); }
/** Already in complex **/
T norm(pt a) { return a.x*a.x + a.y*a.y; }
lf abs(pt a) { return sqrt(norm(a)); }
lf arg(pt a) { return atan2(a.y, a.x); }
ostream& operator << (ostream& os, pt &p) {
 return os << "("<< p.x << "," << p.y << ")";
istream & operator >> (istream & in, pt & p) {
   T x. v: in >> x >> v:
   p = \{x, y\};
   return in:
T dot(pt a, pt b) { return a.x*b.x + a.y*b.y; }
T cross(pt a, pt b) { return a.x*b.y - a.y*b.x; }
T orient(pt a, pt b, pt c) { return cross(b-a,c-a); }
//pt rot(pt p, lf a) { return {p.x*cos(a) - p.v*sin(a), p.x*sin(a) + p.v*cos(a)}; }
//pt rot(pt p, double a) { return p * polar(1.0, a); } /// for complex
//pt rotate to b(pt a, pt b, lf ang) { return rot(a-b, ang)+b; }
pt rot90ccw(pt p) { return {-p.y, p.x}; }
pt rot90cw(pt p) { return {p.v. -p.x}: }
pt translate(pt p, pt v) { return p+v; }
pt scale(pt p, double f, pt c) { return c + (p-c)*f; }
bool are_perp(pt v, pt w) { return dot(v,w) == 0; }
int sign(T x) \{ return (T(0) < x) - (x < T(0)); \}
pt unit(pt a) { return a/abs(a); }
bool in_angle(pt a, pt b, pt c, pt x) {
 assert(orient(a,b,c) != 0);
 if (orient(a,b,c) < 0) swap(b,c);
 return orient(a,b,x) >= 0 && orient(a,c,x) <= 0;
```

```
//If angle(pt a, pt b) { return acos(max(-1.0, min(1.0, dot(a,b)/abs(a)/abs(b)))); }
//lf angle(pt a, pt b) { return atan2(cross(a, b), dot(a, b)); }
/// returns vector to transform points
pt get_linear_transformation(pt p, pt q, pt r, pt fp, pt fq) {
 pt pq = q-p, num{cross(pq, fq-fp), dot(pq, fq-fp)};
 return fp + pt{cross(r-p, num), dot(r-p, num)} / norm(pq);
bool half(pt p) { /// true if is in (0, 180]
 assert(p.x != 0 || p.y != 0); /// the argument of (0,0) is undefined
 return p.y > 0 || (p.y == 0 && p.x < 0);
bool half_from(pt p, pt v = {1, 0}) {
 return cross(v,p) < 0 || (cross(v,p) == 0 && dot(v,p) < 0);
bool polar_cmp(const pt &a, const pt &b) {
 return make tuple(half(a), 0) < make tuple(half(b), cross(a,b));
struct line {
 pt v; T c;
 line(pt v, T c) : v(v), c(c) {}
 line(T a, T b, T c) : v(\{b,-a\}), c(c) {}
 line(pt p, pt q) : v(q-p), c(cross(v,p)) {}
 T side(pt p) { return cross(v,p)-c; }
 lf dist(pt p) { return abs(side(p)) / abs(v); }
 lf sq_dist(pt p) { return side(p)*side(p) / (lf)norm(v); }
 line perp_through(pt p) { return {p, p + rot90ccw(v)}; }
 bool cmp_proj(pt p, pt q) { return dot(v,p) < dot(v,q); }</pre>
 line translate(pt t) { return {v, c + cross(v,t)}; }
 line shift_left(double d) { return {v, c + d*abs(v)}; }
 pt proj(pt p) { return p - rot90ccw(v)*side(p)/norm(v); }
 pt refl(pt p) { return p - rot90ccw(v)*2*side(p)/norm(v); }
}:
bool inter 11(line 11, line 12, pt &out) {
 T d = cross(11.v. 12.v):
 if (d == 0) return false:
 out = (12.v*11.c - 11.v*12.c) / d:
 return true:
line bisector(line 11, line 12, bool interior) {
 assert(cross(11.v, 12.v) != 0); /// 11 and 12 cannot be parallel!
 lf sign = interior ? 1 : -1:
 return {12.v/abs(12.v) + 11.v/abs(11.v) * sign,
         12.c/abs(12.v) + 11.c/abs(11.v) * sign};
}
bool in_disk(pt a, pt b, pt p) {
 return dot(a-p, b-p) <= 0;</pre>
```

```
bool on_segment(pt a, pt b, pt p) {
 return orient(a,b,p) == 0 && in_disk(a,b,p);
bool proper_inter(pt a, pt b, pt c, pt d, pt &out) {
 T oa = orient(c,d,a),
 ob = orient(c,d,b),
 oc = orient(a,b,c),
 od = orient(a,b,d);
 /// Proper intersection exists iff opposite signs
 if (oa*ob < 0 && oc*od < 0) {
   out = (a*ob - b*oa) / (ob-oa):
   return true:
 return false:
set<pt> inter ss(pt a, pt b, pt c, pt d) {
 if (proper inter(a,b,c,d,out)) return {out}:
 set<pt> s:
 if (on_segment(c,d,a)) s.insert(a);
 if (on segment(c.d.b)) s.insert(b);
 if (on_segment(a,b,c)) s.insert(c);
 if (on_segment(a,b,d)) s.insert(d);
 return s;
lf pt_to_seg(pt a, pt b, pt p) {
 if(a != b) {
   line l(a,b);
   if (l.cmp_proj(a,p) && l.cmp_proj(p,b)) /// if closest to projection
     return l.dist(p); /// output distance to line
 return min(abs(p-a), abs(p-b)); /// otherwise distance to A or B
lf seg_to_seg(pt a, pt b, pt c, pt d) {
 pt dummy:
 if (proper inter(a,b,c,d,dummy)) return 0:
 return min({pt_to_seg(a,b,c), pt_to_seg(a,b,d),
            pt to seg(c.d.a), pt to seg(c.d.b)}):
enum {IN, OUT, ON};
struct polygon {
 vector<pt> p:
 polygon(int n) : p(n) {}
 int top = -1, bottom = -1;
 void delete repetead() {
   vector<pt> aux;
   sort(p.begin(), p.end());
   for(pt &i : p)
     if(aux.empty() || aux.back() != i)
       aux.push_back(i);
   p.swap(aux);
```

```
12
```

```
bool is_convex() {
  bool pos = 0, neg = 0;
  for (int i = 0, n = p.size(); i < n; i++) {</pre>
   int o = orient(p[i], p[(i+1)%n], p[(i+2)%n]);
   if (o > 0) pos = 1;
   if (o < 0) neg = 1;
  return ! (pos && neg);
lf area(bool s = false) {
  lf ans = 0:
  for (int i = 0, n = p.size(); i < n; i++)</pre>
    ans += cross(p[i], p[(i+1)%n]);
  ans \neq 2:
  return s ? ans : abs(ans):
lf perimeter() {
  lf per = 0:
  for(int i = 0, n = p.size(); i < n; i++)
    per += abs(p[i] - p[(i+1)\%n]);
  return per;
bool above(pt a, pt p) { return p.y >= a.y; }
bool crosses_ray(pt a, pt p, pt q) {
  return (above(a,q)-above(a,p))*orient(a,p,q) > 0;
int in_polygon(pt a) {
  int crosses = 0:
  for(int i = 0, n = p.size(); i < n; i++) {</pre>
    if(on_segment(p[i], p[(i+1)%n], a)) return ON;
    crosses += crosses_ray(a, p[i], p[(i+1)%n]);
  return (crosses&1 ? IN : OUT);
void normalize() { /// polygon is CCW
  bottom = min element(p.begin(), p.end()) - p.begin();
  vector<pt> tmp(p.begin()+bottom, p.end());
  tmp.insert(tmp.end(), p.begin(), p.begin()+bottom);
  p.swap(tmp);
  bottom = 0:
  top = max_element(p.begin(), p.end()) - p.begin();
int in convex(pt a) {
  assert(bottom == 0 \&\& top != -1):
  if(a < p[0] || a > p[top]) return OUT;
  T orientation = orient(p[0], p[top], a);
  if(orientation == 0) {
    if(a == p[0] || a == p[top]) return ON;
   return top == 1 || top + 1 == p.size() ? ON : IN;
  } else if (orientation < 0) {</pre>
    auto it = lower_bound(p.begin()+1, p.begin()+top, a);
   T d = orient(*prev(it), a, *it);
    return d < 0? IN : (d > 0 ? OUT: ON):
```

```
}
 else {
   auto it = upper_bound(p.rbegin(), p.rend()-top-1, a);
   T d = orient(*it, a, it == p.rbegin() ? p[0] : *prev(it));
   return d < 0 ? IN : (d > 0 ? OUT: ON);
polygon cut(pt a, pt b) {
 line 1(a, b):
 polygon new_polygon(0);
 for(int i = 0, n = p.size(); i < n; ++i) {
   pt c = p[i], d = p[(i+1)\%n];
   lf abc = cross(b-a, c-a), abd = cross(b-a, d-a);
   if(abc >= 0) new_polygon.p.push_back(c);
   if(abc*abd < 0) {
     pt out: inter ll(1, line(c, d), out):
     new_polygon.p.push_back(out);
 return new_polygon;
void convex_hull() {
 sort(p.begin(), p.end());
 vector<pt> ch;
 ch.reserve(p.size()+1);
 for(int it = 0; it < 2; it++) {</pre>
   int start = ch.size();
   for(auto &a : p) {
     /// if colineal are needed, use < and remove repeated points
     while(ch.size() >= start+2 && orient(ch[ch.size()-2], ch.back(), a) <= 0)</pre>
       ch.pop_back();
     ch.push_back(a);
   ch.pop_back();
   reverse(p.begin(), p.end());
 if(ch.size() == 2 \&\& ch[0] == ch[1]) ch.pop back();
 /// be careful with CH of size < 3
 p.swap(ch);
vector<pii> antipodal() {
 vector<pii> ans:
 int n = p.size():
 if (n == 2) ans.push back(\{0, 1\}):
 if(n < 3) return ans:</pre>
 auto nxt = [\&](int x) \{ return (x+1 == n ? 0 : x+1): \}:
 auto area2 = [&](pt a, pt b, pt c) { return cross(b-a, c-a); };
 while (abs(area2(p[n-1], p[0], p[nxt(b0)])) >
       abs(area2(p[n - 1], p[0], p[b0])))
 for(int b = b0, a = 0; b != 0 && a <= b0; ++a) {
   ans.push_back({a, b});
```

```
while (abs(area2(p[a], p[nxt(a)], p[nxt(b)])) >
            abs(area2(p[a], p[nxt(a)], p[b]))) {
       b = nxt(b):
       if(a != b0 || b != 0) ans.push_back({ a, b });
       else return ans;
     if(abs(area2(p[a], p[nxt(a)], p[nxt(b)])) ==
        abs(area2(p[a], p[nxt(a)], p[b]))) {
       if(a != b0 || b != n-1) ans.push_back({ a, nxt(b) });
       else ans.push back({ nxt(a), b }):
   }
   return ans:
 pt centroid() {
   pt c{0, 0}:
   lf scale = 6. * area(true):
   for(int i = 0, n = p.size(); i < n; ++i) {</pre>
     int i = (i+1 == n ? 0 : i+1):
     c = c + (p[i] + p[j]) * cross(p[i], p[j]);
   return c / scale;
 11 pick() {
   11 boundary = 0;
   for(int i = 0, n = p.size(); i < n; i++) {</pre>
     int j = (i+1 == n ? 0 : i+1);
     boundary += _-gcd((ll)abs(p[i].x - p[j].x), (ll)abs(p[i].y - p[j].y));
   return area() + 1 - boundary/2;
 pt& operator[] (int i){ return p[i]; }
struct circle {
 pt c; T r;
}:
circle center(pt a, pt b, pt c) {
 b = b-a, c = c-a:
 assert(cross(b,c) != 0); /// no circumcircle if A,B,C aligned
 pt cen = a + rot90ccw(b*norm(c) - c*norm(b))/cross(b,c)/2:
 return {cen. abs(a-cen)}:
int inter_cl(circle c, line l, pair<pt, pt> &out) {
 1f h2 = c.r*c.r - 1.sq dist(c.c):
 if(h2 >= 0) {
   pt p = 1.proj(c.c);
   pt h = 1.v*sqrt(h2)/abs(1.v);
   out = \{p-h, p+h\};
 return 1 + sign(h2);
```

```
int inter_cc(circle c1, circle c2, pair<pt, pt> &out) {
 pt d=c2.c-c1.c; double d2=norm(d);
 if(d2 == 0) { assert(c1.r != c2.r); return 0; } // concentric circles
 double pd = (d2 + c1.r*c1.r - c2.r*c2.r)/2; // = |0_1P| * d
 double h2 = c1.r*c1.r - pd*pd/d2; // = h2
 if(h2 >= 0) {
   pt p = c1.c + d*pd/d2, h = rot90ccw(d)*sqrt(h2/d2);
   out = \{p-h, p+h\};
 return 1 + sign(h2);
int tangents(circle c1. circle c2. bool inner, vector<pair<pt,pt>> &out) {
 if(inner) c2.r = -c2.r:
 pt d = c2.c-c1.c;
 double dr = c1.r-c2.r, d2 = norm(d), h2 = d2-dr*dr:
 if(d2 == 0 || h2 < 0) { assert(h2 != 0); return 0; }</pre>
 for(double s : {-1.1}) {
   pt v = (d*dr + rot90ccw(d)*sart(h2)*s)/d2:
   out.push_back({c1.c + v*c1.r, c2.c + v*c2.r});
 return 1 + (h2 > 0);
int tangent_through_pt(pt p, circle c, pair<pt, pt> &out) {
 double d = abs(p - c.c);
       if(d < c.r) return 0;</pre>
 pt base = c.c-p;
 double w = sqrt(norm(base) - c.r*c.r);
 pt a = \{w, c.r\}, b = \{w, -c.r\};
 pt s = p + base*a/norm(base)*w;
 pt t = p + base*b/norm(base)*w;
 out = \{s, t\};
 return 1 + (abs(c.c-p) == c.r);
```

4.3 nonTested

```
lf part(pt a, pt b, T r) {
    lf l = abs(a-b);
    pt p = (b-a)/l;
    lf c = dot(a, p), d = 4.0 * (c*c - dot(a, a) + r*r);
    if (d < eps) return angle(a, b) * r * r * 0.5;
    d = sqrt(d) * 0.5;
    lf s = -c - d, t = -c + d;
    if (s < 0.0) s = 0.0; else if (s > l) s = l;
    if (t < 0.0) t = 0.0; else if (t > l) t = l;
    pt u = a + p*s, v = a + p*t;
    return (cross(u, v) + (angle(a, u) + angle(v, b)) * r * r) * 0.5;
}
```

```
lf inter_cp(circle c, polygon p) {
 lf ans = 0;
 int n = p.p.size();
 for (int i = 0; i < n; i++) {
   ans += part(p[i]-c.c, p[(i+1)%4]-c.c, c.r);
 return abs(ans);
}
struct circle{
   point center; double r;
   bool contain(point &p) { return abs(center - p) < r + eps:}
};
T cross(point a, point b) { return a.x*b.y - a.y*b.x; }
point rot90ccw(point p) { return {-p.y, p.x}; }
point get(point a, point b, point c) {
 b = b-a, c = c-a;
 //assert(cross(b,c) != 0): /// no circumcircle if A.B.C aligned
 point cen = a + rot90ccw(b*norm(c) - c*norm(b))/cross(b.c)/2:
 return cen:
}
circle min_circle(vector<point> &cloud, int a, int b){
   point center = (cloud[a] + cloud[b]) / double(2.);
   double rat = abs(center - cloud[a]);
   circle C = {center, rat};
   for (int i = 0; i < b; ++i){
       point x = cloud[i];
       if (C.contain(x)) continue;
       center = get( cloud[a], cloud[b], cloud[i] );
       rat = abs(center - cloud[a]);
       C = {center, rat};
   return C;
}
circle min circle(vector<point> &cloud, int a){
   point center = (cloud[a] + cloud[0]) / double(2.);
   double rat = abs(center - cloud[a]):
   circle C = {center. rat}:
   for (int i = 0; i < a; ++i){</pre>
       point x = cloud[i];
       if (C.contain(x)) continue:
       C = min circle(cloud, a, i):
   return C:
}
circle min_circle(vector<point> cloud){
   // random_shuffle(cloud.begin(), cloud.end());
   int n = (int)cloud.size();
   for (int i = 1; i < n; ++i){
       int u = rand() % i;
       swap(cloud[u], cloud[i]);
```

```
point center = (cloud[0] + cloud[1]) / double(2.);
double rat = abs(center - cloud[0]);
circle C = {center, rat};
for (int i = 2; i < n; ++i){
    point x = cloud[i];
    if (C.contain(x)) continue;
    C = min_circle(cloud, i);
}
return C;
}</pre>
```

5 Graphs

5.1 2-satisfiability

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/y8qhbzn4
struct sat2 {
 int n;
 vector<vector<int>>> g;
 vector<int> tag;
 vector<bool> seen, value;
 stack<int> st;
 sat2(int n) : n(n), g(2, vector < vector < int >> (2*n)), tag(2*n), seen(2*n), value(2*n) {
 int neg(int x) { return 2*n-x-1; }
 void add_or(int u, int v) { implication(neg(u), v); }
 void make_true(int u) { add_edge(neg(u), u); }
 void make_false(int u) { make_true(neg(u)); }
 void eq(int u, int v) {
   implication(u, v):
   implication(v, u);
 void diff(int u, int v) { eq(u, neg(v)); }
 void implication(int u, int v) {
   add edge(u. v):
   add_edge(neg(v), neg(u));
 void add_edge(int u, int v) {
   g[0][u].push_back(v);
   g[1][v].push_back(u);
 void dfs(int id, int u, int t = 0) {
   seen[u] = true;
   for(auto& v : g[id][u])
     if(!seen[v])
       dfs(id, v, t);
   if(id == 0) st.push(u);
   else tag[u] = t;
```

```
void kosaraju() {
   for(int u = 0; u < n; u++) {</pre>
     if(!seen[u]) dfs(0, u);
     if(!seen[neg(u)]) dfs(0, neg(u));
   fill(seen.begin(), seen.end(), false);
   int t = 0;
   while(!st.empty()) {
     int u = st.top(); st.pop();
     if(!seen[u]) dfs(1, u, t++);
 bool satisfiable() {
   kosaraju();
   for(int i = 0: i < n: i++) {</pre>
     if(tag[i] == tag[neg(i)]) return false;
     value[i] = tag[i] > tag[neg(i)];
   return true;
 }
};
```

5.2 Erdos-Gallai theorem

```
/// Complexity: O(|N|*log|N|)
/// Tested: https://tinyurl.com/yb5v9bau
/// Theorem: it gives a necessary and sufficient condition for a finite sequence
           of natural numbers to be the degree sequence of a simple graph
bool erdos(vector<int> &d) {
 11 sum = 0:
 for(int i = 0; i < d.size(); ++i) sum += d[i];</pre>
 if(sum & 1) return false;
 sort(d.rbegin(), d.rend());
 11 1 = 0, r = 0;
 for(int k = 1, i = d.size() - 1; k <= d.size(); ++k) {</pre>
   1 += d[k-1];
   if(k > i) r = d[++i];
   while (i >= k && d[i] < k+1) r += d[i--];
   if(1 > 111*k*(k-1) + 111*k*(i-k+1) + r)
     return false:
 }
 return true;
```

5.3 Eulerian path

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/y85t8e83
```

```
bool eulerian(vector<int> &tour) { /// directed graph
 int one_in = 0, one_out = 0, start = -1;
 bool ok = true;
 for (int i = 0; i < n; i++) {</pre>
   if(out[i] && start == -1) start = i;
   if(out[i] - in[i] == 1) one_out++, start = i;
   else if(in[i] - out[i] == 1) one_in++;
   else ok &= in[i] == out[i];
 ok &= one in == one out && one in <= 1:
 if (ok) {
   function<void(int)> go = [&](int u) {
     while(g[u].size()) {
      int v = g[u].back();
       g[u].pop_back();
      go(v);
     tour.push_back(u);
   };
   go(start);
   reverse(tour.begin(), tour.end());
   if(tour.size() == edges + 1) return true;
 return false;
```

15

5.4 Lowest common ancestor

```
/// Complexity: O(|N|*log|N|)
/// Tested: https://tinyurl.com/y9g2ljv9, https://tinyurl.com/y87q3j93
int lca(int a, int b) {
 if(depth[a] < depth[b]) swap(a, b);</pre>
 //int ans = INT MAX:
 for(int i = LOG2-1: i >= 0: --i)
   if(depth[ dp[a][i] ] >= depth[b]) {
     //ans = min(ans, mn[a][i]):
     a = dp[a][i]:
 //if (a == b) return ans:
 if(a == b) return a;
 for(int i = LOG2-1: i >= 0: --i)
   if(dp[a][i] != dp[b][i]) {
     //ans = min(ans, mn[a][i]);
     //ans = min(ans, mn[b][i]);
     a = dp[a][i],
     b = dp[b][i];
 //ans = min(ans, mn[a][0]);
 //ans = min(ans, mn[b][0]);
 //return ans;
 return dp[a][0];
```

```
void dfs(int u, int d = 1, int p = -1) {
 depth[u] = d;
 for(auto v : g[u]) {
   //int v = x.first;
   //int w = x.second;
   if(v != p) {
     dfs(v, d + 1, u);
     dp[v][0] = u;
     //mn[v][0] = w;
 }
}
void build(int n) {
 for(int i = 0; i <= n; i++) dp[i][0] = -1;
 for(int i = 0: i < n, i++) {
   if(dp[i][0] == -1) {
     dp[i][0] = i;
     //mn[i][O] = INT_MAX;
     dfs(i);
 for(int j = 0; j < LOG2-1; ++j)</pre>
   for(int i = 0; i <= n; ++i) { // nodes</pre>
     dp[i][j+1] = dp[dp[i][j]][j];
     //mn[i][j+1] = min(mn[ dp[i][j] ][j], mn[i][j]);
}
```

5.5 Number of spanning trees

```
/// Tested: not yet
///A -> adjacency matrix
///It is necessary to compute the D-A matrix, where D is a diagonal matrix
///that contains the degree of each node.
///To compute the number of spanning trees it's necessary to compute any
///D-A cofactor
///C(i, i) = (-1)^(i+i) * Mii
///Where Mij is the matrix determinant after removing row i and column j
double mat[MAX][MAX]:
///call determinant(n - 1)
double determinant(int n) {
 double det = 1.0;
 for(int k = 0; k < n; k++) {
   for(int i = k+1; i < n; i++) {</pre>
     assert(mat[k][k] != 0);
     long double factor = mat[i][k]/mat[k][k];
     for(int j = 0; j < n; j++) {</pre>
       mat[i][j] = mat[i][j] - factor*mat[k][j];
```

```
}
  det *= mat[k][k];
}
return round(det);
}
```

5.6 Scc.

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/y8ujj3ws
int scc(int n) {
 vector<int> dfn(n+1), low(n+1), in_stack(n+1);
 stack<int> st;
 int tag = 0;
 function<void(int, int&)> dfs = [&](int u, int &t) {
   dfn[u] = low[u] = ++t;
   st.push(u);
   in_stack[u] = true;
   for(auto &v : g[u]) {
     if(!dfn[v]) {
       dfs(v. t):
       low[u] = min(low[u], low[v]);
     } else if(in stack[v])
       low[u] = min(low[u], dfn[v]);
   if (low[u] == dfn[u]) {
     int v:
     do {
      v = st.top(); st.pop();
       id[v] = tag:
       in_stack[v] = false;
     } while (v != u);
     tag++;
   }
 for(int u = 1, t; u <= n; ++u) {</pre>
   if(!dfn[u]) dfs(u, t = 0);
 return tag;
```

5.7 Tarjan tree

```
/// Complexity: 0(|N|)
/// Tested: https://tinyurl.com/y9g2ljv9, https://tinyurl.com/y87q3j93
struct tarjan_tree {
  int n;
  vector<vector<int>> g, comps;
```

```
vector<pii> bridge;
vector<int> id, art;
tarjan_tree(int n) : n(n), g(n+1), id(n+1), art(n+1) {}
void add_edge(vector<vector<int>> &g, int u, int v) { /// nodes from [1, n]
 g[u].push_back(v);
 g[v].push_back(u);
void add_edge(int u, int v) { add_edge(g, u, v); }
void tarian(bool with bridge) {
 vector<int> dfn(n+1), low(n+1);
 stack<int> st:
 comps.clear();
 function < void(int, int, int&) > dfs = [&](int u, int p, int &t) {
   dfn[u] = low[u] = ++t:
   st.push(u):
   int cntp = 0:
   for(int v : g[u]) {
     cntp += v == p:
     if(!dfn[v]) {
       dfs(v. u. t):
       low[u] = min(low[u], low[v]);
       if(with_bridge && low[v] > dfn[u]) {
         bridge.push_back({min(u,v), max(u,v)});
         comps.push_back({});
         for(int w = -1; w != v; )
          comps.back().push_back(w = st.top()), st.pop();
       if(!with_bridge && low[v] >= dfn[u]) {
         art[u] = (dfn[u] > 1 || dfn[v] > 2);
         comps.push_back({u});
         for(int w = -1; w != v; )
           comps.back().push_back(w = st.top()), st.pop();
     else if (v != p || cntp > 1) low[u] = min(low[u], dfn[v]):
   if(p == -1 && ( with_bridge || g[u].size() == 0 )) {
     comps.push back({}):
     for(int w = -1; w != u;)
       comps.back().push_back(w = st.top()), st.pop();
 };
 for(int u = 1, t: u \le n: ++u)
   if(!dfn[u]) dfs(u, -1, t = 0):
vector<vector<int>> build block cut tree() {
 tarjan(false);
 int t = 0:
 for(int u = 1; u <= n; ++u)</pre>
   if(art[u]) id[u] = t++;
 vector<vector<int>> tree(t+comps.size());
 for(int i = 0; i < comps.size(); ++i)</pre>
   for(int u : comps[i]) {
```

```
if(!art[u]) id[u] = i+t;
    else add_edge(tree, i+t, id[u]);
}
return tree;
}
vector<vector<int>> build_bridge_tree() {
    tarjan(true);
    vector<vector<int>> tree(comps.size());
    for(int i = 0; i < comps.size(); ++i)
        for(int u : comps[i]) id[u] = i;
    for(auto &b : bridge)
        add_edge(tree, id[b.first], id[b.second]);
    return tree;
}
};</pre>
```

5.8 Tree binarization

```
/// Complexity: O(|N|)
/// Tested: Not yet
void add(int u, int v, int w) { ng[u].push_back({v, w}); }
void binarize(int u, int p = -1) {
    int last = u, f = 0;
    for(auto x : g[u]) {
        int v = x.first, w = x.second, node = ng.size();
        if(v == p) continue;
        if(f++) {
            ng.push_back({});
            add(last, node, 0);
            add(node, v, w);
            last = node;
        } else add(u, v, w);
        binarize(v, u);
    }
}
```

5.9 Yen

```
/// Complexity: 0( |K|*|N|^3 )
/// Tested: not yet
int n;
vector<int> graph[ MAXN ];
int cost[ MAXN ][ MAXN ], dist[ MAXN ], connect[ MAXP ], path[ MAXN ];
ll vis = 0, mark = 0, edge[ MAXN ];
vector<int> emp;
struct Path {
  int w;
  vector<int> p;
```

```
Path(): w(0) {}
  Path( int w ) : w(w) { }
 Path( int w, vector<int> p ) : w(w), p(p) { }
  bool operator < ( const Path& other )const {</pre>
   if( w == other.w ) {
     return lexicographical_compare( p.begin(), p.end(), other.p.begin(), other.p.end()
   return w < other.w:</pre>
  bool operator > ( const Path& other )const {
    if( w == other.w ){
     return lexicographical_compare( other.p.begin(), other.p.end(), p.begin(), p.end()
   return w > other.w:
};
void add_edge( int u, int v, int w ) {
  cost[u][v] = w:
  edge[u] |= ( 1LL<<v );
 graph[u].push_back( v );
Path dijkstra( int s, int t ) {
  priority_queue< pii, vector<pii>, greater<pii> > pq;
 fill( dist, dist+n+1, INF );
  pq.push( {0,s} );
  dist[s] = 0;
  while( !pq.empty() ) {
    int u = pq.top().second, c = pq.top().first;
    pq.pop();
    if( u == t ) break;
    if( ((vis>>u)&1) && s != u )
     continue:
    vis |= ( 1LL<<u ):</pre>
    for( int i = 0; i < graph[u].size(); ++i ) {</pre>
     int v = graph[u][i];
     if( ((vis>>v)&1) || ( s == u && !((mark>>v)&1)) ) {
       continue:
     if( cost[u][v] != INF && dist[v] >= c+cost[u][v] ) {
       if( dist[v] > c+cost[u][v] || ( dist[v] == c+cost[u][v] && u < path[v] ) ) {</pre>
         dist[v] = c+cost[u][v];
         path[v] = u:
         pq.push( {dist[v], v} );
  if( dist[t] == INF ) {
    return Path();
```

```
Path ret( dist[t] );
 for( int u = t; u != s; u = path[u] ) {
   ret.p.push_back( u );
 ret.p.push_back( s );
 reverse( ret.p.begin(), ret.p.end() );
 return ret;
vector<int> ven( int s, int t, int k ) {
 priority_queue< Path, vector<Path>, greater<Path> > B;
 vector<vector<int>> A( MAXP );
 vis = 0:
 mark = edge[s]:
 A[0] = dijkstra(s, t).p:
 if( A[0].size() == 0 ) {
   return A[0]:
 for( int it = 1: it < k: ++it ){</pre>
   Path root_path;
   memset( connect, -1, sizeof(connect) );
   vis = 0:
   bool F = true;
   for( int i = 0; i < A[it-1].size()-1; ++i ) {</pre>
     bool flag = false;
     if( F && it > 2 && A[it-1].size() > i+1 &&
         A[it-2].size() > i+1 && A[it-1][i+1] == A[it-2][i+1] ) flag = true;
     else F = false:
     if( i >= A[it-1].size()-1 ) continue;
     int spur_node = A[it-1][i];
     mark = edge[ spur_node ];
     root_path.w += ( i ? cost[ A[it-1][i-1] ][ spur_node ] : 0 );
     root_path.p.push_back( spur_node );
     vis |= ( 1LL<<spur node ):</pre>
     for( int i = 0: i < it: ++i ) {
       if( connect[j] == i-1 && A[j].size() > i && A[j][i] == spur_node ) {
         connect[i] = i:
        if( A[j].size() > i+1 ) {
          mark &= ~( 1LL<<A[j][i+1] );
      }
     if( flag ) continue:
     ll prev vis = vis:
     Path spur_path = dijkstra( spur_node, t );
     vis = prev_vis;
     if( spur_path.p.empty() ) continue;
     Path cur_path = root_path;
     cur_path.w += spur_path.w;
     for( int j = 1; j < spur_path.p.size(); ++j ) {</pre>
       cur_path.p.push_back( spur_path.p[j] );
```

m coUNter

```
B.push( cur_path );
}
if( B.empty() ) return emp;
A[ it ] = B.top().p;
while( !B.empty() && B.top().p == A[it] ) {
    B.pop();
}
}
return A[ k-1 ];
```

6 Math

6.1 Berlekamp-Massey

```
#include<bits/stdc++.h>
using namespace std;
const int mod = 998244353:
inline int pw(int a, int b) {
 int ans = 1:
 while (b) {
   if (b & 1) ans = 1 LL * ans * a % mod;
   a = 1 LL * a * a % mod;
   b >>= 1:
 return ans:
namespace linear_seq {
 int m:
 // a = first m terms
 // p = dependence, length is m
 vector < int > p. a:
 inline vector < int > BM(vector < int > x) { // finds shortest linear recurrence
      given first x terms in O(x^2)
   //ls = last s' recurrence
   vector < int > ls, cur;
   //ld = last t' found
   //lf delta of last found
   int lf, ld;
   for (int i = 0; i < (int) x.size(); ++i) {</pre>
     int t = 0:
     //evaluate at position i
     for (int j = 0; j < (int) cur.size(); ++j)</pre>
      t = (t + 1 LL * x[i - j - 1] * cur[j]) \% mod;
```

```
if ((t - x[i]) \% \text{ mod} == 0) continue;
   if (!cur.size()) { //first non-zero element
     cur.resize(i + 1);
     lf = i;
     1d = (t - x[i]) \% mod;
     continue;
   int k = 1 LL * (t - x[i]) * pw(ld, mod - 2) % mod;
   vector < int > c(i - lf - 1); //add zeroes in front
   c.push back(k): //add '1'
   for (int j = 0; j < (int) ls.size(); ++j) //add minus previous s'</pre>
     c.push back(-1 LL * ls[i] * k % mod);
   if (c.size() < cur.size()) c.resize(cur.size());</pre>
   for (int j = 0; j < (int) cur.size(); ++j)</pre>
     c[j] = (c[j] + cur[j]) \% mod;
   if (i + lf + (int) ls.size() >= (int) cur.size())
     ls = cur. lf = i, ld = (t - x[i]) \% mod;
   cur = c;
 }
 for (int i = 0; i < (int) cur.size(); ++i)</pre>
   cur[i] = (cur[i] % mod + mod) % mod;
 m = cur.size():
 p.resize(m), a.resize(m);
 for (int i = 0; i < m; ++i)</pre>
   p[i] = cur[i], a[i] = x[i];
 return cur;
inline vector < int > mul(vector < int > & a, vector < int > & b) { // a * b mod f; f
    = x ** m - sum{1..m} x**(m-i) * pi
 //mav be optimized using FFT. NTT
 vector < int > r(2 * m);
 for (int i = 0: i < m: ++i)
   if (a[i])
     for (int j = 0; j < m; ++j)
      r[i + j] = (r[i + j] + 1 LL * a[i] * b[j]) % mod;
 for (int i = 2 * m - 1; i \ge m; --i)
   if (r[i])
    for (int j = m - 1; j \ge 0; ---j)
      r[i - j - 1] = (r[i - j - 1] + 1 LL * p[j] * r[i]) % mod;
 r.resize(m);
 return r;
// O(m*m*log(k)) with Fourier O(m*log(m)*log(k))
inline int calc(long long k) { // res = G[ x**k ] = G[ x ** k mod f]
```

if (m == 0) return 0;

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```
vector < int > bs(m), r(m);

if (m == 1) bs[0] = p[0];
else bs[1] = 1;

r[0] = 1;

while (k) {
   if (k & 1) r = mul(r, bs);
   bs = mul(bs, bs);
   k >>= 1;
}
int res = 0;
for (int i = 0; i < m; ++i)
   res = (res + 1 LL * r[i] * a[i]) % mod;
return res;
}
</pre>
```

6.2 Chinese remainder theorem

```
/// Complexity: |N|*log(|N|)
/// Tested: Not yet.
/// finds a suitable x that meets: x is congruent to a i mod n i
/** Works for non-coprime moduli.
Returns {-1,-1} if solution does not exist or input is invalid.
Otherwise, returns \{x,L\}, where x is the solution unique to mod L = LCM of mods
pair<int, int> chinese_remainder_theorem( vector<int> A, vector<int> M ) {
 int n = A.size(), a1 = A[0], m1 = M[0];
 for(int i = 1; i < n; i++) {</pre>
   int a2 = A[i], m2 = M[i];
   int g = \_gcd(m1, m2);
   if( a1 % g != a2 % g ) return {-1,-1};
   int p, q;
   eea(m1/g, m2/g, &p, &q);
   int mod = m1 / g * m2;
   q %= mod; p %= mod;
   int x = ((111*(a1\%mod)*(m2/g))\%mod*q + (111*(a2\%mod)*(m1/g))\%mod*p) \% mod; // if WA
        there is overflow
   a1 = x:
   if (a1 < 0) a1 += mod:
   m1 = mod:
 return {a1, m1}:
```

6.3 Constant modular inverse

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6.4 Extended euclides

```
/// Complexity: O(log(|N|))
/// Tested: https://tinyurl.com/y8yc52gv
ll eea(ll a, ll b, ll& x, ll& y) {
    ll xx = y = 0; ll yy = x = 1;
    while (b) {
        ll q = a / b; ll t = b; b = a % b; a = t;
        t = xx; xx = x - q * xx; x = t;
        t = yy; yy = y - q * yy; y = t;
    }
    return a;
}
ll inverse(ll a, ll n) {
    ll x, y;
    ll g = eea(a, n, x, y);
    if(g > 1)
        return -1;
    return (x % n + n) % n;
}
```

6.5 Fast Fourier transform module

```
/// Complexity: O(|N|*log(|N|))
/// Tested: https://tinyurl.com/yagvw3on
const int mod = 7340033; /// mod = c*2^k+1
/// find g = primitive root of mod.
const int root = 2187; /// (g^c)%mod
const int root_1 = 4665133; /// inverse of root
const int root_pw = 1 << 19; /// 2^k

pii find_c_k(int mod) {
   pii ans;
   for(int k = 1; (1<<k) < mod; k++) {
      int pot = 1<<k;
      if((mod - 1) % pot == 0)
        ans = {(mod-1) / pot, k};
   }
   return ans;
}</pre>
```

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```
int find_primitive_root(int mod) {
  vector<bool> seen(mod);
 for(int r = 2; ; r++) {
   fill(seen.begin(), seen.end(), 0);
    int cur = 1, can = 1;
    for(int i = 0; i <= mod-2 && can; i++) {</pre>
     if(seen[cur]) can = 0;
     seen[cur] = 1:
     cur = (111*cur*r) % mod;
   if(can)
     return r:
  assert(false):
void fft(vector<int> &a, bool inv = 0) {
 int n = a.size():
 for(int i = 1, j = 0; i < n; i++) {</pre>
   int c = n \gg 1:
   for (; j >= c; c >>= 1) j -= c;
   j += c;
   if(i < j) swap(a[i], a[j]);</pre>
  for (int len = 2; len <= n; len <<= 1) {
    int wlen = inv ? root 1 : root:
    for(int i = len; i < root_pw; i <<= 1) wlen = (1 LL * wlen * wlen) % mod;</pre>
    for(int i = 0; i < n; i += len) {</pre>
     int w = 1;
     for(int j = 0; j < (len >> 1); j++) {
       int u = a[i + j], v = (a[i + j + (len >> 1)] * 1 LL * w) % mod;
       a[i + j] = u + v < mod ? u + v : u + v - mod;
       a[i + j + (len >> 1)] = u - v >= 0 ? u - v : u - v + mod;
       w = (w * 1 LL * wlen) \% mod:
  if (inv) {
    int nrev = pow(n):
   for(int i = 0; i < n; i++) a[i] = (a[i] * 1 LL * nrev) % mod;</pre>
}
vector<int> mul(const vector <int> a, const vector <int> b) {
 vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
  while (n < max(a.size(), b.size())) n <<= 1;</pre>
 n <<= 1:
 fa.resize(n); fb.resize(n);
 fft(fa); fft(fb);
 for (int i = 0; i < n; i++) fa[i] = (111 * fa[i] * fb[i]) % mod;</pre>
 fft(fa, 1);
 return fa:
```

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6.6 Fast fourier transform

```
/// Complexity: O(|N|*log(|N|))
/// Tested: https://tinyurl.com/y8g2q66b
namespace fft {
 typedef long long 11;
 const double PI = acos(-1.0):
 vector<int> rev:
 struct pt {
   double r. i:
   pt(double r = 0.0, double i = 0.0) : r(r), i(i) {}
   pt operator + (const pt & b) { return pt(r + b.r. i + b.i); }
   pt operator - (const pt & b) { return pt(r - b.r, i - b.i); }
   pt operator * (const pt & b) { return pt(r * b.r - i * b.i, r * b.i + i * b.r); }
 void fft(vector<pt> &v, int on) {
   int n = v.size();
   for(int i = 1; i < n; i++) if(i < rev[i]) swap(v[i], v[rev[i]]);
   for(int m = 2; m <= n; m <<= 1) {</pre>
     pt wm(cos(-on * 2 * PI / m), sin(-on * 2 * PI / m));
     for(int k = 0; k < n; k += m) {
       pt w(1, 0);
       for(int j = 0; j < m / 2; j++) {</pre>
        pt u = y[k + i];
        pt t = w * v[k + j + m / 2];
        v[k + j] = u + t;
        v[k + j + m / 2] = u - t;
   if(on == -1)
     for(int i = 0; i < n; i++) y[i].r /= n;</pre>
 vector<ll> mul(vector<ll> &a. vector<ll> &b) {
   int n = 1. la = a.size(), lb = b.size(), t;
   for (n = 1, t = 0; n \le (la+lb+1); n \le 1, t++); t = 1 \le (t-1);
   vector<pt> x1(n), x2(n);
   rev.assign(n, 0):
   for(int i = 0; i < n; i++) rev[i] = rev[i >> 1] >> 1 |(i & 1 ? t : 0);
   for(int i = 0; i < la; i++) x1[i] = pt(a[i], 0);</pre>
   for(int i = 0; i < 1b; i++) x2[i] = pt(b[i], 0);
   fft(x1, 1); fft(x2, 1);
   for(int i = 0; i < n; i++) x1[i] = x1[i] * x2[i];</pre>
   fft(x1, -1):
   vector<ll> sol(n);
   for(int i = 0; i < n; i++) sol[i] = x1[i].r + 0.5;
   return sol;
```

.

6.7 Gauss jordan

```
/// Complexity: O(|N|^3)
/// Tested: https://tinyurl.com/y23sh38k
const int EPS = 1;
int gauss (vector<vector<int>> a, vector<int> &ans) {
 int n = a.size(), m = a[0].size()-1;
 vector<int> where(m, -1);
 for(int col = 0, row = 0; col < m && row < n; ++col) {</pre>
   int sel = row;
   for(int i = row; i < n; ++i)</pre>
     if(abs(a[i][col]) > abs(a[sel][col])) sel = i;
   if(abs(a[sel][col]) < EPS) continue;</pre>
   swap(a[sel], a[row]);
   where[col] = row:
   for(int i = 0: i < n: ++i)
     if(i != row) {
       int c = divide(a[i][col], a[row][col]); /// precalc inverses
       for(int j = col; j <= m; ++j)</pre>
         a[i][j] = sub(a[i][j], mul(a[row][j], c));
     }
   ++row;
 ans.assign(m, 0);
 for(int i = 0; i < m; ++i)
   if(where[i] != -1) ans[i] = divide(a[where[i]][m], a[where[i]][i]);
 for(int i = 0; i < n; ++i) {</pre>
   int sum = 0;
   for(int j = 0; j < m; ++j)
     sum = add(sum, mul(ans[j], a[i][j]));
   if(sum != a[i][m]) return 0;
 for(int i = 0; i < m; ++i)</pre>
   if(where[i] == -1) return -1; /// infinite solutions
 return 1;
```

6.8 Integral

- Simpsons rule: $\int_a^b f(x)dx \approx \frac{b-a}{6}[f(a) + 4f(\frac{a+b}{2}) + f(b)]$
- Arc length: $s = \int_a^b \sqrt{1 + [f'(x)]^2} dx$
- Area of a surface of revolution: $A = 2\pi \int_a^b f(x) \sqrt{1 + [f'(x)]^2} dx$
- Volume of a solid of revolution: $V = \pi \int_a^b f(x)^2 dx$

- Note: In case of multiple functions such as g(x) h(x) for a solid of revolution then f(x) = g(x) h(x)
- $f'(x) \approx \frac{f(x+h) f(x-h)}{2h}$
- $f'(x) \approx \frac{f(x-2h)-8f(x-h)+8f(x+h)-f(x+2h)}{12h}$
- $f''(x) \approx \frac{f(x+h)-2f(x)+f(x-h)}{h^2}$

6.9 Lagrange Interpolation

```
/// Complexity: 0(|N|^2)
/// Tested: https://tinyurl.com/y23sh38k
vector<lf> X, F;
If f(lf x) {
    lf answer = 0;
    for(int i = 0; i < (int)X.size(); i++) {
        lf prod = F[i];
        for(int j = 0; j < (int)X.size(); j++) {
            if(i == j) continue;
            prod = mul(prod, divide(sub(x, X[j]), sub(X[i], X[j])));
        }
        answer = add(answer, prod);
    }
    return answer;
}</pre>
```

6.10 Linear diophantine

```
/// Complexity: O(log(|N|))
/// Tested: https://tinyurl.com/y8yc52gv
bool diophantine(ll a, ll b, ll c, ll &x, ll &y, ll &g) {
 x = y = 0;
 if(a == 0 && b == 0) return c == 0;
 if(b == 0) swap(a, b), swap(x, v):
 g = eea(abs(a), abs(b), x, y);
 if(c % g) return false;
 a /= g; b /= g; c /= g;
 if(a < 0) x *= -1;
 x = (x \% b) * (c \% b) \% b;
 if(x < 0) x += b;
 y = (c - a*x) / b;
 return true;
///finds the first k \mid x + b * k / gcd(a, b) >= val
ll greater_or_equal_than(ll a, ll b, ll x, ll val, ll g) {
 lf got = 1.0 * (val - x) * g / b;
 return b > 0 ? ceil(got) : floor(got);
```

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6.11 Matrix multiplication

```
const int MOD = 1e9+7;
struct matrix {
 const int N = 2;
 int m[N][N], r, c;
 matrix(int r = N, int c = N, bool iden = false) : r(r), c(c) {
   memset(m, 0, sizeof m);
   if(iden)
     for(int i = 0; i < r; i++) m[i][i] = 1;</pre>
 matrix operator * (const matrix &o) const {
   matrix ret(r. o.c):
   for(int i = 0; i < r; ++i)</pre>
     for(int i = 0: i < o.c: ++i) {
      11 \&r = ret.m[i][i]:
       for(int k = 0: k < c: ++k)
         r = (r + 111*m[i][k]*o.m[k][j]) % MOD;
     }
   return ret;
 }
};
```

6.12 Miller rabin

```
/// Complexity: ???
/// Tested: A lot.. but no link
ll mul (ll a, ll b, ll mod) {
    ll ret = 0;
    for(a %= mod, b %= mod; b != 0;
        b >>= 1, a <<= 1, a = a >= mod ? a - mod : a) {
        if (b & 1) {
            ret += a;
            if (ret >= mod) ret -= mod;
        }
    }
    return ret;
}
ll fpow (ll a, ll b, ll mod) {
    ll ans = 1;
    for (; b; b >>= 1, a = mul(a, a, mod))
        if (b & 1)
```

```
ans = mul(ans, a, mod);
 return ans;
bool witness (ll a, ll s, ll d, ll n) {
 11 x = fpow(a, d, n);
 if (x == 1 || x == n - 1) return false;
 for (int i = 0; i < s - 1; i++) {
   x = mul(x, x, n);
   if (x == 1) return true:
   if (x == n - 1) return false:
 return true;
ll test[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 0};
bool is_prime (ll n) {
 if (n < 2) return false:
 if (n == 2) return true:
 if (n % 2 == 0) return false:
 11 d = n - 1, s = 0:
 while (d \% 2 == 0) ++s, d /= 2;
 for (int i = 0: test[i] && test[i] < n: ++i)
   if (witness(test[i], s, d, n))
     return false;
 return true;
```

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6.13 Pollard's rho

```
/// Complexity: ???
/// Tested: Not vet
11 pollard_rho(ll n, ll c) {
 11 x = 2, y = 2, i = 1, k = 2, d;
 while (true) {
   x = (mul(x, x, n) + c);
   if (x >= n) x -= n;
   d = \_gcd(x - y, n);
   if (d > 1) return d:
   if (++i == k) y = x, k <<= 1;
 return n;
void factorize(ll n, vector<ll> &f) {
 if (n == 1) return;
 if (is_prime(n)) {
   f.push_back(n);
   return;
 11 d = n;
 for (int i = 2; d == n; i++)
   d = pollard_rho(n, i);
 factorize(d, f);
```

```
factorize(n/d, f);
}
```

6.14 Simplex

```
/// Complexity: O(|N|^2 * |M|) N variables, N restrictions
/// Tested: https://tinyurl.com/ybphh57p
const double EPS = 1e-6;
typedef vector<double> vec;
namespace simplex {
 vector<int> X. Y:
 vector<vec> a:
 vec b. c:
 double z:
 int n. m:
 void pivot(int x, int y) {
   swap(X[y], Y[x]);
   b[x] /= a[x][y];
   for(int i = 0; i < m; i++)</pre>
    if(i != y)
      a[x][i] /= a[x][v];
   a[x][y] = 1 / a[x][y];
   for(int i = 0; i < n; i++)</pre>
     if(i != x && abs(a[i][y]) > EPS) {
     b[i] = a[i][v] * b[x];
     for(int j = 0; j < m; j++)
      if(j != y)
         a[i][j] -= a[i][y] * a[x][j];
     a[i][y] = -a[i][y] * a[x][y];
   z += c[y] * b[x];
   for(int i = 0; i < m; i++)</pre>
    if(i != y)
       c[i] = c[y] * a[x][i];
   c[y] = -c[y] * a[x][y];
 /// A is a vector of 1 and 0. B is the limit restriction. C is the factors of 0.F.
 pair<double, vec> simplex(vector<vec> &A, vec &B, vec &C) {
   a = A: b = B: c = C:
   n = b.size(); m = c.size(); z = 0.0;
   X = vector<int>(m):
   Y = vector<int>(n):
   for(int i = 0; i < m; i++) X[i] = i;</pre>
   for(int i = 0; i < n; i++) Y[i] = i + m;</pre>
   while(1) {
     int x = -1, y = -1;
     double mn = -EPS;
     for(int i = 0; i < n; i++)</pre>
      if(b[i] < mn)
         mn = b[i], x = i;
     if(x < 0) break;
```

```
for(int i = 0; i < m; i++)</pre>
    if(a[x][i] < -EPS) { y = i; break; }</pre>
  assert(y >= 0); // no sol
 pivot(x, y);
while(1) {
  double mx = EPS;
  int x = -1, y = -1;
  for(int i = 0; i < m; i++)</pre>
   if(c[i] > mx)
     mx = c[i], v = i:
  if(v < 0) break:
  double mn = 1e200:
  for(int i = 0; i < n; i++)</pre>
   if(a[i][y] > EPS && b[i] / a[i][y] < mn)</pre>
     mn = b[i] / a[i][y], x = i;
  assert(x >= 0): // unbound
  pivot(x, y);
vec r(m);
for(int i = 0: i < n: i++)</pre>
  if(Y[i] < m)
   r[Y[i]] = b[i];
return make_pair(z, r);
```

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6.15 Simpson

```
/// Complexity: ?????
/// Tested: Not yet
inline lf simpson(lf fl, lf fr, lf fmid, lf l, lf r) {
 return (f1 + fr + 4.0 * fmid) * (r - 1) / 6.0;
lf rsimpson (lf slr, lf fl, lf fr, lf fmid, lf l, lf r) {
       lf mid = (1 + r) * 0.5;
       lf fml = f((1 + mid) * 0.5);
       lf fmr = f((mid + r) * 0.5);
       lf slm = simpson(fl, fmid, fml, l, mid);
       lf smr = simpson(fmid, fr, fmr, mid, r);
       if (fabs(slr - slm - smr) < eps) return slm + smr;</pre>
       return rsimpson(slm, fl, fmid, fml, l, mid) + rsimpson(smr, fmid, fr, fmr, mid,
           r):
lf integrate(lf l.lf r) {
       If mid = (1 + r) * .5, fl = f(1), fr = f(r), fmid = f(mid);
       return rsimpson(simpson(fl, fr, fmid, l, r), fl, fr, fmid, l, r);
```

6.16 Totient and divisors

```
vector<int> count_divisors_sieve() {
 bitset<mx> is_prime; is_prime.set();
  vector<int> cnt(mx, 1);
  is_prime[0] = is_prime[1] = 0;
  for(int i = 2; i < mx; i++) {</pre>
    if(!is_prime[i]) continue;
    cnt[i]++:
    for(int j = i+i; j < mx; j += i) {</pre>
     int n = j, c = 1;
     while(n\%i == 0) n /= i, c++;
     cnt[j] *= c;
     is_prime[j] = 0;
 }
 return cnt:
vector<int> euler phi sieve() {
 bitset<mx> is_prime; is_prime.set();
  vector<int> phi(mx);
  iota(phi.begin(), phi.end(), 0);
  is prime[0] = is prime[1] = 0:
 for(int i = 2: i < mx: i++) {
    if(!is_prime[i]) continue;
   for(int j = i; j < mx; j += i) {</pre>
     phi[j] -= phi[j]/i;
     is_prime[j] = 0;
 }
 return phi;
ll euler_phi(ll n) {
 ll ans = n;
 for(ll i = 2; i * i <= n; ++i) {</pre>
   if(n % i == 0) {
     ans -= ans / i:
     while(n \% i == 0) n /= i;
 }
 if(n > 1) ans -= ans / n:
 return ans:
```

7 Network flows

7.1 Blossom

```
/// Complexity: O(|E||V|^2)
```

```
/// Tested: https://tinyurl.com/oe5rnpk
struct network {
 struct struct_edge { int v; struct_edge * n; };
 typedef struct_edge* edge;
 int n;
 struct_edge pool[MAXE]; ///2*n*n;
  edge top;
 vector<edge> adi;
 queue<int> q;
 vector<int> f, base, inq, inb, inp, match;
 vector<vector<int>> ed:
 network(int n) : n(n), match(n, -1), adj(n), top(pool), f(n), base(n),
                 inq(n), inb(n), inp(n), ed(n), vector < int > (n) {}
 void add_edge(int u, int v) {
   if(ed[u][v]) return;
   ed[u][v] = 1:
   top \rightarrow v = v, top \rightarrow n = adj[u], adj[u] = top ++;
   top->v = u, top->n = adi[v], adi[v] = top++:
 int get_lca(int root, int u, int v) {
   fill(inp.begin(), inp.end(), 0);
   while(1) {
     inp[u = base[u]] = 1;
     if(u == root) break;
     u = f[match[u]];
   while(1) {
     if(inp[v = base[v]]) return v;
     else v = f[ match[v] ];
 void mark(int lca, int u) {
   while(base[u] != lca) {
     int v = match[u];
     inb[ base[u ]] = 1:
     inb[base[v]] = 1:
     u = f[v]:
     if(base[u] != lca) f[u] = v:
   }
 void blossom_contraction(int s, int u, int v) {
   int lca = get lca(s, u, v):
   fill(inb.begin(), inb.end(), 0);
   mark(lca, u): mark(lca, v):
   if(base[u] != lca) f[u] = v;
   if(base[v] != lca) f[v] = u;
   for(int u = 0; u < n; u++)</pre>
     if(inb[base[u]]) {
       base[u] = lca;
       if(!inq[u]) {
          inq[u] = 1;
          q.push(u);
```

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m coUNter}$

```
int bfs(int s) {
   fill(inq.begin(), inq.end(), 0);
   fill(f.begin(), f.end(), -1);
   for(int i = 0; i < n; i++) base[i] = i;</pre>
   q = queue<int>();
   q.push(s);
   inq[s] = 1;
   while(q.size()) {
     int u = q.front(); q.pop();
     for(edge e = adj[u]; e; e = e->n) {
       int v = e \rightarrow v:
       if(base[u] != base[v] && match[u] != v) {
         if((v == s) || (match[v] != -1 && f[match[v]] != -1))
          blossom contraction(s. u. v):
         else if(f[v] == -1) {
          f[v] = u:
           if(match[v] == -1) return v:
           else if(!inq[match[v]]) {
            inq[match[v]] = 1;
            q.push(match[v]);
   return -1;
  int doit(int u) {
   if(u == -1) return 0;
   int v = f[u];
   doit(match[v]);
   match[v] = u; match[u] = v;
   return u != -1:
 /// (i < net.match[i]) => means match
 int maximum matching() {
   int ans = 0:
   for(int u = 0: u < n: u++)
     ans += (match[u] == -1) && doit(bfs(u));
   return ans:
 }
};
```

7.2 Dinic

```
/// Complexity: 0(|E|*|V|^2)
/// Tested: https://tinyurl.com/ya9rgoyk
struct edge { int v, cap, inv, flow; };
struct network {
```

```
int n, s, t;
 vector<int> lvl:
 vector<vector<edge>> g;
 network(int n) : n(n), lvl(n), g(n) {}
 void add_edge(int u, int v, int c) {
   g[u].push_back({v, c, g[v].size(), 0});
   g[v].push_back({u, 0, g[u].size()-1, c});
 bool bfs() {
   fill(lvl.begin(), lvl.end(), -1);
   queue<int> q:
   lvl[s] = 0:
   for(q.push(s); q.size(); q.pop()) {
     int u = q.front();
     for(auto &e : g[u]) {
      if(e.cap > 0 && lvl[e.v] == -1) {
        lvl[e.v] = lvl[u]+1:
        q.push(e.v);
   return lvl[t] != -1;
 int dfs(int u, int nf) {
   if(u == t) return nf;
   int res = 0;
   for(auto &e : g[u]) {
    if(e.cap > 0 \&\& lvl[e.v] == lvl[u]+1) {
       int tf = dfs(e.v, min(nf, e.cap));
      res += tf; nf -= tf; e.cap -= tf;
      g[e.v][e.inv].cap += tf;
      g[e.v][e.inv].flow -= tf;
       e.flow += tf;
       if(nf == 0) return res;
   if(!res) lvl[u] = -1:
   return res:
 int max_flow(int so, int si, int res = 0) {
   s = so: t = si:
   while(bfs()) res += dfs(s, INT MAX):
   return res:
};
```

7.3 Hopcroft karp

```
/// Complexity: 0(|E|*sqrt(|V|))
/// Tested: https://tinyurl.com/yad2g9g9
struct mbm {
```

```
vector<vector<int>> g;
 vector<int> d, match;
 int nil, l, r;
 /// u -> 0 to 1, v -> 0 to r
 mbm(int 1, int r) : 1(1), r(r), nil(1+r), g(1+r),
                   d(1+1+r, INF), match(1+r, 1+r) {}
 void add_edge(int a, int b) {
   g[a].push_back(1+b);
   g[1+b].push_back(a);
 bool bfs() {
   queue<int> q;
   for(int u = 0: u < 1: u++) {
     if(match[u] == nil) {
       d[u] = 0:
       q.push(u);
     } else d[u] = INF:
   d[nil] = INF:
   while(q.size()) {
     int u = q.front(); q.pop();
     if(u == nil) continue;
     for(auto v : g[u]) {
      if(d[ match[v] ] == INF) {
         d[match[v]] = d[u]+1;
         q.push(match[v]);
   return d[nil] != INF;
 bool dfs(int u) {
   if(u == nil) return true;
   for(int v : g[u]) {
     if(d[ match[v] ] == d[u]+1 && dfs(match[v])) {
      match[v] = u: match[u] = v:
      return true:
     }
   d[u] = INF:
   return false;
 int max matching() {
   int ans = 0:
   while(bfs()) {
    for(int u = 0: u < 1: u++) {
       ans += (match[u] == nil && dfs(u));
   return ans;
};
```

7.4 Maximum bipartite matching

```
/// Complexity: O(|E|*|V|)
/// Tested: https://tinyurl.com/yad2g9g9
struct mbm {
 int 1. r:
 vector<vector<int>> g;
 vector<int> match. seen:
 mbm(int 1, int r) : 1(1), r(r), seen(r), match(r), g(1) {}
 void add edge(int 1, int r) { g[1].push back(r); }
 bool dfs(int u) {
   for(auto v : g[u]) {
     if(seen[v]++) continue;
     if(match[v] == -1 || dfs(match[v])) {
      match[v] = u;
       return true;
   return false;
 int max_matching() {
   int ans = 0;
   fill(match.begin(), match.end(), -1);
   for(int u = 0; u < 1; ++u) {
    fill(seen.begin(), seen.end(), 0);
     ans += dfs(u):
   return ans:
};
```

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7.5 Maximum flow minimum cost

```
/// Complexity: 0(|V|*|E|^2*log(|E|))
/// Tested: https://tinyurl.com/ycgpp47z
template <class type>
struct mcmf {
 struct edge { int u, v, cap, flow; type cost; };
 int n:
 vector<edge> ed;
 vector<vector<int>> g;
 vector<int> p;
 vector<type> d, phi;
 mcmf(int n) : n(n), g(n), p(n), d(n), phi(n) {}
 void add_edge(int u, int v, int cap, type cost) {
   g[u].push_back(ed.size());
   ed.push_back({u, v, cap, 0, cost});
   g[v].push_back(ed.size());
   ed.push_back({v, u, 0, 0, -cost});
```

```
bool dijkstra(int s, int t) {
   fill(d.begin(), d.end(), INF);
   fill(p.begin(), p.end(), -1);
   set<pair<type, int>> q;
   d[s] = 0;
   for(q.insert({d[s], s}); q.size();) {
     int u = (*q.begin()).second; q.erase(q.begin());
     for(auto v : g[u]) {
       auto &e = ed[v]:
       type nd = d[e.u]+e.cost+phi[e.u]-phi[e.v];
       if(0 < (e.cap-e.flow) && nd < d[e.v]) {
         q.erase({d[e.v], e.v});
         d[e.v] = nd; p[e.v] = v;
         q.insert({d[e.v], e.v});
     }
   for(int i = 0; i < n; i++) phi[i] = min(INF, phi[i]+d[i]);</pre>
   return d[t] != INF:
 pair<int, type> max_flow(int s, int t) {
   type mc = 0;
   int mf = 0:
   fill(phi.begin(), phi.end(), 0);
   while(dijkstra(s, t)) {
     int flow = INF;
     for(int v = p[t]; v != -1; v = p[ ed[v].u ])
      flow = min(flow, ed[v].cap-ed[v].flow);
     for(int v = p[t]; v != -1; v = p[ ed[v].u ]) {
       edge &e1 = ed[v];
       edge &e2 = ed[v^1];
       mc += e1.cost*flow;
       e1.flow += flow;
       e2.flow -= flow;
     }
     mf += flow:
   return {mf, mc};
};
```

7.6 Stoer Wagner

```
/// Complexity: 0(|V|^3)
/// Tested: https://tinyurl.com/y8eu433d
struct stoer_wagner {
  int n;
  vector<vector<int>> g;
  stoer_wagner(int n) : n(n), g(n, vector<int>(n)) {}
  void add_edge(int a, int b, int w) { g[a][b] = g[b][a] = w; }
  pair<int, vector<int>> min_cut() {
```

```
vector<int> used(n);
   vector<int> cut, best_cut;
   int best_weight = -1;
  for(int p = n-1; p >= 0; --p) {
    vector<int> w = g[0];
    vector<int> added = used;
    int prv, lst = 0;
    for(int i = 0; i < p; ++i) {</pre>
      prv = lst; lst = -1;
      for(int j = 1; j < n; ++j)
        if(!added[i] && (lst == -1 || w[i] > w[lst]))
          lst = j;
      if(i == p-1) {
        for(int j = 0; j < n; j++)
          g[prv][j] += g[lst][j];
        for(int j = 0; j < n; j++)
          g[j][prv] = g[prv][j];
        used[1st] = true:
        cut.push back(1st):
        if(best_weight == -1 || w[lst] < best_weight) {</pre>
          best cut = cut:
          best_weight = w[lst];
      } else {
        for(int j = 0; j < n; j++)
          w[i] += g[lst][i];
        added[lst] = true;
  return {best_weight, best_cut}; /// best_cut contains all nodes in the same set
};
```

7.7 Weighted matching

```
/// Complexity: 0(|V|^3)
/// Tested: https://tinyurl.com/ycpq8eyl problem G
typedef int type;
struct matching_weighted {
  int l, r;
  vector<vector<type>> c;
  matching_weighted(int l, int r) : l(l), r(r), c(l, vector<type>(r)) {
    assert(l <= r);
  }
  void add_edge(int a, int b, type cost) { c[a][b] = cost; }
  type matching() {
    vector<type> v(r), d(r); // v: potential
    vector<int> ml(l, -1), mr(r, -1); // matching pairs
    vector<int> idx(r), prev(r);
    iota(idx.begin(), idx.end(), 0);
```

```
auto residue = [&](int i, int j) { return c[i][j]-v[j]; };
for(int f = 0; f < 1; ++f) {</pre>
 for(int j = 0; j < r; ++j) {
   d[j] = residue(f, j);
   prev[i] = f;
  type w;
  int j, 1;
 for (int s = 0, t = 0;;) {
   if(s == t) {
     1 = s:
     w = d[idx[t++]]:
     for(int k = t: k < r: ++k) {
       j = idx[k];
       type h = d[j];
       if (h <= w) {
        if (h < w) t = s, w = h;
        idx[k] = idx[t]:
        idx[t++] = i:
     for (int k = s; k < t; ++k) {
       j = idx[k];
       if (mr[j] < 0) goto aug;</pre>
   int q = idx[s++], i = mr[q];
   for (int k = t; k < r; ++k) {
     j = idx[k];
     type h = residue(i, j) - residue(i, q) + w;
     if (h < d[i]) {</pre>
       d[i] = h;
       prev[j] = i;
       if(h == w) {
        if(mr[j] < 0) goto aug;</pre>
        idx[k] = idx[t]:
         idx[t++] = j;
   }
  aug: for (int k = 0: k < 1: ++k)
   v[idx[k]] += d[idx[k]] - w:
  int i:
  do {
   mr[i] = i = prev[j];
   swap(j, ml[i]);
 } while (i != f);
type opt = 0;
for (int i = 0; i < 1; ++i)</pre>
 opt += c[i][ml[i]]; // (i, ml[i]) is a solution
return opt;
```

```
}
};
```

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8 Strings

8.1 Aho corasick

```
/// Complexity: O(|text|+SUM(|pattern_i|)+matches)
/// Tested: https://tinyurl.com/y2zq594p
const static int alpha = 26;
int trie[N*alpha][alpha], fail[N*alpha], nodes;
void add(string &s, int i) {
 int cur = 0:
 for(char c : s) {
   int x = c^{-a}:
   if(!trie[cur][x]) trie[cur][x] = ++nodes;
   cur = trie[cur][x]:
 //cnt_word[cur]++;
 //end word[cur] = i: // for i > 0
void build() {
 queue<int> q; q.push(0);
 while(q.size()) {
   int u = q.front(); q.pop();
   for(int i = 0; i < alpha; ++i) {</pre>
     int v = trie[u][i];
     if(!v) continue;
     q.push(v);
     if(!u) continue;
     fail[v] = fail[u];
     while(fail[v] && !trie[ fail[v] ][i]) fail[v] = fail[ fail[v] ];
     fail[v] = trie[ fail[v] ][i];
     //fail_out[v] = end_word[ fail[v] ] ? fail[v] : fail_out[ fail[v] ];
     //cnt_word[v] += cnt_word[ fail[v] ]; // obtener informacion del fail_padre
 }
```

8.2 Hashing

```
/// Tested: https://tinyurl.com/y8qstx97
/// 1000234999, 1000567999, 1000111997, 1000777121
const int MODS[] = { 1001864327, 1001265673 };
const mint BASE(256, 256), ZERO(0, 0), ONE(1, 1);
inline int add(int a, int b, const int& mod) { return a+b >= mod ? a+b-mod : a+b; }
inline int sbt(int a, int b, const int& mod) { return a-b < 0 ? a-b+mod : a-b; }
```

```
inline int mul(int a, int b, const int& mod) { return 111*a*b%mod; }
inline 11 operator ! (const mint a) { return (ll(a.first) << 32) | ll(a.second); }</pre>
inline mint operator + (const mint a, const mint b) {
 return {add(a.first, b.first, MODS[0]), add(a.second, b.second, MODS[1])};
inline mint operator - (const mint a, const mint b) {
 return {sbt(a.first, b.first, MODS[0]), sbt(a.second, b.second, MODS[1])};
inline mint operator * (const mint a, const mint b) {
 return {mul(a.first, b.first, MODS[0]), mul(a.second, b.second, MODS[1])}:
mint base[MAXN]:
void prepare() {
 base[0] = ONE;
 for(int i = 1; i < MAXN; i++) base[i] = base[i-1]*BASE;</pre>
template <class type>
struct hashing {
 vector<mint> code:
 hashing(type &t) {
   code.resize(t.size()+1):
   code[0] = ZER0;
   for (int i = 1; i < code.size(); ++i)</pre>
     code[i] = code[i-1]*BASE + mint{t[i-1], t[i-1]};
 mint query(int 1, int r) {
   return code[r+1] - code[1]*base[r-1+1];
};
```

8.3 Kmp automaton

```
/// Complexity: 0(|N|*alphabet)
/// Tested: not yet
const int alpha = 256;
int aut[102][alpha];
void kmp_automaton(string &t) {
    vector(int> phi = get_phi(t);
    for(int i = 0; i <= t.size(); ++i) {
        for(int c = 0; c < alpha; ++c) {
            if(i == t.size() || (i > 0 && c != t[i])) aut[i][c] = aut[ phi[i-1] ][c];
        else aut[i][c] = i + (c == t[i]);
    }
}
```

8.4 Kmp

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/y7svn3kr
vector<int> get_phi(string &p) {
 vector<int> phi(p.size());
 phi[0] = 0;
 for(int i = 1, j = 0; i < p.size(); ++i ) {</pre>
   while(j > 0 \&\& p[i] != p[j]) j = phi[j-1];
   if(p[i] == p[j]) ++j;
   phi[i] = j;
 return phi:
int get_match(string &t, string &p) {
 vector<int> phi = get_phi(p);
 int matches = 0:
 for(int i = 0, j = 0; i < t.size(); ++i ) {</pre>
   while(j > 0 \&\& t[i] != p[j]) j = phi[j-1];
   if(t[i] == p[j]) ++j;
   if(j == p.size()) {
     matches++;
     j = phi[j-1];
 return matches;
```

8.5 Manacher

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/y6upxbpa
///to = i - from[i];
///len = to - from[i] + 1 = i - 2 * from[i] + 1;
vector<int> manacher(string &s) {
 int n = s.size(), p = 0, pr = -1;
 vector<int> from(2*n-1);
 for(int i = 0; i < 2*n-1; ++i) {</pre>
   int r = i <= 2*pr ? min(p - from[2*p - i], pr) : i/2;
   int 1 = i - r;
   while (1 > 0 \&\& r < n-1 \&\& s[1-1] == s[r+1]) --1, ++r;
   from[i] = 1:
   if (r > pr) {
     pr = r;
     p = i;
 return from;
```

8.6 Minimum expression

```
/// Complexity: 0(|N|)
/// Tested: https://tinyurl.com/y6gfzgsm
int minimum_expression(string s) {
    s = s+s;
    int len = s.size(), i = 0, j = 1, k = 0;
    while(i+k < len && j+k < len) {
        if(s[i+k] = s[j+k]) k++;
        else if(s[i+k] > s[j+k]) i = i+k+1, k = 0;
        else j = j+k+1, k = 0;
        if(i == j) j++;
    }
    return min(i, j);
}
```

8.7 Suffix array

```
/// Complexity: O(|N|*log(|N|))
/// Tested: https://tinyurl.com/y8wdubdw
struct suffix_array {
 const static int alpha = 300;
 int mx, n;
 string s;
 vector<int> pos, tpos, sa, tsa, lcp;
 suffix_array(string t) {
   s = t+"$"; n = s.size(); mx = max(alpha, n)+2;
   pos = tpos = tsa = sa = lcp = vector<int>(n);
 bool check(int i, int gap) {
   if(pos[ sa[i-1] ] != pos[ sa[i] ]) return true;
   if(sa[i-1]+gap < n && sa[i]+gap < n)
    return (pos[ sa[i-1]+gap ] != pos[ sa[i]+gap ]);
   return true;
 void radix sort(int k) {
   vector<int> cnt(mx):
   for(int i = 0: i < n: i++)</pre>
     cnt[(i+k < n) ? pos[i+k]+1 : 1]++;
   for(int i = 1: i < mx: i++)
     cnt[i] += cnt[i-1];
   for(int i = 0; i < n; i++)</pre>
     tsa[cnt[(sa[i]+k < n) ? pos[sa[i]+k] : 0]++] = sa[i];
 void build sa() {
   for(int i = 0; i < n; i++) {</pre>
     sa[i] = i:
    pos[i] = s[i];
```

```
for(int gap = 1; gap < n; gap <<= 1) {</pre>
     radix_sort(gap);
     radix_sort(0);
     tpos[ sa[0] ] = 0;
     for(int i = 1; i < n; i++)</pre>
       tpos[ sa[i] ] = tpos[ sa[i-1] ] + check(i, gap);
     if(pos[ sa[n-1] ] == n-1) break;
 void build lcp() {
   int k = 0:
   1cp[0] = 0:
   for(int i = 0; i < n; i++) {</pre>
     if(pos[i] == 0) continue;
     while(s[i+k] == s[sa[pos[i]-1]+k]) k++:
     lcp[pos[i]] = k;
     k = max(0, k-1):
 int& operator[] ( int i ){ return sa[i]; }
};
```

8.8 Suffix automaton

```
/// Complexity: O(|N|*log(|alphabet|))
/// Tested: https://tinyurl.com/y7cevdeg
struct suffix_automaton {
 struct node {
   int len. link: bool end:
   map<char, int> next;
 }:
 vector<node> sa;
 int last:
 suffix_automaton() {}
 suffix automaton(string s) {
   sa.reserve(s.size()*2);
   last = add node():
   sa[last].len = 0:
   sa[last].link = -1;
   for(int i = 0; i < s.size(); ++i)</pre>
     sa_append(s[i]);
   ///t0 is not suffix
   for(int cur = last; cur; cur = sa[cur].link)
     sa[cur].end = 1;
 int add_node() {
   sa.push_back({});
   return sa.size()-1;
 void sa_append(char c) {
```

```
int cur = add_node();
   sa[cur].len = sa[last].len + 1;
   int p = last;
   while(p != -1 && !sa[p].next[c] ){
     sa[p].next[c] = cur;
    p = sa[p].link;
   if(p == -1) sa[cur].link = 0;
   else {
     int q = sa[p].next[c];
     if(sa[q].len == sa[p].len+1) sa[cur].link = q;
     else {
      int clone = add node():
      sa[clone] = sa[q];
      sa[clone].len = sa[p].len+1;
      sa[q].link = sa[cur].link = clone;
      while(p != -1 && sa[p].next[c] == q) {
        sa[p].next[c] = clone:
        p = sa[p].link;
   last = cur;
 node& operator[](int i) { return sa[i]; }
};
```

8.9 Z algorithm

```
/// Complexity: O(|N|)
/// Tested: https://tinyurl.com/yc3rjh4p
vector<int> z_algorithm (string s) {
    int n = s.size();
    vector<int> z(n);
    int x = 0, y = 0;
    for(int i = 1; i < n; ++i) {
        z[i] = max(0, min(z[i-x], y-i+1));
        while (i+z[i] < n && s[z[i]] == s[i+z[i]])
        x = i, y = i+z[i], z[i]++;
    }
    return z;
}
```

9 Utilities

9.1 Hash STL

9.2 Pragma optimizations

```
#pragma GCC optimize ("03")
#pragma GCC target ("sse4")
#pragma GCC target ("avx,tune=native")
```

9.3 Random

```
// Declare number generator
  mt19937 / mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count())
  // or
  random_device rd
  mt19937 / mt19937_64 rng(rd())

// Use it to shuffle a vector
  shuffle(permutation.begin(), permutation.end(), rng)

// Use it to generate a random number between [fr, to]
  uniform_int_distribution<T> / uniform_real_distribution<T> dis(fr, to);
  dis(rng)
```

9.4 template

```
#include <bits/stdc++.h>
using namespace std;

#define ff first
#define ss second
```

```
#define mp make_pair
#define pb push_back

typedef long long ll;
typedef double lf;
typedef pair<int,int> pii;

const int N = 1e5+10;
const int oo = 1e9;

int main () {
   ios::sync_with_stdio(0);
   cin.tie(0);
   #ifdef LOCAL
        freopen("input.txt", "r", stdin);
   #else
        #define endl '\n'
#endif

return 0;
}
```

9.5 vimsrc

```
//vimrc
noremap <F5> :w <bar> !g++ -DLOCAL -std=c++14 -static -Wall -Wno-unused-result -02
   %:r.cpp -o %:r<CR>
noremap <F6> :w <bar> !g++ -DLOCAL -std=c++14 -static -Wall -Wno-unused-result -02
   %:r.cpp -o %:r && ./%:r<in<CR>
noremap <F9> :<C-U> !./%:r<CR>
set number
set shiftwidth=2
set tabstop=2
set autoindent
set expandtab
//judge
submit() { boca-submit-run TEAM PASSWORD "$1" C++14 "$1.cpp"; }
```

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