Departamento de Sistemas e Computação – FURB Curso de Ciência da Computação Disciplina de Processamento de imagens

Envoltória convexa

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Bibliografia



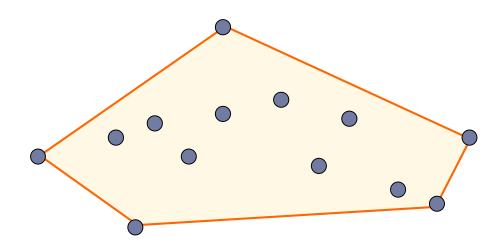
Computational geometry: algorithms and applications

Mark Berg

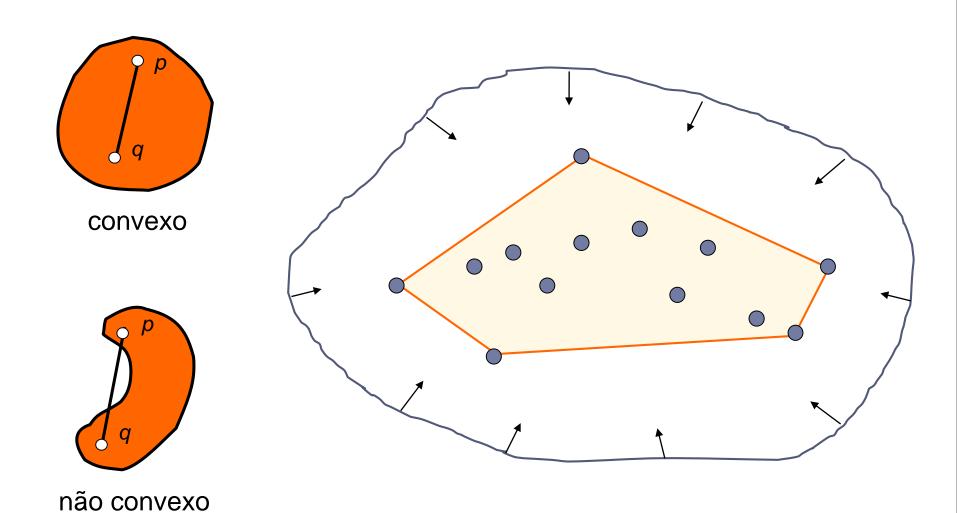
New York: Springer, 2000, 367p.

ENVOLTÓRIA CONVEXA

▶ Enunciado: Dados n pontos queremos encontrar os pontos extremos.

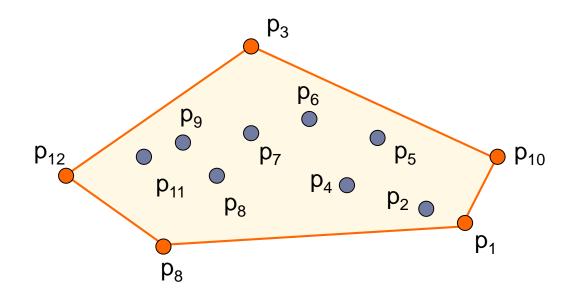


ENVOLTÓRIA CONVEXA



ENVOLTÓRIA CONVEXA

▶ Entrada: p1, p2, p3, p4, p5, p6, p7, p8, p9, p10, p11, p12

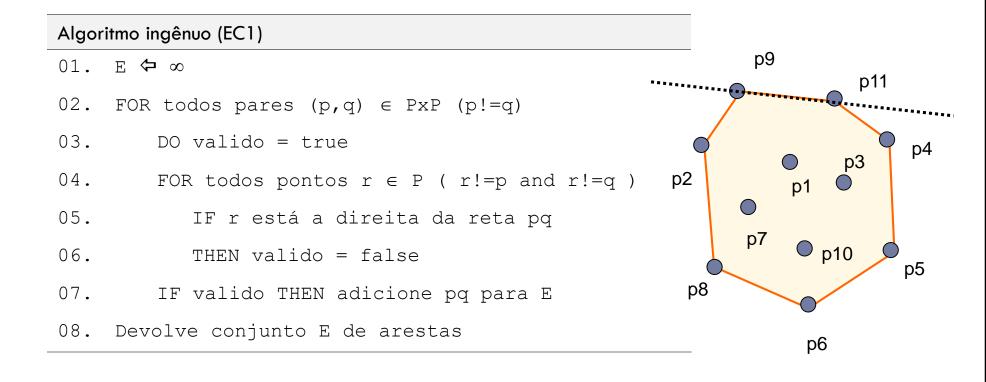


EC é um polígono convexo

Saída: p12, p8, p3, p1, p10

ENVOLTÓRIA CONVEXA [Algoritmo ingênuo]

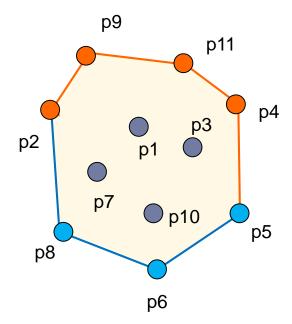
▶ **Entrada:** Conjunto de pontos P no plano



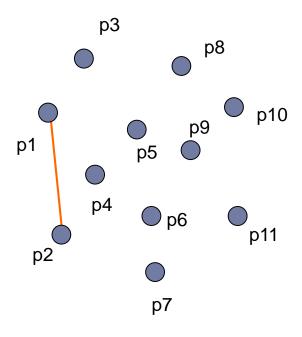
▶ **Saída:** Uma lista de vértices da **EC**

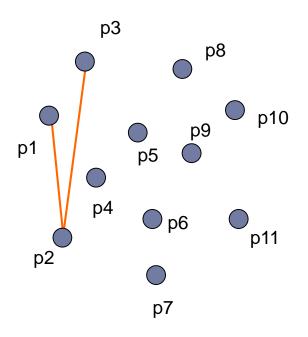
https://doi.org/10.1016%2F0020-0190%2872%2990045-2

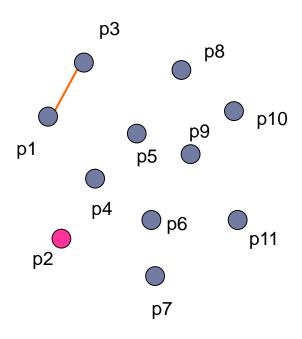
- Classificar os pontos em x, obtendo uma sequencia ordenada p₁, p₂,..., p_n
- Calcular a EC superior percorrendo a lista acima da esquerda para a direita
- Calcular a EC inferior percorrendo a lista acima da direita para a esquerda

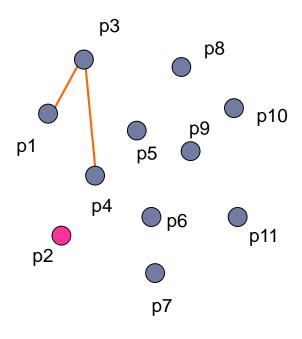


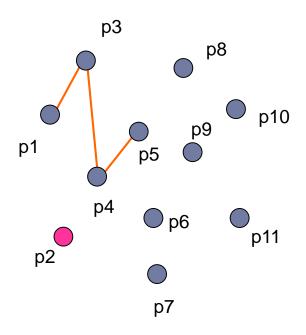
- Definindo curva [virada] para esquerda ou direita
 - Dado três pontos (x1,y1), (x2,y2) e (x3,y3), simplesmente calculando o produto vetorial (x2 x1)(y3 y1) (y2 y1) (x3 x1) dos dois vetores definidos pelos pontos (x1,y1), (x2,y2) e (x2,y2), (x3,y3). Se o resultado for Zero, os três pontos são colineares, se for positivo, os três pontos constituem uma "curva para esquerda", caso contrario uma "curva para direita"

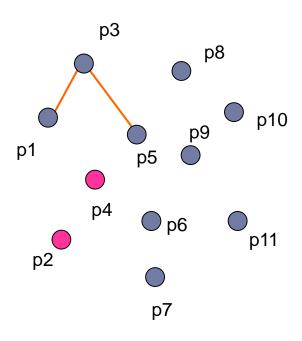


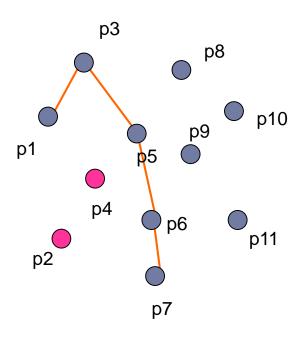


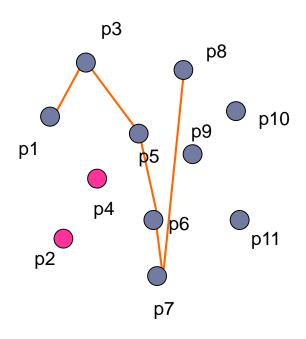


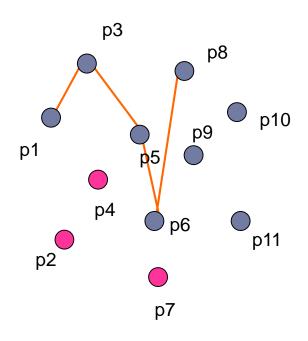


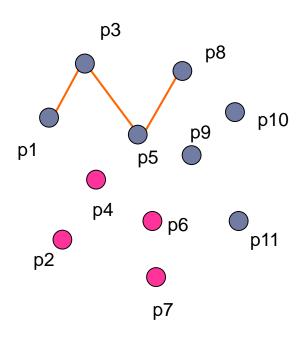


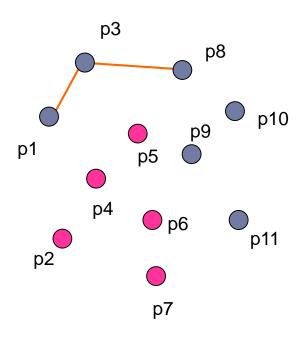


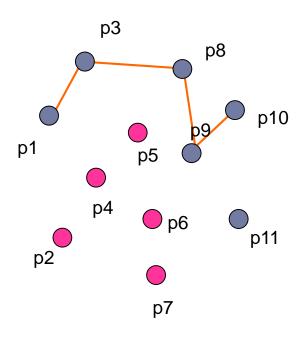


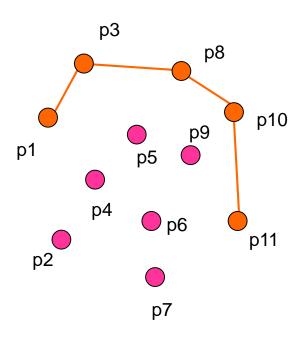












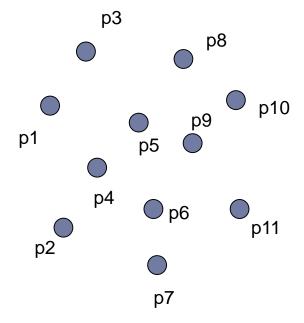
▶ **Entrada:** Conjunto de pontos P no plano

```
Algoritmo Grahan's Scan (EC2)
01.
      Classificar os pontos pela coordenada X (p1, p2,... Pn)
02.
      Colocar p1 e p2 na lista L SUP
03.
      FOR i := 3 TO N
04.
          DO adicionar pi a L SUP
              WHILE L SUP > 2 pontos e últimos três fazem uma
05.
              virada para esquerda
06.
              DO delete o penultimo ponto de L SUP
      Coloque os pontos pn e pn-1 na lista L INF
07.
08.
      FOR i: n-2 DOWNTO 1
09.
              Repita passoas 4-6 com L INF no lugar de L SUP
10.
      Remover o primeiro e ultimo pontos de L_SUP e L_INF
11.
      Devolve concatenacao de L SUP e L INF
```

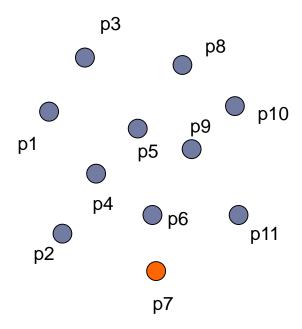
ENVOLTÓRIA CONVEXA [GIFT WRAPPING – JARVIS, 1973]

https://doi.org/10.1016%2F0020-0190%2873%2990020-3

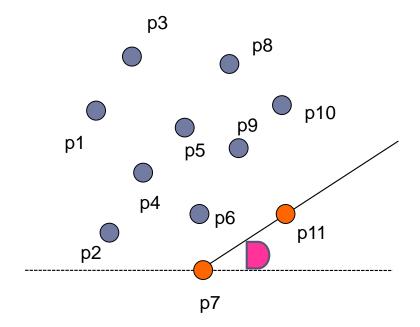
- 1. Find point with smallest y coordinate
- 2. Find edge that minimizes counter-clockwise angle with respect to supporting line
- 3. Repeat until reach last vertex



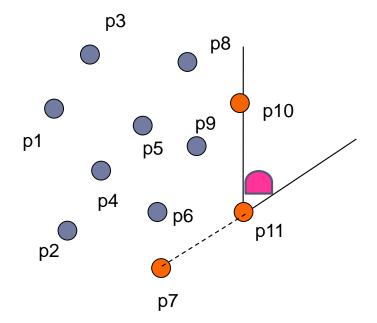
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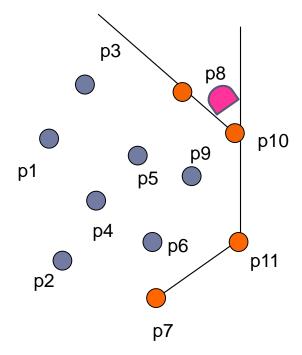
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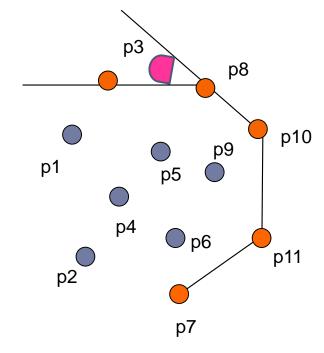
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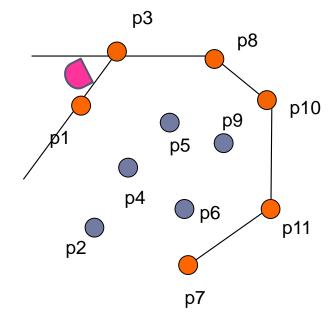
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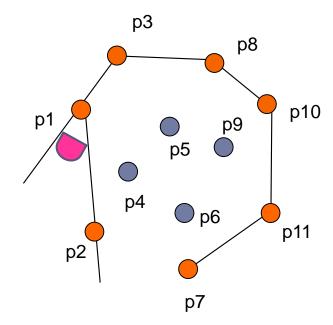
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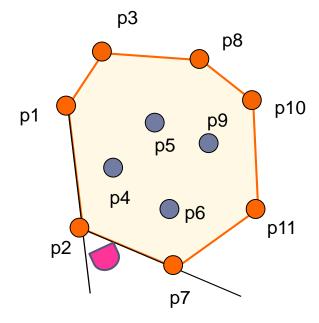
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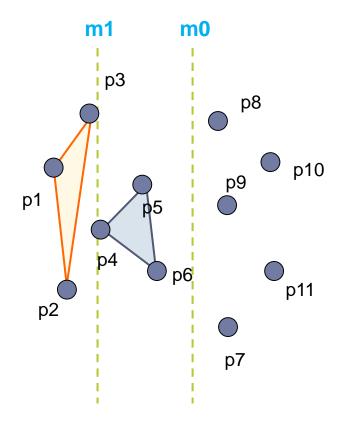
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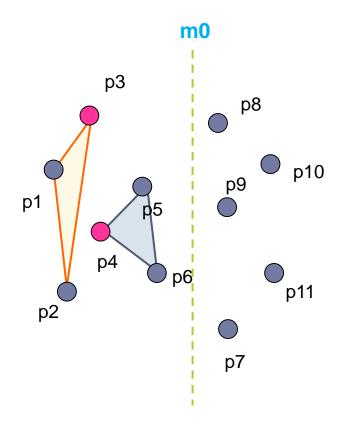
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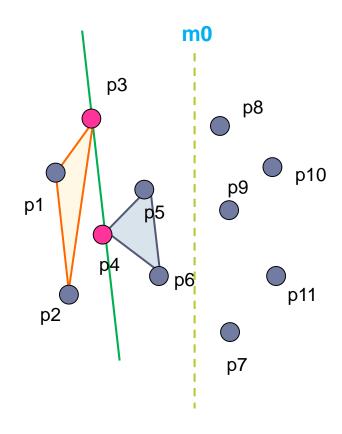
- 1. Sort point in x
- 2. Divide points in two sets A and B with similar number of elements
- 3. Compute Convex Hull of A and B recursively
- 4. Merge A and B to define convex hull



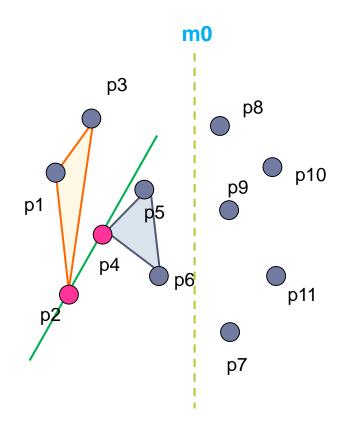
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- 5. Find lower tangent common to A and B
- 6. Find rightmost point a of A and leftmost point b of B
- 7. While ab is not lower tangent to a, decrease a
- 8. While ab is not lower tangent of b, decrease b



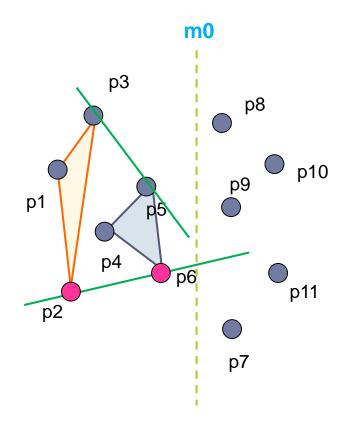
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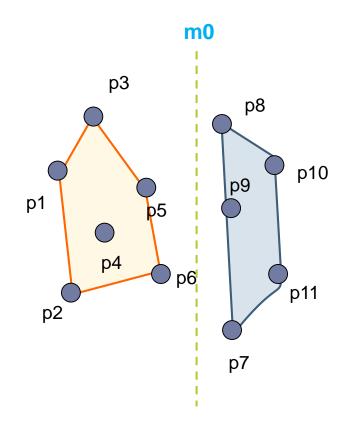
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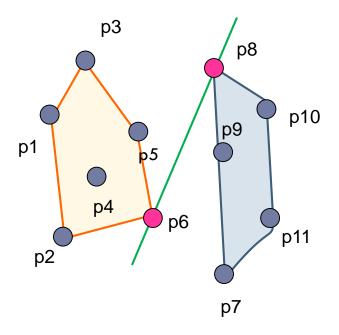
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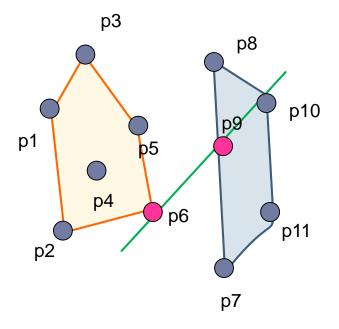
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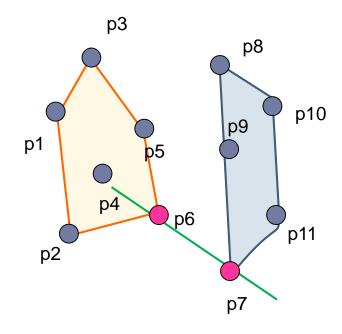
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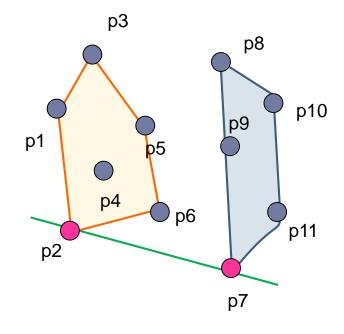
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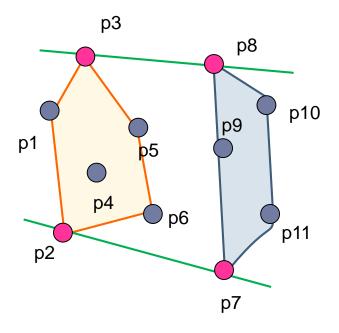


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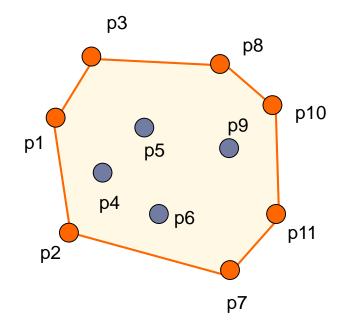


ENVOLTÓRIA CONVEXA [DIVIDE-AND-CONQUER – PREPARATA, 1977]

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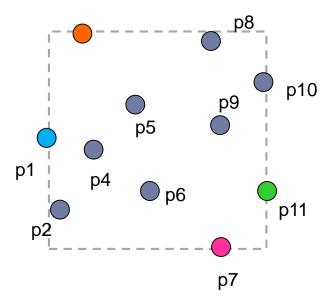
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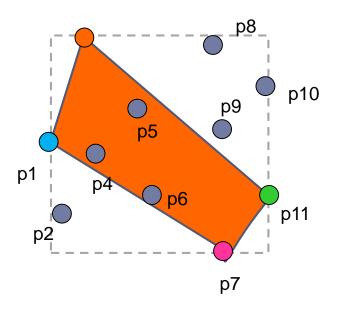
ENVOLTÓRIA CONVEXA [QUICKHULL – BARBER, 1996]

https://doi.org/10.1145%2F235815.235821

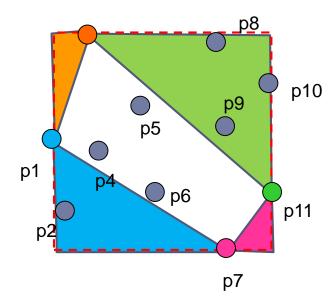
- 1. Find extreme points above, below, left and right
- 2. Define a quadrilateral connecting extreme points, discard internal points
- 3. Process 4 triangles recursively
- 4. For each triangle, find vertex with greater distance from the triangle baseline
- Define another triangle and discard internal triangles
- 6. Process 2 triangles recursively



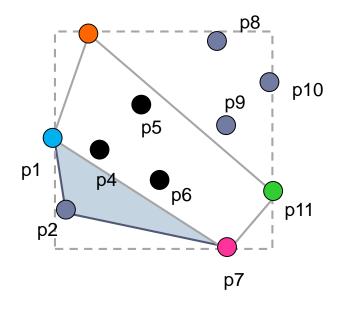
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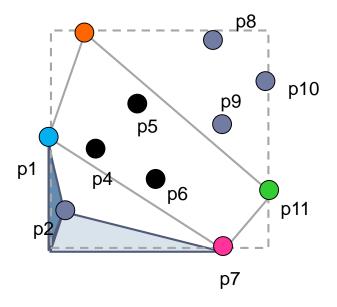
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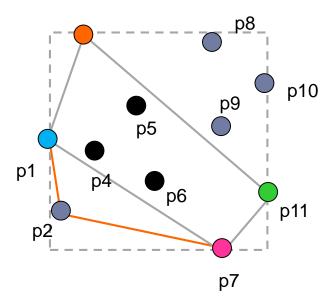
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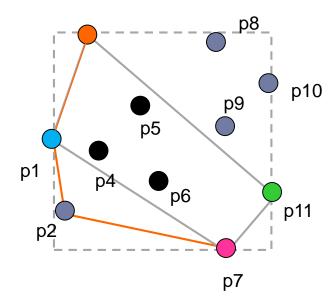
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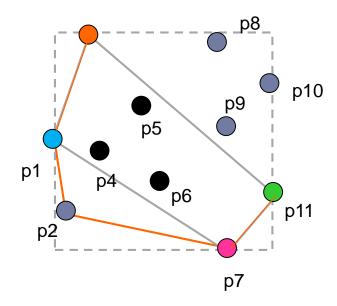
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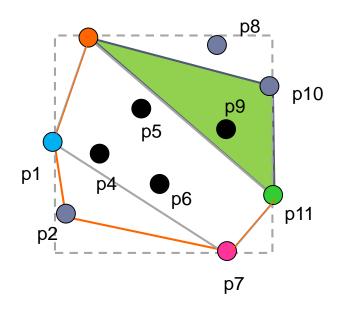
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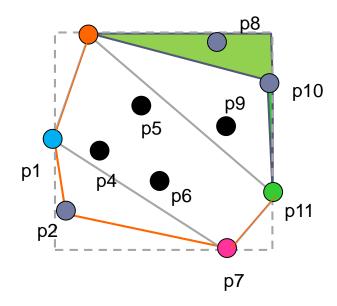
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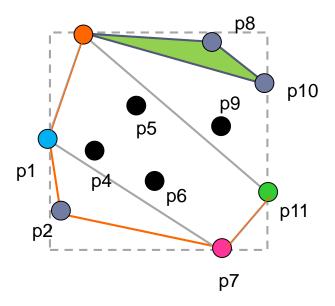
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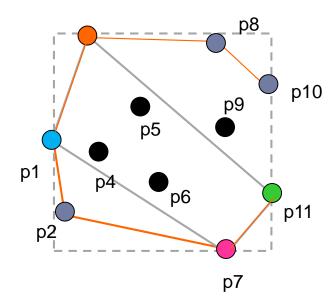
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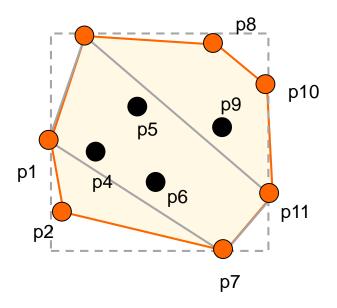
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Exercício

10078 - The Art Gallery

https://uva.onlinejudge.org/index.php?option=onlinejudge&page=show_problem&problem=1019

Cálculo do produto vetorial.