

# 1 Encoding

Let  $t_i$  denote the  $k$  component types,  $c_i$  the components of each type. Let  $C = \{\}$  denote the set of all components. For all components  $c_{k_1}, \dots, c_{k_n}$  of each of the  $k$  sorts it must be specified that the component variables are distinct. Every component slot in an assignment corresponds to one decision variable. Let  $x_{a,s}$  denote these decision variables. Let  $\mathcal{M} : x_{a,s} \rightarrow c_i$  denote the model assignment from each assignment slot variable to exactly one component.

## 1.1 List of Conditions

```
\hline
ComponentIs( AssignmentID , SlotID , Component )
ComponentIn( AssignmentID , SlotID , Component ... )
SameComponent( Assignments ... , SlotID )
InGroup( assignment , slotID , groupID )
TagEqual( Assignment , SlotID , TagID , Value )
TagEqualAbove( assignment , slotID , tagID , value )
TagEqualBelow( assignment , slotID , tagID , value )

Not( Condition )
And( Condition ... )
Or( Condition ... )
Implies( Condition , Condition )
Xor( Condition , Condition )
Iff( Condition , Condition )

MaxAssignment( Assignment ... , Condition ... )
MinAssignment( Assignment ... , Condition ... )
MaxConsecutive( Assignment ... , OrderedComponentSlot , Condition ... )
MinConsecutive( Assignment ... , OrderedComponentSlot , Condition ... )
MaxInSequence( Assignment ... , OrderedComponentSlot , Condition ... )
MinInSequence( Assignment ... , OrderedComponentSlot , Condition ... )

Simultaneous( TimedComponent ... )
Consecutive( TimedComponent1 , TimeComponent2 )

Only( Condition ... )
Exclude( Condition ... )
Include( Condition ... )
```

**ComponentIs(AssignmentID, SlotID, Component)** *component* is assigned to the slot *SlotID* of assignment *AssignmentID*

**ComponentIn(AssignmentID, SlotID, Component...)** one *component* in the set of  $k \geq 1$  components given is assigned to the slot *SlotID* of as-

signment *AssignmentID*. The given components need to be of the type required by that slot.

**SameComponent(Assignments..., SlotID)** The components  $c_1, \dots, c_k$  of slot *slotID* of the  $k$  assignments given are identical. The given assignments need to have a slot *slotID*.

**InGroup(assignment, slotID, groupID)** The component  $c$  assigned to slot *slotID* of assignment *AssignmentID* is in group *groupID*.

**TagEqual(Assignment, SlotID, TagID, Value)** The tag *TagID* of the component  $c$  assigned to slot *SlotID* of assignment *AssignmentID* is equal to *value*.

**TagEqualAbove(assignment, slotID, tagID, value)** The tag *TagID* of the component  $c$  assigned to slot *SlotID* of assignment *AssignmentID* is equal to or greater than *value*.

**TagEqualBelow(assignment, slotID, tagID, value)** The tag *TagID* of the component  $c$  assigned to slot *SlotID* of assignment *AssignmentID* is equal to or smaller than *value*.

**Not(Condition)**

**And(Condition...)**

**Or(Condition...)**

**Implies(Condition, Condition)**

**Xor(Condition, Condition)**

**Iff(Condition, Condition)**

**MaxAssignment(Assignment..., Condition...)**

**MinAssignment(Assignment..., Condition...)**

**MaxConsecutive(Assignment..., ?Ordering?, Condition...)**

**MinConsecutive(Assignment..., ?Ordering?, Condition...)**

**MaxInSequence(Assignment..., ?Ordering?, Condition...)**

**MinInSequence(Assignment..., ?Ordering?, Condition...)**

**Simultaneous(TimedComponent...)**

**Consecutive(TimedComponent1, TimeComponent2)**

**Only(Condition...)**

**Exclude(Condition...)**

## **Include(Condition...)**

Conditions are used to 1. filter assignments, 2. be filled out to be added as rules

Rule: Require(Condition)

These are written out as a language and need to be parsed

By component here I mean the variables in assignments - that should probably be named better.

Basic Condition: Component = Component — Component != Component  
— InGroup(Component, Group) — InSet(Component, Set) — Not(Condition)  
— Condition AND Condition — Condition OR Condition — Condition IMPLIES Condition — Condition IFF Condition — Condition XOR Condition

Numeric Condition = Numeric ; Numeric — Numeric != Numeric — Numeric < Numeric — Numeric <= Numeric — Numeric = Numeric — Numeric != Numeric

Numeric := Int — Tag(Component, Tag) — NumAssigned(Condition) — MaxConsecutive(Condition) — MinConsecutive(Condition) — MaxBreak(Condition) — MinBreak(Condition)

Condition, Component, Group

need: determine whether a thing is a condition or component

2 arguments or a list of arguments?

Operator precedence: NOT AND OR IFF, IMPLIES, XOR —————

Precedence climbing:

primary := '(' expression ')' — variable

————— Instead of numeric: AtLeast(Int, Bitvector) — AtMost(Int, Bitvector) — Exactly(Int, Bitvector)

A bitvector is created with a from an ordered set of assignments and set of conditions.

CreateBitvector(set, set)

from assignments or component slots?

—

When a translator receives a rule, it will need to recursively(?) encode each part.

This means that for each relevant combination of components, a sentence needs to be added to the solver. - build all constraints sequentially or at the same time? but how? -

Example: A1.Shift.StartTime == A2.Shift.StartTime IMPLIES A1.Nurse != A2.Nurse

(A1.Shift.StartTime == A2.Shift.StartTime) IMPLIES (A1.Nurse != A2.Nurse)

A1.Shift.StartTime == A2.Shift.StartTime IMPLIES A1.Nurse != A2.Nurse

Implies condition: only need to add rules for which the premise is true This rule only uses basic logical operations. Instantiate it with all combinations of A1 and A2

Nurses: Patrick Andrea Stefaan Sara Nguyen

InGroup() or

## 1.2 Rules

Rules define constraints between assignments.

```
SameTime(A1, A2) =i A1.Component("Nurse") != A2.Component("Nurse")
"SameTime(A1, A2)" =i "A1.Component("Nurse")" != "A2.Component("Nurse")"
—SameTime(A1, A2)— =i —A1.Component("Nurse")— != —A2.Component("Nurse")—
```

bool op obj op obj

```
bool !=(obj, obj) bool =(obj, obj)
```

```
bool =i(bool, bool)
```

```
bool SameTime(Ass, Ass)
```

```
object Component()
```

A rule is a function taking either 2 or n assignments, returning bool

i want to save parts recursively

every function part of a function needs to depend on assignments

LEFT OPERATOR RIGHT

Left and right can be functions or rules

the rule object provides a function for

```
rule := expression op expression expression := function — rule
```

Formula: wrapper class for some proper formula or: function to invoke for every translatable unit

information that needs to be kept: assignment between external and internal

ID

elements might be added to groups after group-wide constraints are created

- When a create constraint is called, create a constraint object

- Go through all constraint objects and call generate on them all - Create

component objects and constraint objects - generate constraints

-

1. create a problem object with a name/id
2. create a class for each type of component (can use using)
3. create a class for each type of assignment (must derive, is abstract)
4. create groups and tags
5. create components, with appropriate groups and tags
- 6.

## 2 Problem

Let there be  $k$  types of values,  $n = (n_1, \dots, n_k)$  values, which are assumed to be distinct and  $m$  slots which may be optional and associated with a penalty. Each value and each slot are of exactly one type. A set of constraints may apply that either limit the set of possible solutions or impose a penalty on solutions violating them. The problem considered here consists of finding an assignment

of a value to each slot so that both the type requirement and any additional hard constraint is satisfied and the cost incurred by violated soft constraints is minimized.

/\*

## 2.1 Model

### Problem

The problem is assigning to each of  $n$  assignment slots one of  $m$  components. Components have a type and the component needs to fit that of the assignment slot.

Problem

```

Tags
Groups
ComponentTypes

Assignments
  ComponentSlots
    ComponentType
    Fixed , Component
  Optionality
  Weight

Components
  ComponentType
  Groups
  Tags

Rules
  AssignmentSet
  Weight
  Optionality
  TopCondition
    ConditionType
    Subconditions
    Attributes

```

#### 2.1.1 Components

The basic elements of each scheduling problem are the **components**. In a university timetabling problem these could be events, rooms and teachers. Each component has a *component type*  $\mathcal{T}_C$  and a finite domain  $\mathcal{D}_C$ . Component types can be boolean, discrete and finite numeric types (?) or custom types. Components are assumed to be unique:  $i \neq j \implies c_i \neq c_j$ . The domain

of a component is, by default, the set of all values of that type. A problem specifies a set of groups and a set of tags. Each component has for each tag a corresponding value  $tag_{j,i} \in \mathbb{N}_0$ . For each group  $group_j$  and component  $c_i$ , it holds that either  $group_{j,i}$  or  $\neg group_{j,i}$ . As an example, a staff scheduling problem could have a component type  $\mathcal{T}_{nurse}$  with the corresponding domain  $\mathcal{D}_{nurse} = \{Anna, Ben, Charlie\}$ . Component types are called *ordered* if they define a strict weak ordering on their domain.

### 2.1.2 Assignments

**Assignments** are sub-problems of the scheduling problem, in which a number of components of different types need to be combined: for example, an assignment of a room, a time, a class and students.

Each assignment  $a_i$  defines a sub-problem, in which . It consists of slots  $as_{i,1}, \dots, as_{i,k}$  that each specify a component type  $type_{i,j}$ , a number of components to fill that slot  $\in \{1, k \in \mathbb{N}, n\} \subset \mathbb{N}$ . Additionally it specifies whether that slot is optional  $optional_{i,j}$  and, if so, a weight  $weight_{i,j}$ .

An assignment  $a_i$  consists of  $k \geq 1$  component slots  $as_{i,1}, \dots, as_{i,j}$ .

Slots can be *fixed*, meaning

In a nurse rostering problem, an assignment  $a_i$  could consist of a  $type_{i,1} =$  Assignments are called *ordered* if they have at least one component slot of an ordered component type.

In the university course timetabling as presented in ITC19 [?], an assignment could consist of a single fixed *course* slot, a single non-optional *roomTime* slot and as many optional *student* slots as the course capacity allows.

### 2.1.3 Condition

A condition as an object holds information that a . A set of assignments can be evaluated over a condition. The number of assignments depend on the condition type.

A **basic condition** is one that does not

A **composite condition** is a condition the evaluation of which depends on at least one other condition. Such a chain of conditions will always end with basic

### 2.1.4 Rules

Additionally to the assignments, a problem has **rules**, specifying the relationship between individual assignments.

To apply a condition, they need to be associated with a set of assignments over which they are meant to hold, as well as information necessary for th

A rule connects a chain of conditions to a set of combinations of assignment over which it should be enforced.

When a rule is generated, it must be generated for each viable combination of assignments.

### 2.1.5 Model

A model is an assignment of components to assignment slots such that no hard rule is violated and the sum of penalties of .

### 2.1.6 Basic Conditions

**ComponentIs(Assignment, Slot, Value)** **formal**  $x_{a,s} = c_j$   
**smt2**  $(= x_{a,s} c_j)$

**ComponentIn(Assignment, Slot, Value...)** **formal**  
**smt2**

$$\bigvee_{c_j \in values} s_{a_i} = c_j$$

**SameComponent(Assignment..., Slot)** For each combination of two assignments  $(a_i, a_j), i \neq j$ , add the constraint

$$x_{a_i, slot} = component1 \wedge x_{a_j, slot} = component2 \implies component1 \neq component2$$

**formal**

**smt2**

**InGroup(Assignment, Slot, Group)** This condition can be encoded by a constraint that limits the domain of the slot variable to the component variables that fulfil the group condition. Let *comps* be the component variables corresponding to exactly those components

$$\bigvee_{c \in allComponents}$$

**TagEqual(assignment *a*, slot *s*, tag, value)** Like InGroup, this condition can be handled by passing a constraint that limits the domain of assignment variable. Let *comps* be the component variables corresponding to exactly those components  $c_i$  where  $tag_{c_i} = value$ .

$$\bigvee_{c \in comps} x_{a,s} = c$$

**TagEqualAbove(assignment, slotID, tagID, value)** Let *comps* be the component variables corresponding to exactly those components  $c_i$  where  $tag_{c_i} \geq value$ .

$$\bigvee_{c \in comps} x_{a,s} = c$$

**TagEqualBelow(assignment, slotID, tagID, value)** Let *comps* be the component variables corresponding to exactly those components  $c_i$  where  $tag_{c_i} \leq value$ .

$$\bigvee_{c \in comps} x_{a,s} = c$$

### 2.1.7 Composite Conditions

*Condition* used inside *..* signifies the evaluation of the *.*

**Not(Condition)**

$$\neg Condition$$

**And(Condition...)**

$$\bigwedge_{c \in Conditions} c$$

**formal**

**smt2**

**Or(Condition...)**

$$\bigvee_{c \in Conditions} c$$

**Implies(Condition1, Condition2)**

$$Condition1 \implies Condition2$$

**formal**

**smt2**

**Xor(Condition1, Condition2)**

$$(\neg Condition1 \wedge Condition2) \vee (Condition1 \wedge \neg Condition2)$$

**formal**

**smt2**

**Iff(Condition1, Condition2)**

$$(Condition1 \implies Condition2) \wedge (Condition2 \implies Condition1)$$

**formal**

**smt2**



**MaxAssignment(Assignment..., Condition...)**

**MinAssignment(Assignment..., Condition...)** Add a bitvector  $bv_i$  of the length  $\#Assignments$  For each assignment  $a_k$ , add a rule of the form

$$\bigwedge_{c \in \text{Condition}} c(a_i) \implies \text{extract}(a_k)$$

[?] describe a declarative , the number of set bits can be checked by

$$bv = bv \& (bv - 1)$$

If this constraint is violated, the penalty should be multiplied by the Hamming-Weight of the resulting vector.

**MaxConsecutive(Assignment..., OrderedComponentSlot, Condition...)**

**MinConsecutive(Assignment..., OrderedComponentSlot, Condition...)**

All given assignments must have a slot of the *OrderedComponentSlot* type. This type needs to be an ordered type. The  $i$ th bit of the generated bitvector corresponds to the  $i$ th assignment with respect to the given order. Assignments that are equivalent with respect to the given order are *collapsed* into a bit that is set if the conditions hold for at least one of the equivalent assignments. Let  $Conditions(A)$  denote the conjunction of the given conditions resolved for the assignment  $A$ . For every given assignment, add:

$$Conditions(A) \implies bv.\text{extract}(i)$$

As described in [?], the following equation will hold iff there exists no block of a length more than  $max$ :

$$0 = bv \& (bv >> 1) \& (bv >> 2) \& \dots \& (bv >> max)$$

The same paper describes a three-step algorithm to check minimal block length by first

**Simultaneous(Assignment..., OrderedComponentSlot)** Each of the given assignments needs to have a slot of the *OrderedComponentSlot* type and that type needs to be ordered.

**Consecutive(Assignment..., OrderedComponentSlot)**

**Only(Assignment, Condition...)**

### 3 Modelling The Second International Nurse Rostering Competition (2014)

In short, the scheduling problem presented in [?] consists of assigning nurses to shifts in multiple consecutive planning periods. The competition was held in 2015, on the website the benchmarks and a solution validator are still available.

This is an example for a scenario (global part of the problem description) from a test instance:

```
SCENARIO = n005w4
```

```
WEEKS = 4
```

```
SKILLS = 2  
HeadNurse  
Nurse
```

```
SHIFT_TYPES = 3  
Early (2,5)  
Late (2,3)  
Night (4,5)
```

```
FORBIDDEN_SHIFT_TYPES_SUCCESSIONS  
Early 0  
Late 1 Early  
Night 2 Early Late
```

```
CONTRACTS = 2  
FullTime (15,22) (3,5) (2,3) 2 1  
PartTime (7,11) (3,5) (3,5) 2 1
```

```
NURSES = 5  
Patrick FullTime 2 HeadNurse Nurse  
Andrea FullTime 2 HeadNurse Nurse  
Stefaan PartTime 2 HeadNurse Nurse  
Sara PartTime 1 Nurse  
Nguyen FullTime 1 Nurse
```

Only one time unit is necessary: a counter of days since the start of the planning horizon. Create groups: HeadNurse, Nurse, Early, Late, Night, Saturday1, Sunday1, Saturday2, Sunday2, Saturday3, Sunday3, Saturday4, Sunday4. Nurse can be used as an alias for Agent, Shift for TimedTask.

```
using Nurse = Agent<std::string, std::string, std::string>;  
using Shift = TimedTask<std::string, std::string, std::string, int>;
```

Create each nurse, add the appropriate groups for roles. When reading from a file, this can be done in a loop; here only exemplary for a single nurse:

```
Nurse patrick("Patrick");
patrick.addGroup("HeadNurse");
patrick.addGroup("Nurse");
```

There is only a single type of assignment for this problem. Since the problem is to assign nurses *to* shifts, each shift will be fixed in one assignment.

```
class Shift : public Assignment {

    Expression generate() override;
    Agent nurse;
    const TimedTask shift;
    std::vector<Rule> requirements;

};
```

Create a TimedTask for every role required for each shift. Add the appropriate groups for role requirements, shift type and day of the week. Add these requirements to all Shifts *s*:

```
s.Require(shift.inGroup(HeadNurse) IMPLIES nurse.inGroup(HeadNurse));
s.Require(shift.inGroup(Nurse) IMPLIES nurse.inGroup(Nurse));
```

Ideas to model the constraints as described in [?] as rules:

- H1. Single assignment per day**  $A1.Simultaneous(A2) \text{ IMPLIES } A1.Nurse \neq A2.Nurse$ , or  
 $A1.Shift.StartTime == A2.Shift.StartTime \text{ IMPLIES } A1.Nurse \neq A2.Nurse$
- H2. Under-staffing** This is already enforced by making required tasks non-optional.
- H3. Illegal shift type successions** For every illegal combination (prec, succ):  
 $InGroup(A1.Shift, prec) \text{ AND } InGroup(A2.Shift, succ) \text{ AND } ImmediateSuccessor(A1.Shift, A2.Shift) \text{ IMPLIES } A1.Nurse \neq A2.Nurse$
- H4. Missing required skill** Enforced by making required role tasks non-optional.
- S1. Insufficient staffing for optimal coverage** Enforced by adding the appropriate weight (30) as penalty to optional shifts.
- S2. Consecutive assignments** Define  
 $SameType(A1, A2) := (A1.shift.inGroup(Early) \text{ AND } A2.shift.inGroup(Early))$   
 $OR (A1.shift.inGroup(Late) \text{ AND } A2.shift.inGroup(Late)) \text{ OR } (A1.shift.inGroup(Night) \text{ AND } A2.shift.inGroup(Night))$   
 $MinConsecutive(A1.nurse == A2.nurse \text{ AND } SameType(A1, A2)) \text{ } i = MIN$
- S3. Consecutive days off**  $MaxConsecutive(A1.nurse == A2.nurse \text{ AND } SameType(A1, A2)) \text{ } i = MAX$

S4. Preferences

S5. Complete week-end

S6. Total assignments (only evaluated at the end of the planning period)

S7. Total working week-ends (only evaluated at the end of the planning period)