

Load balancing of optimization algorithms

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Motivation

- large-scale optimization problems require a significant amount of computational resources
- multiple optimization algorithms should be able to share the same computational resources
- hardware cost reduction



Thesis goals

- formalize load balancing problem of the optimization algorithms
- propose a load balancing algorithm
- design and implement the load balancer



State-of-the-art solutions

- mainly generic
- static/dynamic load balancing
- focused on load balancing of short-running tasks
 - handling HTTP request

 no specific domain knowledge

Formalization



Problem formalization

- formalized as an integer linear programming problem in section 3.1
- specified input/output variables
- resources assignment defined as binary assignment ${}^r x_t^j = \{0, 1\}$
- various constraints
 - execution cost maximum $\forall t, j : P^j \leq C_t^j \implies \sum_{t+1}^{\infty} \sum_{r \in R} {}^r x_t^j = 0$
 - execution time maximum $\forall t, j : D^j \leq t \implies \sum_{t+1}^{\infty} \sum_{r \in R} {}^r x_t^j = 0$
 - resources reallocation $\sum_{r \in R} {}^r x_t^j = 1 \implies \forall r \in R : {}^r x_t^j - {}^r x_{t+1}^j \geq 0$
- optimization criteria $\max crit_t = \alpha \sum_{j \in J} S_t^j - (1 - \alpha) \sum_{j \in J} C_t^j \quad 0 \leq \alpha \leq 1$



Proposed load balancing algorithm

Input: Q - queue with jobs to schedule

```
1  $Q$ : jobs queue;  
2  $J$ : set of jobs;  
3  $P$ : predictions;  
4  $E$ : execution plan;  
5  $J \leftarrow Q.poll()$ ;  
6  $J' \leftarrow J.filter(job \rightarrow job.D > job.T \wedge job.P > job.C)$ ;  
7  $J'' \leftarrow convert(J')$ ;  
8  $P \leftarrow predict(J'')$ ;  
9  $E \leftarrow schedule(J'', P)$ ;  
10  $E' \leftarrow convert(E)$ ;
```

Result: Q is empty

Output: E' - execution plan

Algorithm 1: Load balancing algorithm



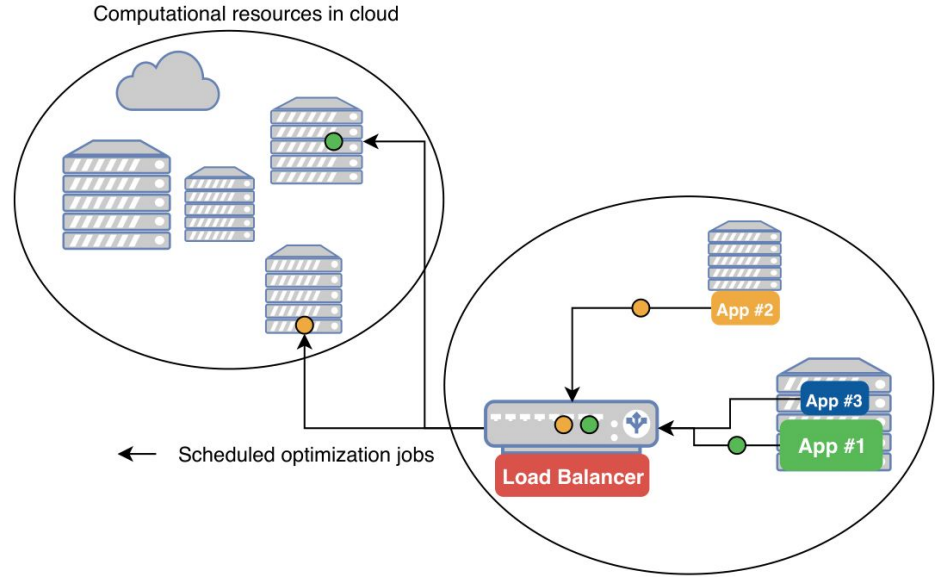
Optimization algorithm

Load balancing decisions

- heuristic algorithm
- OptaPlanner

Algorithm values prediction

- ${}^r\Delta_t^j = {}^r|v_t^j - v_{t-1}^j| \cdot {}^rx_t^j$
- hyperbola time-series fitting
- Levenberg–Marquardt algorithm

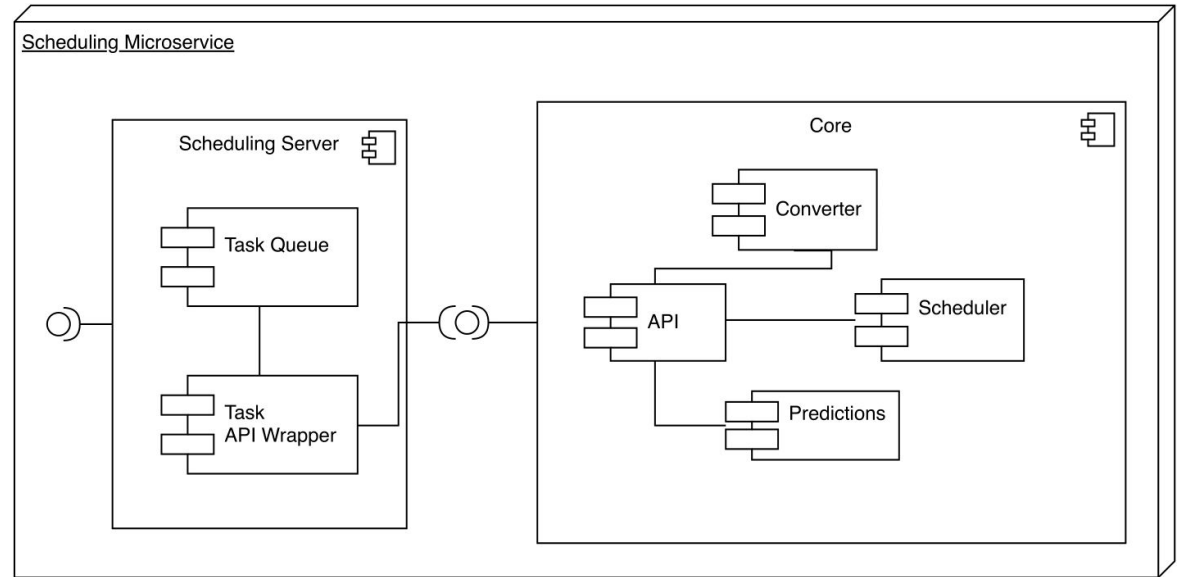


Architecture



Application architecture

- microservice design
- separation of the scheduling core
- core
 - custom data representation
 - optimization engine
- server
 - uses core API
 - exposes scheduling API





Implementation

- Java Virtual Machine
- Kotlin
- Docker
- Docker Compose
- Ktor

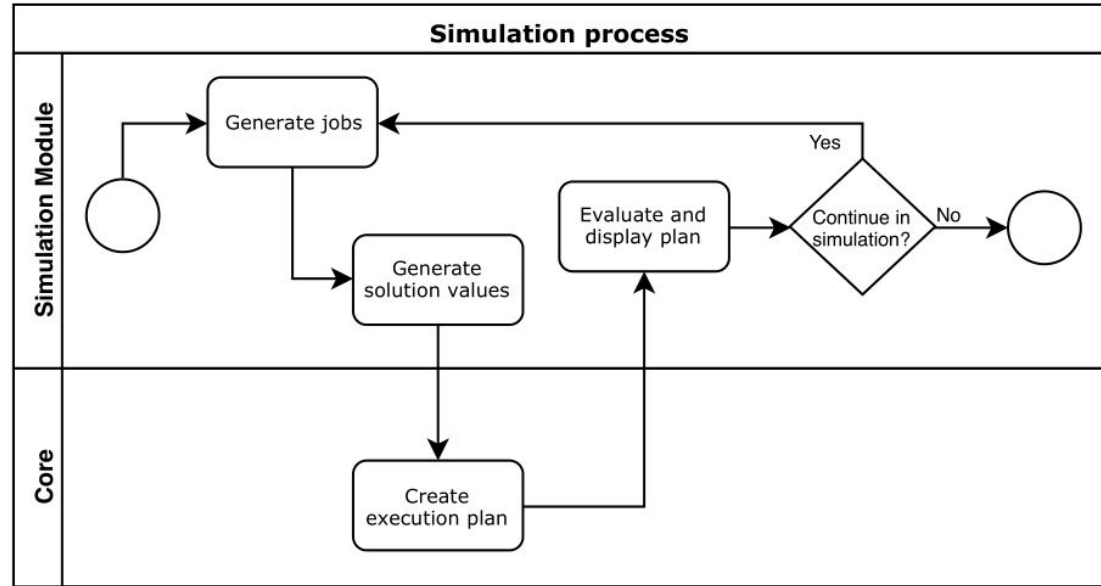


Routes Discovery Library



Simulations

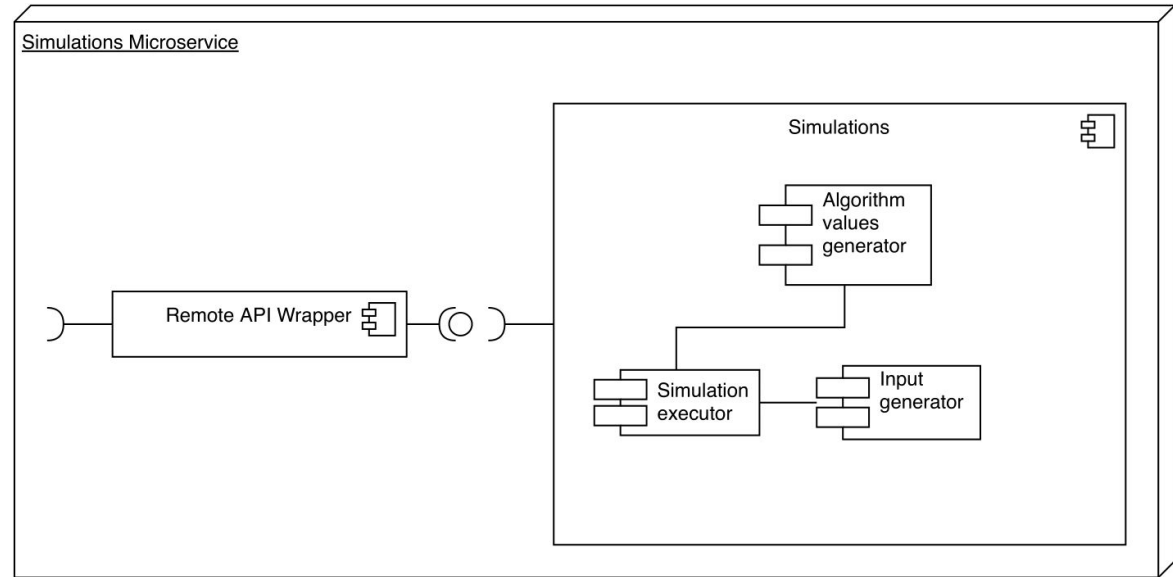
- implemented in separate module
- simulation control
- input generators
- local/remote simulations
- random parameters generation
- objective function values created by instances of heuristics algorithms





Simulations

- simulation microservice
- simulates real-world usage
- uses the same simulation core
- executed inside Docker Compose



Conclusion



Future work

Infrastructure

- execution module development
- optimize resource usage
- fine-tuning of the optimization engine

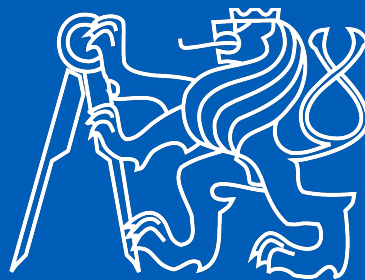
Routes discovery library

- refactoring
- usage of the generic dependency injection framework
- publish under an open-source license



Conclusion

- traditional load balancing strategies are not optimal
- domain specific load balancing of optimization algorithms
- problem formalization
- design and implementation of the load balancer
- simulations and evaluation of the load balancer



Thank you for your attention!

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Jaký je výpočetní overhead load-balancing platformy samotné, na kolik zatěžuje rozvrhovací algoritmus platformu samotnou?

- jedná se o další optimalizační úlohu navíc
 - tedy zatěžuje minimálně jeden zdroj dle definice
 - zatížení pouze ve chvíli, kdy algoritmus plánuje
- algoritmus samotný je paralelizovatelný
 - počet vláken, které může algoritmus použít je vstupem



V jaké škále a od jaké velikosti rozvrhovaných jobů dává platforma smysl?

- klasické techniky rozvrhování zátěže selhávají
- není možné uspokojit všechny minimální nároky algoritmů najednou
 - v jednoduchých případech možné řešit prioritní frontou

$$\exists t : \sum_J j_t^{min} > |R_t| \text{ kde } |R_t| \geq 5$$



Objective function over iterations

