Specifications of Implemented Refactorings

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This document collects the pseudo-code specifications of all refactoring implemented in our engine. **Note:** This is work in progress; some specifications are missing, and not all implementations agree completely with the specifications.

1 Pseudocode Conventions

We give our specifications in generic, imperative pseudocode. Parameters and return values are informally typed, with syntax tree nodes having one of the types from Fig. ??. Additionally, we use an ML-like `option` type with constructors `None` and `Some` for functions that may or may not return a value.

Where convenient, we make use of ML-like lists, with list literals of the form $[1; 2; 3]$ and $|xs|$ indicating the length of list xs.

The names of refactorings are written in SMALL CAPS, whereas utility functions appear in `monospace`. A list of utility functions with brief descriptions is given in Fig. ??. An invocation of a refactoring is written with floor-brackets ⌊LIKE THIS⌋() to indicate that any language extensions used in the output program produced by the refactoring should be eliminated before proceeding.

We write $A <: B$ to mean that type A extends or implements type B, and $m <: m'$ to mean that method m overrides method m'.

2 The Refactorings

2.1 Convert Anonymous to Local

This refactoring converts an anonymous class to a local class. Implemented in `TypePromotion/AnonymousClassToLocalClass.jrag`; see Algorithm ??.

2.2 Convert Anonymous to Nested

This refactoring converts an anonymous class to a member class. Implemented in `TypePromotion/AnonymousClassToMemberClass.jrag`; see Algorithm ??.

Note: the implementation additionally handles the case where A occurs in a field initialiser.

Algorithm 1 Convert Anonymous to Local(A : *AnonymousClass*, n : *Name*) : *LocalClass*

Require: Java
Ensure: Java $\cup$ locked names

$c \leftarrow$ class instance expression containing A
$d \leftarrow \lfloor$Extract Temp$\rfloor(c,$ `unCapitalise`$(n))$ – not possible to do!!!
$b \leftarrow$ enclosing body declaration of d
`lockNames`(b, n)
convert A to class named n, remove it from c
Insert Type(b, A)
lock type access of c to A
Inline Temp(d) – but without checks (TODO?)
return A

Algorithm 2 Convert Anonymous to Nested(A : *AnonymousClass*, n : *Name*) : *MemberType*

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Require: Java
Ensure: Java

$L \leftarrow$ Convert Anonymous to Local(A, n)
return Convert Local to Member Class(L)

2.3 Convert Local to Member Class

This refactoring converts a local class to a member class. Implemented in `TypePromotion/LocalClassToMemberClass.jrag`; see Algorithms **??**, **??**, **??**.

2.4 Extract Class

This refactoring extracts some fields of a class into a newly created member class. Implemented in `ExtractClass/ExtractClass.jrag`; see Algorithm **??**.

Initializer evaluation is order independent for example if we can invert their order without breaking name and dataflow dependencies.

This is only a bare-bones specification. The implementation additionally allows to encapsulate the extracted fields, and to move the wrapper class W to the toplevel.

2.5 Extract Constant

This refactoring extracts a constant expression into a field. Implemented in `ExtractTemp/ExtractConstant.jrag`; see Algorithm **??**.

An expression is extractible if its type is not `void`, it is not a reference to a type or package, and it is not the keyword `super`; furthermore, it cannot be on the right-hand side of a dot.

Algorithm 3 CONVERT LOCAL TO MEMBER CLASS(L : *LocalClass*) : *MemberType*

Require: Java
Ensure: Java $\cup$ locked names, fresh variables

```
A ← enclosing type of L
closeOverTypeVariables(L)
closeOverLocalVariables(L)
if L is in static context then
   make L static
end if
lockNames(name(L))
lock all names in L
remove L from its declaring method
INSERT TYPE(A, L)
```

Algorithm 4 `closeOverTypeVariables`(L : *LocalClass*)

```
m ← empty map
U ← accesses to L
for all accesses V to type variables T of the enclosing body declaration do
   if m(T) undefined then
      create new type variable T′ with same bounds as T
      add T′ as type parameter to L
      m(T) ← T′
      for all u ∈ U do
         add locked access to T as type argument to u
      end for
   end if
   lock V onto m(T)
end for
```

Algorithm 5 `closeOverLocalVariables`(L : *LocalClass*)

```
m ← empty map
for all accesses v to local variables x of enclosing body declaration do
  if m(x) undefined then
    create private final field f of same type as x
    add f to L
    m(v) ← f
    for all constructors c of L do
      create new parameter p of same type and name as x
      insert p as first parameter of c
      if c is chaining constructor then
        add access to p as parameter to chaining invocation
      else
        insert assignment from p to f as first statement in c
      end if
    end for
    for all instantiations i of L do
      insert access to x as first argument to i
    end for
  end if
  lock v onto m(x)
end for
```

The *effective type* of an expression e is the same as the type of e, except when the type of e is an anonymous class, in which case the effective type is its superclass, or when the type of e is a captured type variable, in which case the effective type is its upper bound.

2.6 Extract Method

See ECOOP 2009 publication. (TODO)

2.7 Extract Temp

This refactoring extracts an expression into a local variable. Implemented in `ExtractTemp/ExtractTemp.jrag`; see Algorithms **??**, **??**, **??**, **??**.

2.7.1 Insert Local Variable

The refactoring inserts a local variable before a given statement. Implemented in `ExtractTemp/IntroduceUnusedLocal.jrag`.

Algorithm 6 EXTRACT CLASS(C : *Class*, *fs* : `list` *Field*, n : *Name*, *fn* : *Name*)

Require: Java

Ensure: Java ∪ locked names, locked dataflow, first-class array init

$v \leftarrow$ maximum visibility of any of the *fs*
$W \leftarrow$ new `static` class of name n with visibility v
INSERT TYPE(C, W)
$w \leftarrow$ new field of type W and name *fn*, initialised to a new instance of W
INSERT FIELD(C, w)
for all $f \in fs$ **do**
    assert f is not static
    for all uses v of f **do**
        qualify v with a locked access to w
    end for
    if f has initialiser **then**
        split field declaration and initializer, leaving initializer in initializer block after
    end if
    remove f
    INSERT FIELD(W, f)
end for
inits $\leftarrow$ {initializers of *fs*}
for all *init* $\in$ *inits* **do**
    lock names and dataflow
    move *init* after already moved initializers (possibly w)
    `try` unlocking names and dataflow
    if unlocking was successful **then**
        `continue`
    else
        move *init* back and `break`
    end if
end for
in W create default constructor and constructor for initializing all fields
if all *inits* were moved **and** initializer evaluation is order independent **then**
    change the constructor call for w to initialize the fields and remove *inits*
else
    merge consecutive *inits* to common initializer blocks
end if

Algorithm 7 EXTRACT CONSTANT($e : \mathit{Expr}, n : \mathit{Name}$)

Require: Java
Ensure: Java $\cup$ locked names, locked dataflow

assert e is extractible
$A \leftarrow$ enclosing type of e
$t \leftarrow$ effective type of e
$f \leftarrow$ new `private` (`public` if A is an interface) `static` `final` field of type t and name n
INSERT FIELD(A, f)
lock names, flow, and synchronisation of e
set initialiser of f to e
replace e with locked access to f

Algorithm 8 EXTRACT TEMP($e : \mathit{Expr}, n : \mathit{Name}$)

Require: Java
Ensure: Java $\cup$ locked names, locked dataflow

$t \leftarrow$ effective type of e
$v \leftarrow$ new local variable of type t and name n
$s \leftarrow$ enclosing statement of e
INSERT LOCAL VARIABLE(s, v)
EXTRACT ASSIGNMENT(v, e)
MERGE DECLARATION(v)

Algorithm 9 INSERT LOCAL VARIABLE($s : \mathit{Stmt}, v : \mathit{LocalVar}$)

Require: Java
Ensure: Java $\cup$ locked names

$b \leftarrow$ enclosing block of s
assert variable v can be introduced into block b
`lockNames`(b, n)
insert v before s

2.7.2 Extract Assignment

This refactoring extracts an expression into an assignment to a local variable. Implemented in `ExtractTemp/ExtractAssignment.jrag`.

Algorithm 10 Extract Assignment(v : *LocalVar*, e : *Expr*) : *Assignment*

Require: Java
Ensure: Java $\cup$ locked dependencies

```
assert e is extractible
a ← new assignment from e to v
if e is in expression statement then
    replace e with a
else
    s ← enclosing statement of e
    lock all names in e
    insert a before s
    replace e with locked access to v
end if
return a
```

2.7.3 Merge Variable Declaration

This refactoring merges a variable declaration with the assignment immediately following it, if that assignment is an assignment to the same variable. Implemented in `ExtractTemp/MergeVarDecl.jrag`.

Algorithm 11 Merge Variable Declaration(v : *LocalVar*)

Require: Java \ multi-declarations – TODO no checks for this in the implementation and no test for this, refactorings are not stand alone enough imo
Ensure: Java

```
if v has initialiser then
    return
end if
s ← statement following v
if s is assignment to v then
    make RHS of s the initialiser of v
    remove s
end if
```

2.8 Inline Constant

This refactoring inlines a constant field into all its uses. Implemented in `InlineTemp/InlineConstant.jrag`; see Algorithms **??**, **??**, **??**.

Algorithm 12 Inline Constant(f : *Field*)

Require: Java \ implicit assignment conversion
Ensure: Java

```
for all uses u of f do
    Inline Constant(u)
end for
Remove Field(f)
```

Algorithm 13 Inline Constant(u : *FieldAccess*)

Require: Java
Ensure: Java ∪ locked dependencies

```
f ← field accessed by u
assert f is final and static, and has an initialiser
e ← locked copy of the initialiser of f
assert if u is qualified, then its qualifier is a pure expression
replace u with e, discarding its qualifier if any
```

Algorithm 14 Remove Field(f : *Field*)

Require: Java
Ensure: Java

```
if f is not used and if it has an initialiser, it is pure then
    remove f
end if
```

2.9 Inline Method

See ECOOP 2009 publication. (TODO)

2.10 Inline Temp

This refactoring inlines a local variable into all its uses. Implemented in `InlineTemp/InlineTemp.jrag`; see Algorithms ??, ??, ??, ??

Algorithm 15 INLINE TEMP(d : *LocalVar*)

Require: Java
Ensure: Java

1: $a \leftarrow \lfloor$SPLIT DECLARATION$\rfloor(d)$
2: $\lfloor$INLINE ASSIGNMENT$\rfloor(a)$
3: $\lfloor$REMOVE DECL$\rfloor(v)$

Algorithm 16 SPLIT DECLARATION(d : *LocalVar*) : `option` *Assignment*

Require: Java $\setminus$ compound declarations
Ensure: Java $\cup$ locked names, first-class array init

```
if d has initialiser then
    x ← variable declared in d
    a ← new assignment from initialiser of d to x
    insert a as statement after d
    remove initialiser of d
    return Some a
else
    return None
end if
```

2.11 Introduce Factory

This refactoring introduces a static factory method as a replacement for a given constructor, and updates all uses of the constructor to use this method instead. Implemented in `IntroduceFactory/IntroduceFactory.jrag`; see Algorithm ??

We use `createFactoryMethod` (implemented in `util/ConstructorExt.jrag`) to create the factory method corresponding to constructor *cd* and insert it into the host type of *cd*. The factory method has the same signature as *cd*, but it has its own copies of all type variables of the host type used in *cd*.

Algorithm 17 INLINE ASSIGNMENT(a : *Assignment*)

Require: Java $\setminus$ implicit assignment conversion
Ensure: Java $\cup$ locked dependencies

```
x ← LHS of a
assert x refers to local variable
U ← all u such that a is a reaching definition of u
for all u ∈ U do
   assert a is the only reaching definition of u
   assert u is not an lvalue
   assert u, a are in same body declaration
   replace u with a locked copy of the RHS of a
end for
if U ≠ ∅ then
   remove a
end if
```

Algorithm 18 REMOVE DECL(d : *LocalVar*)

Require: Java $\setminus$ compound declarations
Ensure: Java

```
if d is not used and has no initialiser then
   remove d
end if
```

Algorithm 19 INTRODUCE FACTORY(cd : *ConstructorDecl*)

Require: Java
Ensure: Java $\cup$ locked names

```
f ← static factory method for cd
for all uses u of cd and its parameterised copies do
   if u is a class instance expression without anonymous class and it is not
   in f then
      replace u with a call to f
   end if
end for
```

2.12 Introduce Indirection

This refactoring creates a static method m' in type B that delegates to a method m in type A. Implemented in `IntroduceIndirection/IntroduceIndirection.jrag`; see Algorithm **??**.

Algorithm 20 INTRODUCE INDIRECTION(m : *Method*, B : *ClassOrInterface*)

Require: Java
Ensure: Java ∪ locked names, `return void`

assert B is non-library
$fn \leftarrow$ fresh method name
$m' \leftarrow$ copy of m with locked names and empty body
set name of m' to fn
$xs \leftarrow$ locked accesses to parameters of m'
set body of m' to `return` m(*xs*)`;`
INSERT METHOD(`hostType`(m), m')
MAKE METHOD STATIC(m')
MOVE STATIC METHOD(m', B)

2.13 Introduce Parameter

This refactoring turns an expression into a parameter of the surrounding method. Implemented in `ChangeMethodSignature/IntroduceParameter.jrag`; see Algorithm **??**.

Algorithm 21 INTRODUCE PARAMETER(e : *Expr*, n : *Name*)

Require: Java
Ensure: Java ∪ locked names

assert n is a valid name
assert e is extractible and constant
assert e appears within a method m
assert m is not overridden by and does not override any other methods
assert m has no parameter or local variable n
`lockMethodCalls`(`name`(m))
$t \leftarrow$ effective type of e
$p \leftarrow$ new parameter of type t and name n
insert p as the first parameter of m
replace e with locked access to p
for all calls c to m **do**
    insert a locked copy of e as first argument of c
end for

2.14 Introduce Parameter Object

This refactoring wraps a set P of parameters of a method m into a single parameter n of type w, where w is a newly created wrapper class containing fields corresponding to all the parameters in P. Implemented in `IntroduceParameterObject/IntroduceParameterObjec` see Algorithm **??**.

Algorithm 22 INTRODUCE PARAMETER OBJECT(m : *Method*, P : `set` *Parameter*, w : `set` Name, n : `set` Name)

Require: Java \ variable arity parameters
Ensure: Java ∪ locked names

assert m has a body
assert the parameters in P are in contiguous positions $i, \ldots, i+k$
$W \leftarrow$ new class containing fields for all the P and a standard constructor to initialise them
INSERT TYPE(`hostType`(m), W)
`lockMethodCalls`(`name`(m))
for all relatives r of m **do**
  assert r has no parameter or local variable with name n
  $[p_1; \ldots; p_n] \leftarrow$ parameters of r
  $p \leftarrow$ new parameter of type W and name n
  replace parameters $p_i, \ldots, p_{i+k}$ with p
  for all $j \in \{i, \ldots, i+k\}$ **do**
    $v_j \leftarrow$ new variable of same name, type, and finality as p_j
    insert assignment from $p.f_j$ to v_j at beginning of m
  end for
  for all calls c to r **do**
    $[a_1; \ldots; a_n] \leftarrow$ arguments of c
    replace arguments $a_i, \ldots, a_{i+k}$ with `new` W`(`a_i`,...,`a_{i+k}`)`
  end for
end for

Note that we need to perform the transformation for all relatives of m, *i.e.* for all methods r such that there exists a method m' with $m <:^* m'$ and $r <:^* m'$. We also lock all calls to methods of the same as m in the whole program; this ensures that if overloading resolution changes due to the transformation, the name binding framework will insert appropriate casts to rectify the situation.

Note: the implementation actually: eliminates variable arity parameter for this method and adjusts all calls; does not require p_i to be contiguous and adds new argument at the beginning. (This can be unsound for parameters with side effects!!!)

2.15 Move Inner To Toplevel

This refactoring converts a member type to a toplevel type. Implemented in `TypePromotion/MoveMemberTypeToToplevel.jrag`; see Algorithms **??**, **??**, **??**.

Algorithm 23 MOVE MEMBER TYPE TO TOPLEVEL(M : *MemberType*)

Require: Java
Ensure: Java $\cup$ locked names

if M is not static **then**
  $\lfloor$MAKE TYPE STATIC$\rfloor(M)$
end if
$p \leftarrow$ `hostPkg`(M)
lock all names in M
remove M from its host type
INSERT TYPE(p, M)

Algorithm 24 INSERT TYPE(p : *Package*, T : *ClassOrInterface*)

Require: Java
Ensure: Java $\cup$ locked names

assert no type or subpackage of same name as T in p
`lockNames`(`name`(T))
remove modifiers `static`, `private`, `protected` from T
insert T into p

2.16 Move Instance Method

This refactoring moves a method into a variable, which is either a parameter of that method or an accessible field. Implemented in `Move/MoveMethod.jrag`.

2.17 Move Members

In order to move Field, static methods, and member types, we simply lock all references to them, as well as all names contained in them, and (for fields) the flow dependencies of their initialiser, and then move them inside the AST.

2.18 Promote Temp to Field

This refactoring turns a local variable into a field. Implemented in `PromoteTempToField/PromoteTempToField.`

2.19 Pull Up

This refactoring pulls up a method m from its host class B to the super class A. Implemented in `PullUp/PullUpMethod.jrag`.

TODO: explain translation of type variables; this is basically a right-inverse of the type variable substitution that happens when inheriting a method

Note that INSERT METHOD ensures that the inserted method is not called from anywhere.

Algorithm 25 MAKE TYPE STATIC(M : *MemberType*)

Require: Java

Ensure: Java ∪ `with`, locked names

```
[A_n; ...; A_1] ← enclosing types of M
for all i ∈ {1, ..., n} do
   f ← new field of type A_i with name this$i
   INSERT FIELD(M, f)
   for all constructors c of M do
      p ← parameter of type A_i with name this$i
      assert no parameter or variable this$i in c
      insert p as first parameter of c
      if c is chaining then
         add this$i as first argument of chaining call
      else
         a ← new assignment of p to f
         insert a after super call
      end if
   end for
end for
for all constructors c of M do
   for all non-chaining invocations u of c do
      es ← enclosing instances of u
      assert |es| = n
      insert es as initial arguments to u
      discard qualifier of u, if any
   end for
end for
if M not in inner class then
   put modifier static on M
end if
for all non-static callables m of M do
   if m has a body then
      surround body of m by
      with(this$n, ..., this$1, this) {...}
   end if
end for
```

Algorithm 26 Move Method(m : *InstanceMethod*, v : *Variable*)

Require: Java

Ensure: Java ∪ locked names, `return void`, fresh variables, demand `final`

assert v is either a parameter of m or a field
$T \leftarrow$ type of v
assert T is a non-library class
assert m has a body and is not from library
$m' \leftarrow$ copy of m with `synchronized` removed and all names locked
$xs \leftarrow$ list of locked accesses to parameters of m
if v is a parameter **then**
    $i \leftarrow$ position of v in parameter list of m
    remove ith parameter from m'
    remove ith element of xs
else
    $i \leftarrow 0$
end if
$v' \leftarrow$ `final` local variable declaration with same name and type as v, initialised to `this`
insert v' as first statement into m'
lock all uses of v inside m' to v'
$qs \leftarrow []$
for all enclosing classes C of m **do**
    $p_C \leftarrow$ demand `final` parameter with fresh name, of type C
    make p_C the ith parameter of m'
    $e \leftarrow$ access to C`.this`
    insert e as ith element into xs
    $qs \leftarrow [\![p_C]\!] :: qs$
end for
wrap body of m' into `with(`qs`) {...}`
set body of m to `return` $[\![v]\!]$`.`$[\![m]\!]$`(`xs`);`
Insert Method(T, m')
eliminate `with` statement in m'
Inline Temp(v')
for all p_C **do**
    Remove Parameter(p_C) **or** Id()
end for

Algorithm 27 PROMOTE TEMP TO FIELD(d : *LocalVar*)

Require: Java
Ensure: Java $\cup$ locked dependencies

```
⌊SPLIT DECLARATION⌋(d)
d′ ← new private field of same type and name as d
make d′ static if d is in static context
⌊INSERT FIELD⌋(hostType(d), d′)
for all uses u of d do
    lock u onto d′
    lock reaching definitions of u
end for
REMOVE DECL(d)
```

Algorithm 28 INSERT FIELD(T : *ClassOrInterface*, d : *Field*)

Require: Java
Ensure: Java $\cup$ locked names

```
assert T has no local field with same name as d
assert d has no initialiser
assert if T is inner and d is static, then d is a constant
lockNames(name(d))
insert field d into T
```

Algorithm 29 PULL UP METHOD(m : *Method*)

Require: Java
Ensure: Java $\cup$ locked names

```
assert the host type of m B is a non-library class
assert the superclass A of B is also non-library
m′ ← copy of m with locked names
translate type variables in m′ from B to A
Insert Method(A, m′)
remove m from B
```

2.20 Push Down

This refactoring pushes a method down to all subclasses of its defining class. Implemented in `PushDown/PushDownMethod.jrag`.

Algorithm 30 TRIVIALLY OVERRIDE(B : *Type*, m : *VirtualMethod*) : `option` *MethodCall*

Require: Java \ implicit method modifiers
Ensure: Java + locked names, `return void`

assert m is not `final`
if m not a member method of B **then**
  return `None`
end if
$m' \leftarrow$ copy of m with locked names
if m is `abstract` **then**
  insert method m' into B
  return `None`
else
  $xs \leftarrow$ list of locked accesses to parameters of m'
  $c \leftarrow$ `super`.m(xs)
  set body of m' to `return` c;
  insert method m' into B
  return `Some` c
end if

Algorithm 31 REMOVE METHOD(m : *Method*)

Require: Java
Ensure: Java

assert $(\texttt{uses}(m) \cup \texttt{calls}(m)) \setminus \texttt{below}(m) = \emptyset$
$o \leftarrow \{m' \mid m <: m'\}$
if $o \neq \emptyset \wedge \forall m' \in o.m'$ is abstract **then**
  for all types B that inherit m **do**
    MAKE TYPE ABSTRACT(B)
  end for
end if
remove m

2.21 Rename

This family of refactorings is used for renaming named program entities. Implemented in `Renaming/`.

Refactoring RENAME FIELD changes the name of a field f to n. It ensures that n is indeed a valid name and that the host type of f contains no other

Algorithm 32 MAKE METHOD ABSTRACT(m : *Method*)

Require: Java
Ensure: Java

assert $\texttt{calls}(m) \setminus \texttt{below}(m) = \emptyset$
for all types B that inherit m **do**
  MAKE TYPE ABSTRACT(B)
end for
make m `abstract`

Algorithm 33 MAKE TYPE ABSTRACT(T : *Type*)

Require: Java
Ensure: Java

if T is interface **then**
  return
end if
assert T is class and never instantiated
make T `abstract`

Algorithm 34 PUSH DOWN VIRTUAL METHOD(m : *VirtualMethod*)

Require: Java
Ensure: Java $\cup$ locked names

for all types $B <: \texttt{hostType}(m)$ **do**
  $c \leftarrow \lfloor$TRIVIALLY OVERRIDE$\rfloor(B, m)$
  if $c \neq \texttt{None}$ **then**
    INLINE METHOD(c)
  end if
end for
REMOVE METHOD(m)
    or MAKE METHOD ABSTRACT(m)
    or ID()

Algorithm 35 RENAME FIELD(f : *Field*, n : *Name*)

Require: Java
Ensure: Java $\cup$ locked names

assert n is a valid name
assert host type of f contains no other field of name n
$\texttt{lockNames}(\{n, \texttt{name}(f)\})$
set name of f to n

field called n. It then globally locks all accesses to variables, types, or packages named either n or $\texttt{name}(f)$, and changes the name of f to n.

Algorithm 36 RENAME LOCAL(v : *Local*, n : *Name*)

Require: Java
Ensure: Java $\cup$ locked names

```
assert n is a valid name
assert scope of v does not intersect scope of any other Local named n
lockNames(block(v), {n, name(f)})
set name of v to n
```

Refactoring RENAME LOCAL changes the name of a local variable or parameter v to n. It ensures that n is indeed a valid name and that the renaming v to n will not violate the rule that scopes of local variables of the same name cannot be nested. It then again locks all accesses to variables, types, or packages named either n or $\texttt{name}(v)$, but only within the enclosing block of v, and changes the name of v to n.

Algorithm 37 RENAME METHOD(m : *Method*, n : *Name*

Require: Java
Ensure: Java $\cup$ locked names, locked overriding

```
assert n is a valid name
lockMethodNames({name(m), n})
lockOverriding({name(m), n})
for all m′ such that ∃m″.m <:* m″ ∧ m′ <:* m″ do
    assert m′ is not native
    s ← signature of m′ after renaming
    assert host type of m′ contains no local method of signature s
    assert m′ can override or hide any ancestor method of signature s
    assert m′ can be overridden or hidden by any descendant method of
        signature s
    set name of m′ to n
    remove any static import of m′ if it would become vacuous
end for
```

Refactoring RENAME METHOD changes the name of a method m to n. It ensures that n is a valid name, then locks all calls to methods of name $\texttt{name}(m)$ or n, and their overriding dependencies. Now it changes the names of all methods m' related to m (*i.e.*, such that m and m' both transitively override the same method), checking that the resulting program will be well-formed: in particular, there cannot be another local method with the same signature, and any methods that the renamed m' would override or hide must, in fact, be overridable or hidable by m', and vice versa for methods that would override or hide m'. If there is a static import that only imports m' (and not also another static member of the surrounding class), then remove that import. We could, of course, try to adjust it, but changing imports is a tricky business.

Algorithm 38 Rename Type(T : *Type*, n : *Name*

Require: Java
Ensure: Java $\cup$ locked names

assert n is a valid name
assert no native method is nested in T
assert there is no nesting or enclosing type of name n
assert if T is a toplevel type, there is no other toplevel type n in the enclosing package, and it has no subpackage of name n
assert if T is a type parameter, there is no type parameter of name n in the parameter list where it occurs
`lockNames`($\{$`name`$(T), n\}$)
set name of T to n
set names of constructors of T to n
if T is public, change the name of its compilation unit to match
remove any single type import declaration of T that would clash with a visible type or with another import declaration
remove any static import of T if it would become vacuous

Refactoring Rename Type changes the name of a type T to n. It is fairly straightforward, except for the well-formedness checks and the treatment of import declarations.

2.22 Self-Encapsulate Field

This refactoring makes a field private, rerouting all accesses to it through getter and setter methods. Implemented in `SelfEncapsulateField/SelfEncapsulateField.jrag`.

By “abbreviated assignment” we mean `x+=y` and friends, as well as increment and decrement expressions. The language restriction tries to expand these into normal assignments, but may fail if the data flow is too complicated. If it succeeds, every lvalue will appear on the left hand side of a (simple) assignment.

Note that even when f is final there may still be assignments to f from within constructors; we cannot encapsulate these assignments, so we skip them.

3 Node Types

See Fig. **??**. We also use the non-node type *Name* to represent names.

4 Utility Functions

See Fig. **??**.

Algorithm 39 SELF-ENCAPSULATE FIELD(f : *Field*)

Require: Java $\setminus$ abbreviated assignments
Ensure: Java $\cup$ locked names

create getter method g for f
if f is not `final`, create setter method s for it
for all all uses u of f and its substituted copies **do**
  if $u \notin \texttt{below}(g) \cup \texttt{below}(s)$ **then**
    if u is an rvalue **then**
      replace u with locked access to g
    else
      if f is not `final` **then**
        $q \leftarrow$ qualifier of u, if any
        $r \leftarrow$ RHS of assignment for which u is LHS
        replace u with locked access to s on argument r, qualified with q if applicable
      end if
    end if
  end if
end for

Node Type	Description
ClassOrInterface	either a class or an interface; is a *Type*
Field	field declaration
LocalVar	local variable declaration
MemberType	type declared inside another type; is a *Type*
Method	method declaration
MethodCall	method call
Package	package
Type	type declaration
VirtualMethod	non-`private` instance method; is a *Method*

Figure 1: Node Types

Name	Description
`below`(n)	returns the set of all nodes below n in the syntax tree
`calls`(m)	returns all calls that may dynamically resolve to method m; can be a conservative over-approximation
`hostPkg`(e)	returns the package of the compilation unit containing e
`hostType`(e)	returns the closest enclosing type declaration around e
`lockMethodCalls`(n)	locks all calls to methods named n anywhere in the program
`lockNames`(n)	locks all names anywhere in the program that refer to a declaration with name n
`name`(e)	returns the name of program entity e
`uses`(m)	returns all calls that statically bind to method m

Figure 2: Utility Functions

List of Algorithms