# ${\bf Homework~2}$ CSC 445-01: Theory of Computation

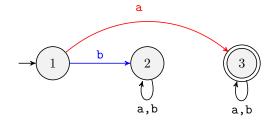
Matthew Mabrey, Luke Kurlandski February 18, 2021

#### 1.4

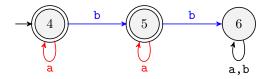
We use blue for 'b' and red for 'a' when space is limited

#### $\mathbf{e}$

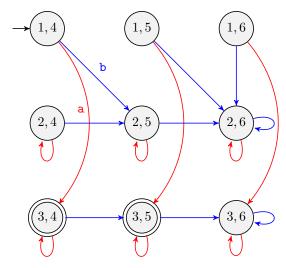
 $\boldsymbol{w}$  starts with an 'a'



w has at most one 'b'

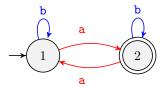


 $\boldsymbol{w}$  starts with an 'a' and has at most one 'b'

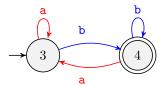


## $\mathbf{f}$

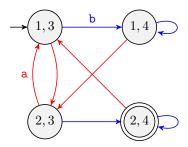
 $\boldsymbol{w}$  has an odd number of 'a'



w ends in a 'b'

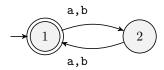


 $\boldsymbol{w}$  has an odd number of 'a' and ends in a 'b'

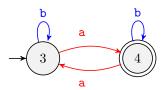


## $\mathbf{g}$

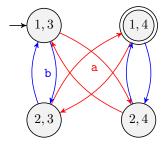
 $\boldsymbol{w}$  has an even length



w has an odd number of 'a'



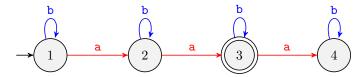
 $\boldsymbol{w}$  has an even length and has an odd number of 'a'



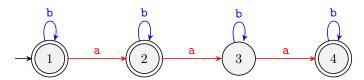
#### 1.5

#### g

w is any string that contains exactly two 'a'



w is any string that does not contain exactly two 'a'



#### 1.7

b

 $\mathbf{c}$ 

 $\mathbf{e}$ 

#### 1.31

#### 1.33

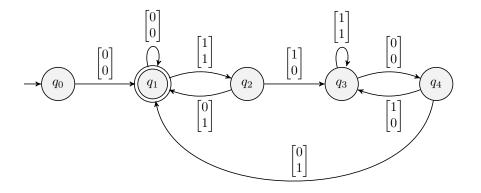
From problem 1.31, we know we are allowed to "wind the tape backwards" i.e., we will demonstrate that  $C^R$  is a regular language, thus prove that C itself is regular. To prove that  $C^R$  is regular, we will construct a nondeterministic finite automata that recognizes  $C^R$ .

For each symbol, we will read its top bit and judge whether or not the bottom bit is correct such that  $3\text{row}_1 = \text{row}_2$ . To do so, we track two conditions: 1) whether or not  $3\text{row}_1$  results in a carry-out in base two and 2) what the previous top symbol is. With these two pieces of information, given a top bit, we can check to see if the bottom bit is correct.

These two pieces of information result in an automata with 4 states to keep track of, but we had an additional start state.

- $q_0$ : the start state, to prevent the automata from accepting the empty string
- $q_1$ : the state where there is no carry-in from previous step and the previous top symbol is a 0
- $q_2$ : the state where there is a carry-in and the previous top value is a 1
- $q_3$ : the state where there is no carry-in and the previous top value is a 1
- $q_4$ : the state where there is a carry-in from previous step and the previous top value is a 0

Our transition function merely Recall that we read the string backwards, so the binary number is read from the  $2^0$ th place first.



# 1.34