

File for the algorithms

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1 Introduction

Algorithm 1 DFS

```
1: visited  $\leftarrow$  emptySet
2: procedure DFS(state)
3:   if state = goal then
4:     return state
5:   for neighbour in neighbours(state) do
6:     if neighbour not in visited then
7:       DFS(neighbour)
8:   visited.add(state)
```

Algorithm 2 A*

```
1:  $g(state)$  ▷ The cost to reach the current state
2:  $h(state)$  ▷ Estimated cost of the cheapest path from state to goal
3:  $f(state) \leftarrow g(state) + h(state)$ 
4:  $neighbours(state)$  ▷ Expands possible moves from current state ordered by  $g + h$ 
5: procedure AStar( $startState$ )
6:    $closedSet \leftarrow \text{empty set}$ 
7:    $openSet \leftarrow \text{empty set}$ 
8:    $openSet.add(startState)$ 
9:   while  $openSet$  is not empty do
10:     $currentState \leftarrow null$ 
11:    for  $state$  in  $openSet$  do
12:      if  $f(state) < f(currentState)$  then
13:         $currentState \leftarrow state$ 
14:       $openSet.remove(currentState)$ 
15:       $closedSet.add(currentState)$ 
16:      if  $currentState = goal$  then
17:        return  $currentState$ 
18:      for  $neighbour$  in  $neighbours(currentState)$  do
19:        if  $neighbour$  not in  $closedSet$  then
20:          if  $neighbour$  in  $openSet$  then
21:            if  $g(currentState) < g(neighbour)$  then
22:               $g(neighbour) \leftarrow g(currentState)$ 
23:          else
24:             $openSet.add(neighbour)$ 
```

Algorithm 3 Manhattan Distance

```
1: procedure MANDIST( $state$ ) ▷ The current puzzle configuration
2:    $total \leftarrow 0$ 
3:    $puzzleLength \leftarrow state.size()$ 
4:    $dimensions \leftarrow \sqrt{puzzleLength}$ 
5:   for  $i \leftarrow 1, puzzleLength$  do ▷ Loops through each tile of the puzzle
6:      $tileValue \leftarrow state[i]$ 
7:      $expectedRow \leftarrow (tileValue - 1) \div dimensions$ 
8:      $expectedCol \leftarrow (tileValue - 1) \bmod dimensions$ 
9:      $rowNum \leftarrow i \div dimensions$ 
10:     $rowNum \leftarrow i \bmod dimensions$ 
11:     $total \leftarrow total + |expectedRow - rowNum| + |expectedCol - colNum|$ 
12:   return  $total$  ▷ The heuristic is the total
```

Algorithm 4 Iterative Deepening A Star

```
1: state                                ▷ The current puzzle configuration
2:  $g(state)$                             ▷ The cost to reach the current state
3:  $h(state)$                             ▷ Estimated cost of the cheapest path from state to goal
4:  $f(state) \leftarrow g(state) + h(state)$ 
5:  $neighbours(state)$   ▷ Expands possible moves from current state ordered by  $g + h$ 

6: procedure IDASTAR(state)
7:    $bound \leftarrow f(state)$ 
8:   while not solved do                ▷ Loops until a solution is found
9:      $bound \leftarrow DFS(state, bound)$   ▷ Performs a bounded depth-first search

10: procedure DFS(state, bound)
11:   if  $f(state) > bound$  then
12:     return  $f(state)$ 
13:   if  $h(state) = 0$  then                ▷ No more moves needed to reach goal state
14:     return solved
15:    $min \leftarrow \infty$ 
16:   for neighbour in  $neighbours(state)$  do
17:      $temp \leftarrow DFS(neighbour, bound)$ 
18:     if  $temp < min$  then
19:        $min \leftarrow temp$ 
20:   return  $min$                         ▷ Returns the smallest of the neighbours
```

Algorithm 5 Breadth-First Search

```
1: procedure BFS(state)
2:    $s \leftarrow \text{empty set}$ 
3:    $q \leftarrow \text{empty queue}$ 
4:    $q.add(state)$ 
5:    $q.enqueue(state)$ 
6:   while  $q.size() > 0$  do
7:      $currentState \leftarrow q.dequeue()$ 
8:     if  $currentState = goal$  then
9:       return current
10:    for neighbour in  $neighbours(currentState)$  do
11:      if neighbour is not in  $s$  then
12:         $s.add(neighbour)$ 
13:         $q.enqueue(neighbour)$ 
```

Algorithm 6 Is Current State Solvable

```
1: procedure ISSOLVABLE(state)
2:    $puzzleLength \leftarrow state.size()$ 
3:    $gridWidth \leftarrow \sqrt{puzzleLength}$ 
4:    $blankRowEven \leftarrow true$ 
5:   for  $i \leftarrow 1, puzzleLength$  do
6:     if  $state[i] = 0$  then
7:        $blankRowEven \leftarrow (i/gridWidth) \bmod 2 = 0$ 
8:       continue
9:       for  $j \leftarrow i + 1, puzzleLength$  do
10:        if  $state[i] > state[j]$  and  $state[j] \neq 0$  then
11:           $parity \leftarrow !parity$ 
12:   if  $gridWidth \bmod 2 = 0$  and  $blankRowEven$  then
13:     return  $!parity$ 
14:   return  $parity$ 
```

Algorithm 7 Iterative Deepening Depth-First Search

```
1:  $goal$  ▷ the completed goal state of the puzzle

2: procedure IDDFS(state)
3:   for  $depth \leftarrow 0, \infty$  do
4:      $found \leftarrow DLS(state, depth)$ 
5:     if  $found \neq null$  then
6:       return  $found$ 
7: procedure DLS(state, depth)
8:   if  $depth = 0$  and  $state = goal$  then
9:     return  $found$ 
10:  if  $depth > 0$  then
11:    for  $neighbour$  in  $neighbours(state)$  do
12:       $found \leftarrow DLS(neighbour, depth - 1)$ 
13:      if  $found \neq null$  then
14:        return  $found$ 
15:  return  $null$ 
```

Progress report

Description of project: aims, motivation	First	2.1	2.2	3	Fail
Description and understanding of issues and problems addressed in the project	First	2.1	2.2	3	Fail
Achievement so far according to what is reasonably expected for the type of project	First	2.1	2.2	3	Fail
Discussion and justification of changes to project aims, scope, workplan	First	2.1	2.2	3	Fail

Quality of writing

Clarity, structure correctness of writing	First	2.1	2.2	3	Fail
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Comments

Markers should circle the appropriate level of performance in each section. Report and evaluation sheet should be collected by the student from the supervisor.