SMA* Pseudo

Wikipedia

```
function SMA-star(problem): path
1
      queue: set of nodes, ordered by f-cost;
 2
 3
    begin
      queue.insert(problem.root-node);
4
 5
      while True do begin
6
        if queue.empty() then return failure;
 7
        //there is no solution that fits in the given memory
8
        node := queue.begin(); // min-f-cost-node
9
        if problem.is-goal(node) then return success;
10
11
12
        s := next-successor(node)
13
        if !problem.is-goal(s) && depth(s) == max_depth then
14
             f(s) := inf;
            // there is no memory left to go past s,
15
            // so the entire path is useless
16
17
        else
             f(s) := max(f(node), g(s) + h(s));
18
             // f-value of the successor is the maximum of
19
            // f-value of the parent and
20
            // heuristic of the successor + path length to the successor
21
        endif
22
        if no more successors then
23
           update f-cost of node and those of its ancestors if needed
24
25
        if node.successors ⊆ queue then queue.remove(node);
26
        // all children have already been added to the queue
27
28
        // via a shorter way
        if memory is full then begin
29
30
          badNode := shallowest node with highest f-cost;
          for parent in badNode.parents do begin
31
             parent.successors.remove(badNode);
32
             if needed then queue.insert(parent);
33
          endfor
34
        endif
35
36
37
        queue.insert(s);
      endwhile
38
39
    end
```

https://cis.temple.edu/

```
function SMA*(problem) returns a solution sequence
1
       Queue, a queue of nodes ordered by f-cost
2
           {Queue is a static local variable}
3
    { Queue <-- MAKE-QUEUE(MAKE-NODE(INITIAL-STATE[problem]))
4
5
          if Queue is empty then return failure
6
          n <-- deepest least f-cost node in Queue
7
          if GOAL-TEST(n) then return success
8
          s <-- NEXT-SUCCESSOR(n)
9
          if s is not a goal and is at maximum depth then
10
             f(s) <-- INFINITY
11
12
          else
13
             f(s) \leftarrow MIN(f(n),g(s)+h(s))
          if all of n's successors have been generated then
14
             update n's f-cost and those of its ancestors if necessary
15
          if SUCCESSORS(n) all in memory then remove n from Queue
16
          if memory is full then
17
            {delete shallowest, highest f-cost node in Queue
18
             remove it from its parent's successor list
19
             insert its parent on Queue if necessary}
20
          insert s in Queue}}
21
```

Efficient memory-bounded search methods S.Russel

```
Algorithm SMA*(start):
put start on OPEN; USED \leftarrow 1;
loop
  if empty(OPEN) return with failure;
  best \leftarrow deepest least-f-cost leaf in OPEN;
  if goal(best) then return with success;
  succ \leftarrow next\text{-}successor(best);
  f(succ) \leftarrow max(f(best), g(succ) + h(succ));
  if completed(best), BACKUP(best);
  if S(best) all in memory, remove best from OPEN.
  USED \leftarrow USED+1;
  if USED > MAX then
     delete shallowest, highest-f-cost node in OPEN;
     remove it from its parent's successor list;
     insert its parent on OPEN if necessary;
     USED \leftarrow USED-1;
  insert succ on OPEN.
Procedure BACKUP(n):
```

```
if n is completed and has a parent then
   f(n) \leftarrow \text{least } f\text{-cost of all successors};
   if f(n) changed, BACKUP(parent(n)).
```

Norgiv 1st edition

```
function SMA*(problem) returns a solution sequence
  inputs: problem, a problem
  local variables: Queue, a queue of nodes ordered by f-cost
  Queue \leftarrow MAKE-QUEUE(\{MAKE-NODE(INITIAL-STATE[problem])\})
  loop do
      if Queue is empty then return failure
      n \leftarrow deepest least-f-cost node in Queue
      if GOAL-TEST(n) then return success
      s \leftarrow \text{Next-Successor}(n)
      if s is not a goal and is at maximum depth then
          f(s) \leftarrow \infty
      else
          f(s) \leftarrow MAX(f(n), g(s)+h(s))
      if all of n's successors have been generated then
          update n's f-cost and those of its ancestors if necessary
      if SUCCESSORS(n) all in memory then remove n from Queue
      if memory is full then
          delete shallowest, highest-f-cost node in Queue
          remove it from its parent's successor list
          insert its parent on Queue if necessary
      insert s on Queue
```

end