## CS 228: Probabilistic Graphical Models

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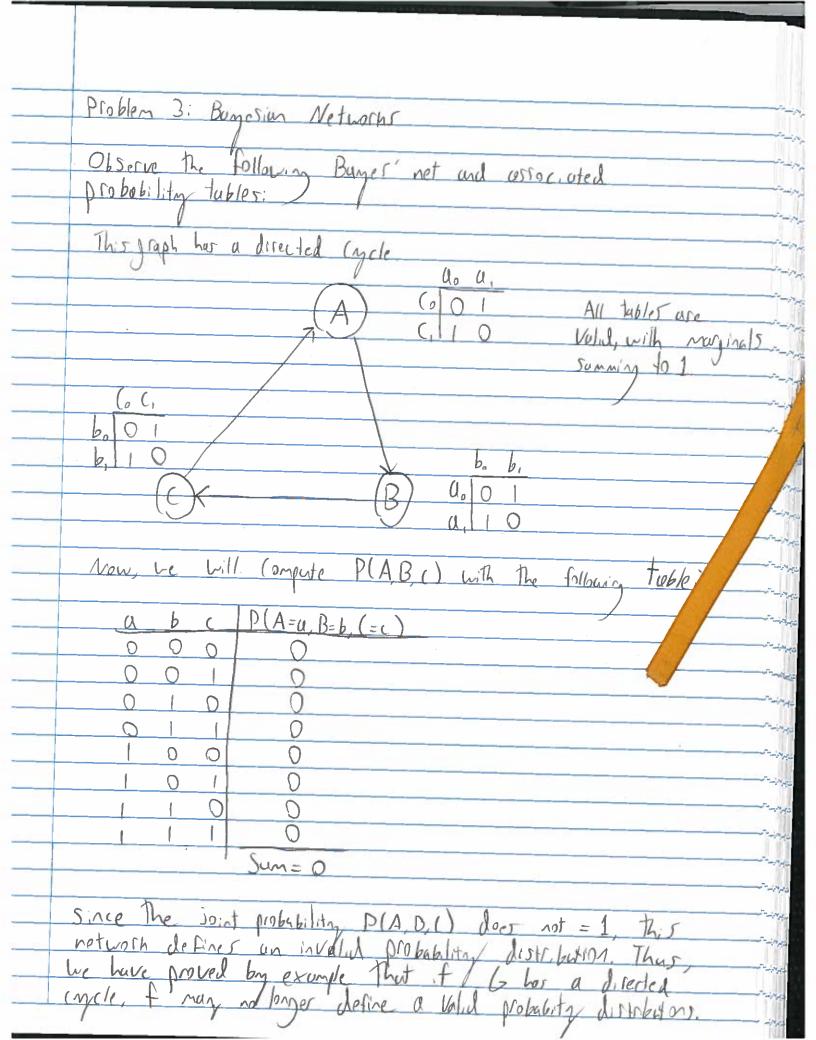
Due: 01/24/17

Submitted: 01/26/17

(using 2 late days)

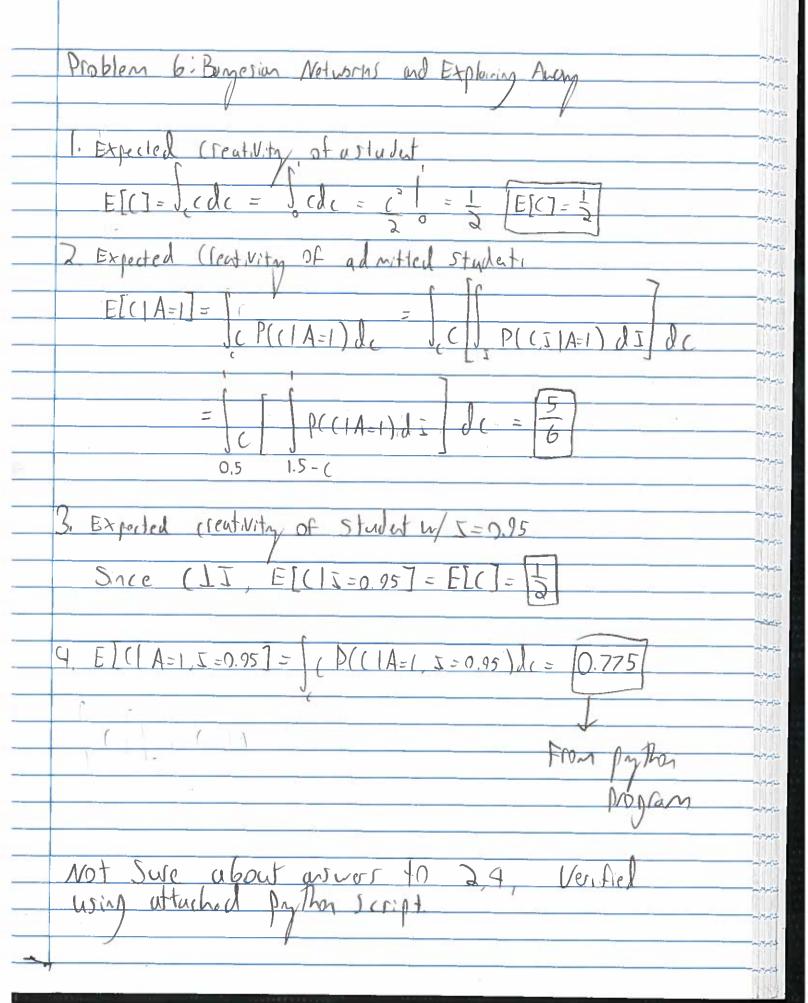
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	Problem 1: Probability Theory	
	Theory	
	I D D D	
	Let D= Have Disease	
_	Let T = Test Positive.	
	Civen:	
	P(T D) = 0.99	
_	$P(\overline{\tau} \overline{\rho}) = 0.19$	
-	P(D) = 0.0001	180
	Questias:	
	Question: What is P(DIT)? Since The disease of the D	
	what is P(VII)! Since The disease is late the	
-	Chance V Oct // his t	
_	Using Buyer Rule: Juite low.	
1	P(DIT) = P(TID) P(D) , MISSIM P(T)	
	P(DII) = P(TID) P(D), MISSIN P(T)	
+	We also have P(TID) = 1 - P(TID)	-
4	and $P(T \overline{p}) = 1 - P(\overline{+} \overline{p})$	
	Using Buyes' Rule:	
1	osing corps rate.	
+	NATE OF THE PARTY	
+	PIDIT) = PITID) PID), PIDIT) = PITID) PID	
	P(T) $P(T)$	
	P(DIT) = 1-P(DIT) -> 1-P(TID)P(D) = P(TID)P(D)	1
	TO P(V) = P(TID) P(D)	
1	P(T) $P(T)$	91411
-		
	-> P(T) = P(TID)P(D) + P(TID)P(D) = 0.99.0.0001+0.01.0999) = 0.001	J-1413
1		-44
	P(DIT) = P(TID) P(D) = 0.19.0.0001 = 0.0098	فتهابس
	$P(T) = 0.99 \cdot 0.0001 = 0.0098$	
	0,0101	100
	P(DIT) = 0.00 98	-444
		11.1151.05

	Problem 2: Review of Programming
	· Give on O(min) algorithm for solving: max P(x, x,)
	Nothing depends on Xn besides the last term Plx, IX,
_	So the only in entries we care about in the last tuble we
- 12	the methics where max Plxnlx.) for all Xn. ES. This
	tunes O(m2) to find and store there extres
	Now that we have calculated the The may of the last term
	for all Volues of Xm & 5, let's look at term N-1: P(xx 1xxx).
	Since we already know which is entries of the last table we
	will use we just need to use these Values to culculate
	The maximum product of P(X-1/x )P(x/x) for all possible
	Vulner of Xn-1 ES
	Pseudocode: Runtime analysis next to each line
	Mux ur = m Zeror
	For each V in S: O(m)
	V OFF = M Zero S
- 0	For each $V_i$ in $S$ : $V_{arr}[s] = P(x_i = V_i   x_i - V_i) \cdot \max_{i \in I} P(x_i   x_i = V_i) O(i)$
	$V_{\text{acc}[3]} = P(x = V,   y - V) \cdot m_{y} P(x   y - 1E) O(1)$
	max arr[:] = max (V_arr) D(m) to
	Butine of loop: O(m. (m.1+m)) = O(m2)
	The time of time of the time of time of the time of the time of ti
	Now iterate bachwoods Through all remaining tables connection This
	March and Stories the court for use by the aut court tita
	Now iterate backwards through all renoining tables repeating this process, and storing the results for use by the next computation. Those will be not a iterations before treaching X, for Olmin) time
	Finally ofter (culin & do a next block flow) all D.
	Finally after leaching & do a pass back. Through all the stock tables, following the soute for each 1x, 45.  This pass takes O(mx) time.
	The part to be a country to the
	IND PUD PUMPER DESTRICTIONS
	Find 5: 4: 0(13) 10(13)
	First for time = O(m2n) + O(m) = O(m2n)



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lem 4: Conditional Independence			Ç.
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	_[-	<del></del>	
(I. P(B, y), P(a), P(Bla), P(yla)			
		ĺ	Ì
P(a B,y) = P(B,y a)P(a) P(a)		—~~	
11013, M = 10, M (1) V			
P(By) P(By)			Į
P(B, K la) X			1
			Ì
Wo we need to know Blayla to compute P(B, y/a)		~~	-
Wo he need to know Bly la to compute P(B, 2/a) wing P(B/a)	-1-	<u> </u>	Š
1 Win MBlu), //(3/4)			
			ļ
(2, P(B,n), P(a), P(B,n/a)			Ì
Distriction of the second			
P(u B,y) = P(B,yla) P(u) / P(B,yla) /			
$P(D, y) \rightarrow P(u) \vee$			
$P(B_{\alpha}) \sqrt{ Y_{E}S }$			
			i
12 D(+10) D(014) D(0)		~~	÷
X3, P(bla), P(gla), P(a)		— <u>`</u> ~-	4
			122
MS This has strictly less information than 1, so			
it is also insufficient			-
The state of the s		~~	i
	- [-	-,~,	ث
1			Ŀ
VI. If Blyla, Then P(B, yla) = P(Bla). P(yla).			ļ
Now we have exough into to compute P(U/B, m) YES			Ť
J. 1110 J. GMWT 17 18/19/ 7E3			-
V2. TYES We have more information, so Still yes.			ļ
X3. We can mu compute P(Bylu) = P(Blu), P(yla)			Ī
Dr. Sill ( to the DID )			i
But we Still (anot compute PIB, m) sol there	- 1		+
is still insufficient information. / Wo?			+
			Ī
		10000	52

Problem 5: Bungsian Networks a. First, ALB, from proph structure, d-Sep (A:B) = True Proof: - A > D = B is blocked by V-structure -A+(+F+H+E+B:5 bb ched by F+H+E P(A=0,B=0)=0.8.0.3=0.24 D. Since d-Sep (AiE) = The ALE Proof: - E - B -> D - A blocked by B -> D - A V- Struiture - E + B - D - F + C + A / blocked by D - F + C - B - H ( F. blocked by V structure P(E=1 | A=1) = P(E=1 | B) = P(E=1 | B=0). P(B=0) + P(E=1 | B=1). P(B=1)  $= 0.9 \cdot 0.3 + 0.1 \cdot 0.7 = 0.34$ 7. Q. U-Sep. (A: E/EDH3) = False - A - C -> F -XH CE is an active Puly A -> ( -> F = open ( ) F -) ( ) = spen F -> PIE = Open V-structure W middle pres b. d-sop (LiElD) = False - GEFECEA - DE 15 an active polh befel = open F+(+A=Open ( A - D= Open A +D+B= Open, V-Structure w/ middle given D+B+E=Open (, O- Sep, (\$A;B3 36;H3/F) = Fulse - all pairs FA:63, FA:43, FB:63, FB:43 must be d-sep -B->E->H is an active path just takes one active path for one pair to make statement false

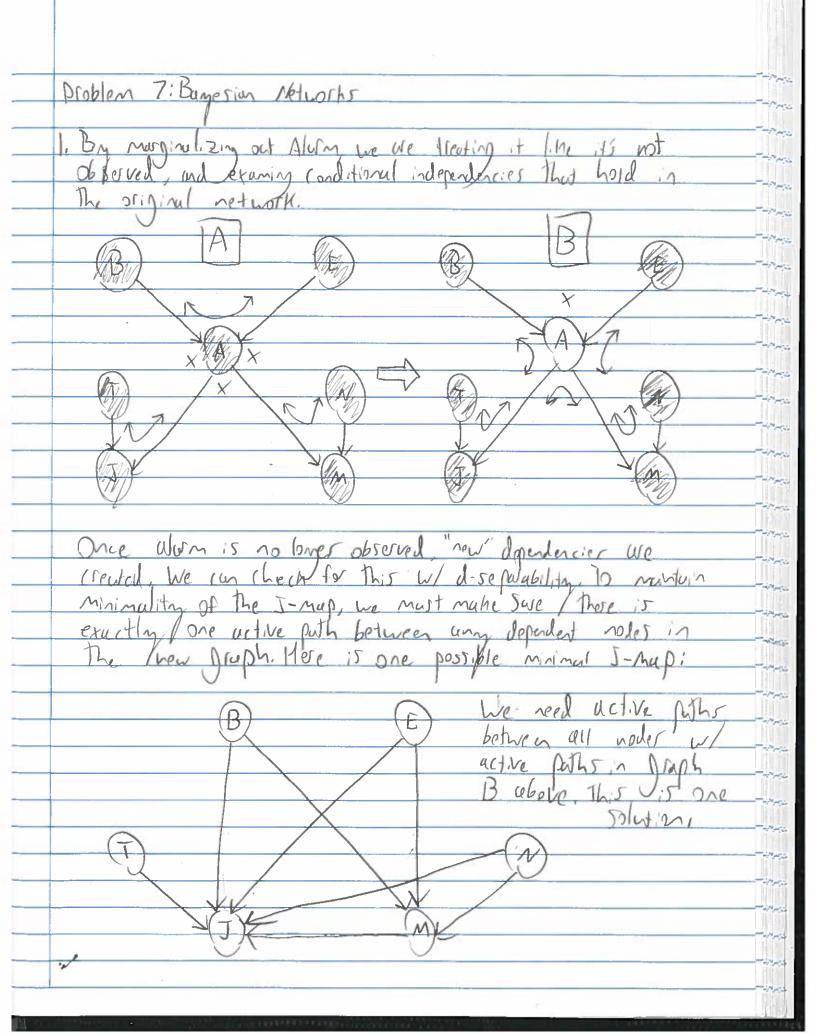


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Python script to check for problem 6
#!/usr/bin/env python
import numpy as np
n = 100000000
c = np.random.rand(n)
i = np.random.rand(n)
# Calculate E(C)
print "1. E(C) = \{\}".format(c.mean())
# Calculate P(A=1)
s = c + i
print "P(A=1) = {}".format(len((s>1.5).nonzero()[0])/float(n))
# Calculate E(C|A=1)
aidx = s > 1.5
print "2. E(C|A=1) = \{\}".format(c[aidx].mean())
print "P(C|A=1) = {}".format(len(c[aidx])/float(n))
# Calculate P(C|I=0.95)
lb,ub = 0.95-0.00005, 0.95+0.00005
lidx = (i>lb).nonzero()
uidx = (i<ub).nonzero()</pre>
iidx = lidx[0][np.in1d(lidx, uidx)]
```

print "3.  $E(C|I=0.95) = {}$ ".format(c[iidx].mean())

print "4.  $E(C|A=1,I=0.95) = {}$ ".format(c[idx].mean())

# Calculate P(C|A=1, I=0.95) aidx = aidx.nonzero()[0] idx = iidx[np.in1d(iidx, aidx)]



Problem 7. Bungerian Networks 2. To generalize the procedure: of the plant as unobserved as observed d- sepuability The previous The previous for two nodes that is not present in Inode-remove an edge to create adding ed

Droblem &: Towards inferce in Buyelon Networks 1. This problem can be solved with a recursive algorithm: function Compute (x): If X has no purents: return P(X) For all Parets of x, store (ompute (Pa(x)) return P(x / Pa(x)) 2. For any joint distribution P(+, 1) you can sample by This follows from the chain rule: P(x, y)= P(x) P(x) Alnorithmi I First draw a sample x ~ P(x) for all not voter Of the network much These nodes as covered . Then for all children of covered moder draw a somple from any roder that depend solely on previously Covered funder using the random Sample multingmind function juwanteed by the problem lared on the assertion above cover that a sampled order. · Refeat this process till all order in the Arath been covered, and you will have a sample for each parameter.

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	Problem 9: Programming Assignment	~
		~~:
	1. There we 2 = 2 binory images.	~~~
		~~~
	2. You would need 2 -1 = 289-1 palameters to Steering an arbitrary distribution over all 28x28 binory images	~~~
	in action distribution over all alx 28 birum imandi	~~
-		~~
	3. The joint poly for the Banks' Not in figure 1 can	~~
	3. The joint paf for the Bayer Net in figure 1 can be written asi	~~+
		~~~
	P(Z, Z, X, X, ) = P(Z,)P(Z,).P(X, 1Z, Z,)P(X, 1Z, Z)	7 7
	The conditional terms have 3 parameters each, 50	7 F 2
	# Pulanders = 784.3+2 = 2354	~~~ ~~~
	5. The intuitive role of the Z, Z parameters in this model is as prior beliefs for the model.	1 1 1 1
		1
		74
		1000
		7.2
	7. The relation that between the images is that They	
1	7. The relationship between The manes is that Then correspond spatially, except fliffed over the line yex	7
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