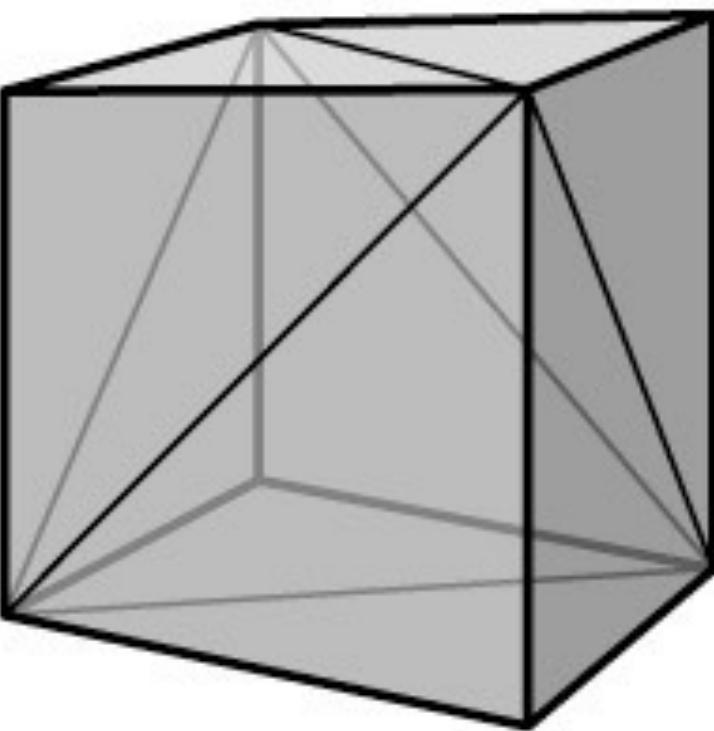


Lecture 7:

Digital Geometry Processing

**Interactive Computer Graphics
Stanford CS248, Winter 2019**

A small triangle mesh



8 vertices, 12 triangles

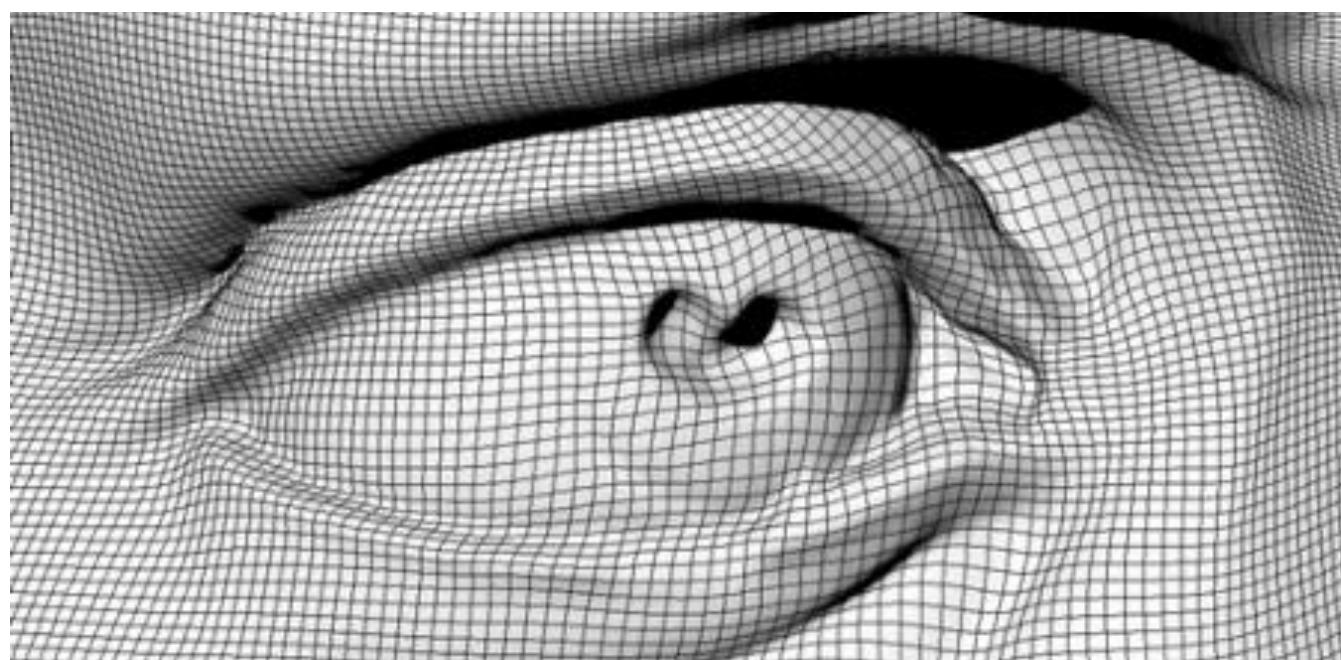
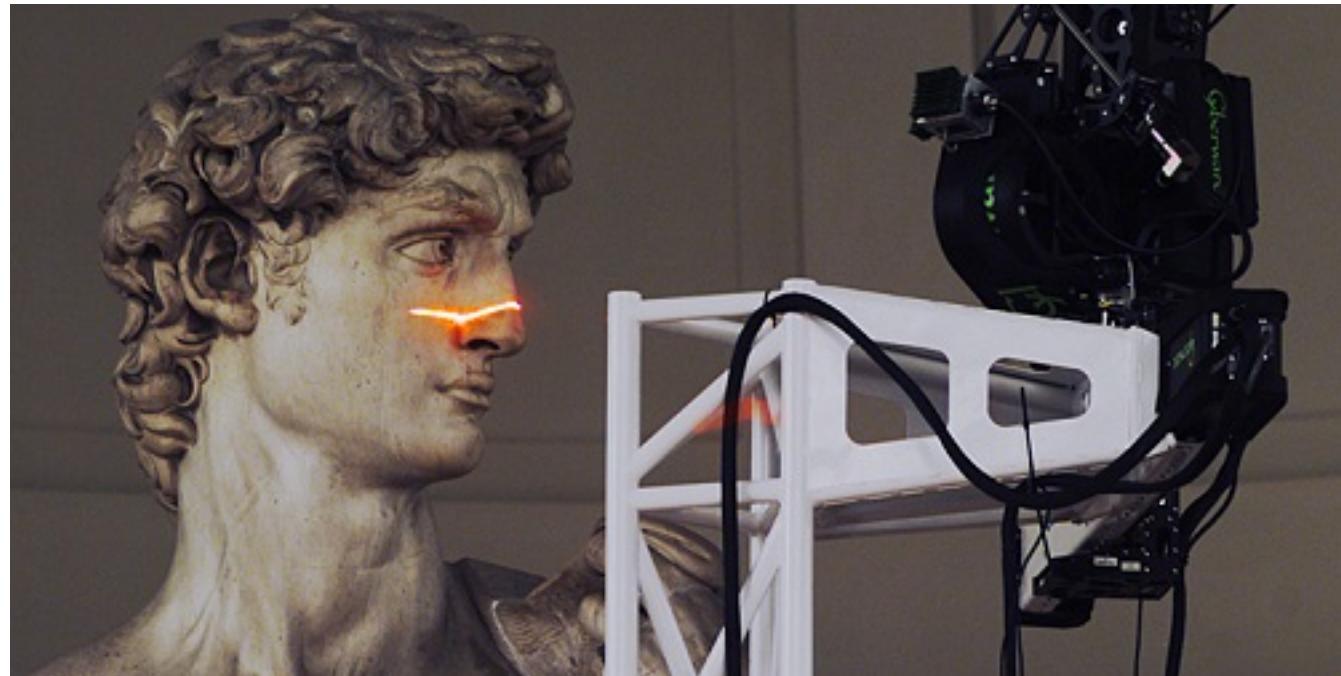
A large triangle mesh

David

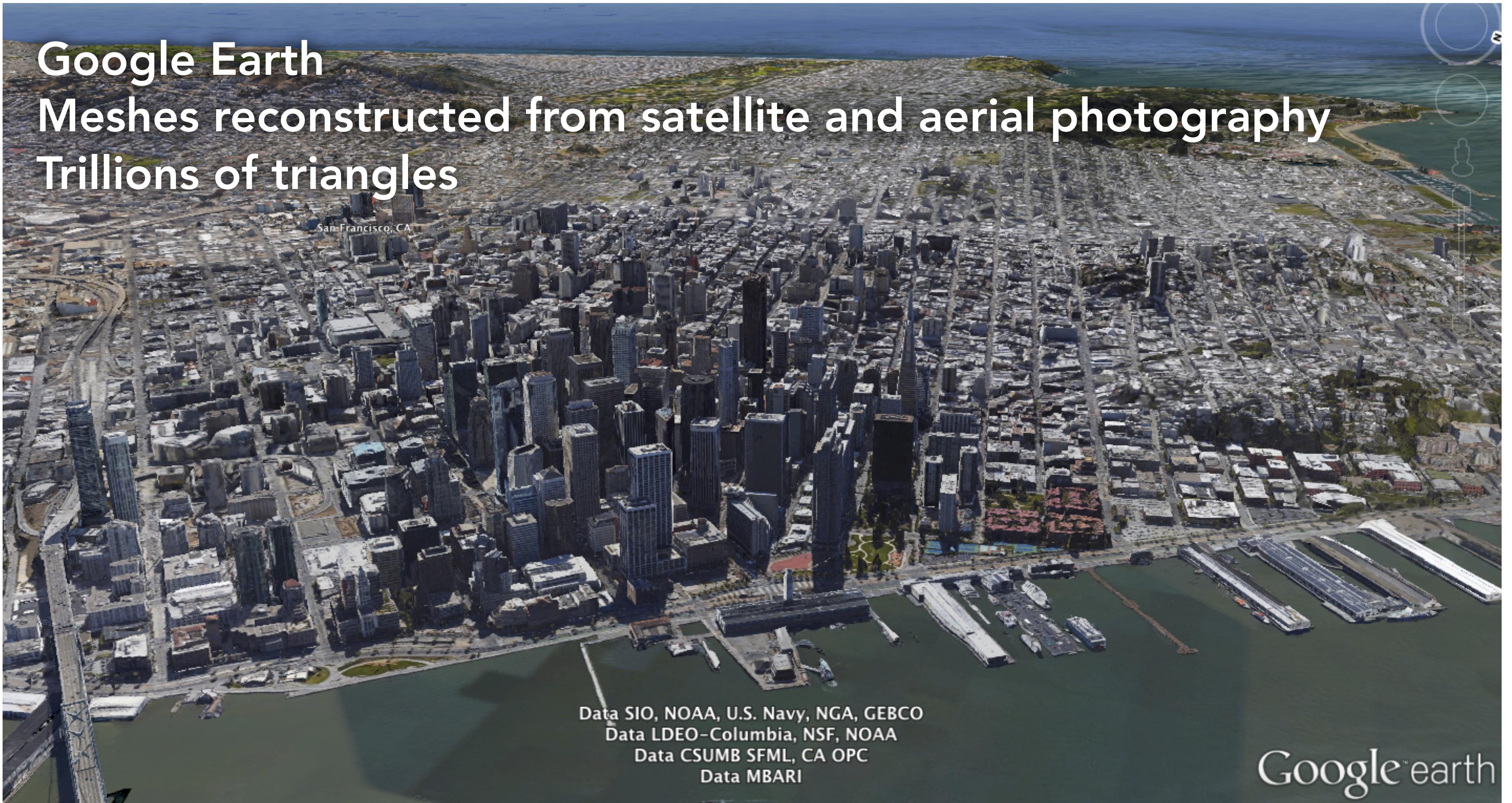
Digital Michelangelo Project

28,184,526 vertices

56,230,343 triangles

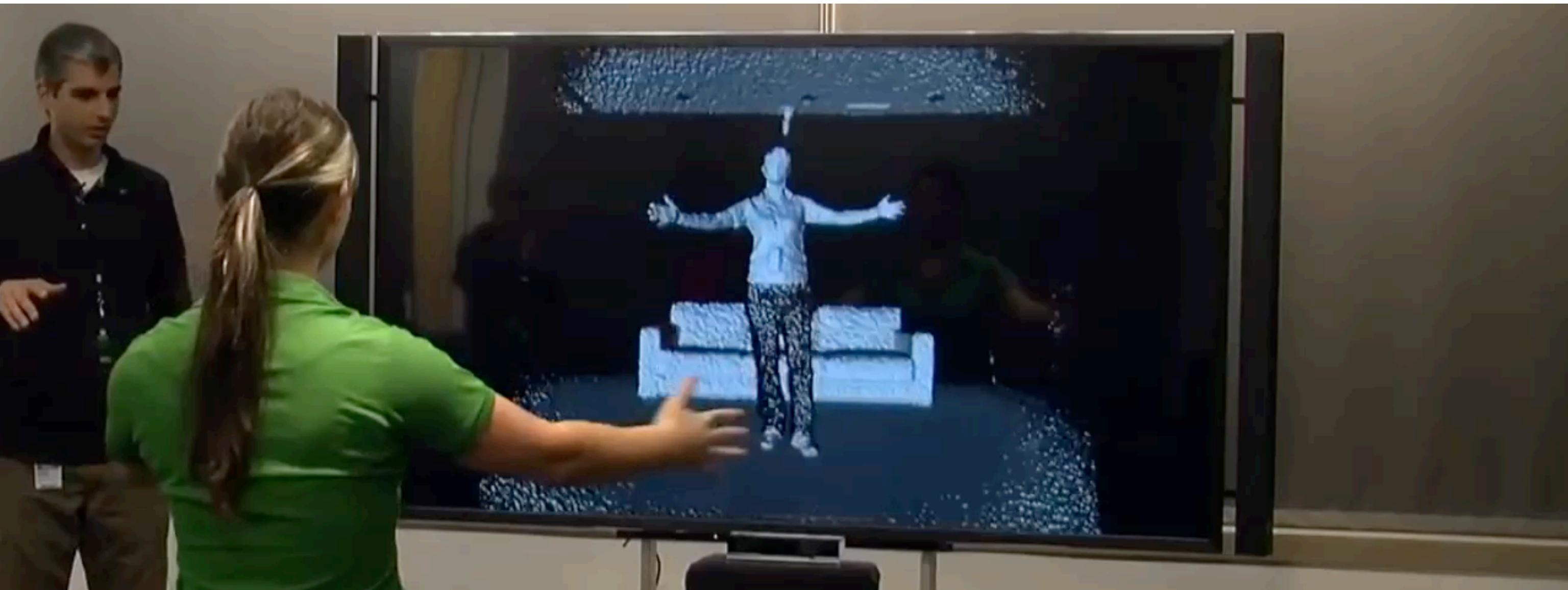


Even larger meshes

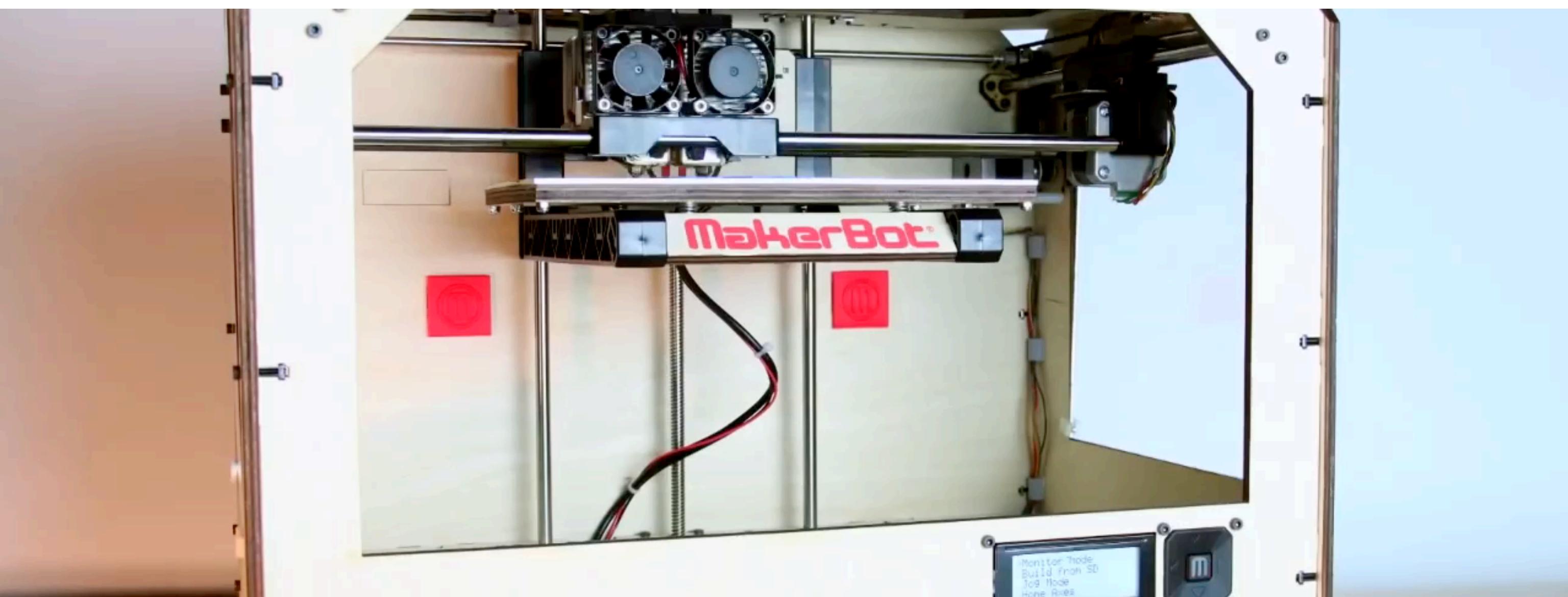


Digital geometry processing: motivations

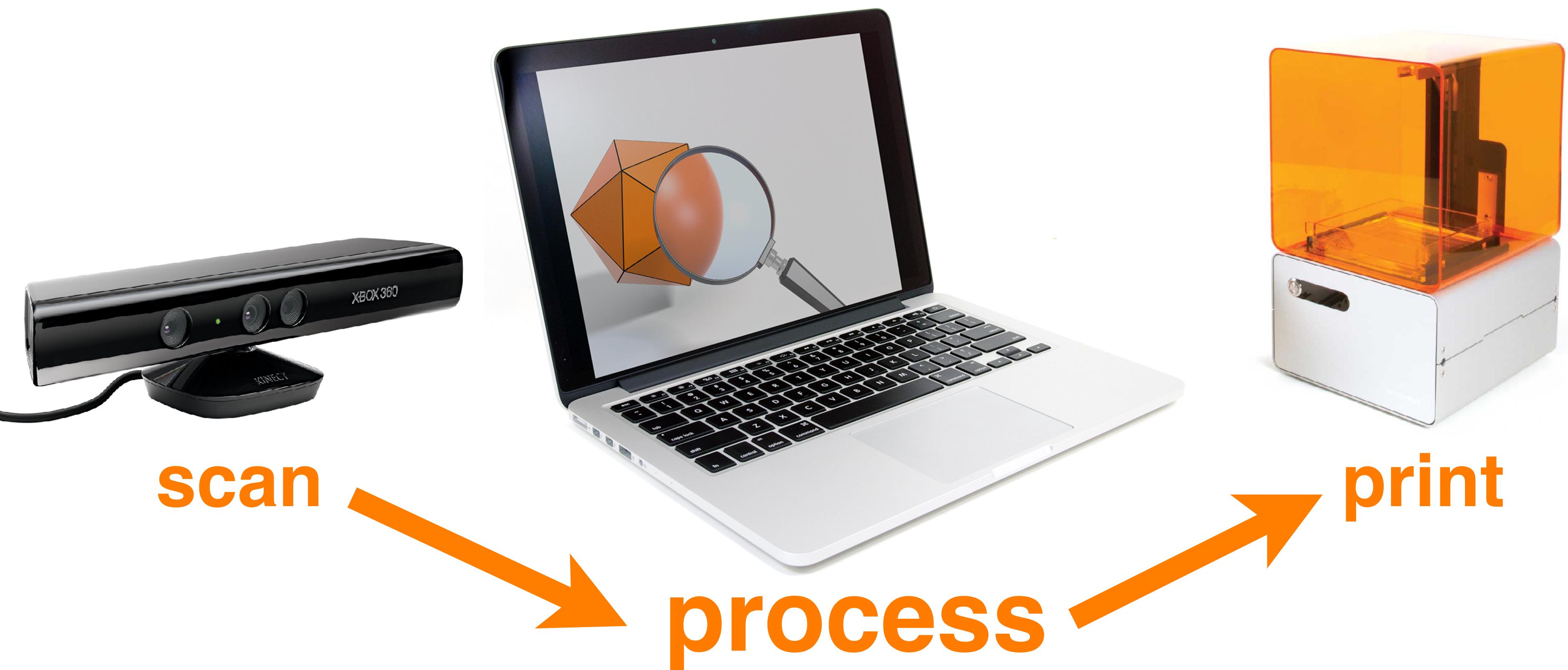
3D Scanning



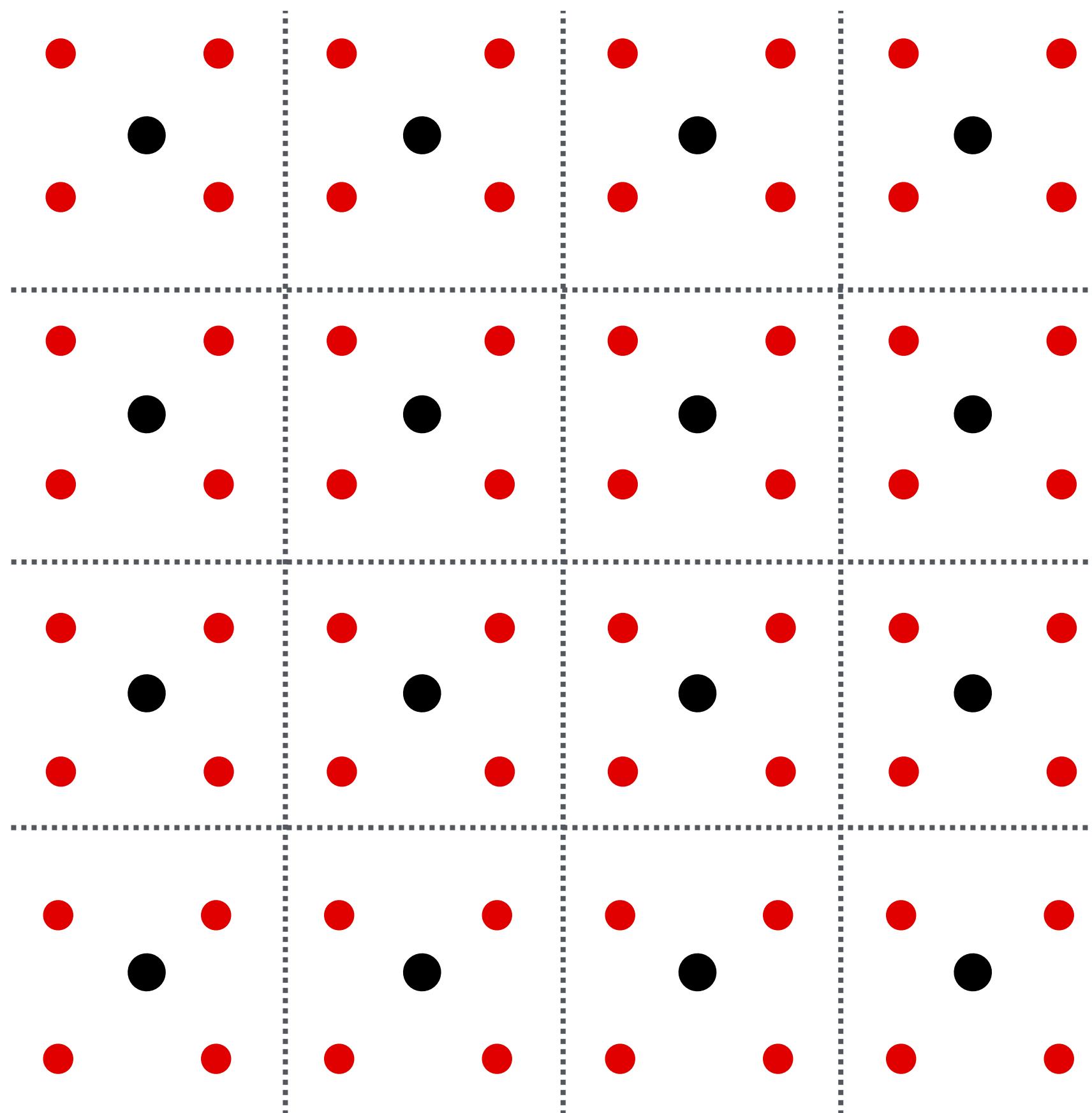
3D Printing



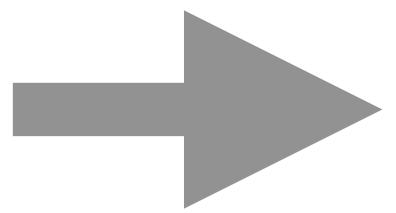
Geometry processing pipeline



Recall: image upsampling



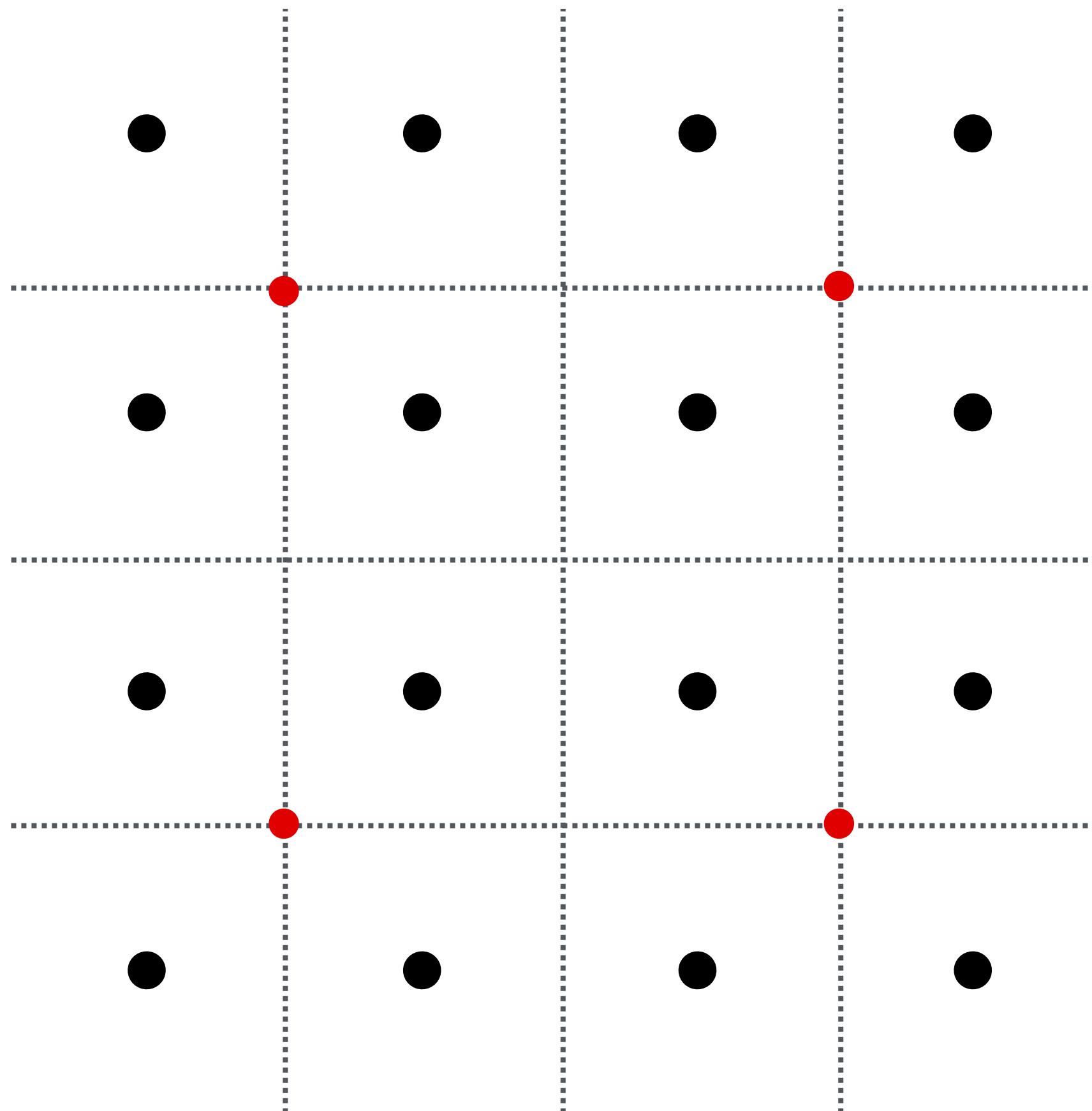
Recall: image upsampling



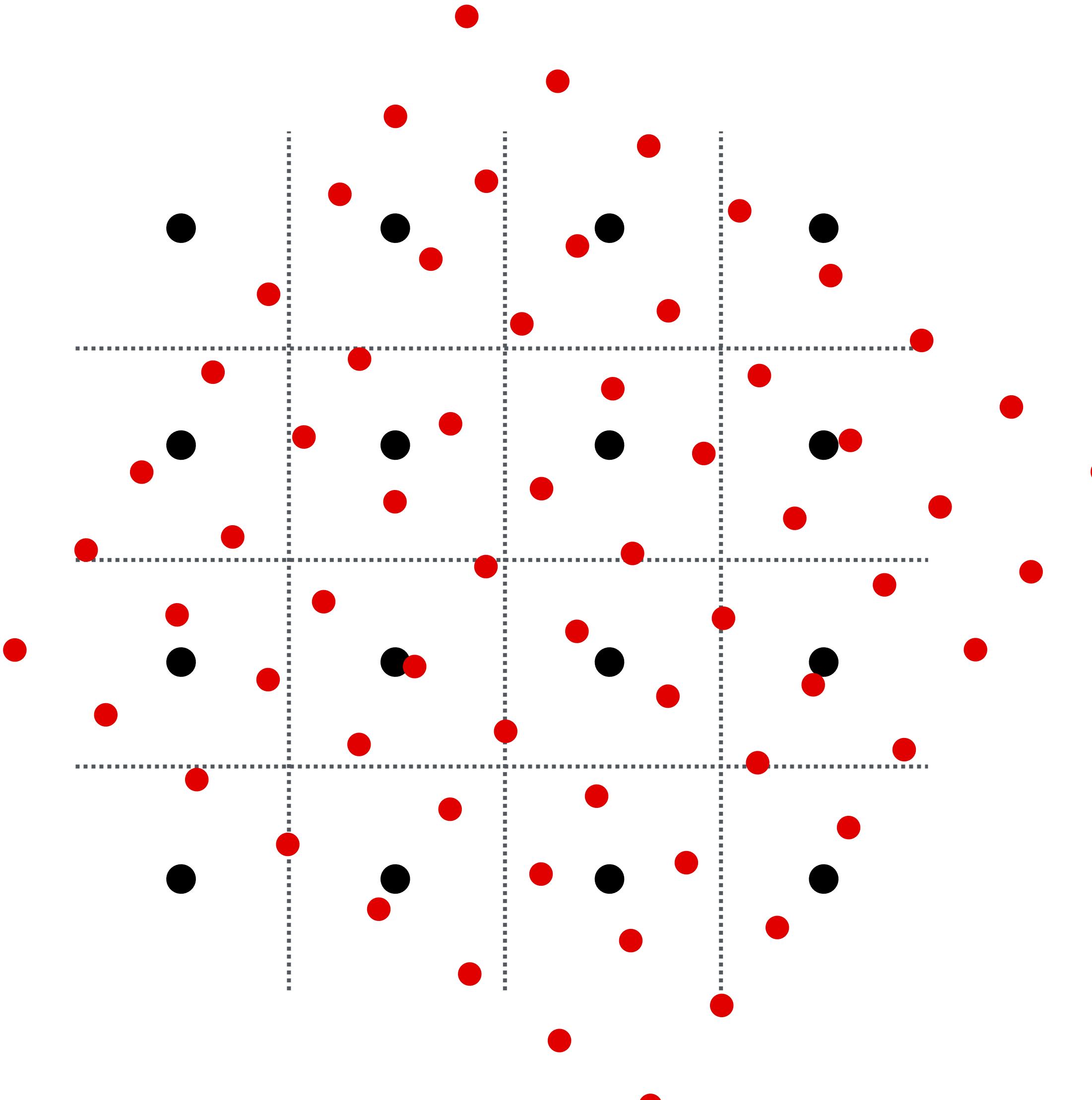
**Upsampling via
bilinear interpolation**



Recall: image downsampling

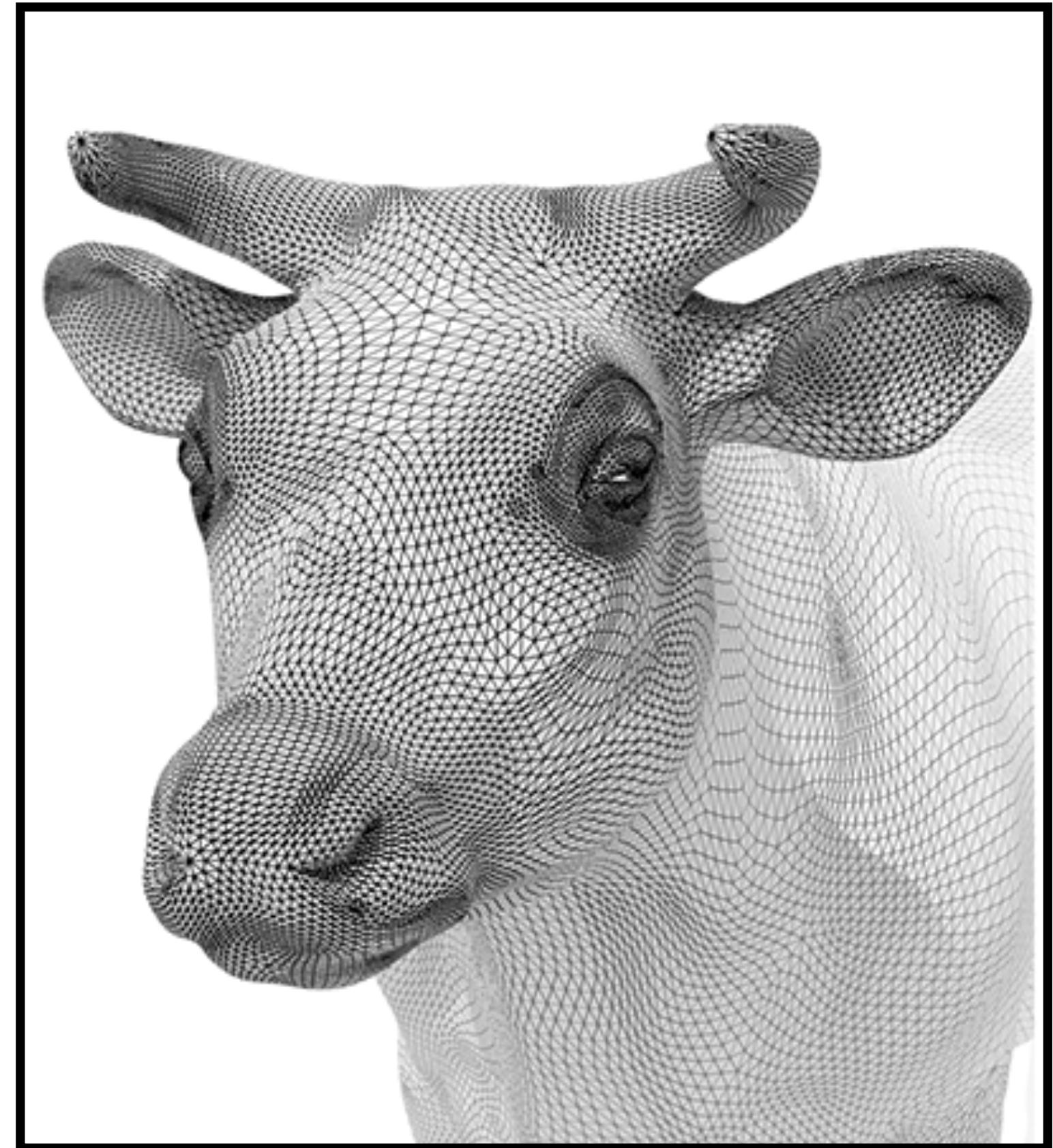
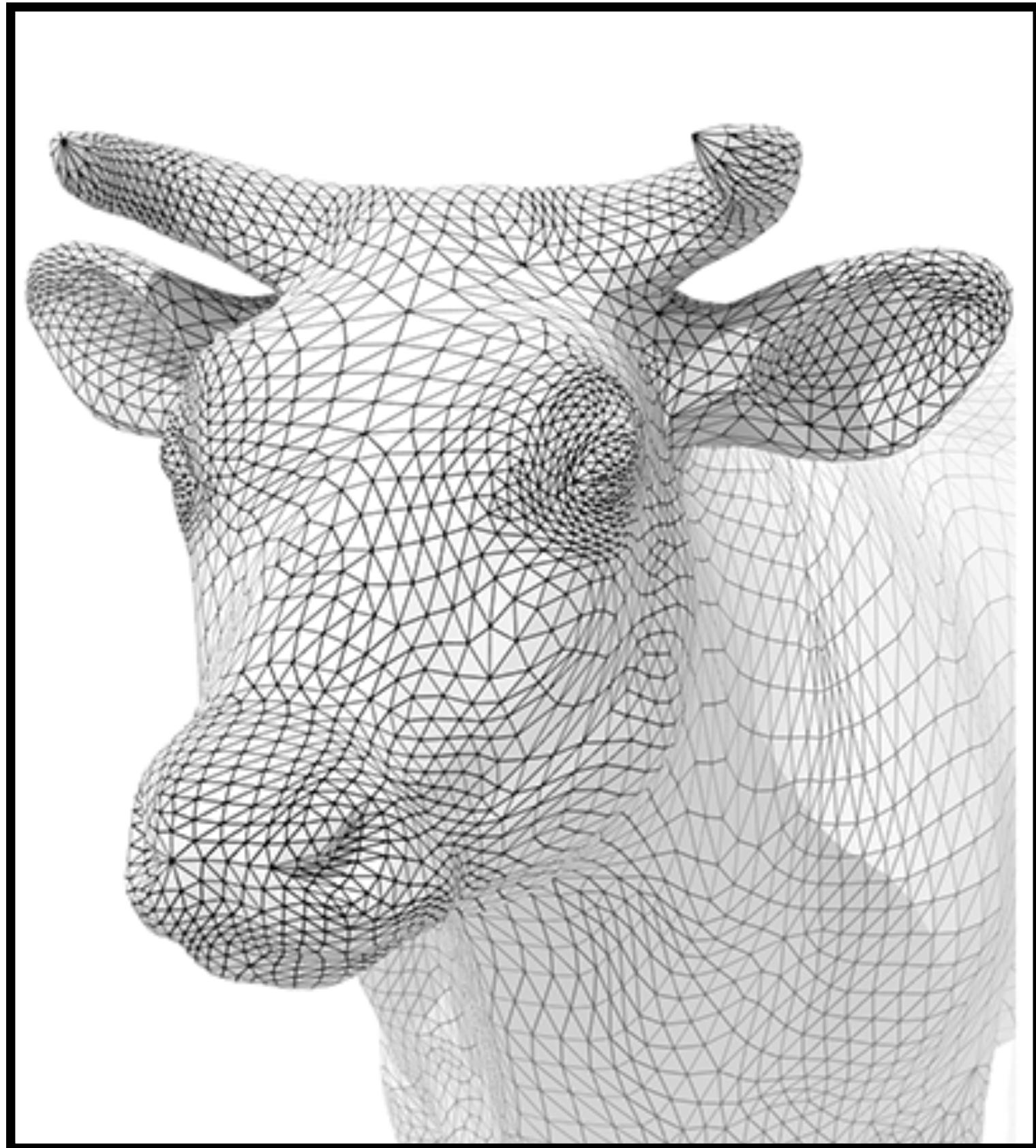


Recall: image resampling



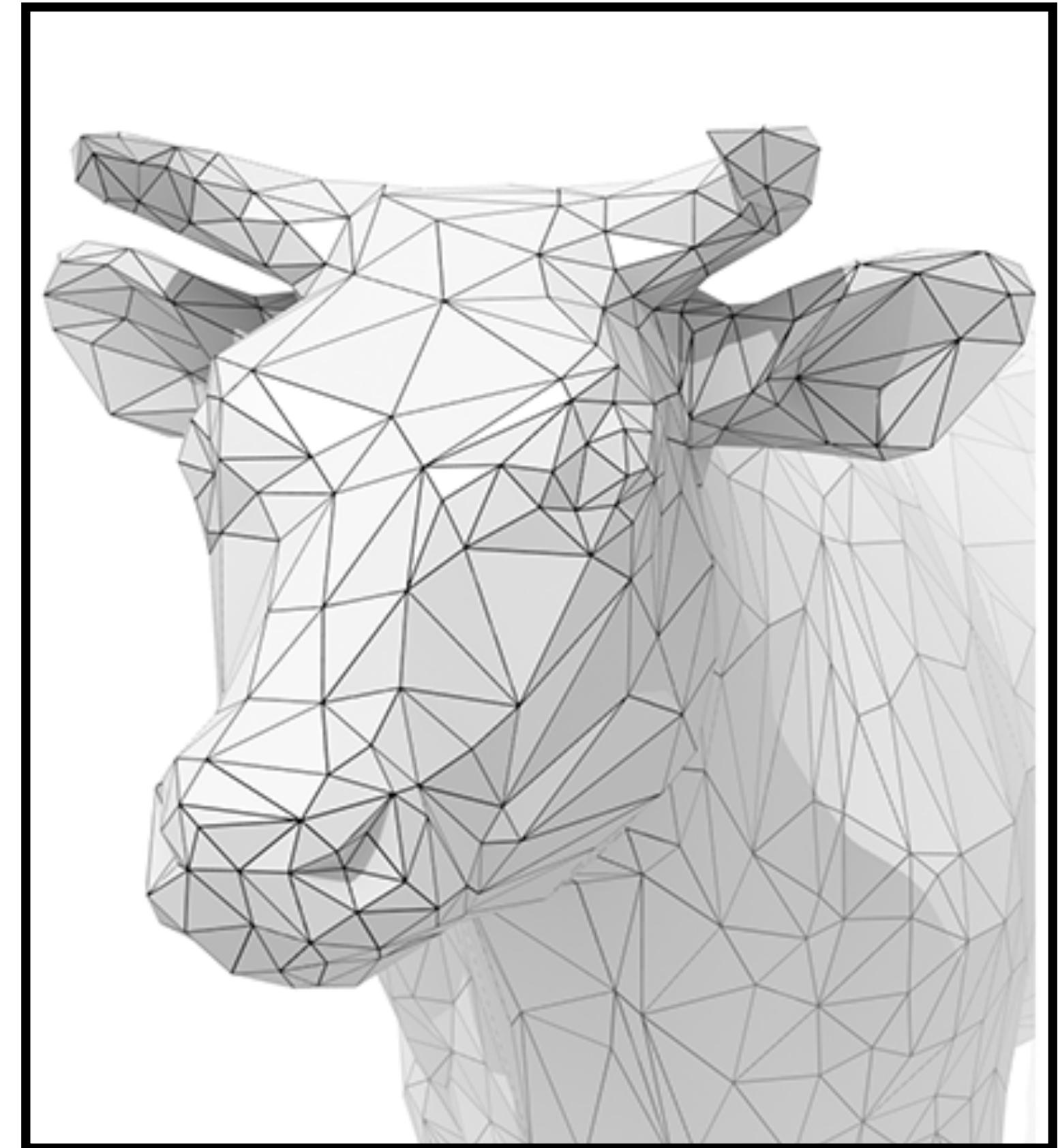
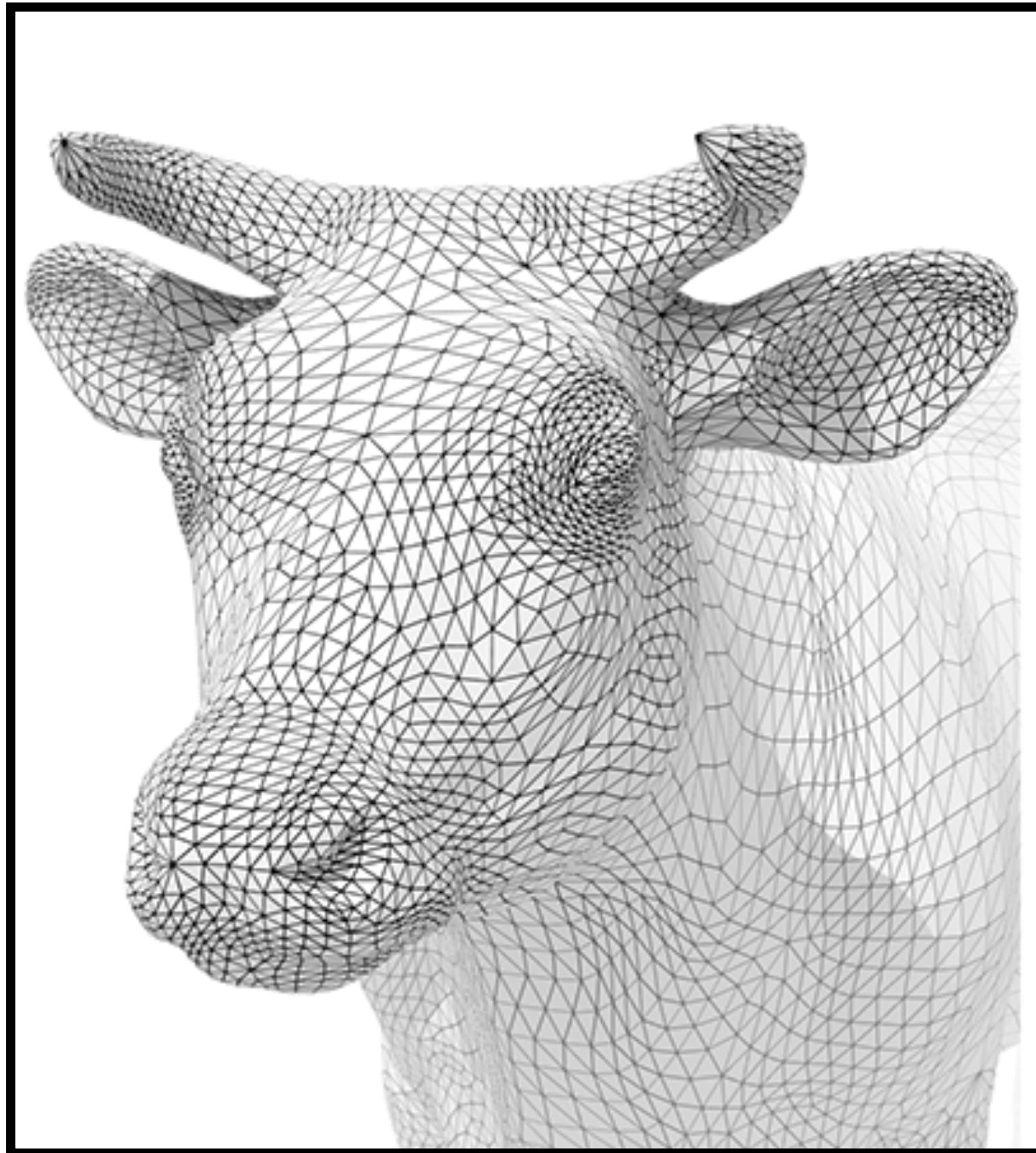
Examples of geometry processing

Mesh upsampling — subdivision



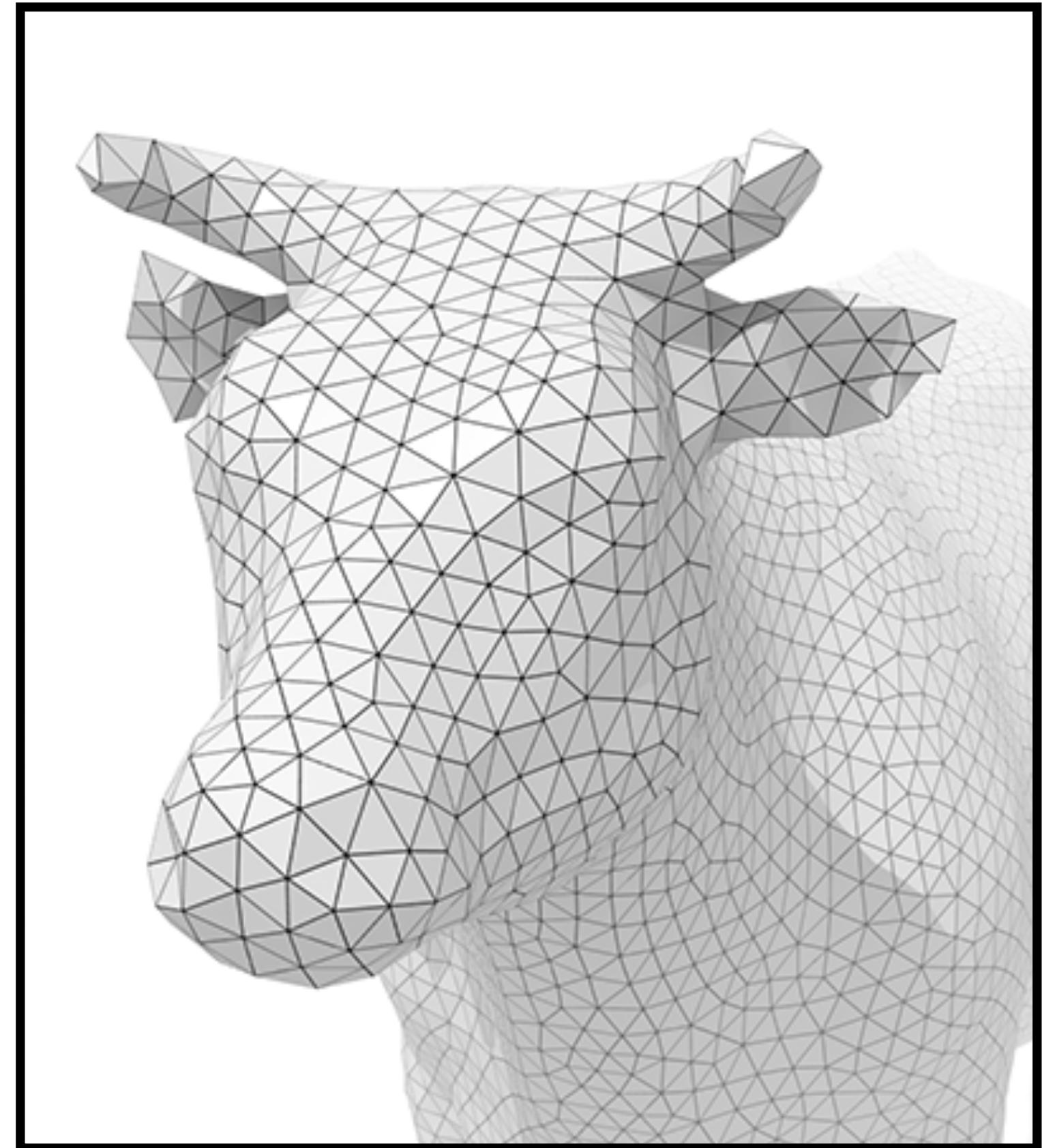
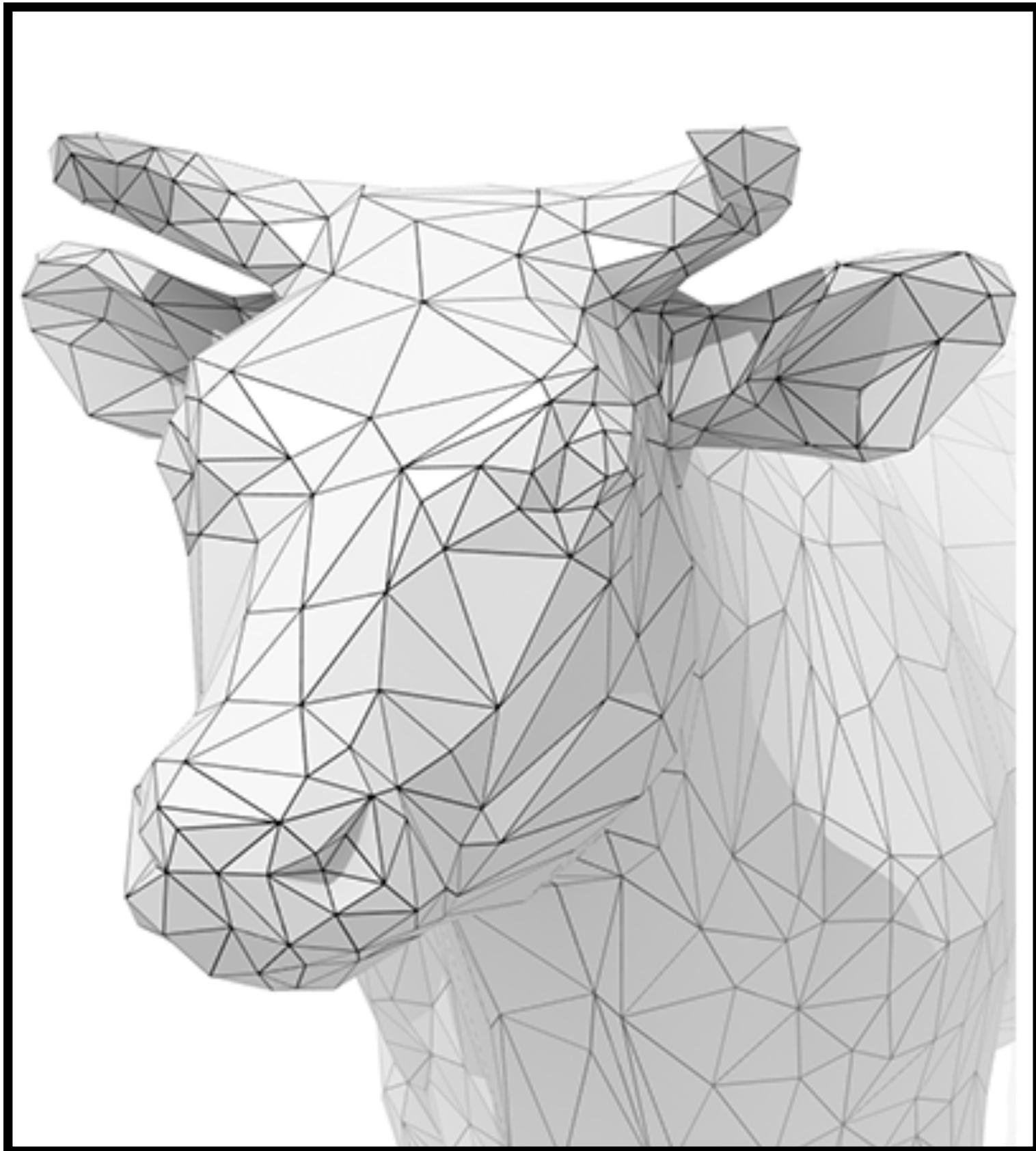
Increase resolution via interpolation

Mesh downsampling — simplification



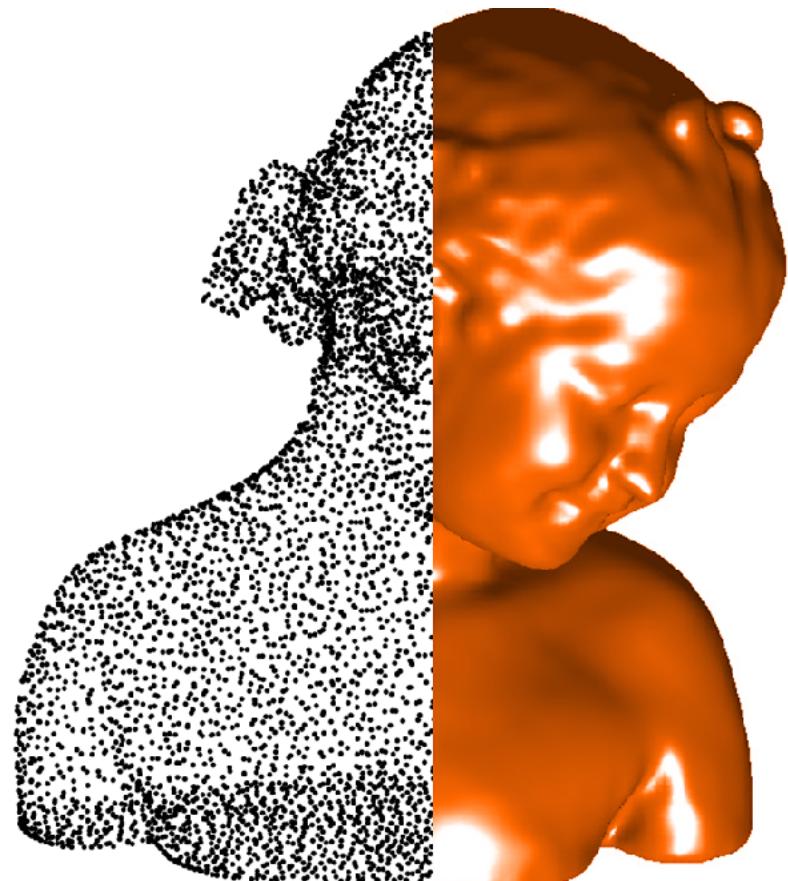
Decrease resolution; try to preserve shape/appearance

Mesh resampling — regularization

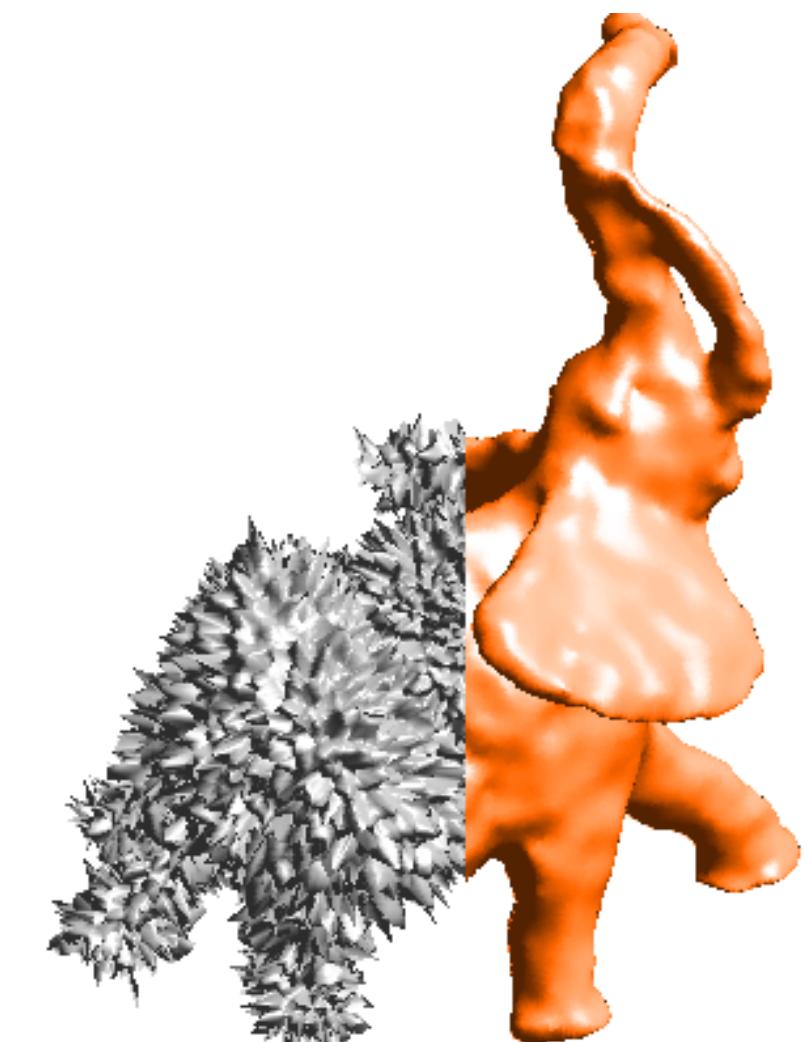


Modify sample distribution to improve quality

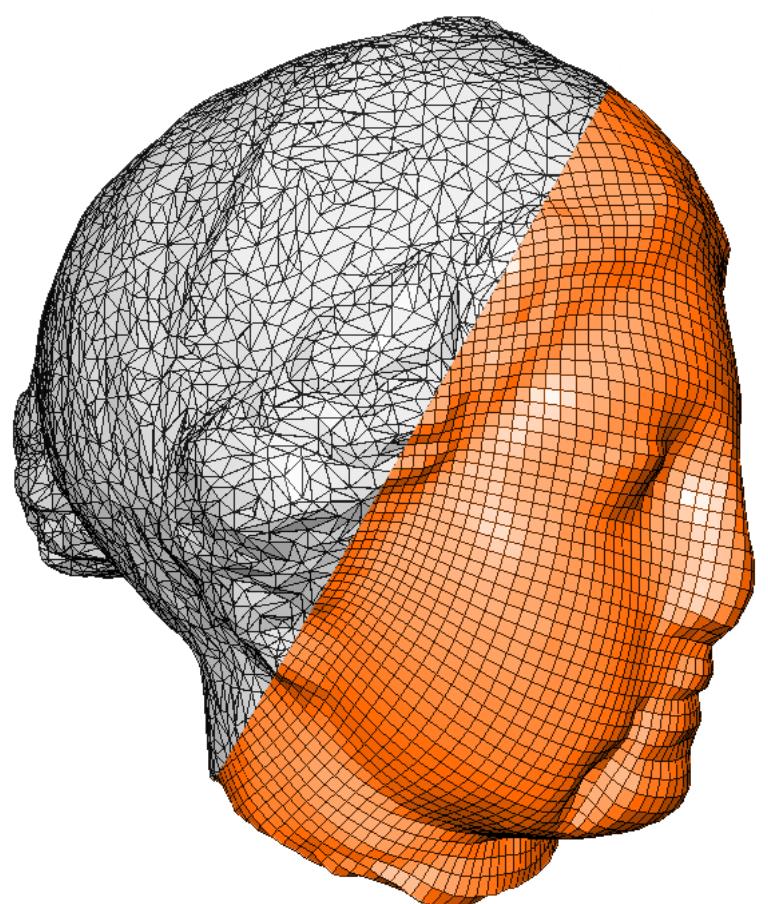
More geometry processing tasks



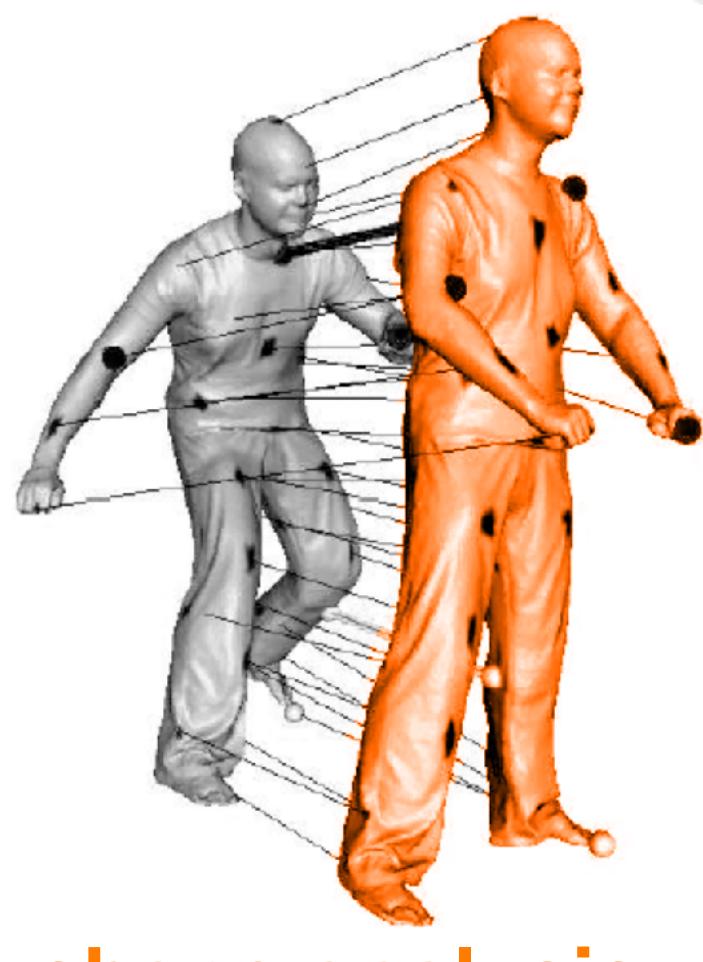
reconstruction



filtering



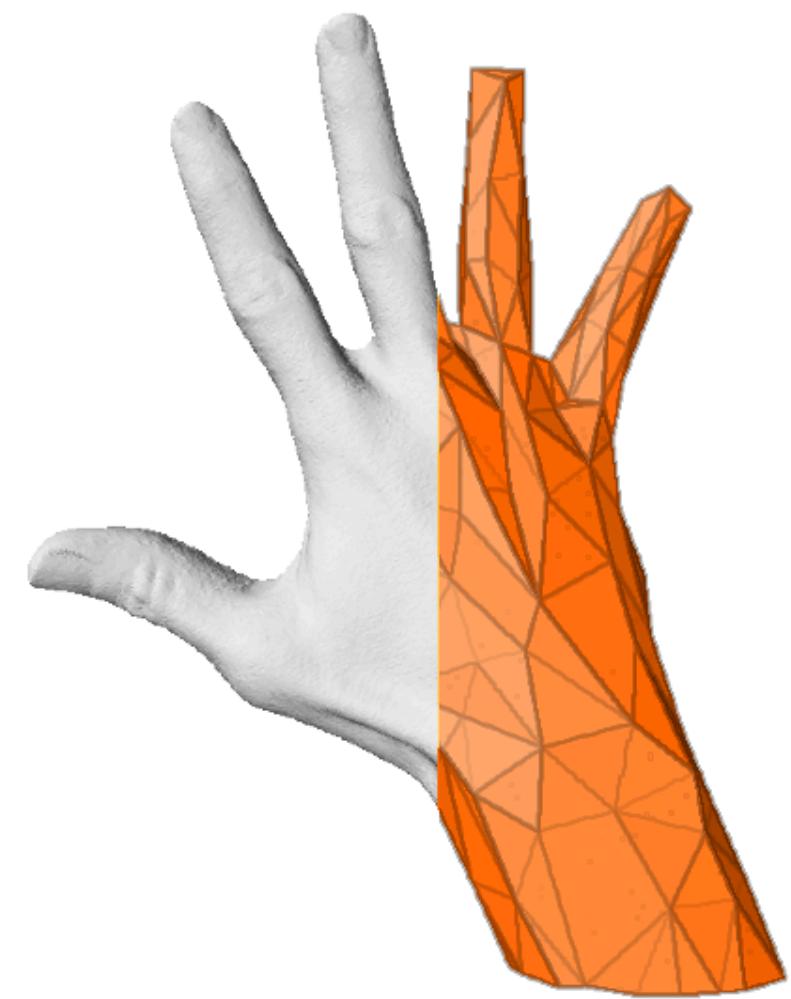
remeshing



shape analysis



parameterization



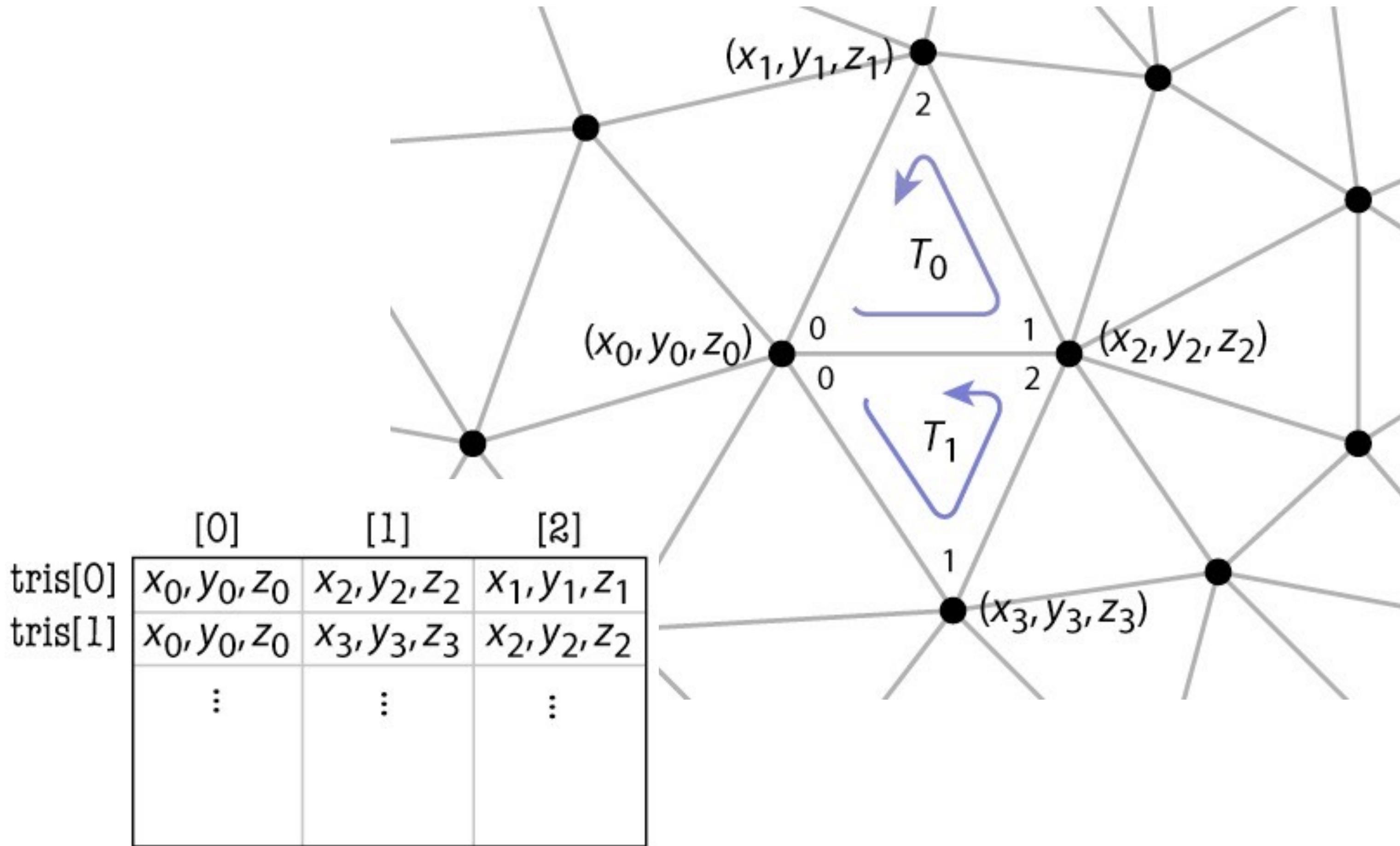
compression

Today

- **Study how to represent meshes (data structures)**
- **Study how to process meshes (basic geometry processing)**
 - **Subdivision**
 - **Mesh simplification**
 - **Mesh resampling**

Mesh representations

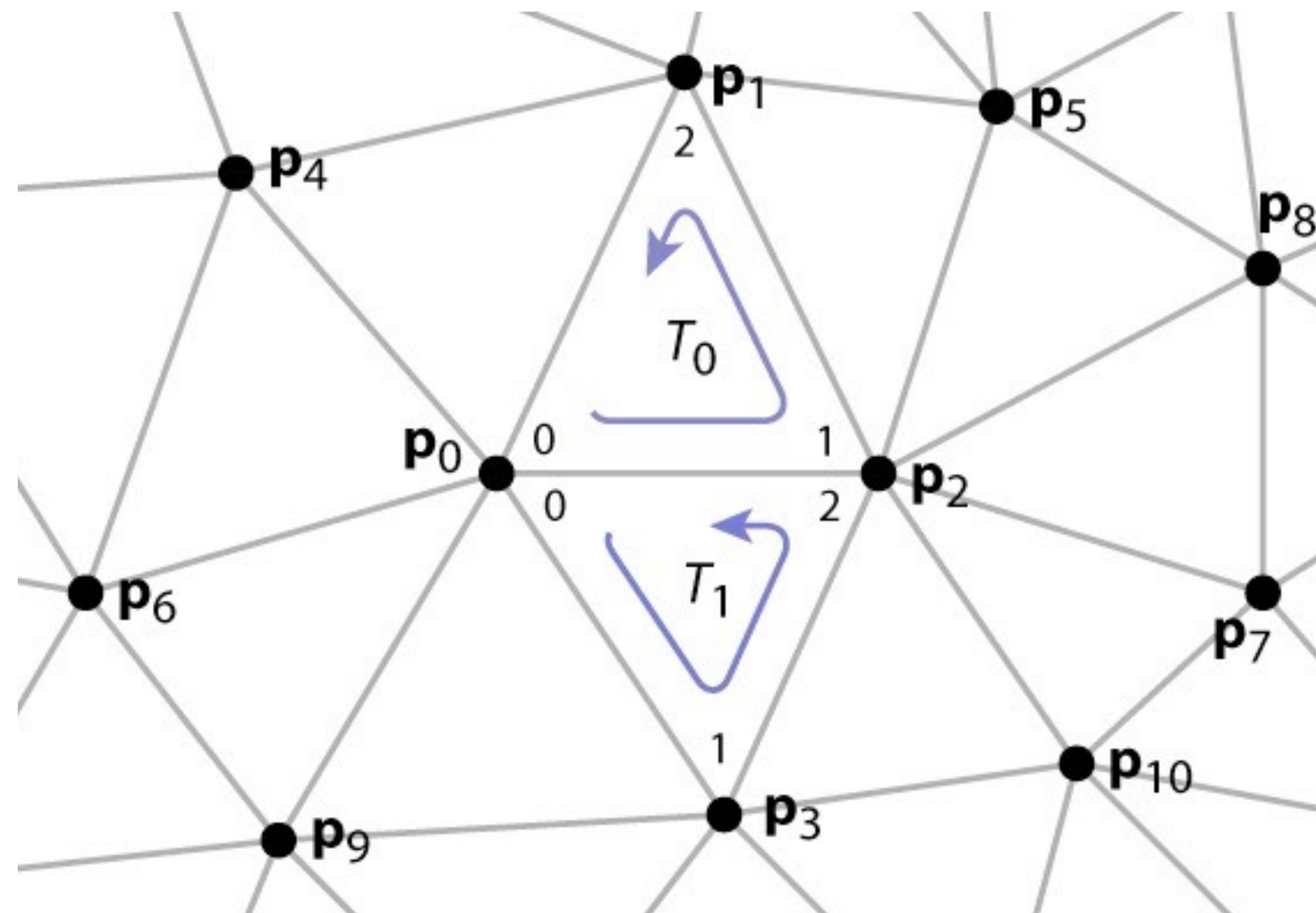
List of triangles



Lists of vertexes / indexed triangle

verts[0]	x_0, y_0, z_0
verts[1]	x_1, y_1, z_1
	x_2, y_2, z_2
	x_3, y_3, z_3
:	

tInd[0]	0, 2, 1
tInd[1]	0, 3, 2
:	

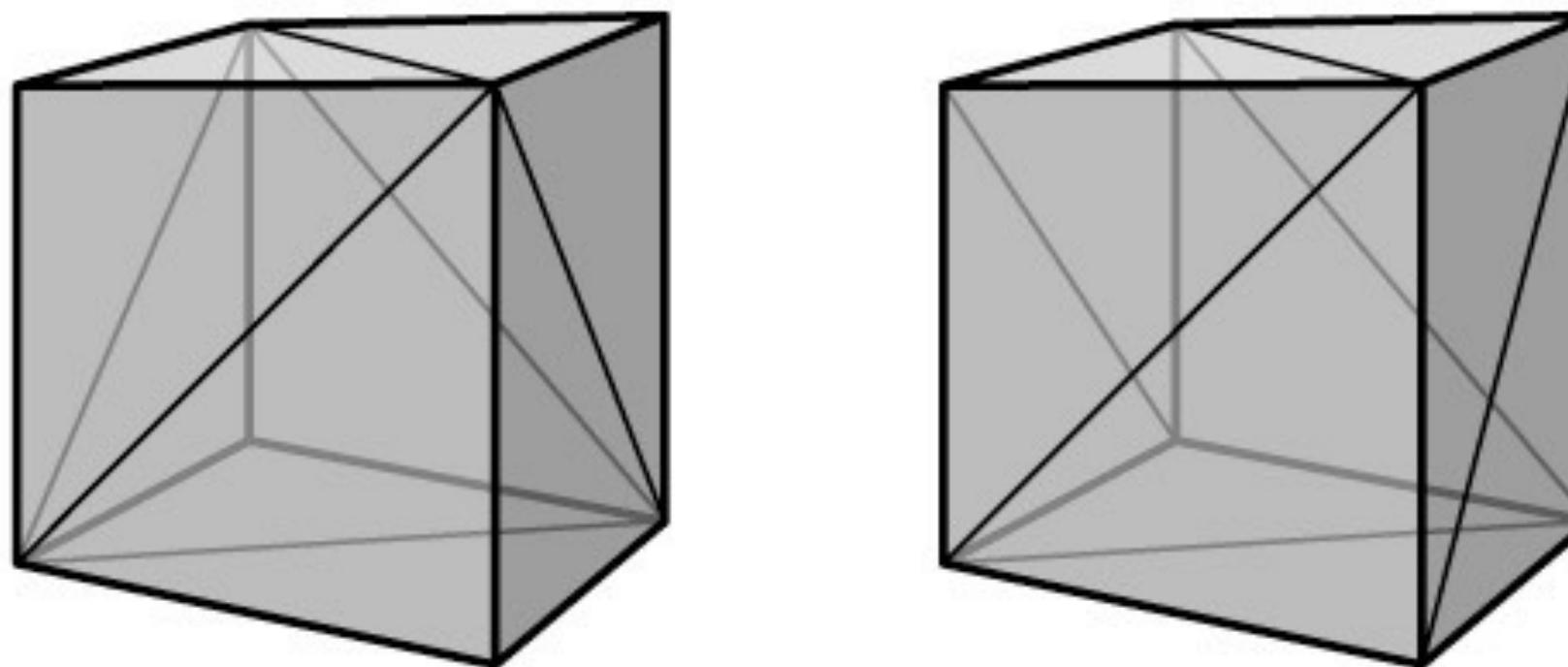


Comparison

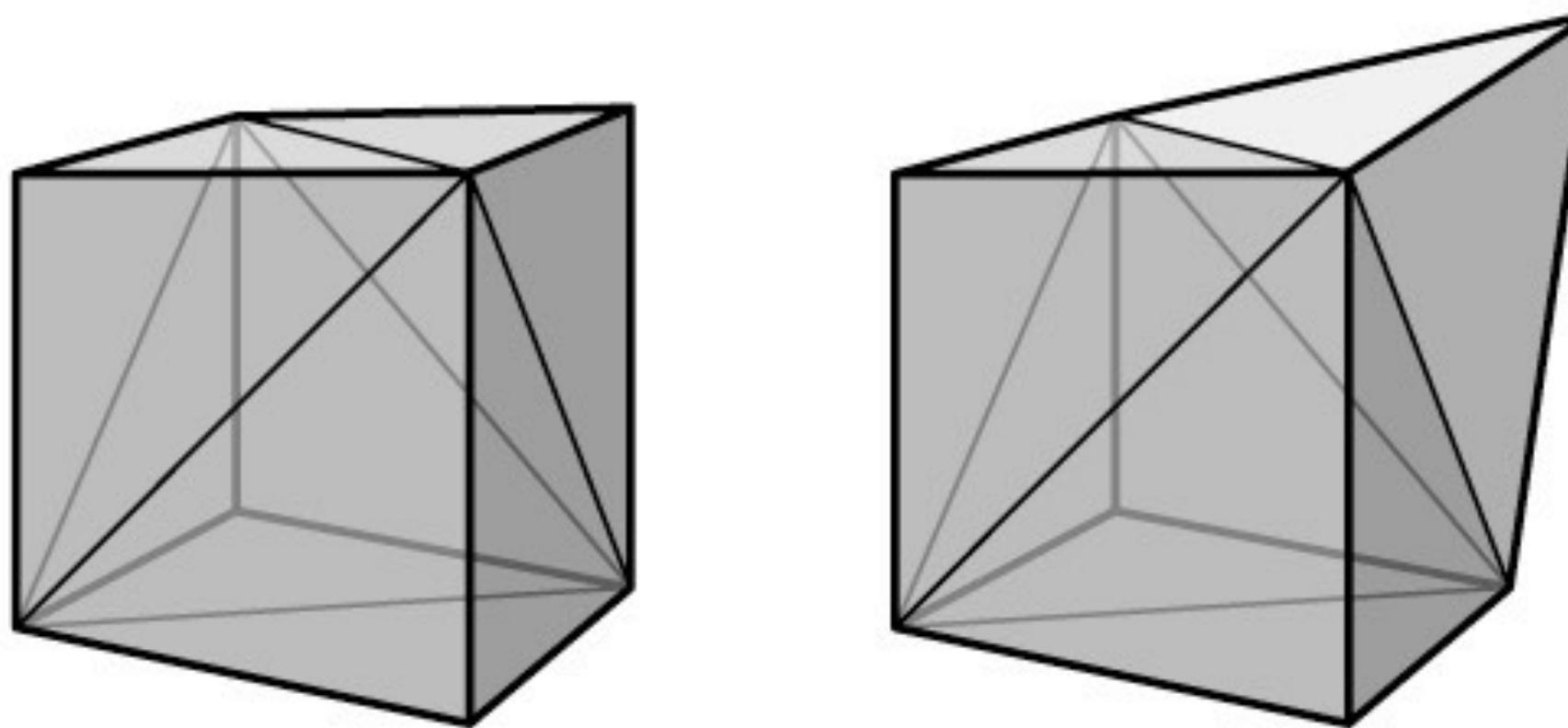
- **List of triangles**
 - + Simple
 - Contains redundant vertex information
- **Vertextes + indexed triangles**
 - + Sharing vertices reduces memory usage
 - + Ensure integrity of the mesh (moving a vertex causes that vertex in all the polygons to move)

Mesh topology vs surface geometry

Same vertex positions, different mesh topology



Same topology, different vertex positions



Topological mesh information

■ Applications:

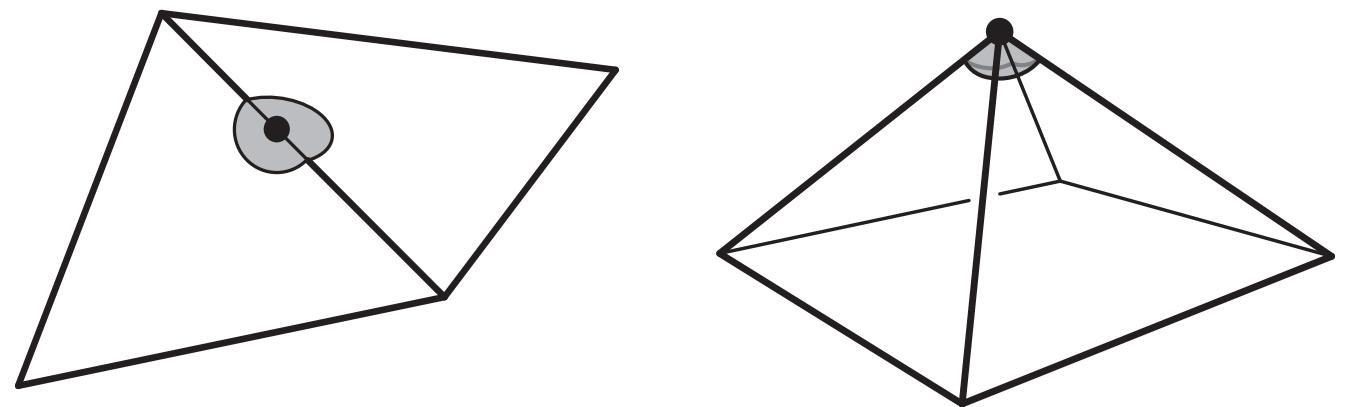
- Constant time access to neighbors
 - e.g. surface normal calculation, subdivision
- Editing the geometry
 - e.g. adding/removing vertices, faces, edges, etc.

■ Solution: topological data structures

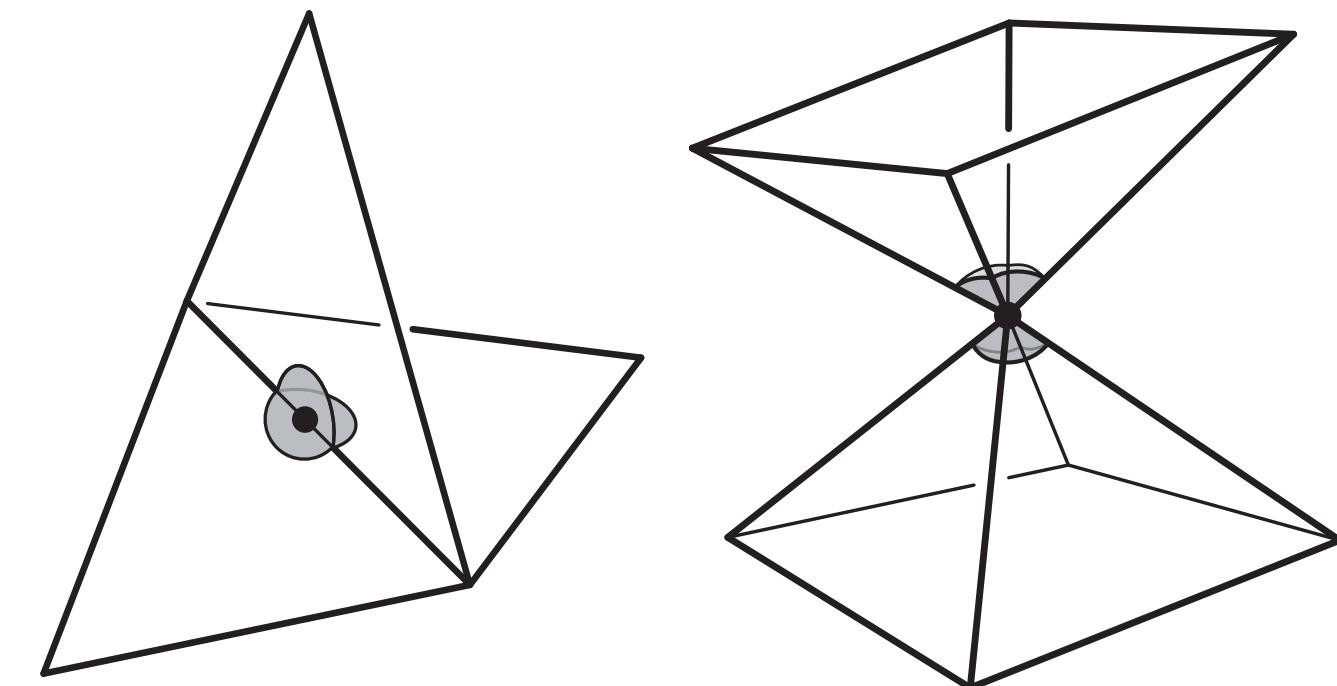
Topological validity: manifold

- Recall, a 2D manifold is a surface that when cut with a small sphere always yields a disk

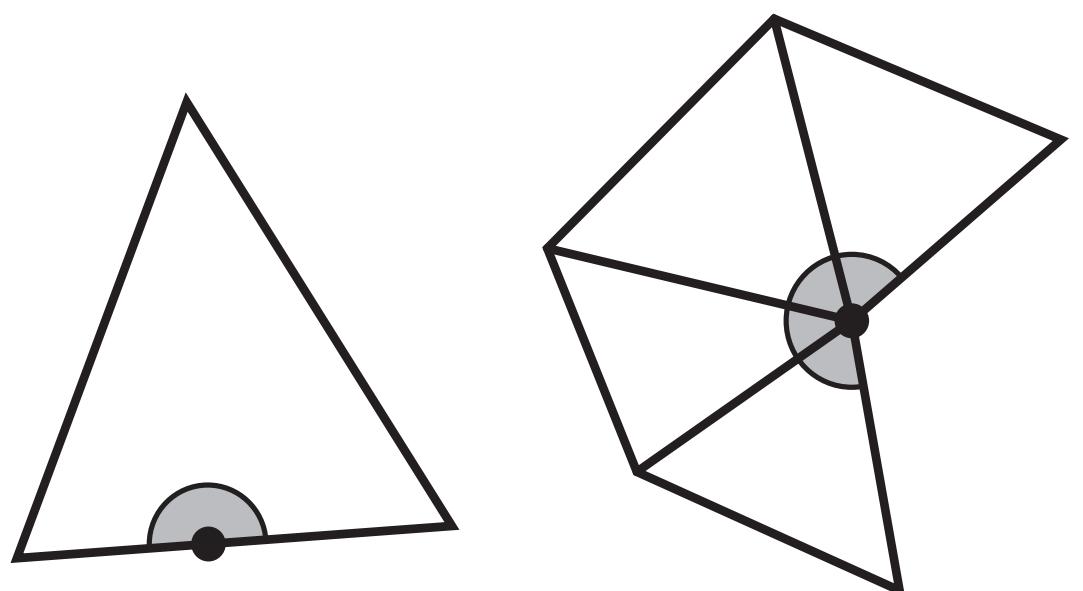
Manifold



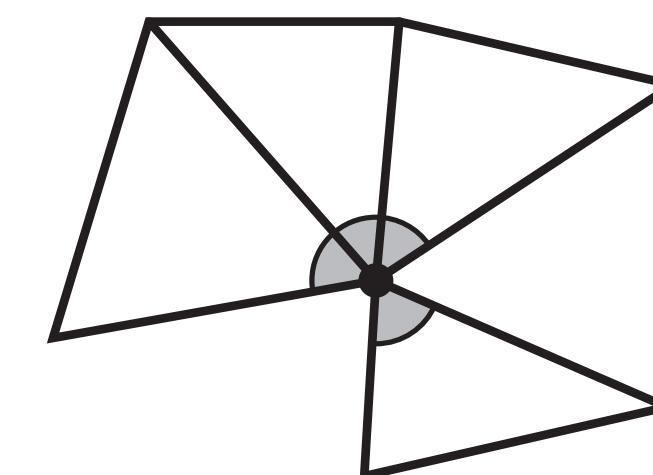
Not manifold



With border



With border



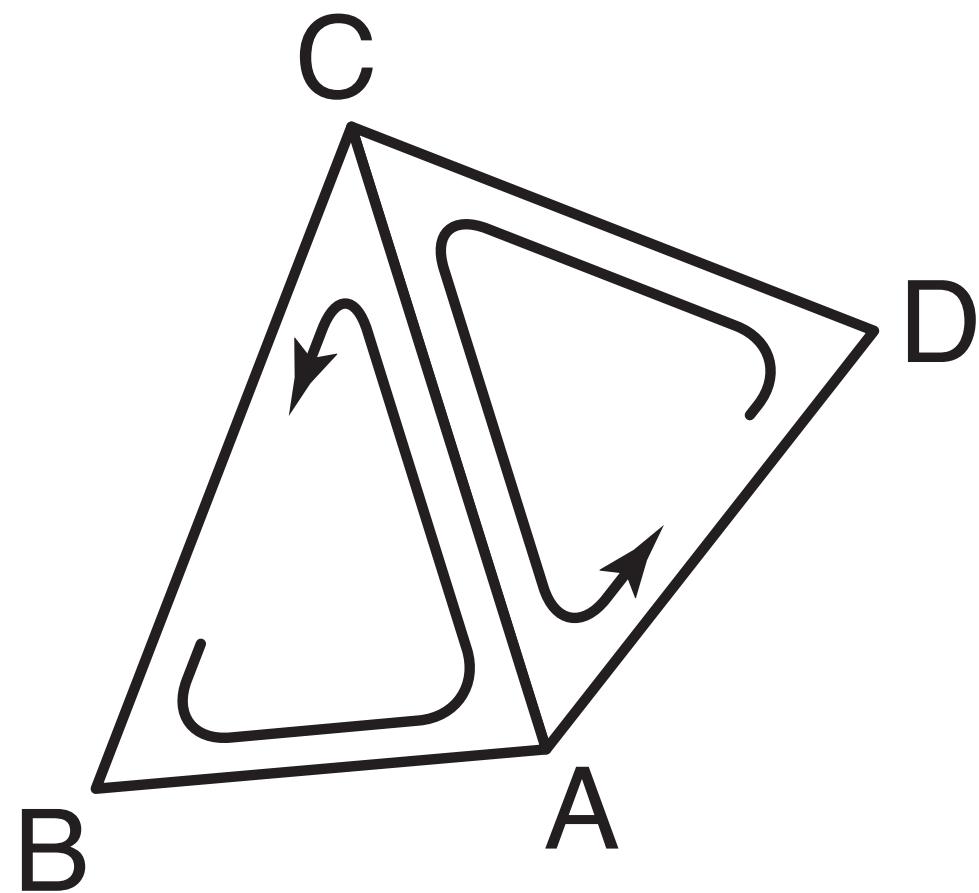
Manifolds have useful properties

- A 2D manifold is a surface that when cut with a small sphere always yields a disk
- If a mesh is manifold, we can rely on these useful properties: *
 - An edge connects exactly two faces
 - An edge connects exactly two vertices
 - A face consists of a ring of edges and vertices
 - A vertex consists of a ring of edges and faces
 - Euler's polyhedron formula holds: $\#f - \#e + \#v = 2$
(for a surface topologically equivalent to a sphere)
(Check for a cube: $6 - 12 + 8 = 2$)

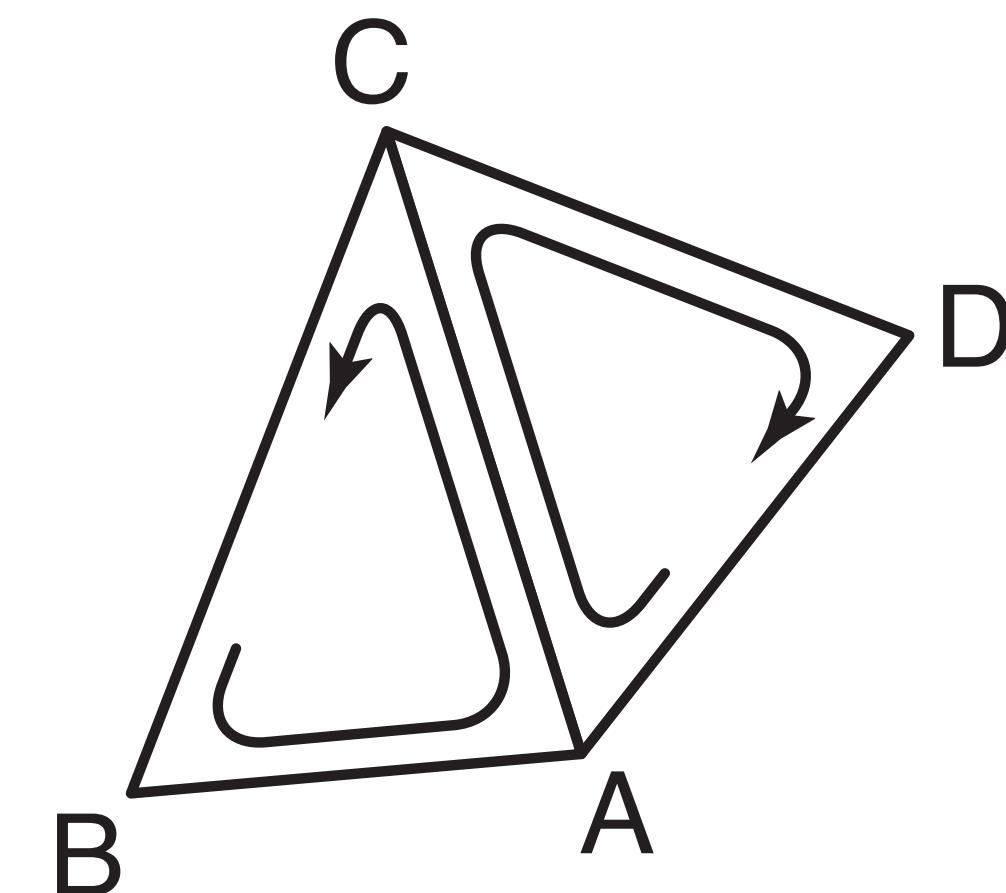
* Some of these properties only apply to non-border mesh regions

Topological validity: orientation consistency

Both facing front



Inconsistent orientations



Non-orientable
(e.g., Moebius strip)

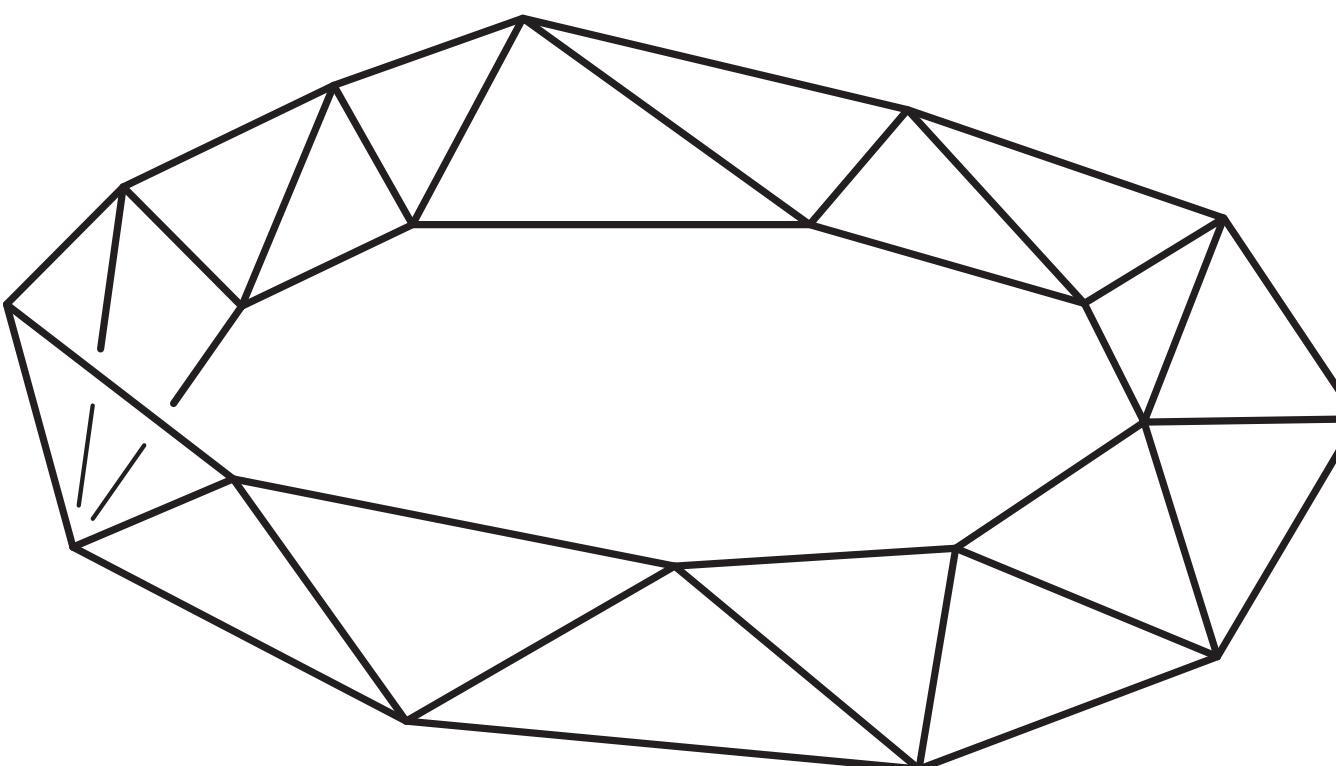
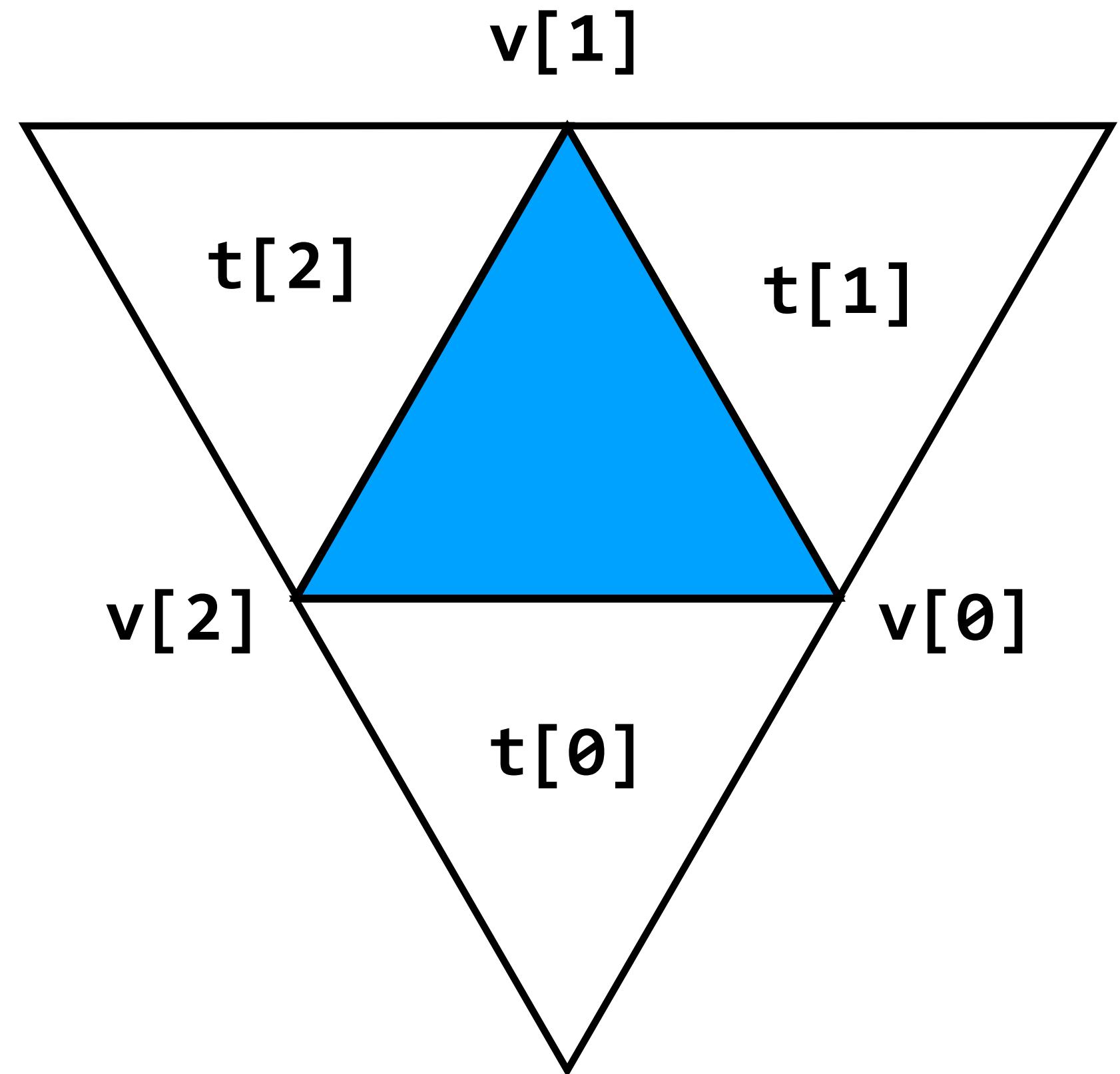


Image credit: Wikipedia

Simple example: triangle-neighbor data structure

```
struct Tri {  
    Vert   * v[3];  
    Tri     * t[3];  
}
```

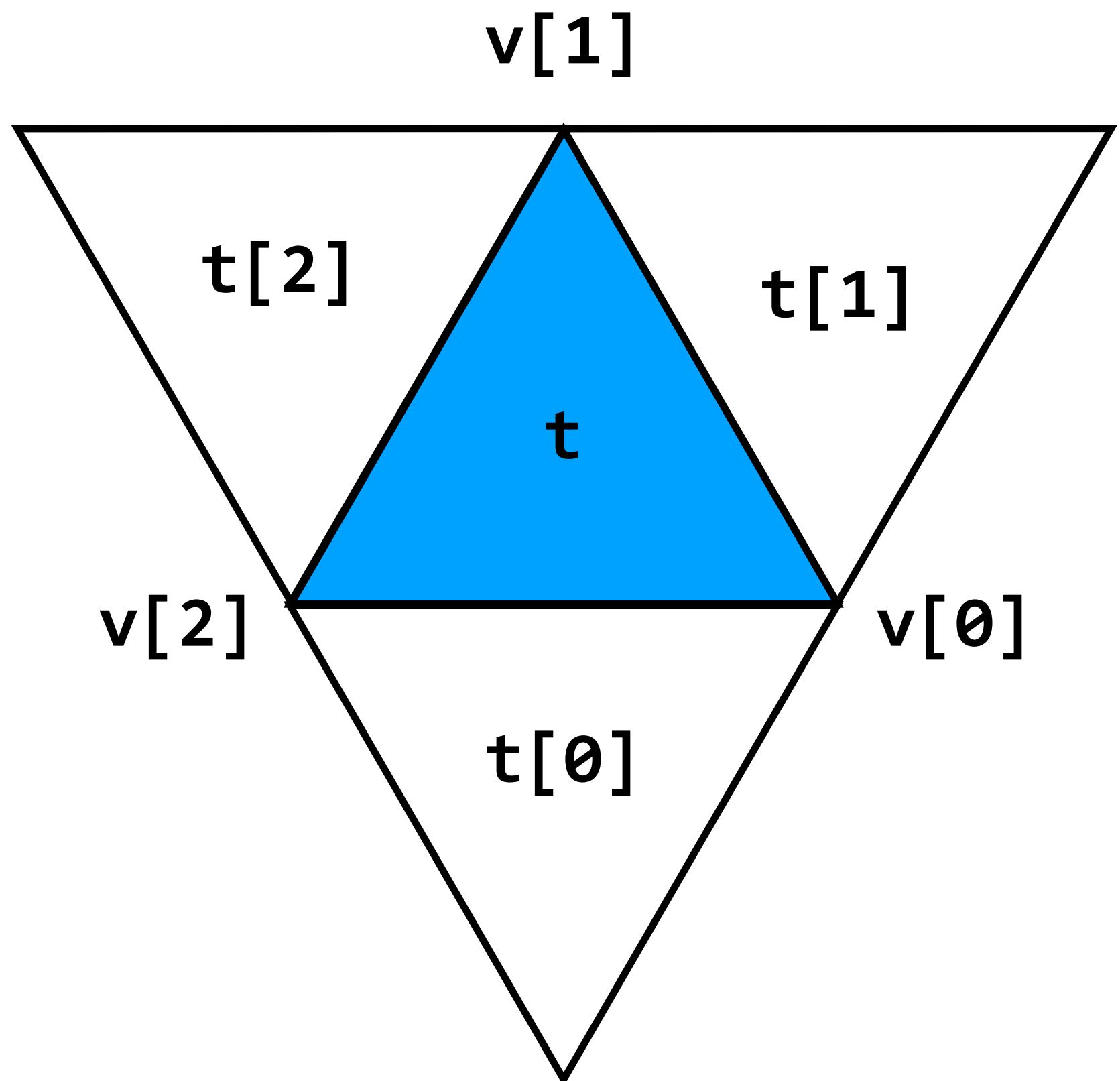
```
struct Vert {  
    Point  pt;  
    Tri    *t;  
}
```



Triangle-neighbor – mesh traversal

Find next triangle counter-clockwise around vertex v from triangle t

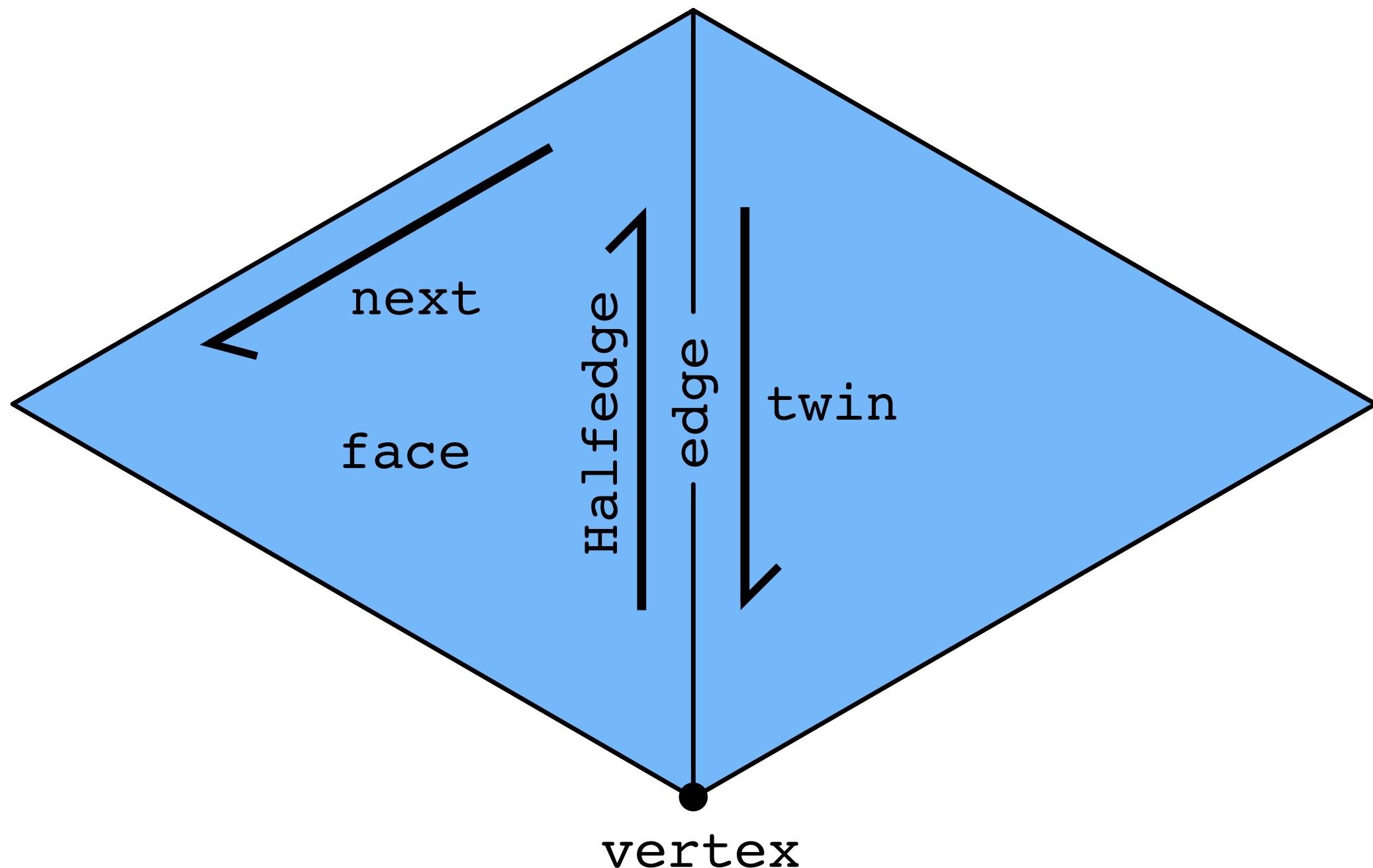
```
Tri* tccwvt(Vert *v, Tri *t)
{
    if (v == t->v[0])
        return t[0];
    if (v == t->v[1])
        return t[1];
    if (v == t->v[2])
        return t[2];
}
```



Half-edge data structure

```
struct Halfedge {  
    Halfedge *twin,  
    Halfedge *next;  
    Vertex *vertex;  
    Edge *edge;  
    Face *face;  
}  
  
struct Vertex {  
    Point pt;  
    Halfedge *halfedge;  
}  
  
struct Edge {  
    Halfedge *halfedge;  
}  
  
struct Face {  
    Halfedge *halfedge;  
}
```

Key idea: two half-edges act as “glue” between mesh elements



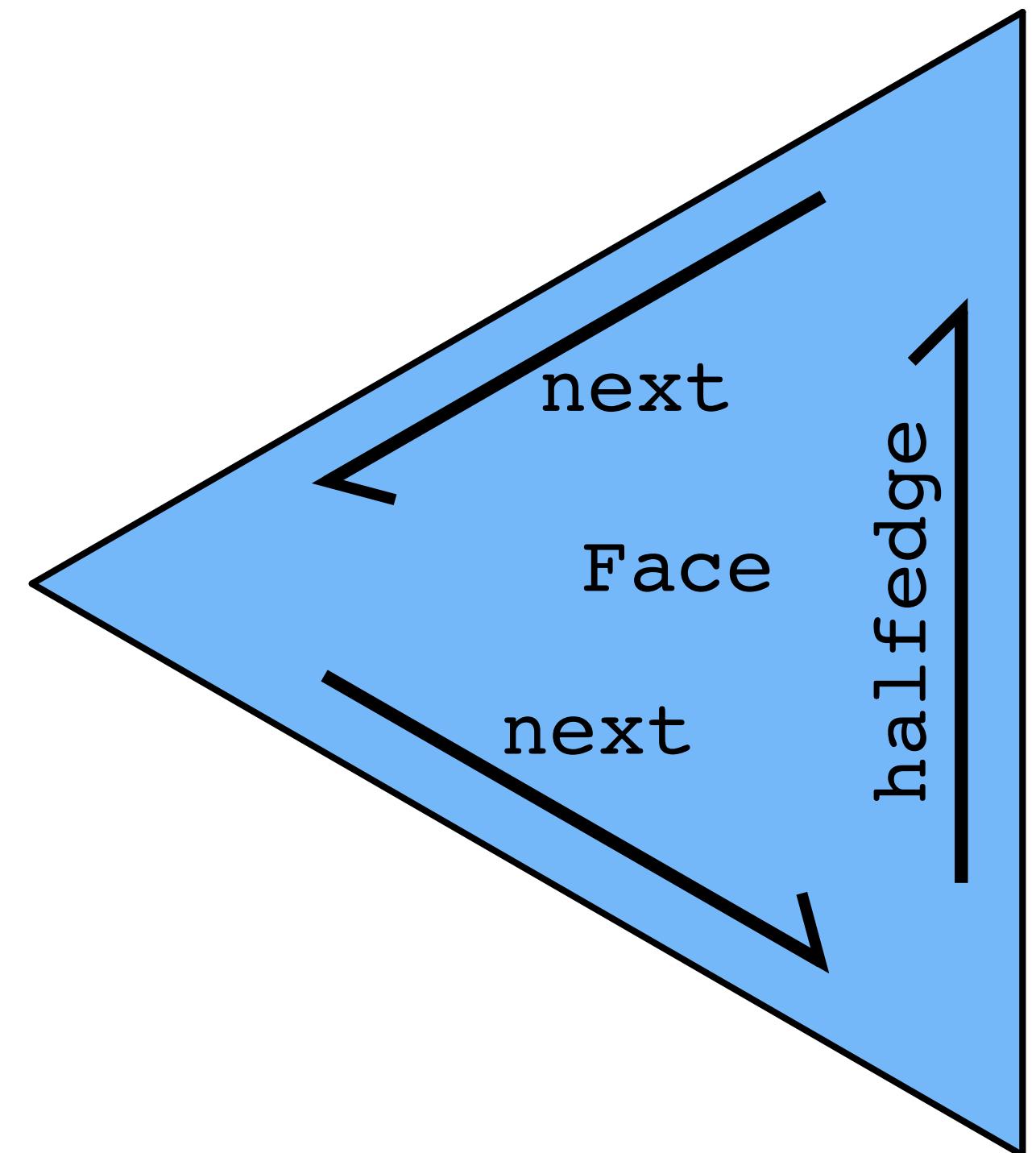
Each vertex, edge and face points to one of its half edges

Half-edge structure facilitates mesh traversal

- Use twin and next pointers to move around mesh
- Process vertex, edge and/or face pointers

Example 1: process all vertices of a face

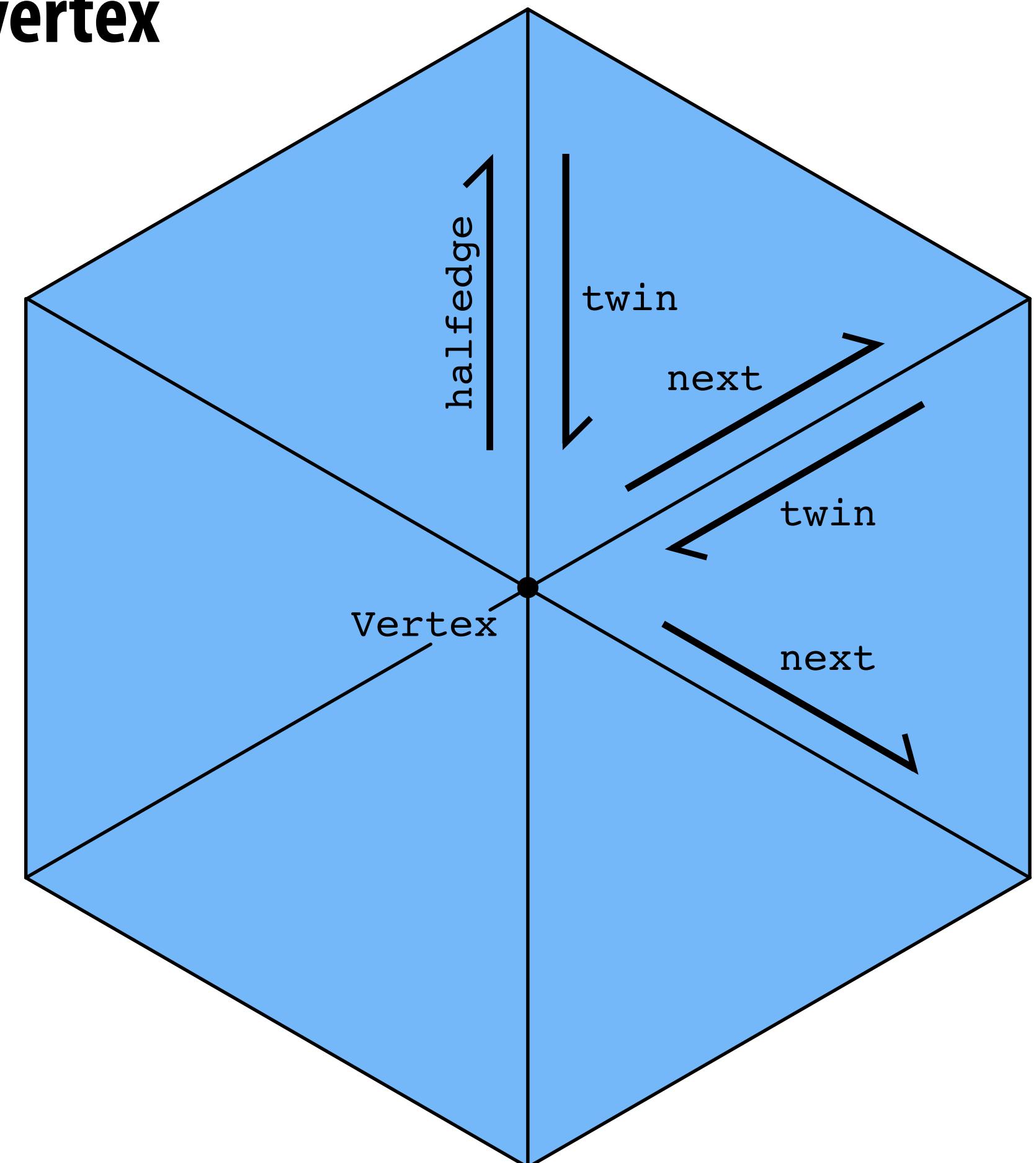
```
Halfedge* h = f->halfedge;  
do {  
    process(h->vertex);  
    h = h->next;  
}  
while( h != f->halfedge );
```



Half-edge structure facilitates mesh traversal

Example 2: process all edges around a vertex

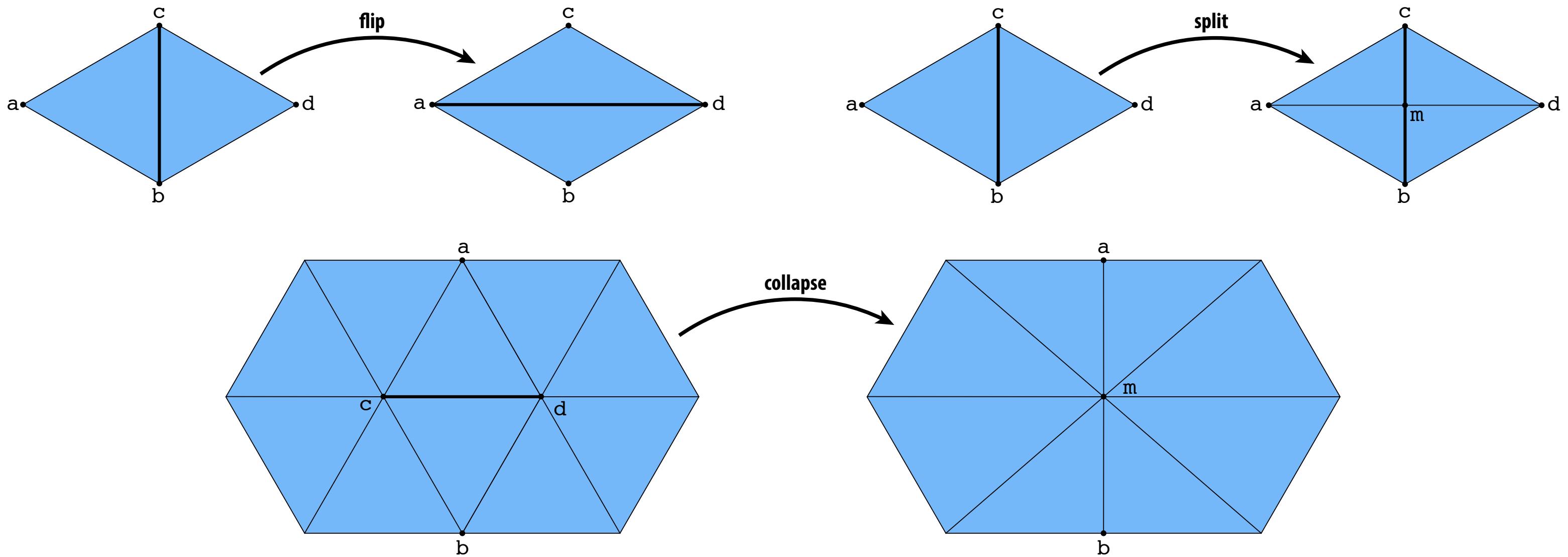
```
Halfedge* h = v->halfedge;  
do {  
    process(h->edge);  
    h = h->twin->next;  
}  
while( h != v->halfedge );
```



Local mesh operations

Half-Edge – local mesh editing

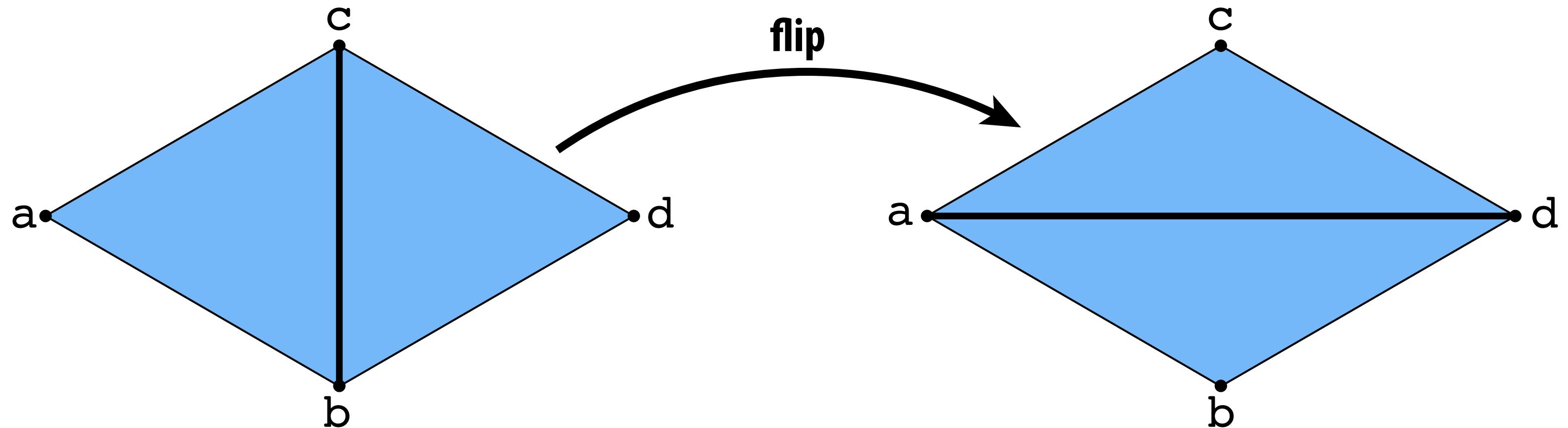
- Consider basic operations for linked list: insert, delete
- Basic ops for half-edge mesh: flip, split, collapse edges



Allocate / delete elements; reassign pointers
(Care is needed to preserve mesh manifold property)

Half-edge – edge flip

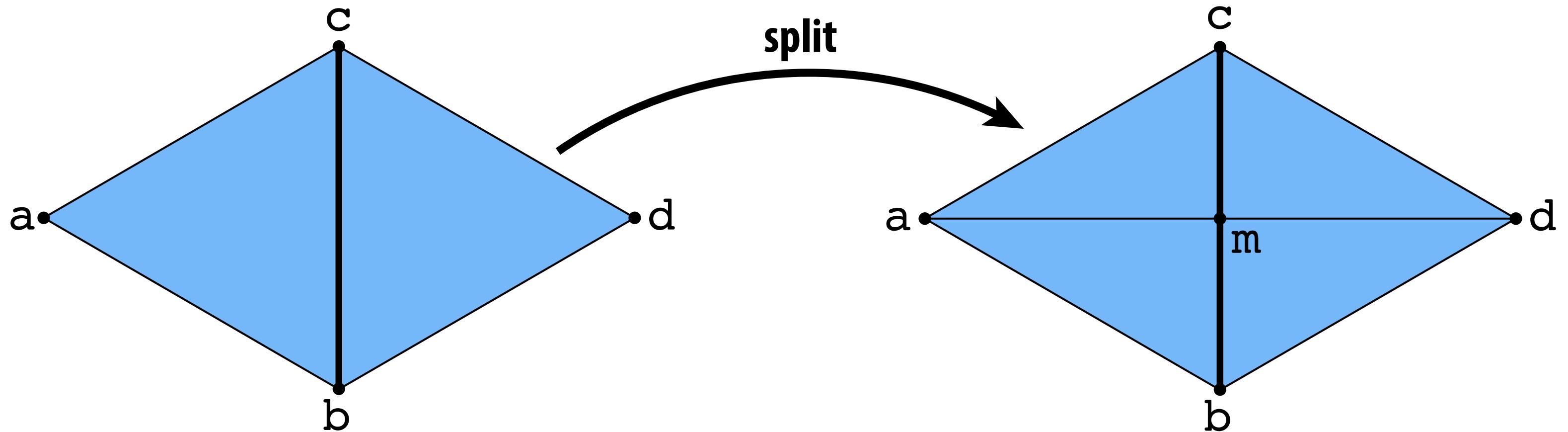
Triangles (a,b,c) , (b,d,c) become (a,d,c) , (a,b,d) :



- Long list of half-edge pointer reassessments
- However, no mesh elements created/destroyed

Half-edge – edge split

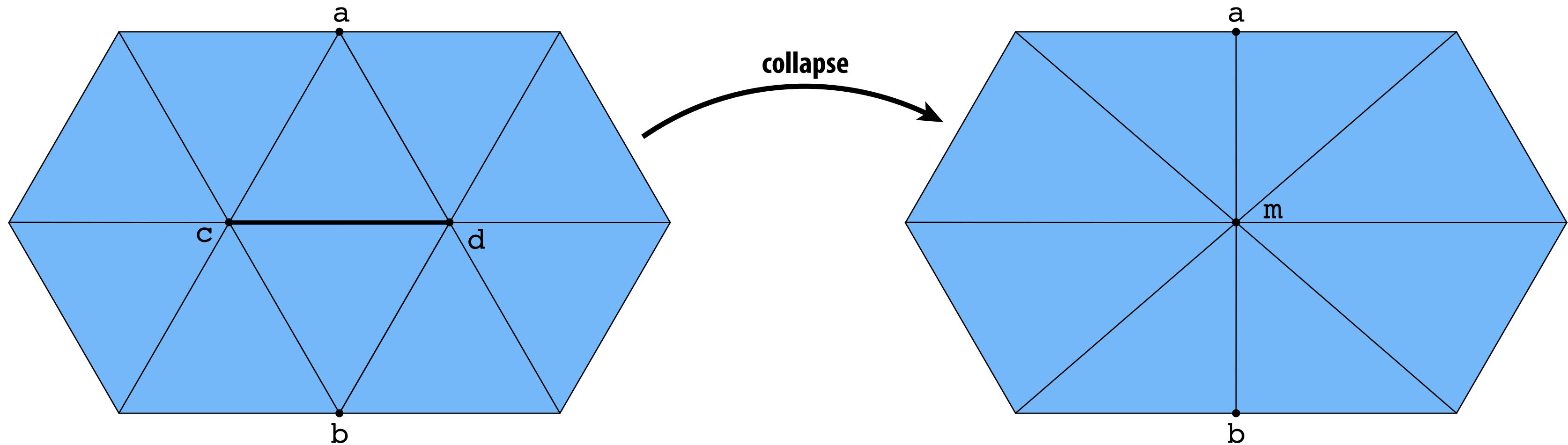
Insert midpoint m of edge (c,b) , connect to get four triangles:



- Must add elements to mesh (new vertex, faces, edges)
- Again, many half-edge pointer reassessments

Half-edge – edge collapse

Replace edge (c,d) with a single vertex m:



- Must delete elements from the mesh
- Again, many half-edge pointer reassessments

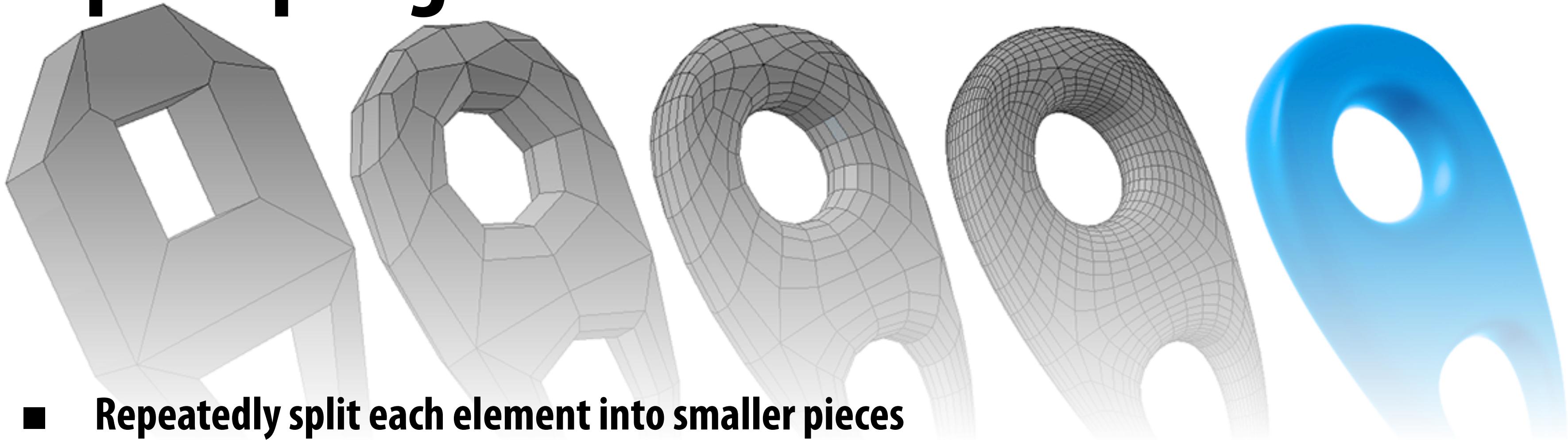
Global mesh operations: geometry processing

- Mesh subdivision (form of subsampling)
- Mesh simplification (form of downsampling)
- Mesh regularization (form of resampling)

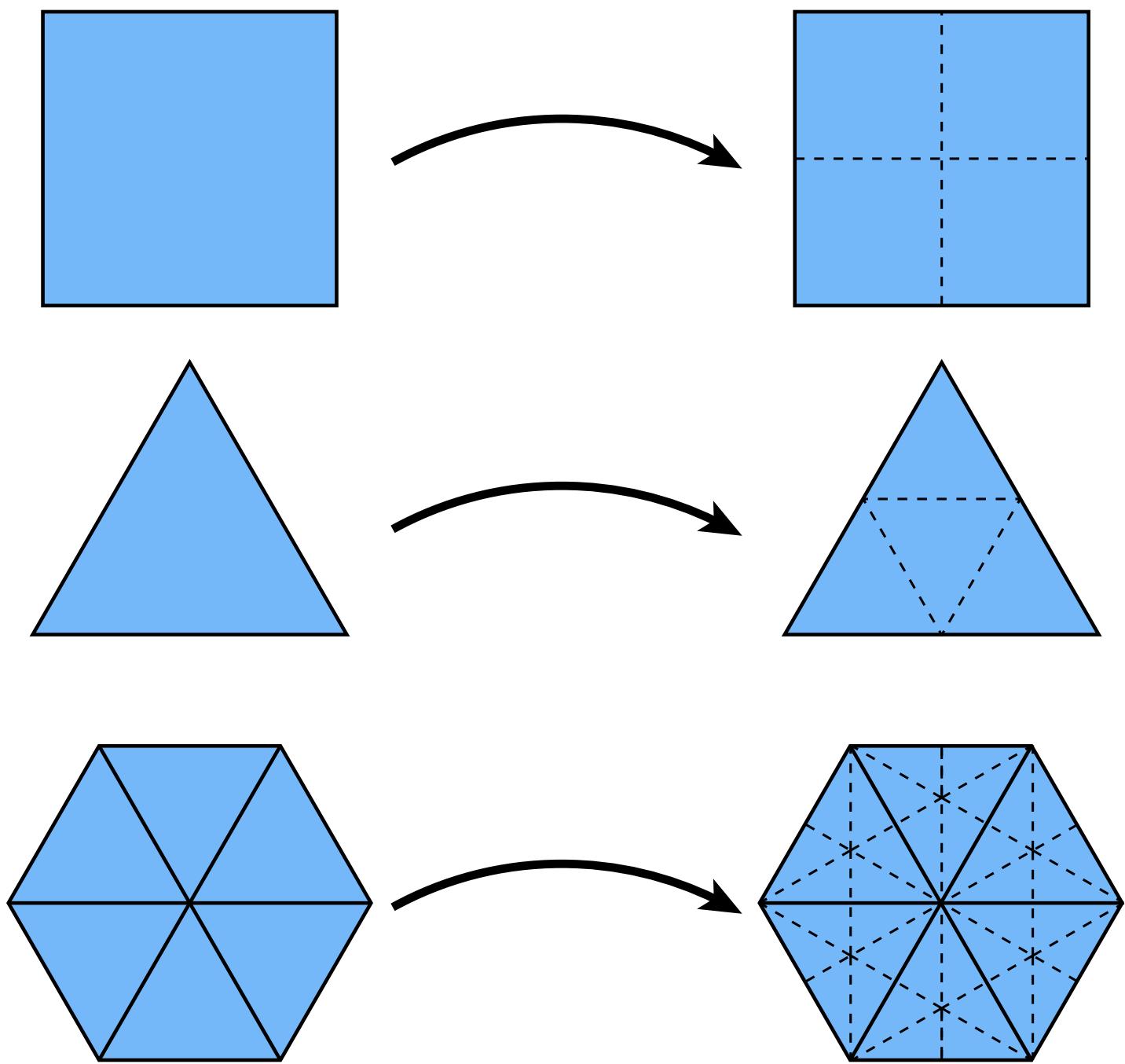


Upsampling a mesh — subdivision

Upsampling via subdivision



- Repeatedly split each element into smaller pieces
- Replace vertex positions with weighted average of neighbors
- Main considerations:
 - interpolating vs. approximating
 - limit surface continuity (C^1, C^2, \dots)
 - behavior at irregular vertices
- Many options:
 - Quad: Catmull-Clark
 - Triangle: Loop, butterfly, $\sqrt{3}$

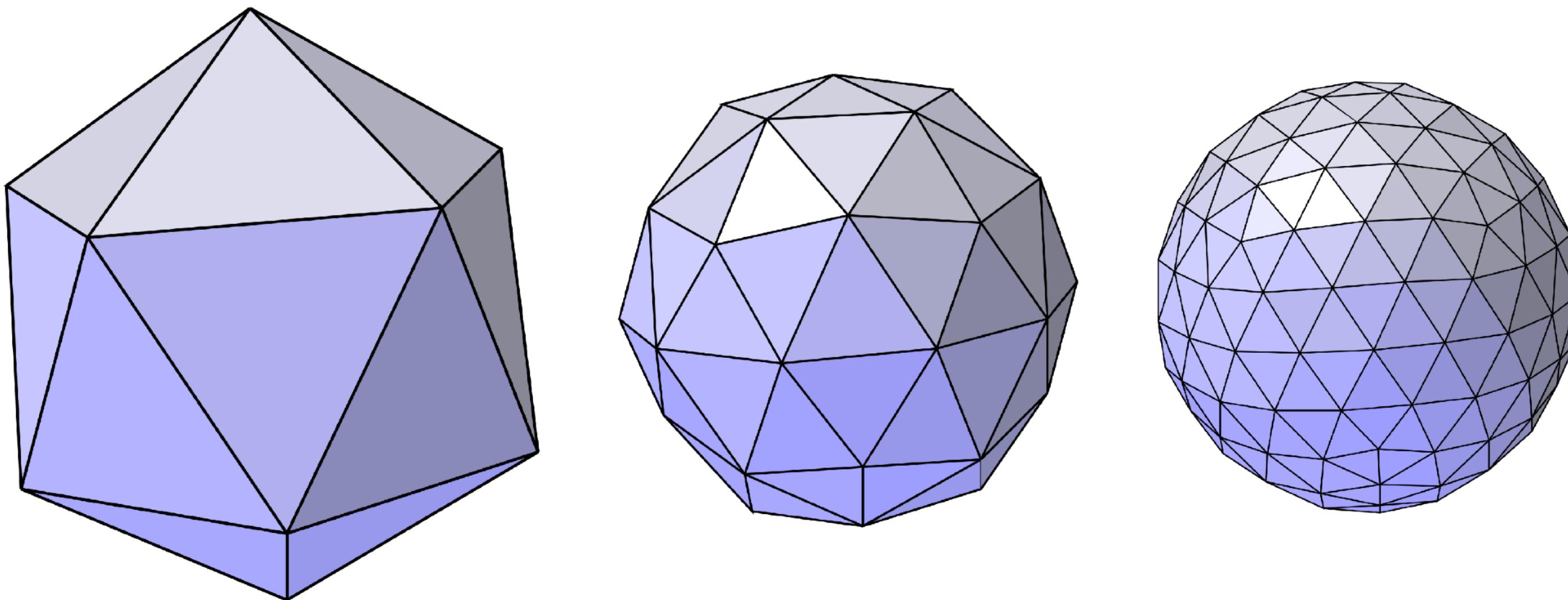


Loop subdivision

Common subdivision rule for triangle meshes

“C2” smoothness away from irregular vertices

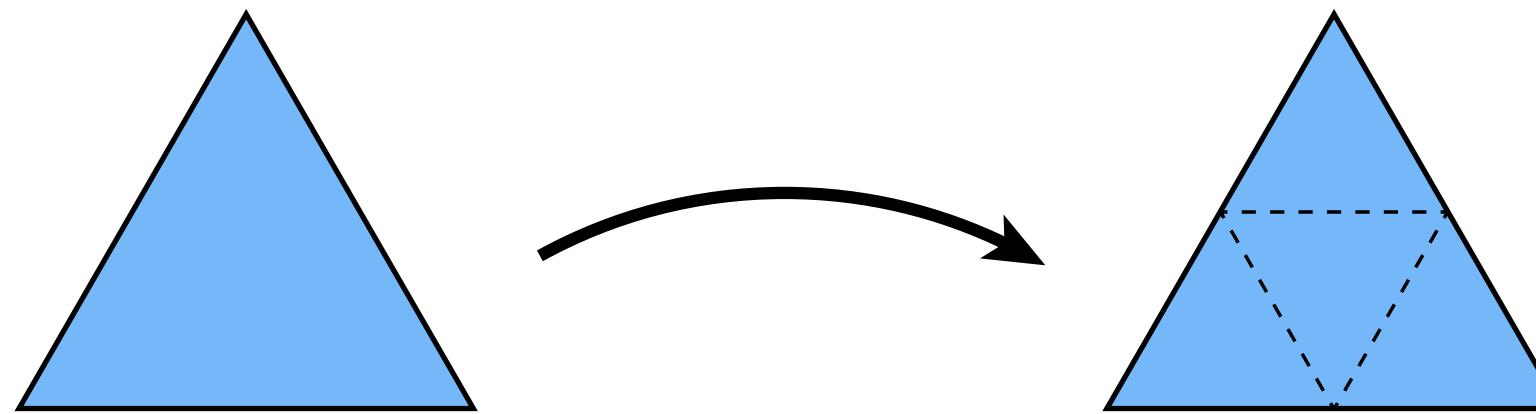
Approximating, not interpolating



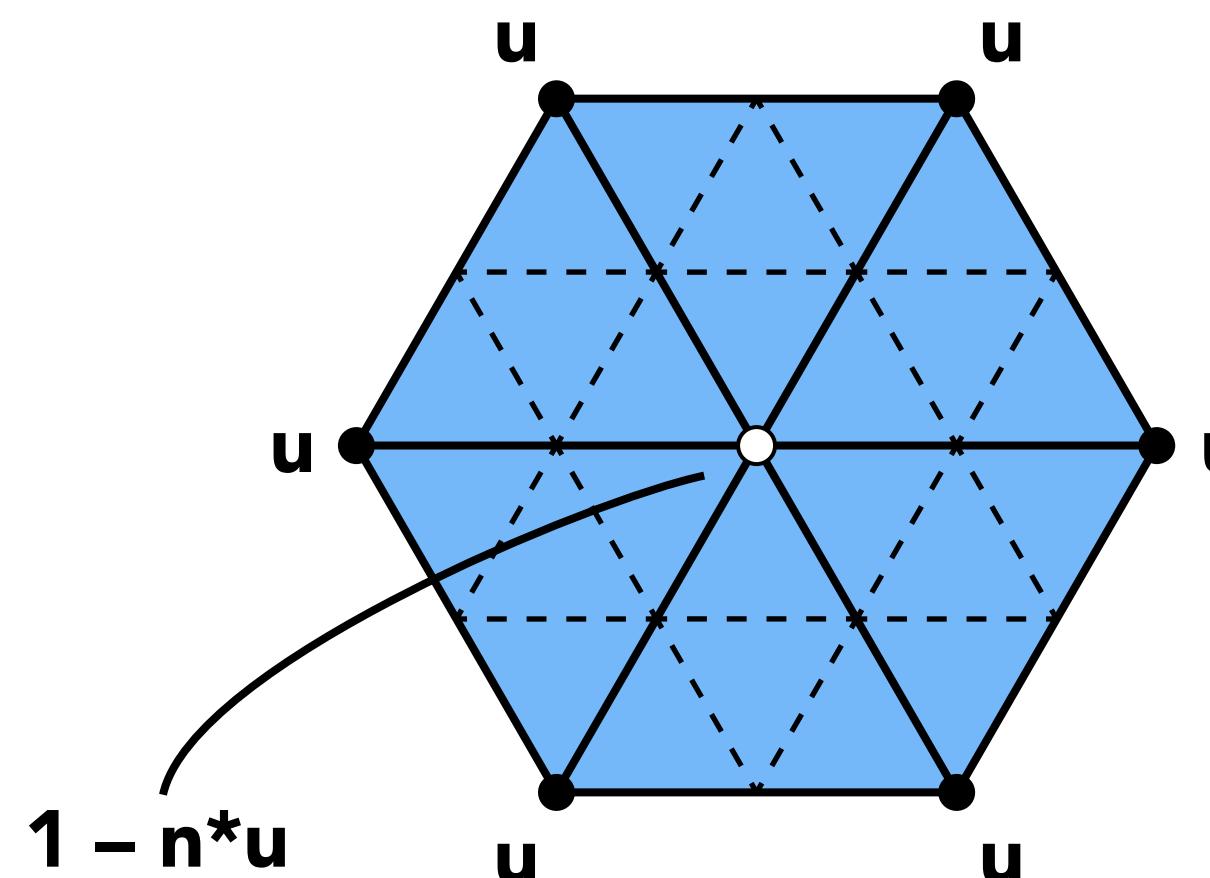
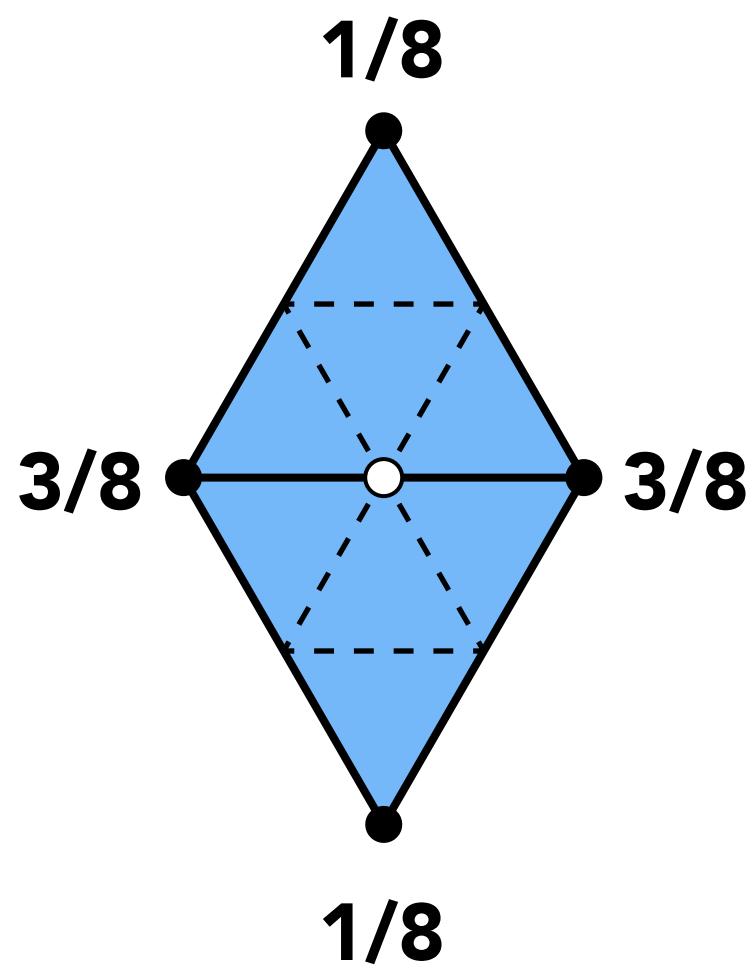
Simon Fuhrman

Loop subdivision algorithm

- Split each triangle into four



- Compute new vertex positions using weighted sum of prior vertex positions:



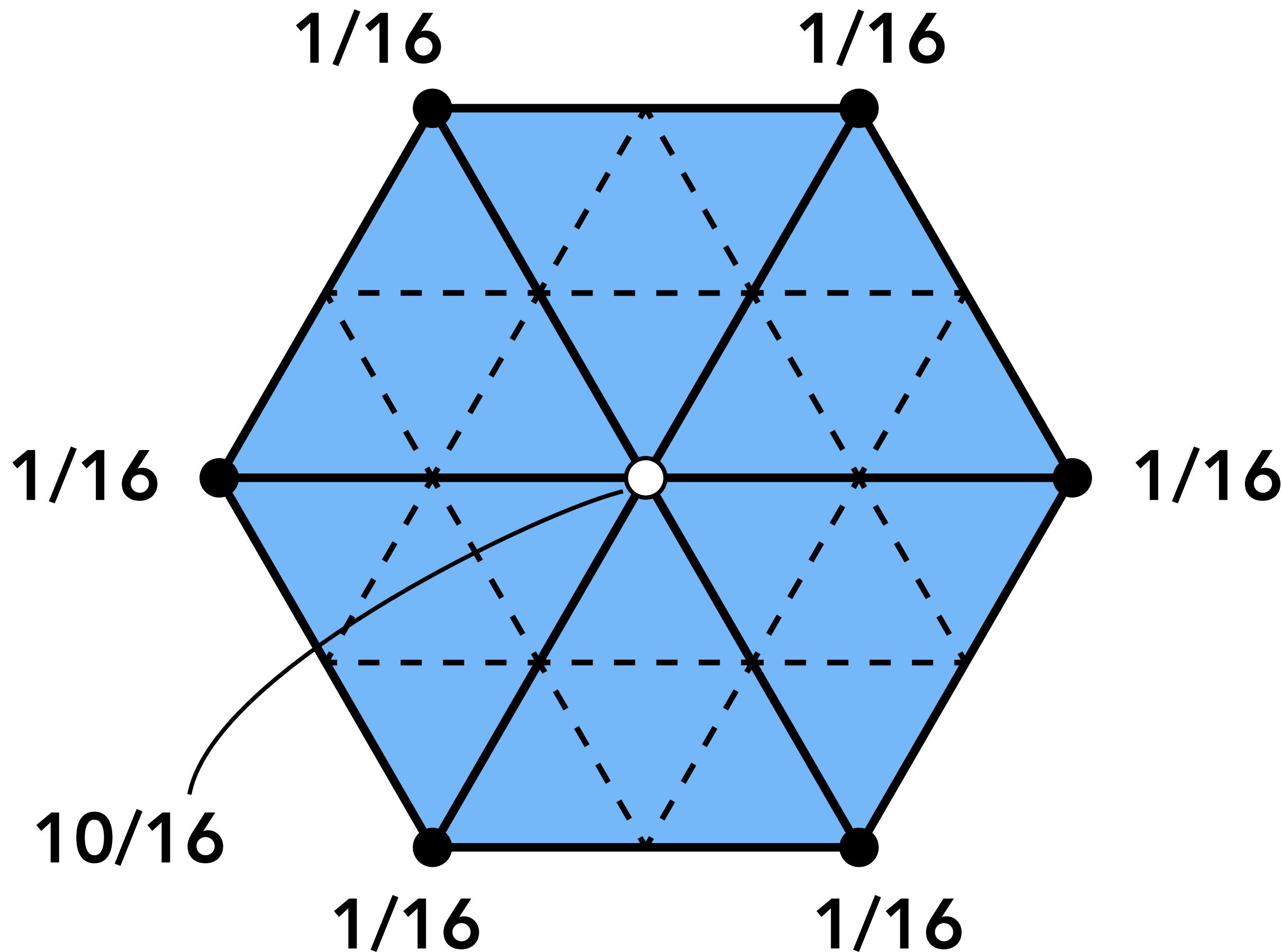
n : vertex degree
 u : $3/16$ if $n=3$, $3/(8n)$ otherwise

New vertices
(weighted sum of vertices on split edge, and vertices "across from" edge)

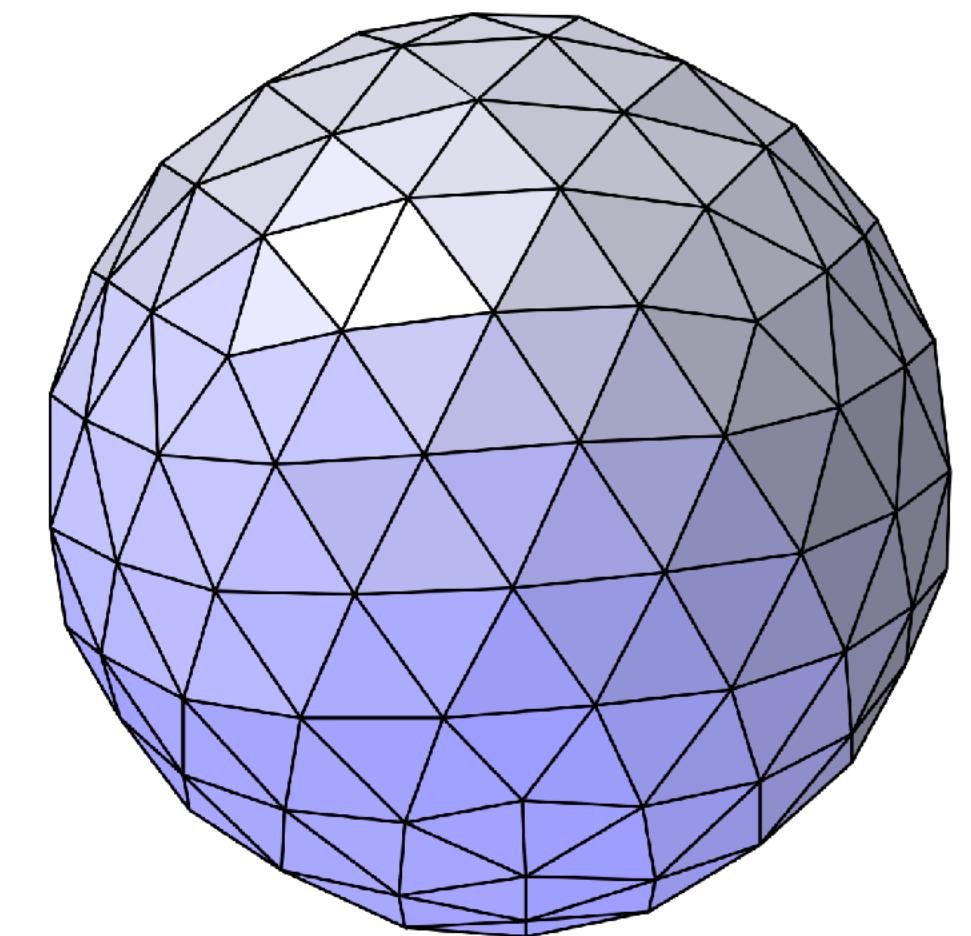
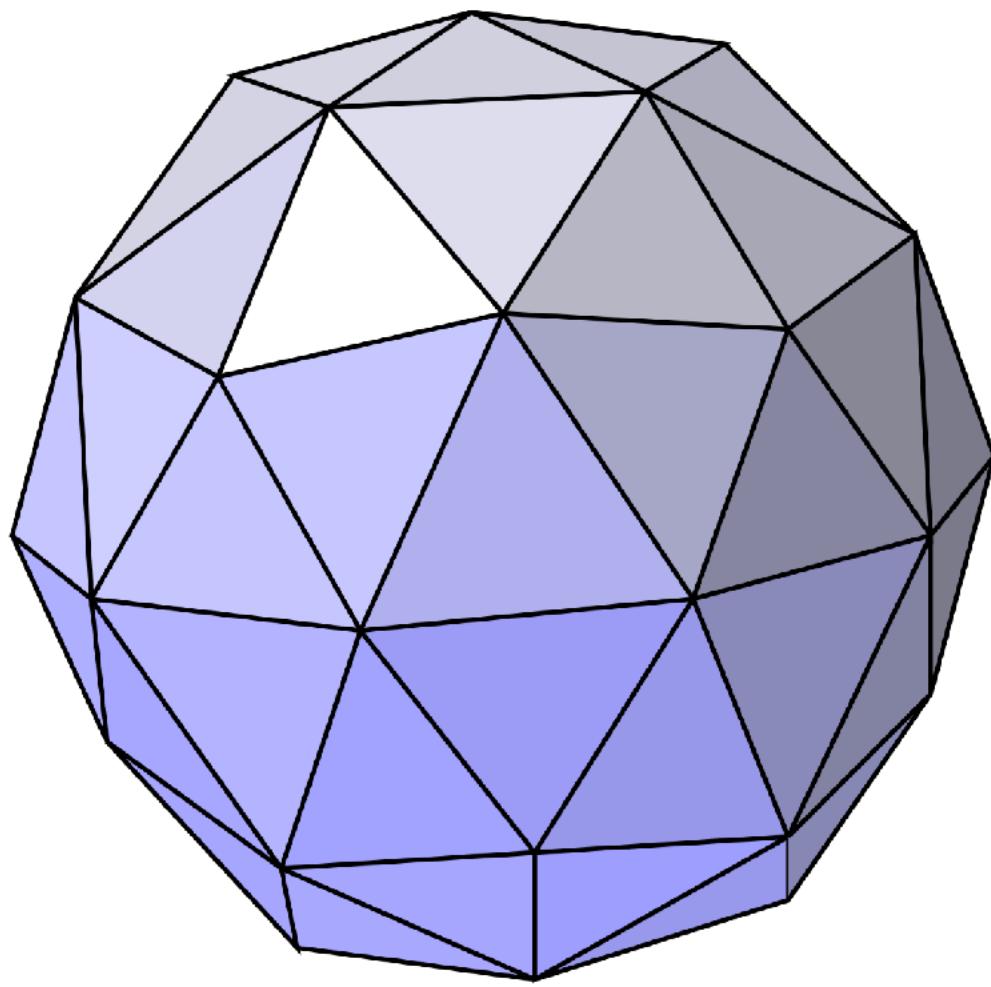
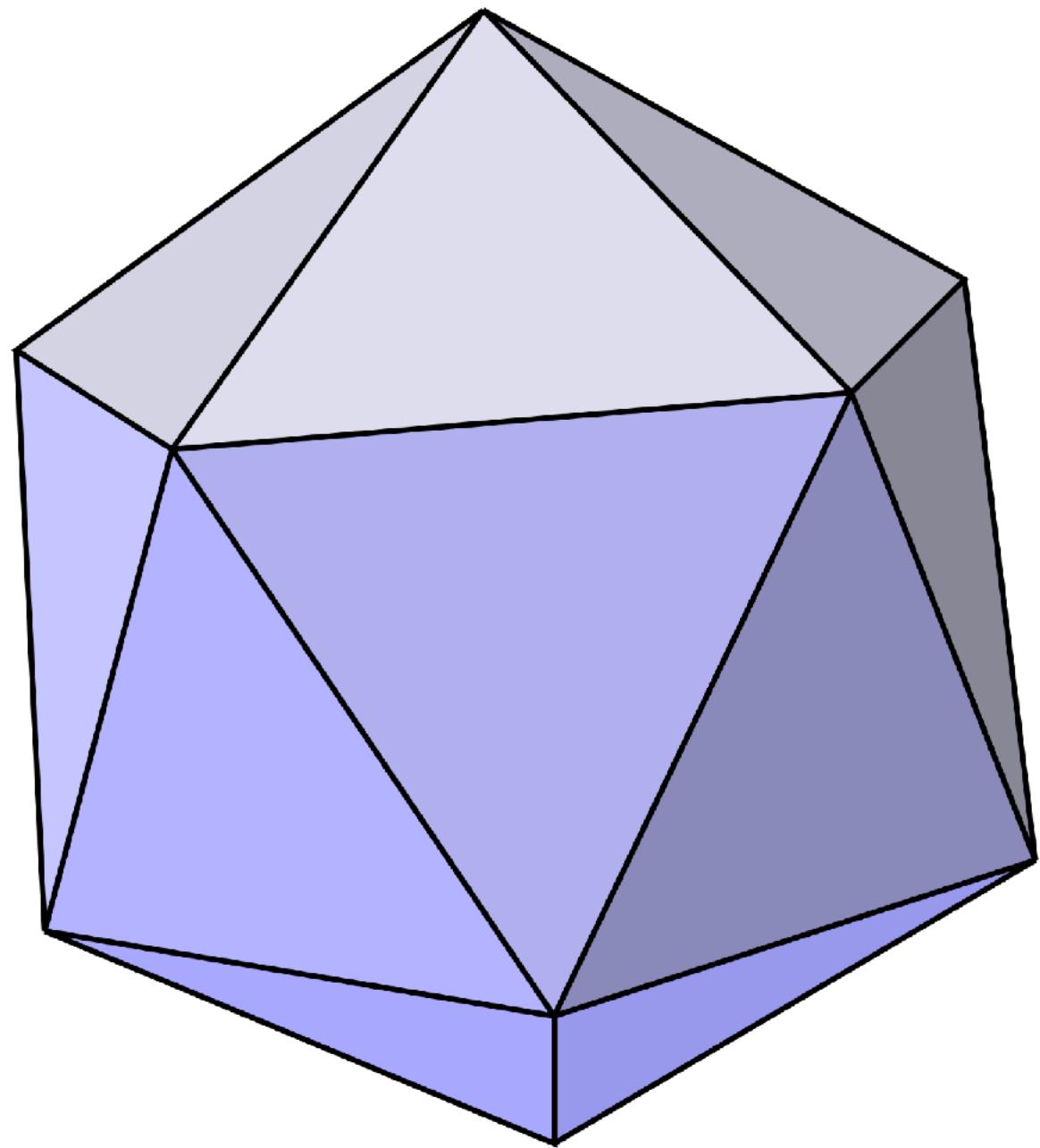
Old vertices
(weighted sum of edge adjacent vertices)

Loop subdivision algorithm

- Example, for degree 6 vertices (“regular” vertices)



Loop subdivision results



Simon Fuhrman

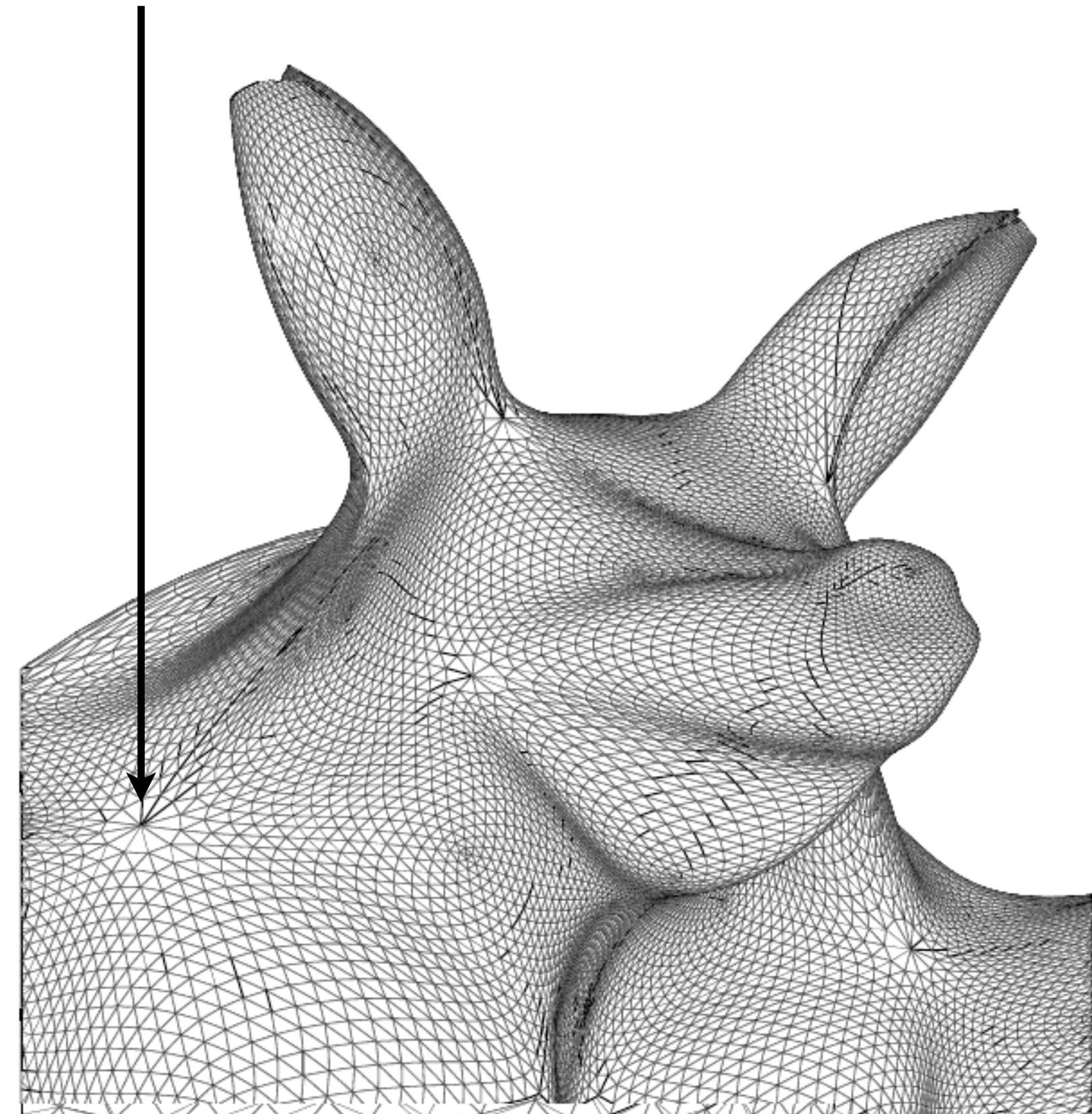
Semi-regular meshes

Most of the mesh has vertices with degree 6

But if the mesh is topologically equivalent to a sphere, then not all the vertices can have degree 6

Must have a few extraordinary points (degree not equal to 6)

Extraordinary vertex



Proof: always an extraordinary vertex

Our triangle mesh (topologically equivalent to sphere) has V vertices, E edges, and T triangles

$$E = \frac{3}{2} T$$

- There are 3 edges per triangle, and each edge is part of 2 triangles
- Therefore $E = \frac{3}{2}T$

$$T = 2V - 4$$

- Euler Convex Polyhedron Formula: $T - E + V = 2$
- $\Rightarrow V = \frac{3}{2}T - T + 2 \Rightarrow T = 2V - 4$

If all vertices had 6 triangles, $T = 2V$

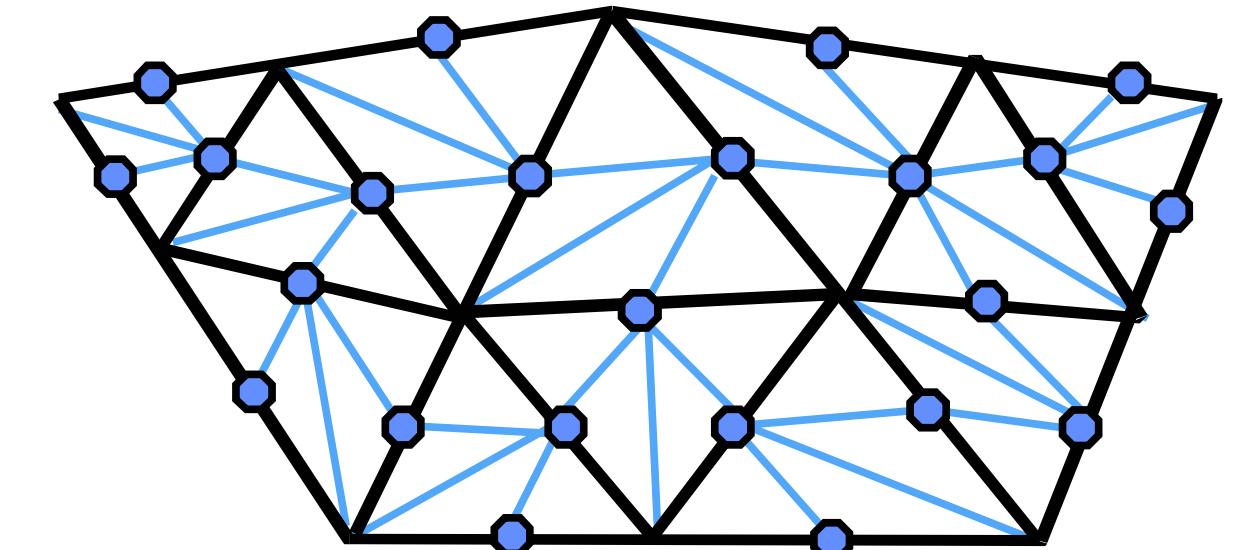
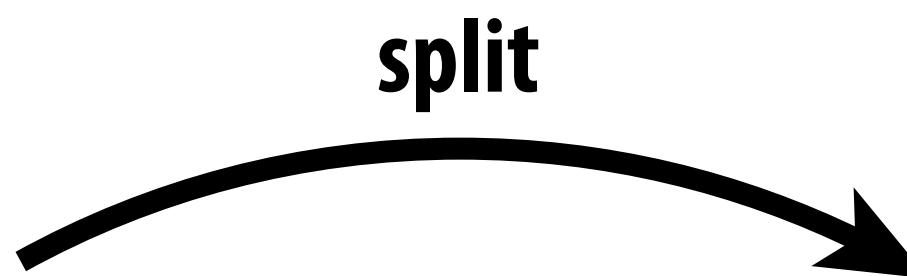
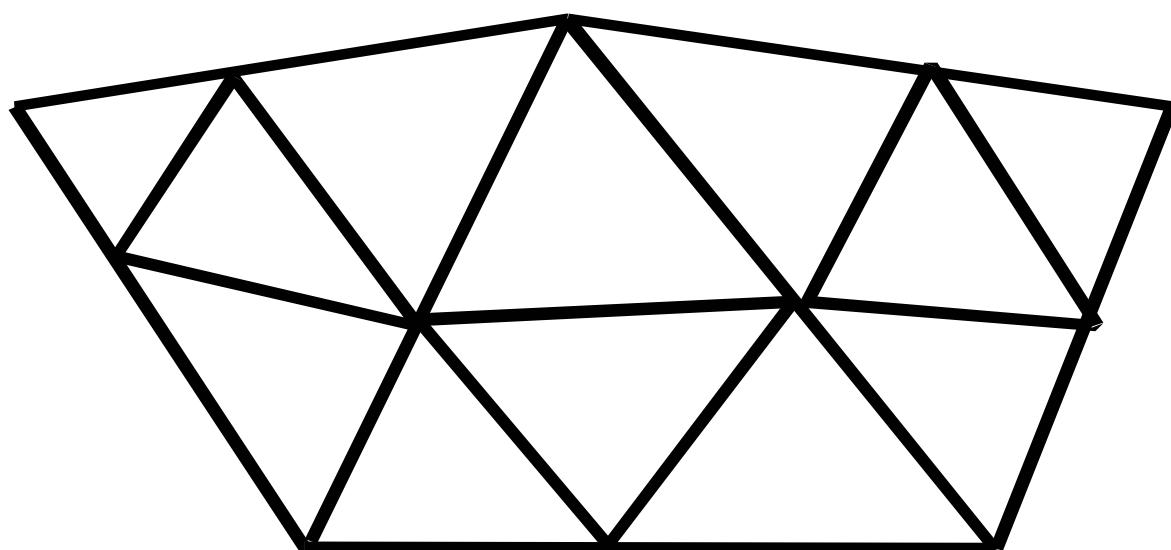
- There are 6 edges per vertex, and every edge connects 2 vertices
- Therefore, $E = 6/2V \Rightarrow 3/2T = 6/2V \Rightarrow T = 2V$

T cannot equal both $2V - 4$ and $2V$, a contradiction

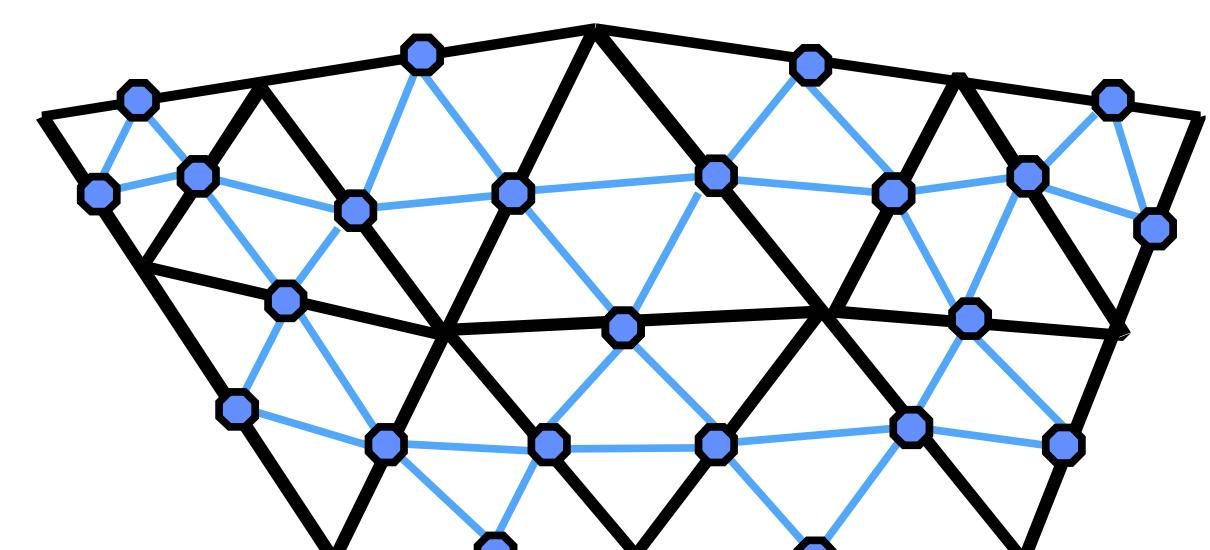
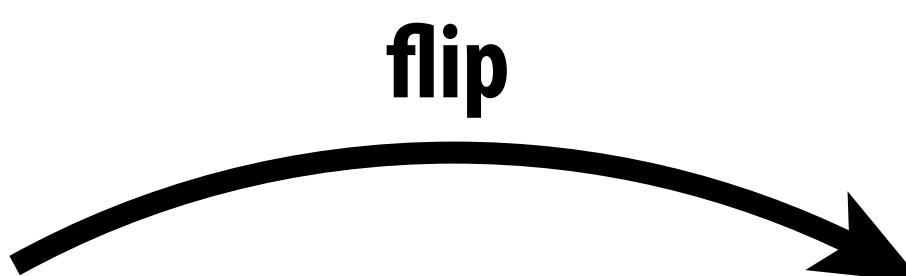
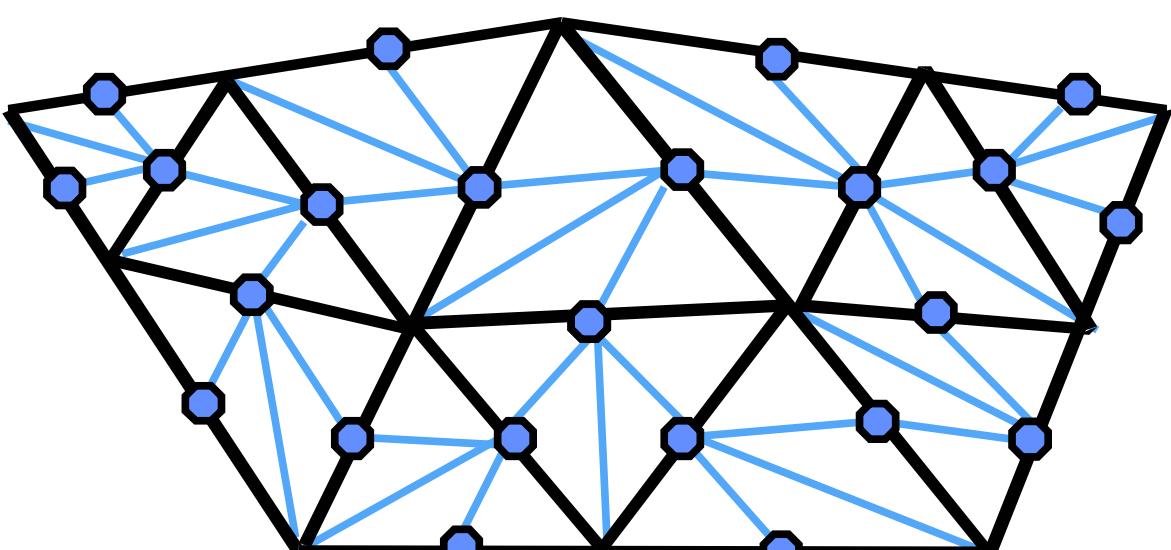
- Therefore, the mesh cannot have 6 triangles for every vertex

Loop subdivision via edge operations

First, split edges of original mesh in any order:



Next, flip new edges that touch a new and old vertex:

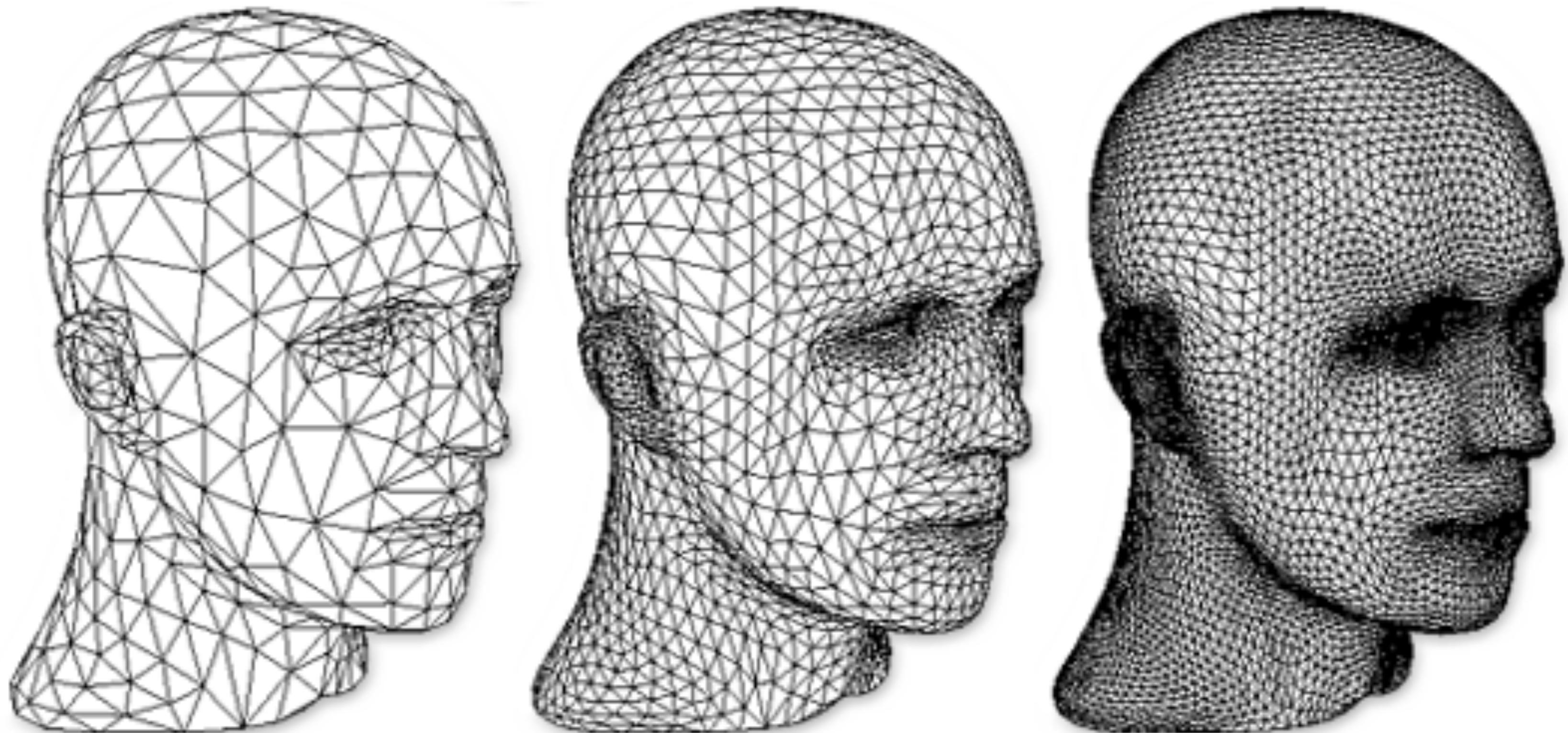


(Don't forget to update vertex positions!)

Continuity of loop subdivision surface

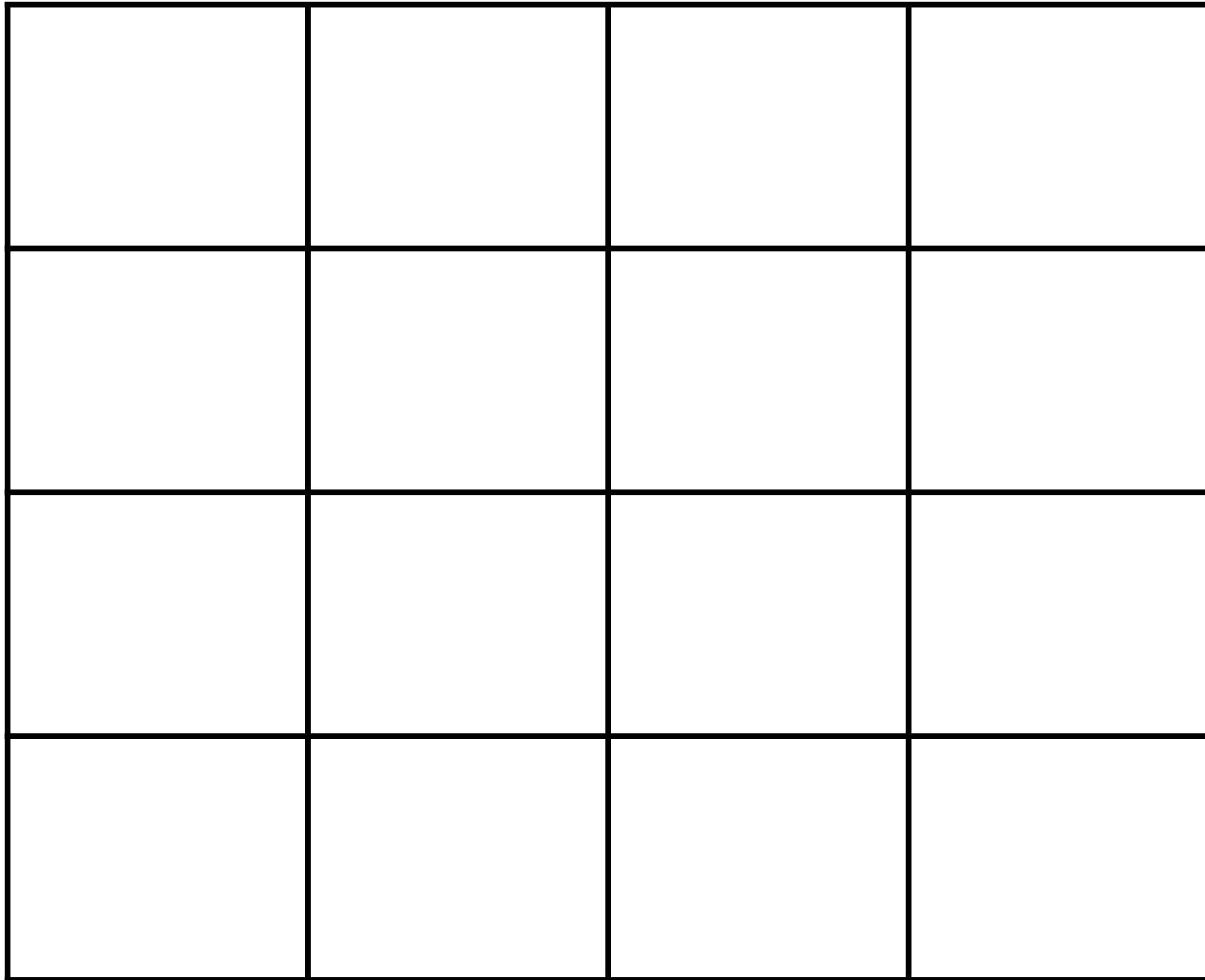
- At extraordinary vertices
 - Surface is at least C^1 continuous
- Everywhere else (“ordinary” regions)
 - Surface is C^2 continuous

Loop subdivision results

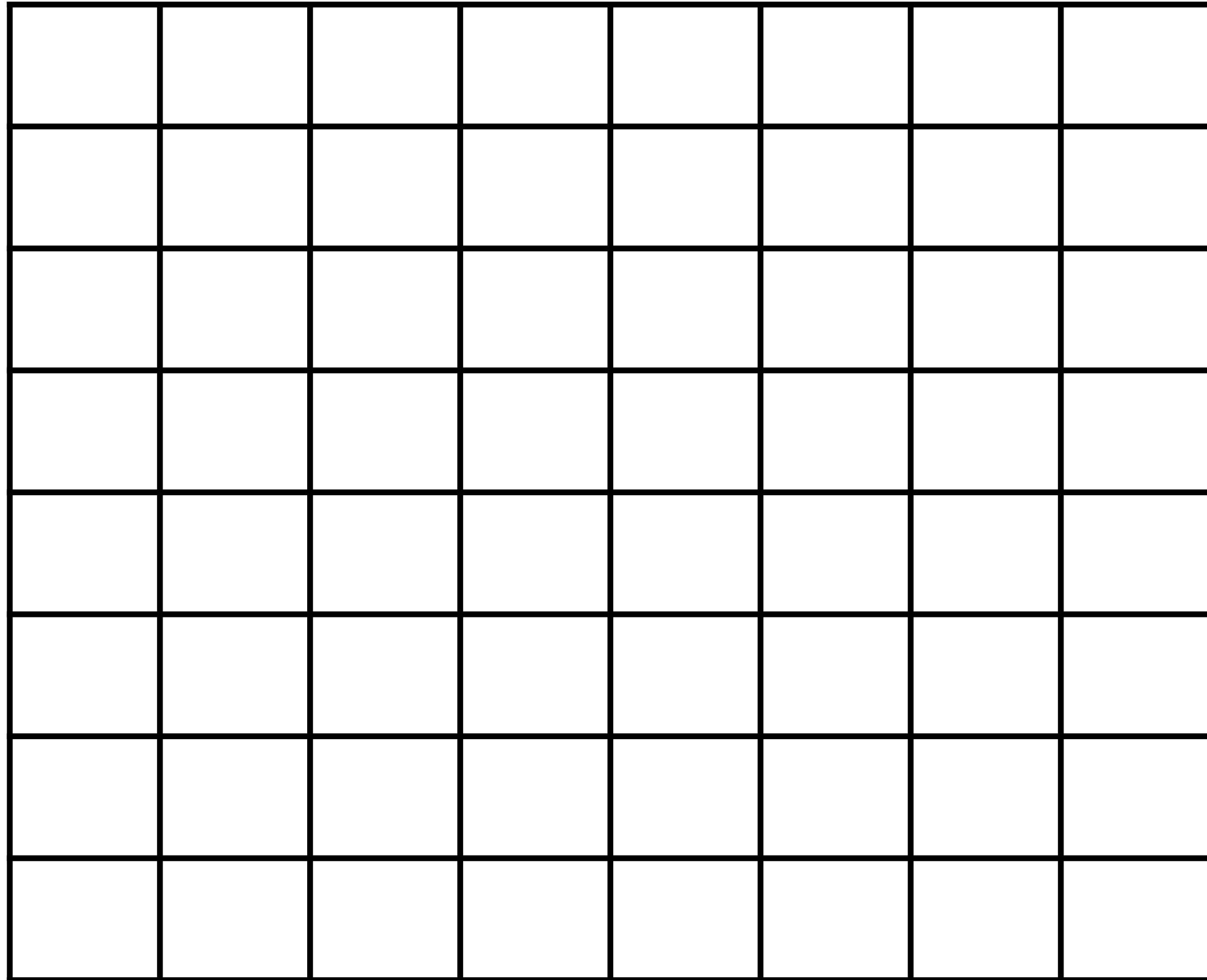


Catmull-Clark subdivision

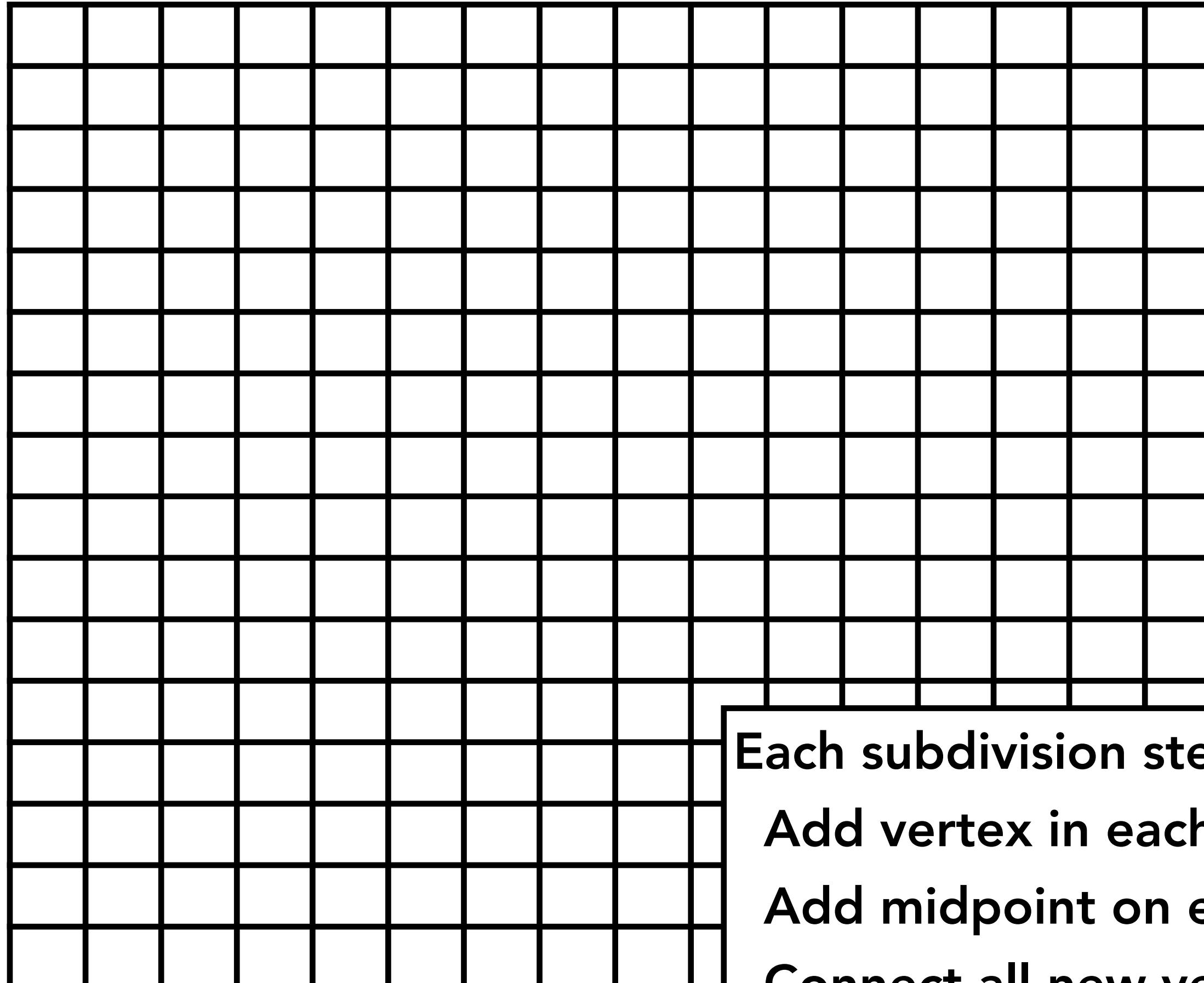
Catmull-Clark subdivision (regular quad mesh)



Catmull-Clark subdivision (regular quad mesh)

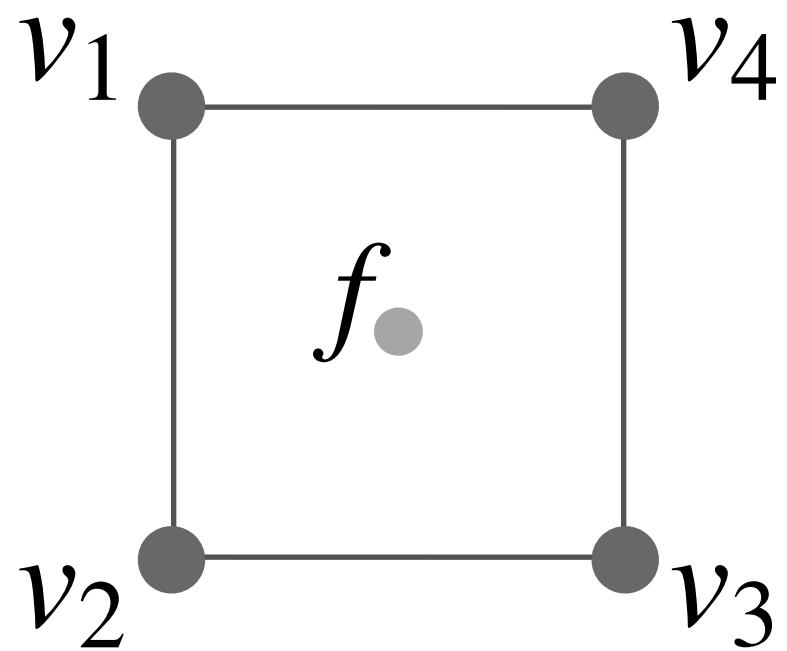


Catmull-Clark subdivision (regular quad mesh)



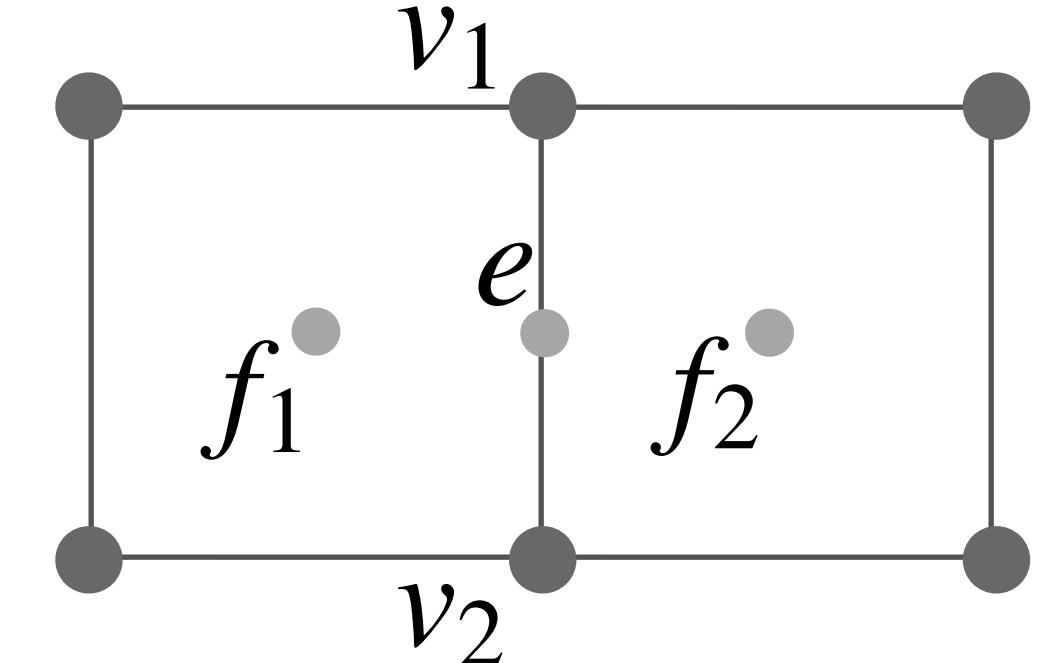
Catmull-Clark vertex update rules (quad mesh)

Face point



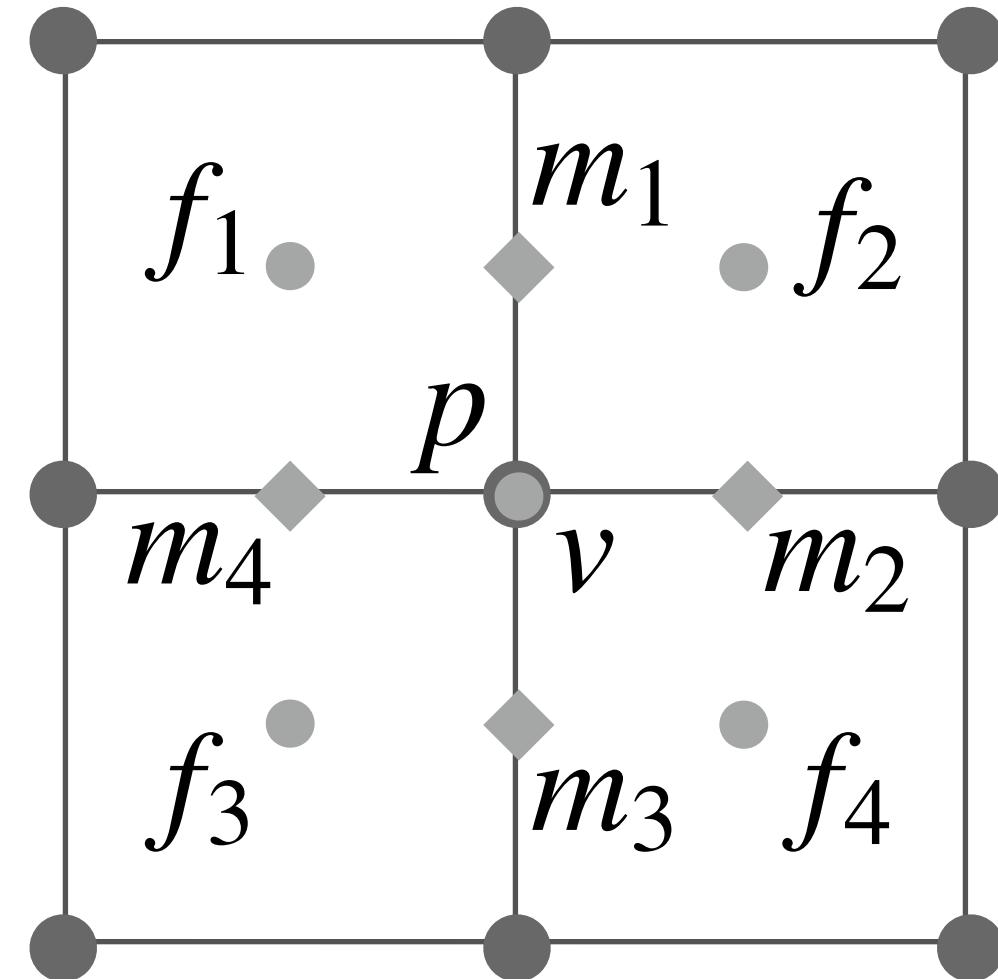
$$f = \frac{v_1 + v_2 + v_3 + v_4}{4}$$

Edge point



$$e = \frac{v_1 + v_2 + f_1 + f_2}{4}$$

Vertex point



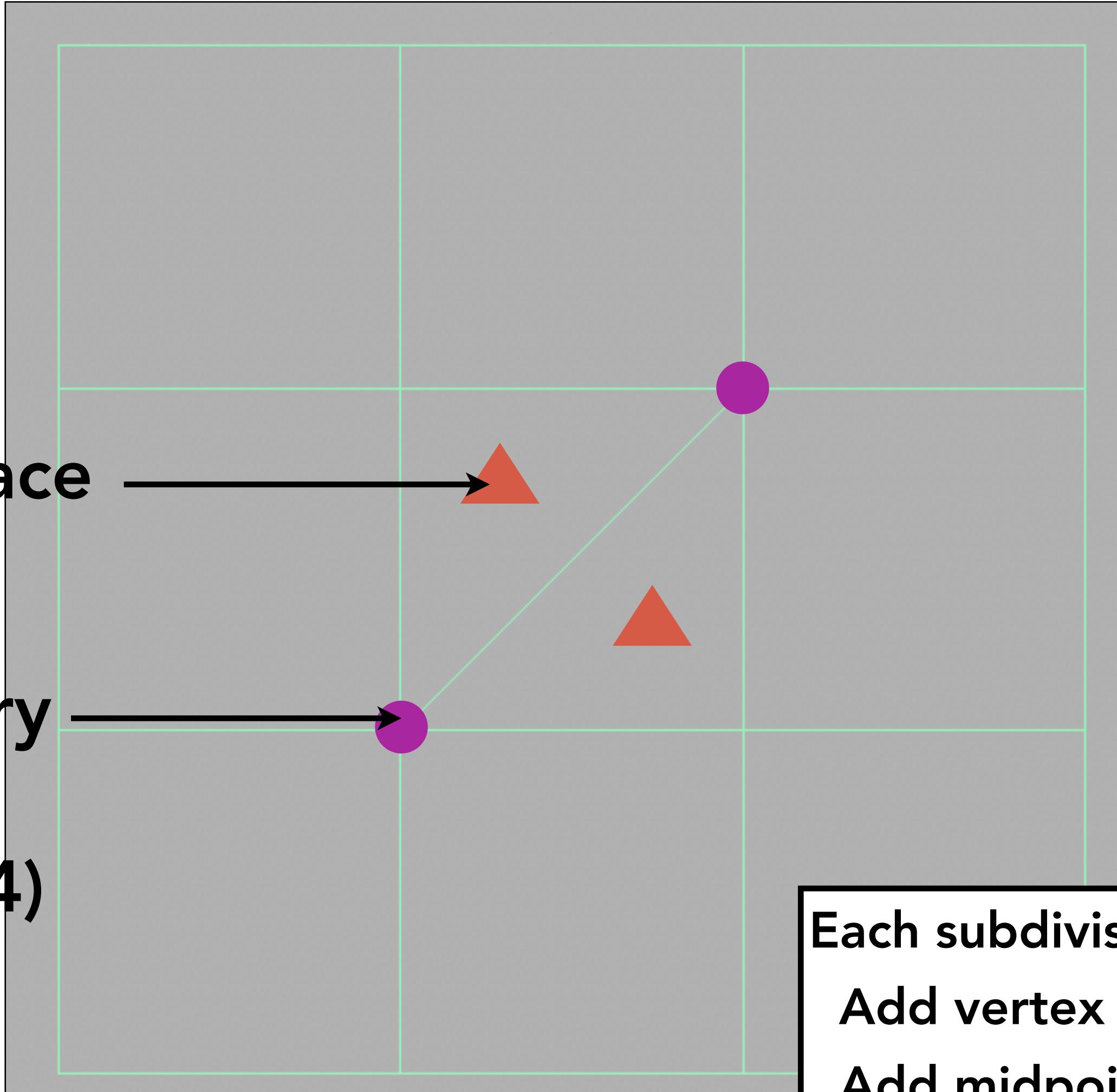
$$v = \frac{f_1 + f_2 + f_3 + f_4 + 2(m_1 + m_2 + m_3 + m_4) + 4p}{16}$$

m midpoint of edge, not "edge point"

p old "vertex point"

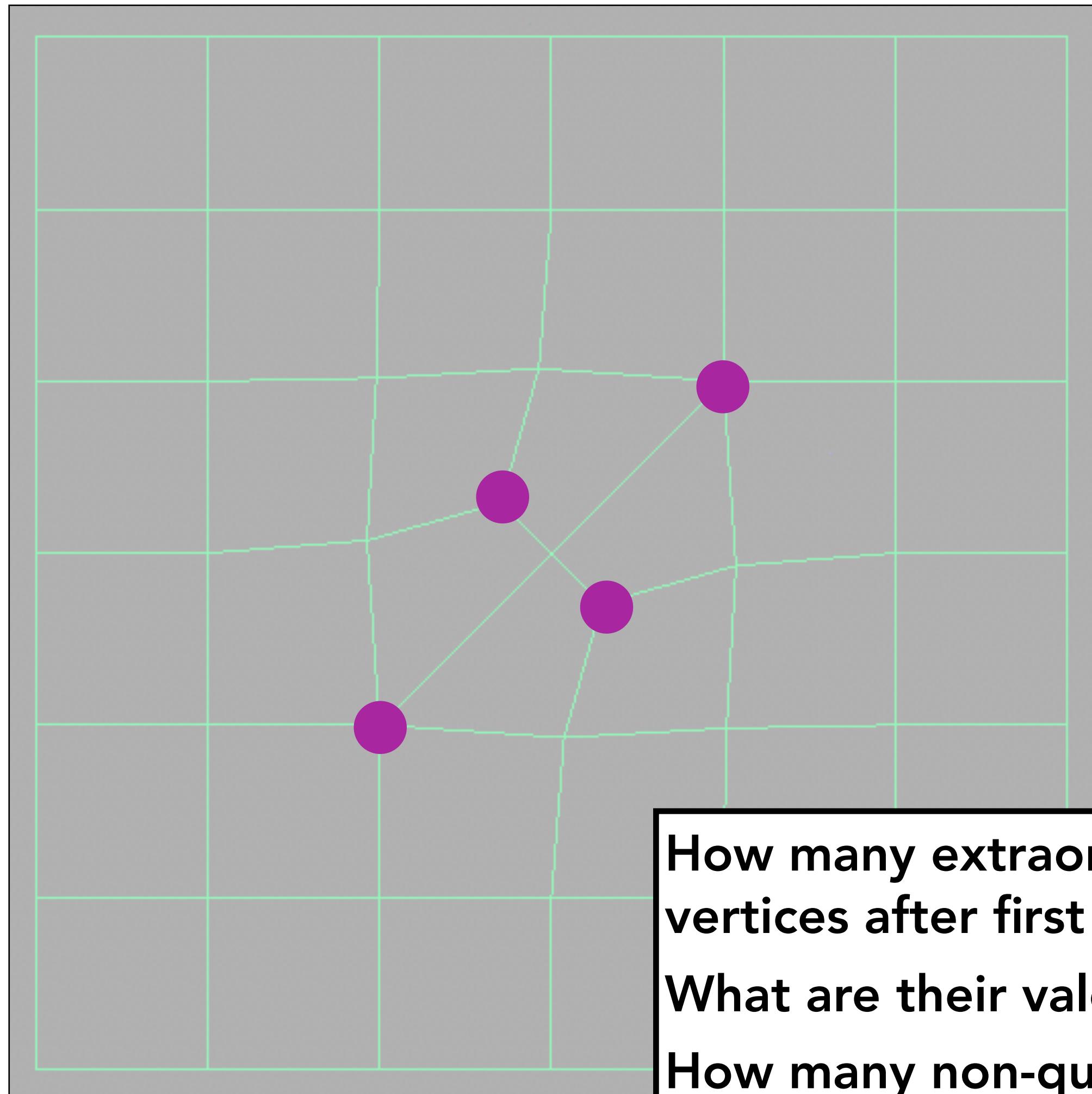
Catmull-Clark subdivision (general mesh)

Non-quad face
Extraordinary vertex
(valence $\neq 4$)

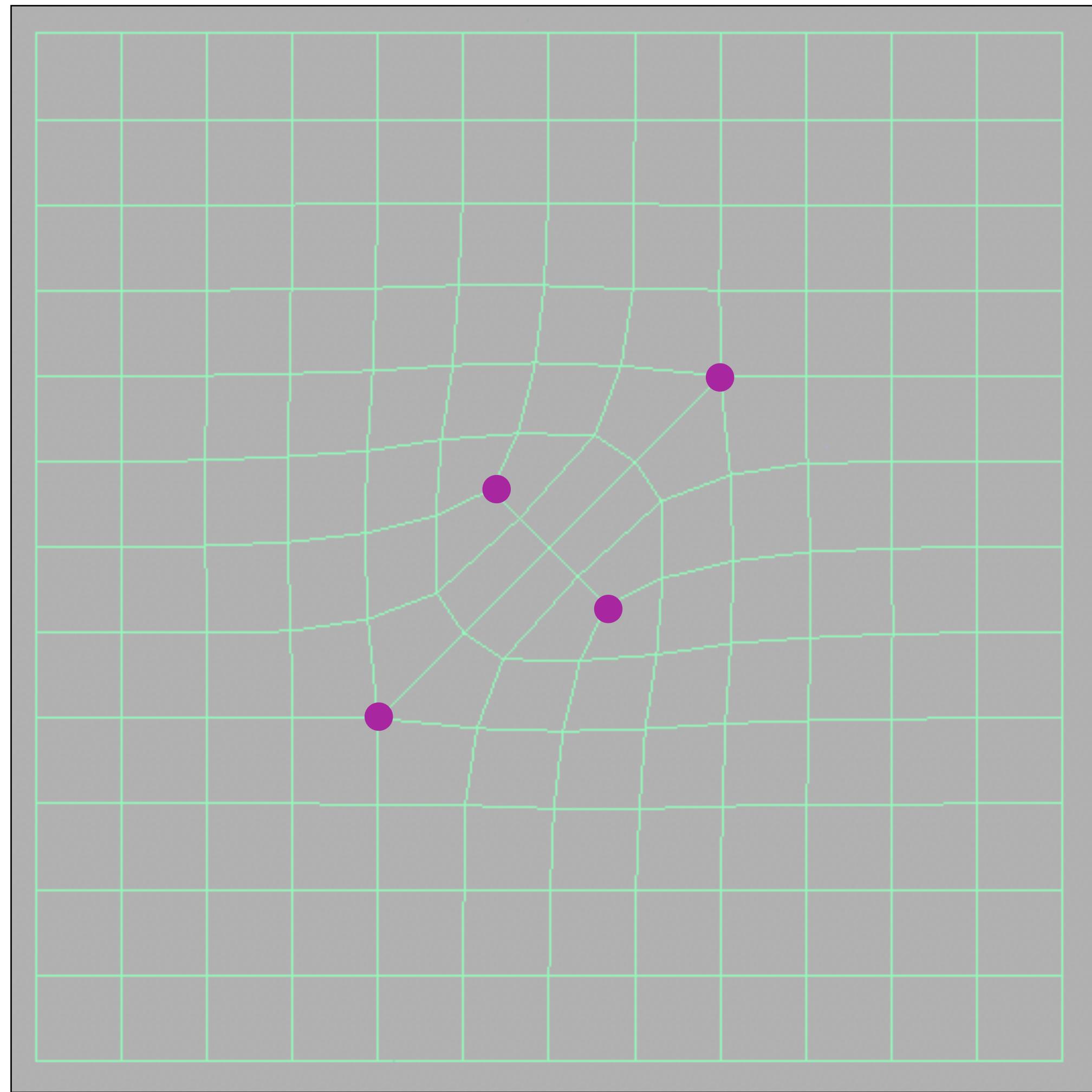


Each subdivision step:
Add vertex in each face
Add midpoint on each edge
Connect all new vertices

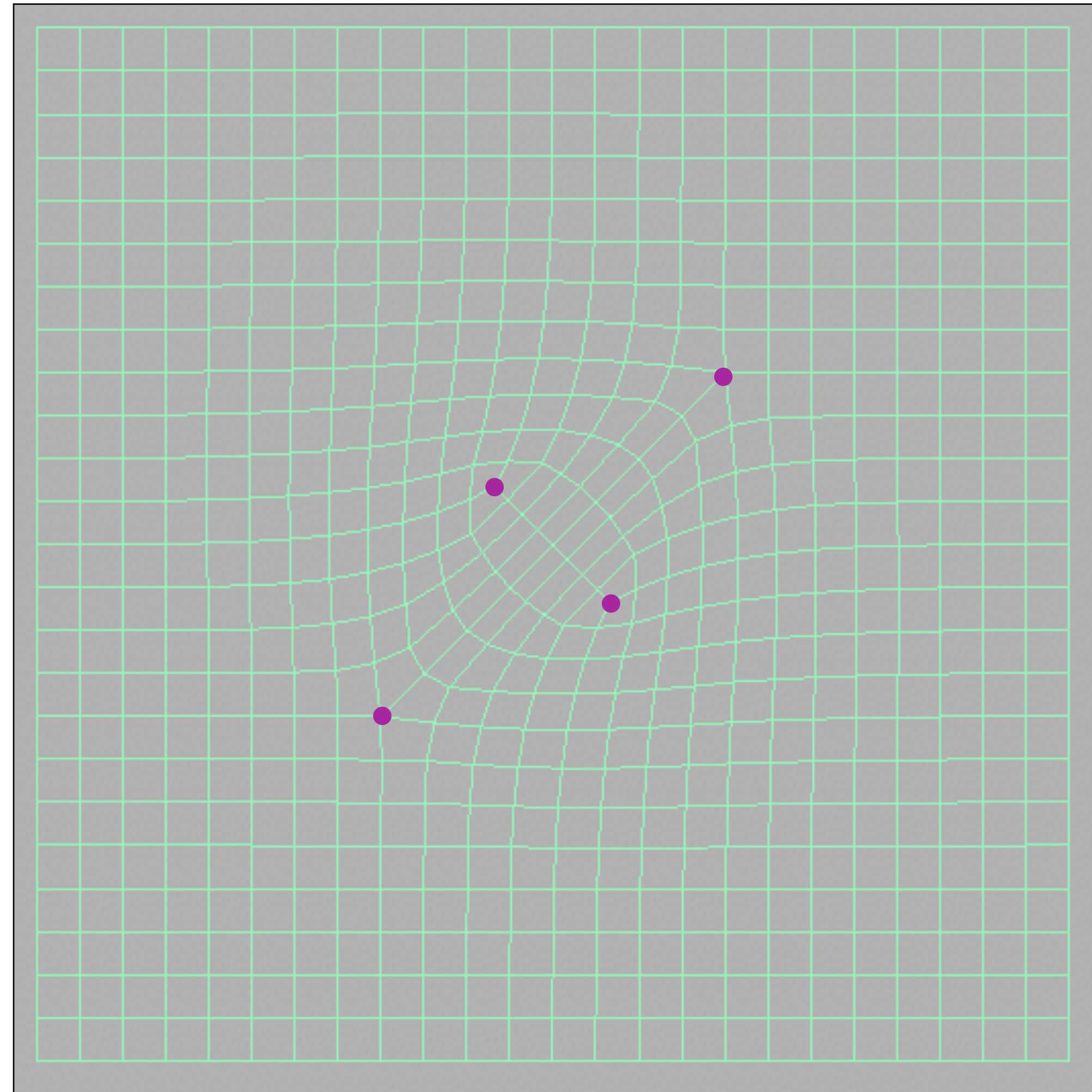
Catmull-Clark subdivision (general mesh)



Catmull-Clark subdivision (general mesh)



Catmull-Clark subdivision (general mesh)



Catmull-Clark vertex update rules (general mesh)

f = average of surrounding vertices

$$e = \frac{f_1 + f_2 + v_1 + v_2}{4}$$

**These rules reduce to earlier quad
rules for ordinary vertices / faces**

$$v = \frac{\bar{f}}{n} + \frac{2\bar{m}}{n} + \frac{p(n-3)}{n}$$

\bar{m} = average of adjacent midpoints

\bar{f} = average of adjacent face points

n = valence of vertex

p = old "vertex" point

Continuity of Catmull-Clark surface

- At extraordinary points
 - Surface is at least C^1 continuous
- Everywhere else (“ordinary” regions)
 - Surface is C^2 continuous

What about sharp creases?

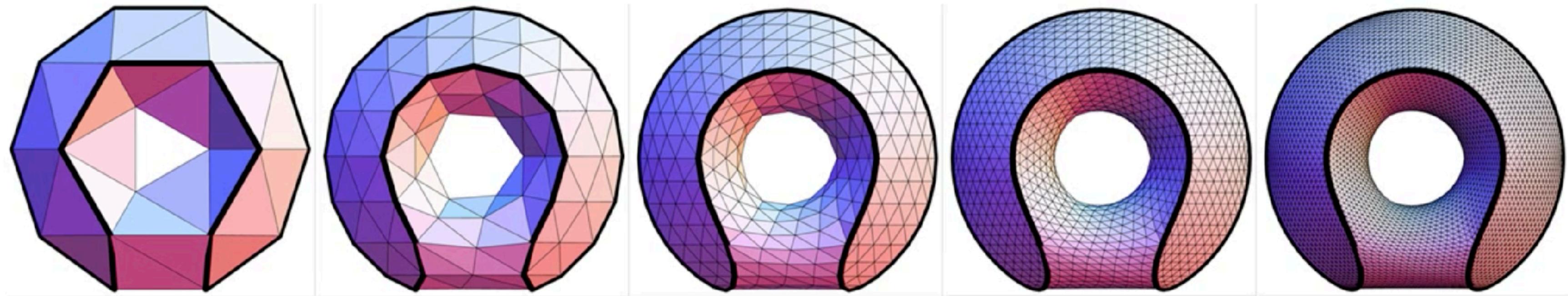


From Pixar Short, "Geri's Game"

Hand is modeled as a Catmull Clark surface with creases between skin and fingernail

What about sharp creases?

Loop with Sharp Creases



Catmull-Clark with Sharp Creases

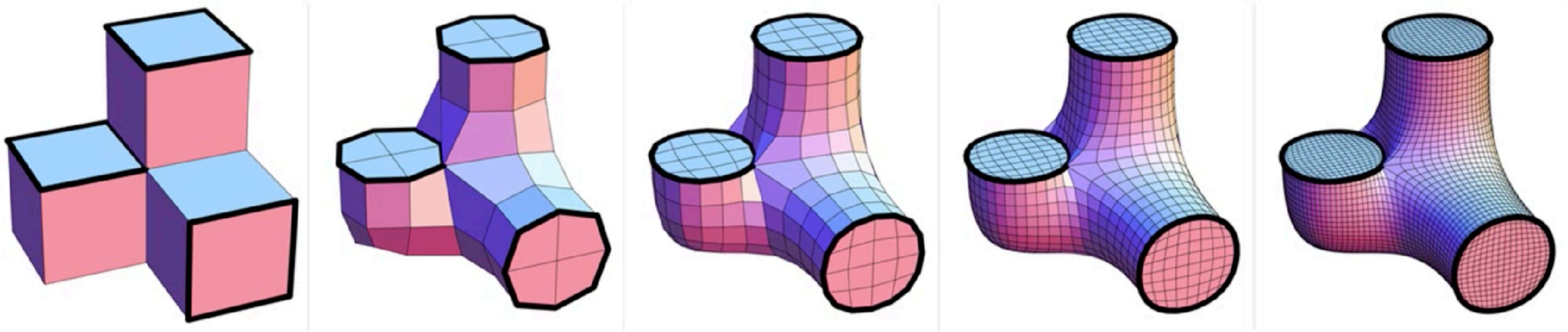
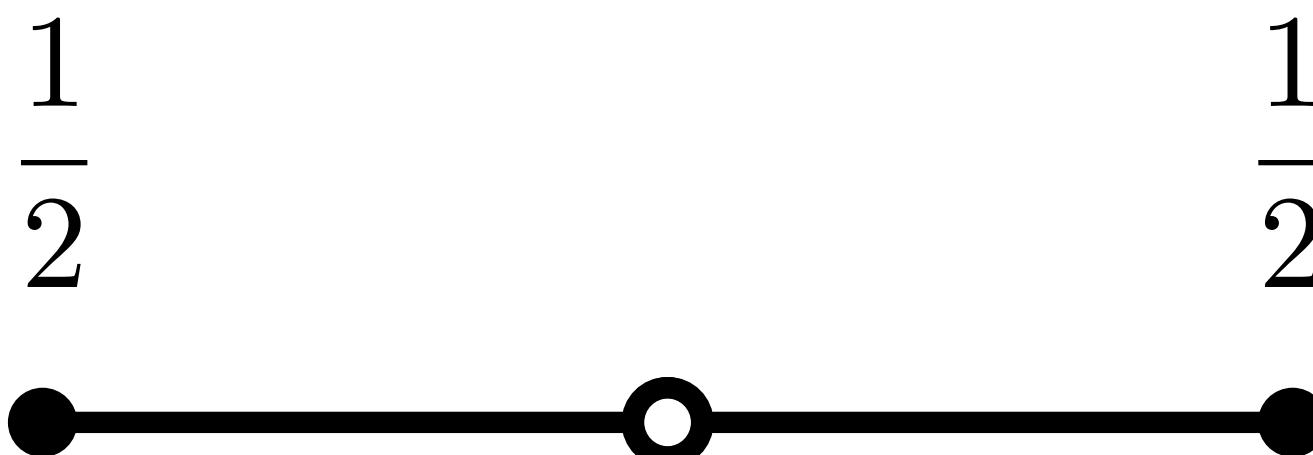


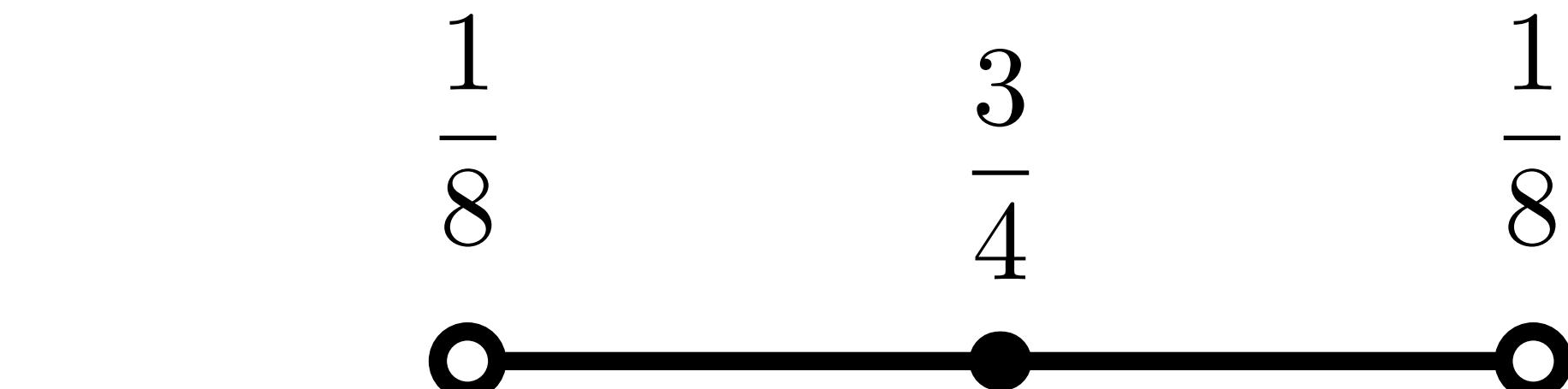
Figure from: Hakenberg et al. Volume Enclosed by Subdivision Surfaces with Sharp Creases

Creases and boundaries

- Can create creases in subdivision surfaces by marking certain edges as “sharp”. Surface boundary edges can be handled the same way
 - Use different subdivision rules for vertices along these “sharp” edges



Insert new midpoint vertex,
weights as shown



Update existing vertices,
weights as shown

Subdivision in action (“Geri’s Game”, Pixar)

- Subdivision used for entire character:

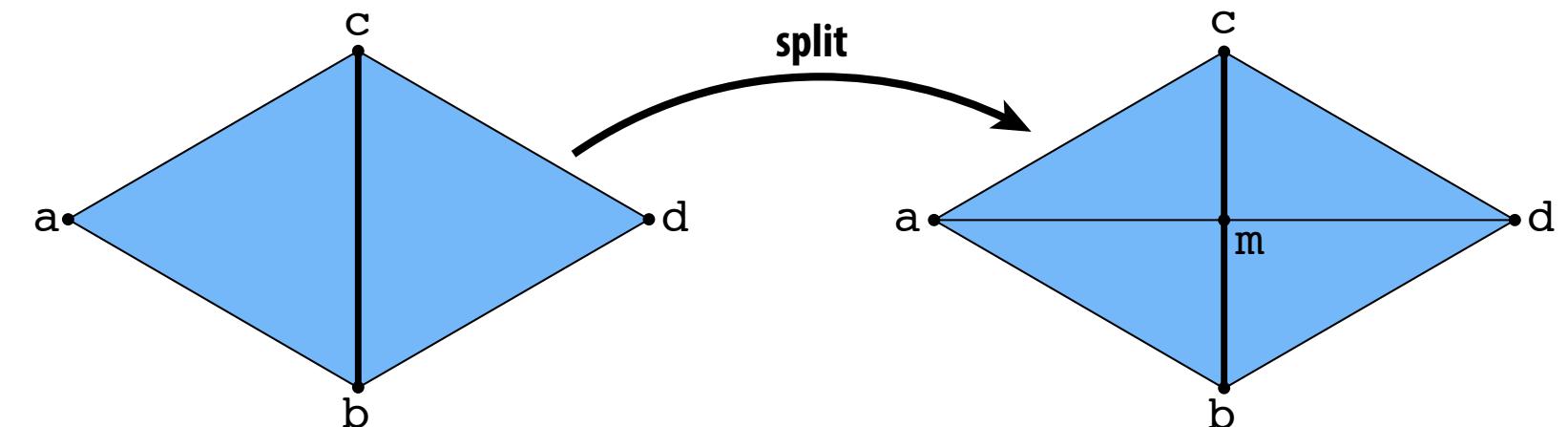
- Hands and head
- Clothing, tie, shoes



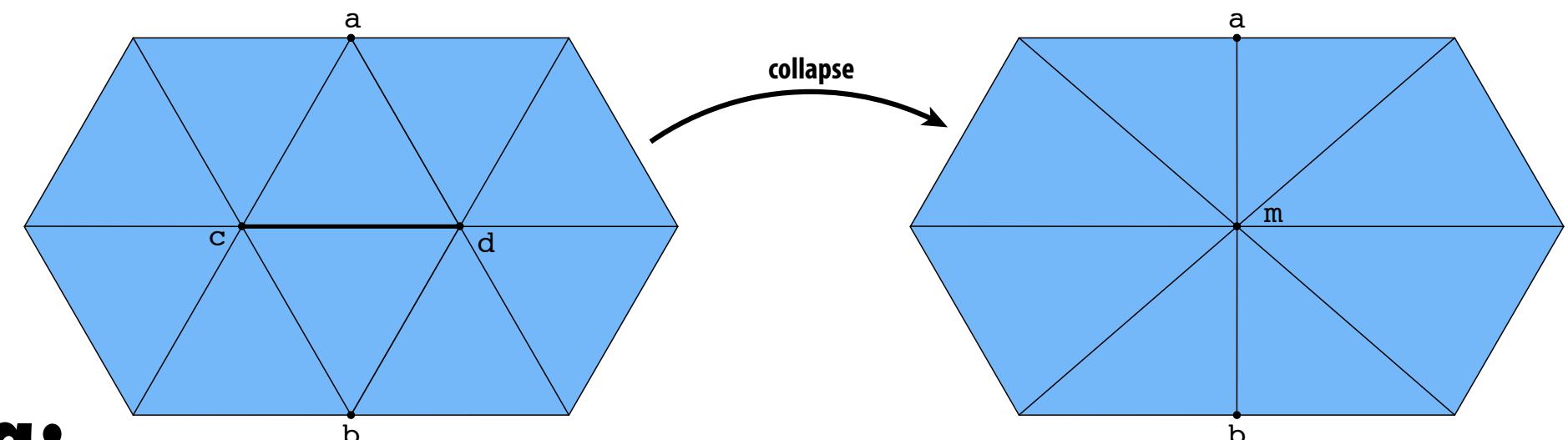
Mesh simplification — downsampling

How do we resample meshes? (reminder)

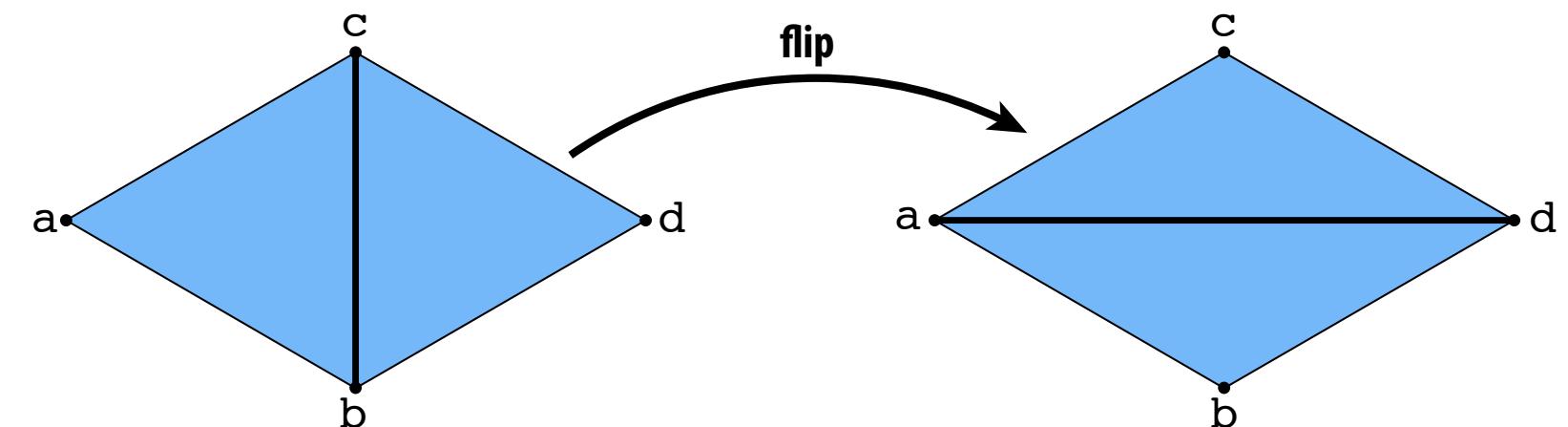
- Edge split is (local) upsampling:



- Edge collapse is (local) downsampling:



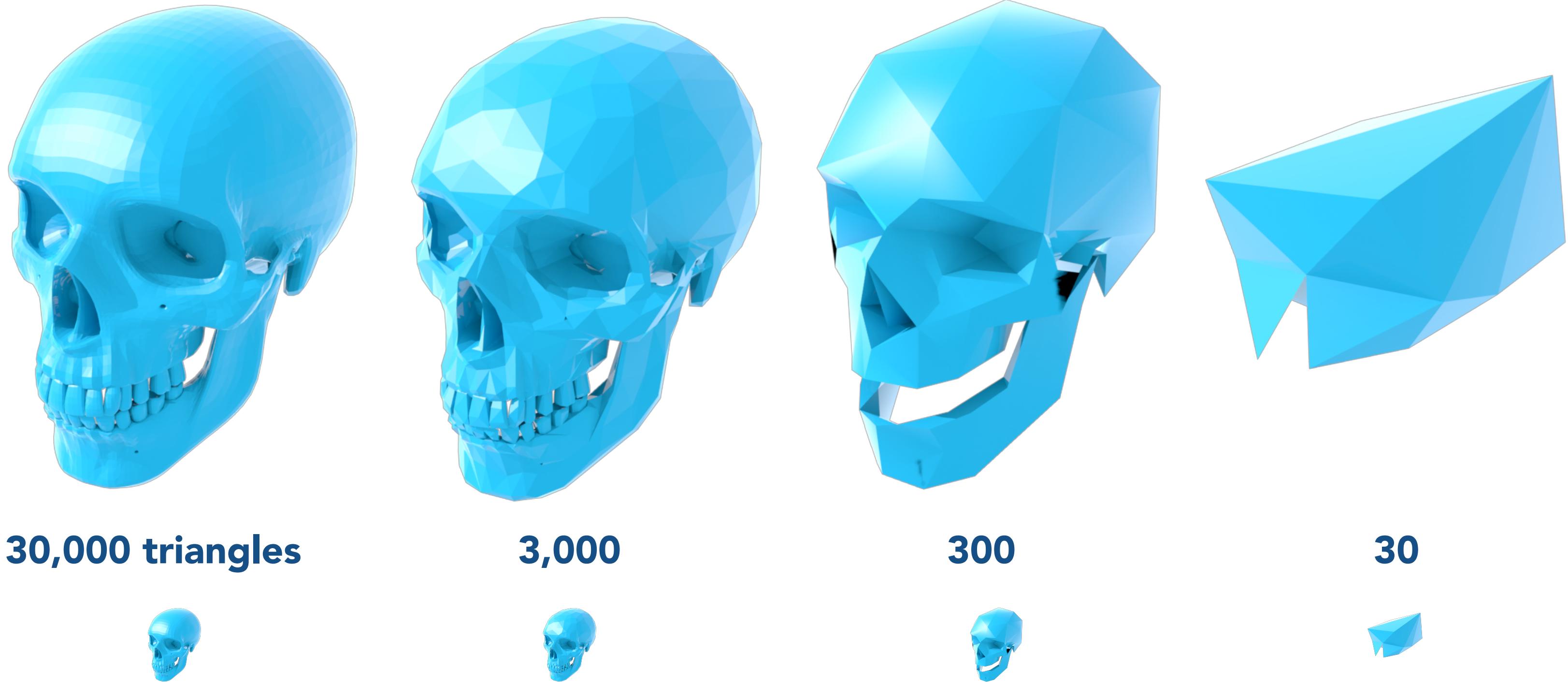
- Edge flip is (local) resampling:



- Still need to intelligently decide which edges to modify!

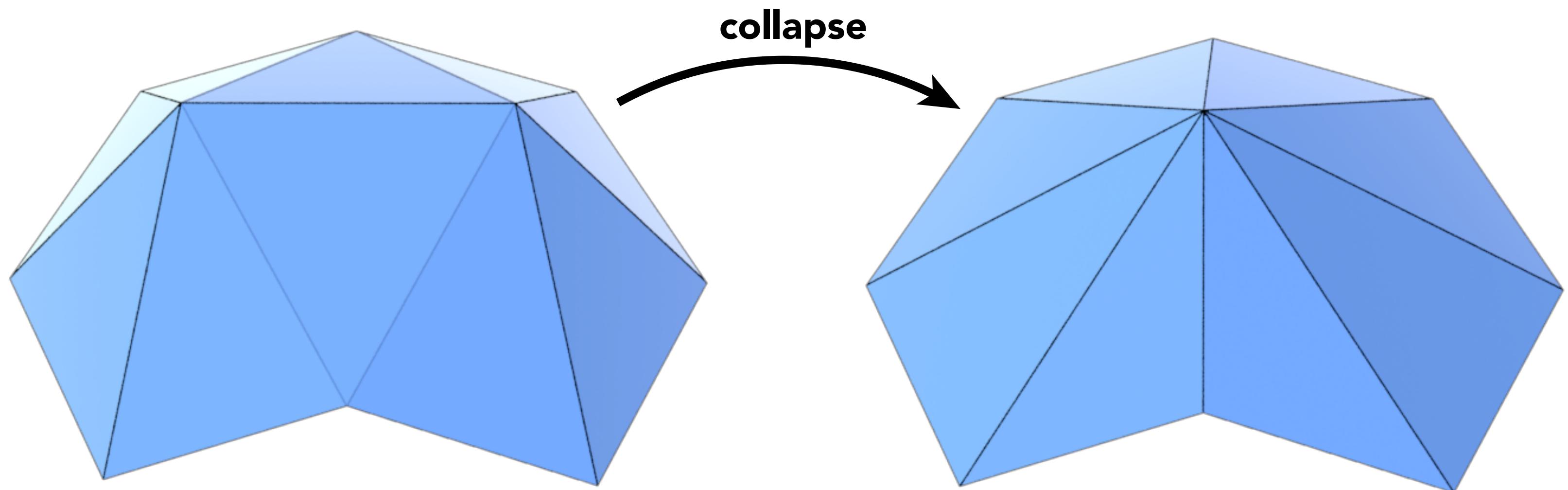
Mesh simplification

- Goal: reduce number of mesh elements while maintaining overall shape



Estimate: error introduced by collapsing an edge?

How much geometric error is introduced by collapsing an edge?



Sketch of Quadric Error Mesh Simplification

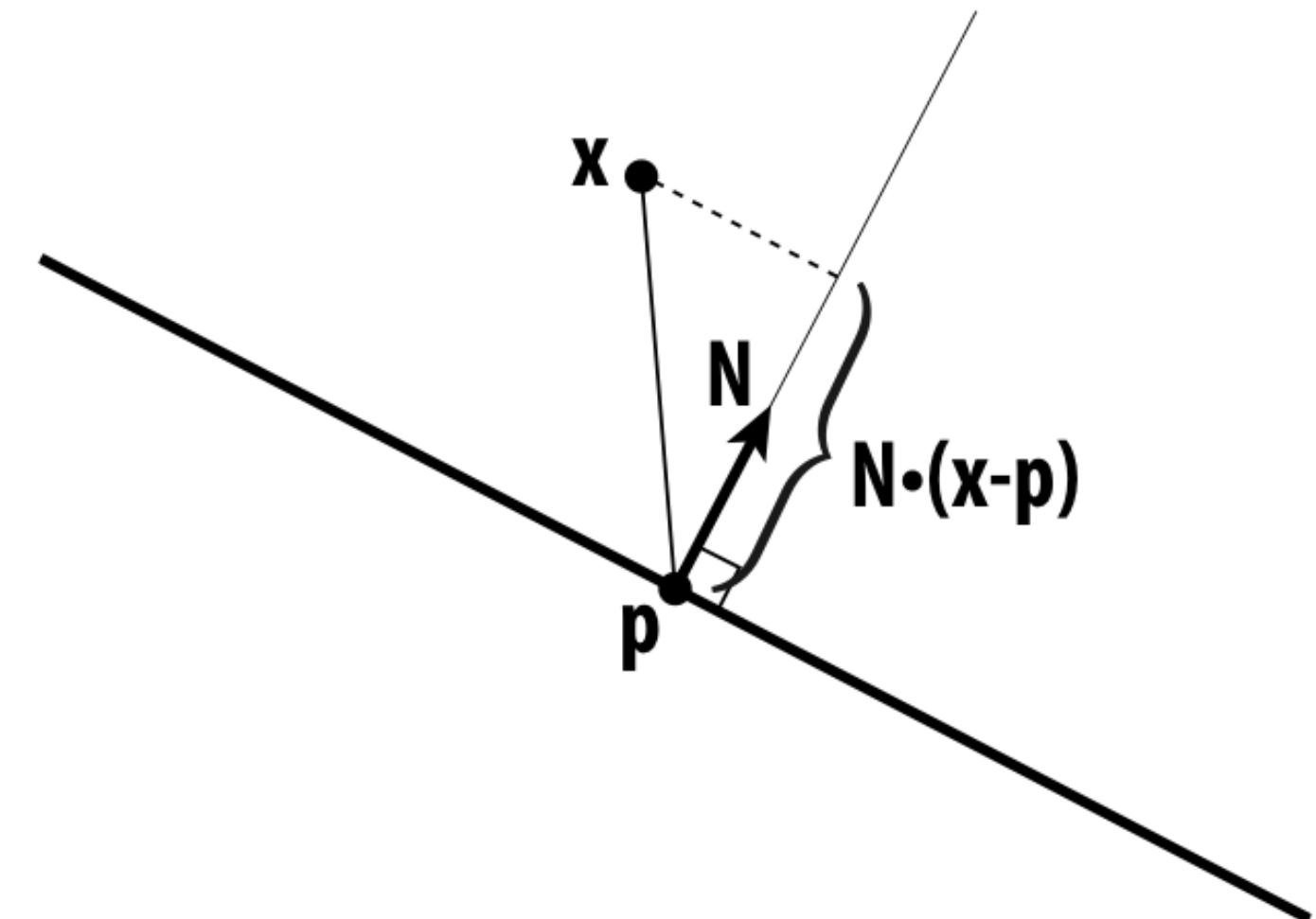
Simplification via quadric error

- Iteratively collapse edges
- Which edges? Assign score with quadric error metric*
 - Approximate distance to surface as sum of squared distances to planes containing nearby triangles
 - Iteratively collapse edge with smallest score
 - Greedy algorithm... great results!

* (Garland & Heckbert 1997)

Review: point-to-plane distance

Signed distance to plane with normal N
passing through point p ?
 $\Rightarrow N \cdot (x - p)$



Quadric error matrix (encodes squared distance)

- Suppose we have:
 - a query point (x, y, z)
 - a normal (a, b, c)
 - an offset $d := -(x_p, y_p, z_p) \cdot (a, b, c)$

$$Q = \begin{bmatrix} a^2 & ab & ac & ad \\ ab & b^2 & bc & bd \\ ac & bc & c^2 & cd \\ ad & bd & cd & d^2 \end{bmatrix}$$

- Then in homogeneous coordinates, let

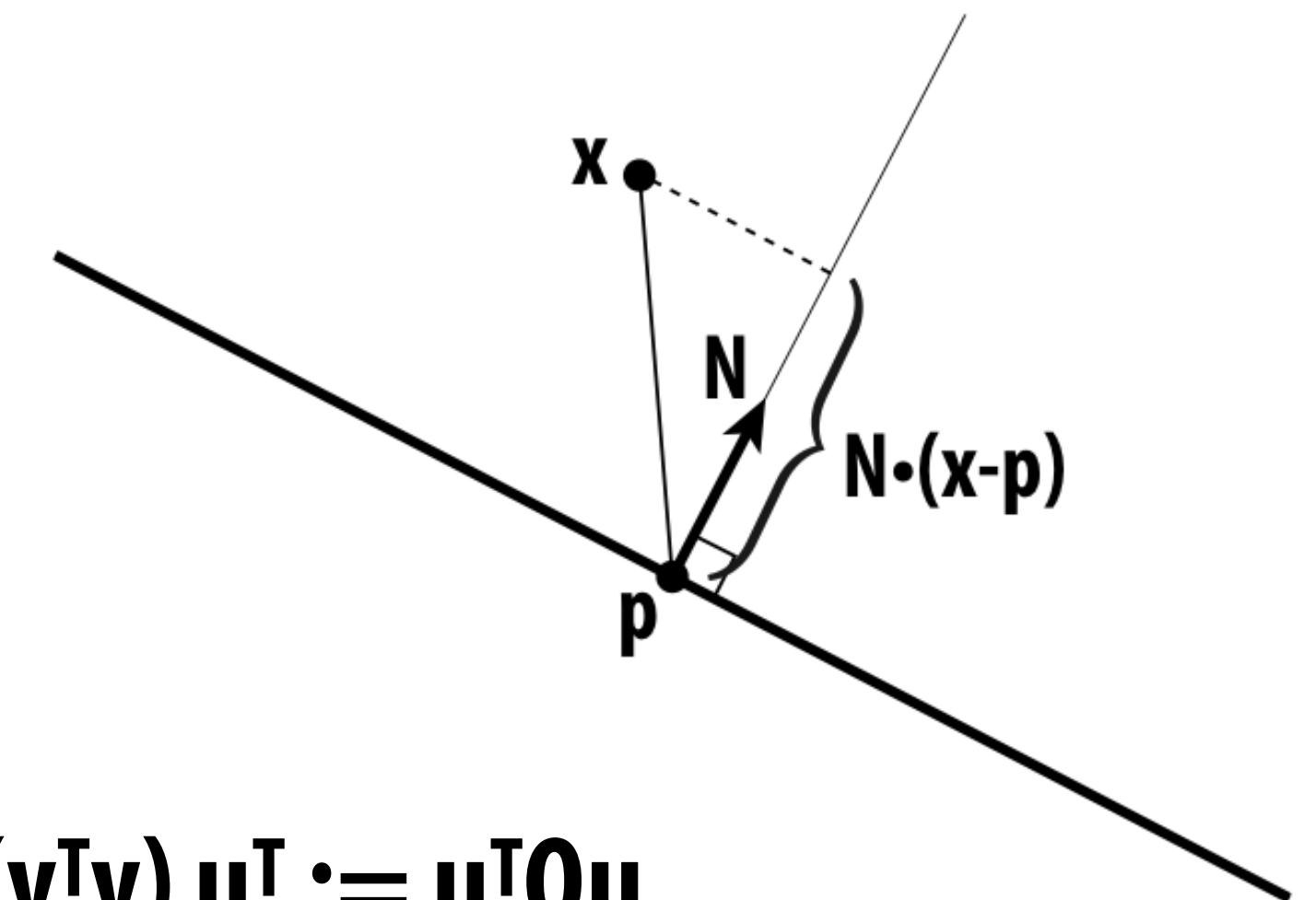
- $u := (x, y, z, 1)$
- $v := (a, b, c, d)$

- Signed distance to plane is then

$$D = uv^T = vu^T = ax + by + cz + d$$

- Squared distance is $D^2 = (uv^T)(vu^T) = u(v^Tv)u^T := u^TQu$

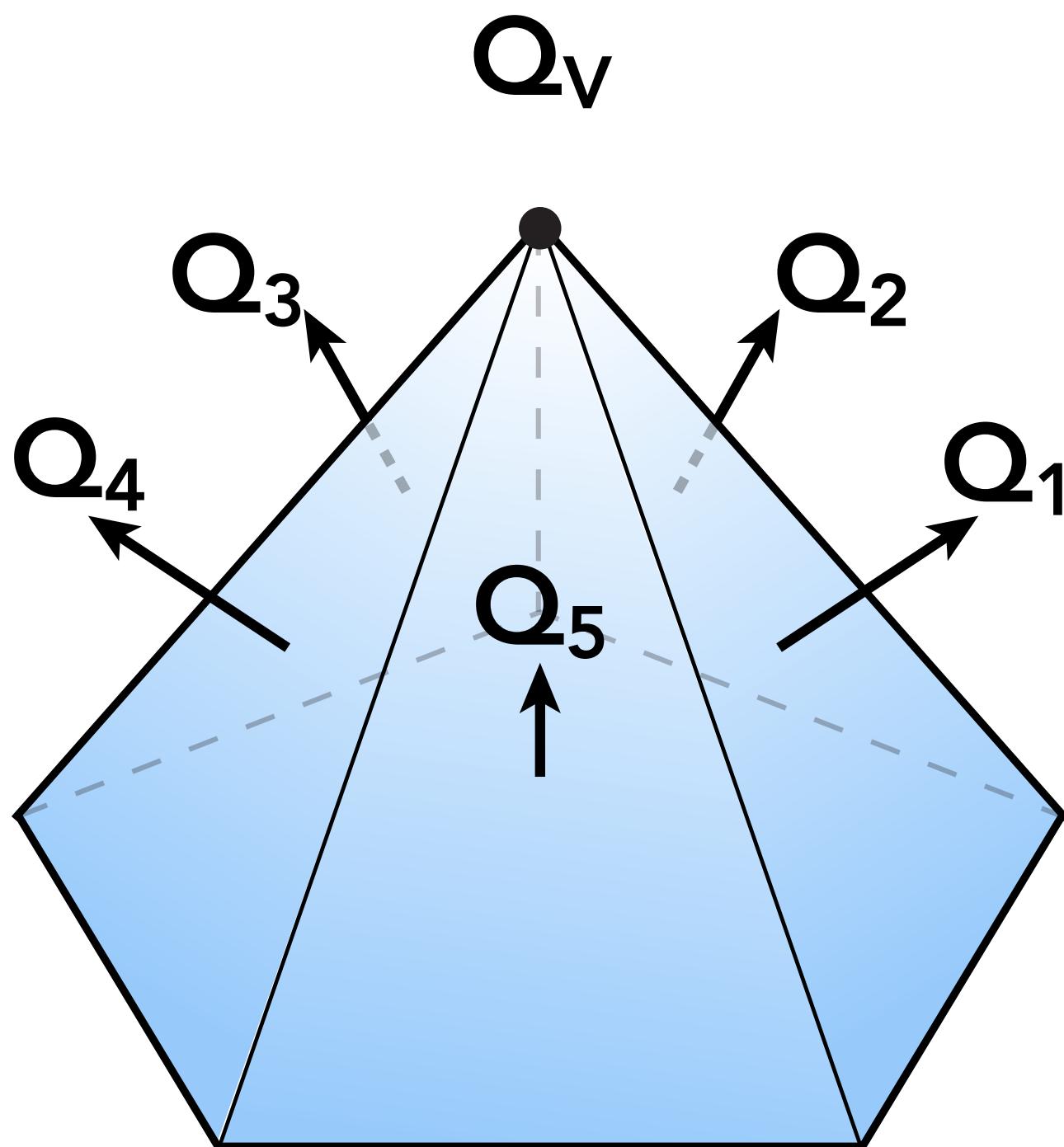
- Distance is 2nd degree ("quadric") polynomial in x, y, z



Quadric error at mesh vertex

Heuristic: error at vertex V is sum of squared distances to triangles connected to V

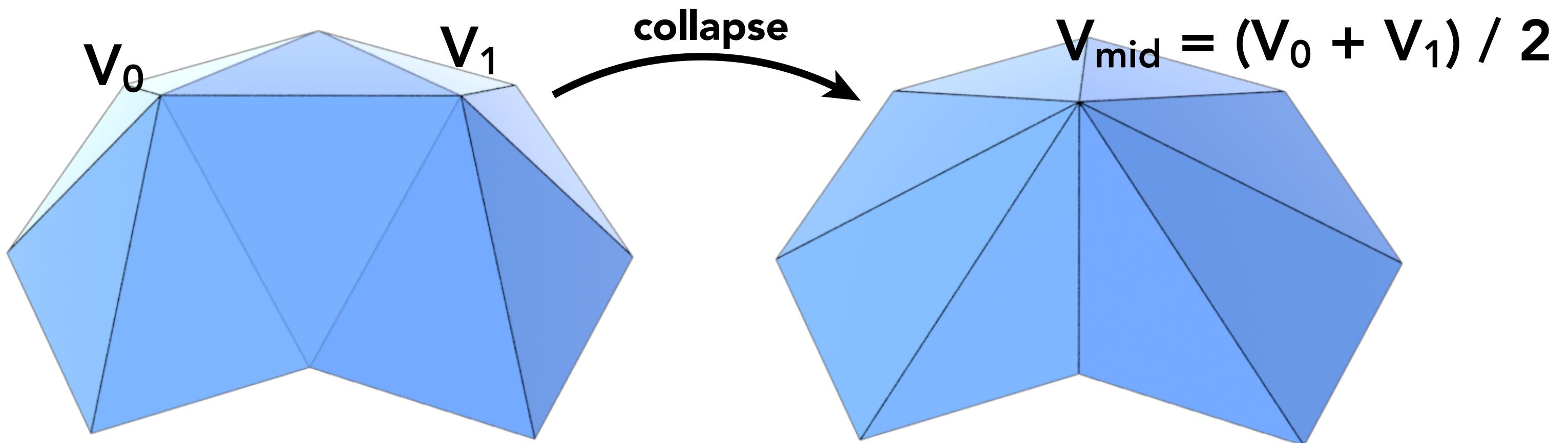
Encode this as a single quadric matrix per vertex that is the sum of quadric error matrices for all triangles



$$Q_V = \sum_{i=1}^N Q_i$$

Cost of edge collapse

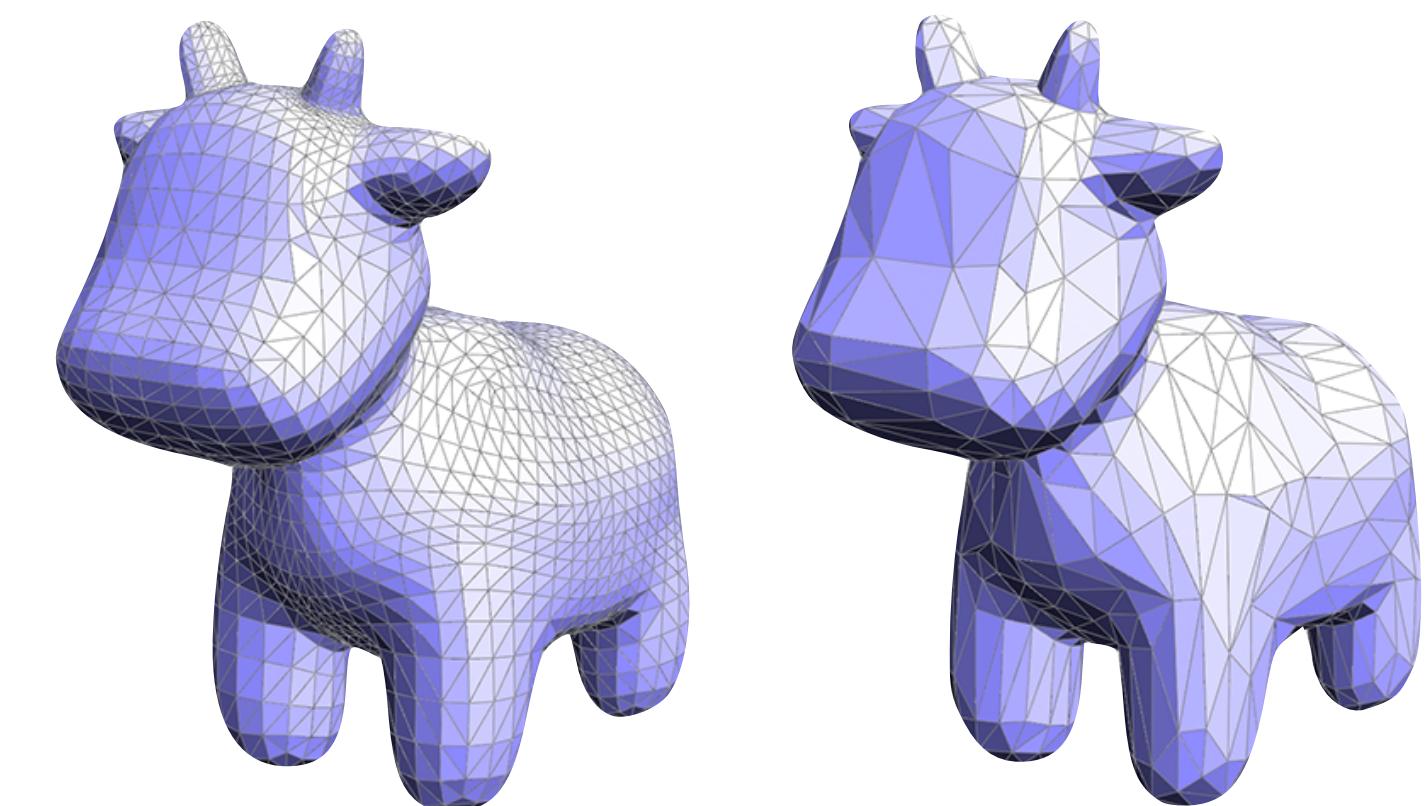
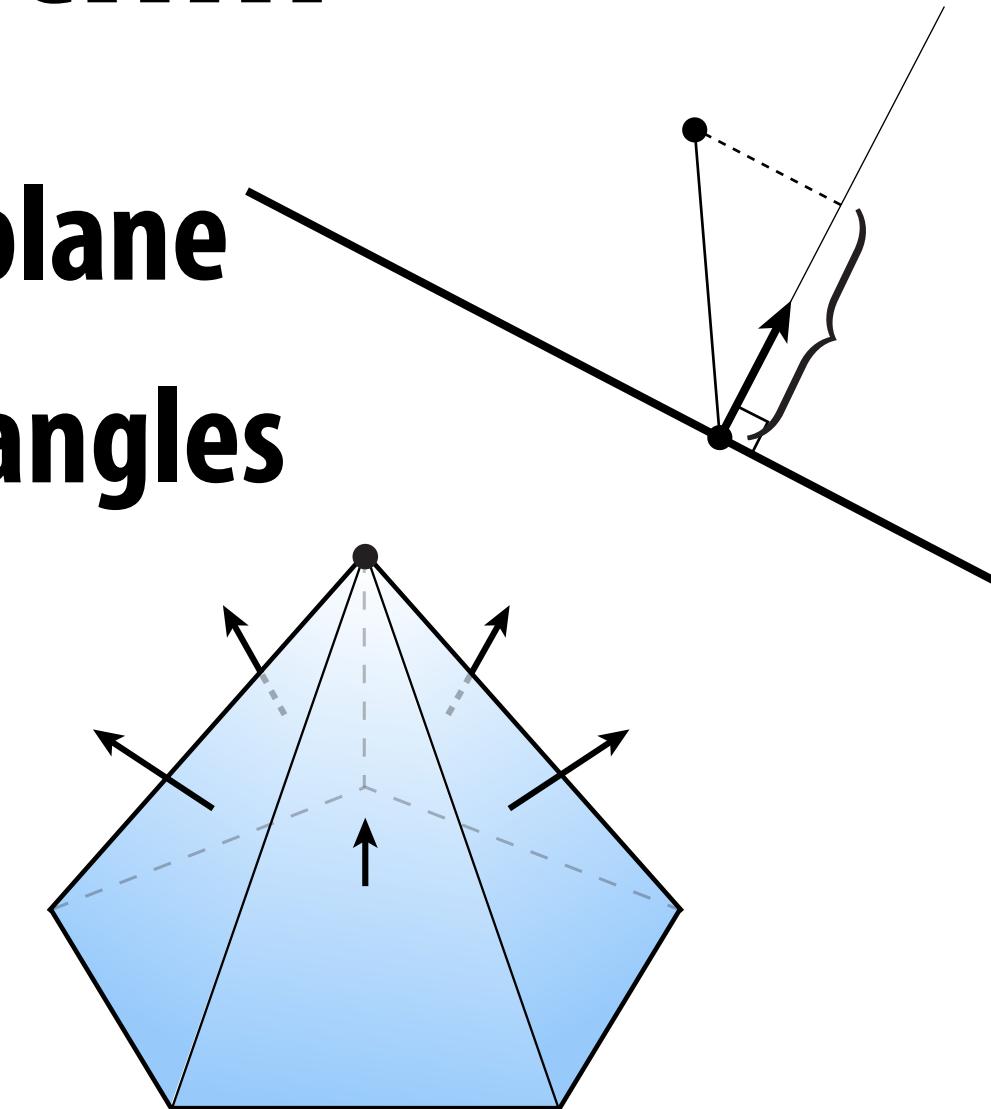
- How much does it cost to collapse an edge?
- Idea: compute edge midpoint V_{mid} , measure quadric error at this point
- Error at V_{mid} given by $v_{\text{mid}}^T(Q_0 + Q_1)v_{\text{mid}}$
- Intuition: cost is sum of squared differences to original position of triangles now touching V_{mid}



- Better idea: choose point on edge (not necessarily the midpoint) that minimizes quadric error
- More details: Garland & Heckbert 1997

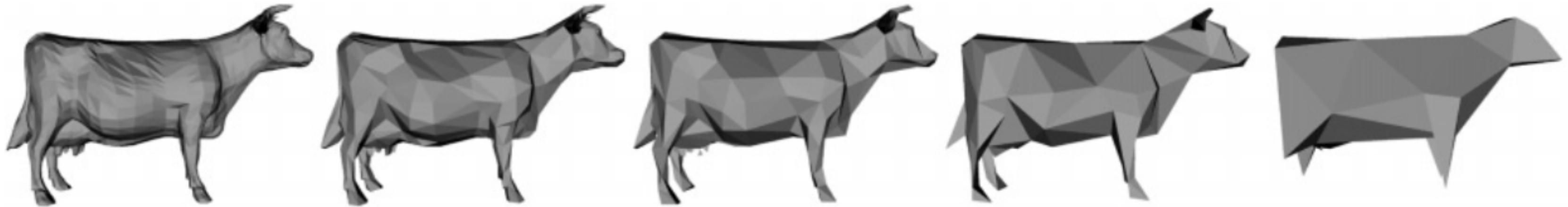
Quadric error simplification: algorithm

- Compute quadric error matrix Q for each triangle's plane
- Set Q at each vertex to sum of Q 's from neighbor triangles
- Set Q at each edge to sum of Q 's at endpoints
- Find point at each edge minimizing quadric error
- Until we reach target # of triangles:
 - collapse edge (i,j) with smallest cost to get new vertex m
 - add Q_i and Q_j to get quadric Q_m at vertex m
 - update cost of edges touching vertex m



Quadric error mesh simplification

Garland and Heckbert '97



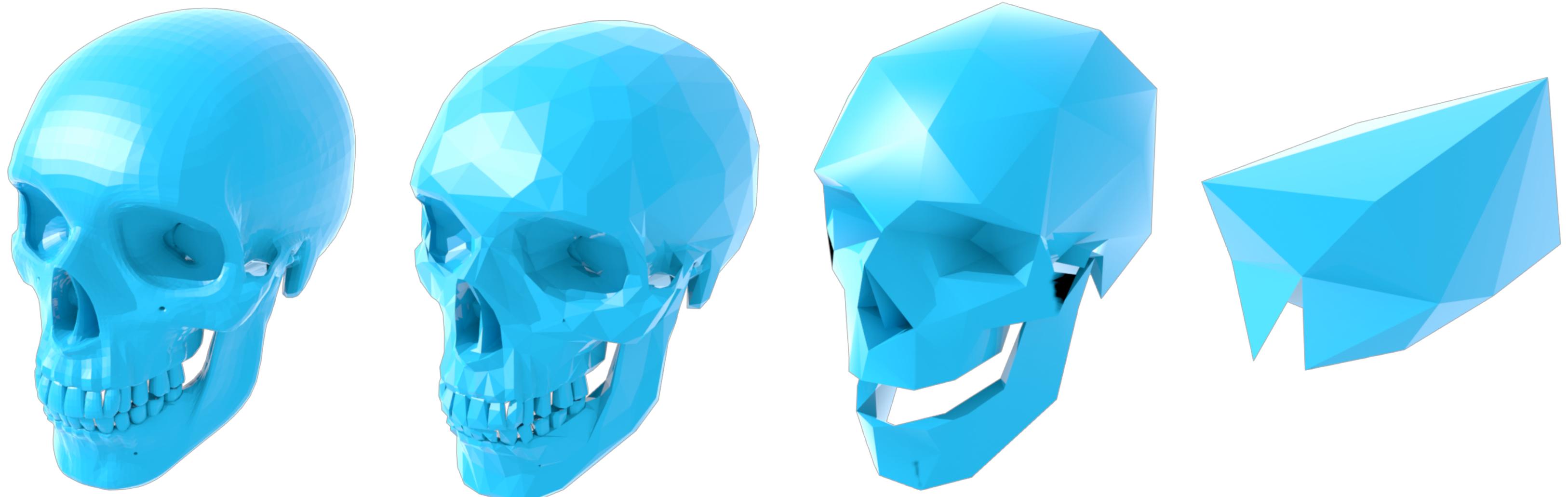
5,804

994

532

248

64



30,000 triangles

3,000

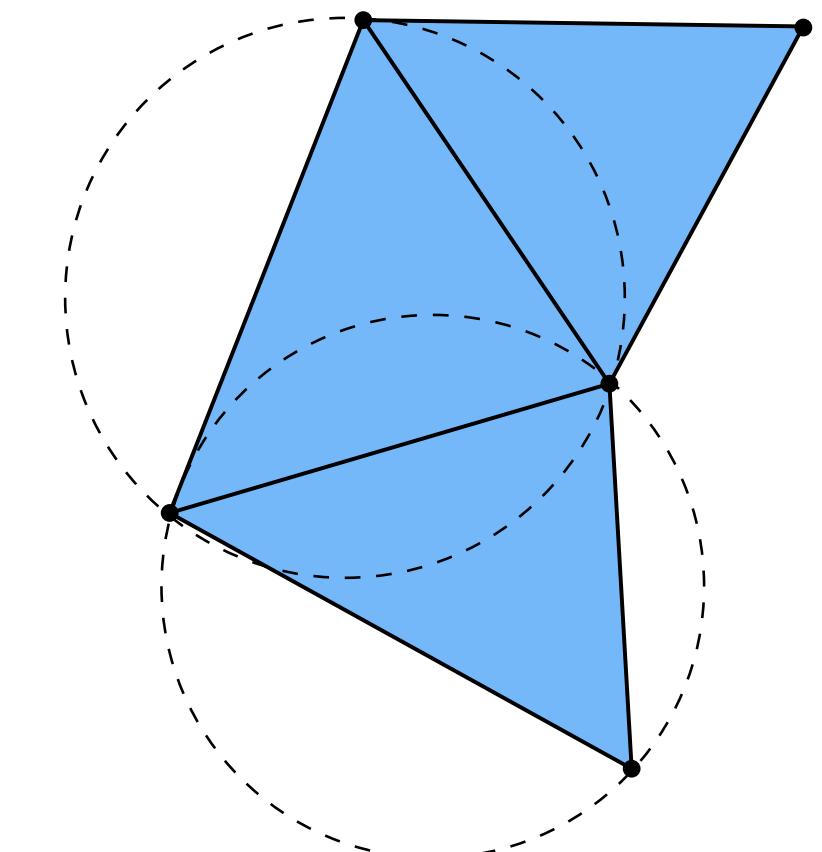
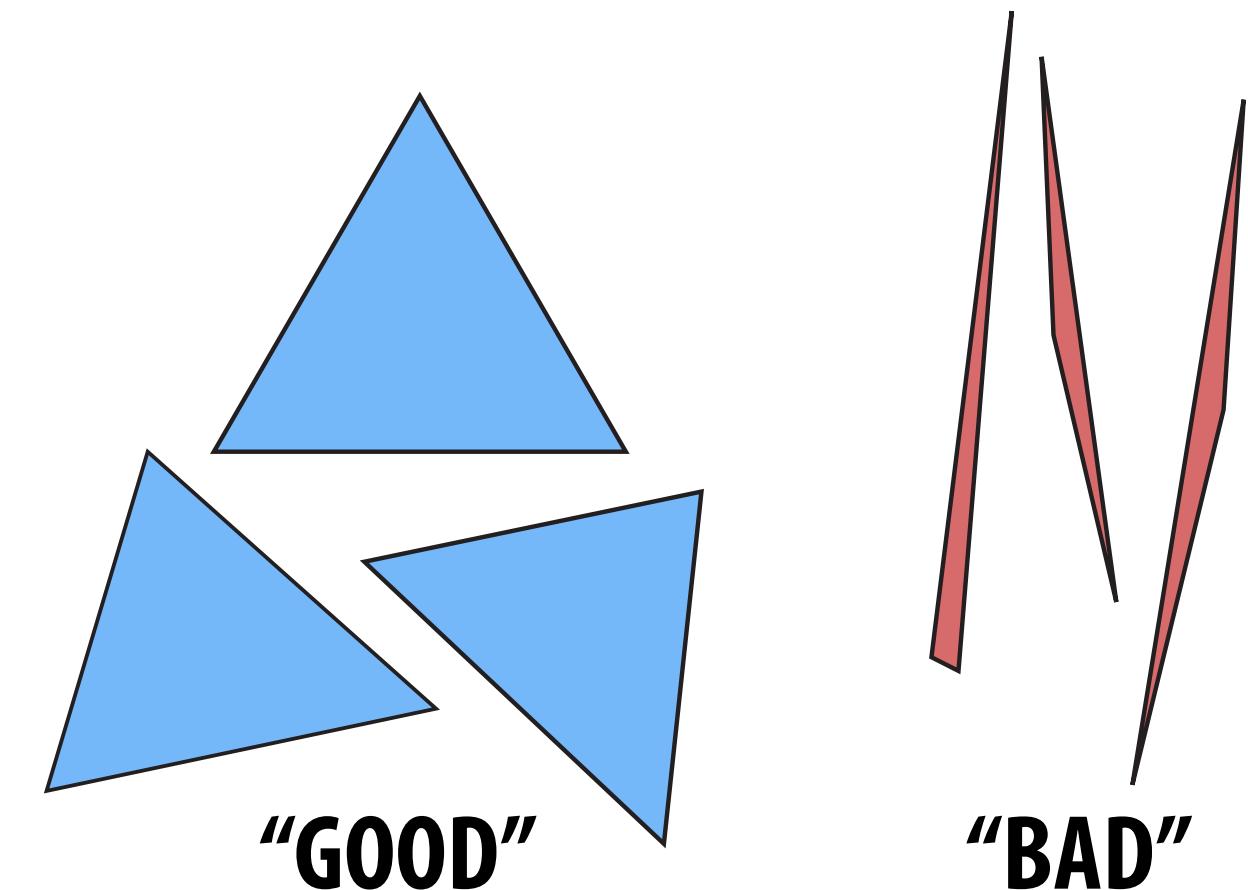
300

30

Mesh Regularization

What makes a “good” triangle mesh?

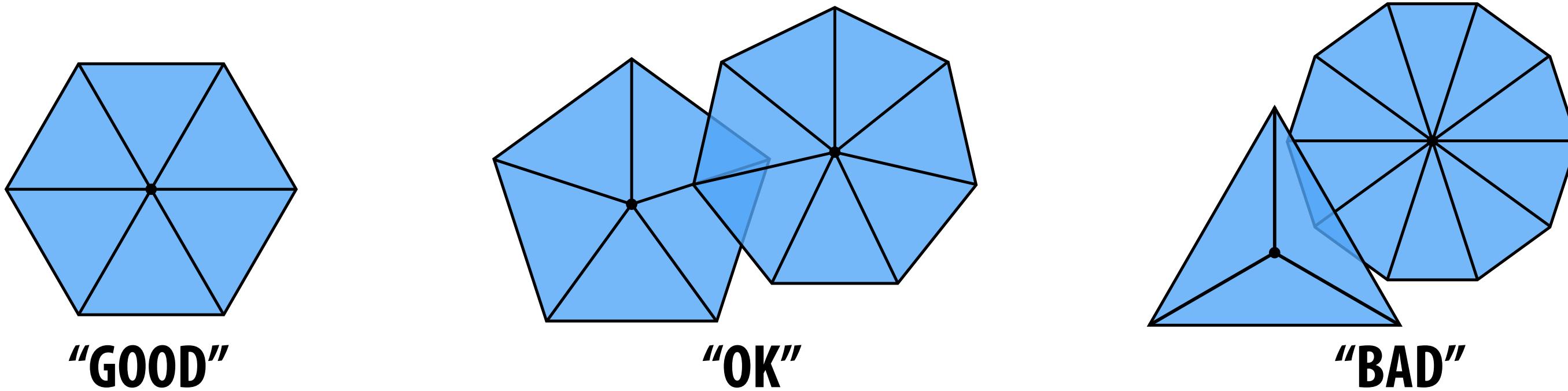
- One rule of thumb: triangle shape
- More specific condition: Delaunay
 - “Circumcircle interiors contain no vertices.”
- Not always a good condition, but often*
 - Good for simulation
 - Not always best for shape approximation



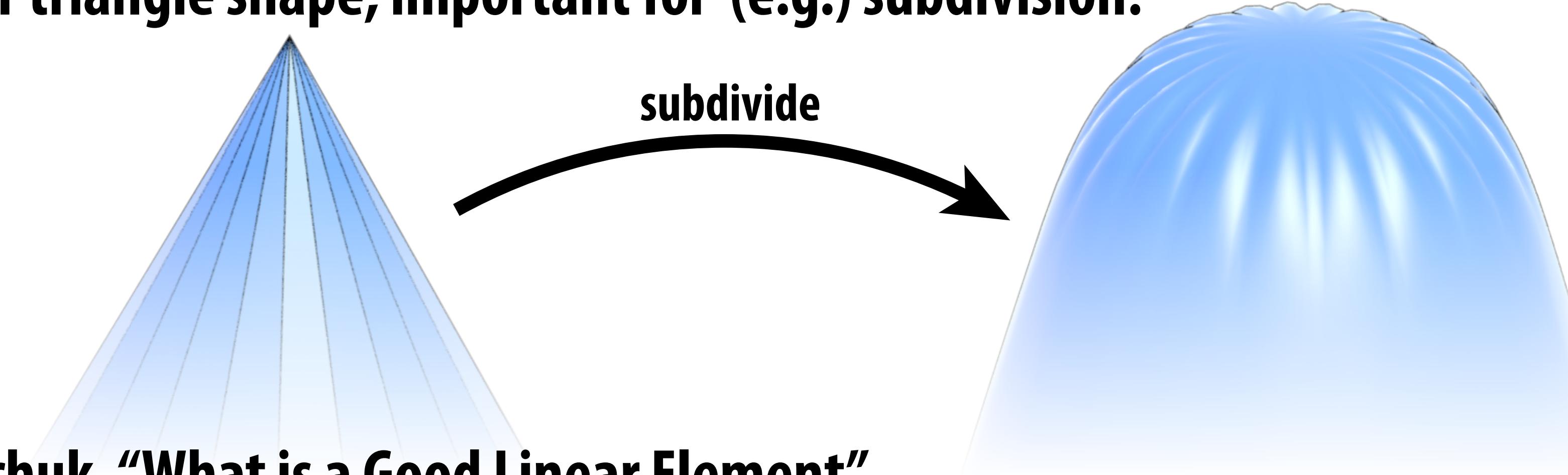
*See Shewchuk, “What is a Good Linear Element”

What else constitutes a good mesh?

- Rule of thumb: regular vertex degree
- Triangle meshes: ideal is every vertex with valence 6:



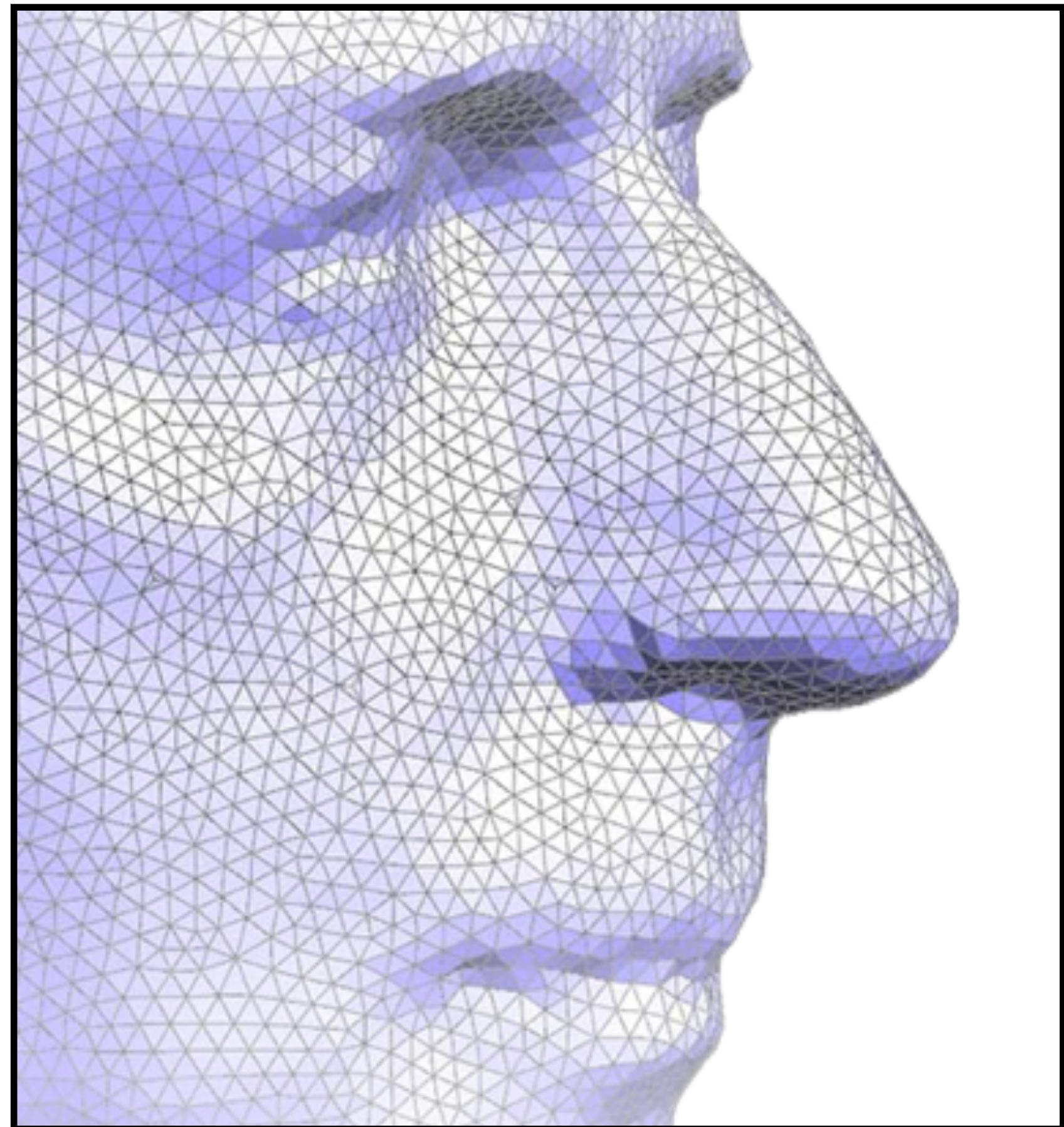
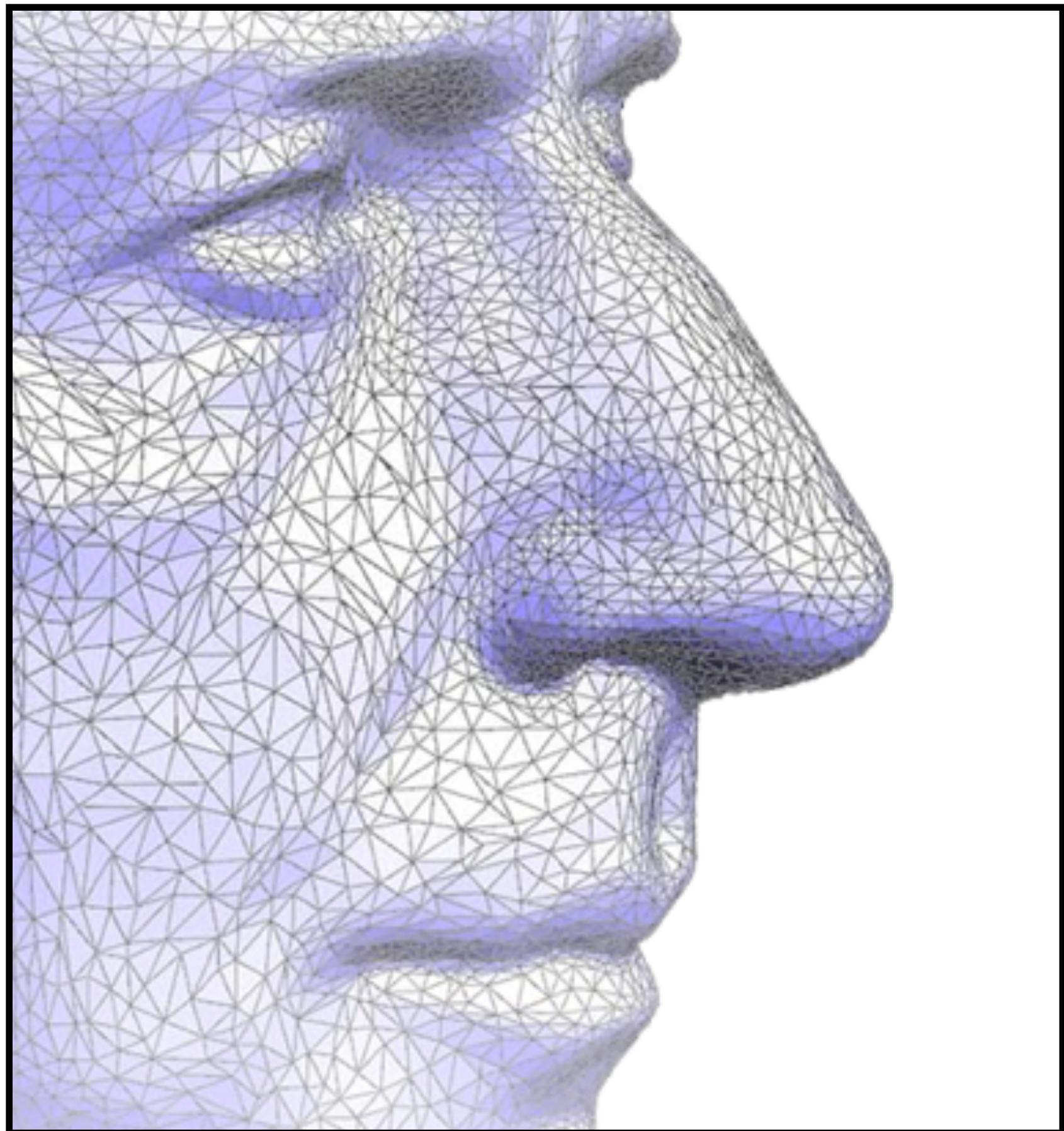
Why? Better triangle shape, important for (e.g.) subdivision:



*See Shewchuk, "What is a Good Linear Element"

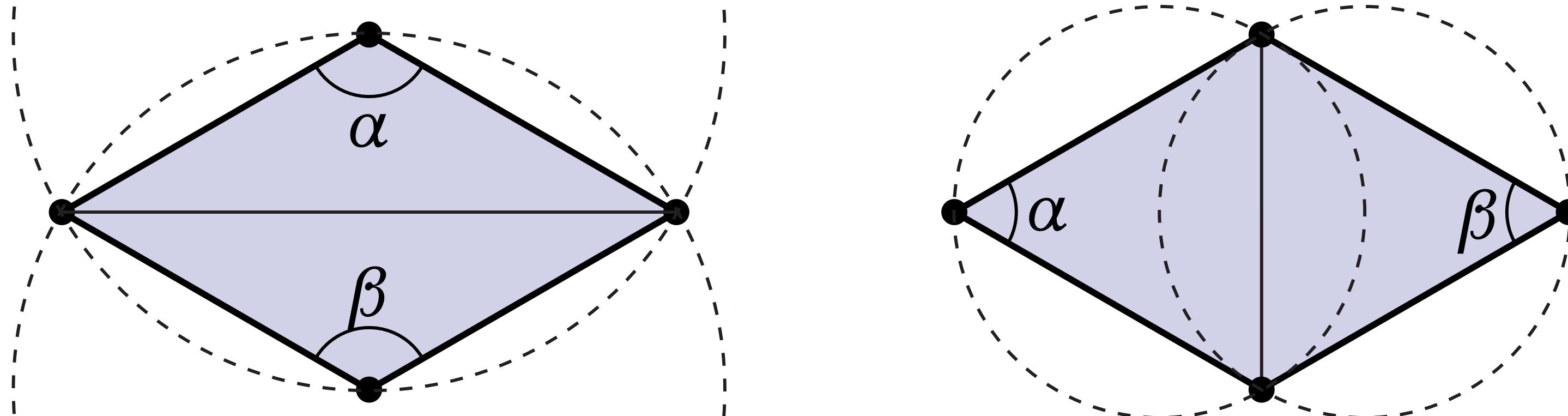
Isotropic remeshing

Goal: try to make triangles uniform in shape and size



How do we make a mesh “more delaunay”?

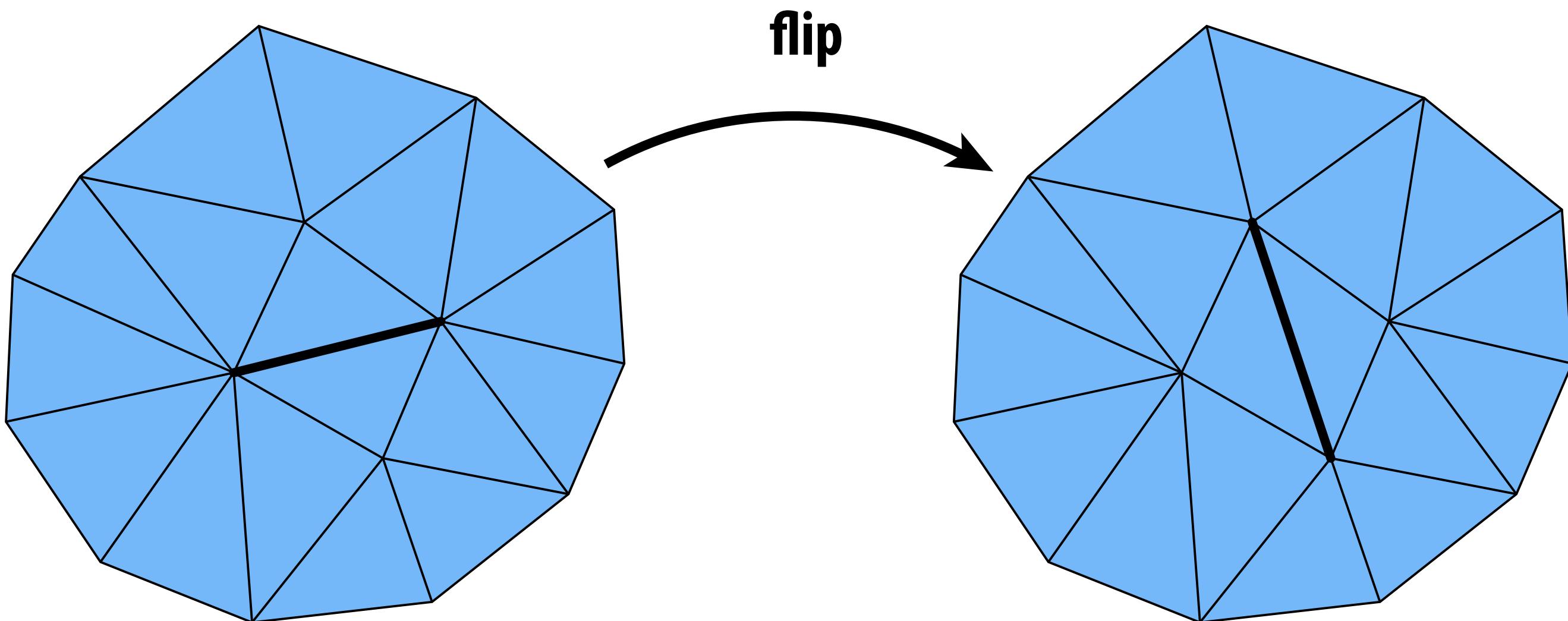
- Already have a good tool: edge flips!
- If $\alpha + \beta > \pi$, flip it!



- In practice: a simple, effective way to improve mesh quality

How do we improve degree?

- Edge flips!
- If total deviation from degree 6 gets smaller, flip it!

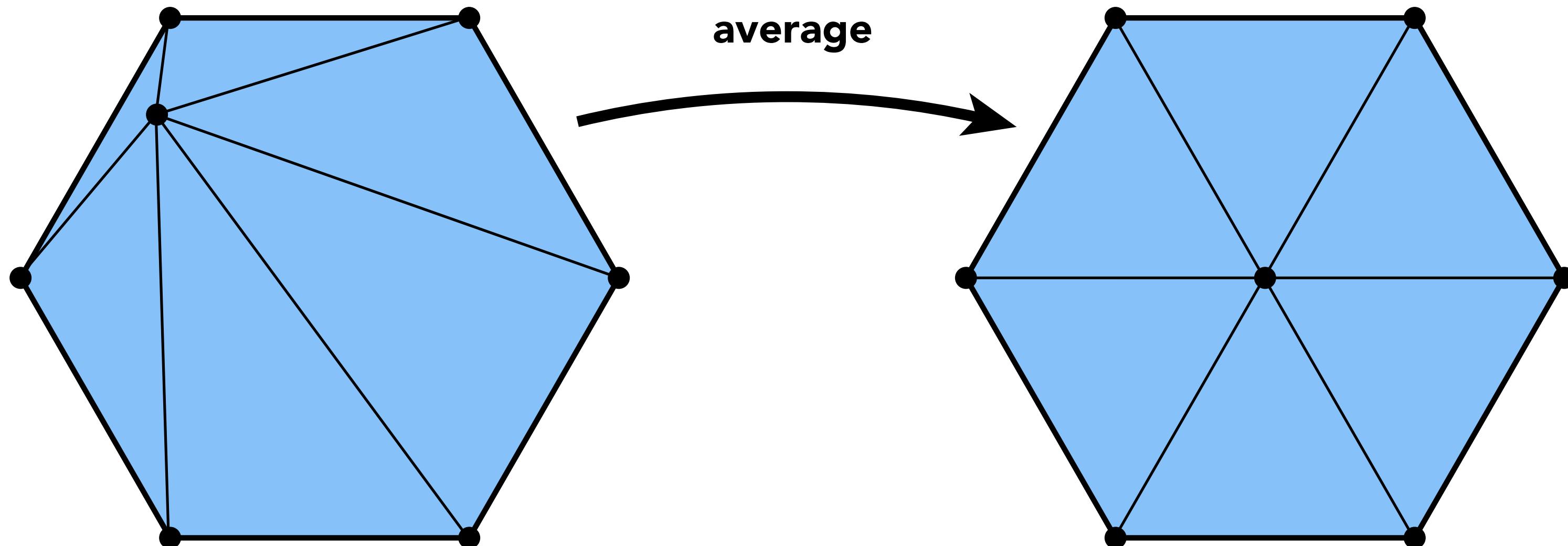


Iterative edge flipping acts like “discrete diffusion” of degree

No (known) guarantees; works well in practice

How do we make triangles “more round”?

- Delaunay doesn’t mean equilateral triangles
- Can often improve shape by centering vertices:

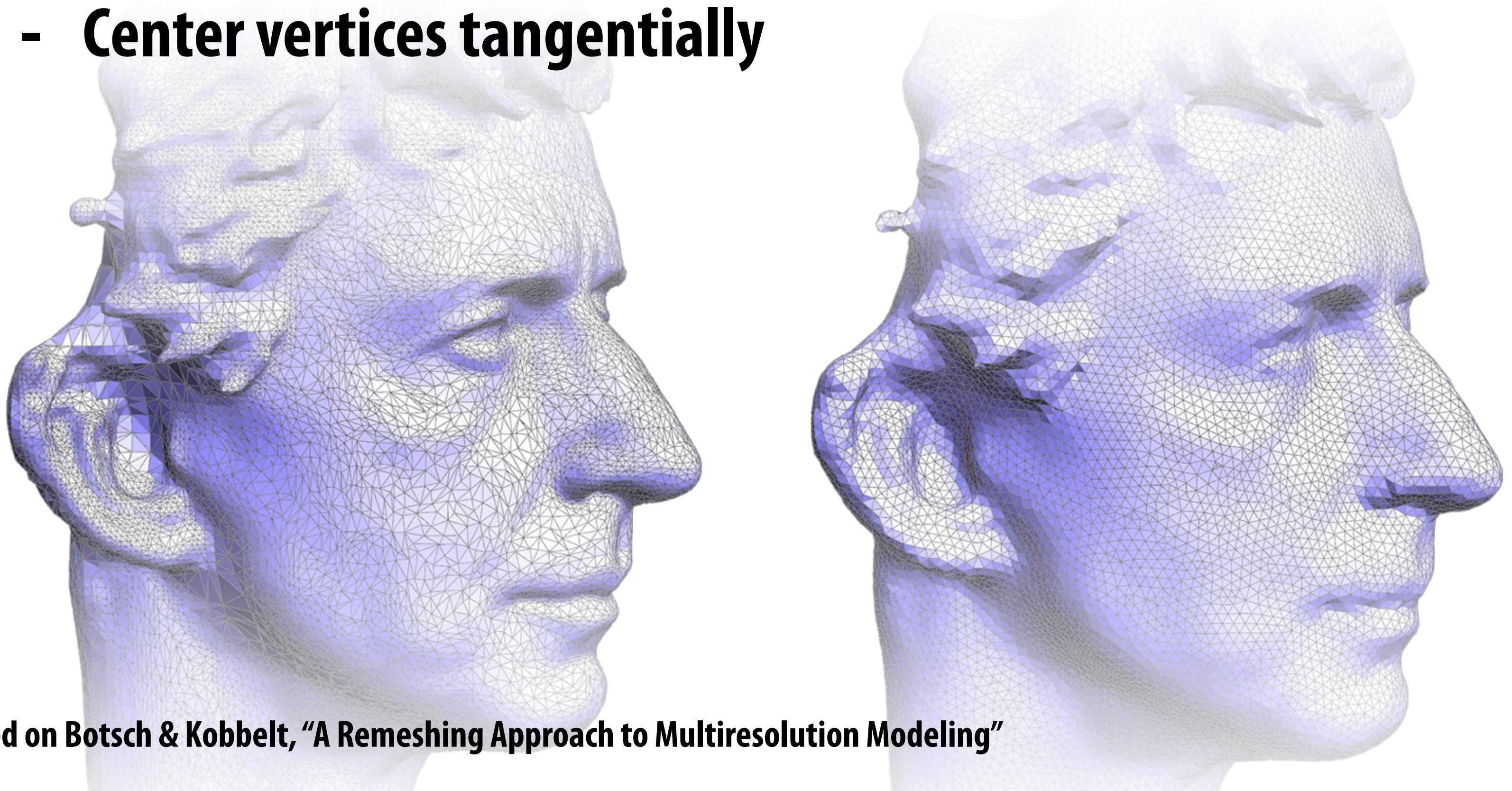


[See Crane, “Digital Geometry Processing with Discrete Exterior Calculus”]

Isotropic remeshing algorithm*

■ Repeat four steps:

- Split edges over 4/3rds mean edge length
- Collapse edges less than 4/5ths mean edge length
- Flip edges to improve vertex degree
- Center vertices tangentially



* Based on Botsch & Kobbelt, "A Remeshing Approach to Multiresolution Modeling"

Things to remember

- **Triangle mesh representations**
 - **Triangles vs points+triangles**
 - **Half-edge structure for mesh traversal and editing**
- **Geometry processing basics**
 - **Local operations: flip, split, and collapse edges**
 - **Upsampling by subdivision (Loop, Catmull-Clark)**
 - **Downsampling by simplification (Quadric error)**
 - **Regularization by isotropic remeshing**

Acknowledgements

- **Thanks to Keenan Crane, Ren Ng, Pat Hanrahan, James O'Brien, Steve Marschner for presentation resources**