CSCI 3104 PS1a

Luna Mcbride

TOTAL POINTS

8.5 / 10

QUESTION 1

13/3

√ + 3 pts Correct

- + 2 pts Something wrong about maintenance step. Please check solution, you need to clarify that LI is true before/after each iteration of loop.
 - + 1 pts only one component is correct
 - + 1 pts Generally in right direction, but incorrect.
 - + 1 pts No one-sentence description for each one.
 - + 2 pts Your Maintenance section is not correct
 - + 2 pts Termination part is not clear
 - + 0 pts Not try or totally incorrect
 - + 0 pts Click here to replace this description.

QUESTION 2

6 pts

2.1 2 / 2

√ + 2 pts Correct

- + 2 pts See the solution to phrase the LI better. Specifically give a context to 'i'. For ex - you can start with "At the beginning of ith iteration..".
- + **1.5 pts** In the right direction but a Loop Invariant is stated differently. See solutions.
 - + 0 pts Not attempted
- √ + 0 pts (Non-scored rubric item) Good work on attempting the complete proof. But you were good when you stated the Loop Invariant. We are not looking at that portion of your solution now.
- + **0 pts** Incorrect answer please check the solutions provided.
- + $\bf 2$ pts Please also mention that at the start of the ith iteration ret has the maximum the subarray A[0, 1,, i-1]
- + 2 pts Please mention when (start or end of the iteration) does the variable ret contains the greatest

element and in which subarray.

- + 2 pts There is no variable max, variable ret stores the maximum value.
- + 1 pts The Loop invariant representation is not right. Please check the solution.
- + 2 pts See the solution to see how Loop Invariant's are written
- + 1 pts ret holds the maximum value and not it's index.

2.2 0.5 / 2

+ 2 pts Correct

√ + 0 pts Incorrect

- + 2 pts Technically correct, but does not define loop invariant precisely
- \checkmark + 0.5 pts The loop invariant should reflect the logic through which the algorithm works. Your answer is technically true, but it does not serve as a real loop invariant in terms of the solution to the problem.
- + **0.5 pts** Partially correct. ret = -1 condition not specified correctly when n is not in the array.
- + **0.5 pts** Partially correct. index value of n condition not specified correctly

2.3 2/2

√ + 2 pts Correct

- + 0 pts Not attempted
- + 1 pts Need more explanation
- + 2 pts Phrase it better. Please see the solution provided.
- + **0 pts** Incorrect answer. Please read and understand the concept again
- + **0 pts** variable 'sum' is the sum of A[0....i-1], not A[0....i], before the start of each iteration. Please read again and understand the subarray concept.
 - + 0 pts Can be written more clearly

- + **0 pts** The sum is equal to the sum of the elements in the subarray A[0:i-1], not just the elements in the subarray.
 - + **0 pts** Please write maintenance step properly.
- + 1 pts Answer seems to be in the right direction, but it is not clear and enough.

QUESTION 3

3 1/1

- + **0 pts** Incomplete or Incorrect Base case and Inductive Hypothesis, or Empty Solution submitted. Please check the solution for reference.
 - + 1 pts Correct and Great Work !!!
- + **0.5 pts** Missing or Incorrect Base Case. Please Check Solution File.
- + **0.5 pts** Missing or Incorrect Inductive Hypothesis. Please Check Solution File for reference.
- $\sqrt{+1}$ pts In the right direction. But please compare your solution with the actual solution file so that you do perfect the next time.
- + 1 pts In the right direction. Missing what LHS and RHS are, but proved both the sides to be equal. Please check solution file for your reference and do better the next time.
- + 1 pts In the right direction. Please check the corrected question and solution file for reference so that you do perfect the next time.
- + **0.5 pts** In the right direction. But please compare your solution with the actual solution file so that you do perfect the next time.
- + **0.5 pts** In the right direction, but incorrect question assumed and hence tried to prove the wrong statement.

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ID: 107607144

CSCI 3104, Algorithms Problem Set 1a (10 points) Profs. Hoenigman & Agrawal Fall 2019, CU-Boulder

Advice 1: For every problem in this class, you must justify your answer: show how you arrived at it and why it is correct. If there are assumptions you need to make along the way, state those clearly.

Advice 2: Verbal reasoning is typically insufficient for full credit. Instead, write a logical argument, in the style of a mathematical proof.

Instructions for submitting your solution:

- The solutions **should be typed** and we cannot accept hand-written solutions. Here's a short intro to Latex.
- You should submit your work through **Gradescope** only.
- If you don't have an account on it, sign up for one using your CU email. You should have gotten an email to sign up. If your name based CU email doesn't work, try the identikey@colorado.edu version.
- Gradescope will only accept .pdf files (except for code files that should be submitted separately on Gradescope if a problem set has them) and try to fit your work in the box provided.
- You cannot submit a pdf which has less pages than what we provided you as Gradescope won't allow it.

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- 1. (3 pts) What are the three components of a loop invariant proof? Write a one-sentence description for each one.
 - 1) Initialization; This is the step where you make sure the invariant holds before even entering the loop.
 - 2) Maintenance; This is where you check if the invariance holds for each individual loop. It does not matter if you check immediately after the last loop or immediately before the next, but it must be checked to make sure this applies for all values.
 - 3) Termination; At this point the beginning and middle has been checked, however, it must also end by following the invariance to truly fit. Once we have the end value, that makes all values checked and we can then definitively prove the relation.

2. (6 pts total) Identify the loop invariant in the following algorithms.

```
(a) FindMaxElement(A) : //suppose array A is not empty
ret = A[0]
for i = 1 to length(A)-1 {
    if A[i] > ret{
        ret = A[i]
    }}
return ret
```

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Well, a maximum is the highest value, so the invariance should be that a value is the highest up until now.

Invariance: $(A[0],...,A[i-1]) \le ret$ for $0 \le i \le length(A)-1$ (deincremented as the for loop increments automatically)

Initialization: We start with ret=A[0], which is the highest of all the values we have looked at thus far, being itself.

Maintenance: The if statement takes sets the ret to the new highest if A[i] (before increment) is bigger, forcing it to be the highest of the values seen such far. That means all others could be less than or equal to the current highest ret, but never greater as the highest was just taken by the if statement. Termination: The final value (length(A)-1 value, as we start at 0), once checked either takes the last highest value or holds a previous value that is highest for ret. There is no chance of a value greater (if coded it right, that is), and thus is the ret value can only be greater than or equal to absolutely everything else in array A.

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(b) FindElement(A, n) : //suppose no duplicates in array A and array A is not empty
ret = -1 //index -1 implies the element haven't been found yet
for i = 0 to length(A)-1 {
 if A[i] == n{
 ret = i
 }}
return ret

Note: I did parts a and c before checking piazza and thus seeing to just type the invariance, but I do not want to delete all that typing and it is how I thought it out. sorry.

Invariance: ret \geq -1

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```
(c) SumArray(A) : //suppose array A is not empty
sum = 0
for i = 0 to length(A)-1 {
    sum += A[i]
}
return sum
```

Invariance: sum=A[0]+A[1]...+A[i-1] for $0 \le i \le length(A)-1$ (Assuming a +0 for the base, however, that is unnecessary)

Initialization: sum=0, which takes the case of the assumed 0, not including the A[i-1] as -1 is outside the range and would probably break the computer

Maintenance: Sum adds the array value at each value i in the loop. Following the idea listed in C++ logic (i incremented at the end of the loop as part of the loop deinition (i=0;i;x;i++)), the i needs to be deincremented to get the correct index when before or after an iteration to count it.

Termination: the last value is at the length of A (-1 as to account for the index start at 0). The i still increments, but the for loop stops it after the increment (the i_ix in (i=0;i_ix;i++)), making the last i value length(A) and the last added value A[i-1], the total amount of indecies to account the whole array.

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3. (1 pt) If r is a real number not equal to 1, then for every $n \geq 0$,

$$\sum_{i=0}^{n} r^{i} = \frac{(1 - r^{n+1})}{(1 - r)}.$$

Provide the first two steps of a proof by induction i.e. base case and the inductive hypothesis. You will be asked to complete this proof later in **PS1b**.

Base case: i=0base case: 1=0 $\sum_{i=0}^{0} r^{i} = \frac{(1-r^{n+1})}{(1-r)}$ $r^{0} = 1$ $\frac{(1-r^{0+1})}{(1-r)}$ $\frac{(1-r)}{(1-r)}$ $\frac{1-r}{1-r} = 1$

 $r^0 = \frac{(1-r^{0+1})}{(1-r)}$

1 = 1 (Plugged in equivalent values shown above)

Therefore, since 1 does equal 1, the equation holds for the base case

Inductive Hypothesis: If the equation works for the base case, it should be fair to assume this should work for some integer k. $\sum_{i=0}^{k} r^i = \frac{(1-r^{k+1})}{(1-r)}$ where $k \geq 0$ and r != 1