

## Homework 2

1. (a) Global CPI for P1:  $1 \times 10\% + 2 \times 20\% + 3 \times 50\% + 3 \times 20\% = \boxed{2.6}$

Global CPI for P2:  $2 \times 10\% + 2 \times 20\% + 2 \times 50\% + 2 \times 20\% = \boxed{2}$

(b) Class A:  $1.0 \times 10^6 \times 10\% = 1 \times 10^5$  instructions

Class B, C, D have  $2 \times 10^5$ ,  $5 \times 10^5$ ,  $2 \times 10^5$  instructions

Clock Cycles for P1:  $1 \times 10^5 \times 1 + 2 \times 10^5 \times 2 + 5 \times 10^5 \times 3 + 2 \times 10^5 \times 3 = \boxed{2.6 \times 10^6}$

Clock Cycles for P2:  $1 \times 10^6 \times 2 = \boxed{2 \times 10^6}$

(c) Time cost of P1:  $\frac{2.6 \times 10^6}{2.5 \times 10^9} = 1.04 \times 10^{-3} \text{ s}$

Time cost of P2:  $\frac{2 \times 10^6}{3 \times 10^9} \approx 6.67 \times 10^{-4} \text{ s}$

Hence  $\boxed{P2}$  is faster.

2. (a) x30 is  $\boxed{0x50000000}$ .

It is overflow, because MSB of x5 and x6 are both 1, but MSB of x30 is 0, it must overflow.

(b) x30 is  $\boxed{0xB0000000}$ .

It is NOT overflow, because MSB of x5 and x6 are the same.

3. 23 in signed 8-bit: ~~0000~~ 0001011<sub>2</sub>

112 in signed 8-bit: 01110000<sub>2</sub>

(a)  $00010111_2 + 01110000_2 = 10000111_2$ , MSB of 23 and 112 are both 0, but MSB of result is 1, overflow. Saturating arithmetic result in  $\boxed{127}$  in decimal.

(b)  $00010111_2 - 01110000_2 = 00010111_2 + 10010000_2 = 10100111_2$ , MSB of 23 and 112 are both 0, no overflow. Saturating arithmetic result is  $\boxed{-89}$  in decimal.

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4.	iter	multiplicand	product	operation
	0	00010000	00000000 01100100	
	1		00000000 00110010	shift right
	2		00000000 00011001	shift right
	3		00010100 00011001 00001010 00001100	add multiplicand to left shift right
	4		00000101 00000110	shift right
	5		00000010 10000011 00010110 10000011	shift right add multiplicand to left
	6		00001011 01000001	shift right
	7		00011110 10000001 00011110 10000000	add multiplicand to left shift right
	8		<u>000001111010000</u>	shift right

Hence,  $0x62 \times 0x14 = 0x07D0$  ( $100 \times 20 = 2000$  in decimal).

5.	iter	Q	Divisor	R	operation
	0	000000	011100000000		
	1	000000	001010100000	000000111110 101011111110 000000111110	$R = R - Div$ $R < 0$ , $R = R + Div$ , shift 0 into Q, <del>shift Div</del> Shift Div right $R = R - Div$
	2	000000	000101010000	000000111110	$R < 0$ , $R = R + Div$ , shift 0 into Q Shift Div right $R = R - Div$
	3	000000	000010101000	000000111110	$R < 0$ , $R = R + Div$ , shift 0 into Q Shift Div right
	4	000000	000010101000	000000111110	Same as 1
	5	000000	000001010100	000000111110 000000101000	Same as 1 $R = R - Div$
	6	000001	000000010101		$R \geq 0$ , shift 1 into Q Shift Div right
	7	<u>000010</u>	000000010101	<u>000000010100</u> 11 20	Same as 1

Result:  $62 \div 21 = 2 \dots 20$

6. (a)  $S=0$ ,  $\text{exp} = 00011000_2$ ,  $\text{frac} = 0$ ,  $\text{bias} = 127_{10}$

$$\text{num} = (-1)^S \times (1 + \text{frac}) \times 2^{(\text{exp} - \text{bias})} = (-1)^0 \times (1+0) \times 2^{-103} = \boxed{2^{-103} \approx 9.86 \times 10^{-32}}$$

(b)  $63.25_{10} = 11111.01_2 = 1.111101 \times 2^5$

$$S=0, \text{exp} = 5 - \text{bias} = 122 = \overset{10000100_2}{\cancel{0111010_2}}, \text{frac} = 1111010000000000000000_2$$

IEEE 754 representation is ~~0x3D7D0000~~ 0x427D0000