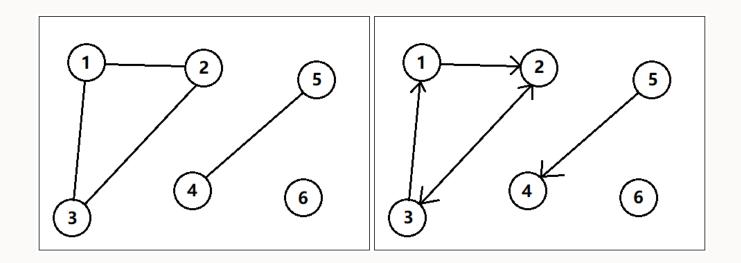
# Data Structures and Algorithm Analysis Lab 13, Graph

#### **Contents**

- . Implementing graphs.
- . Depth-first search and breadth-first search.

# Undirected graph and directed graph



The picture on the left is an undirected graph. The picture on the right is a directed graph.

Note that a tree is an undirected graph such that all node are connected and have no cycles.

# Represent a graph in JAVA

There are many ways to represent a graph in JAVA. Let's first define what a graph should contains.

```
public interface Graph {
    // return number of vertices.
    int size();

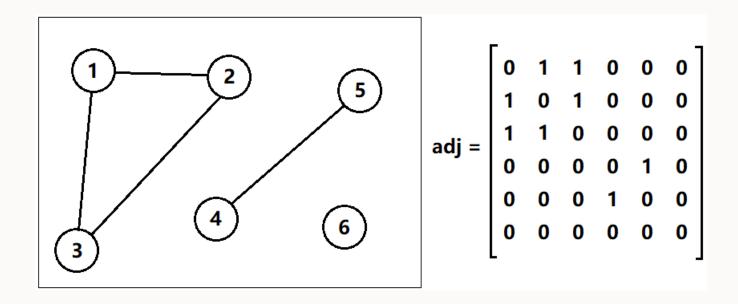
    // add an edge between v1 and v2, where 0 <= v1, v2 <=
        size()-1
    void addEdge( int v1, int v2 );

    // return all vertices that vertex v is connected to.
    Iterable < Integer > adjacency( int v );

    // return whether there is an edge from v1 to v2.
    boolean hasEdge( int v1, int v2 );
}
```

# Represent a graph in JAVA

Let's first implement a graph with adjacency matrix.



If there is a edge from v1 to v2, set adj[v1][v2] to 1.

# Graph with adjacency matrix

When implementing graph with adjacency matrix, we usually need to store number of vertices and an adjacency matrix.

```
public class GraphAdjacency implements Graph {
   private int num;
   private boolean[][] adj;

   public GraphAdjacency( int verticesNumber ) {
      num = verticesNumber;
      adj = new boolean[num][num];
   }
```

# **Graph with adjacency matrix**

```
. . .
 public int size() {
      return num;
 public void addEdge(int v1, int v2) {
      adj[v1][v2] = true;
 public Iterable < Integer > adjacency(int v) {
      LinkedList < Integer > list = new LinkedList <>();
      for( int i = 0; i < num; i ++ )</pre>
          if( adj[v][i] )
              list.add(i);
      return list;
 }
 public boolean hasEdge(int v1, int v2) {
      return adj[v1][v2];
 }
```

# Graph with adjacency matrix

#### Advantage:

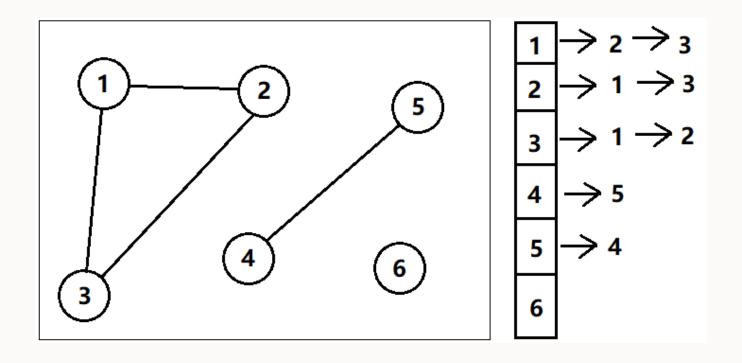
- . Very easy to implement.
- . Very fast to add a new edge.
- . Very fast to know whether two vertices are connected.

### Disadvantage:

- $N^2$  storage requirement.
- Linear time to know all vertices connected to a vertex.

In conclusion, it is very suitable to be used in a dense graph.

Let's implement a graph with adjacency list.



If there is a edge from v1 to v2, the list adj[v1] contains v2.

When using linked list to implement this graph, be aware of the generics.

```
public class GraphAdjList implements Graph {
  private int num;
  private LinkedList < Integer > [] adjList;
  @SuppressWarnings (value="unchecked")
  public GraphAdjList( int verticesNumber ) {
      num = verticesNumber;
      adjList = new LinkedList[num];
      for( int i = 0; i < adjList.length; i ++ )</pre>
        adjList[i] = new LinkedList <>();
  }
  public int size() {
      return num;
```

```
public void addEdge(int v1, int v2) {
   if( adjList[v1].contains(v2) )
      return;
   adjList[v1].add(v2);
}

public Iterable < Integer > adjacency(int v) {
   return (LinkedList < Integer >) adjList[v].clone();
}

public boolean hasEdge(int v1, int v2) {
   return adjList[v1].contains(v2);
}
```

#### Advantage:

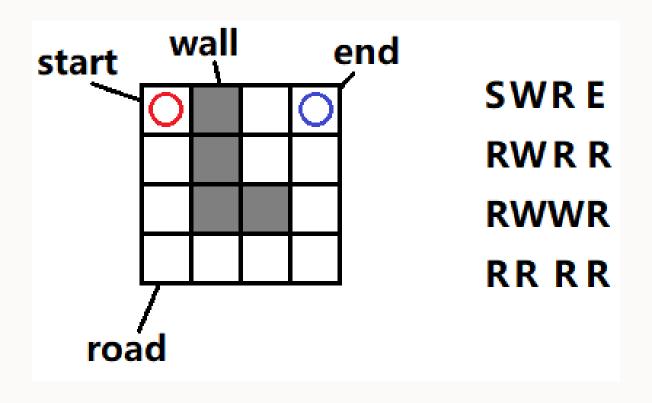
- . Using (Vertices+Edges) space.
- . Very fast to know all vertices connected to one vertex.

### Disadvantage:

. Need to iterate the list to know if two vertices are connected.

In conclusion, it is very suitable to be used in a sparse graph.

Let's solve a searching problem. We use characters to represent a maze.



Given a maze, can you tell me whether you can reach the "end" from the "start"?

- .Build a graph.
- . Depth-first search.
- . Breadth-first search.

```
public boolean solveMaze( char[][] maze ) {
   // ????
}
```

The first thing we need to do is to convert char[][] array into a graph. In this problem it is trivial. We just view each grid (or character) as a vertex. And there are edges between adjacent grids.

### **Depth-first search**

Now we have the graph, we can use depth-first search to search from the "start" vertex to every other vertices.

```
boolean dfs (Graph graph, int current, int end, boolean []
 visited ) {
  for( Integer adj : graph.adjacency(current) ) {
      if( visited[adj] )
          continue;
     visited[adj] = true;
      if( current == end )
          return true;
      if( dfs(graph, adj, end, visited) )
          return true;
 return false;
boolean dfs( Graph graph, int start, int end ) {
  boolean[] visited = new boolean[graph.size()];
 return dfs(graph, start, end, visited);
```

#### **Breadth-first search**

We can also use the breadth-first search.

```
boolean bfs( Graph graph, int start, int end ) {
  boolean[] visited = new boolean[graph.size()];
  Queue < Integer > queue = new LinkedList <>();
  visited[start] = true;
  queue.add(start);
  while( !queue.isEmpty() ) {
      int current = queue.poll();
      if( current == end )
          return true;
      for( int adj : graph.adjacency(current) ) {
          if( visited[adj] )
              continue;
          visited[adj] = true;
          queue.add(adj);
  return false;
```

We may not need to explicitly define a "graph" structure every time. We may just use the "graph" concept to write algorithms.

```
private static boolean bfs( char[][] maze ) {
  int height = maze.length;
  int width = maze[0].length;
  boolean[][] visited = new boolean[height][width];
  Queue < Point > queue = new LinkedList <>();
  // add start point in queue and set visited
  int[] dh = new int[] { -1, 1, 0, 0 };
  int[] dw = new int[] { 0, 0, -1, 1 };
  while( !queue.isEmpty() ) {
      Point current = queue.poll();
      int h = current.x;
      int w = current.y;
      if ( maze[h][w] == 'E')
          return true;
      for( int i = 0; i < 4; i ++ ) {</pre>
          int h2 = h+dh[i];
          int w2 = w + dw[i];
```

```
// if (h2,w2) is good add it into the queue
}

return false;
```