

NF
1st normalization (INF)

- 数据有原子性 (每行要有具有原子性的值，也要有主键)
- 没有重复数据组

2NF

- 所有列都是主键的一部分，对主键起到补充说明作用，且符合 INF
若符合 INF 的表只有一列主键，则也符合 2NF

3NF

- 符合 2NF，且没有传递函数依赖性。每一列只依赖主键。

SQL (Structured Query Language)

DDL, DML, DQL, DCL

↓
Data Control Language

1. Data Definition Language (DDL)

primary key

① `create table emp(
 id int primary key,
 name varchar(30)
)`

② `create table emp(
 id int,
 name varchar(30),
 constraint pk-con primary key(id)
)`

③ `create table emp(
 name varchar(30),
 dep varchar(20),
 salary numeric(9,2),
 primary key(name, dep)
)`

foreign key

① `create table emp(
 ...
 depId int references dep(id)
)`

② `create table emp(
 ...
 constraint fk-dep foreign key(depId) references dep(id)
)`

not null

`create table emp(
 ...
 name varchar(30) not null
)`

unique

① `create table emp(
 name varchar(30) unique
)`

② `— — (③ — — (`
 ...
 unique(name) ...
 unique(a, b) ...
)

default

`— — (...
 salary numeric(9,2) default 0.0
)`

constraint

① `price numeric constraint positive-price check (price > 0)`
② `— — (...
 check (price > 0)
)`

2. Data Manipulation Language (DML)

update
`insert into table (col1, col2, col3) values (..., ..., ...);
update table set col1 = ____ where ____;
delete from table where ____;`

drop `drop table table-name;`
alter `alter table Y add column name varchar(30);
alter table Y drop column name
(checks, foreign key, indexes 也可 drop)`

id	name	color
5	bear	w, y, blue
6	hat	g, y
9	kite	r, blue, g
10	yoyo	w, y

依然非 INF，
各列储存了相同分类
的数据

id	name	color1	color2	color3
5	bear	w	y	blue
6	hat	g	y	
9	kite	r	blue	g
10	yoyo			

Data Type

Text data types

char(len) 固定长度

varchar(len) 不定长度

clob 很长文本 (text 在 mysql)

Date types

date YYYY-MM-DD

datetime YYYY-MM-DD HH:mm:ss

timestamp 同上

范围：1970-01-01 00:00:01 UTC

到 2038-01-19 03:14:07 UTC

alter (cont.)

加列 alter table emp add column office varchar(30) [check (office <> '')];
 增列 alter table emp drop column deptId [cascade];
 加限制 alter table emp add primary key(id)
 {foreign key (deptId) references dept(id)
 constraint phone-unique unique (phone)}

删除限制 alter table emp drop constraint cons_name; alter table emp alter column col drop not null
 修改默认 alter table emp alter column salary set default 0.00
 alter table emp alter column salary drop default
 重命名 alter table emp rename column phone to tel_no
 alter table emp rename to employee

alter table emp alter column col set not null
 修改数据类型 alter table emp alter column salary type numeric(9,4);

Numerical

不是小数点后几位

int bigint 整型
 float(n) 用户指定精度,共n位
 real 固定6位有效
 numeric(p,d) 共p位,小数点后d位.
 Binary data types
 varbinary(max length)
 varbinary(max length)
 blob binary large object
 bytestream used in pgsql

- (2) ① in, not in
 ② like, %, _
 where name like '张%'
 %: 匹配0个、1个、多个 _: 匹配1个
 ③ null, not null
 where name is null
 where name is not null

Some functions

- ① upper(_), lower(_): 大/小写 ② trim(_): 去除头尾空格 ③ substr(_,-,-): 子串
 ④ replace(_,-,-): 替换 ⑤ length(_): 长度 ⑥ round, trunc:
 replace('sheep', 'ee', 'i') returns 'Ship'
 substr('Nanshan', 3, 4) returns 'nsha'
 round(3.141592, 3) -- 3.142
 trunc(3.141592, 3) -- 3.141
- ⑦ cast (col as type)

Join

1. Cross Join 笛卡尔积

select * from t1 cross join t2

$t_1 \text{ cross join } t_2$
 $t_1 \text{ inner join } t_2 \text{ on true}$

t_1, t_2 } are same

2. Inner Join

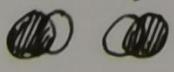


select * from $\left\{ \begin{array}{l} t_1 \text{ inner join } t_2 \text{ on } t_1.\text{num} = t_2.\text{num} \\ t_1 \text{ inner join } t_2 \text{ using (num)} \end{array} \right.$

inner 可省略

t_1, t_2 nature join 会把所有 t_1 和 t_2 公共的列加入 using (...) 中

3. Left/Right Outer Join



$t_1 \text{ left join } t_2 \text{ on } t_1.\text{num} = t_2.\text{num}$
 $t_1 \text{ left join } t_2 \text{ using (num)}$

4. Full Outer Join



Combining Query

query 1 union [all] query 2 并集
 query 1 intersect [all] query 2 交集
 query 1 except [all] query 2 差集
 加 all 不去重
 不加 all 等于加 distinct

About NULL

col + null
 col - null
 col * null
 col / null

col > null
 col = null
 col < null

true and null → null
 false and null → false
 true or null → true
 false or null → null
 col is null → true false

Ordering

用 case 排序

用 order by case

when role='A' then 1
 when role='B' then 2
 else 3
 end

排序规则:
 string 字典序
 number 大小

boolean false < true
 dates & times 越晚越大
 SQL Server, MySQL, null 最小
 Oracle, PostgreSQL null 最大

Window Function

① over () 开窗为全局

② over (partition by ... order by ...)

select country, title, runtime, year_release,
rank() over (order by year_release)
from movie where year_release > 2015;

③ select country, runtime, year_release,
avg(runtime) over (partition by country
order by year_release)
from movie where year_release > 2015;

country	runtime	y-r	avg
bd	160	2016	160
bd	140	2017	75.5
bd	0	2017	75.5
bd	2	2017	75.5

select rank() over (partition by country)
from movie where year_release > 2015;

select country, runtime, year_release,
avg(runtime) over (partition by country, year_release)
from movie where year_release > 2015;

country	runtime	y-r	avg
bd	160	2016	160
bd	140	2017	47.33...
bd	0	2017	47.33...
bd	2	2017	47.33...

④ Multi-Window

rank(): 1 2 2 2 5 5 7 8 8 8 11 12 12

dense_rank(): 1 2 2 2 3 3 4 5 5 5 6 7 7

row_number(): 1 2 3 4 5 6 7 8 9 10 11 12 13

⑤ lag(col, n, default) 填补空白(可无)
偏移量(可无)

select movieid, title, country, y-r, runtime,
sum(runtime) over w sum-all,

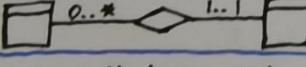
max(runtime) over w max,

avg(runtime) over w avg

from movie where y-r > 2010

window w as (partition by country order by y-r);

E-R Diagram

- ① "→" means "one"
 entity 实体
- ② "—" means "many"
 relation 关系
- ③ "—" 全参与, 全包含 
- ④ "L..h" L..h为参与的最小、最大 cardinality(unique 元素の个数)


① Weak entity(弱实体):

一个实体的存在依赖另一实体的存在为前提.



 [instr]: 表示1个学生有多个(0~...)辅导员

Relation Algebra

6 basic operator: select σ , project Π , union \cup , set difference $-$, Cartesian product \times , rename ρ

① $\sigma_p(r)$: 从表 r 中根据条件 p 选择

$\Pi_{A_1, A_2, \dots, A_k}(r)$: 从表 r 选出 A_1, A_2, \dots, A_k 列

$\sigma_{dept='phd'} \wedge salary < 90000$ (instructor) $\Pi_{id, name, salary}$ (instructor)

③ Join: $\sigma_{instr.id = teach.id}$ (instr x teach)

④ \cup union, \cap intersect, $-$ difference

另一种表示: $r \bowtie_s S = \sigma_{s: predicate}(r \times S)$ \bowtie : predicate(谓词)

$r \cup S$ $r \cap S$ $r - S$

⑤ " \leftarrow " Assignment(赋值)

⑥ Rename ρ

$Physics \leftarrow \sigma_{dept='phy'}(instr)$

$\rho_x(E)$: 给表达式 E 的结果命名为 x

$Music \leftarrow \sigma_{dept='music'}(instr)$

$\rho_{x(A_1, A_2, \dots, A_n)}(E)$: 也可以给列命名为 A_1, A_2, \dots, A_n

$Physics \cup Music$

Durability: 提交的事务由系统保存, 即使故障或重启,

事务也正常可用.

Transaction

ACID:

Atomicity: 要么全提交, 要么不提交.

Consistency: 要么创建一个新有效数据状态, 要么回滚

Isolation: 隔离性

dirty read: 读取其他事务修改未提交的数据 (脏)

nonrepeatable read: 重新读取它以前读的脏,

发现被已提交的事务修改

phantom read: 重新执行一个查询, 发现结果集

由于最近另一个事务的提交改变

serialization anomaly: 同一事务的结果可能因

每个事务提交顺序不同而结果不同.

A allowed, but not in pgsql

P possible

N not possible

default in pgsql

Isolation level | DR NR PR SA

Read uncommitted | A P P P

Read committed | N P P P

Repeatable read | N N A P

Serializable | N N N N

Serializability

一个 schedule(并发的)

等价于一个串行的 scheme

则它可序列化.

简化为只有 read, write

两个过程都 read(R), 则

其它三种

操作对

不冲突.

SQL about transaction:

begin; rollback to s1;

commit; show transaction-isolation;

rollback; set default_transaction-isolation = '—';

savepoint s1;

Function & Procedure

在pgsql里, procedure是返回值为void(没返回值)的函数
Function分类:
 - Query Language func (written in SQL)
 - Procedure Language func (pgsql/Tcl/Perl/Python)
 - Internal function
 - C-language function

2. PostgreSQL Func

```
create or replace function func-name(...) returns re-type
as $$  
declare  
var-name var-type := value;  
begin  
--;  
end  
$$ language plpgsql;
```

4. 返回值类型也可以是表
returns table
(
 col1 col-type,
 col2 col-type,
 ...)

1. SQL Func
create [or replace] function add (x integer, y integer)
returns integer as \$\$
\$1 select x+y; ← select \$1+\$2; ↑
language sql;
有输出, 则是: (x int, y int, out sum int)
必须有缺省
有默认值: (x int, y int default 2, z int default 3)

3. 条件语句

```
① if (...) then  
...;  
else if (...) then  
...;  
else  
...;  
end if;  
② for var-name in start..end loop  
...;  
end loop;  
③ for i in 1..num loop  
res = res * i;  
end loop;  
④ while condition loop  
...;  
end loop;
```

Trigger

Trigger Function返回值为trigger, 返回null或new, old等.

```
create trigger trigger-name  
before/after insert or update or delete (一个或多个)
```

on table-name

for each row/statement

```
execute procedure trigger-func();
```

4. Variable

	Data Type	说明
new	record	new为新的行, 若为row的I/U, new为null, 若为stat或row&D
old	record	old为旧的行, 若为row的U/D, old为null, 若为stat或row的I
TG_name	name	name of trigger
TG_when	text	before/after/instead of 定义
TG_level	text	row/statement 定义
TG_op	text	insert/update/delete/truncate

1. Trigger 可作用于:

1. tables

2. 级别: 1. row 每改变一行触发一次

2. views

3. foreign tables

执行一条语句触发1次, 即使语句改变了多行

3. before, after, instead of (级别)

	when event	row-level	statement-level
before	I/U/D	T, F	T, V, F
before	Trun	-	T
after	I/U/D	T, F	T
after	Trun	-	T

T: tables
V: views
F: foreign tables

5. Event trigger

```
create event trigger name  
on event-name
```

```
execute function func-name();
```

3.1: create event trigger tr-fun-ev on ddl-command_st
execute function event_trl();

Indexing

```
create index idx-name on table-name (col1, col2, ...);  
drop index [if exists] idx-name [cascade];
```

分类: 1. integrated index: PK index in MySQL InnoDB, SQL Server

external index: indexes in PostgreSQL, MySQL MyISAM

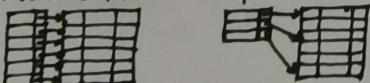
2. clustered index: primary index

non-clustered index: secondary index (have to be dense)

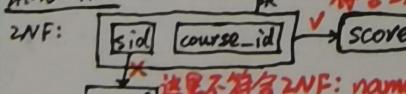
3. dense index: sparse index:

4. multi-key index

single-key index



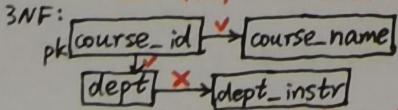
补充: NF



符合2NF: SCORE与pk都有关

* 这里不完全符合2NF: name只与pk的一部分 sid有关

符合1NF且为单列pk的, 一定符合2NF

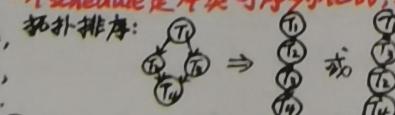


dept_instr P与dept有关, 与pk准直接关系, 只有传递关系

→ schedule S可通过交换不冲突操作得到 S', 则 S 与 S' 冲突等价.
S 冲突等价于一个串行 schedule, 则 S 冲突可序列化.

Precedence graph: 若 T_i, T_j 冲突, T_i 必须在 T_j 前, 则图 T_i → T_j

→ schedule 是冲突可序列化的, 如果 precedence graph 是无环的.

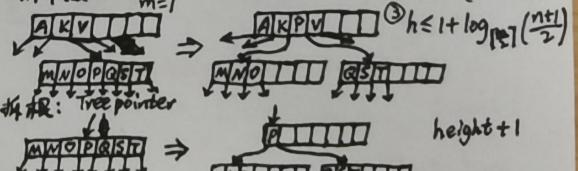


操作对数不同, 不冲突.

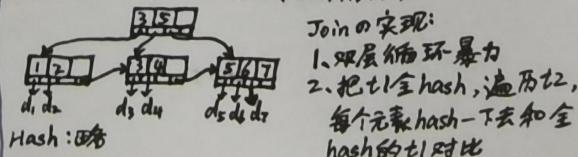
Recoverable schedule: T_j 读取 T_i 之前写入的数据, 则 T_i 的提交在 T_j 之前.

B-Tree 实现

① # of children = # of keys + 1
② $m \geq \lceil m/2 \rceil \leq \# \text{ of children} \leq m$ (root 除外)



B+ Tree: 数据只存叶子节点, 有后继



Join 的实现:

1. 双层循环暴力
2. 把 t1 全 hash, 遍历 t2, 每个元素 hash 下去和全 hash 的 t1 对比

3. sort-merge join
排序, 然后双指针扫描

With 用法

1. 表示一个表, 简化子查询

```
with t1 as (  
    select ...  
)  
select * from t1  
where ...
```

2. 表示迭代

```
with recursive t(n) as (  
    values (1)  
    union all  
    select n+1 from t where n < 100;  
)  
select sum(n) from t limit 100;
```