

CS201: Discrete Math for Computer Science
2024 Spring Semester Written Assignment #3
Due: Apr. 16th, 2023

The assignment needs to be written in English. Assignments in any other language will get zero point. Any plagiarism behavior will lead to zero point.

Q. 1. Compute the following without calculator and explain your answer.

(1) $(33^{15} \bmod 32)^3 \bmod 15$

(2) $\gcd(210, 1638)$

(3) $34x \equiv 77 \pmod{89}$

(4) The last decimal digit of 3^{1000} (Hint: Fermat's little theorem)

Solution:

(1) This is mainly computed based on Corollary 2 on page 242, i.e., $ab \bmod m = ((a \bmod m)(b \bmod m)) \bmod m$. It is perfectly fine if the student does not mention this corollary.

$$\begin{aligned} & (33^{15} \bmod 32)^3 \bmod 15 \\ &= ((33 \bmod 32)^{15} \bmod 32)^3 \bmod 15 \\ &= (1 \bmod 32)^3 \bmod 15 \\ &= 1 \bmod 15 \\ &= 1 \end{aligned}$$

(2) Using Euclidean Algorithm

$$\begin{aligned} 1638 &= 210 \cdot 7 + 168 \\ 210 &= 168 \cdot 1 + 42 \\ 168 &= 42 \cdot 4 \end{aligned}$$

Thus, $\gcd(210, 1638) = 42$.

- (3) Consider the inverse \bar{a} such that $\bar{a} \cdot 34 \equiv 1 \pmod{89}$. We use the extended Euclidean to solve \bar{a} . In particular,

$$\begin{aligned} 89 &= 34 \cdot 2 + 21 \\ 34 &= 21 \cdot 1 + 13 \\ 21 &= 13 \cdot 1 + 8 \\ 13 &= 8 \cdot 1 + 5 \\ 8 &= 5 \cdot 1 + 3 \\ 5 &= 3 \cdot 1 + 2 \\ 3 &= 2 \cdot 1 + 1 \\ 2 &= 1 \cdot 1 \end{aligned}$$

Thus,

$$\begin{aligned} 1 &= 3 - 2 \cdot 1 \\ &= 3 - (5 - 3 \cdot 1) \cdot 1 &= -5 \cdot 1 + 3 \cdot 2 \\ &= -5 \cdot 1 + (8 - 5 \cdot 1) \cdot 2 &= 8 \cdot 2 - 5 \cdot 3 \\ &= 8 \cdot 2 - (13 - 8 \cdot 1) \cdot 3 &= -13 \cdot 3 + 8 \cdot 5 \\ &= -13 \cdot 3 + (21 - 13 \cdot 1) \cdot 5 &= 21 \cdot 5 - 13 \cdot 8 \\ &= 21 \cdot 5 - (34 - 21 \cdot 1) \cdot 8 &= -34 \cdot 8 + 21 \cdot 13 \\ &= -34 \cdot 8 + (89 - 34 \cdot 2) \cdot 13 &= 89 \cdot 13 - 34 \cdot 34 \end{aligned}$$

Thus, we have $-34 \cdot 34 \pmod{89} = 1$, which implies that $55 \cdot 34 \pmod{89} = 1$. Thus, $\bar{a} = 55$. As a result, we have $x \equiv 55 \cdot 77 \pmod{89} \equiv 52 \pmod{89}$.

- (4) The last decimal digit of 3^{1000} is equivalent to computing $3^{1000} \pmod{10}$. By Fermat's little theorem, we have $3^4 \equiv 1 \pmod{5}$. Thus, $3^{1000} \equiv 3^{4 \times 250} \equiv 1 \pmod{5}$. In addition, $3^{1000} \equiv 1 \pmod{2}$, because 3^{1000} has only 3 as its factor and hence is an odd number. Then, since system $3^{1000} \equiv 1 \pmod{5}$ and $3^{1000} \equiv 1 \pmod{2}$ is equivalent to $3^{1000} \equiv 1 \pmod{10}$, we have $3^{1000} \pmod{10} = 1 \pmod{10} = 1$.

Q. 2. Use extended Euclidean algorithm to express $\gcd(561, 234)$ as a linear combination of 561 and 234.

Solution: By Euclidean algorithm, we have

$$\begin{aligned} 561 &= 2 \cdot 234 + 93 \\ 234 &= 2 \cdot 93 + 48 \\ 93 &= 1 \cdot 48 + 45 \\ 48 &= 1 \cdot 45 + 3. \end{aligned}$$

Thus, $\gcd(561, 234) = 3$. Accordingly, we can derive the linear combination:

$$\begin{aligned}
 3 &= 1 \cdot 48 - 1 \cdot 45 \\
 &= 1 \cdot 48 - 1 \cdot (93 - 48) \\
 &= 2 \cdot 48 - 1 \cdot 93 \\
 &= 2 \cdot (234 - 2 \cdot 93) - 1 \cdot 93 \\
 &= 2 \cdot 234 - 5 \cdot 93 \\
 &= 2 \cdot 234 - 5 \cdot (561 - 2 \cdot 234) \\
 &= 12 \cdot 234 - 5 \cdot 561.
 \end{aligned}$$

Q. 3. Let a , b , and c be integers. Suppose m is an integer greater than 1 and $ac \equiv bc \pmod{m}$. Prove $a \equiv b \pmod{m/\gcd(c, m)}$.

Solution: Let $m' = m/\gcd(c, m)$. Because all the common factors of m and c are divided out of m to obtain m' , it follows that m' and c are relatively prime. Since $ac \equiv bc \pmod{m}$, we have m divides $ac - bc = (a - b)c$, which follows that m' divides $(a - b)c$. Since m' and c are relatively prime, we see that m' divides $a - b$, which leads to $a \equiv b \pmod{m'}$.

Q. 4. For two integers a, b , suppose that $\gcd(a, b) = 1$ and $b \geq a$. Prove that $\gcd(b + a, b - a) \leq 2$.

Solution: Now suppose that $d|(b + a)$ and $d|(b - a)$. Then $d|(b + a) + (b - a) = 2b$ and $d|(b + a) - (b - a) = 2a$. Thus, $d|\gcd(2b, 2a) = 2\gcd(a, b) = 2$. Thus, $d \leq 2$ and so $\gcd(b + a, b - a) \leq 2$.

[Alternate solution.] Since $\gcd(b, a) = 1$, then by Bezout's identity, there exist integers s and t such that $sb + ta = 1$. This gives us

$$\begin{aligned}
 (s + t)(b + a) + (s - t)(b - a) &= sb + sa + tb + ta + sb - sa - tb + ta \\
 &= 2sb + 2ta \\
 &= 2,
 \end{aligned}$$

from which we conclude that $\gcd(b + a, b - a)$ cannot exceed 2.

Q. 5. Given an integer a , we say that a number n passes the "Fermat primality test (for base a)" if $a^{n-1} \equiv 1 \pmod{n}$.

(a) For $a = 2$, does $n = 561$ pass the test?

- (b) Did the test give the correct answer in this case?

Solution:

- (a) We have

$$\begin{aligned} 2^{560} &\equiv 2^{20 \cdot 28} \pmod{561} \\ &\equiv (2^{20})^{28} \pmod{561} \\ &\equiv (67)^{28} \pmod{561} \\ &\equiv (67^4)^7 \pmod{561} \\ &\equiv 1^7 \pmod{561} \\ &\equiv 1. \end{aligned}$$

Thus, $2^{560} \equiv 1 \pmod{561}$. So 561 passes the Fermat test with test value 2.

- (b) We have $561 = 3 \cdot 11 \cdot 17$. So, 561 is not a prime, and thus the test failed.

Q. 6. Solve the following linear congruence equations.

- (a) $778x \equiv 10 \pmod{379}$.
(b) $312x \equiv 3 \pmod{97}$.

Solution:

- (a) Note that 379 is a prime. To find the modular inverse of 778, we first apply Euclidean algorithm.

$$\begin{aligned} 778 &= 2 \cdot 379 + 20 \\ 379 &= 18 \cdot 20 + 19 \\ 20 &= 1 \cdot 19 + 1. \end{aligned}$$

Reading backwards we have $1 = 19 \cdot 778 - 39 \cdot 379$. Thus, we have $x \equiv 10 \cdot 10 \equiv 190 \pmod{379}$.

(b) Applying Euclidean algorithm, we have

$$\begin{aligned}312 &= 3 \cdot 97 + 21 \\97 &= 4 \cdot 21 + 13 \\21 &= 1 \cdot 13 + 8 \\13 &= 1 \cdot 8 + 5 \\8 &= 1 \cdot 5 + 3 \\5 &= 1 \cdot 3 + 2 \\3 &= 1 \cdot 2 + 1.\end{aligned}$$

Reading Euclidean algorithm backwards we have $1 = 37 \cdot 312 - 119 \cdot 97$.
So, $312 \cdot 37 \equiv 1 \pmod{97}$. Thus, $x \equiv 37 \cdot 3 \equiv 111 \equiv 14 \pmod{97}$.

Q. 7. Find all solutions, if any, to the system of congruences $x \equiv 5 \pmod{6}$, $x \equiv 3 \pmod{10}$, and $x \equiv 8 \pmod{15}$.

Solution: We cannot apply the Chinese remainder theorem directly, since the moduli are not pairwise relatively prime. However, we can use the Chinese remainder theorem, translate these congruences into a set of congruences that together are equivalent to the given congruence. Since we want $x \equiv 5 \pmod{6}$, we must have $x \equiv 5 \equiv 1 \pmod{2}$ and $x \equiv 5 \equiv 2 \pmod{3}$. Similarly, from the second congruence we must have $x \equiv 1 \pmod{2}$ and $x \equiv 3 \pmod{5}$; and from the third congruence we must have $x \equiv 2 \pmod{3}$ and $x \equiv 3 \pmod{5}$. Since these six statements are consistent, we see that our system is equivalent to the system $x \equiv 1 \pmod{2}$, $x \equiv 2 \pmod{3}$, $x \equiv 3 \pmod{5}$. These can be solved using the Chinese remainder theorem to yield $x \equiv 23 \pmod{30}$. Therefore the solutions are all integers of the form $23 + 30k$, where k is an integer.

Q. 8. (a) Show that if n is an integer, then $n^2 \equiv 0$ or $1 \pmod{4}$.

(b) Use (a) to show that if m is a positive integer of the form $4k + 3$ for some nonnegative integer k , then m is not the sum of the squares of two integers.

Solution: There are two cases. If n is even, then $n = 2k$ for some integer k , so $n^2 = 4k^2$, which means that $n^2 \equiv 0 \pmod{4}$. If n is odd, then $n = 2k + 1$ for some integer k , so $n^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$, which means that $n^2 \equiv 1 \pmod{4}$.

By (a), the sum of two squares must be either $0 + 0 = 0$, $0 + 1 = 1$, or $1 + 1 = 2$, modulo 4, never 3, and therefore not of the form $4k + 3$.

Q. 9. Prove that if a and m are positive integers such that $\gcd(a, m) \neq 1$ then a does not have an inverse modulo m .

Solution: We prove this by contrapositive. Assume that a has an inverse modulo m , i.e., there exists an integer b such that

$$ab \equiv 1 \pmod{m}.$$

This is equivalent to $m \mid (ab - 1)$, which means that there is an integer k such that

$$ab - 1 = mk,$$

which is

$$ba + (-k)m = 1.$$

Suppose that d is any common divisor of a and m , i.e., $d \mid a$ and $d \mid m$. Since b and k are integers, it follows that $d \mid (ba - km)$, so $d \mid 1$. Thus, we must have $d = 1$, which completes the proof.

Q. 10. Find counterexamples to each of these statements about congruences.

- (a) If $ac \equiv bc \pmod{m}$, where a, b, c , and m are integers with $m \geq 2$, then $a \equiv b \pmod{m}$.
- (b) If $a \equiv b \pmod{m}$ and $c \equiv d \pmod{m}$, where a, b, c, d , and m are integers with c and d positive and $m \geq 2$, then $a^c \equiv b^d \pmod{m}$.

Solution:

- (a) Let $m = c = 2$, $a = 0$ and $b = 1$. Then $0 = ac \equiv bc = 2 \pmod{2}$, but $0 = a \not\equiv b = 1 \pmod{2}$.
- (b) Let $m = 5$, $a = b = 3$, $c = 1$, and $d = 6$. Then $3 \equiv 3 \pmod{5}$ and $1 \equiv 6 \pmod{5}$, but $3^1 = 3 \not\equiv 4 \equiv 3^6 = 729 \pmod{5}$.

Q. 11. Show that we can easily factor n when we know that n is the product of two primes, p and q , and we know the value of $(p - 1)(q - 1)$.

Solution: Suppose that we know both $n = pq$ and $(p - 1)(q - 1)$. To find p and q , first note that $(p - 1)(q - 1) = pq - p - q + 1 = n - (p + q) + 1$. From this we can find $s = p + q$. Then with $n = pq$, we can use the quadratic formula to find p and q .

Q. 12. Consider the RSA encryption method. Let our public key be $(n, e) = (65, 7)$, and our private key be d .

- (a) What is the encryption \hat{M} of a message $M = 8$?
- (b) To decrypt, what value d do we need to use?
- (c) Using d , run the RSA decryption method on \hat{M} .

Solution:

- (a) To encrypt $M = 8$, we have

$$\begin{aligned}
 \hat{M} &= M^e \bmod n \\
 &= 8^7 \bmod 65 \\
 &= 8^{2 \cdot 3 + 1} \bmod 65 \\
 &= 64^3 \cdot 8 \bmod 65 \\
 &= (-1)^3 \cdot 8 \bmod 65 \\
 &= -8 \bmod 65 \\
 &= 57 \bmod 65.
 \end{aligned}$$

So the encrypted message is $\hat{M} = 57$.

- (b) From $n = 65 = 5 \times 13$, we have $(p-1)(q-1) = 48$. Recall we can find d by running Euclidean algorithm.

$$\begin{aligned}
 \gcd(\phi(n), e) &= \gcd(48, 7) \\
 &= \gcd(7, 6) \quad \text{as } 48 = 6 \cdot 7 + 6 \\
 &= \gcd(6, 1) \quad \text{as } 7 = 1 \cdot 6 + 1 \\
 &= 1.
 \end{aligned}$$

Thus $d = \gcd(48, 7) = 1$. Reading backwards we get $1 = 7 \cdot 7 - 1 \cdot 48$. Then the private key $d = 7$.

- (c) To complete the RSA decryption, we calculate

$$\begin{aligned}
 \hat{M}^d \bmod n &= 57^7 \bmod 65 \\
 &= (-8)^7 \bmod 65 \\
 &= (-8)^{2 \cdot 3 + 1} \bmod 65 \\
 &= (64)^3 \cdot (-8) \bmod 65 \\
 &= 8 \bmod 65.
 \end{aligned}$$

Therefore, the original message is $M = 8$ as desired.