

INDIAN INSTITUTE OF TECHNOLOGY INDORE

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# Report for Lab Assignment 1

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Parallel Computing Lab (CS 359)

*Author*

Aniket Sangwan (cse180001005@iiti.ac.in)

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# Chapter 1

## Report

### 1.1 Problem Statement

Write a parallel program to parse a string of symbols. The inputs are a context-free grammar  $G$  in Chomsky Normal Form and a string of symbols. In the end, the program should print yes if the string of symbols can be derived by the rules of the grammar and no otherwise. Write a sequential program (no OpenMP directives at all) as well. Compare the running time of the sequential program with the running time of the parallel program and compute the speedup for different grammars and different string lengths.

### 1.2 Approach

Cocke–Younger–Kasami (CYK) Algorithm is a dynamic programming based algorithm for parsing strings for a given context free grammar in Chomsky Normal Form. The algorithm is based on the principle that the solution to problem  $[i, j]$  can be constructed from solution to subproblem  $[i, k]$  and solution to subproblem  $[k, j]$ . It uses the result of previously parsed substring. Thus, it is based on DP.

- We can parallelize the substrings of same length as they only depend on previously parsed substrings, which are always shorter.
- We can also parallelize every loop iterating through all the productions because they are independent of each other.

### 1.3 Input

- The first line of input contains the number of productions  $n$  in the given CNF grammar.
- The next  $n$  contains each production in a new line. For Example:

4  
S  $\rightarrow$  AC | AB  
C  $\rightarrow$  SB

A -> a

B -> b

aabb.

- The third line contains the string to be parsed.

## 1.4 Output

- The first line states whether the string can be parsed or not.
- The second line outputs the time taken by program.
- You can view the parsing table using `print_table()` function. ( Uncomment the line 206 in the OPENMP code ).

## 1.5 Complexity

- The complexity of the serial program is  $O(n^3 * |G|)$ , where  $n$  is the length of the string to be parsed and  $|G|$  is the length of the grammar.
- The complexity of the parallel program (OPENMP) is  $O(n^3 * |G|/\text{no\_threads})$ , where  $n$  is the length of the string to be parsed and  $|G|$  is the length of the grammar.

## 1.6 Runtime Comparsion

To evaluate the program for different cases, I used **three different grammars** (available in the files g1.txt, g2.txt and g3.txt) and generated **three different test files of length 200, 800 and 1500**. I compared the runtime of each test case for serial and parallel programs and generated graphs to visualise the improvement.

### 1.6.1 Grammar 1 (g1.txt)

The grammar used is:

S->AC|AB

C->SB

A->a

B->b.

The images below show the input and output for above grammar using different number of threads.

```
> g++ A1_serial.cpp
~/Documents/5_sem/cs309-parallel/labs/Lab1 .....%
> ./a.out
4
S->AC|AB
C->SB
A->a
B->b
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
String can be parsed using this grammar
Time Taken by serial program: 0.28346 s
```

Figure 1.1: Serial Program

```

> g++ A1_openmp.cpp -fopenmp

~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 2
4
S->AC|AB
C->SB
A->a
B->b
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
String can be parsed using this grammar
Time Taken by openmp program: 0.12829 s

~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 3
4
S->AC|AB
C->SB
A->a
B->b
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
String can be parsed using this grammar
Time Taken by openmp program: 0.11161 s

~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 4
4
S->AC|AB
C->SB
A->a
B->b
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
String can be parsed using this grammar
Time Taken by openmp program: 0.09739 s

```

Figure 1.2: OPENMP Program

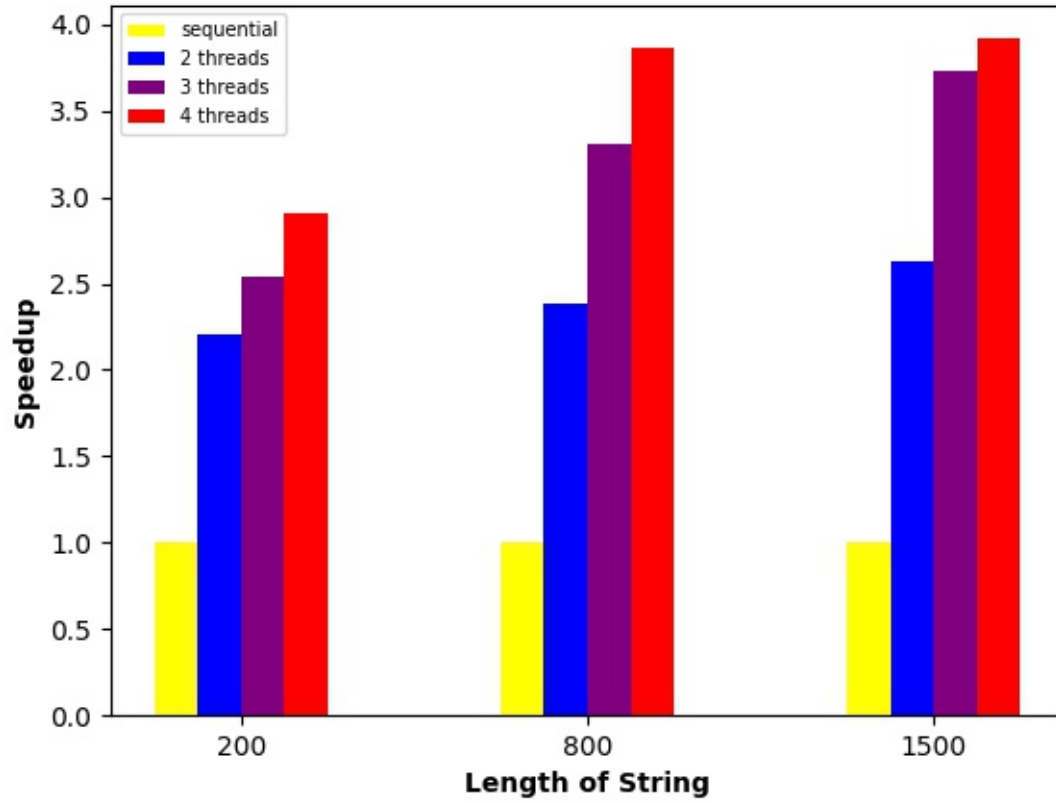


Figure 1.3: Speedup vs Length of string for Grammar 1

The below table displays the runtime comparison for different tests.

String Length	Time (in sec)			
	Sequential	2 Threads	3 Threads	4 Threads
200	0.28346	0.12829	0.11161	0.09739
800	10.97589	4.60321	3.32380	2.84105
1500	94.29656	35.81603	25.27589	24.07729

Table 1.1: Run-Time

### 1.6.2 Grammar 2 (g2.txt)

The grammar used is:

$S \rightarrow BA \mid DC \mid c$

$A \rightarrow a$

$B \rightarrow AS$

$C \rightarrow b$

$D \rightarrow CS$

The images below show the input and output for above grammar using different number of threads.

```
> g++ A1_serial.cpp
~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out
5
S -> BA | DC | c
A -> a
B -> AS
C -> b
D -> CS
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbcbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
String can be parsed using this grammar
Time Taken by serial program: 0.28552 s
```

Figure 1.4: Serial Program

```
~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> g++ A1_openmp.cpp -fopenmp
~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 2
5
S -> BA | DC | c
A -> a
B -> AS
C -> b
D -> CS
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbcbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
String can be parsed using this grammar
Time Taken by openmp program: 0.21176 s

~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 3
5
S -> BA | DC | c
A -> a
B -> AS
C -> b
D -> CS
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbcbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
String can be parsed using this grammar
Time Taken by openmp program: 0.11971 s

~/Documents/5_sem/cs309-parallel/labs/Lab1 ..... %
> ./a.out 4
5
S -> BA | DC | c
A -> a
B -> AS
C -> b
D -> CS
aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaabbbbbbbbbbbbbbbbbbbbbbb
bbbbbbbbbbbbbbbbbbbbbbbbbbbbcbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbbb
bbbaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
String can be parsed using this grammar
Time Taken by openmp program: 0.10752 s
```

Figure 1.5: OPENMP Program



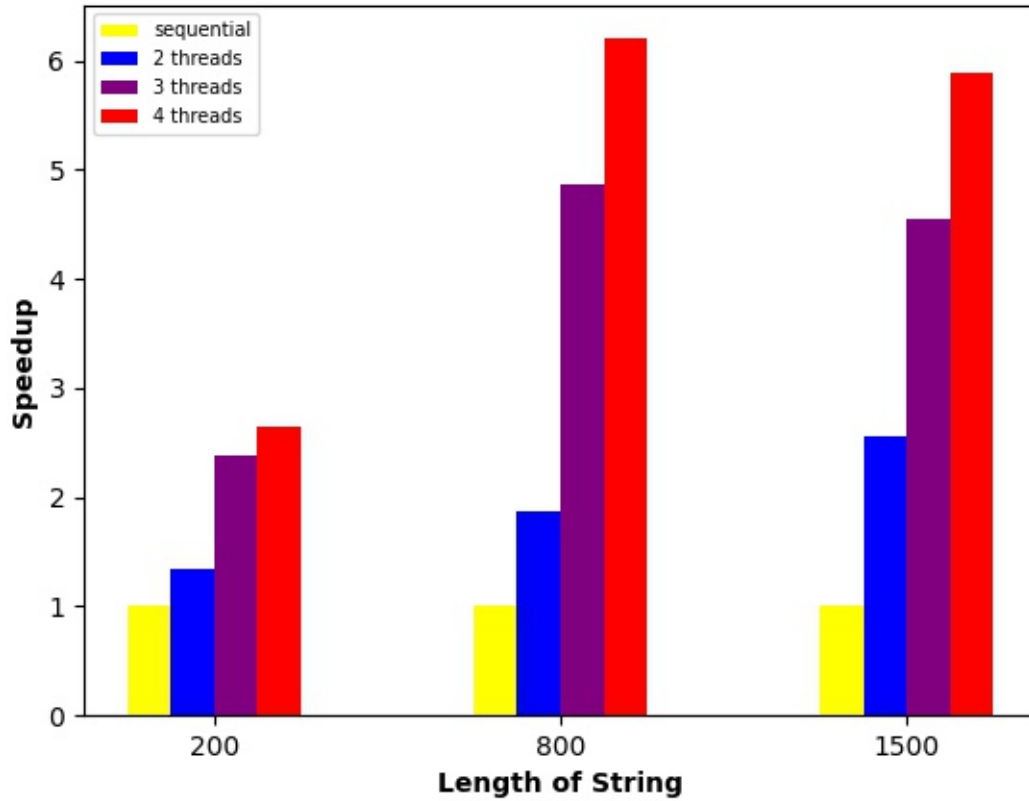


Figure 1.6: Speedup vs Length of string for Grammar 2

The below table displays the runtime comparison for different tests.

String Length	Time (in sec)			
	Sequential	2 Threads	3 Threads	4 Threads
200	0.28852	0.21176	0.11971	0.10752
800	14.18116	7.60084	2.90939	2.28607
1500	96.42192	37.58093	21.19945	16.35429

Table 1.2: Run-Time

### 1.6.3 Grammar 3 (g3.txt)

The grammar used is:

$S \rightarrow AB$

$A \rightarrow a$

$B \rightarrow a$

The images below show the input and output for above grammar using different number of threads.



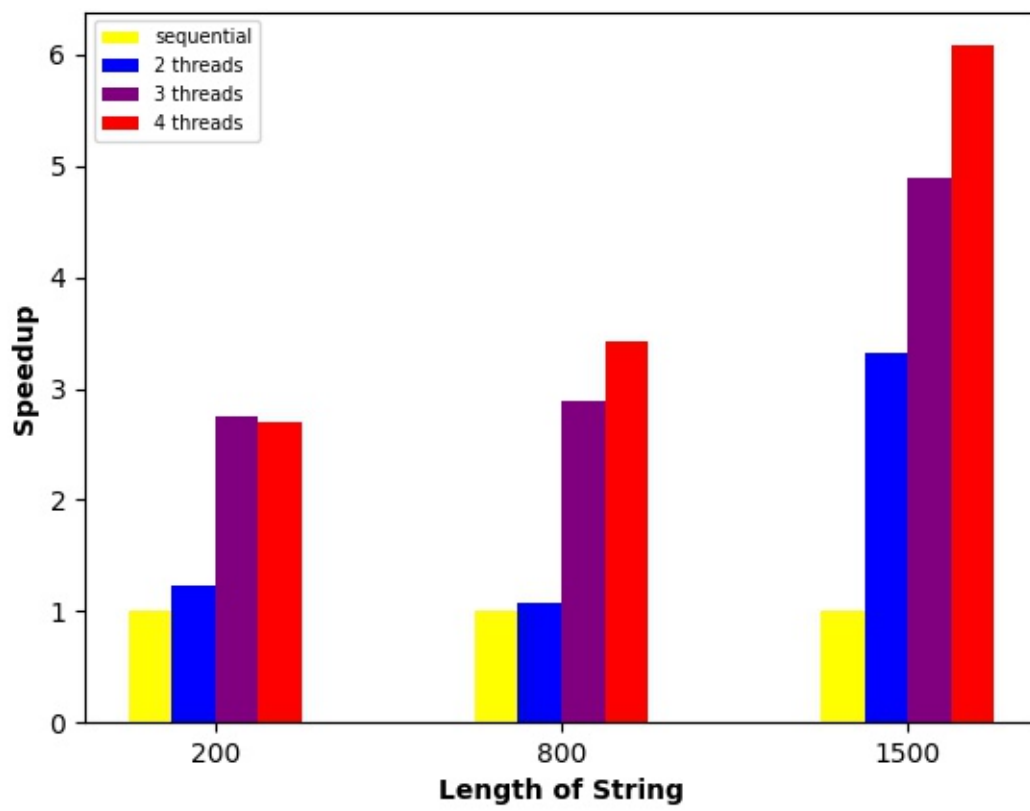


Figure 1.9: Speedup vs Length of string for Grammar 2

## 1.6.4 Conclusion

We conclude that the runtime reduces with a increase in the number of threads. The speedup also increases with increase in input size. Hence, this algorithm performs better for larger test cases.

## 1.6.5 Implementation

```
1 // This code is inspired from https://github.com/ahmadshafique/CYK-Parser/  
  blob/master/CYK%20parser.cpp  
2  
3 #include<bits/stdc++.h>  
4 #include <omp.h>  
5  
6 using namespace std;  
7  
8 #define pb push_back  
9 #define endl "\n"  
10  
11 #define N 1600  
12  
13 string grammar[N][N]; // Stores the grammar provided by the user  
14 string temp[N];  
15 string parse_matrix[N][N];  
16 int p,no_prod; // no_prod stores the no of productions  
17 int size; // No of threads  
18  
19 void break_grammar(string a){ //Stores the RHS of grammar in the string  
  array temp. Used to separate productions.  
20  
21  int i;  
22  p=0;  
23  while(a.size()){  
24  
25    i=a.find(" ");  
26  
27    int w=i;  
28    while(a[w-1]!=' ')w--; // Ignores space  
29  
30    if(w>a.size()){  
31  
32      temp[p++] = a;  
33      a="";  
34    }  
35    else{  
36  
37      temp[p++] = a.substr(0,w);  
38      w=i;  
39      while(a[w+1]!=' ')w++; // Ignores space  
40      a=a.substr(w+1,a.size());  
41    }  
42  }  
43 }  
44  
45 string join(string a,string b){ //Joins different non terminals to get new  
  parsing matrix value  
46
```

```

47     int i;
48     string curr=a;
49     for(i=0;i<b.size();i++)
50         if(curr.find(b[i]) > curr.size())
51             curr+=b[i];
52     return (curr);
53 }
54
55 string check_red(string p){ //Iterates through complete grammar to check if
56     the string p can be reduced or not
57
58     int j,k;
59     string curr1="";
60     #pragma omp parallel num_threads(size)
61     {
62         string curr="";
63
64         #pragma omp for // Parallelizes the next loop
65         for(j=0;j<no_prod;j++){
66
67             int k=1;
68
69             while(grammar[j][k]!=""){
70
71                 if(grammar[j][k]==p){
72                     curr=join(curr,grammar[j][0]);
73                 }
74
75                 k++;
76             }
77         }
78         #pragma omp critical // To prevent concurrent writes
79         curr1=join(curr1,curr);
80     }
81     return curr1;
82 }
83
84 string get_all_comb(string a, string b){ //Creates every possible
85     combination of strings a and b
86
87     int i,j;
88     string templ=a, re="";
89     for(i=0;i<a.size();i++)
90         for(j=0;j<b.size();j++){
91
92             templ="";
93             templ=templ+a[i]+b[j];
94             re=re+check_red(templ); //Checks if the current string can be
95             obtained or not
96         }
97     return re;
98 }
99
100 void print_table(string str){ // Prints the parsing table
101     for(int i=0;i<str.size();i++)
102     {
103         int k=0;

```

```

102     int l=str.size()-i-1;
103     for(int j=1;j<str.size();j++)
104     {
105         if(parse_matrix[k][j]!="")
106             cout<<parse_matrix[k][j]<<"\t";
107         else cout<<"phi\t";
108         k++;
109     }
110     cout<<endl;
111 }
112 }
113
114 int main(int argc, char **argv){
115
116     int i,u,j,l,k;
117     string a,str,curr,temp1,start;
118     size = atoi(argv[1]);
119
120     start="S";
121
122     // cout<<"\nEnter the number of productions\n";
123     // cout<<"Also enter each production in the format (S->AB|CD|a)\n";
124     cin >> no_prod;
125     cin.ignore();
126     for(i=0;i<no_prod;i++){
127
128         getline(cin,a);
129
130         u=a.find(">");
131         int w=u;
132
133         while(a[w-1]==' ')w--; // Ignores space in grammar
134
135         grammar[i][0] = a.substr(0,w);
136
137         w=u+1;
138         while(a[w+1]==' ')w++; // Ignores space in grammar
139
140         a = a.substr(w+1, a.size());
141
142         break_grammar(a);
143         for(j=0;j<p;j++)
144         {
145             grammar[i][j+1]=temp[j];
146         }
147     }
148
149     string st;
150
151     // cout<<"\nEnter the string of symbols you need to check: ";
152     getline(cin,str);
153
154     double start_time=omp_get_wtime();
155
156     for(i=0;i<str.size();i++){ //Assigns values to principal diagonal of
157         matrix
158         st = "";

```

```

159     st+=str[i];
160
161     #pragma omp parallel num_threads(size)
162     {
163         string curr="";
164
165         #pragma omp for // Parallelizes the next loop
166
167         for(j=0;j<no_prod;j++){
168
169             int k=1;
170             while(grammar[j][k] != ""){
171
172                 if(grammar[j][k] == st){
173
174                     curr=join(curr,grammar[j][0]);
175                 }
176                 k++;
177             }
178         }
179         #pragma omp critical // To prevent concurrent writes
180         parse_matrix[i][i]+=curr;
181     }
182 }
183
184 for(k=1;k<str.size();k++){ //Assigns values to upper half of the matrix
185
186     for(j=k;j<str.size();j++){
187
188         #pragma omp parallel num_threads(size)
189         {
190             string curr="";
191
192             #pragma omp for // Parallelizes the next loop
193
194             for(l=j-k;l<j;l++){
195
196                 string temp1 = get_all_comb(parse_matrix[j-k][l],parse_matrix[l
197 +1][j]);
198                 curr = join(curr,temp1);
199             }
200             #pragma omp critical // To prevent concurrent writes
201             parse_matrix[j-k][j]+= curr;
202         }
203     }
204
205     // Prints the Parsing Table
206     // print_table(str);
207
208     int f=0;
209
210     if(parse_matrix[0][str.size()-1].find(start)<=parse_matrix[0][str.size()
211 -1].size()){ //Checks if last element of first row contains a Start
212         //variable
213         cout<<"String can be parsed using this grammar\n";
214     }else{
215         cout<<"String cannot be parsed using this grammar\n";
216     }
217 }

```

```

214 }
215
216
217 double final_tot=omp_get_wtime()-start_time;
218
219 cout<<fixed<<setprecision(5) <<"Time Taken by openmp program: "<<
    final_tot <<" s "<<endl;
220
221 return 0;
222 }

```