INDIAN INSTITUTE OF TECHNOLOGY INDORE

Report for Lab Assignment 4 - RDA

Parallel Computing Lab (CS 359)

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Chapter 1

Report

1.1 Problem Statement

Write a parallel program to solve Travelling Salesman Problem using Red Deer algorithm. TSP is given as:

Given a list of cities and the distances between each pair of cities, what is the shortest possible route that visits each city and returns to the origin city?

- Each city needs to be visited exactly one time.
- We must return to the starting city, so our total distance needs to be calculated accordingly

1.2 Approach

TSP is a NP hard problem, i.e. no accurate solution can be found for this problem in polynomial time complexity. Red Deer algorithm is a population based meta heuristic algorithm which can be used to find near optimal solutions to difficult problems. It continues refining the search space to reach the minimum value while considering the exploration and exploitation phases satisfactorily.

The Red Deer algorithm for TSP proceeds in the following steps:

• Creating population: We select random numbers between 0 and 1 to proceed with our algorithm and generate 'popSize' number of individuals (Red Deer). Encoding scheme of meta-heuristic algorithms is used to generate city routes using these values. The loop generating RDs is parallelized.

- **Determining fitness:** We convert the current Red Deer to route using encoding scheme and calculate the distance of the route taken.
- Ranking Red Deers: We sort the population in the order of increasing fitness and return it in a new vector. We can **parallelize the two for loops** in this function which are used to calculate the fitness and to store it in ascending order.

```
#pragma omp parallel for
for (int i=0;i<population.size();i++){
    fitnessResults[i].F=routeFitness(population[i],index);
    fitnessResults[i].S=i;
}
```

```
#pragma omp parallel for
for(int i=0;i<fitnessResults.size();i++){

sorted_population[i]=population[fitnessResults[i].S];
fitness_scores[i]=(fitnessResults[i].F);

}
```

- Separating Male from Hinds: First, 'no_of_males' number of individuals with minimum distance are selected as males and the rest of the RDs are considered to be hinds. Notably, male RDs are the good solutions in this algorithm. Here the male RDs maintain the intensification properties, while hind RDs consider the diversification phase of the algorithm.
- Roaring male RDs: Male RDs try to increase their grace by roaring. We find the neighbors of the male RD, and if the objective functions of the neighbors are better than the male RD, we replace them with the prior ones. The following equation is used to update the position of males:

$$male_{new} = male_{old} + a_1 * ((UB - LB) * a_2) + LB : ifa_3 >= 0.5$$

 $male_{new} = male_{old} - a_1 * ((UB - LB) * a_2) + LB : ifa_3 >= 0.5,$

where a1, a2 and a3 are generated randomly by a uniform distribution between zero and one. If the new position is better than the previous one, the position is updated.

• Select commanders and stags: Male RDs are divided into two types, namely commanders and stags. The number of commander RDs are given by:

$$N_{comm} = \gamma * (N_{male})$$

$$N_{stags} = N_{male} - N_{comm}$$

• **Fight between commanders and stags:** Each commander fights with every stag and two new solutions are generated. The best among the commander and generated solutions is assigned the new commander. The new solutions are generated using following equations:

$$New_1 = (Comm + Stag)/2 + b_1 * ((UB - LB) * b_2) + LB$$

 $New_2 = (Comm + Stag)/2 - b_1 * ((UB - LB) * b_2) + LB,$

where b1 and b2 are generated randomly by a uniform distribution between zero and one.

- Forming harems: Here, we form the harems. A harem is a group of hinds in which a male commander seized them. The number of hinds in harems depends on the power of male commanders
- Mating process for commanders: Commanders mate with α percent of hinds in his harem and β percent of hinds in some randomly selected harem. The offspring is generated using the following formula:

$$offs = (Comm + Hind)/2 + (UB - LB) * c,$$

where c is generated randomly by a uniform distribution between zero and one. This process is **parallelized** using the below technique:

```
#pragma omp parallel
                vector < vector < double >> offs temp;
                for (int i=0; i < commanders. size(); i++){
                     for (int j=0; j<alpha*harems[i].size(); j++){
                         c=get rand();
                         vector < double > child (sz);
                         for (int k=0; k < sz; k++)
                              child [k]=(commanders [i] [k]+harems [i] [j] [k])
      /2 + (UB-LB) *c;
                         offs temp.pb(child);
                    }
13
               #pragma omp critical
16
                offs.insert(offs.end(),offs temp.begin(),offs temp.end());
17
           }
```

- Mating of stags with their nearest hind: Each stag mates with its closest hind. The offspring is generated using the same formula as above. This process is also parallelized using the same technique as above.
- Selecting the next generation: To select the next generation, two different strategies have been followed. In the first one, we keep all the male RD in the next generation. The second strategy refers to the remainder of the population in the

next generation. We choose hinds out of all hinds and offspring generated by mating process regarding the fitness value by using the roulette wheel mechanism.

```
#pragma omp parallel for reduction(+:tot_fitness)
for(int i=0;i<offs.size();i++){
    fitnessscores[i]=routeFitness(offs[i], index);
    tot_fitness+=fitnessscores[i];
}</pre>
```

```
#pragma omp parallel
           vector < vector < double >> nextGen temp;
           #pragma omp for nowait
           for(int i=0; i < pop size-males.size(); i++){
                double pick=get rand();
                pick*=100;
                for(int j=0; j< fitness scores. size(); j++){
                     if(j = fitnessscores.size()-1){
                         nextGen temp.pb(offs[j]);
                         break;
                     if (pick > 100-fitness scores[j]) {
                          // \text{ cout} << j << \text{endl};
                         nextGen_temp.pb( offs [j]);
                         break;
20
                     }
21
                }
22
23
24
           #pragma omp critical
25
           nextGen.insert(nextGen.end(),nextGen_temp.begin(),nextGen_temp
26
      . end());
```

1.3 Input

The input is provided as variables in the program.

• The user can select the number of cities, population size, number of males, alpha, beta, gamma and no of generations.

1.4 Output

• The first line outputs the initial total distance.

- The second line outputs the final minimized total distance of the solution.
- The final line outputs the time taken (in sec) by the program.

1.5 Runtime and Convergence Analysis

To evaluate the performance for different cases, I used **three different values for population sizes** (100, 200 and 400) while keeping number of cities = 40, Number of males=30, alpha=0.9, beta=0.4, gamma=0.7 and number of generations=1000. I compared the runtime of each test case for serial and parallel programs and generated graphs to visualise the performance.

1.5.1 Test Case 1 (Population size=100)

```
309-parallel/labs/Lab4
 g++ RDA serial.cpp
Initial Distance was: 4171.73
Final Distance is: 3143.08
Time Taken by serial program: 7.66850 s
 g++ RDA parallel.cpp -fopenmp
 ./a.out 2
Initial Distance was: 4235.09
Final Distance is: 3289.16
Time Taken by serial program: 5.35218 s
 ./a.out 3
Initial Distance was: 3956.24
Final Distance is: 3295.3
Time Taken by serial program: 6.08324 s
~/Documents/5 sem/cs309-parallel/labs/Lab4
 ./a.out 4
Initial Distance was: 4064.15
Final Distance is: 3237.65
Time Taken by serial program: 5.87276 s
```

Figure 1.1: Output when population=100

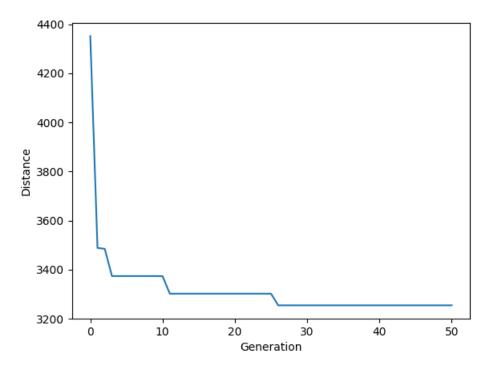


Figure 1.2: Convergence Graph when population=100

1.5.2 Test Case 2 (Population size=200)

```
309-parallel/labs/Lab4
  g++ RDA serial.cpp
Initial Distance was: 3428.08
Final Distance is: 3196.64
Time Taken by serial program: 13.31405 s
 g++ RDA parallel.cpp -fopenmp
 ./a.out 2
Initial Distance was: 3932.9
Final Distance is: 2987.16
Time Taken by serial program: 11.22058 s
 ./a.out 3
Initial Distance was: 4167.99
Final Distance is: 3019.71
Time Taken by serial program: 10.69063 s
 ./a.out 4
Initial Distance was: 4380.34
Final Distance is: 3229.53
Time Taken by serial program: 8.08117 s
```

Figure 1.3: Output when population=200

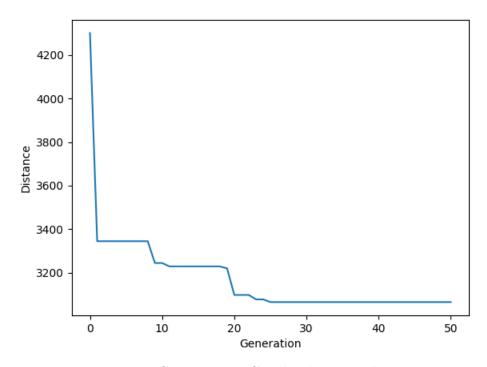


Figure 1.4: Convergence Graph when population=200

1.5.3 Test Case 3 (Population size=400)

```
g++ RDA_serial.cpp
Initial Distance was: 3940.33
Final Distance is: 3226.89
Time Taken by serial program: 25.39734 s
~/Documents/5 sem/cs309-parallel/labs/Lab4
 g++ RDA parallel.cpp -fopenmp
 ./a.out 2
Initial Distance was: 3843.28
Final Distance is: 3122.69
Time Taken by serial program: 20.90214 s
~/Documents/5 sem/cs309-parallel/labs/Lab4 .....
 ./a.out 3
Initial Distance was: 3836.46
Final Distance is: 3101
Time Taken by serial program: 19.68539 s
  ./a.out 4
Initial Distance was: 4168.96
Final Distance is: 3224.71
Time Taken by serial program: 19.88506 s
```

Figure 1.5: Output when population=400

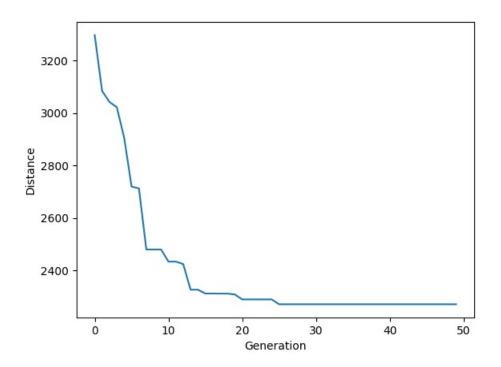


Figure 1.6: Convergence Graph when population=400

	Runtime (in sec)			
Population	1 Process	2 Processes	3 Processes	4 Processes
100	7.66850	5.35218	6.08324	5.87276
200	13.31405	11.22058	10.69063	8.08117
400	25.39734	20.90214	19.68539	19.88506

Table 1.1: Run-Times

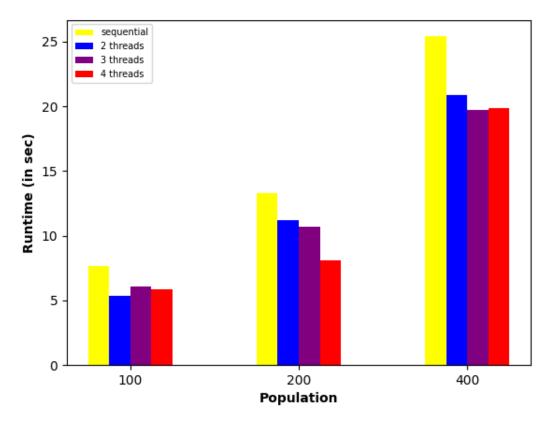


Figure 1.7: Change in runtime with change in population for serial and parallel programs

1.6 Conclusion

We conclude that the runtime usually reduces when the number of threads are increased while keeping population constant. The speedup is not as expected because of the precalculations and overhead involved in the algorithm. Along with this, there are certain steps in the algorithm that cannot be parallelized due to data dependency and critical regions. We also observe an increase in the speedup with increase in population size due to the dependency of parallelized loops on 'popSize'.

1.7 Implementation

```
Author: Aniket Sangwan (180001005)
5 #include < bits / stdc++.h>
6 | \# include < omp.h >
7 using namespace std;
9 #define double long double
10 #define pb push_back
#define F first
#define S second
14 const int inf=1e6;
16 mt19937 rnd(chrono::system clock::now().time since epoch().count());
const double RANDOM MAX = 4294967295;
19 // Returns random value
20 double get rand(){
     return (double)rnd() / RANDOM MAX;
22 }
23
24
  vector<double>fitness scores; // Stores the fitness values of every Red
25
     Deer
26
 double LB=-1, UB=1; // Lower and Upper Bound used in the algorithm
27
28
  // Structure for City in TSP problem
  struct City{
30
31
      int id; // Stores the id of the city (Unique for every city)
32
      int x,y; // Coordinates of the city
33
34
      // Calculates distance between two cities
35
      double distance (City c) {
36
37
           double xdis=abs(x-c.x);
38
           double ydis=abs(y-c.y);
39
           double distance=sqrt(xdis*xdis+ydis*ydis);
40
41
           return distance;
42
43
      }
45
      // Prints the coordinates of the city
46
      void print coord(){
47
          cout <<" ( "<< x<<" , "<< y<<" ) \n";
48
49
50
51 };
  // Calculates the total distance of the route taken
54 double routeDistance(vector<double>route, City index[]) {
```

```
double path dist=0, fitness=0.0;
56
57
        vector < double > sorted RD = route;
58
        sort(sorted_RD.begin(), sorted RD.end());
59
        map<double, int>RD to city;
60
61
        for (int i=0; i < sorted_RD.size(); i++){
62
             RD_to_city[sorted_RD[i]]=i;
63
64
65
        // TSP starts and ends at the same place. So the initial city is
66
       inserted again
        route.pb(route[0]);
67
68
        // Calculates pairwise distance
69
        for (int i=0; i < route.size()-1; i++){
70
71
             City from City = index [RD to city [route [i]]];
72
             City to City = index [RD_to_city [route[i+1]]];
73
74
             path dist+=fromCity.distance(toCity);
75
76
        }
77
78
        return path dist;
79
80
81
82
   /* Returns fitness value of given route.
84 We aim to minimize the distance. Here, I am trying to minimize the fitness
       too
   double routeFitness(vector<double>route, City index[]){
86
        double fitness=routeDistance(route, index);
        return fitness;
89
90
91
92
93
   // Randomly creates initial population of Red deer using randomizer
   vector < vector < double >> initialPopulation(int pop_size, int no_of_cities) {
95
96
        vector < vector < double >> RD list;
97
98
        RD list.resize(pop size);
99
100
        #pragma omp parallel for
101
        \begin{array}{ll} \textbf{for} \; (\; \textbf{int} \quad i = 0; i \! < \! pop\_size \; ; \; i \! + \! + \! ) \{ \\ \end{array}
             RD_list[i].resize(no_of_cities);
103
             \begin{tabular}{ll} for (int j = 0; j < no\_of\_cities; j + +) & \\ \end{tabular}
104
                  RD_list[i][j] = get_rand();
106
        }
107
108
        return RD list;
109
110
```

```
111 }
      Sorts the population in increasing fitness score and returns in a new
113
   vector < vector < double >> rankRDs (vector < vector < double >> population, City index
      []){
11
       vector < pair < double, int >> fitness Results; // Stores fitness value and
116
      inital route id
       fitnessResults.resize(population.size());
117
       fitness scores.resize(population.size());
118
119
       #pragma omp parallel for
12
       for (int i=0; i < population.size(); i++){
12
            fitnessResults[i].F=routeFitness(population[i],index);
122
            fitnessResults[i].S=i;
123
       }
124
125
       // Sorts in ascending order wrt fitness value
126
       sort(fitnessResults.begin(), fitnessResults.end());
127
12
       vector < vector < double >> sorted_population; // Will store population
129
      sorted in increasing fitness value
       sorted_population.resize(population.size());
130
       #pragma omp parallel for
       for(int i=0;i<fitnessResults.size();i++){</pre>
133
13
            sorted population[i]=population[fitnessResults[i].S];
            fitness scores[i]=(fitnessResults[i].F);
133
       }
138
139
       return sorted population;
140
141
142
143
      Separates males and hinds by separating 'no of males' Red deers with
144
      minimum fitness values
   vector < vector < double >>> Separate male and hinds (vector < vector < double
      >>population, int no of males, int no of hinds, City index[]) {
146
       vector < vector < double >>> Separate;
14
       Separate.resize(2);
14
14
       Separate [0]. resize (no_of_males);
       Separate [1]. resize (no of hinds);
15
152
       for (int i=0; i< no of males; i++)
153
            Separate [0][i]=population[i]; // Store males
156
       for (int i=no of males; i < population.size(); i++)
            Separate [1][i-no_of_males]=population[i]; // Stores Hinds
158
159
160
       return Separate;
161 }
162
```

```
/* Simulating the roaring process.
   Finds the neighbours of male RD and if their OF is better, we replace them
       with the male new
   */
165
   vector < vector < double >> roar (vector < vector < double >> males, City index []) {
166
16
        double a1=get rand(), a2=get rand(), a3=get rand();
169
        for (int i=0; i < males. size(); i++){
170
171
             vector < double > male new ( males [ i ]. size () );
172
17
             // Updating the position of males
17
             for (int j=0; j < males[i]. size(); j++){
17
                  if (a3>=0.5)
176
                      male\_new\,[\;j\,]\!=\!males\,[\;i\;]\,[\;j\,]\!+\!a1*(\,(U\!B\!-\!\!LB)*a2\,)\!+\!\!LB\,;
17
                 else
17
                      male \text{new}[j] = \text{males}[i][j] - a1 * ((UB-LB) * a2) + LB;
17
180
18
             if (routeFitness (male new, index) < routeFitness (males [i], index))
18
                 males [i] = male new;
183
184
185
        return males;
186
187
188
18
       Separates commanders and stags by separating gamma*(no of males) with
190
       minimum fitness into commanders
   vector < vector < double >>> Separate _ commanders _ and _ stags (vector < vector <
191
       double >> males, double gamma) {
192
        int no of commanders=gamma*males.size();
193
        vector<vector<double>>>Separate;
19
        Separate.resize(2);
19
19
        Separate [0]. resize (no of commanders);
191
        Separate [1]. resize (males. size () -no of commanders);
198
199
        for (int i=0; i< no of commanders; i++)
200
             Separate [0][i] = males [i];
20
20
        for (int i=no of commanders; i < males.size(); i++)
204
             Separate [1] [i-no of commanders] = males [i];
205
206
        return Separate;
   }
208
209
210
   Each commander fights with every stag and two new solutions are generated.
211
       The best among the
   commander and generated solutions is assigned the new commander
212
213 */
   void Fight(vector<vector<double>>>&commanders, vector<vector<double>>>stags,
214
       City index[]) {
215
```

```
216
                int no of stags=stags.size();
                int sz=commanders[0].size();
217
21
                for (int i=0; i<commanders. size(); i++){
219
22
                         for (int j=0; j < stags. size(); j++){
22
                                                                                                            // Two new solutions
                         vector < double > New1(sz), New2(sz);
222
                         double b1=get rand(), b2=get rand();
223
224
                                  for (int k=0; k < sz; k++){
225
                                           New1 [k] = (commanders [i] [k] + stags [j] [k]) /2.0 + b1*((UB-LB)*b2) +
226
             LB;
                                           New2 [k] = (commanders [i][k] + stags [j][k]) /2.0 - b1 * ((UB-LB) * b2) +
             LB;
                                  }
229
                                  // Calculating minimum fitness and
23
                                  double fitness1=routeFitness(commanders[i], index);
23
                                  double fitness2=routeFitness(New1, index);
233
                                  double fitness3=routeFitness(New2, index);
                                  double minfitness=min(fitness1, min(fitness2, fitness3));
23
235
                                  if (fitness2=minfitness) commanders [i]=New1;
236
                                  else if (fitness3=minfitness) commanders [i]=New2;
237
                         }
238
               }
239
240
241
242
             Forming harems. Hinds are divided proportionally among male commanders
243
              according to their power
      vector < vector < double >>> form\_harems(vector < double >>> & vector < double >>> & v
              commanders, vector < vector < double >> hinds, City index []) {
245
                double tot_fitness=0; // Stores total fitness
246
                shuffle(hinds.begin(), hinds.end(), std::default random engine(rand()));
24
24
                // Stores fitness values
249
                vector < double > fitness;
250
                fitness.resize(commanders.size());
25
252
                for (int i=0; i < commanders. size(); i++){
                         fitness [i]=routeFitness (commanders [i], index);
25
                         tot fitness+=fitness[i];
                }
256
251
258
                vector < vector < double >>> harems; // Stores the hinds in each harem
25
               harems.resize(commanders.size());
260
26
                int sz=commanders[0].size();
26
                int hinds_taken=0; // No of hinds assigned to harems
26
                int no_of_hinds=hinds.size(); // Total no of hinds
264
                                                                                                                                                 // No of hinds
               int harem_size=(fitness [0] / tot_fitness)*no_of_hinds;
265
              in next harem
266
                // Proportionally assigning harems
267
                for (int i=0; i < commanders. size(); i++){
268
```

```
harems[i].resize(harem size);
269
            for (int j=hinds taken; j<hinds taken+harem size; j++){
27
27
                 harems[i][j-hinds taken]=(hinds[j]);
27
27
            hinds taken+=harem size;
27
            if (i<commanders.size()-2)
276
                 harem_size = (fitness[i+1]/tot_fitness)*no_of_hinds;
27
            else
27
                 harem size=no of hinds-hinds taken;
279
        }
280
   return harems;
282
283
284
285
286
   Mating commander of a harem with alpha percent of hinds in his harem
287
   Mating commander of a harem with beta percent of hinds in another random
28
       harem
289
   vector\!<\!vector\!<\!double\!>\!> Mate\_alpha\_beta\big(vector\!<\!vector\!<\!double\!>\!> commanders\,,
290
       vector < vector < double >>> harems, double alpha, double beta) {
29
        int sz=commanders[0].size();
292
        int c=get rand();
293
        vector < vector < double >> offs;
29
29
       #pragma omp parallel
29
            vector < vector < double >> offs temp;
29
            for (int i=0; i < commanders. size(); i++)
29
                 for (int j=0; j<alpha*harems[i].size(); j++){
30
30
                      c=get rand();
                      vector < double > child (sz);
30
                      for (int k=0; k \le z; k++)
30
                           child[k] = (commanders[i][k] + harems[i][j][k])/2 + (UB-LB)*c
30!
306
                      offs_temp.pb(child);
30
                 }
30
            }
31
            #pragma omp critical
313
            offs.insert(offs.end(),offs temp.begin(),offs temp.end());
312
313
314
       #pragma omp parallel
315
31
            vector < vector < double >> offs temp;
317
            for (int i=0; i < commanders. size(); i++){
319
                 int w=rand()% harems. size(); // Randomly selected harem
320
321
                 while (w=i) w=rand ()%harems.size();
322
                 for (int j=0; j < beta*harems[w]. size(); j++){
323
```

```
c = get_rand();
324
                      vector < double > child (sz);
32
                      for (int k=0; k < sz; k++){
32
                          child [k]=(commanders [i][k]+harems [w][j][k])/2+(UB-LB)*c
32
32
                      offs temp.pb(child);
                 }
330
            }
331
332
            #pragma omp critical
333
            offs.insert(offs.end(),offs_temp.begin(),offs_temp.end());
334
       }
336
   return offs;
338
339
   // Mating stag with the nearest hind
340
   vector<vector<double>>> Mate stag hind(vector<vector<double>>>stags, vector<
341
       vector < double >> hinds, City index []) {
34
        vector < vector < double >> offs;
343
344
       // Finding hind with minimum distance from 'stags[i]' stag and mating
345
       with it
       #pragma omp parallel
346
347
            vector < vector < double >> offs temp;
34
            #pragma omp for nowait
            for (int i=0; i < stags.size(); i++){
35
                 int cindx = -1, cdist = inf;
35
                 for (int j=0; j<hinds. size (); j++){
35
35
                      vector < double > diff(stags[i].size());
35.
35
                      for (int k=0; k < stags[i]. size(); k++)
357
                           diff[k]=abs(stags[i][k]-hinds[j][k]);
358
359
                      int diff fitness=routeFitness(diff, index);
36
                      if (diff fitness < cdist)
36
                          cdist=diff fitness , cindx=j ;
360
                 }
36
                 // Generating offspring with the selected hind
366
                 int j=cindx;
367
                 double c=get rand();
36
                 vector < double > child (stags [i]. size ());
369
                 for (int k=0; k < stags[i]. size(); k++){}
370
                      child[k] = (stags[i][k] + hinds[j][k])/2 + (UB-LB)*c;
37
372
                 offs temp.pb(child);
373
374
            }
375
376
            #pragma omp critical
            offs.insert(offs.end(), offs temp.begin(), offs temp.end());
377
       }
378
```

```
379
   return offs;
380
381
38
   // Selecting the next generation using elitism and roulette wheel mechanism
38
   vector < vector < double >> NextGen (vector < vector < double >> offs, vector < vector <
38
       double>>males, int pop size, City index[]) {
38!
        vector < vector < double >> nextGen;
386
381
        // All the males are selected in the next generation
388
        for (int i=0; i < males. size(); i++){
389
            nextGen.pb(males[i]);
39
392
        vector < double > fitness scores;
393
        fitnessscores.resize(offs.size());
39
39
        double tot fitness=0; // Stores total fitness
396
39
       #pragma omp parallel for reduction (+:tot fitness)
39
        for (int i=0; i < offs. size(); i++){
39
             fitness scores [i] = route Fitness (offs [i], index);
400
401
             tot_fitness+=fitnessscores[i];
        }
402
403
        // Assigning fitness weighed probability
404
        for (int i=0; i < fitnessscores.size(); i++){
40
40
             fitnessscores [i]=100*((fitnessscores [i]/tot_fitness));
40
             if(i!=0) fitnessscores[i]+=fitnessscores[i-1];
408
409
410
        }
41
        // Using fitness proportionate selection (Roulette Wheel mechanism)
412
       #pragma omp parallel
41
414
             vector < vector < double >> nextGen temp;
413
416
            #pragma omp for nowait
417
             for (int i=0; i < pop size-males. size(); i++){
418
419
                 double pick=get rand();
42
                 pick*=100;
42
                 for (int j=0; j< fitness scores. size(); j++){
422
423
                      if (j = fitnessscores. size()-1){
424
                           nextGen temp.pb(offs[j]);
42
                           break;
426
42"
                      if (pick > 100-fitnessscores[j]) {
42
                           // \text{ cout} << j << \text{endl};
429
                           nextGen temp.pb(offs[j]);
430
                           break;
431
                      }
432
433
                 }
434
            }
435
```

```
436
           #pragma omp critical
43
            nextGen.insert(nextGen.end(),nextGen temp.begin(),nextGen temp.end
438
       ());
43
44
44
   return nextGen;
442
443 }
444
   // Using all the above functions to get next generation
445
   vector < vector < double >> nextGeneration (vector < vector < double >> population, int
446
        popSize, int no_of_males, double alpha, double beta, double gamma, City
        index []) {
447
       vector < vector < double >> popRanked = rankRDs (population, index);
448
44
       // cout << routeFitness (popRanked [0], index) << ", ";
450
45
       vector < vector < double >>> separate = Separate and hinds (
45
       popRanked, no of males, popSize-no of males, index);
       vector < vector < double >> Males = separate [0];
453
       vector < vector < double >> Hinds = separate [1];
454
455
       Males=roar (Males, index);
456
       Males=rankRDs (Males, index);
457
458
       separate=Separate commanders and stags (Males, gamma);
45
       vector < vector < double >> Commanders = separate [0];
46
       vector < vector < double >> Stags = separate [1];
46
465
       Fight (Commanders, Stags, index);
46
46
       vector < vector < double >>> harems = form harems (Commanders, Hinds,
465
       index);
46
       vector < vector < double >> offs = Mate alpha beta (Commanders, harems, alpha,
46
       beta):
468
       vector < vector < double >> offs 1=Mate stag hind (Stags, Hinds, index);
469
470
       offs.insert(offs.end(),offs1.begin(),offs1.end());
47
47
       offs.insert(offs.end(),Commanders.begin(),Commanders.end());
       offs.insert(offs.end(),Stags.begin(),Stags.end());
47
       offs.insert(offs.end(), Hinds.begin(), Hinds.end());
47
476
       vector < vector < double >> nextGen=NextGen( offs , Males , popSize , index );
47
47
   return nextGen;
479
48
48
   void RD_Algo( vector<City>cityList , int popSize , int no_of_males , int
482
       generations, double alpha, double beta, double gamma, City index[]) {
483
       // Creating initial population from city list
484
       vector < vector < double >> population;
485
486
```

```
population=initialPopulation(popSize, cityList.size());
487
488
       cout << "Initial Distance was: " << route Distance (population [0], index) <<
489
       endl;
       for (int i=0; i < generations; i++){
49
492
            fitness_scores.clear();
495
            population=nextGeneration(population, popSize, no of males, alpha,
494
       beta, gamma, index);
495
       }
496
       cout << "Final Distance is: "<< routeDistance (population [0], index) << endl;
49
498
499
500
   int main(int argc, char **argv){
501
       ios_base::sync_with_stdio(false); cin.tie(NULL); cout.tie(NULL);
502
503
       omp set num threads(atoi(argv[1]));
50
50
       int no_of_cities=40; // No of cities
506
                            // Population Size
       int popSize=100;
507
       int no_of_males=30; // No of males
       int generations=100; // No of generations
509
510
       double alpha = 0.9;
511
       double beta=0.4;
512
       double gamma=0.7;
513
514
       int no of hinds = popSize-no of males;
515
516
517
518
       vector < City > city List; // Stores the initial list of cities
519
       City index [no of cities];
52
521
       // Assigning random coordinates to cities
522
       for(int i=0;i< no\_of\_cities;i++){
523
            double x=200*((double) rand()/RAND MAX);
524
            double y=200*((double) rand()/RAND MAX);
            City c;
526
            c.x=x, c.y=y, c.id=i;
52
            cityList.pb(c);
            index[i]=c; // Assigning index to cities
52
       }
       auto begin = chrono::high resolution clock::now();
533
       RD Algo(cityList, popSize, no of males, generations, alpha, beta, gamma
534
       , index);
       auto end = chrono::high_resolution_clock::now();
       auto duration = chrono::duration_cast<chrono::microseconds>(end-begin);
537
       cout << fixed << set precision (5) << "Time Taken by serial program: "<<
538
       duration.count()/1000000.0 << "s" << endl;
539
540 return 0;
```