Types and pattern matching

Pure Object Oriented Language

- In a pure object oriented language, every value is an object
- If based on classes, then the type of each value is a class
- Is Scala a pure object oriented language?
 - What about functions?
 - What about primitives?
- Is Java? (No)

How are primitives treated

- Primitive types like Int or Boolean, treated like any other class
- Compiler represents scala.Int to JVM's primitive type 32bit int, scala.Boolean to JVM's primitive type boolean etc

Example

Let's look at how natural numbers can be constructed as objects, ie 0, 1, 2, 3, ...

types_and_pattern_matching.naturals

Functions as objects

- Function values are treated as objects in Scala
- Function type A=>B is abbreviation for the trait scala.Function1[A,B]
 - This has the apply method: Class[_ <: Int => Int]
- Functions are objects with an apply method
- Traits Function2, etc for multiple parameters
- See more: http://www.scala-lang.org/api/2.9.2/scala/Function1.html

```
scala> def next(n:Int):Int = n+1
next: (n: Int)Int
scala> next _ getClass
res3: Class[_ <: Int => Int] = class $$Lambda$1123/1125499532
```

Expansion behind the scenes

 What does the expansion of an anonymous function look like?

```
(x: Int) => x * x
// expands into

{ class AnonFun extends Function1[Int, Int] {
   def apply(x: Int) = x * x
   }
   new AnonFun
}
```

Anonymous class syntax

- Anonymous classes are unnamed classes
- How can we create an instance?
- By using the reserved word new and defining the body with braces
- Also, instances can be created from traits
- This is syntax sugar.

```
//anonymous class
val myPoint = new{ val x = 1; val y = 2 }

// instance from trait
trait AnonymousHero {
  def superpower: String
}

val myHero = new AnonymousHero {
  def superpower = "I can compile Scala with my brain"
}
```

Expansion, even shorter

Using the anonymous class syntax, the expanded function is even shorter

```
(x: Int) => x * x
// expands into

new Function1[Int, Int] {
  def apply(x: Int) = x * x
}
```

Methods and functions

- In many situations, you can ignore the difference between functions and methods
- But, methods such as def f(x: Int): Boolean = ... are not functions
 - A Scala method, as in Java, is a part of a class. It has a name, a signature, optionally some annotations, and some bytecode
 - A function in Scala is a complete object. There are a series of traits in Scala to represent functions with various numbers of arguments: Function0, Function1
- A method is converted to a function value, if it's name is used in a place, where a function value is expected
- When we treat a method as a function, such as by assigning it to a variable, Scala actually creates a
 function object whose apply method calls the original method, and that is the object that gets assigned to
 the variable

```
def f(x: Int): Boolean = ...
// is converted to
(x: Int) => f(x)
// in expanded form
new Function1[Int, Boolean] {
  def apply(x: Int) = f(x)
}
```

Example

 List revisited - how to use the function parameter syntax on an object?

polymorphism.List

00 + FP Polymorphism

- OO subtyping
- FP generics
- Let us look at their interactions
 - bounds constraints for type parameters
 - variance how parametrized types behave when subtyped

Type Bounds

- Let us assume that a method returns a value, but can also throw an exception
- What is the most precise return type?

def test(list: List): Boolean

Upper type bounds

- type parameters and abstract types may be constrained by a type bound
- An upper type bound T <: A declares that type variable T refers to a subtype of type A
- Type bound let us assume some knowledge about the paremetrized type
- Without the upper type bound annotation, it would not be possible to call method isSimilar in method findSimilar.

```
trait Similar {
  def isSimilar(x: Any): Boolean
}
object UpperBoundTest extends Application {
  def findSimilar[T <: Similar](e: T, xs: List[T]): Boolean =
    if (xs.isEmpty) false
    else if (e.isSimilar(xs.head)) true
    else findSimilar[T](e, xs.tail)
}</pre>
```

Lower type bounds

- Lower type bounds declare a type to be a supertype of another type
- The term T >: A expresses that the type parameter T or the abstract type T refer to a supertype of type A.

[S >: NonEmpty]

Mixed bounds

- It is possible to mix a lower bound with an upper bound
- We could express, that any type between NonEmpty and IntSet are allowed

```
[ S >: NonEmpty <:IntSet]
```

Covariance

- Int is a subtype of object : Int <: Object
- How about List[Int] <: List[Object]

Lab Time!

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