

# A Noise-tolerant Differentiable Learning Approach for Single Occurrence Regular Expression with Interleaving



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#### Motivation

**Problem.** Learning single occurrence regular expression with interleaving (SOIRE) in full expressive power from a set of text strings with noise.

#### Challenge.

- 1. heavy computation in searching to get the full expressive power
- 2. wrong search bias resulting from noisy data

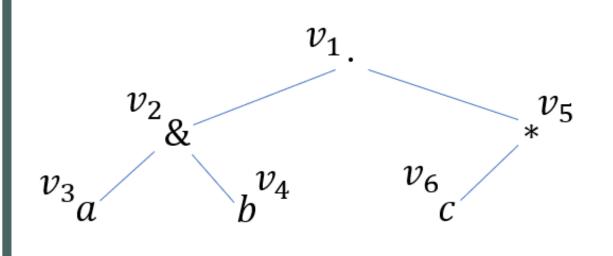
#### Contributions

- 1. We have proposed a noise-tolerant differentiable learning approach SOIREDL. We first train the neural network that simulates SOIRE matching to classify the given strings, and then interpret the target SOIRE from the parameters.
- 2. Theoretically, the faithful encodings learnt by SOIREDL one-to-one correspond to SOIREs for a bounded size.
- 3. Experimental results have demonstrated higher performance compared with the SOTA approaches.

### Filter Matching

Filter matching for a SOIRE r and a string s is to check if r matches  $filter(s,\alpha(r))$ , where  $\alpha(r)$  denotes the set of symbol in r, and function filter(s,V) returns a string that only retains symbols in V, where  $V\subseteq\Sigma$ . Let  $g_{i,j}^t \in \{0,1\}$  denote whether  $r^t$  matches  $filter(s_{i,j},\alpha(r^t))$ , where  $s_{i,j}$  denotes the substring of s from i to j, and where  $s_{1,0}=\epsilon$  specially.

Example 1. For  $(a\&b)c^*$  and s = dbac,



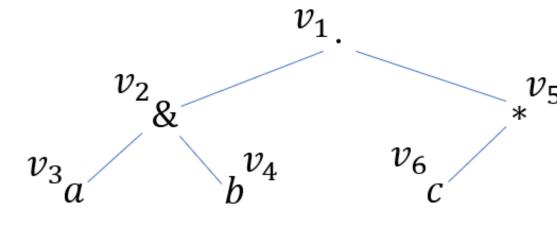
- filter matching is to check if  $(a\&b)c^*$  $matches \quad filter(dbac, \{a, b, c\})$  $bac, \ as \ \alpha((a\&b)c^*) = \{a, b, c\}.$
- $g_{1,2}^2$  denotes if a&b  $(r^2)$  matches baand  $g_{1,2}^2 = 1$ .

#### From SOIRETM to Neural Network

The trainable parameters  $\theta = (w, u)$ :

- ullet  $w\in[0,1]^{T imes|\mathbb{B}|}$ ,  $w_a^t$  denotes the probability of vertex t representing a symbol in  $\Sigma$  or an ordinary operator or the none operator.
- $\bullet$   $u \in [0,1]^{T imes T}$ ,  $u_{t'}^t$  denotes the probability of vertex t choosing vertex t'as its right child.

where  $\mathbb{B} = \Sigma \cup \{?, *, +, \cdot, \&, |, none\}$ , T the bounded size of the target SOIRE. Example 2.



When T = 6,  $w^1$ ,  $w^2$ ,  $w^3$ ,  $w^4$ ,  $w^5$ ,  $w^6$ ,  $u^1_5$ ,  $v_5$   $u_4^2$  are 1s, whereas other parameters are 0s. When T = 8,  $w_{none}^7$ ,  $w_{none}^8$  are 1s in addition to the above parameters.

Four main parts in conversion.

•  $\rho_a^t$ , the differentiable version of  $\alpha(r^t)$ .  $(\sigma_{01}(x) = \min(\max(x,0),1))$ 

$$\rho_a^t = \sigma_{01}(w_a^t + \sum_{o \in \{?, *, +, \cdot, \&, |\}} w_o^t \rho_a^{t+1} + \sum_{o \in \{\cdot, \&, |\}} w_o^t \sum_{t'=t+2}^T u_{t'}^t \rho_a^{t'})$$

ullet  $flag_{i,j}^{t,t'}$ , the probability that there does not exist a symbol occurring in both  $s_{i,j}$  and  $\alpha(r^t)$  but not occurring in  $\alpha(r^{t'})$ .

$$flag_{i,j}^{t,t'} = 1 - \sigma_{01}(\sum_{a \in \Sigma} \sigma_{01}(1[a \in s_{i,j}] + (\rho_a^t - \rho_a^{t'}) - 1))$$

•  $g_{i,j}^t$ , whether  $r^t$  matches  $filter(s_{i,j}, \alpha(r^t))$ 

$$g_{i,j}^t = \sum_{a \in \Sigma} w_a^t \cdot 1[filter(s_{i,j}, a) = a]$$

$$+ \sum_{o \in \{?,*,+\}} w_o^t p_{i,j}^t(o) + \sum_{o \in \{\cdot,\&,|\}} w_o^t \sum_{t'=t+2}^{1} u_{t'}^t p_{i,j}^t(o,t')$$

• return value of SOIRETM, combining  $filter(s, \alpha(r)) = s$  and  $g_{1,|s|}^1$ .

$$\hat{y} = g_{1,|s|}^1 - \max_{a \in \Sigma} \sigma_{01}(1[a \in s] - \rho_a^1)$$

Objective function:  $\frac{1}{2}(\hat{y}-y)^2$ , where y is ground-truth label for r matching s.

#### SOIRETM

**Theorem 1.** Given a SOIRE r and a string s,  $r \models s$  iff  $filter(s, \alpha(r)) = s \ and \ r \models filter(s, \alpha(r)).$ 

The steps of SOIRETM:

- 1. build the syntax tree of r
- 2. check if  $filter(s, \alpha(r)) = s$
- 3. calculate  $g_{i,j}^t$  from shorter substrings to longer ones and from bottom to top of the syntax tree (dynamic programming)
- 4. return  $g_{1,|s|}^1$  (filter matching)

**Theorem 2.** Given a SOIRE r and a string s,  $r \models s$  iff SOIRETM(r,s) = 1.

#### Faithful Encoding

length T is said to be faithful if it satisfies all the following conditions:

- 3.  $\forall 1 \leq t \leq T, \sum_{t'=t+2}^{T} u_{t'}^t + \sum_{a \in \Sigma \cup \{?,*,+,none\}} w_a^t = 1.$
- 6.  $\forall 3 \leq t \leq T, \forall 1 \leq p \leq t-2, (t-1-p)u_t^p + \sum_{n'=n+1}^{t-1} \sum_{t'=t+1}^T u_{t'}^{p'} \leq t$
- 7.  $\forall a \in \Sigma, \sum_{t=1}^{T} w_a^t \leq 1$ .

### Interpretation

**Theorem 3.** Given a bounded size  $T \in \mathbb{Z}^+$ , prefix notations of SOIREs  $r \ with \ |r| \leq T \ and \ faithful \ encodings \ \theta \ with \ length \ T \ have \ a \ one-to-one$  $correspondence, i.e., Enc2Pre(\theta) = PreForm(r).$ 

Based on this correspondence, we apply beam search to find a faithful encoding nearby the learnt encoding and then interpret it to the target SOIRE.

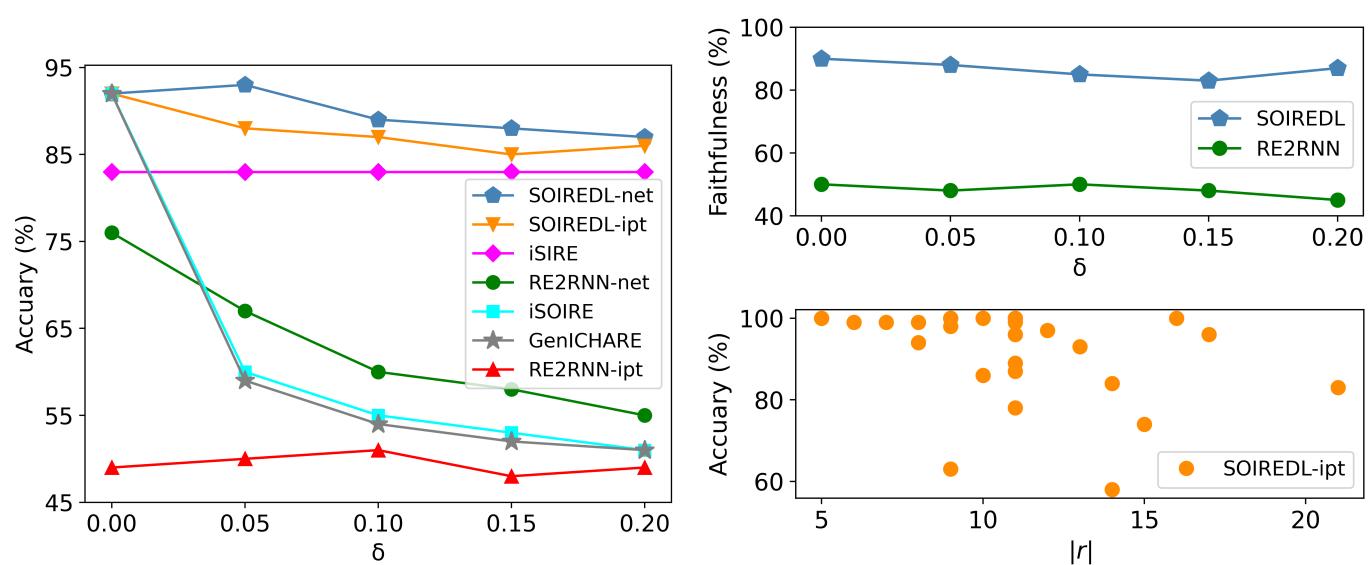
- The interpretation steps are conducted from bottom to top of the syntax tree. The score of each SOIRE is defined as the geometric mean of the probabilities of all operators and symbols.
- calculate the accuracy of each SOIRE in the last step on the training set and pick out the SOIRE with the highest accuracy.

## Result Analysis

Competitor.

- Positive strings only: iSOIRE (Li et al. 2019) for RSOIREs, GenICHARE (Zhang et al. 2018) for ICHAREs
- Both positive and negative strings: iSIRE (Li et al. 2020) for SIREs, RE2RNN (Jiang et al. 2020) for automatons, SOIREDL (Ours) for SOIREs

Comparisons on Noisy Data, Faithfulness and Scalability.



- SOIREDL gets the SOTA results and is the most robust on noisy data.
- The neural network of SOIREDL and its interpreted SOIRE are more consistent in performing SOIRE matching.
- The accuracy of SOIREDL decreases when the size of the ground-truth SOIRE increases.
  - This may be due to the difficulty for a neural network to capture the long-distance dependency in SOIRE matching.

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**Definition 1** (Faithful encoding). An encoding  $\theta = (w, u)$  of SOIREs with

- 1.  $\forall 1 \leq t \leq T, w^t \text{ is a one-hot vector.}$
- 2.  $\forall 1 \leq t \leq T, u^t$  is either a one-hot vector or an all-zero vector.
- 4.  $\forall 1 < t < T 1, w_{none}^{t+1} w_{none}^{t} > 0.$
- 5.  $\forall 2 \leq t \leq T, \sum_{a \in \{?,+,*,\cdot,\&,|\}} w_a^{t-1} + \sum_{t'=1}^{t-2} u_t^{t'} + w_{none}^t = 1.$