

Section 1: Database Fundamentals and Theory (Q1 & Q3)

Q1.1 & Q1.2: Definitions and Advantages (10 Marks)

| Guideline | Principle | Perfect Answer Example |
|-----------------------------|-------------------------------------------------------------------------------------------------------------------------------------------|------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| Data vs. Information | Use precise, contrasting terms. Data is raw. Information is processed data that provides meaning and context for decision-making. | Data is raw, unprocessed facts. Information is data that has been processed, organised, and structured to provide context and meaning. |
| DBMS Advantages | Memorise the five core benefits that solve traditional file system problems: Sharing, Security, Integrity, Redundancy, and Access. | 1. Minimized Data Redundancy. 2. Increased Data Integrity. 3. Enhanced Data Security. 4. Improved Data Sharing. 5. Faster Data Access and Retrieval. |
| Access Language | Name the purpose and the standard language. | Database Access Language is code used to define, manipulate (insert, change, delete), and retrieve data. Example: Structured Query Language (SQL). |

Q3: Keys and Data Analysis (10 Marks)

| Guideline | Principle | Perfect Answer Example |
|-----------------------------------------|------------------------------------------------------------------------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| Q3.1 Foreign Key (FK) | Identify the column in the child table that references the Primary Key (PK) of a parent table. | The Foreign Key in the Character table is ClassID . |
| Q3.3.1 Primary Key Suitability | PKs must be unique and non-null . If data shows a single column value repeats, it cannot be a PK. | No , the column is not suitable as a Primary Key. The value Modise is repeated for different clients, violating the uniqueness requirement. |
| Q3.3.2 Candidate Key Suitability | Candidate keys are potential Primary Keys, meaning they must also be unique . Use the same principle as the PK check. | No , the column is not suitable as a Candidate Key. The value Johannesburg is repeated for different clients, violating the uniqueness requirement. |



Section 2: ERD Interpretation (Q2 & Q3.2)

The most common error is misreading the multiplicities.

| Guideline | Step-by-Step Method for Reading Business Rules | Perfect Answer Example (Q2.2.3 Country/Province) |
|--------------------------------|--------------------------------------------------------------------------------------------|------------------------------------------------------------------------------|
| 1. Identify Cardinality | Look at the OPPOSITE end of the relationship line to the entity you are describing. | <i>To describe Province:</i> Look at the Country side: it shows 1..1. |
| 2. Translate Notation | Translate the notation into mandatory (mandated by | 1..1 means "must be in one |

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| | the straight line) or optional (circle) and the quantity (1 or many). | and only one Country." |
| 3. Write the Rule | Write the full sentence for the relationship: Entity A (Cardinality) Action Entity B. | A Province must belong to one and only one Country . |
| 4. Reverse the Rule | Now describe the opposite direction. | <i>To describe Country:</i> Look at the Province side: it shows 1..*. A Country has one or more Provinces . |

| Q | Relationship Type Identification | Perfect Answer Example |
|------------------------------|----------------------------------------------------------------|-------------------------------------------------------------------------------------------------------|
| Q2.2.1 (1..1 to 1..1) | If both sides are 1..1, it's One-to-One (1:1) . | One-to-One (1:1): One Person owns one Car, and one Car is owned by one Person. |
| Q2.2.2 (1..* to 1..*) | If both sides are 1..* or *, it's Many-to-Many (M:N) . | Many-to-Many (M:N): A Driver can drive many Trucks, and a Truck can be driven by many Drivers. |

Section 3: Normalisation (Q5)

Normalisation is a strict process. Follow the steps sequentially by identifying and resolving dependencies.

Q5.2: Identifying Normal Forms (9 Marks)

| Guideline | Rule for Identification | Q5.2.2 Perfect Motivation (1NF) |
|-------------------|------------------------------------------------------------------------------------------------------------------------------------------------------|----------------------------------------------------------------------------------------------------------------------|
| 1NF Check | Are all columns atomic? (Single values, no repeating groups). If Yes, it is at least 1NF. | Yes, all attributes are atomic. |
| 2NF Check | Is there a Partial Dependency (PD) ? (A non-key attribute depends on <i>part</i> of a composite key). If Yes, it is NOT in 2NF . | NO , it is not in 2NF because Partial Dependencies exist (e.g., Player Name -> Player Date Joined). |
| 3NF Check | Is there a Transitive Dependency (TD) ? (A non-key attribute depends on <i>another non-key attribute</i>). If Yes, it is NOT in 3NF . | NO , it is not in 3NF because Transitive Dependencies exist (e.g., Spaceship Name -> Spaceship Value). |
| Conclusion | State the highest form the relation <i>satisfies</i> . | First Normal Form (1NF) (because it violates both 2NF and 3NF). |

Q5.3: Normalising to 2NF (Resolve Partial Dependencies) (10 Marks)

1NF Relation: \$R\$ (<u>Player Name</u>, <u>Spaceship Registration</u>, Player Date Joined, Spaceship Name, Spaceship Value)

| Step | Action | Resulting Relation |
|--------------------|-------------------------------------------------------|--------------------|
| 1. Find PDs | Identify attributes determined by only part of | |

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| | the composite key: Player Name -> Player Date Joined and Spaceship Registration -> Spaceship Name, Spaceship Value. | |
| 2. Decompose PDs | Create a new table for each partial key and its dependent non-key attributes. | PLAYER (<u>Player Name</u>, Player Date Joined) |
| 3. Decompose PDs | Create a new table for the other partial key and its dependent attributes. | SPACESHIP (<u>Spaceship Registration</u>, Spaceship Name, Spaceship Value) |
| 4. Maintain Link | Create a link table using the full composite key of the original 1NF relation to preserve the relationship. | PLAYER_SPACESHIP (<u>Player Name</u>, <u>Spaceship Registration</u>) |

Q5.4: Normalising to 3NF (Resolve Transitive Dependencies) (9 Marks)

2NF Relations:

1. **PLAYER** (<u>Player Name</u>, Player Date Joined)
2. **SPACESHIP** (<u>Spaceship Registration</u>, Spaceship Name, Spaceship Value)
3. **PLAYER_SPACESHIP** (<u>Player Name</u>, <u>Spaceship Registration</u>)

| Step | Action | Resulting Relation |
|-------------------|---------------------------------------------------------------------------------------------------------------------------------------|--------------------|
| 1. Find TD | Check non-key to non-key dependencies. In SPACESHIP, assume Spaceship Name -> Spaceship Value (a non-key attribute determines another | |

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| | non-key attribute). | |
| 2. Decompose TD | Create a new table for the non-key determinant and its dependent attribute. | SPACESHIP_VALUE_LOOKUP (<u>Spaceship Name</u>, Spaceship Value) |
| 3. Update Parent | Remove the transitively dependent non-key attribute from the original table, leaving the non-key determinant (which becomes a FK). | SPACESHIP_DETAIL (<u>Spaceship Registration</u>, Spaceship Name \rightarrow FK to <i>Lookup</i>) |
| 4. Final 3NF | All relations are now in 3NF. | Final Relations: PLAYER, PLAYER_SPACESHIP, SPACESHIP_DETAIL, SPACESHIP_VALUE_LOOKUP |



Section 4: SQL Application (Q6)

The key to perfect SQL is correct syntax and using the right clause for the job.

Q6.1.1: CREATE TABLE (5 Marks)

| Guideline | Component | Perfect Code Snippet |
|---------------|-----------------------------------------------------------|----------------------------------------------------------------------------------------------------------|
| Schema | Define all columns, data types, and NOT NULL constraints. | PresidentID INT AUTO_INCREMENT NOT NULL, Name VARCHAR(250) NOT NULL, Year YEAR NOT NULL, ... |

| | | |
|-------------|-------------------------------------------------------------------|-----------------------------------------------------------------------------------|
| Keys | Define the primary key (PK) and foreign key (FK) with references. | PRIMARY KEY (PresidentID), FOREIGN KEY (CountryID) REFERENCES Country (CountryID) |
|-------------|-------------------------------------------------------------------|-----------------------------------------------------------------------------------|

SQL

```
CREATE TABLE President (
  PresidentID INT AUTO_INCREMENT NOT NULL,
  CountryID INT NOT NULL,
  Name VARCHAR(250) NOT NULL,
  Surname VARCHAR(250) NOT NULL,
  Year YEAR NOT NULL,
  PRIMARY KEY (PresidentID),
  FOREIGN KEY (CountryID)
    REFERENCES Country (CountryID)
);
```

Q6.1.2 - Q6.1.5: SELECT, COUNT, INSERT, FILTER (4+4+3+3 Marks)

| Question | SQL Command Guide | Perfect Code Snippet |
|----------------------------------|-----------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------|
| Q6.1.2 (Count after 2009) | Use COUNT() and the WHERE clause for filtering individual rows. | SELECT COUNT(PresidentID) FROM President WHERE Year > 2009; |
| Q6.1.3 (INSERT) | Specify the table and list the columns and values in order. | INSERT INTO Country (CountryID, Name, Abbreviation, CallingCode) VALUES (5, 'Botswana', 'BW', '267'); |

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|------------------------------|--------------------------------------------------------------------------|-----------------------------------------------------|
| Q6.1.4 (SELECT/ORDER) | Use SELECT * for all columns and ORDER BY for sorting. | SELECT * FROM Country ORDER BY Name ASC; |
| Q6.1.5 (LIKE Filter) | Use WHERE with the LIKE operator and the % wildcard for partial matches. | SELECT * FROM President WHERE Surname LIKE 'R%'; |

Q6.1.6: JOIN (5 Marks)

| Guideline | Principle | Perfect Code Snippet |
|---------------|---------------------------------------------------------------------------------------------|--------------------------------------------------------------------|
| JOIN | Link tables using the PK=FK relationship. Use table aliases (P and C) for brevity. | FROM President P JOIN Country C ON P.CountryID = C.CountryID |
| SELECT | Select only the required columns, using AS to rename joined attributes (like Country Name). | SELECT P.Name, P.Surname, C.Name AS CountryName |

SQL

```

SELECT
  P.Name,
  P.Surname,
  C.Name AS CountryName
FROM
  President P
JOIN
  Country C ON P.CountryID = C.CountryID;

```


Q6.2 & Q6.4: WHERE vs HAVING and Query Results (13 Marks)

| Q | Guideline for Result/Explanation | Perfect Answer |
|----------------------------------------------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| Q6.2 WHERE vs HAVING | WHERE filters rows <i>before</i> grouping. HAVING filters groups <i>after</i> grouping (must be used with GROUP BY). | WHERE filters individual rows based on a condition applied to non-aggregated columns. HAVING filters groups of rows based on a condition applied to aggregate functions (e.g., COUNT(), SUM()). |
| Q6.4.2 (SELECT J.*, COUNT()... HAVING >= 2) | The GROUP BY clause groups all cases by judge. The COUNT() function counts the cases per group. The HAVING clause filters out groups (judges) with less than 2 cases. | The query returns all columns for every Judge who has handled two or more cases, along with the total count of cases they handled. |
| Q6.4.3 (DROP TABLE Judge;) | DROP TABLE removes the database object entirely, not just the data. | The entire Judge table, including its structure (schema) and all data, is permanently deleted from the database. |