Lecture 06: Applications

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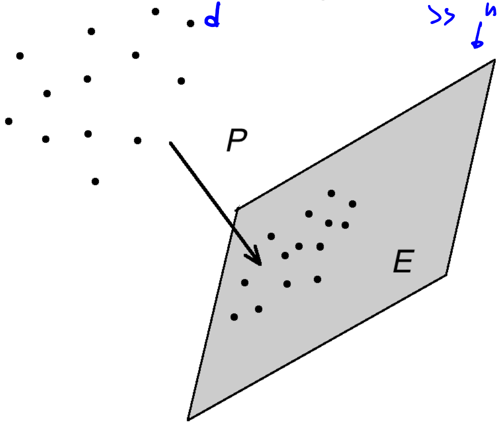
Dimension Reduction

- We started this course by observing that high dimensions could be a major problem, due to the curse of dimensionality.
- Why don't we then fix the problem due to high-dimensions as follows: Take the high-dimensional data and transform it into a low-dimensional data.
- Let there be N high-dimensional vectors $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$, where $\boldsymbol{x}_j^T = [x_1, x_2, ..., x_d]$, for $j = 1, 2, ..., N$, where d is very high.
- We would like to make a transformation of $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$ into $\boldsymbol{y}_1, \boldsymbol{y}_2, ..., \boldsymbol{y}_N$, where $\boldsymbol{y}_j^T = [y_1, y_2, ..., y_n]$, for $j = 1, 2, ..., N$, where $n \ll d$, and yet the geometry of $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$ are preserved in $\boldsymbol{y}_1, \boldsymbol{y}_2, ..., \boldsymbol{y}_N$.

$$[\underbrace{\cdots\cdots\cdots}_{d}] \underset{\gg}{\longrightarrow} [\underbrace{\quad\quad}_{n}]$$

Dimension Reduction

- First we have to define what do we mean by “the geometry of the data”: By “the geometry of the data”, we mean the pairwise distances between the original data.



Is this possible? It turns out it is possible if $n = O(\ln(N)) \ll d$.

Dimension Reduction

- Thm (Johnson-Lindenstrauss Lemma): $\forall$ vectors $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$, where $\boldsymbol{x}_j \in \mathbb{R}^d$, for $j = 1, 2, ..., N$, there exists a linear map $T : \boldsymbol{x}_j \to \boldsymbol{y}_j$, where $\boldsymbol{y}_j \in \mathbb{R}^n$ and $n \ll d$ such that the following holds

$y_j = T \cdot x_k$

$$\Pr\left\{(1-\delta)||\boldsymbol{x}_k - \boldsymbol{x}_j||_2 \le ||\boldsymbol{y}_k - \boldsymbol{y}_j||_2 \le (1+\delta)||\boldsymbol{x}_k - \boldsymbol{x}_j||_2\right\} \ge 1 - \epsilon \quad (1)$$

$y_k = T \cdot x_k$

for any $j \neq k$, where $j = 1, 2, ..., N$ and $k = 1, 2, ..., N$, and small $\delta > 0$ and $\epsilon > 0$ if

$$n > \frac{1}{c}\left(\ln(N) + \frac{1}{2}\ln\left(\frac{1}{\epsilon}\right) + \frac{1}{2}\ln(2)\right)$$

holds where

$$c = \frac{\delta(2-\delta)}{k}\min\left\{\frac{\delta(2-\delta)}{k}, 1\right\} \quad (2)$$

Dimension Reduction

Proof:

- The vectors $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$ and $\boldsymbol{y}_1, \boldsymbol{y}_2, ..., \boldsymbol{y}_N$ are all deterministic.
- However, we will use a probabilistic method to find the mapping $T : \boldsymbol{x}_j \to \boldsymbol{y}_j$.
- Specifically, we will choose a linear map $T(\cdot)$ at random, and then we will prove that the linear map $T(\cdot)$ satisfies the properties that we seek.
- Now, a linear map is simply a multiplication of $\boldsymbol{x}_1, \boldsymbol{x}_2, ..., \boldsymbol{x}_N$ by a matrix $\boldsymbol{T}$, of size $\underset{n \times d}{\cancel{d \times n}}$, to obtain $\boldsymbol{y}_1, \boldsymbol{y}_2, ..., \boldsymbol{y}_N$.
- Hence, if we will choose a linear map $T(\cdot)$ at random, this means that we should choose a matrix $\boldsymbol{T}$ at random.
- How is it possible that a randomly chosen matrix $\boldsymbol{T}$ would work?

$$\underset{n \cdot d}{T} \cdot \underset{d}{X_1} = \underset{n}{Y_1}$$

Dimension Reduction

- Let's have a random matrix $\boldsymbol{G}$ of size $n \times d$ populated by i.i.d. Gaussian entries, i.e., the (i,j)-th element of $\boldsymbol{G}$, denoted by G_{ij}, is generated i.i.d. according to the Gaussian distribution $N(0,1)$ for $i = 1, 2, ..., n$ and $j = 1, 2, ..., d$.
- Let us have a fixed vector $\boldsymbol{z} \in \mathbb{R}^d$, with norm $||\boldsymbol{z}||_2$.
- Now, let's investigate the distribution of $\boldsymbol{Gz}$. The i-th element of $\boldsymbol{Gz}$, denoted by $(\boldsymbol{Gz})_i$ is given by

$$(\boldsymbol{Gz})_i = \sum_{j=1}^{d} G_{ij} z_j \sim N\left(0, \sum_{j=1}^{d} z_j^2\right) = N\left(0, ||\boldsymbol{z}||_2^2\right) \stackrel{(a)}{=} N(0,1) \quad (3)$$

where (a) holds if and only if $||\boldsymbol{z}||_2 = 1$.

$x_1 \; x_2 \; \ldots \; x_N$

$x_k \quad x_l$

Dimension Reduction

On the other hand, let's fix two vectors, $\boldsymbol{x}_k$ and $\boldsymbol{x}_l$. Then, the vector $\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l$, according to (3), has i.i.d. zero-mean Gaussian elements each with variance $||\boldsymbol{x}_k - \boldsymbol{x}_l||_2^2$.
Hence, if we normalize $\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l$ by $||\boldsymbol{x}_k - \boldsymbol{x}_l||_2$, the vector

$[\cdot] \sim N(0, ||x_k - x_l||^2$

$$\frac{||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2}$$

$[\,0\,] \rightarrow N(0,1)$

will have i.i.d. zero-mean Gaussian elements each with variance one.
As a result, we can use the Thin-Shell Theorem, which states that

$(1+\delta)\sqrt{n}$

$(1-\delta)\sqrt{n}$

$$\Pr\left\{\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2} - \sqrt{n}\right| \leq \delta\sqrt{n}\right\} \geq 1 - 2\exp\left(-n\frac{\delta(2-\delta)}{k}\min\left\{\frac{\delta(2-\delta)}{k}, 1\right\}\right) \qquad (4)$$

Dimension Reduction

We will now prove the main theorem as follows. We will select a pair of vectors $\boldsymbol{x}_k$ and $\boldsymbol{x}_l$. We will prove (1) for the selected $\boldsymbol{x}_k$ and $\boldsymbol{x}_l$. Then, we will use the union bound to prove that the theorem holds for all pairs satisfying the given condition.

We start with the selected pair $\boldsymbol{x}_k$ and $\boldsymbol{x}_l$:

$x_1 \quad x_2 \quad \cdot \quad x_n$

$$
\begin{aligned}
&\Pr\left\{\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2} - \sqrt{n}\right| \le \delta\sqrt{n}\right\} \\
&= \Pr\left\{\left|||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2 - ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\sqrt{n}\right| \le ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\delta\sqrt{n}\right\} \\
&= \Pr\left\{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\sqrt{n}(1-\delta) \le ||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2 \le ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\sqrt{n}(1+\delta)\right\} \\
&= \cancel{\Pr\left\{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\sqrt{n}(1-\delta) \le ||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2 \le ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2\sqrt{n}(1+\delta)\right\}} \\
&= \Pr\left\{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2(1-\delta) \le \frac{1}{\sqrt{n}}||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2 \le ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2(1+\delta)\right\} \\
&= \text{continuation on next page} \qquad T = \frac{G}{\sqrt{n}} \qquad (5)
\end{aligned}
$$

Dimension Reduction

Continuation of (5)

$$\stackrel{(a)}{=} \Pr\left\{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2(1-\delta) \le ||\boldsymbol{T}\boldsymbol{x}_k - \boldsymbol{T}\boldsymbol{x}_l||_2 \le ||\boldsymbol{x}_k - \boldsymbol{x}_l||_2(1+\delta)\right\}$$

$$\stackrel{(b)}{\ge} 1 - 2\exp\left(-n\frac{\delta(2-\delta)}{k}\min\left\{\frac{\delta(2-\delta)}{k}, 1\right\}\right) \stackrel{(c)}{=} 1 - 2e^{-nc} \tag{6}$$

where (a) comes by setting

$$\boldsymbol{T} = \frac{1}{\sqrt{n}}\boldsymbol{G},$$

and (c) comes by setting

$$c = \frac{\delta(2-\delta)}{k}\min\left\{\frac{\delta(2-\delta)}{k}, 1\right\} \tag{7}$$

Dimension Reduction

We now take the union bound over all pairs of N vectors

$$
\begin{aligned}
&\Pr\left\{\bigcap_{k=1}^{N}\bigcap_{l=k+1}^{N}\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k-\boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k-\boldsymbol{x}_l||_2}-\sqrt{n}\right|\le\delta\sqrt{n}\right\} \qquad P(\cup_i A_i) \le \sum_i P(A_i)\\
&=1-\Pr\left\{\bigcup_{k=1}^{N}\bigcup_{\substack{l=k+1\\ l=1}}^{N}\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k-\boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k-\boldsymbol{x}_l||_2}-\sqrt{n}\right|\ge\delta\sqrt{n}\right\}\\
&\ge 1-\sum_{k=1}^{N}\sum_{l=1}^{N}\Pr\left\{\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k-\boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k-\boldsymbol{x}_l||_2}-\sqrt{n}\right|\ge\delta\sqrt{n}\right\}\\
&=1-N^2\Pr\left\{\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k-\boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k-\boldsymbol{x}_l||_2}-\sqrt{n}\right|\ge\delta\sqrt{n}\right\}\\
&\ge 1-2N^2e^{-nc}\\
&=1-e^{\ln(2)+2\ln(N)-nc}
\end{aligned}
$$

Dimension Reduction

Now, we want to set n, N, and c such that

$$\Pr\left\{\left|\frac{||\boldsymbol{G}\boldsymbol{x}_k - \boldsymbol{G}\boldsymbol{x}_l||_2}{||\boldsymbol{x}_k - \boldsymbol{x}_l||_2} - \sqrt{n}\right| \leq \delta\sqrt{n}\right\} \geq 1 - e^{\ln(2)+2\ln(N)-nc} \geq 1-\epsilon$$

$n \leq d$

$N > d$

which occurs if

$$1 - e^{\ln(2)+2\ln(N)-nc} \geq 1-\epsilon$$

or equivalently if

$$\ln(2) + 2\ln(N) - nc < \ln(\epsilon)$$

or equivalently if

$n > O(\ln(N))$ for d

$$\ln(N) < nc - \frac{1}{2}\ln\left(\frac{1}{\epsilon}\right) - \frac{1}{2}\ln(2),$$

$d = e^{100}$

$n = 1000600$

where c is given in (7) as function of δ. Q.E.D.

Note! We have lost d. Where is d?