

```

procedure coinrow(input C: array [1..n] of integer,
input/output F: array [0..n] of integer)
Algoritma
F[0] ← 0
F[1] ← C[1]
for i ← 2 to n do
    F[i] ← max(C[i] + F[i-2], F[i-1])

```

Pertanyaan :

- apa input size dari algoritma di atas $\rightarrow n$
- apa operasi dasarnya? $\rightarrow +$
- hitung $T(n)$ -ya \rightarrow
- masuk ke dalam kelas apa kompleksitas waktu dari algoritma tersebut

\hookrightarrow linear

$$C(n) \rightarrow \sum_{i=2}^n 1 = 2 - 1 + 1 = 1 - n \in O(n)$$

Algoritma (1) :

```

int sum(Array A, int n) /* ass n>=1 */
{
    int i=0, S=0;
    while (i<n){
        S+=A[i];
        i+=2;
    }
    return S;
}

```

- Tentukan input sizenya $\rightarrow n$
- Tentukan operasi dasarnya $\rightarrow + S += A[i]; i += 2$
- Hitung $T(n)$ algoritma tersebut dan tentukan kelas kompleksitasnya

\hookrightarrow linear $O(n)$

$$T(0) = 0$$

$$\begin{aligned} \text{loop } k &= n = \text{genap} \Rightarrow \frac{n}{2} \\ n &= \text{ganjil} \Rightarrow \frac{n+1}{2} \end{aligned}$$

$$T(n) = k \times O(n)$$

$$\therefore T(n) \begin{cases} \frac{n}{2} \times O(n) \\ \frac{n+1}{2} \times O(n) \end{cases}$$

$$\text{Fib}(n) = \begin{cases} 0, & n=0 \\ 1, & n=1 \\ T(n-1) + T(n-2) \end{cases}$$

$$T(n) = T(n-1) + T(n-2)$$

$$T(n) - T(n-1) - T(n-2) = 0$$

$$a_0 t_n - a_1 t_{n-1} - a_2 t_{n-2}$$

$$t^2 - t - 1$$

$$r = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

$$= \frac{-(-1) \pm \sqrt{1 - 4(-1)(-1)}}{2} = \frac{1 \pm \sqrt{5}}{2}$$

$$= \frac{1 - \sqrt{5}}{2}$$

$$T(n) = C_1 \left(\frac{1 + \sqrt{5}}{2} \right)^n + C_2 \left(\frac{1 - \sqrt{5}}{2} \right)^n$$

$$T(0) = C_1 \left(\frac{1 + \sqrt{5}}{2} \right)^0 + C_2 \left(\frac{1 - \sqrt{5}}{2} \right)^0 \Rightarrow C_1 + C_2 = 0$$

$$T(1) = C_1 \left(\frac{1 + \sqrt{5}}{2} \right)^1 + C_2 \left(\frac{1 - \sqrt{5}}{2} \right)^1 \Rightarrow C_1 \left(\frac{1 + \sqrt{5}}{2} \right) + C_2 \left(\frac{1 - \sqrt{5}}{2} \right)$$

$$C_1 = -C_2$$

$$T(n) = -C_2 \frac{1 + \sqrt{5}}{2} + C_2 \frac{1 - \sqrt{5}}{2}$$

$$= \frac{-C_2 + C_2 \sqrt{5} + C_2 + C_2 \sqrt{5}}{2} = 1$$

$$= C_2 \sqrt{5} = 1$$

$$= C_2 = \frac{1}{\sqrt{5}}$$

$$C_1 = -C_2$$

$$C_1 = -\frac{1}{\sqrt{5}}$$

$$T(n) = -\frac{1}{\sqrt{5}} \left(\frac{1 + \sqrt{5}}{2} \right)^n + \frac{1}{\sqrt{5}} \left(\frac{1 - \sqrt{5}}{2} \right)^n \in O \left(\left(\frac{1 + \sqrt{5}}{2} \right)^n \right)$$

Diketahui algoritma berikut ini:

Solusi:

```
Function Fib(n:integer) -> integer
  if n == 0
    return 0
  else if n == 1
    return 1
  else
    return Fib(n-1) + Fib(n-2)
```

Hitung $T(n)$ dengan persamaan karakteristik dan tentukan kelas kompleksitasnya

```

Soal 2:
Function Q(n:integer)
//Input: A positive integer n
if n = 1
  return 1
else
  return Q(n - 1) + 2 * n - 1
Pertanyaan: Hitung T(n) dengan cara substitusi

```

$$\begin{aligned}
 m(n) &= m(n-1) + 1 \\
 &= m(n-2) + 2
 \end{aligned}$$

$$m(n) = m(n-1) + n$$

$$= m(n-1-1) + n-1$$

$$= m(1) + n-1$$

$$= 1 + n-1 \Rightarrow n \in \theta(n)$$

$$m(n) = \begin{cases} 0, & n=0 \\ 1, & n=1 \\ 3 + m(n-1) \end{cases}$$

$$m(n) = m(n-1) + 3$$

$$= (m(n-2) + 3) + 3 \Rightarrow m(n-2) + 6$$

$$= (m(n-3) + 3) + 6 \Rightarrow m(n-3) + 9$$

$$m(n-1) + 12$$

$$m(n) = m(n \cdot n) + 3n$$

$$= m(n-n-1) + 3(n-1)$$

$$= m(1) + 3n-3$$

$$= 1 + 3n-3$$

$$= 3n-2 \in O(n)$$