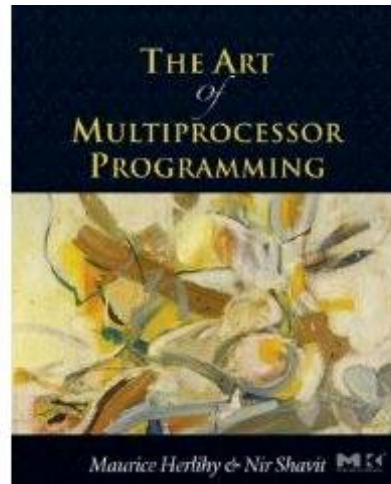


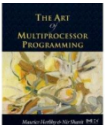
Programming Language Basics



Companion slides for
**The Art of Multiprocessor
Programming**
by Maurice Herlihy & Nir Shavit
With some very minor changes by APS

Languages for Multiprocessor Programming

- Java
- PThreads
 - C and C++
- C#
- MPI
- Etc...



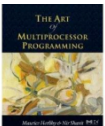
Threads

- Execution of a sequential program
- You can tell a thread
 - What to do
 - When to start
- You can
 - Wait for it to finish
- Other stuff:
 - Interrupt it, give it priority, etc.



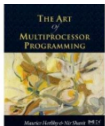
Threads in Java

- **Class `java.lang.Thread`**
- **Each thread has a method**
 - **`Void run()`**
- **Executes when it starts**
- **Thread vanishes when it returns**
- **You must provide this method**



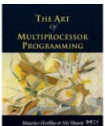
Creating a Thread

- **Create a Runnable object**
 - Runnable is an interface
 - Provides run() method
- **Pass Runnable object to thread constructor**



A Runnable Class

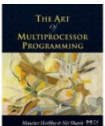
```
public class Hello implements Runnable {  
    String message;  
    public Hello(String m) {  
        message = m;  
    }  
    public void run() {  
        System.out.println(message);  
    }  
}
```



A Runnable Class

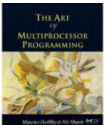
```
public class Hello implements Runnable {  
    String message;  
    public Hello(String m) {  
        message = m;  
    }  
    public void run() {  
        System.out.println(message);  
    }  
}
```

Runnable interface



Creating a Thread

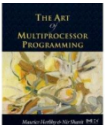
```
String m = "Hello from " + i;  
Runnable h = new Hello(m);  
Thread t = new Thread(h);
```



Creating a Thread

```
String m = "Hello from " + i;  
Runnable h = new Hello(m);  
Thread t = new Thread(h);
```

Create a Runnable object



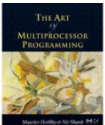
Creating a Thread

```
String m = "Hello from " + i;
```

```
Runnable h = new Hello(m);
```

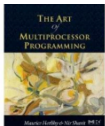
```
Thread t = new Thread(h);
```

Create the thread



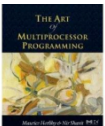
Syntactic Help

- **Defining a single-use class like Hello can be a nuisance**
- **Java provides special syntax**
- **Anonymous inner classes**
 - **May be more trouble than it's worth**
 - **You should recognize it**



Anonymous Inner Class

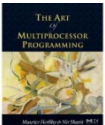
```
t = new Thread(  
    new Runnable() {  
        public void run() {  
            System.out.println(m);  
        }  
    }  
);
```



Anonymous Inner Class

```
t = new Thread(  
    new Runnable() {  
        public void run() {  
            System.out.println(m);  
        }  
    }  
);
```

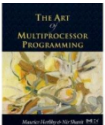
**Creates object of
anonymous Runnable
class**



Anonymous Inner Class

```
t = new Thread(  
    new Runnable() {  
        public void run() {  
            System.out.println(m);  
        }  
    }  
);
```

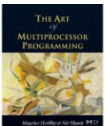
**Calls Thread constructor
with anonymous object**



Starting a Thread

```
t.start();
```

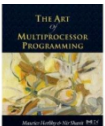
- **Starts the new thread**
- **Caller returns immediately**
- **Caller & thread run in parallel**



Joining a Thread

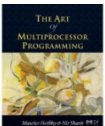
```
t.join();
```

- **Blocks the caller**
- **Waits for the thread to finish**
- **Returns when the thread is done**



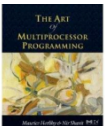
Monitors

- Each object has an implicit lock
- Managed by **synchronized modifier**
 - Methods
 - Code blocks
- OK for easy cases
- Not always for hard cases



Call Center Scenario

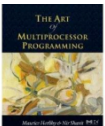
- **Calls arrive faster than they can be answered**
 - **Play recorded message**
 - “your call is very important to us ...”
 - **Put call in queue**
 - Play insipid music ...
 - **Operators dequeue call when ready**
 - **Single enqueueer, multiple dequeuers**



Bad Queue Implementation

```
class Queue<T> {  
    int head = 0, tail = 0;  
    T[] items = new T[QSIZE];  
    public enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    public T deq() {  
        return items[(head++) % QSIZE];  
    }  
}
```

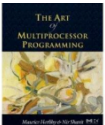
In practice, can't create array of generic type, so use `ArrayList<T>` instead



Bad Queue Implementation

```
class Queue<T> {  
    int head = 0, tail = 0;  
    T[] items = new T[QSIZE];  
    public enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    public T deq() {  
        return items[(head++) % QSIZE];  
    }  
}
```

**Works for
generic type T**



Bad Queue Implementation

```
class Queue<T> {  
    int head = 0, tail = 0;  
    T[] items = new T[QSIZE];  
    public enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    public T deq() {  
        return items[(head++) % QSIZE];  
    }  
}
```

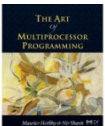
Array of T items



Bad Queue Implementation

```
class Queue<T> {  
    int head = 0, tail = 0;  
    T[] items = new T[QSIZE];  
    public enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    public T deq() {  
        return items[(head++) % QSIZE];  
    }  
}
```

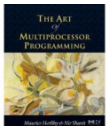
next slot to dequeue, 1st empty slot
#items in queue = tail - head



Bad Queue Implementation

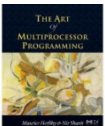
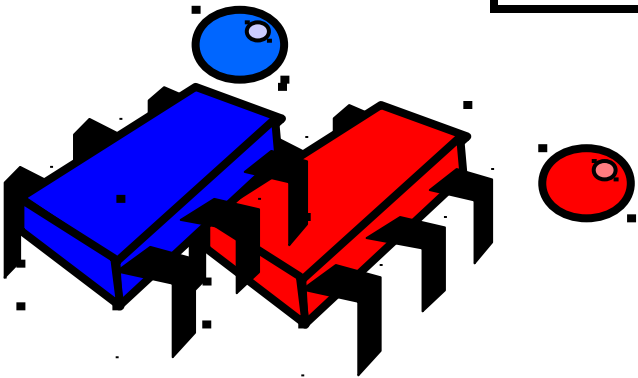
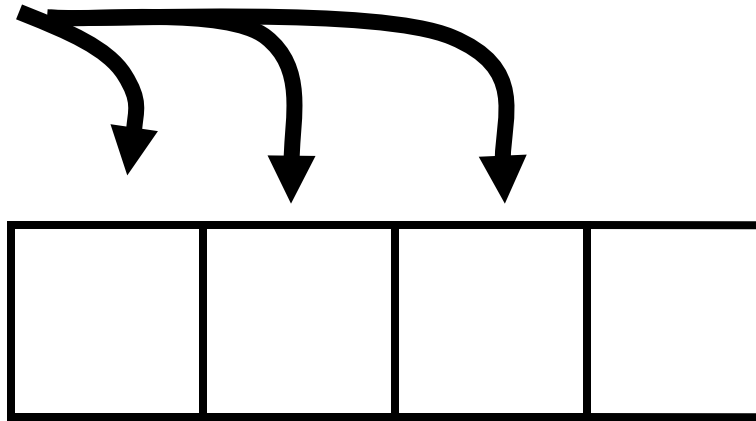
```
class Queue<T> {  
    int head = 0, tail = 0;  
    T[] items = new T[QSIZE];  
    public void enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    public T deq() {  
        return items[(head++) % QSIZE];  
    }  
}
```

Put in empty slot,
advance head



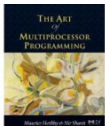
Of course, this doesn't work

head



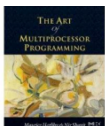
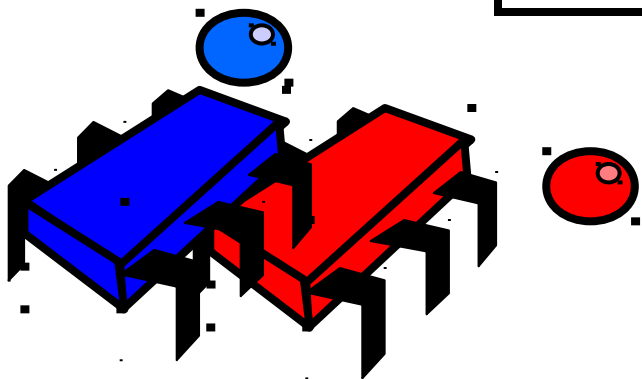
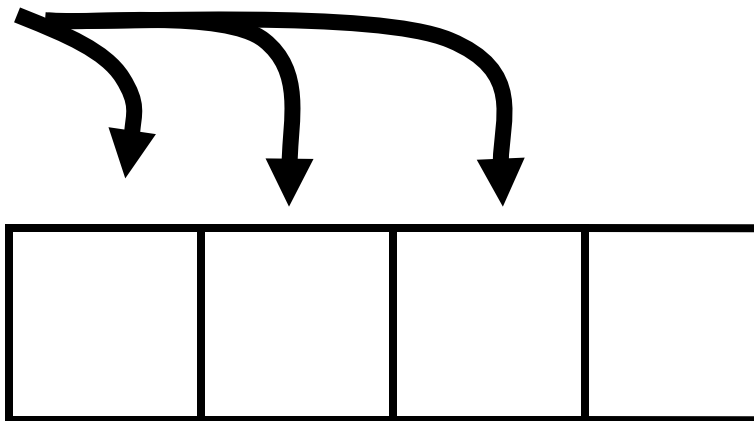
Mutual Exclusion

- Only one thread modifying queue fields at a time
- Use **synchronized methods**
 - Locks object on call
 - Releases lock on return



Mutual Exclusion

head



Synchronized Method

```
class Queue<T> {  
    ...  
    public synchronized void enq(T x) {  
        items[(tail++) % QSIZE];  
    }  
    ...  
}
```



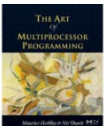
Synchronized Method

```
class Queue<T> {  
    ...  
    public synchronized enq(T x) {  
        items[(tail++) % QSIZE];  
    }  
    ...  
}
```

**Lock acquired on entry,
released on exit**

Syntactic Sugar

```
class Queue<T> {  
    ...  
    public void enq(T x) {  
        synchronized (this) {  
            items[(tail++) % QSIZE];  
        }  
    }  
    ...  
}
```



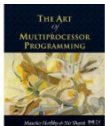
Syntactic Sugar

```
class Queue<T> {  
    ...  
    public void enq(T x) {  
        synchronized (this) {  
            items[(tail++) % QSIZE];  
        }  
    }  
    ...  
}
```

**Same meaning, more
verbose**

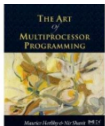
Vocabulary

- **A synchronized method locks the object**
- **No other thread can call another synchronized method for that same object**
- **Code in middle is critical section**



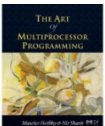
Re-entrant Locks

- **What happens if you lock the same object twice?**
 - In Java, no deadlock
 - Keeps track of number of times locked and unlocked
 - Unlock occurs when they balance out



Still Doesn't Work

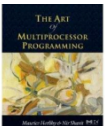
```
class Queue<T> {  
    ...  
    public synchronized void enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    ...  
}
```



Still Doesn't Work

```
class Queue<T> {  
    ...  
    public synchronized void enq(T x) {  
        items[(tail++) % QSIZE] = x;  
    }  
    ...  
}
```

What if the array is full?



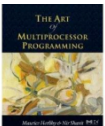
Waiting

- **What if**
 - Enqueueer finds a full array?
 - Dequeueer finds an empty array?
- **Throw an exception?**
 - What can caller do?
 - Repeated retries wasteful
- **Wait for something to happen**



Waiting Synchronized Method

```
class Queue<T> {  
    ...  
    public synchronized void enq(T x) {  
        while (tail - head == QSIZE) {};  
        items[(tail++) % QSIZE] = x;  
    }  
    ...  
}
```



Waiting Synchronized Method

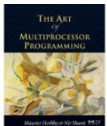
```
class Queue<T> {  
    ...  
    public synchronized enq(T x) {  
        while (tail - head == QSIZE) {};  
        items[(tail++) % QSIZE] = x;  
    }  
    ...  
}
```

Spin while the array is full



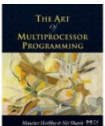
Deadlock

- **Enqueuer is**
 - Waiting for a dequeuer
 - While holding the lock
- **Dequeuer**
 - Waiting for enqueuer to release lock
- **Nothing will ever happen**



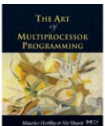
Waiting Thread

- Release lock while waiting
- When “something happens”
 - Re-acquire lock
 - Either
 - Re-release lock & resume waiting
 - Finish up and return



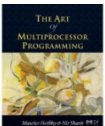
Styles of Waiting

- **Spinning**
 - Repeatedly retest condition
- **Blocking**
 - Ask OS to run someone else



Styles of Waiting

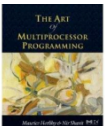
- **Spinning**
 - Good for very short intervals
 - Expensive to call OS
 - Works **only** on multiprocessors!
- **Blocking**
 - Good for longer intervals
 - Processor can do work
- **Clever libraries sometimes mix**



The wait() Method

```
q.wait();
```

- Releases lock on q
- Sleeps (gives up processor)
- Awakens (resumes running)
- Reacquires lock & returns
- (note: wait() throws InterruptedException)



The wait() Method

```
class Queue<T> {  
    ...  
    public synchronized void enq(T x) {  
        while (tail - head == QSIZE) {  
            wait();  
        };  
        items[(tail++) % QSIZE] = x;  
    }  
    ...  
}
```



Waiting Synchronized Method

```
class Queue<T> {
```

```
...
```

```
public synchronized enq(T x) {
```

```
    while (tail - head == QSIZE) {
```

```
        wait();
```

```
    };
```

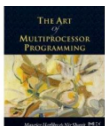
```
    items[(tail++) % QSIZE] = x;
```

```
}
```

```
...
```

```
}
```

Keep retesting condition



Waiting Synchronized Method

```
class Queue<T> {
```

```
...
```

```
public synchronized enq(T x) {
```

```
    while (tail - head == QSIZE) {
```

```
        wait();
```

```
    };
```

```
    items[(head++) % QSIZE] = x;
```

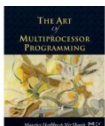
```
}
```

```
...
```

```
}
```

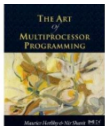
Keep retesting condition

Release lock & sleep



Wake up and Smell the Coffee

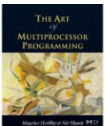
- **When does a waiting thread awaken?**
 - **Must be notified by another thread**
 - **when something has happened**
- **Failure to notify in a timely way is called a “lost wakeup”**



The wait() Method

```
q.notify();
```

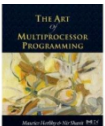
- **Awakens one waiting thread**
- **Which will reacquire lock & returns**



The wait() Method

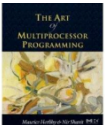
```
q.notifyAll();
```

- Awakens **all** waiting threads
- Which will reacquire lock & return



The wait() Method with Problem

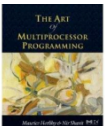
```
public synchronized enq(T x) {  
    while (tail - head == QSIZE) {  
        wait();  
    };  
    items[(tail++) % QSIZE] = x;  
    if (tail - head == QSIZE - 1) {  
        notify();  
    }  
}
```



The wait() Method with Problem

```
public synchronized enq(T x) {  
    while (tail - head == QSIZE) {  
        wait();  
    };  
    items[(tail++) % QSIZE] = x;  
    if (tail - head == QSIZE - 1) {  
        notify();  
    }  
}
```

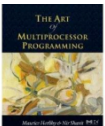
Wait for empty slot



The wait() Method with Problem

```
public synchronized enq(T x) {  
    while (tail - head == QSIZE) {  
        wait();  
    };  
    items[(tail++) % QSIZE] = x;  
    if (tail - head == QSIZE - 1) {  
        notify();  
    }  
}
```

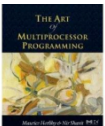
Stuff item into array



The wait() Method with Problem

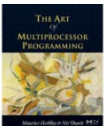
```
public synchronized enq(T x) {  
    while (tail - head == QSIZE) {  
        wait();  
    };  
    items[(tail++) % QSIZE] = x;  
    if (tail - head == QSIZE - 1) {  
        notify();  
    }  
}
```

If the queue was empty,
wake up a dequeuer

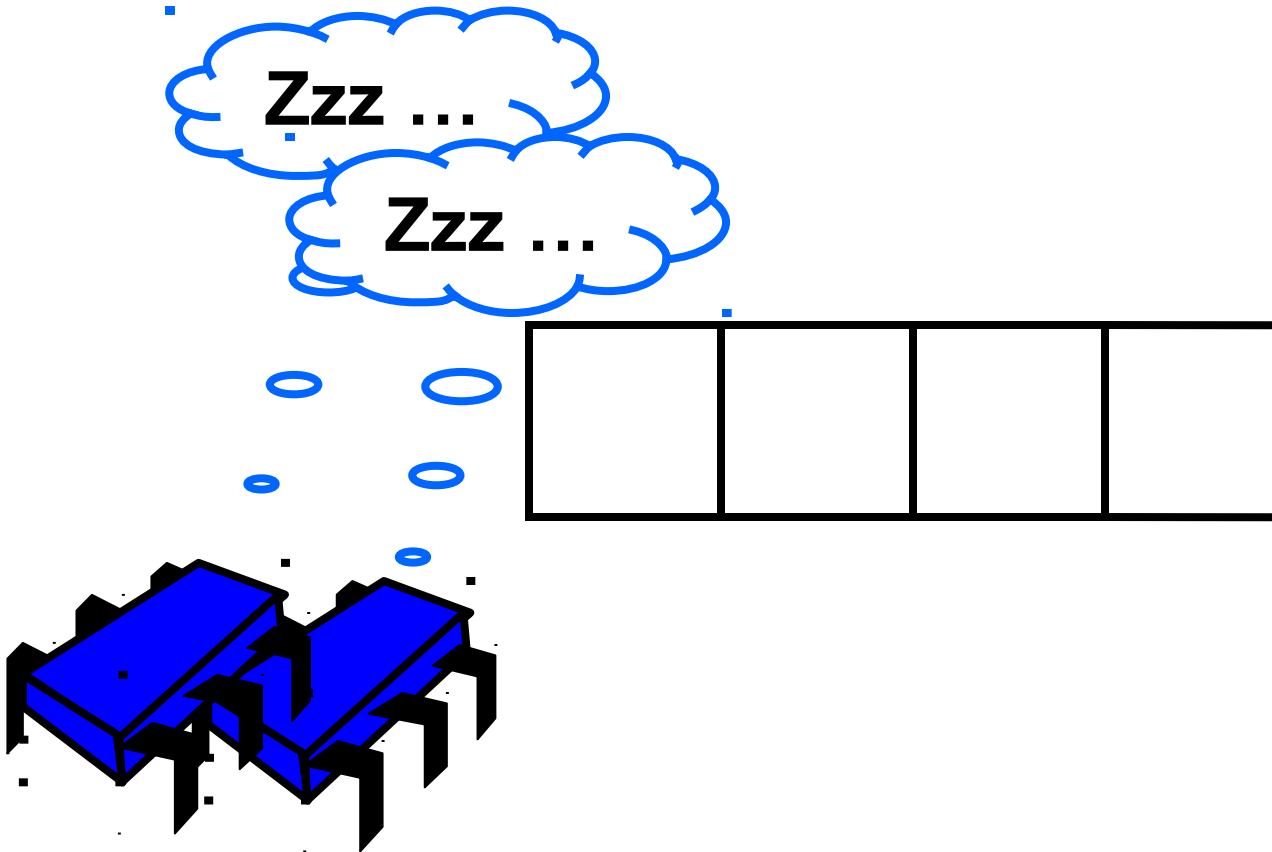


Lost Wakeup

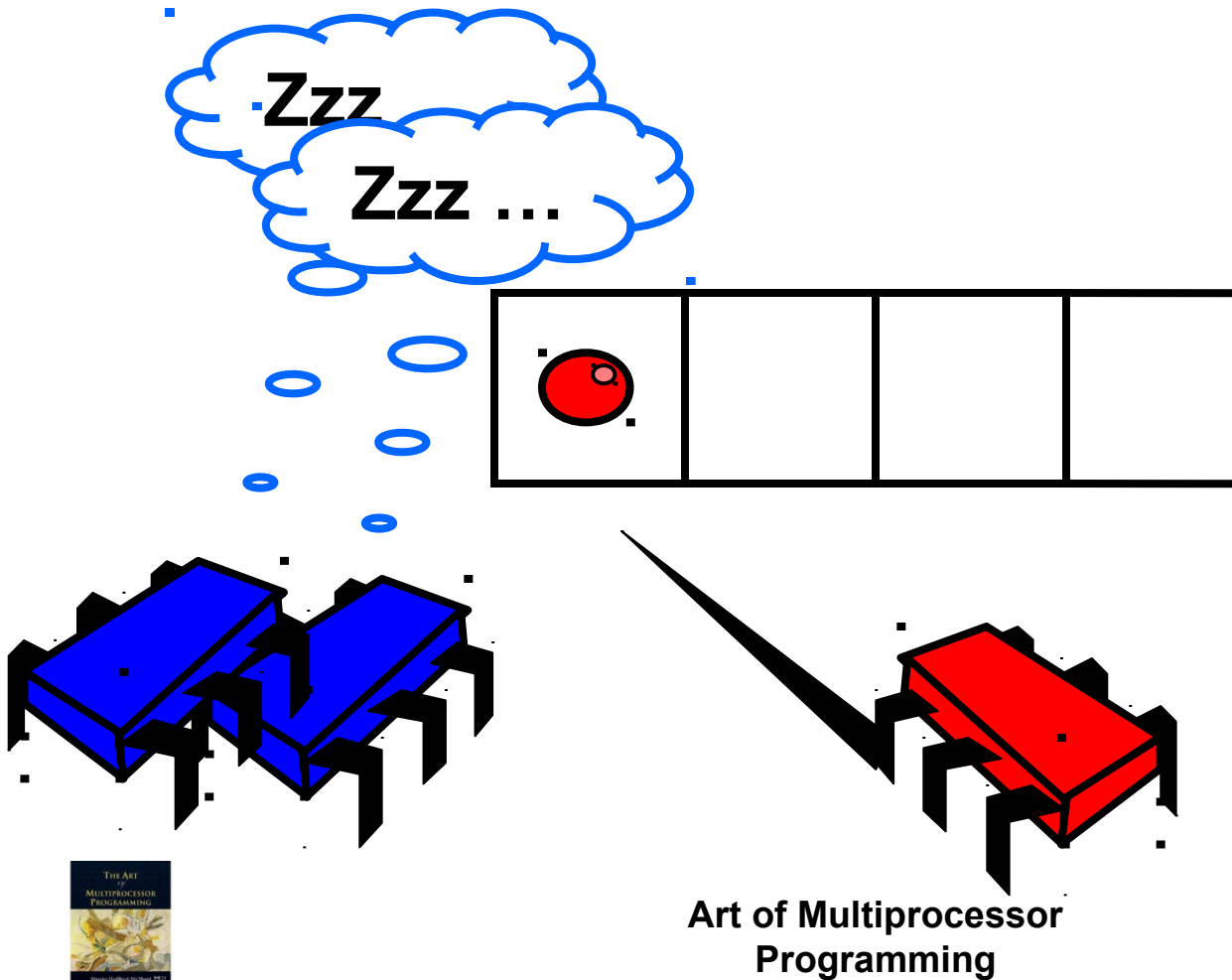
- This code has a lost wakeup bug
- Possible to have
 - Waiting dequeuer
 - Non-empty queue
- Because not enough threads awakened



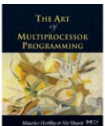
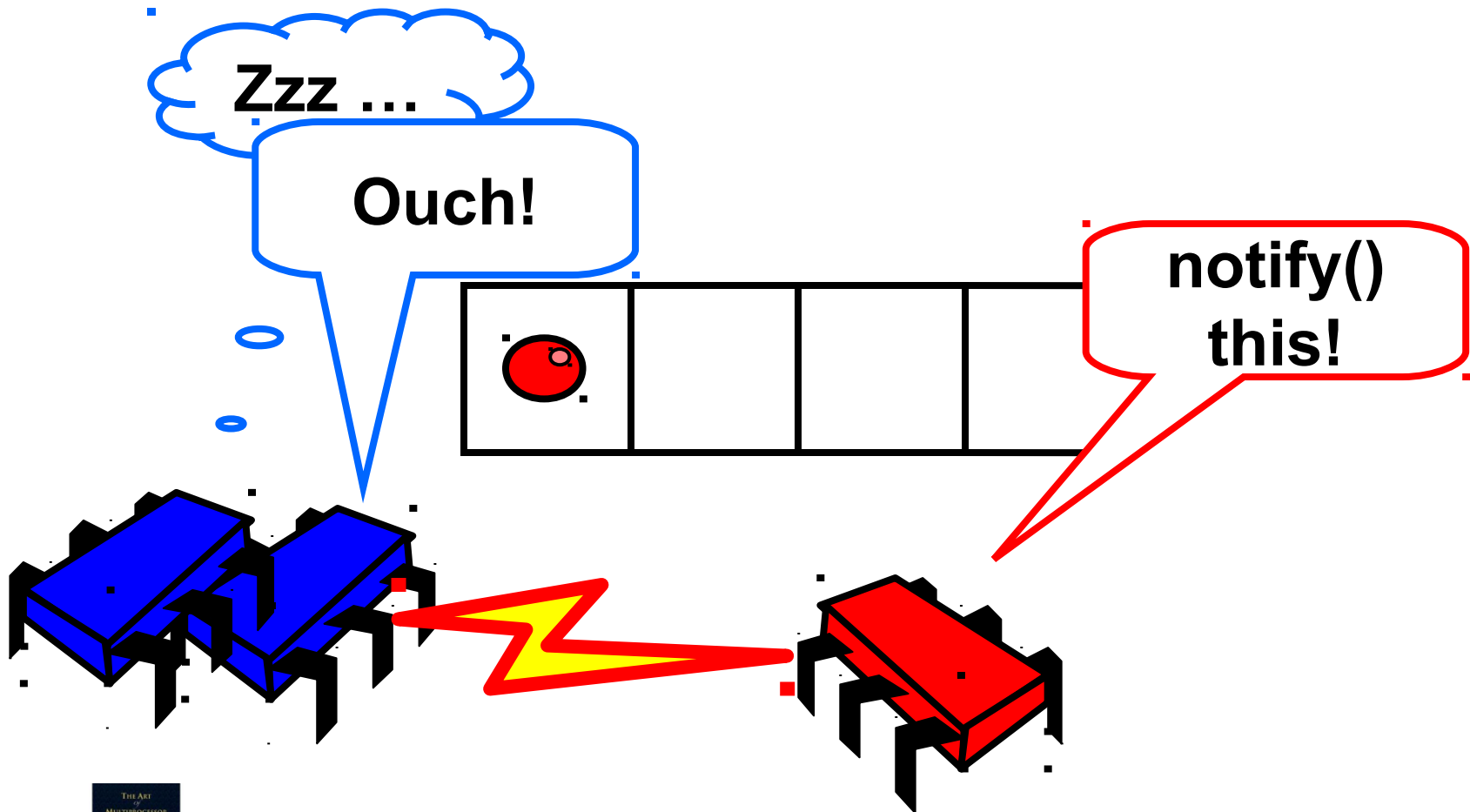
Empty queue, waiting dequeuers



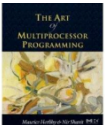
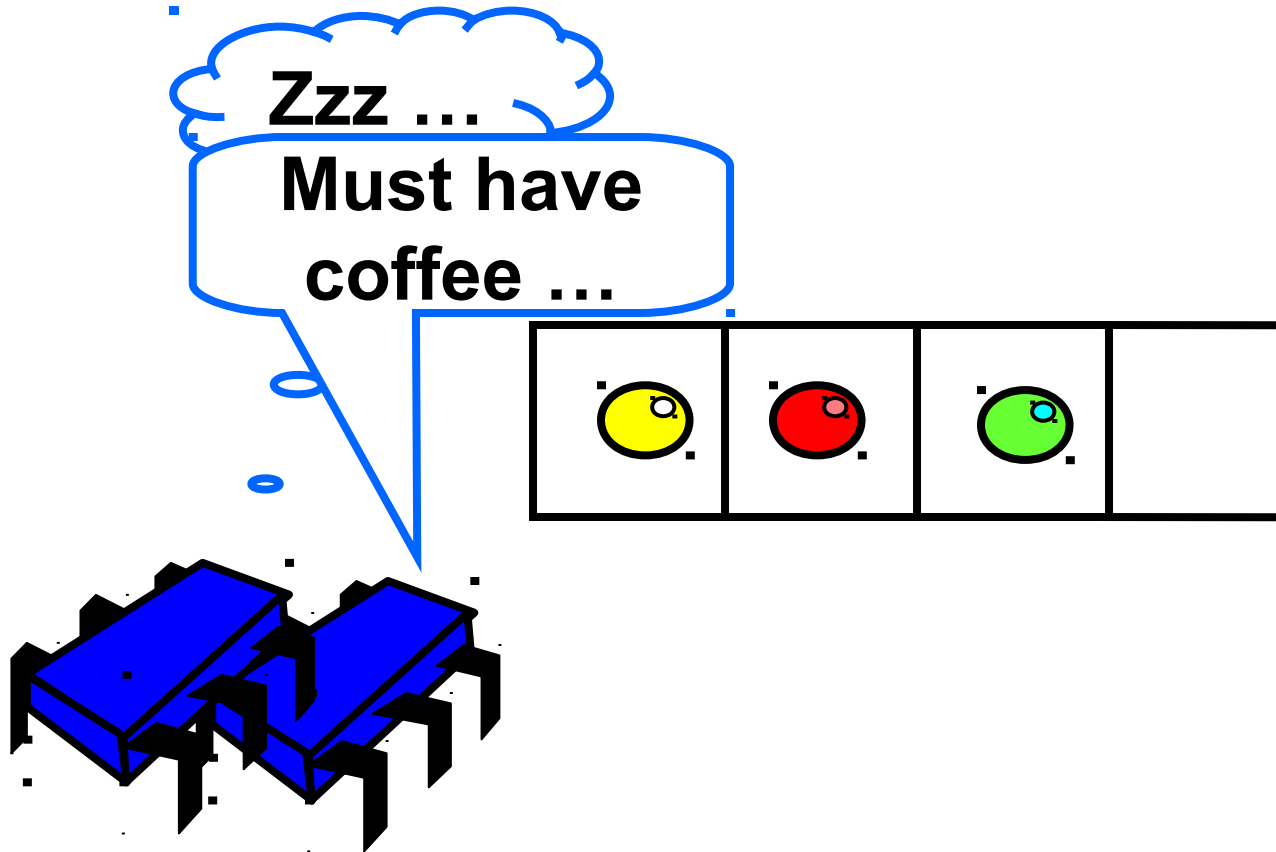
Enqueueer puts item in queue



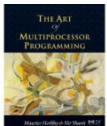
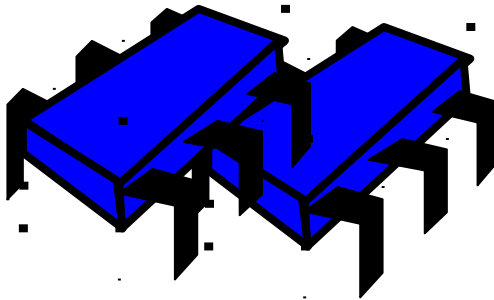
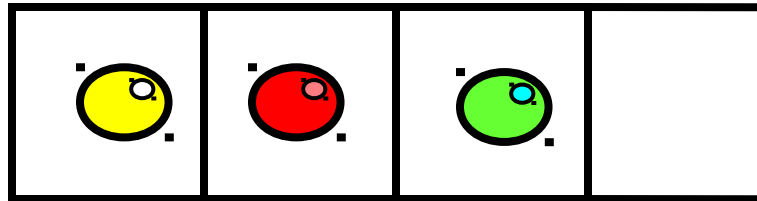
Since queue was empty, wakes dequeuer



1st Dequeueer slow, overtaken by enqueueers

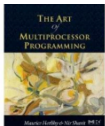


1st Dequeueur finishes



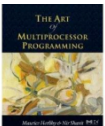
Solutions

- Don't write buggy code d'oh!
- Always call `notifyAll()`
- Can also use timed waits
 - Wake up after specified time



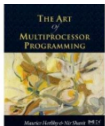
The wait() Method Solution

```
public synchronized enq(T x) {  
    while (tail - head == QSIZE) {  
        wait();  
    };  
    items[(tail++) % QSIZE] = x;  
    if (tail - head == QSIZE - 1) {  
        notifyAll();  
    }  
}
```



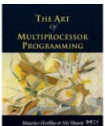
Thread-Local Data

- In many of our examples we assume
 - Threads have unique ids
 - In range $0, \dots, n-1$
- Where do they come from?
 - Passed in to Runner Constructor?
 - Many threads from same Runner?
 - Long-lived data
 - Unique to a thread



Thread-Local Data in Java

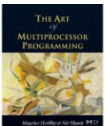
- **ThreadLocal<T> class**
- **No built-in language support**
- **Library classes**
 - **Syntax is awkward**
 - **Very useful anyway**
- **Note: for instance variables**
 - **Local variables in methods are on the thread's stack, so not shared**



ThreadLocal methods

```
ThreadLocal<T> local;  
T x = ...;  
local.set(x);
```

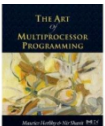
- **Changes calling thread's version of object**
- **Other threads' versions unaffected**



ThreadLocal methods

```
ThreadLocal<T> local;  
T x = local.get();
```

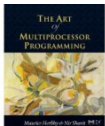
- Returns calling thread's version of object



Initializing ThreadLocals

```
T x = local.initialValue();
```

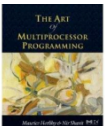
- Called by **get()** method the first time the thread-local variable is accessed.



Example

```
int me = ThreadID.get()
```

- Return unique thread id
- Take a number first time called



Thread-Local IDs

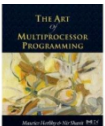
```
public class ThreadID {  
    private static volatile int nextID = 0;  
    private static LocalID threadID =  
        new LocalID();  
    public static int get() {  
        return threadID.get();  
    }  
    ... // define LocalID here  
}
```



Thread-Local IDs

```
public class ThreadID {  
    private static volatile int nextID = 0;  
    private static LocalID threadID =  
        new LocalID();  
    public static int get() {  
        return threadID.get();  
    }  
    ... // define LocalID here  
}
```

Next ID to assign



Thread-Local IDs

```
public class ThreadID {  
    private static volatile int nextID = 0;  
    private static LocalID threadID =  
        new LocalID();  
    public static int get() {  
        return threadID.get();  
    }  
    ... // define LocalID here  
}
```

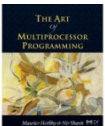
Declare & initialize thread-local ID



Thread-Local IDs

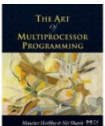
```
public class ThreadID {  
    private static volatile int nextID = 0;  
    private static LocalID threadID =  
        new LocalID();  
    public static int get() {  
        return threadID.get();  
    }  
    ... // define LocalID here  
}
```

Return value of thread-local ID



The Inner Class

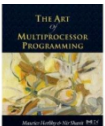
```
private static class LocalID
    extends ThreadLocal<Integer> {
    protected synchronized Integer
        initialValue() {
            return nextID++;
        }
}
```



The Inner Class

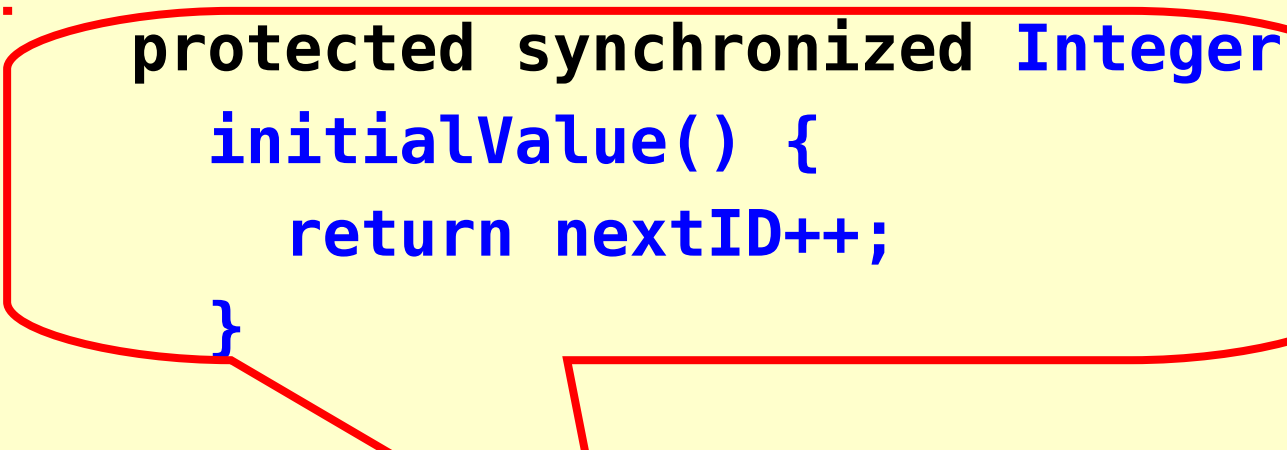
```
private static class LocalID
{
    extends ThreadLocal<Integer> {
        protected synchronized Integer
        initialValue() {
            return nextID++;
        }
    }
}
```

Subclass of ThreadLocal<Integer>

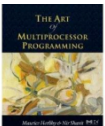


The Inner Class

```
private static class LocalID
    extends ThreadLocal<Integer> {
    protected synchronized Integer
        initialValue() {
            return nextID++;
        }
}
```

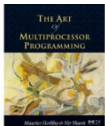


Overrides initialValue()



Summary

- **Threads**
 - And how to control them
- **Synchronized methods**
 - Wait, notify, and NotifyAll
- **Thread-Local objects**



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