

# Assignment 3

April 6, 2022

## 1 Q1

By definition, since  $a|b$ , we have  $\exists c \in \mathbb{Z}, b = ca$ . Similarly, we have  $\exists d \in \mathbb{Z}, a = db$ . So  $b = ca = cdb \implies (cd - 1)b = 0$ .

For  $b = 0$ , it is clear that  $a = 0$  such that  $a = b$  or  $a = -b$  is satisfied.

For  $b \neq 0$ , we have  $cd = 1$ , which has a solution of  $c = 1, d = 1$  or  $c = -1, d = -1$ . Therefore,  $a = b$  or  $a = -b$ . ■

## 2 Q2

1. Since  $1768 = 16 \times 110 + 8, 110 = 16 \times 6 + 14, 6 = 16 \times 0 + 6$ , then  $(1768)_{10} = (6E8)_{16}$ .
2. Since  $010_2 = 2_8, 101_2 = 5_8$ , then  $(10101)_2 = (25)_8$ .
3. Since  $3_{16} = 0011_2, B_{16} = 1011_2, 5_{16} = 0101_2, A_{16} = 1010_2$ , then  $(3B5A)_{16} = (11101101011010)_2$ .

## 3 Q3

1.  $256 = 2^8$
2.  $1890 = 2 \times 3^3 \times 5 \times 7$
3.  $5! = 2^3 \times 3 \times 5$

## 4 Q4

1. By Euclidean algorithm,

$$\begin{aligned} 267 &= 3 \times 79 + 30 \\ 79 &= 2 \times 30 + 19 \\ 30 &= 1 \times 19 + 11 \\ 19 &= 1 \times 11 + 8 \\ 11 &= 1 \times 8 + 3 \\ 8 &= 2 \times 3 + 2 \\ 3 &= 1 \times 2 + 1 \\ 2 &= 2 \times 1 \end{aligned}$$

we can get  $\gcd(267, 79) = 1$

2. From Euclidean algorithm, we have  $q_1 = 3, q_2 = 2, q_3 = q_4 = q_5 = 1, q_6 = 2, q_7 = 1, q_8 = 2$ . By extended Euclidean algorithm, from  $s_0 = 1, s_1 = 0, t_0 = 0, t_1 = 1$ , and for  $j = 2, 3, 4, 5, 6, 7, 8$ , we have

$$s_j = s_{j-2} - q_{j-1}s_{j-1}, \quad t_j = t_{j-2} - q_{j-1}t_{j-1}$$

Finally, we can get  $s_8 = 29, t_8 = -98$ . Thus,  $1 = \gcd(267, 79) = 267 \times 29 + 79 \times (-98)$

## 5 Q5

In this problem, we need to use a lemma: if  $d|a$  and  $d|b$ , then  $d|\gcd(a, b)$ . Here is the proof:

By Bezout's theorem,  $\gcd(a, b) = sa + tb$  for some integer  $s$  and  $t$ . Since  $d|a$  and  $d|b$ ,  $s$  and  $t$  are integers, by the property of division, then we have  $d|sa + tb = \gcd(a, b)$ . ■

Then we can start the proof as requested in the problem.

Let  $p = \gcd(\gcd(a, b), y)$ ,  $q = \gcd(d_1, d_2)$  and  $r = \gcd(a, b)$ .

By the definition of greatest common divisor, we have

$$\begin{array}{ll} p|\gcd(a, b), \text{ i.e., } p|r, & p|y \\ r|a, & r|b \end{array}$$

By the property of division, we have

$$p|a, \quad p|b, \quad p|y$$

By the lemma we had introduced, we have

$$\begin{array}{l} p|\gcd(a, y) = d_1 \text{ (since } p|a \text{ and } p|y) \\ p|\gcd(b, y) = d_2 \text{ (since } p|b \text{ and } p|y) \end{array}$$

Thus (also the lemma),

$$p|\gcd(d_1, d_2) = q \tag{1}$$

Conversely, since  $q = \gcd(d_1, d_2)$ , by the definition of greatest common divisor, we have

$$\begin{array}{l} q|d_1 = \gcd(a, y) \\ q|d_2 = \gcd(b, y) \end{array}$$

Since

$$\begin{array}{ll} \gcd(a, y)|a, & \gcd(a, y)|y \\ \gcd(b, y)|b, & \gcd(b, y)|y \end{array}$$

by the property of division, we have

$$q|a, \quad q|b, \quad q|y$$

By the lemma we had introduced, we have

$$q|\gcd(a, b) \text{ (since } q|a \text{ and } q|b)$$

Thus (also the lemma),

$$q|\gcd(\gcd(a, b), y) = p \tag{2}$$

By expression (1) and (2), we have

$$p|q, \quad q|p$$

As proved in Q1, we have  $p = q$  or  $p = -q$ . Since  $p$  and  $q$  are positive integers, we have  $p = q$ , i.e.,

$$\gcd(\gcd(a, b), y) = \gcd(d_1, d_2)$$

■

## 6 Q6

Let  $c = \gcd(a + b, a - b)$ , then by definition we have

$$c|(b + a) \implies \exists m \in \mathbb{Z}, b + a = cm$$

$$c|(b - a) \implies \exists n \in \mathbb{Z}, b - a = cn$$

So we can get

$$2a = (m - n)c, \quad 2b = (m + n)c$$

i.e.,  $c|2a, c|2b$ .

Suppose  $\gcd(a, c) = k$  and  $\gcd(b, c) = l$ . Since

$$b - a = cn, \text{ i.e., } b = cn + a$$

we have  $\gcd(b, c) = \gcd(a, c)$  by the lemma of Euclidean algorithm. Thus, we have  $k|a$  and  $l = k|b$ , i.e.,  $k$  is a common divisor of  $a$  and  $b$ .

Since  $0 < k \leq \gcd(a, b) = 1$  and  $k \in \mathbb{Z}$ , we have  $k = 1$ . Thus,  $\gcd(a, c) = \gcd(b, c) = 1$ . Then we have

$$c|2$$

since  $\gcd(a, c) = \gcd(b, c) = 1$  and  $c|2a, c|2b$ .

Therefore,  $\exists q \in \mathbb{Z}^+$ ,

$$2 = cq \geq c = \gcd(b + a, b - a)$$

■

## 7 Q7

1. Take  $p = 4$  (not a prime) and  $a = 2$  (not divisible by  $p$ ). We have  $a^{p-1} = 2^3 \not\equiv 1 \pmod{p} = 1 \pmod{4}$  since  $2^3 \pmod{4} = 0$  but  $1 \pmod{4} = 1$ .

2. • By Fermat's little theorem, we have  $302^{10} \equiv 1 \pmod{11}$  since 11 is a prime and 302 cannot be divided by 11. So

$$302^{302} = 302^{30 \times 10 + 2} = (302^{10})^{30} \times 302^2 \equiv 302^2 \pmod{11}$$

Since  $302^2 = (11 \times 27 + 5)^2 \equiv 5^2 \equiv 3 \pmod{11}$ , we have  $302^{302} \pmod{11} = 3$ .

- By Fermat's little theorem, we have  $4762^{12} \equiv 1 \pmod{13}$  since 13 is a prime and 4762 cannot be divided by 13. So

$$4762^{5367} = 4762^{447 \times 12 + 3} = (4762^{12})^{447} \times 4762^3 \equiv 4762^3 \pmod{13}$$

Since  $4762^3 = (13 \times 366 + 4)^3 \equiv 4^3 = 12 \pmod{13}$ , we have  $4762^{5367} \pmod{13} = 4762^3 = 12$ .

- By Fermat's little theorem, we have  $4^{522} \equiv 1 \pmod{523}$  since 523 is a prime and 4 cannot be divided by 523. So

$$2^{39674} = 4^{83 \times 239} = 4^{38 \times 522 + 1} = (4^{522})^{38} \times 4 \equiv 4 \pmod{523}$$

that is,  $2^{39674} \pmod{523} = 4$ .

## 8 Q8

1. By Euclidean algorithm,

$$267 = 3 \times 79 + 30$$

$$79 = 2 \times 30 + 19$$

$$30 = 1 \times 19 + 11$$

$$19 = 1 \times 11 + 8$$

$$11 = 1 \times 8 + 3$$

$$8 = 2 \times 3 + 2$$

$$3 = 1 \times 2 + 1$$

$$2 = 2 \times 1$$

we can get  $\gcd(267, 79) = 1$ . By extended Euclidean algorithm, from  $s_0 = 1, s_1 = 0, t_0 = 0, t_1 = 1$ , and for  $j = 2, 3, 4, 5, 6, 7, 8$ , we have

$$s_j = s_{j-2} - q_{j-1}s_{j-1}, \quad t_j = t_{j-2} - q_{j-1}t_{j-1}$$

Finally, we can get  $s_8 = 29, t_8 = -98$ . Thus,  $1 = \gcd(267, 79) = 267 \times 29 + 79 \times (-98)$ . So 29 is the inverse of  $267 \bmod 79$ . Therefore,  $29 \times 267x \equiv 29 \times 3 \pmod{79} \implies x \equiv 87 \pmod{79}$ .

2. By Euclidean algorithm,

$$312 = 3 \times 97 + 21$$

$$97 = 4 \times 21 + 13$$

$$21 = 1 \times 13 + 8$$

$$13 = 1 \times 8 + 5$$

$$8 = 1 \times 5 + 3$$

$$5 = 1 \times 3 + 2$$

$$3 = 1 \times 2 + 1$$

$$2 = 2 \times 1$$

we can get  $\gcd(312, 97) = 1$ . By extended Euclidean algorithm, from  $s_0 = 1, s_1 = 0, t_0 = 0, t_1 = 1$ , and for  $j = 2, 3, 4, 5, 6, 7, 8$ , we have

$$s_j = s_{j-2} - q_{j-1}s_{j-1}, \quad t_j = t_{j-2} - q_{j-1}t_{j-1}$$

Finally, we can get  $s_8 = 37, t_8 = -119$ . Thus,  $1 = \gcd(312, 97) = 312 \times 37 + 97 \times (-119)$ . So 37 is the inverse of  $312 \bmod 97$ . Therefore,  $37 \times 312x \equiv 37 \times 3 \pmod{97} \implies x \equiv 111 \pmod{97}$ .

## 9 Q9

Let domain  $A = \{0, \dots, m-1\}$  and codomain  $B = \{0, \dots, m-1\}$

• Injective.

Suppose for any  $x, y \in A$ , we have  $f(x) = f(y)$ , then

$$(ax) \bmod m = (ay) \bmod m$$

$$\implies ax \equiv ay \pmod{m} \text{ (by definition of congruence)}$$

$$\implies x \equiv y \pmod{m} \text{ (since } \gcd(a, m) = 1)$$

$$\implies x = y \text{ (since } x < m \text{ and } y < m)$$

So  $f$  is injective.

• Surjective.

Since  $|A| = |\{0, \dots, m-1\}| = m$  and  $f : A \mapsto B$  is injective, then  $|f(A)| = |\{f(0), \dots, f(m-1)\}| = m$ .

Since  $|B| = |\{0, \dots, m-1\}| = m$  and  $f(A) \subseteq B$ , we have  $f(A) = B$ , i.e.,  $f$  is surjective.

## 10 Q10

For  $n$  being even, i.e.,  $n = 2k$  with  $k$  an integer. Then  $n^2 = 4k^2$ , so  $n^2 \equiv 0 \pmod{4}$ .

For  $n$  being odd, i.e.,  $n = 2k + 1$  with  $k$  an integer. Then  $n^2 = 4k^2 + 4k + 1 = 4(k^2 + k) + 1$ , so  $n^2 \equiv 1 \pmod{4}$ .  
After all,  $n^2 \equiv 0$  or  $1 \pmod{4}$ . ■

## 11 Q11

Suppose there are two integers  $a$  and  $b$ , then from Q10, we have  $n^2 \equiv 0 \pmod{4}$  if  $n$  is even and  $n^2 \equiv 1 \pmod{4}$  if  $n$  is odd, where  $n$  is an integer. We can start from the following situations:

- Both  $a$  and  $b$  are even. Then  $a^2 \equiv 0 \pmod{4}$  and  $b^2 \equiv 0 \pmod{4}$ . By the property of congruence, we have  $a^2 + b^2 \equiv 0 + 0 \pmod{4}$ , i.e.,  $a^2 + b^2 \equiv 0 \pmod{4}$ .
- Both  $a$  and  $b$  are odd. Then  $a^2 \equiv 1 \pmod{4}$  and  $b^2 \equiv 1 \pmod{4}$ . By the property of congruence, we have  $a^2 + b^2 \equiv 1 + 1 \pmod{4}$ , i.e.,  $a^2 + b^2 \equiv 2 \pmod{4}$ .
- One is even and the other is odd in  $a$  and  $b$ . There is no change that we suppose  $a$  is odd and  $b$  is even (if not in real situation, then exchange the value of  $a$  and  $b$ ), then  $a^2 \equiv 1 \pmod{4}$  and  $b^2 \equiv 0 \pmod{4}$ . By the property of congruence, we have  $a^2 + b^2 \equiv 1 + 0 \pmod{4}$ , i.e.,  $a^2 + b^2 \equiv 1 \pmod{4}$ .

Thus,  $a^2 + b^2 \equiv 0$  or  $1$  or  $2 \pmod{4}$ . However, for  $m = 4k + 3$  with  $k \in \mathbb{Z} \wedge k \geq 0$ , i.e.,  $m \equiv 3 \pmod{4}$ , it is clear that  $a^2 + b^2$  and  $m$  are not in the same residue class  $\pmod{4}$ . So,  $m$  will not be the sum of the squares of two integers. ■

## 12 Q12

To prove the proposition in the problem, we only need to prove its contrapositive: for positive integers  $a$  and  $m$ , if  $a$  has the inverse modulo  $m$ , i.e.,  $ax \equiv 1 \pmod{m}$  has solution of  $x$ , then  $\gcd(a, m) = 1$ .

Suppose there is a solution  $x_0$  of  $ax \equiv 1 \pmod{m}$ , then  $ax_0 - 1 = km$  for  $k \in \mathbb{Z}$ . Thus,  $1 = ax_0 - km$ . By the property of division, we have  $\gcd(a, m) | (ax_0 - km)$ . So,  $\gcd(a, m) | 1$ , i.e.,  $\gcd(a, m) = 1$ . ■

## 13 Q13

The same as Q2.

## 14 Q14

By definition, since  $a \equiv b \pmod{m}$ , we have  $m | (a - b)$ , i.e.  $\exists k \in \mathbb{Z}, (a - b) = km$ , which means  $a = km + b$ . By the lemma of Euclidean algorithm, we have  $\gcd(a, m) = \gcd(m, b) = \gcd(b, m)$ . ■

## 15 Q15

Since  $x \equiv 3 \pmod{6}$ , we have  $\exists t \in \mathbb{Z}^+, x = 3 + 6t$ . Substituting into  $x \equiv 4 \pmod{7}$ , we have

$$3 + 6t \equiv 4 \pmod{7}$$

which can be simplified as  $t \equiv 6 \pmod{7}$ , i.e.,  $\exists k \in \mathbb{Z}^+, t = 6 + 7k$ . Therefore,  $x = 3 + 6t = 3 + 6(6 + 7k) = 39 + 42k$ , i.e.,  $x \equiv 39 \pmod{42}$ .

## 16 Q16

Since  $x \equiv 5 \pmod{6}$ , we have  $\exists t \in \mathbb{Z}^+, x = 5 + 6t$ . Substituting into  $x \equiv 3 \pmod{10}$ , we have

$$5 + 6t \equiv 3 \pmod{10}$$

which can be simplified as  $t \equiv 3 \pmod{5}$ , i.e.,  $\exists k \in \mathbb{Z}^+, t = 3 + 5k$ . Thus,  $x = 5 + 6t = 5 + 6(3 + 5k) = 23 + 30k$ . Similarly, substituting into  $x \equiv 8 \pmod{15}$ , we have

$$23 + 30k \equiv 8 \pmod{15}$$

which can be simplified as  $k \equiv 0 \pmod{1}$ . By back substituting, we have  $x \equiv 23 \pmod{30}$ .

## 17 Q17

Since  $de \equiv 1 \pmod{(p-1)(q-1)}$ , we have

$$\exists k \in \mathbb{Z}, de = 1 + k(p-1)(q-1)$$

Since  $C \equiv M^e \pmod{pq}$ , we have

$$C^d \equiv (M^e)^d \equiv M^{de} \equiv M^{1+k(p-1)(q-1)} \pmod{pq}$$

Since  $p$  and  $q$  are primes, and  $\gcd(M, pq) > 1$ , then we have  $M = lp$  for  $l \in \mathbb{Z}$  and  $\gcd(M, q) = 1$ . (Or  $M = lq$  for  $l \in \mathbb{Z}$  and  $\gcd(M, p) = 1$ , which has the same analysis).

By Fermat's little theorem, we have

$$M^{q-1} \equiv 1 \pmod{q}$$

Thus,

$$C^d \equiv M^{ed} \equiv M^{1+k(p-1)(q-1)} \equiv M \cdot (M^{q-1})^{k(p-1)} \equiv M \pmod{q}$$

By the property of congruence, we have

$$\exists t \in \mathbb{Z}, M^{ed} = M + tq, \text{ i.e., } (lp)^{ed} = lp + tq$$

Since  $p|lp$  and  $p|(lp)^{ed}$ , we have  $p|tq$ . Since  $\gcd(p, q) = 1$ , by the property of division, we have  $p|t$ . So we can let  $t = rp$  for  $r \in \mathbb{Z}$ . Then

$$M^{ed} = M + tq = M + rpq$$

Therefore,  $C^d \equiv M^{ed} \equiv M \pmod{pq}$ . ■