

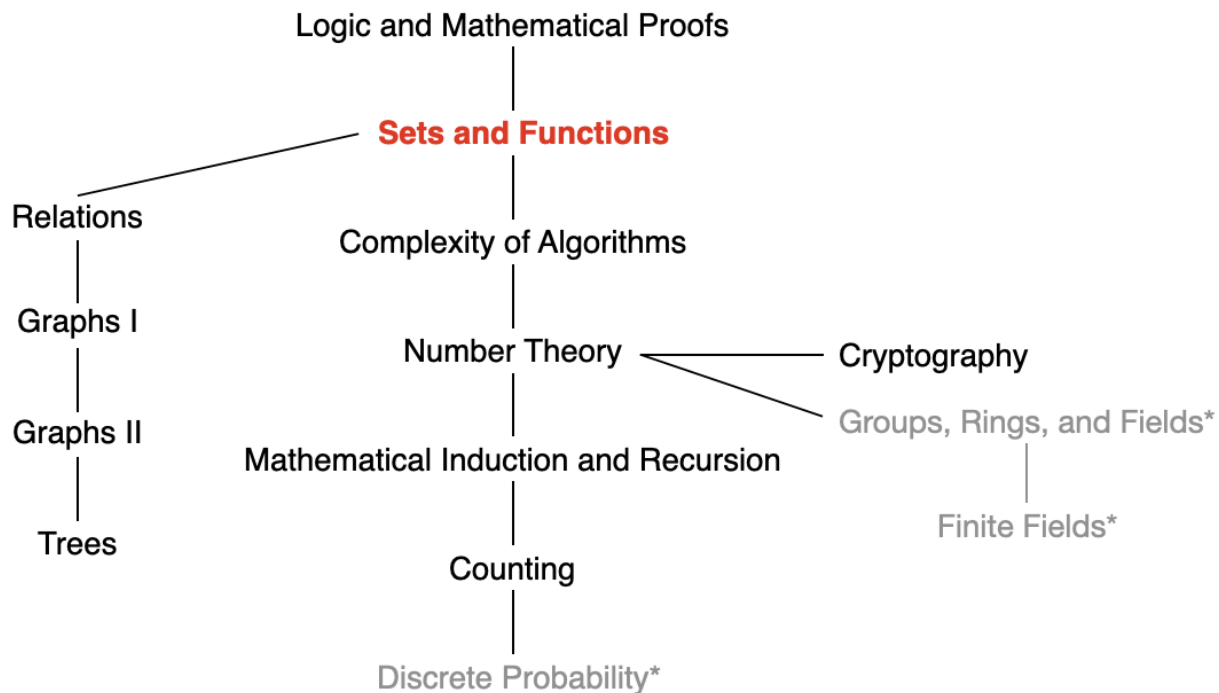
Discrete Mathematics for Computer Science

Lecture 5: Set and Function

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This Lecture



Set and Functions: set, set operations, functions,
sequences and summation, cardinality of sets

Inverse Functions

Let f be a **one-to-one correspondence (bijection)** from the set A to the set B . The **inverse function** of f is the function that assigns an element b belonging to B to the unique element a in A such that $f(a) = b$.

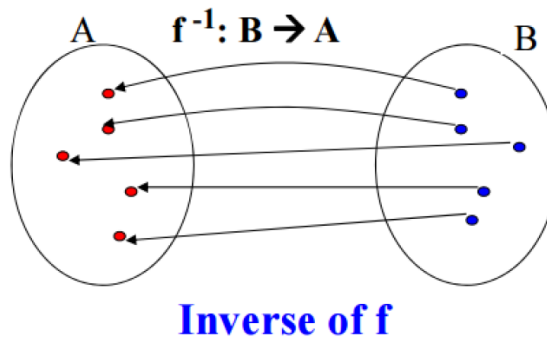
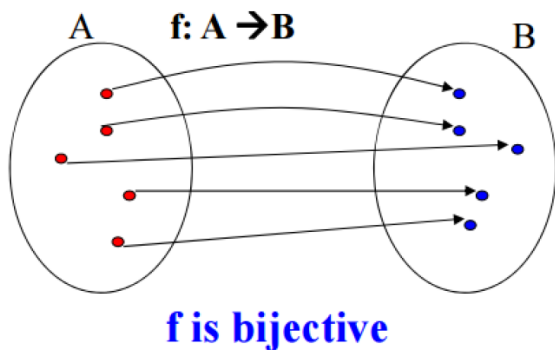
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Hence, $f^{-1}(b) = a$ when $f(a) = b$.

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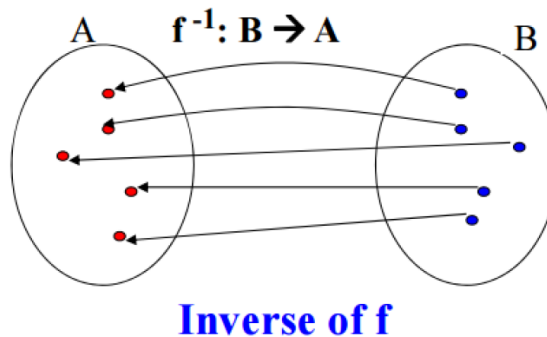
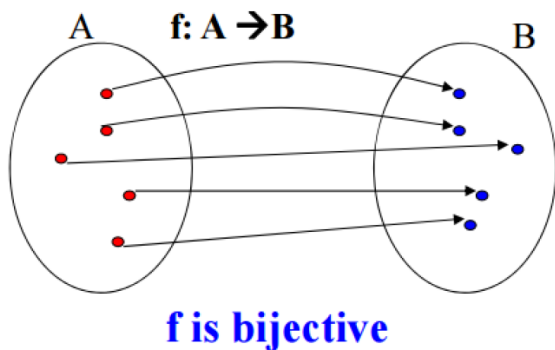


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A bijection is called **invertible**.

Inverse Functions

If f is **not a one-to-one correspondence (bijection)**, it is impossible to define the inverse function of f . Why?

Inverse Functions

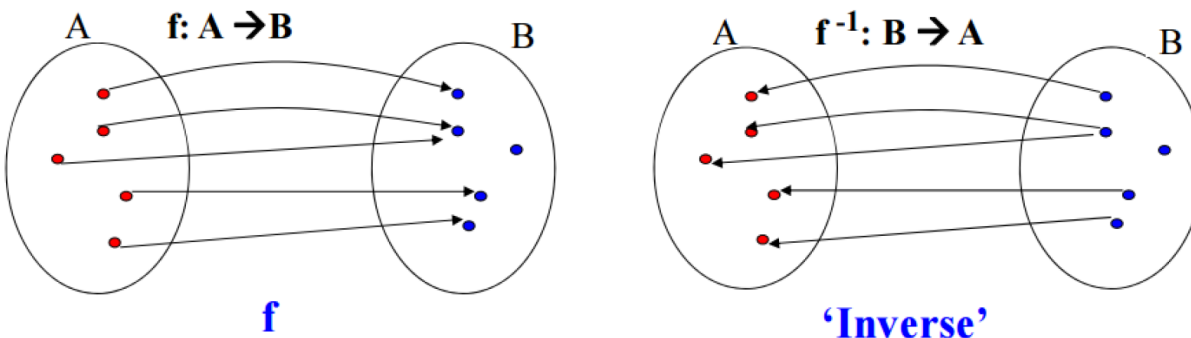
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The inverse is **not a function**: one element of B is mapped to **two different** elements of A .

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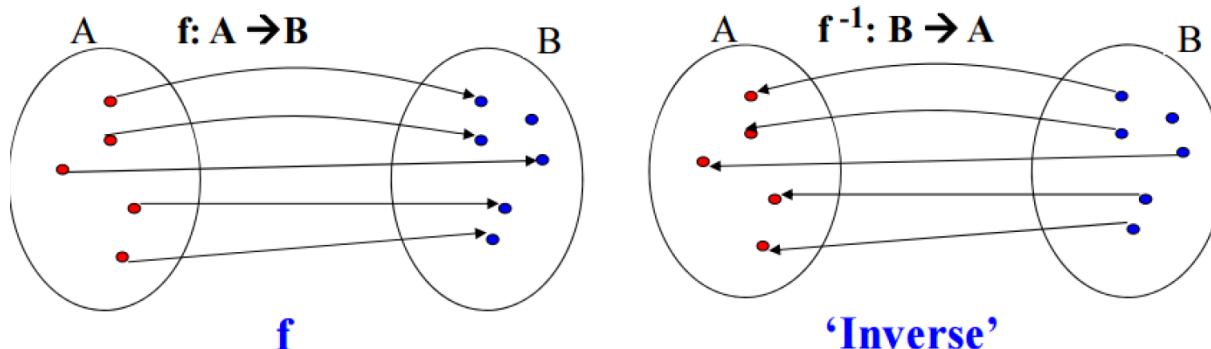
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Assume f is not onto (surjective):

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The inverse is not a function: one element of B is **not assigned** an element of A .

Proof for Inverse Function

1 Prove function f is a bijection: injective, surjective

To show that f is <i>injective</i>	Show that if $f(x) = f(y)$ for all $x, y \in A$, then $x = y$
To show that f is not <i>injective</i>	Find specific elements $x, y \in A$ such that $x \neq y$ and $f(x) = f(y)$
To show that f is <i>surjective</i>	Consider an arbitrary element $y \in B$ and find an element $x \in A$ such that $f(x) = y$
To show that f is not <i>surjective</i>	Find a specific element $y \in B$ such that $f(x) \neq y$ for all $x \in A$

2 If f is a bijection, then it is invertible

3 Determine the inverse function

Inverse Functions: Example 1

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To determine the inverse function, suppose that y is the image of x , so that $y = x + 1$. Then, $x = y - 1$. This means that $y - 1$ is the unique element of \mathbf{Z} that is sent to y by f . Consequently, $f^{-1}(y) = y - 1$.

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Proof: It is invertible, as it is a bijection:

- **Injective:** Consider x and x' . If $f(x) = f(x')$ (i.e., $x^2 = (x')^2$), then we have $x^2 - (x')^2 = (x + x')(x - x') = 0$. Since we consider the set of all nonnegative real numbers, we must have $x = x'$.
- **Surjective:** Consider an arbitrary nonnegative real number y . There exists a nonnegative real number x such that $x = \sqrt{y}$, which means that $x^2 = y$.

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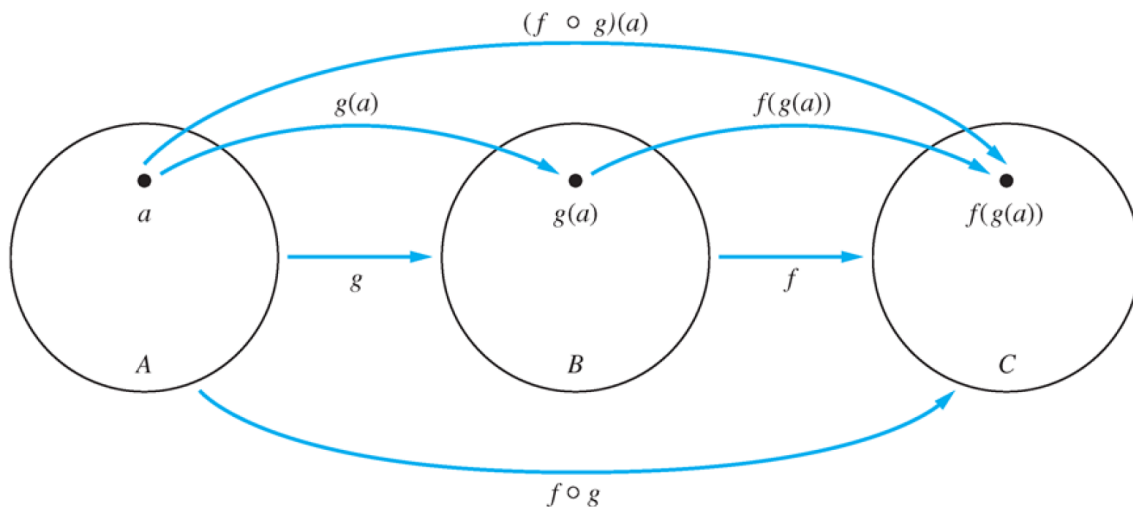
To reverse the function, suppose that y is the image of x , so that $y = x^2$. Then, $x = \sqrt{y}$. Consequently, $f^{-1}(y) = \sqrt{y}$.

Composition of Functions

Let f be a function from B to C and let g be a function from A to B .
The **composition** of the functions f and g , denoted by $f \circ g$, is defined by
 $(f \circ g)(x) = f(g(x))$.

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Composition of Functions

■ Example 1:

Let $A = \{1, 2, 3\}$ and $B = \{a, b, c, d\}$.

$g : A \rightarrow A$ $f : A \rightarrow B$

$1 \mapsto 3$ $1 \mapsto b$

$2 \mapsto 1$ $2 \mapsto a$

$3 \mapsto 2$ $3 \mapsto d$

What is $f \circ g$?

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$f \circ g : A \rightarrow B$

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Composition of Functions

■ Example 2:

Let $f : \mathbf{Z} \rightarrow \mathbf{Z}$ and $g : \mathbf{Z} \rightarrow \mathbf{Z}$, where $f(x) = 2x$ and $g(x) = x^2$.

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Note: In general, the order of composition **matters**.

Composition of Functions

- Suppose that f is a bijection from A to B . Then $f \circ f^{-1} = I_B$ and $f^{-1} \circ f = I_A$, Since

$$(f^{-1} \circ f)(a) = f^{-1}(f(a)) = f^{-1}(b) = a$$

$$(f \circ f^{-1})(b) = f(f^{-1}(b)) = f(a) = b,$$

where I_A , I_B denote the *identity functions* on the sets A and B , respectively.

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Note: Identity function is sometimes denoted by $\iota_A(\cdot)$:

$$\iota_A(x) = x$$

Floor and Ceiling Functions

- The **floor function** assigns a real number x the **largest integer that is $\leq x$** , denoted by $\lfloor x \rfloor$. E.g., $\lfloor 3.5 \rfloor = 3$.
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$$(1a) \lfloor x \rfloor = n \text{ if and only if } n \leq x < n + 1$$

$$(1b) \lceil x \rceil = n \text{ if and only if } n - 1 < x \leq n$$

$$(1c) \lfloor x \rfloor = n \text{ if and only if } x - 1 < n \leq x$$

$$(1d) \lceil x \rceil = n \text{ if and only if } x \leq n < x + 1$$

$$(2) \quad x - 1 < \lfloor x \rfloor \leq x \leq \lceil x \rceil < x + 1$$

$$(3a) \lfloor -x \rfloor = -\lceil x \rceil$$

$$(3b) \lceil -x \rceil = -\lfloor x \rfloor$$

$$(4a) \lfloor x + n \rfloor = \lfloor x \rfloor + n$$

$$(4b) \lceil x + n \rceil = \lceil x \rceil + n$$

Note: n is an integer, x is a real number.

Floor and Ceiling Functions: Example 1

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- $0 \leq \epsilon < \frac{1}{2}$: In this case, $2x = 2n + 2\epsilon$. Since $0 \leq 2\epsilon < 1$, we have $\lfloor 2x \rfloor = 2n$. Similarly, $x + \frac{1}{2} = n + \frac{1}{2} + \epsilon$. Since $0 \leq \frac{1}{2} + \epsilon < 1$, we have $\lfloor x + \frac{1}{2} \rfloor = n$. Thus, $\lfloor 2x \rfloor = 2n$, and $\lfloor x \rfloor + \lfloor x + \frac{1}{2} \rfloor = 2n$.

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- $\frac{1}{2} \leq \epsilon < 1$: In this case, $2x = 2n + 2\epsilon = (2n + 1) + (2\epsilon - 1)$. Since $0 \leq 2\epsilon - 1 < 1$, we have $\lfloor 2x \rfloor = 2n + 1$

Floor and Ceiling Functions: Example 2

Prove or disprove that $\lceil x + y \rceil = \lceil x \rceil + \lceil y \rceil$ for all real numbers x and y .

Floor and Ceiling Functions: Example 2

Prove or disprove that $\lceil x + y \rceil = \lceil x \rceil + \lceil y \rceil$ for all real numbers x and y .

Proof: This statement is false. Consider a counterexample $x = \frac{1}{2}$ and $\frac{1}{2}$. We can find that $\lceil x + y \rceil = 1$, but $\lceil x \rceil + \lceil y \rceil = 2$.

Factorial Function

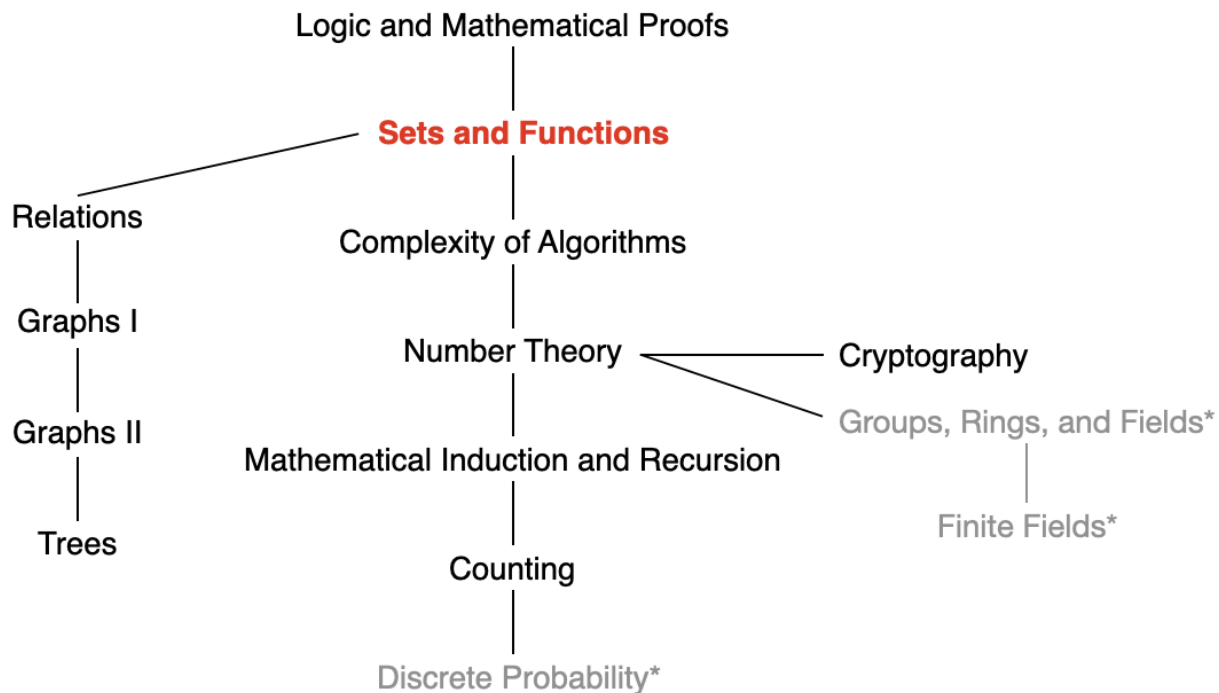
The **factorial function** $f : \mathbf{N} \rightarrow \mathbf{Z}^+$ is the product of the first n positive integers when n is a nonnegative integer, denoted by $f(n) = n!$.

Summary of Function

- Function $f : A \rightarrow B$: an assignment of **exactly one** element of B to **each** element of A
- One-to-one function
- Onto function
- One-to-one correspondence: one-to-one function and onto
- Inverse function
- Floor function, ceiling function, factorial function



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Sequences

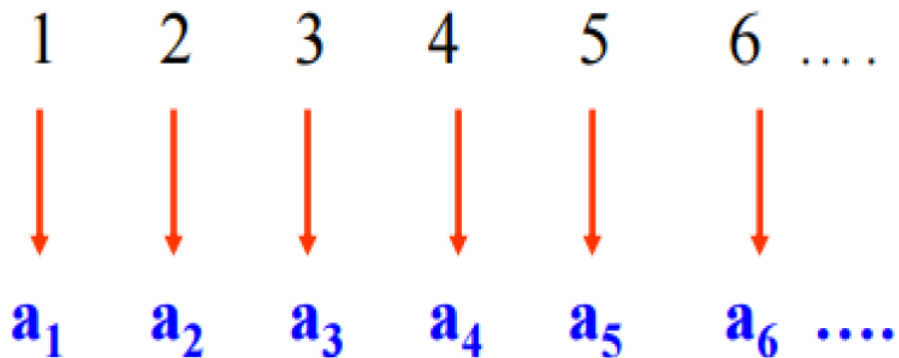
A **sequence** is a **function** from a subset of the set of integers (typically the set $\{0, 1, 2, \dots\}$ or $\{1, 2, 3, \dots\}$) to a set S .

We use the notation a_n to denote the image of the integer n . $\{a_n\}$ represents the ordered list $\{a_1, a_2, a_3, \dots\}$

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Sequences

Examples:

- $a_n = n^2$, where $n = 1, 2, 3, \dots$
- $a_n = (-1)^n$, where $n = 1, 2, 3, \dots$
- $a_n = 2^n$, where $n = 1, 2, 3, \dots$

Geometric Progression

A **geometric progression** is a sequence of the form

$$a, ar, ar^2, \dots, ar^n, \dots$$

where the **initial term** a and the **common ratio** r are real numbers.

Example: $a_n = 3 \times \left(\frac{1}{2}\right)^n$, where $n = 0, 1, 2, 3, \dots$

Arithmetic Progression

An **arithmetic progression** is a sequence of the form

$$a, a + d, a + 2d, a + 3d, \dots, a + nd, \dots$$

where the **initial term** a and **common difference** d are real numbers.

Example: $a_n = -1 + 4n$, where $n = 0, 1, 2, 3, \dots$

Recursively Defined Sequences

1 Providing explicit formulas, e.g., $a_n = -1 + 4n$, where $n = 0, 1, 2, 3, \dots$

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- one or more **initial terms**
- a **rule** for determining **subsequent terms** **from** those that precede them.

The n -th element of the sequence $\{a_n\}$ is defined recursively in terms of the **previous elements** of the sequence and the **initial elements** of the sequence.

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Examples:

- $a_0 = 1$, $a_n = a_{n-1} + 2$ for $n = 1, 2, 3, \dots$
- $f_0 = 0$, $f_1 = 1$, $f_n = f_{n-1} + f_{n-2}$ for $n = 2, 3, 4, \dots$ (Fibonacci sequence)

Summations

The summation of the terms of a sequence is

$$\sum_{j=m}^n = a_m + a_{m+1} + \dots + a_n$$

- j : the index of summation; the choice of the letter is arbitrary
- m : the lower limit of the summation
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$$\sum_{j=1}^n (ax_j + by_j) = a \sum_{j=1}^n x_j + b \sum_{j=1}^n y_j$$

$$\sum_{i=1}^m \sum_{j=1}^n a_i b_j = \sum_{i=1}^m a_i \sum_{j=1}^n b_j$$



Summations

The sum of the first n terms of the arithmetic progression:

$$S_n = \sum_{j=0}^n (a + jd) = (n+1)a + d \sum_{j=0}^n j = (n+1)a + d \frac{n(n+1)}{2}$$

The sum of the first n terms of the geometric progression:

- $r \neq 1$

$$S_n = \sum_{j=0}^n (ar^j) = a \sum_{j=0}^n r^j = \frac{ar^{n+1} - a}{r - 1}$$

- $r = 1$

$$S_n = \sum_{j=0}^n (ar^j) = (n+1)a$$

Summations: Example

■ Examples:

$$\diamond S = \sum_{j=1}^5 (2 + 3j)$$

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$$\diamond S = \sum_{i=1}^4 \sum_{j=1}^2 (2i - j)$$

$$\diamond S = \sum_{j=0}^3 2(5)^j$$

$$\diamond S = \sum_{i=1}^4 \sum_{j=1}^3 ij$$

Summations: Example

■ Examples:

$$\diamond S = \sum_{j=1}^5 (2 + 3j) \quad 55$$

$$\diamond S = \sum_{j=3}^5 (2 + 3j) \quad 42$$

$$\diamond S = \sum_{i=1}^4 \sum_{j=1}^2 (2i - j) \quad 28$$

$$\diamond S = \sum_{j=0}^3 2(5)^j \quad 312$$

$$\diamond S = \sum_{i=1}^4 \sum_{j=1}^3 ij \quad 60$$

Infinite Series

Infinite geometric series can be computed in the closed form for $|x| < 1$.

$$\sum_{k=0}^{\infty} x^k = \lim_{n \rightarrow \infty} \sum_{k=0}^n x^k = \lim_{n \rightarrow \infty} \frac{x^{n+1} - 1}{x - 1} = \frac{1}{1 - x}$$

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$$\sum_{k=0}^{\infty} kx^{k-1} = \frac{1}{(1 - x)^2}$$

Some Useful Summation Formulas

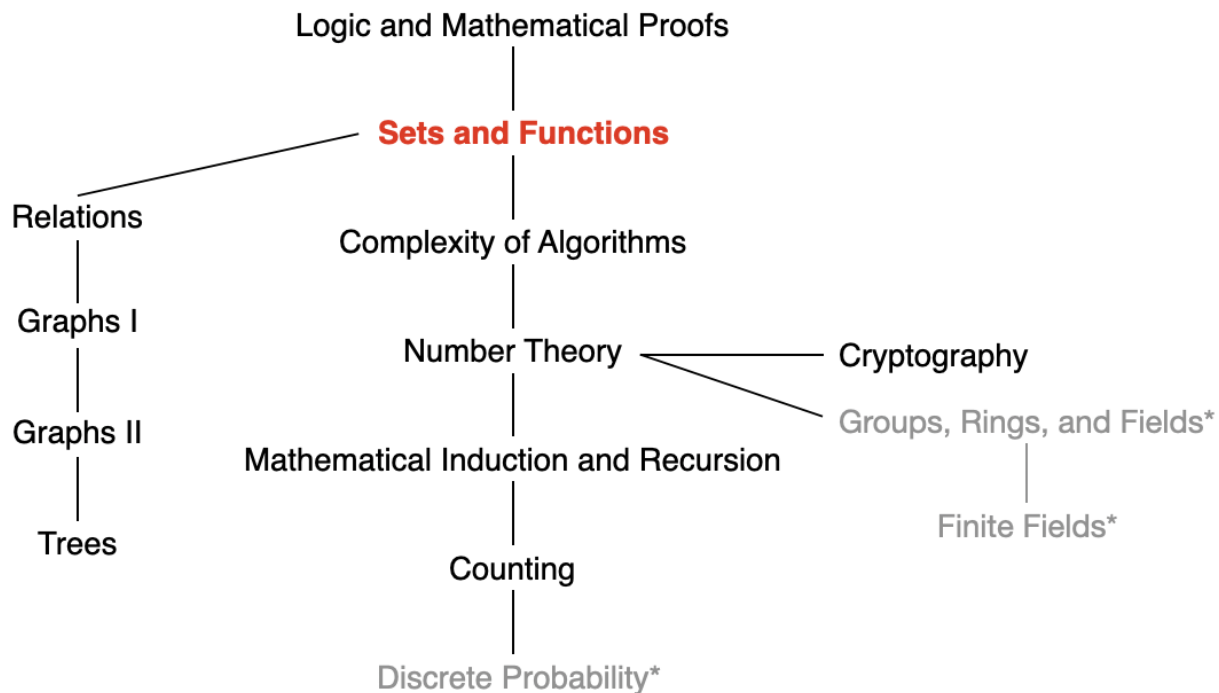
$\sum_{k=0}^n ar^k \ (r \neq 0)$	$\frac{ar^{n+1} - a}{r - 1}, r \neq 1$
$\sum_{k=1}^n k$	$\frac{n(n+1)}{2}$
$\sum_{k=1}^n k^2$	$\frac{n(n+1)(2n+1)}{6}$
$\sum_{k=1}^n k^3$	$\frac{n^2(n+1)^2}{4}$
$\sum_{k=0}^{\infty} x^k, x < 1$	$\frac{1}{1-x}$
$\sum_{k=1}^{\infty} kx^{k-1}, x < 1$	$\frac{1}{(1-x)^2}$



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Cardinality of Sets

Recall: the cardinality of a finite set is defined by the number of the elements in the set.

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The sets A and B have the **same cardinality** if there is a **one-to-one correspondence** between elements in A and B .

If there is a **one-to-one function** from A to B , the cardinality of A is **less than or equal to** the cardinality of B , denoted by $|A| \leq |B|$.

Cardinality of Sets

Recall: the cardinality of a finite set is defined by the number of the elements in the set.

The sets A and B have the **same cardinality** if there is a **one-to-one correspondence** between elements in A and B .

If there is a **one-to-one function** from A to B , the cardinality of A is **less than or equal to** the cardinality of B , denoted by $|A| \leq |B|$.

Moreover, when $|A| \leq |B|$ and A and B have different cardinalities, we say that the cardinality of A is less than the cardinality of B , denoted by $|A| < |B|$.

Countable Sets

A set that is either **finite** or has the **same cardinality as the set of positive integers \mathbb{Z}^+** is called **countable**. A set that is not countable is called uncountable.

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The elements of the set can be **enumerated** and listed.

Hilbert's Paradox: Grand Hotel

The Grand Hotel has countably infinite number of rooms, each occupied by a guest. We can always accommodate a new guest at this hotel.

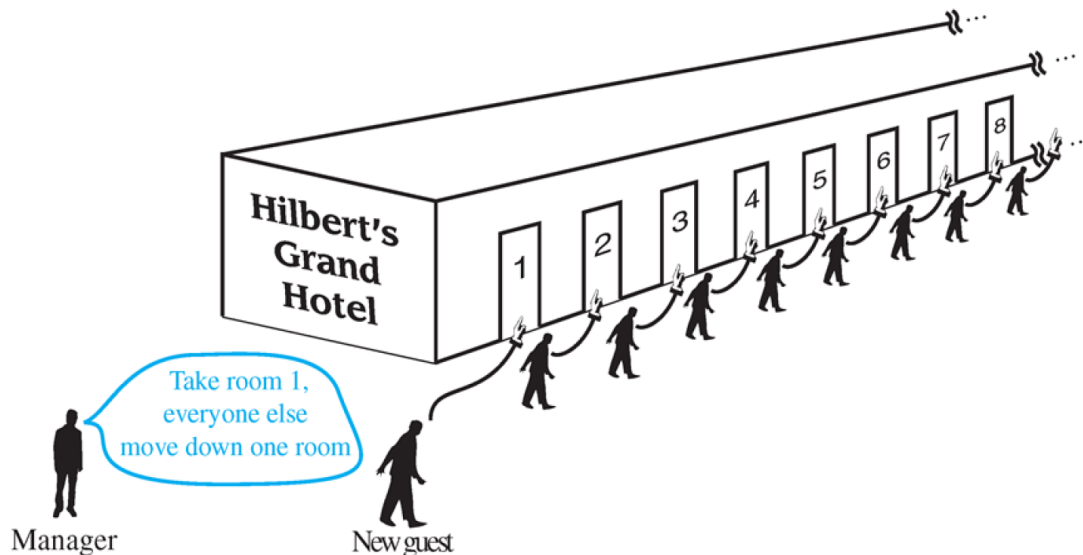


FIGURE 2 A New Guest Arrives at Hilbert's Grand Hotel.

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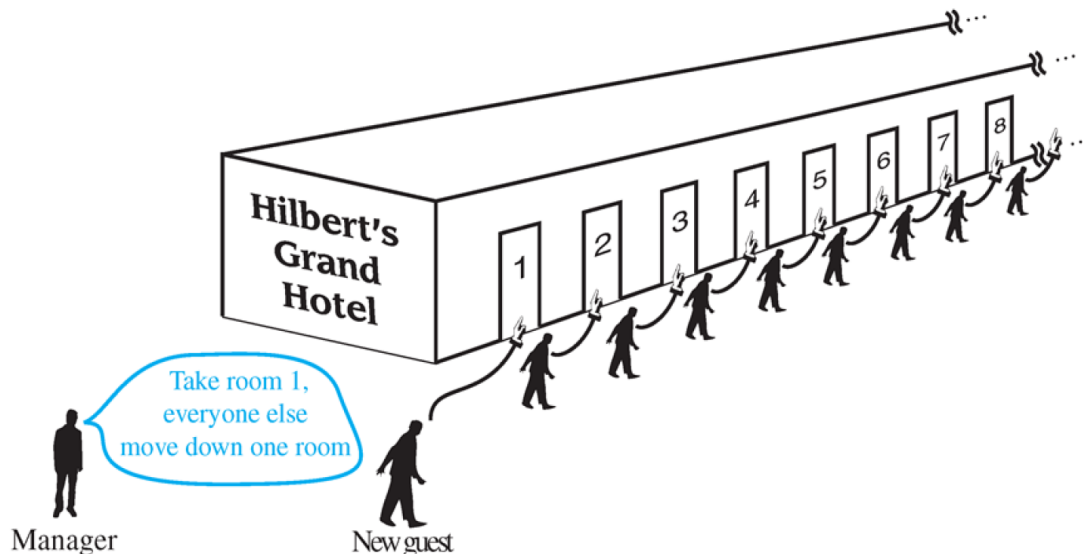


FIGURE 2 A New Guest Arrives at Hilbert's Grand Hotel.

Finitely many room: "All rooms are occupied" is equivalent to "no new guests can be accommodated".

Infinitely many room: This equivalence no longer holds.

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- Onto: For any arbitrary element in $t \in A$, we have an $n = (t + 1)/2 \in \mathbb{Z}^+$ such that $f(n) = t$.

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- All elements are listed

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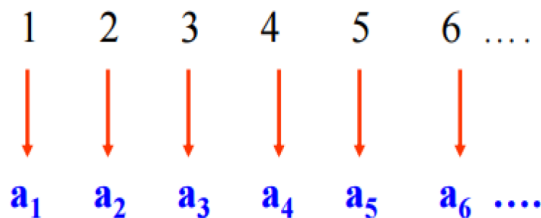
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A **sequence** is a **function** from a subset of the set of integers to a set S .



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- when n is even: $f(n) = n/2$
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Do \mathbf{Z}^+ and \mathbf{Z} have the same cardinality? **Yes**, because there is a one-to-one correspondence between \mathbf{Z}^+ and \mathbf{Z} .

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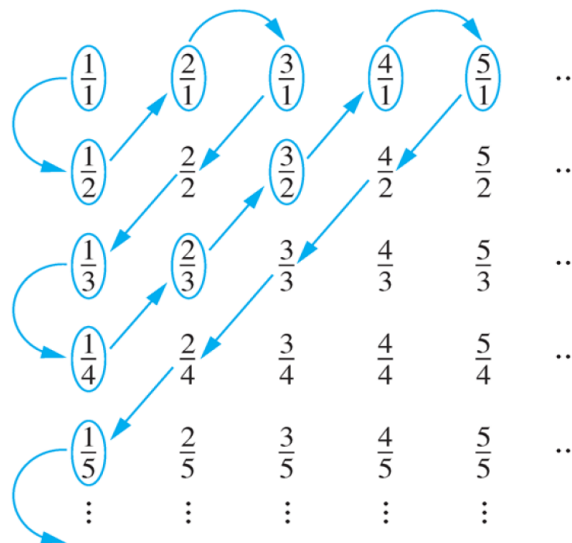
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Solution:

Constructing the list: first list p/q with $p + q = 2$, next list p/q with $p + q = 3$, and so on.

$1, 1/2, 2, 3, 1/3, 1/4, 2/3, \dots$



Countable Sets: Example 4

Theorem: The set of finite strings S over a finite alphabet A is countably infinite. (Assume an alphabetical ordering of symbols in A)

For example, let $A = \{ 'a', 'b', 'c' \}$. Then, set $S = \{ '', 'a', 'b', 'c', 'ab' \dots, 'aaaaa', \dots \}$

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Solution:

We show that the strings can be listed in a sequence. First list

- (i) all the strings of length 0 in alphabetical order.
- (ii) then all the strings of length 1 in lexicographic order.
- (iii) and so on.

This implies a bijection from \mathbb{Z}^+ to S .

Countable Sets: Example 5

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Solution:

Let S be the set of strings constructed from the characters which may appear in a Java program. Use the ordering from the previous example. Take each string in turn

- feed the string into a Java compiler
- if the compiler says YES, this is a syntactically correct Java program, we add this program to the list
- we move on to the next string

In this way, we construct a bijection from \mathbb{Z}^+ to the set of Java programs.

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Theorem: If A and B are countable sets, then $A \cup B$ is also countable.

Uncountable Sets: Example 1

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Theorem: The set of real numbers \mathbf{R} is uncountable.

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Proof by Contradiction: Suppose \mathbf{R} is countable. Then, the interval from 0 to 1 is countable. This implies that the elements of this set can be listed as r_1, r_2, r_3, \dots , where

- $r_1 = 0.d_{11}d_{12}d_{13}d_{14}$

- $r_2 = 0.d_{21}d_{22}d_{23}d_{24}$

- $r_3 = 0.d_{31}d_{32}d_{33}d_{34}$

where all $d_{ij} \in \{0, 1, 2, \dots, 9\}$.

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Proof by Contradiction:

We want to show that not all real numbers in the interval between 0 and 1 are in this list. Form a new number called $r = 0.d_1d_2d_3d_4$, where $d_i = 2$ if $d_{ii} \neq 2$, and $d_i = 3$ if $d_{ii} = 2$.

Example: suppose $r_1 = 0.75243\dots$	$d_1 = 2$
$r_2 = 0.524310\dots$	$d_2 = 3$
$r_3 = 0.131257\dots$	$d_3 = 2$
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Form a new set called $R = b_0b_1b_2b_3\dots$, where $b_i = 0$ if $b_{ii} = 1$, and $b_i = 1$ if $b_{ii} = 0$. R is different from each set in the list. Each bit string is unique, and R and S_i differ in the i -th bit for all i .

Schroder-Bernstein Theorem

Theorem: If A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$.

In other words, if there are one-to-one functions f from A to B and g from B to A , then there is a one-to-one correspondence between A and B , and hence $|A| = |B|$.

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Example: Show that $|(0, 1)| = |(0, 1)|$

$$f(x) = x, g(x) = x/2$$

Computable vs Uncomputable

Definition: We say that a function is **computable** if there is a computer program in some programming language that finds the values of this function. If a function is not computable, we say it is **uncomputable**.

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Cantor's theorem: If S is a set, then $|S| < |P(S)|$.

This Lecture

