## HW5

25. Shannon code. Consider the following method for generating a code for a random variable X which takes on m values  $\{1, 2, \ldots, m\}$  with probabilities  $p_1, p_2, \ldots, p_m$ . Assume that the probabilities are ordered so that  $p_1 \ge p_2 \ge \cdots \ge p_m$ . Define

$$F_i = \sum_{k=1}^{i-1} p_i \,, \tag{5.155}$$

the sum of the probabilities of all symbols less than i. Then the codeword for i is the number  $F_i \in [0, 1]$  rounded off to  $l_i$  bits, where  $l_i = \lceil \log \frac{1}{P_i} \rceil$ .

(a) Show that the code constructed by this process is prefix-free and the average length satisfies

$$H(X) \le L < H(X) + 1$$
. (5.156)

(b) Construct the code for the probability distribution (0.5, 0.25, 0.125, 0.125).

2. How many fingers has a Martian? Let

$$S = \begin{pmatrix} S_1, \ldots, S_m \\ p_1, \ldots, p_m \end{pmatrix}.$$

The  $S_i$ 's are encoded into strings from a D-symbol output alphabet in a uniquely decodable manner. If m=6 and the codeword lengths are  $(l_1,l_2,\ldots,l_6)=(1,1,2,3,2,3)$ , find a good lower bound on D. You may wish to explain the title of the problem.