

Stock Market

- Initial investment Y_0 , daily return ratio r_i , in t -th day, your money is

$$Y_t = Y_0 r_1 \cdot \dots \cdot r_t.$$

- Now if returns ratio r_i are i.i.d., with

$$r_i = \begin{cases} 4, & \text{w.p. } 1/2 \\ 0, & \text{w.p. } 1/2 \end{cases}$$

- So you think the expected return ratio is $E[r_i] = 2$.
- And then

$$E[Y_t] = E[Y_0 r_1 \cdot \dots \cdot r_t] = Y_0 (E[r_i])^t = Y_0 2^t ???$$

- With $Y_0 = 1$, actual return Y_t goes like

1 4 16 0 0 ...

- Why?
 - The 'typical' sequences will end up with 0 return.
 - Occasionally, we got high return.
 - The expected return is increasing.
 - Expectation does not show the typical feature of this random sequence. We can turn to **typical set**.

只看期望的话是纯增的
期望无法体现随机序列的典型特征。
需关注典型集。

Weak Law of Large Numbers

Theorem (Weak Law of Large Numbers)

Suppose that X_1, X_2, \dots, X_n are n independent, identically distributed (i.i.d.) random variables, then

$$\frac{1}{n} \sum_{i=1}^n X_i \rightarrow E[X] \quad \text{in probability,}$$

i.e. for every number $\epsilon > 0$,

$$\lim_{n \rightarrow \infty} \Pr \left[\left| \frac{1}{n} \sum_{i=1}^n X_i - E[X] \right| \leq \epsilon \right] = 1.$$

Asymptotic Equipartition Property (AEP)

Definition (Convergence of random variables)

Given a sequence of random variables, X_1, X_2, \dots , we say that the sequence X_1, X_2, \dots **converges** to a random variable X :

- In probability** if for every $\epsilon > 0$, $\Pr[|X_n - X| \geq \epsilon] \rightarrow 0$
- In mean square** if $E[(X_n - X)^2] \rightarrow 0$
- With probability 1** (a.k.a. **almost surely**) if $\Pr[\lim_{n \rightarrow \infty} X_n = X] = 1$

Theorem 3.1.1 (AEP)

If X_1, X_2, \dots are i.i.d. $\sim p(x)$, then

$$-\frac{1}{n} \log p(X_1, X_2, \dots, X_n) \rightarrow H(X) \quad \text{in probability.}$$

Proof.

Since X_i are i.i.d., so are $\log p(X_i)$. Hence, by the **weak law of large numbers**,

$$\begin{aligned} -\frac{1}{n} \log p(X_1, X_2, \dots, X_n) &= -\frac{1}{n} \sum_i \log p(X_i) \\ &\rightarrow -E[\log p(X)] \quad \text{in probability} \\ &= H(X) \end{aligned}$$

□

Typical Set

Definition

A **typical set** $A_\epsilon^{(n)}$ contains all sequence realizations $(x_1, x_2, \dots, x_n) \in \mathcal{X}^n$ with

$$2^{-n(H(X)+\epsilon)} \leq p(x_1, x_2, \dots, x_n) \leq 2^{-n(H(X)-\epsilon)}.$$

Consequences of AEP

Theorem 3.1.2

- If $(x_1, x_2, \dots, x_n) \in A_\epsilon^{(n)}$, then $H(X) - \epsilon \leq -\frac{1}{n} \log p(x_1, x_2, \dots, x_n) \leq H(X) + \epsilon$.
- $\Pr[(X_1, X_2, \dots, X_n) \in A_\epsilon^{(n)}] > 1 - \epsilon$ for n sufficiently large.
- $|A_\epsilon^{(n)}| \leq 2^{n(H(X)+\epsilon)}$, where $|A|$ denotes the cardinality of the set A .
- $|A_\epsilon^{(n)}| \geq (1 - \epsilon)2^{n(H(X)-\epsilon)}$ for n sufficiently large.

Proof.

- Immediate from the definition of $A_\epsilon^{(n)}$. □

The number of bits used to describe sequences in typical set is approximately $nH(X)$.



Proof.

- By Theorem 3.1.1, the probability of the event $(X_1, X_2, \dots, X_n) \in A_\epsilon^{(n)}$ tends to 1 as $n \rightarrow \infty$. Thus, for any $\delta > 0$, there exists an n_0 such that for all $n \geq n_0$, we have

$$\Pr \left\{ \left| -\frac{1}{n} \log p(X_1, X_2, \dots, X_n) - H(X) \right| < \epsilon \right\} > 1 - \delta.$$

Setting $\delta = \epsilon$, the conclusion follows. □

Proof.

$$\begin{aligned} 3. \quad 1 &= \sum_{\mathbf{x} \in \mathcal{X}^n} p(\mathbf{x}) \geq \sum_{\mathbf{x} \in A_\epsilon^{(n)}} p(\mathbf{x}) \\ &\geq \sum_{\mathbf{x} \in A_\epsilon^{(n)}} 2^{-n(H(X)+\epsilon)} \\ &= 2^{-n(H(X)+\epsilon)} |A_\epsilon^{(n)}|. \end{aligned}$$

□

Proof.

- For sufficiently large n , $\Pr[A_\epsilon^{(n)}] > 1 - \epsilon$, so that

$$\begin{aligned} 1 - \epsilon &< \Pr[A_\epsilon^{(n)}] \\ &\leq \sum_{\mathbf{x} \in A_\epsilon^{(n)}} 2^{-n(H(X)-\epsilon)} \\ &= 2^{-n(H(X)-\epsilon)} |A_\epsilon^{(n)}|. \end{aligned}$$

□

$\forall \epsilon > 0$

$$\lim_{n \rightarrow \infty} \Pr[|-\frac{1}{n} \log p(X_1, \dots, X_n) - H(X)| \leq \epsilon] = 1$$

$$\Rightarrow -\epsilon \leq -\frac{1}{n} \log p(X_1, \dots, X_n) - H(X) \leq \epsilon$$

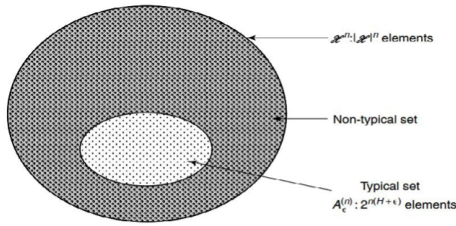
$$\begin{aligned} -n[H(X)-\epsilon] &\geq \log p(X_1, \dots, X_n) \geq -n[H(X)+\epsilon] \\ 2^{-n[H(X)-\epsilon]} &\geq p(X_1, \dots, X_n) \geq 2^{-n[H(X)+\epsilon]} \end{aligned}$$

上限是任何情况都成立, 下限只在 n 是够大时成立.

Typical set diagram

This enables us to divide all sequences into two sets

- Typical set: high probability to occur, sample entropy is close to true entropy
so we will focus on analyzing sequences in typical set
- Non-typical set: small probability, can ignore in general



发生概率高, 采样熵接近于真实熵

- 一般忽略非典型集

Asymptotic Equipartition Property (AEP)

Theorem 3.2.1

Let X_1, X_2, \dots, X_n be i.i.d. random variables with distribution $p(x)$, and $X^n = X_1 X_2 \dots X_n$. For **arbitrarily small** $\epsilon > 0$, there exists a **code** that maps every realization $x^n = x_1 x_2 \dots x_n$ of X^n into one binary string, such that the mapping is one-to-one (and therefore invertible) and

$$E \left[\frac{1}{n} \ell(X^n) \right] \leq H(X) + \epsilon$$

for **a sufficiently large** n .

Proof.

Description in typical set requires no more than $n(H(X) + \epsilon) + 1$ bits (correction of 1 bit because of integrality).

Description in atypical set $A_\epsilon^{(n)c}$ requires no more than $n \log |\mathcal{X}| + 1$ bits.

Add **another bit** to indicate whether in $A_\epsilon^{(n)}$ or not to get whole description. □

Proof.

Let $\ell(x^n)$ be the length of the binary description of x^n . Then, $\forall \epsilon > 0$, there exists n_0 s.t. $\forall n > n_0$,

$$\begin{aligned} E(\ell(X^n)) &= \sum_{x^n} p(x^n) \ell(x^n) \\ &= \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) \ell(x^n) + \sum_{x^n \in A_\epsilon^{(n)c}} p(x^n) \ell(x^n) \\ &\leq \sum_{x^n \in A_\epsilon^{(n)}} p(x^n) (n(H + \epsilon) + 2) + \sum_{x^n \in A_\epsilon^{(n)c}} p(x^n) (n \log |\mathcal{X}| + 2) \\ &= \Pr[A_\epsilon^{(n)}] (n(H + \epsilon) + 2) + \Pr[A_\epsilon^{(n)c}] (n \log |\mathcal{X}| + 2) \\ &\leq n(H + \epsilon) + \epsilon n (\log |\mathcal{X}|) + 2 \\ &= n(H + \epsilon') \end{aligned}$$

where $\epsilon' = \epsilon + \epsilon \log |\mathcal{X}| + \frac{2}{n}$ can be made arbitrarily small by choosing n properly. □