

Probability

- **Random variable** (r.v.): variable that takes on (discrete) values with certain probabilities
- **Probability distribution**: for an r.v. specifies the probabilities with which the variable takes on each possible value
 - Each probability must be between 0 and 1
 - The probabilities must sum to 1
- **Event**: a particular occurrence in some experiment
 - $\Pr[E]$: probability of event E
- **Conditional probability**: probability that one event occurs, given that some other even occurred
 - $\Pr[A \mid B] = \Pr[A \text{ and } B] / \Pr[B]$
- Two r.v.'s X, Y are **independent** if for all x, y :
 $\Pr[X = x \mid Y = y] = \Pr[X = x]$
- **Law of total probability**: say E_1, \dots, E_n are a partition of all possibilities. Then for any A :
 $\Pr[A] = \sum_i \Pr[A \text{ and } E_i] = \sum_i \Pr[A \mid E_i] \cdot \Pr[E_i]$
- **Bayes's theorem**
 $\Pr[A \mid B] = \Pr[B \mid A] \cdot \Pr[A] / \Pr[B]$

Number Theory

费马小定理 (Fermat's Law)

设 p 为素数, a 是任意整数, 则 $a^p \equiv a \pmod{p}$

若 a 不是 p 的倍数, 则 $a^{p-1} \equiv 1 \pmod{p}$

证明:

Lemma 1: $a, b, c \in \mathbb{Z}, m \in \mathbb{Z}^*$ 且 $\gcd(m, c) = 1$. 则当 $ac \equiv bc \pmod{m}$ 时, $a \equiv b \pmod{m}$

$$\because ac \equiv bc \pmod{m} \quad \therefore (a-b)c \equiv 0 \pmod{m}$$

$$\because \gcd(m, c) = 1 \quad \therefore a-b \equiv 0 \pmod{m} \Rightarrow a \equiv b \pmod{m}$$

Lemma 2:

贝祖定理 (Bezout's Law):

设 a, b 为不全为 0 的整数, 则存在整数 x, y , 使得 $ax + by = \gcd(a, b)$.

证明: ① 若 a, b 其一为 0, 设 $b = 0$, 则 $\gcd(a, b) = a$ 满足

② 若 $a, b \neq 0$. 设 $a, b > 0$ 且 $a \geq b$, $\gcd(a, b) = d$. 则有

$$ax + by = d \Rightarrow a'x + b'y = 1, \text{ 其中 } \gcd(a', b') = 1.$$

(贝祖定理的一个引理: 若 $\gcd(a, b) = 1$, 则 $\exists x, y \in \mathbb{Z}, ax + by = 1$)

由辗转相除法可知: $\gcd(a, b) = \gcd(b, r_1) = \gcd(r_1, r_2) = \dots = \gcd(r_{n-1}, r_n)$
也即: $a = q_1 b + r_1$
 $b = q_2 r_1 + r_2$
 \vdots
 $r_{n-2} = q_{n-1} r_{n-1} + r_n$
 $r_{n-1} = q_n r_n$
 $= \gcd(r_n, 0) = r_n$

假设在 r_{n-1} 与 r_n 互质时退出, 即

$$r_{n-2} = x_n r_{n-1} + 1 \Rightarrow 1 = r_{n-2} - x_n r_{n-1}$$

将 $r_{n-1} = r_{n-3} - x_{n-1} r_{n-2}$ 代回上式, 有

$$1 = (1 + x_n x_{n-1}) r_{n-2} - x_n r_{n-3}$$

不断回代可得 x, y 的值从而 $ax + by = 1$.

We can use [extended Euclidean algorithm](#) to find Bezout's identity.

Example: Express $\gcd(252, 198) = 18$ as a linear combination of 252 and 198.

Solution: To show that $\gcd(252, 198) = 18$, the Euclidean algorithm uses these divisions:

$$\begin{aligned} 252 &= 1 \cdot 198 + 54 \\ 198 &= 3 \cdot 54 + 36 \\ 54 &= 1 \cdot 36 + 18 \\ 36 &= 2 \cdot 18. \end{aligned}$$

Substituting the above expressions:

$$18 = 54 - 1 \cdot 36 = 54 - 1 \cdot (198 - 3 \cdot 54) = 4 \cdot 54 - 1 \cdot 198.$$

$$18 = 4 \cdot (252 - 1 \cdot 198) - 1 \cdot 198 = 4 \cdot 252 - 5 \cdot 198.$$

逆元:

对正整数 a, m , 若有 $ax \equiv 1 \pmod{m}$, 则 x 的最小正整数解为 a 模 m 的逆元.

$$\gcd(a, m) = 1, m > 1 \Leftrightarrow ax \equiv 1 \pmod{m} \text{ 有解}$$

逆元求法: ① 费马小定理: $a \cdot a^{p-2} \equiv 1 \pmod{p} \Rightarrow a^{p-2} \pmod{p}$ 为 a 的逆元

② 扩展欧几里得算法 (贝祖定理利用): 求解 $ax + by = \gcd(a, b)$ 的 x, y .

Using **extended Euclidean algorithm**:

Example: Find an inverse of 101 modulo 4620. That is, find \bar{a} such that $\bar{a} \cdot 101 \equiv 1 \pmod{4620}$.

$$\begin{array}{ll} 4620 = 45 \cdot 101 + 75 & 1 = 3 - 1 \cdot 2 \\ 101 = 1 \cdot 75 + 26 & 1 = 3 - 1 \cdot (23 - 7 \cdot 3) = -1 \cdot 23 + 8 \cdot 3 \\ 75 = 2 \cdot 26 + 23 & 1 = -1 \cdot 23 + 8 \cdot (26 - 1 \cdot 23) = 8 \cdot 26 - 9 \cdot 23 \\ 26 = 1 \cdot 23 + 3 & 1 = 8 \cdot 26 - 9 \cdot (75 - 2 \cdot 26) = 26 \cdot 26 - 9 \cdot 75 \\ 23 = 7 \cdot 3 + 2 & 1 = 26 \cdot (101 - 1 \cdot 75) - 9 \cdot 75 \\ 3 = 1 \cdot 2 + 1 & = 26 \cdot 101 - 35 \cdot 75 \\ 2 = 2 \cdot 1 & 1 = 26 \cdot 101 - 35 \cdot (4620 - 45 \cdot 101) \\ & = -35 \cdot 4620 + 1601 \cdot 101 \end{array}$$

That $-35 \cdot 4620 + 1601 \cdot 101 = 1$ tells us that -35 and 1601 are Bezout coefficients of 4620 and 101 . We have

$$1 \pmod{4620} = 1601 \cdot 101 \pmod{4620}$$

Thus, 1601 is an inverse of 101 modulo 4620 .

欧拉函数:

$\varphi(n)$, 表示小于等于 n 和 n 互质的数的个数.

当 n 为质数时, $\varphi(n) = n - 1$

若 $\gcd(a, b) = 1$, 则 $\varphi(a \times b) = \varphi(a) \times \varphi(b)$

当 n 是奇数时, $\varphi(2n) = \varphi(n)$

欧拉定理: 若 $\gcd(a, m) = 1$, 则 $a^{\varphi(m)} \equiv 1 \pmod{m}$

中国剩余定理:

$$\text{同余方程组} \begin{cases} x \equiv a_1 \pmod{m_1} \\ x \equiv a_2 \pmod{m_2} \\ \vdots \\ x \equiv a_n \pmod{m_n} \end{cases} \text{ 中 } m_i \text{ 两两互质}$$

有唯一解 $x \pmod{m}$ ($m = m_1 m_2 \dots m_n$)

步骤: ① 计算模数之积: $m = \prod_i m_i$

② 对第 i 个方程, 计算

i) $M_i = \frac{m}{m_i}$

ii) M_i 模 m_i 下的逆元 M_i^{-1}

iii) $C_i = M_i M_i^{-1}$

③ 解为 $x = \sum_{i=1}^k a_i C_i \pmod{m}$

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

① Let $m = 3 \cdot 5 \cdot 7 = 105$, $M_1 = m/3 = 35$, $M_2 = m/5 = 21$, and $M_3 = m/7 = 15$.

② Compute the inverse of M_k modulo m_k :

▶ $35 \cdot 2 \equiv 1 \pmod{3}$ $y_1 = 2$

▶ $21 \equiv 1 \pmod{5}$ $y_2 = 1$

▶ $15 \equiv 1 \pmod{7}$ $y_3 = 1$

③ Compute a solution x :

$$x = 2 \cdot 35 \cdot 2 + 3 \cdot 21 \cdot 1 + 2 \cdot 15 \cdot 1 \equiv 233 \equiv 23 \pmod{105}$$

④ The solutions are all integers x that satisfy $x \equiv 23 \pmod{105}$.

回代法:

(1) $x \equiv 1 \pmod{5}$

(2) $x \equiv 2 \pmod{6}$

(3) $x \equiv 3 \pmod{7}$

According to (1), $x = 5t + 1$, where t is an integer.

Substituting this expression into (2), we have $5t + 1 \equiv 2 \pmod{6}$, which means that $t \equiv 5 \pmod{6}$. Thus, $t = 6u + 5$, where u is an integer.

Substituting $x = 5t + 1$ and $t = 6u + 5$ into (3), we have $30u + 26 \equiv 3 \pmod{7}$, which implies that $u \equiv 6 \pmod{7}$. Thus, $u = 7v + 6$, where v is an integer.

Thus, we must have $x = 210v + 206$. Translating this back into a congruence,

$$x \equiv 206 \pmod{210}.$$

