



# Solving Problems by Searching

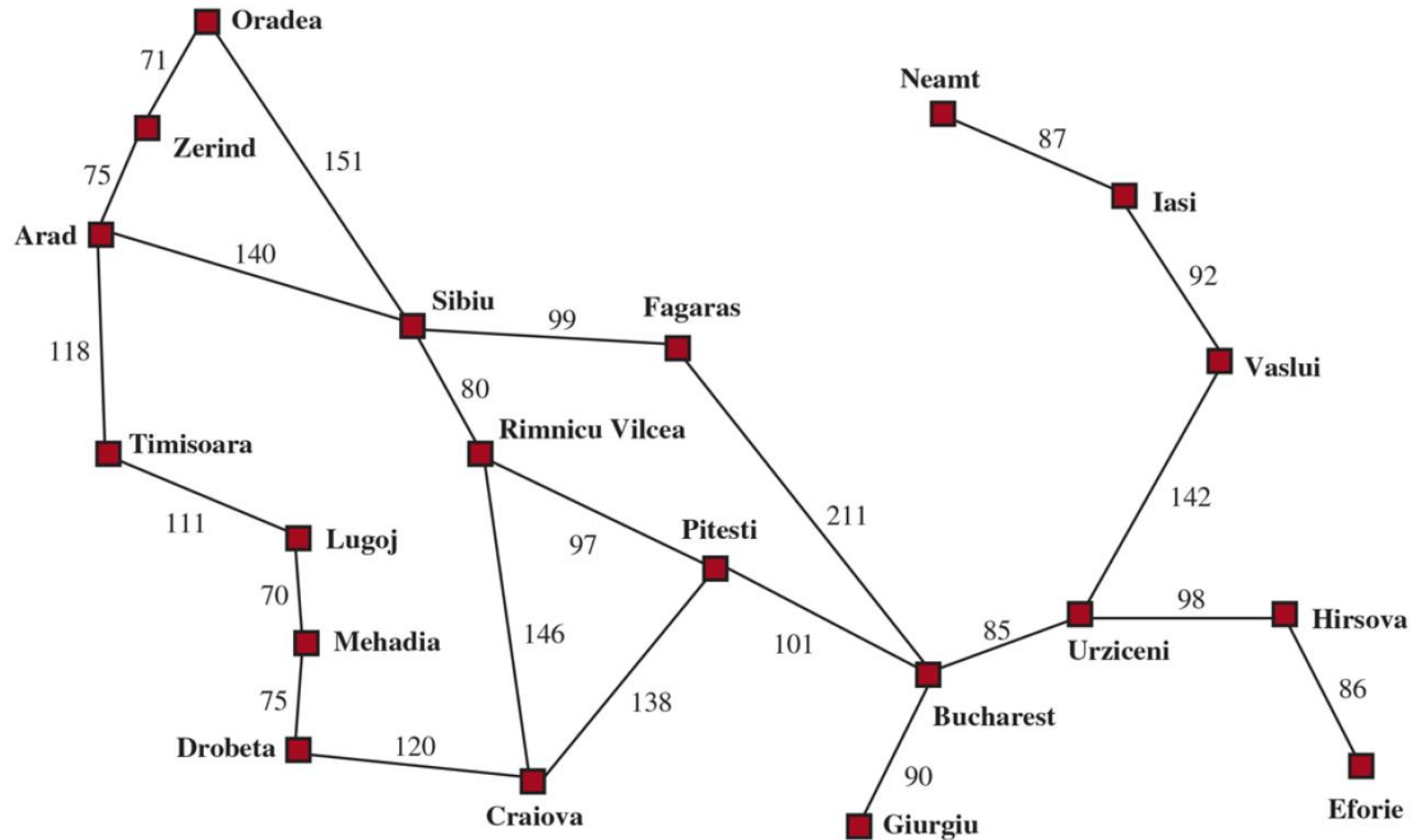
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# Problem-solving Agent

- ▶ When the correct action to take is not immediately obvious,
  - ▶ an agent may need to **plan ahead**:
  - ▶ to consider a sequence of actions that form a path to a *goal state*
  - ▶ The computational process it undertakes is called **search**.

# A Touring Vacation in Romania



A simplified road map of part of Romania, with road distances in miles.



# Map: The Information About The World

- ▶ With that information, the agent can follow this four-phase problem-solving process:
  - ▶ GOAL FORMULATION : e.g. Bucarest
  - ▶ PROBLEM FORMULATION :
    - ▶ considers the actions of traveling from one city to an adjacent city
    - ▶ The only fact about the state of the world (that will change due to an action) is *the current city*
  - ▶ SEARCH :
    - ▶ simulates sequences of actions in its model
    - ▶ searching until it finds a sequence of actions that reaches the goal (a **solution**)
  - ▶ EXECUTION (the actions the solution)



# A Search Problem

- **State space** : A set of possible states that the environment can be in
- **The initial state** that the agent starts in
- A set of one or more **goal states**
- **ACTION(s)** returns a finite set of actions that can be executed in s
- A **transition model** :  $\text{RESULT}(\text{Arad}, \text{ToZerind}) = \text{Zerind}$
- An **action cost function**
  - reflects performance measure (e.g. distance, time)
- An **optimal solution** has the lowest cost among all solutions



# Formulating problems

- The problem of getting to Bucharest is a **model**
  - not the real thing.
  - The process of removing detail from a representation is called **abstraction**

# 8-puzzle

7	2	4
5		6
8	3	1

Start State

	1	2
3	4	5
6	7	8

Goal State

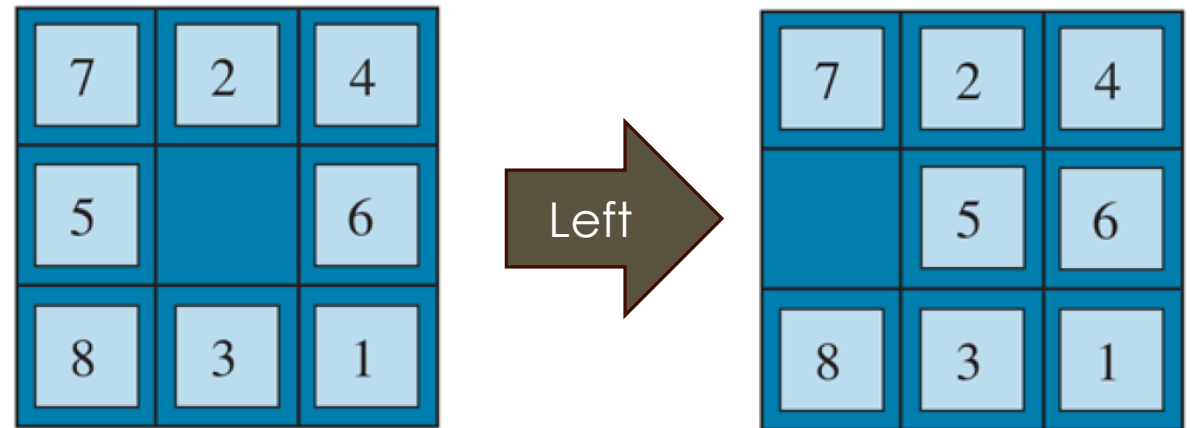
A typical instance of the 8-puzzle.

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# 8-puzzle

- **STATES:** A state description specifies the location of each of the tiles.
- **INITIAL STATE:** Any state can be designated as the initial state
  - State space is partitioned into two halves.
- **ACTIONS:** blank space moving Left, Right, Up, or Down.
  - If the blank is at an edge or corner then not all actions will be applicable.
- **TRANSITION MODEL:** Maps a state and action to a resulting state
- **GOAL STATE**
- **ACTION COST:** Each action costs 1.





# Search Algorithms

- Algorithms that superimpose a search tree over the state-space graph
  - forming various paths from the initial state
  - trying to find a path that reaches a goal state

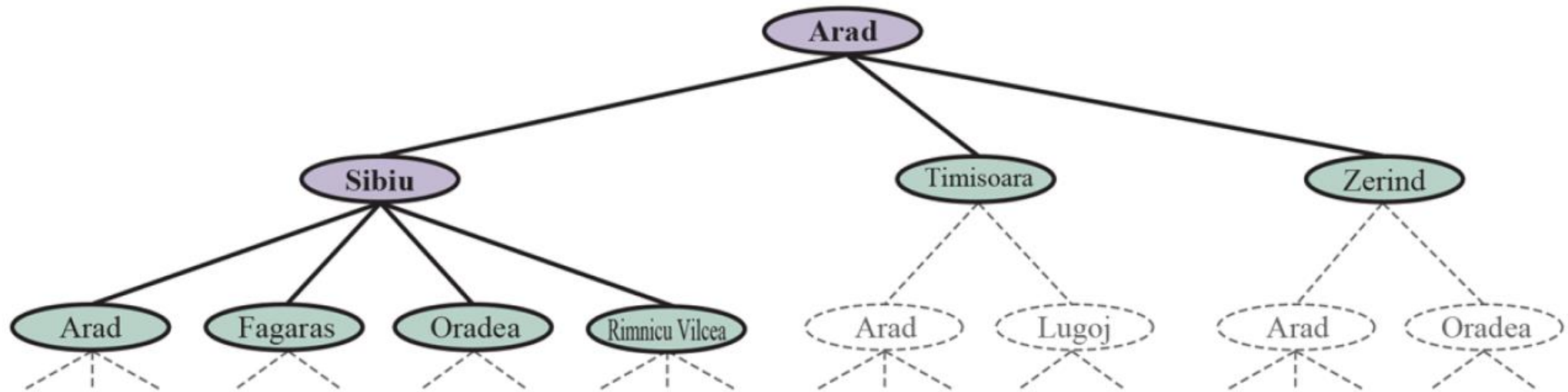
## **The state space**

- describes the (possibly infinite) set of states in the world, and
- the actions that allow transitions from one state to another

## **The search tree**

- describes paths between these states, reaching towards the goal
- may have multiple paths to any given state
- but each node in the tree has a unique path back to the root

# A Touring Vacation in Romania



# Best-first search : A general approach

```
function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure  
  node  $\leftarrow$  NODE(STATE=problem.INITIAL)  
  frontier  $\leftarrow$  a priority queue ordered by f, with node as an element  
  reached  $\leftarrow$  a lookup table, with one entry with key problem.INITIAL and value node  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    if problem.IS-GOAL(node.STATE) then return node  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if s is not in reached or child.PATH-COST < reached[s].PATH-COST then  
        reached[s]  $\leftarrow$  child  
        add child to frontier  
  return failure
```

```
function EXPAND(problem, node) yields nodes  
  s  $\leftarrow$  node.STATE  
  for each action in problem.ACTIONS(s) do  
    s'  $\leftarrow$  problem.RESULT(s, action)  
    cost  $\leftarrow$  node.PATH-COST + problem.ACTION-COST(s, action, s')  
    yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)
```

# Measuring problem-solving performance

- ▶ **COMPLETENESS:**

- ▶ Is the algorithm guaranteed to find a solution when there is one, and to correctly report failure when there is not?

- ▶ **COST OPTIMALITY:**

- ▶ Does it find a solution with the lowest path cost of all solutions?

- ▶ **TIME COMPLEXITY**

- ▶ **SPACE COMPLEXITY**

- ▶ In many AI problems, the state-space graph is represented only *implicitly* by the initial state, actions, and transition model.

- ▶ For an implicit state space, complexity can be measured in terms of

- ▶ the **depth** or number of actions in an optimal solution;  $d$
- ▶ the **maximum number of actions** in any path;  $m$
- ▶ and the **branching factor** or number of successors of a node that need to be considered;  $b$

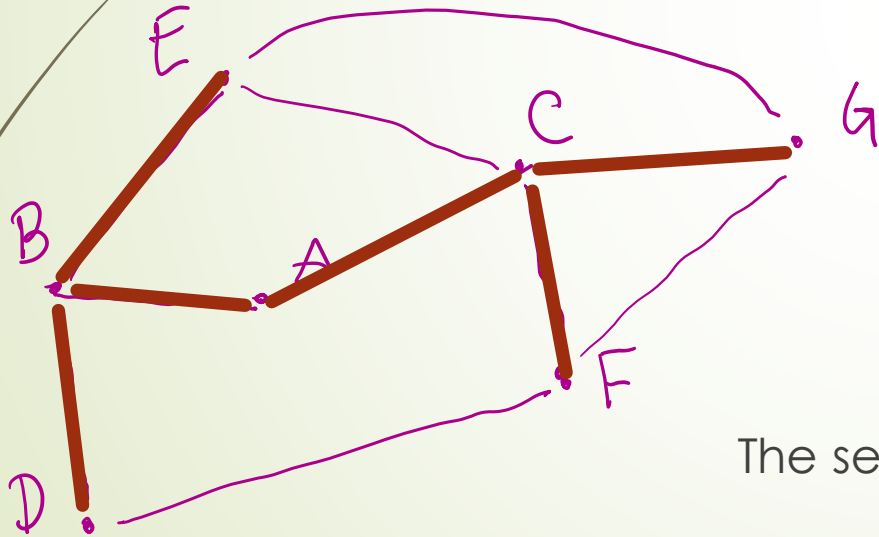
# Uninformed Search Strategies

- An **uninformed** search algorithm is given *no clue about how close a state is to the goal(s)*.
  - For example, consider our agent in **Arad** with the goal of reaching **Bucharest**.
  - An **uninformed agent** with no knowledge of Romanian geography has no clue whether going to **Zerind** or **Sibiu** is a better first step.
- In contrast, an **informed agent** who knows the location of each city knows that **Sibiu** is much closer to **Bucharest** and thus more likely to be on the *shortest path*.

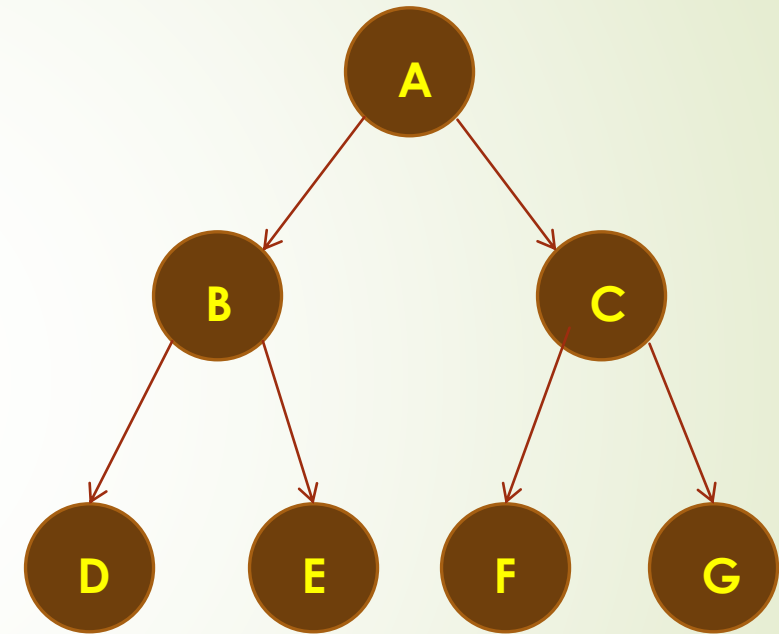


# Breadth First Search

- When all actions have the same cost
- This is a systematic search strategy that is therefore complete even on infinite state spaces.



The sequence of states being searched



# Breadth-first Search

- always finds a solution with a minimal number of actions
  - When it is generating nodes at depth  $d$ , it has already generated all the nodes at depth  $d-1$ .
  - If one of them were a solution, it would have been found.
- It is complete.
- Suppose that the solution is at depth  $d$ , with branching factor  $= b$ .
  - The total number of nodes generated is  $1 + b + b^2 + b^3 + \dots + b^d = O(b^d)$
  - All the nodes remain in memory
    - so both time and space complexity are  $O(b^d)$



# Breadth-first Search

```
function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure  
  node  $\leftarrow$  NODE(problem.INITIAL)  
  if problem.IS-GOAL(node.STATE) then return node  
  frontier  $\leftarrow$  a FIFO queue, with node as an element  
  reached  $\leftarrow$  {problem.INITIAL}  
  while not IS-EMPTY(frontier) do  
    node  $\leftarrow$  POP(frontier)  
    for each child in EXPAND(problem, node) do  
      s  $\leftarrow$  child.STATE  
      if problem.IS-GOAL(s) then return child  
      if s is not in reached then  
        add s to reached  
        add child to frontier  
  return failure
```

avoiding repeated  
states



# Simplest Algorithm Structure of BFS

**s = initial\_state**

**while** not Goal(s)

**for each** successor\_state x of s

    enqueue(x)

s = dequeue()

# Tracing back for path

PRINT-PATH( $G, s, v$ )

```
1  if  $v == s$ 
2      print  $s$ 
3  elseif  $v.\pi == \text{NIL}$ 
4      print “no path from”  $s$  “to”  $v$  “exists”
5  else PRINT-PATH( $G, s, v.\pi$ )
6      print  $v$ 
```

NOTE:

- $G$  is state space graph
- $s$  is initial state
- $v$  is the current state
- $\pi$  is the parent state according to Breadth-first Search