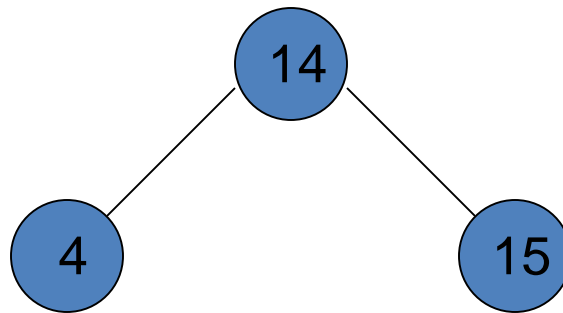


# Data Structures & Algorithms

## Week 10: Binary Search Tree

# Traversing a Binary Tree

- Ways to print a 3 node tree:



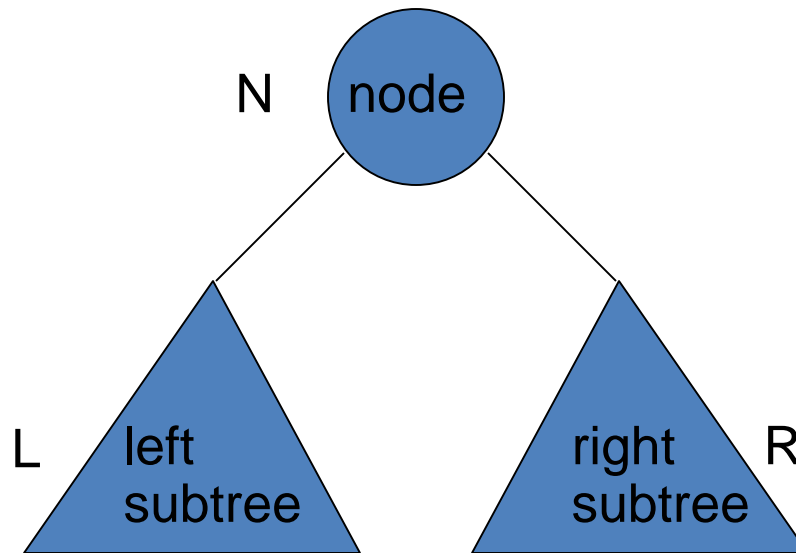
(4, 14, 15), (4,15,14)

(14,4,15), (14,15,4)

(15,4,14), (15,14,4)

# Traversing a Binary Tree

- In case of the general binary tree:



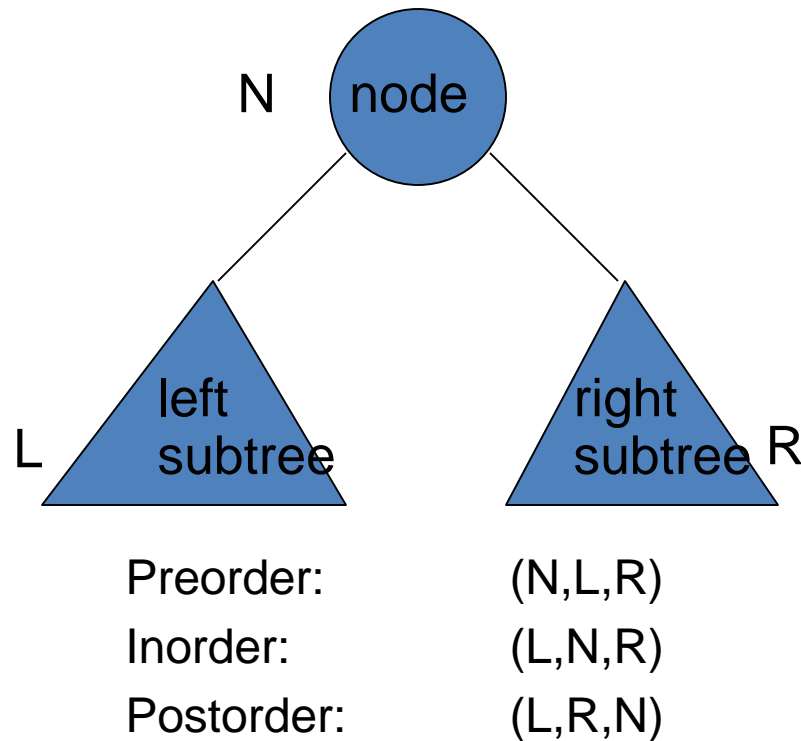
(L,N,R), (L,R,N)

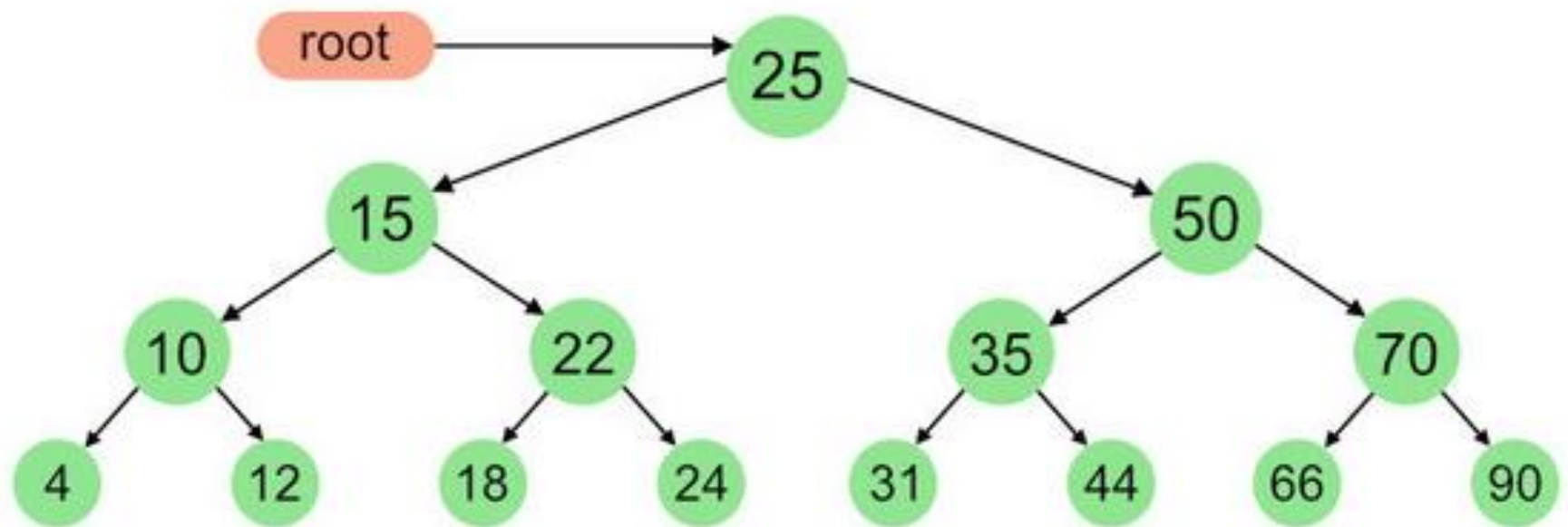
(N,L,R), (N,R,L)

(R,L,N), (R,N,L)

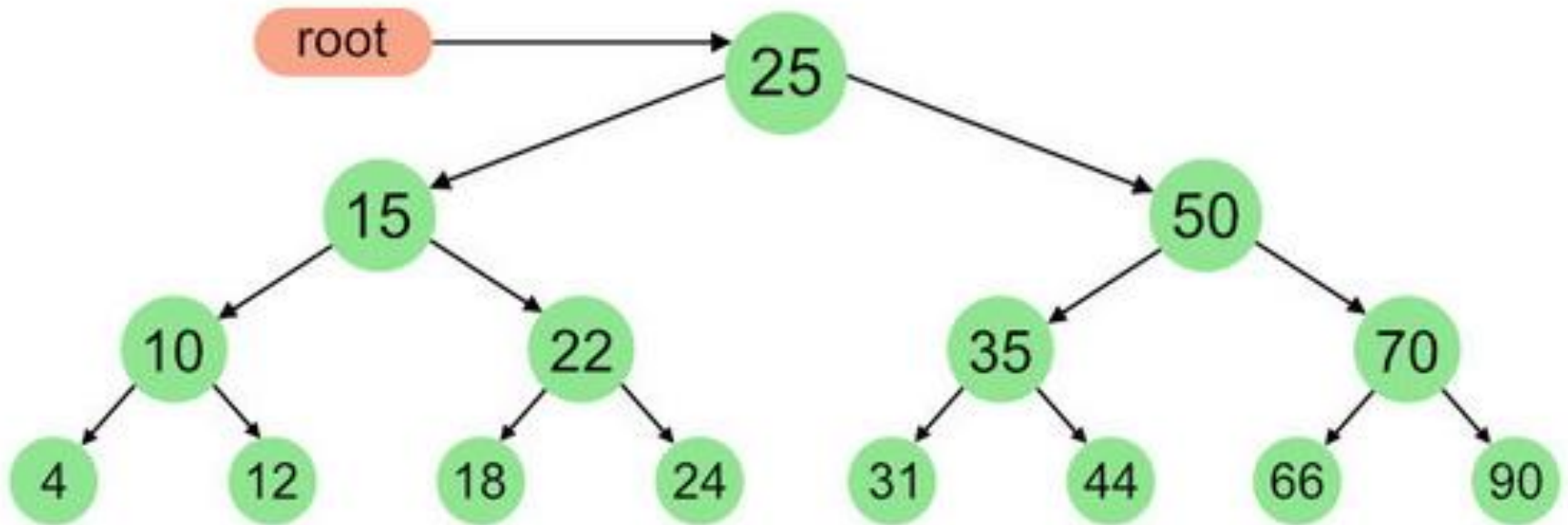
# Traversing a Binary Tree

- Three common ways

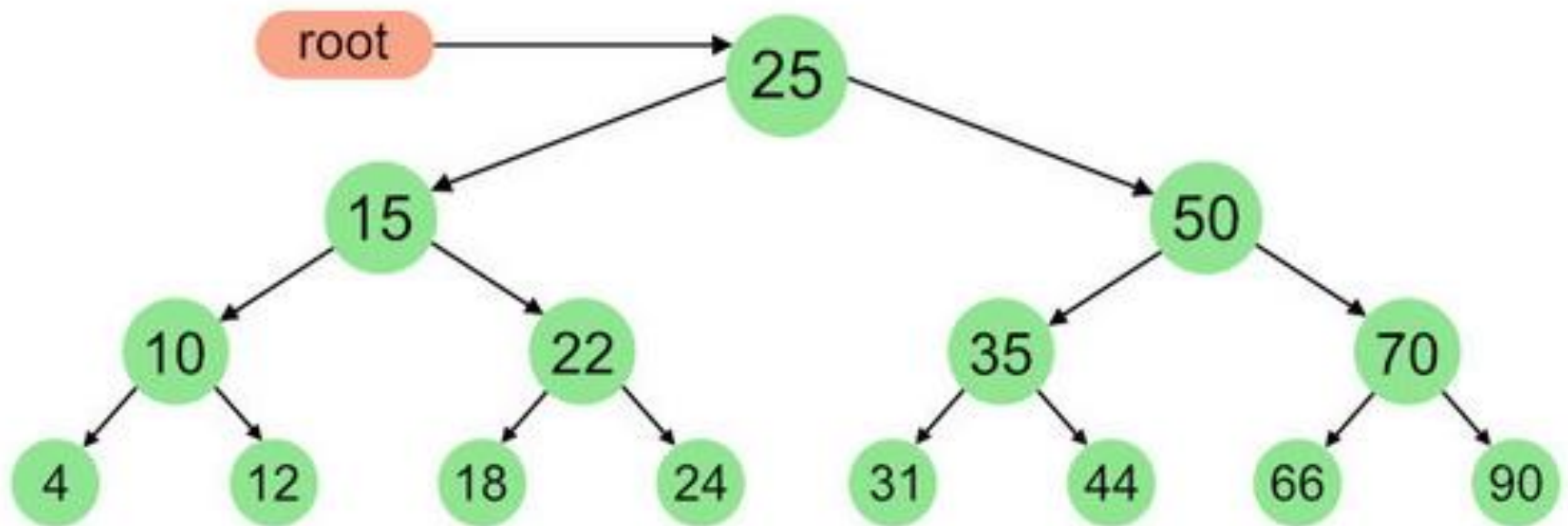




A Pre-order traversal visits nodes in the following order:  
25, 15, 10, 4, 12, 22, 18, 24, 50, 35, 31, 44, 70, 66, 90



A Post-order traversal visits nodes in the following order:  
4, 12, 10, 18, 24, 22, 15, 31, 44, 35, 66, 90, 70, 50, 25



# Traversing a Binary Tree

```
void preorder(BSTNode* BSTNode)
{
    if( BSTNode != NULL )
    {
        cout << BSTNode->getInfo() << " ";
        preorder(BSTNode->getLeft());
        preorder(BSTNode->getRight());
    }
}
```



# Traversing a Binary Tree

```
void inorder(BSTNode* BSTNode)
{
    if( BSTNode != NULL )
    {
        inorder(BSTNode->getLeft());
        cout << BSTNode->getInfo()<<" ";
        inorder(BSTNode->getRight());
    }
}
```

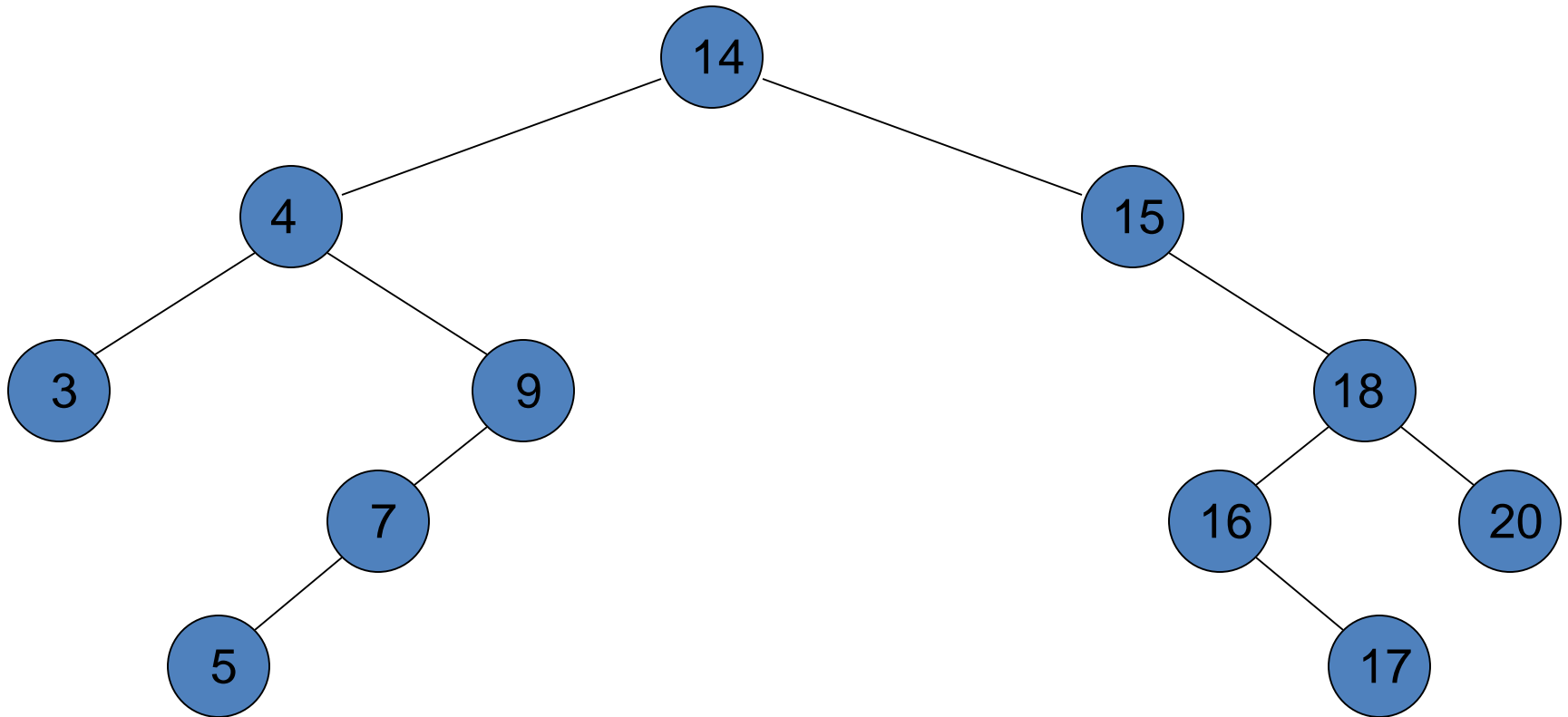
# Traversing a Binary Tree

```
void postorder(BSTNode* BSTNode)
{
    if( BSTNode != NULL )
    {
        postorder(BSTNode->getLeft());
        postorder(BSTNode->getRight());
        cout << BSTNode->getInfo() << " ";
    }
}
```

# Traversing a Binary Tree

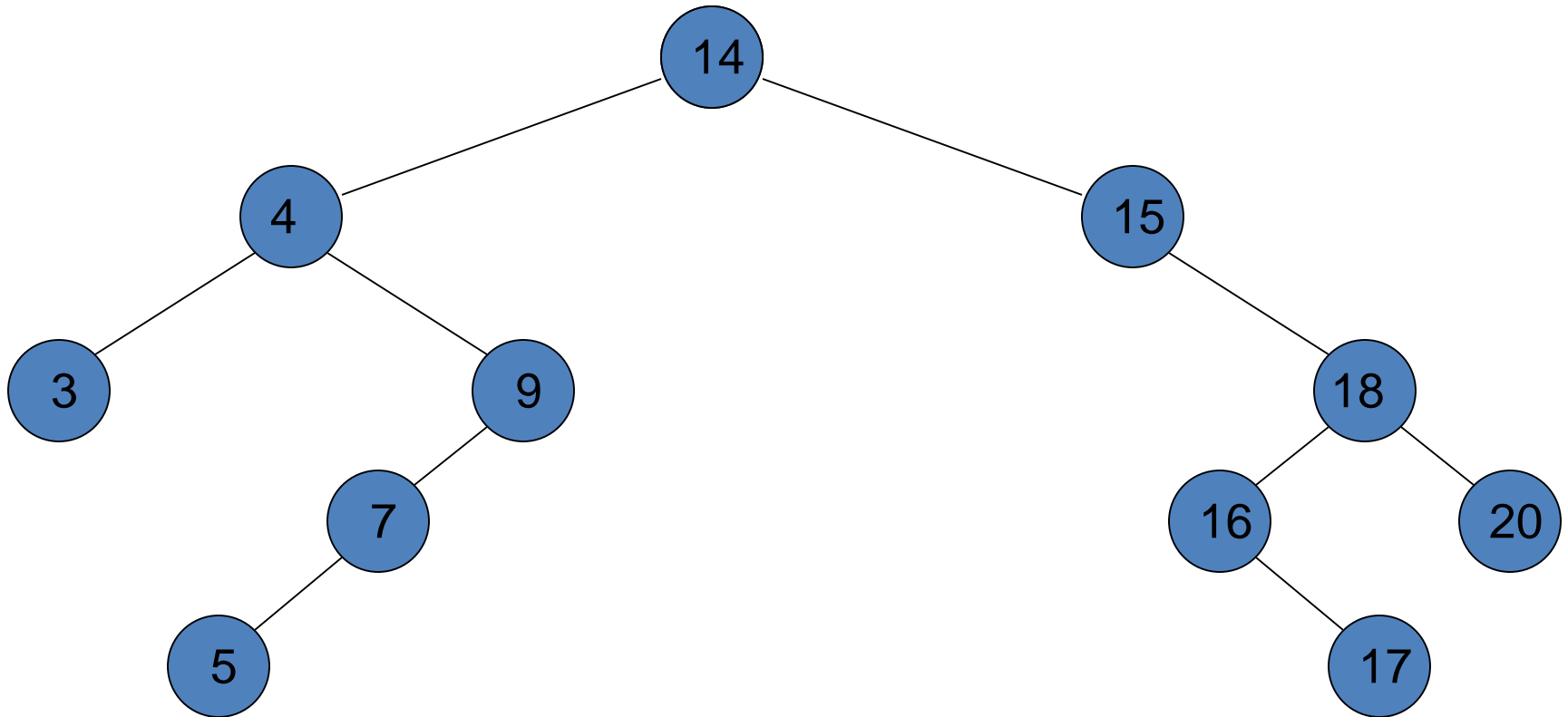
```
cout << "inorder: ";    preorder( root );  
cout << "inorder: ";    inorder( root );  
cout << "postorder: ";  postorder( root );
```

# Traversing a Binary Tree



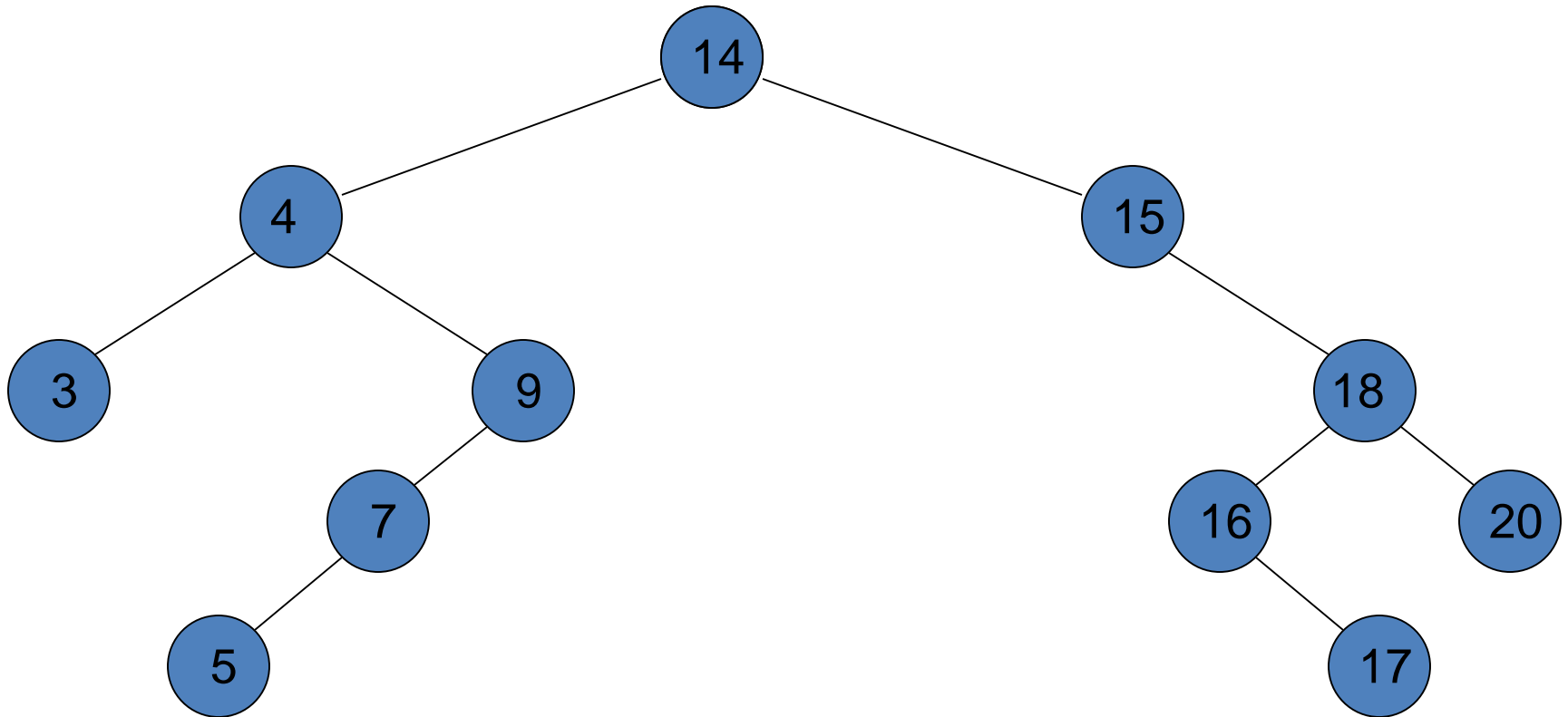
Preorder: 14 4 3 9 7 5 15 18 16 17 20

# Traversing a Binary Tree



Inorder: 3 4 5 7 9 14 15 16 17 18 20

# Traversing a Binary Tree



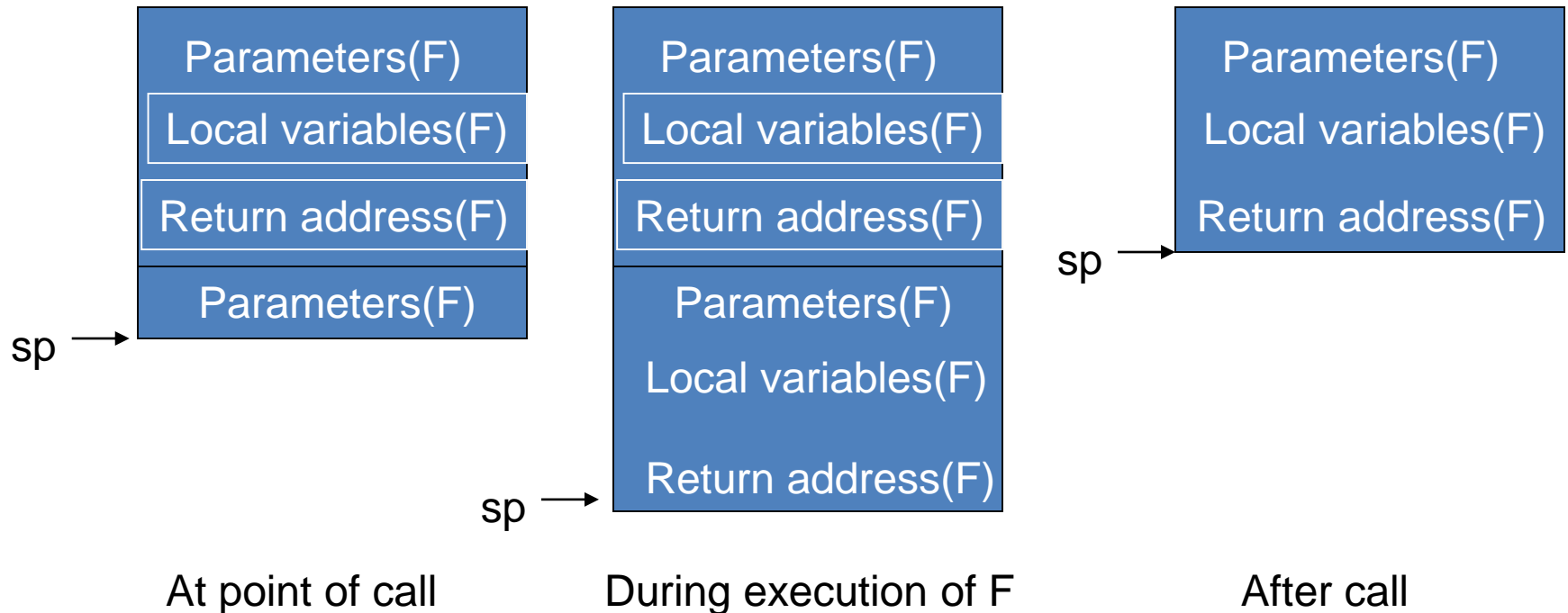
Postorder: 3 5 7 9 4 17 16 20 18 15 14

# Recursive Call

- Recall that a stack is used during function calls.
- The caller function places the arguments on the stack and passes control to the called function.
- Local variables are allocated storage on the call stack.
- Calling a function itself makes no difference as far as the call stack is concerned.

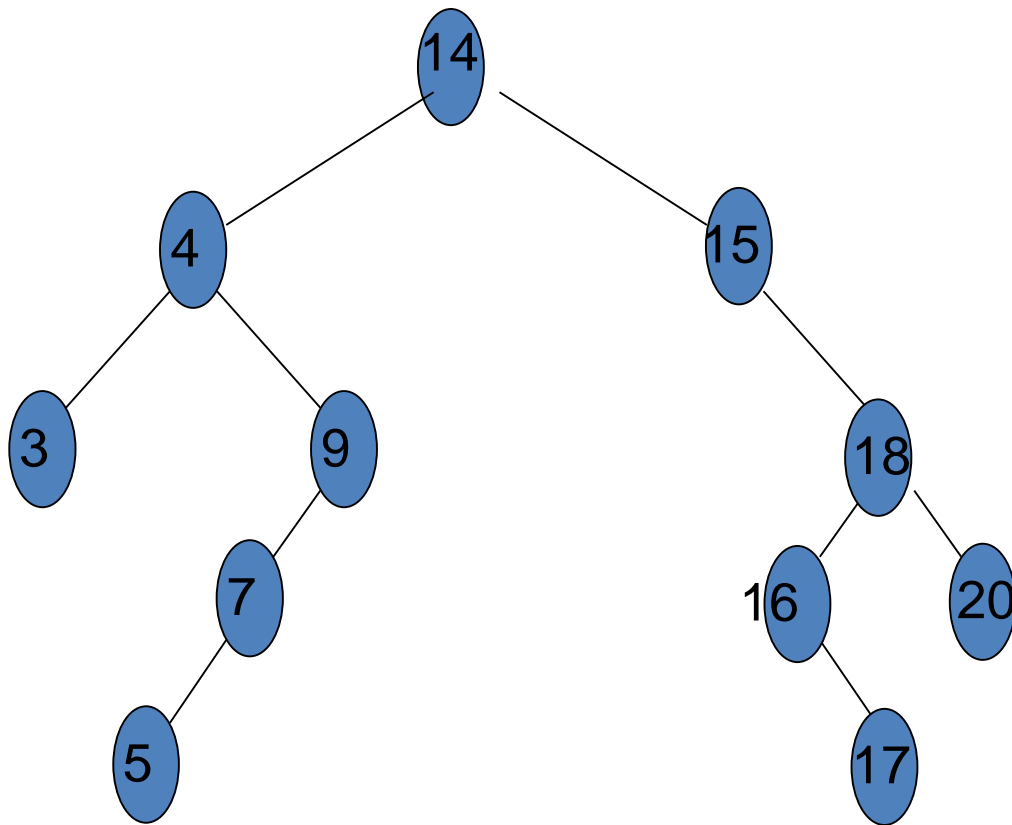
# Stack Layout during a call

- Here is stack layout when function F calls function F (recursively):



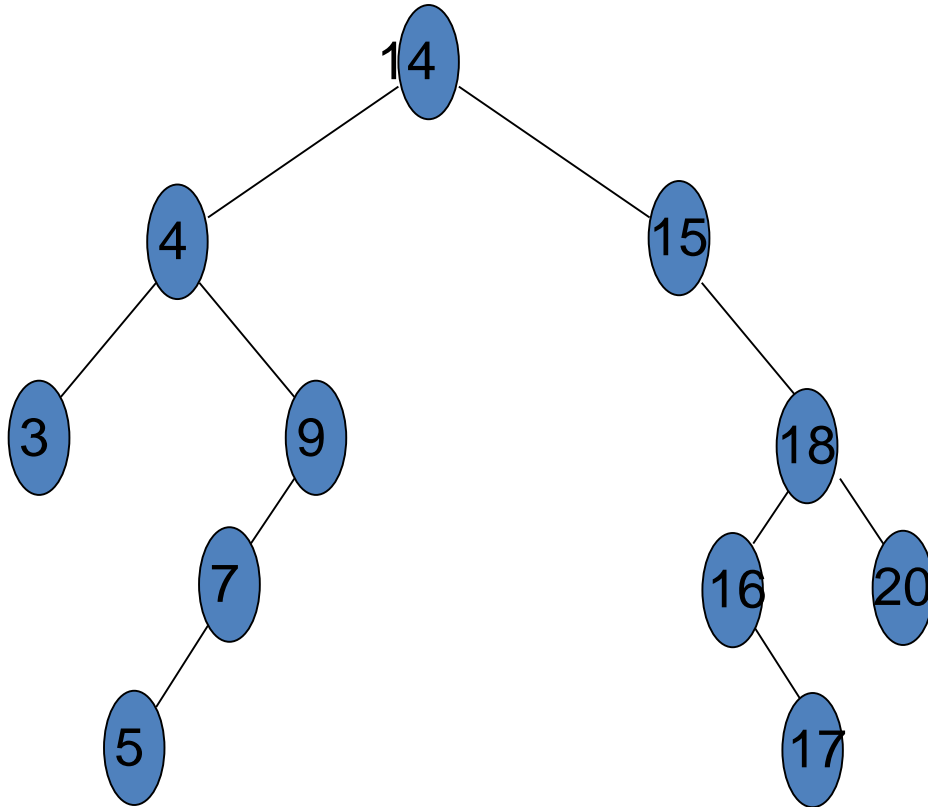


# Recursion: preorder



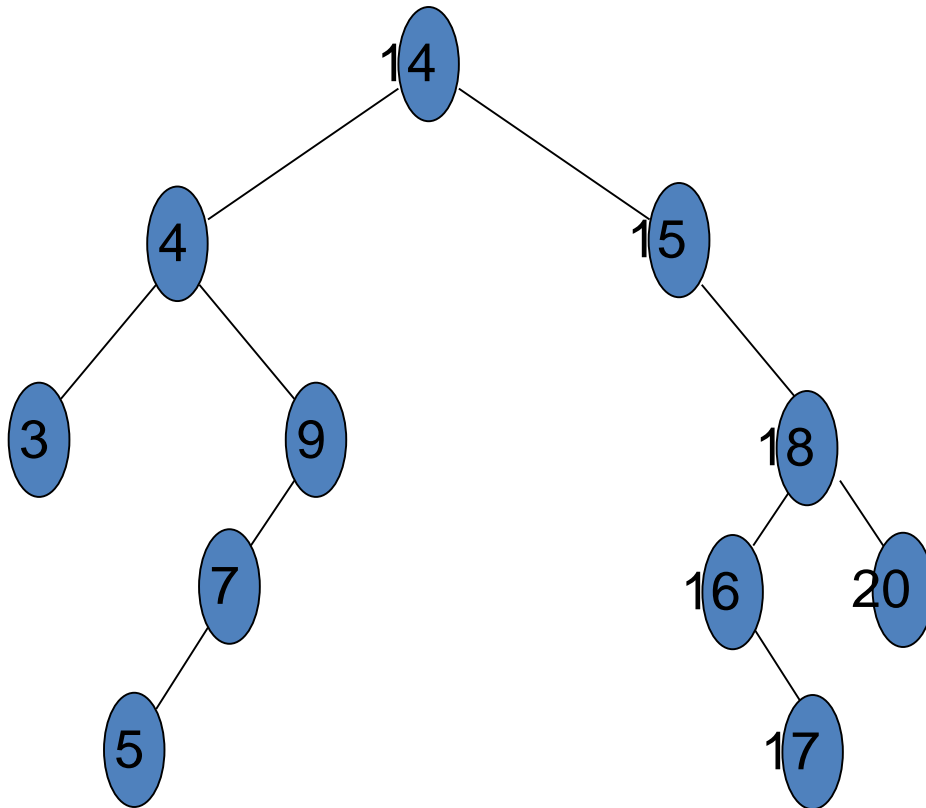
```
preorder(14)
14
  ..preorder(4)
  4
    ....preorder(3)
    3
      .....preorder(null)
      .....preorder(null)
      ....preorder(9)
      9
        .....preorder(7)
        7
          .....preorder(5)
          5
            .....preorder(null)
            .....preorder(null)
            .....preorder(null)
            .....preorder(null)
```

# Recursion: preorder



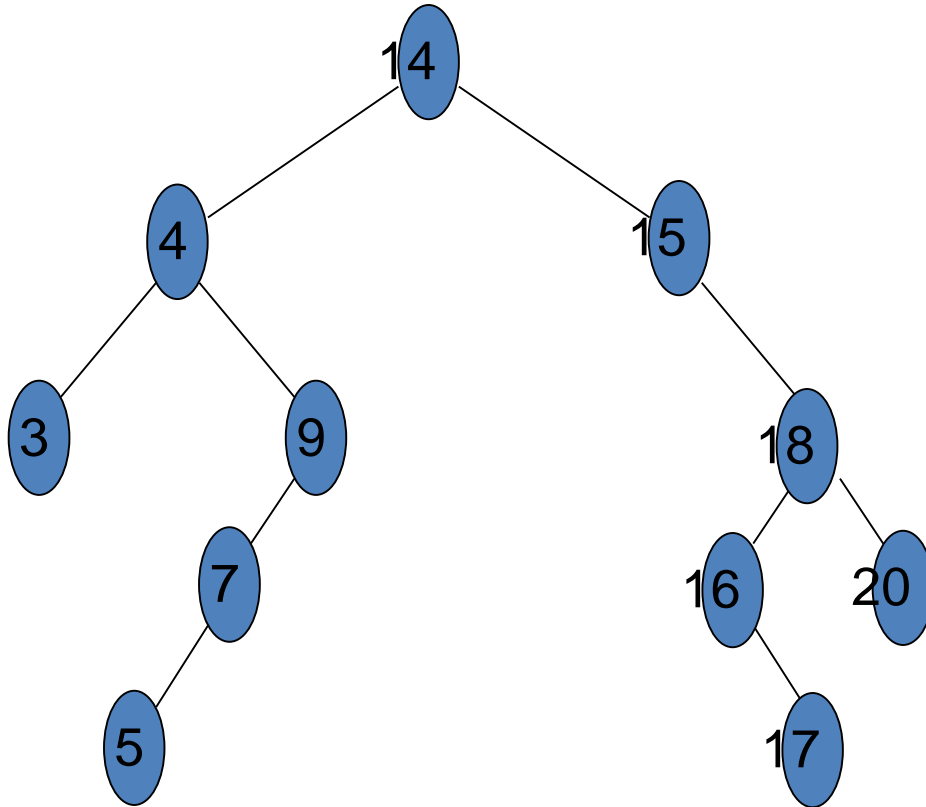
```
..preorder(15)
15
....preorder(null)
....preorder(18)
18
.....preorder(16)
16
.....preorder(null)
.....preorder(17)
17
.....preorder(null)
.....preorder(null)
.....preorder(20)
20
.....preorder(null)
.....preorder(null)
```

# Recursion: inorder



```
inorder(14)
..inorder(4)
....inorder(3)
.....inorder(null)
3
.....inorder(null)
4
....inorder(9)
.....inorder(7)
.....inorder(5)
.....inorder(null)
5
.....inorder(null)
7
.....inorder(null)
9
.....inorder(null)
14
```

# Recursion: inorder

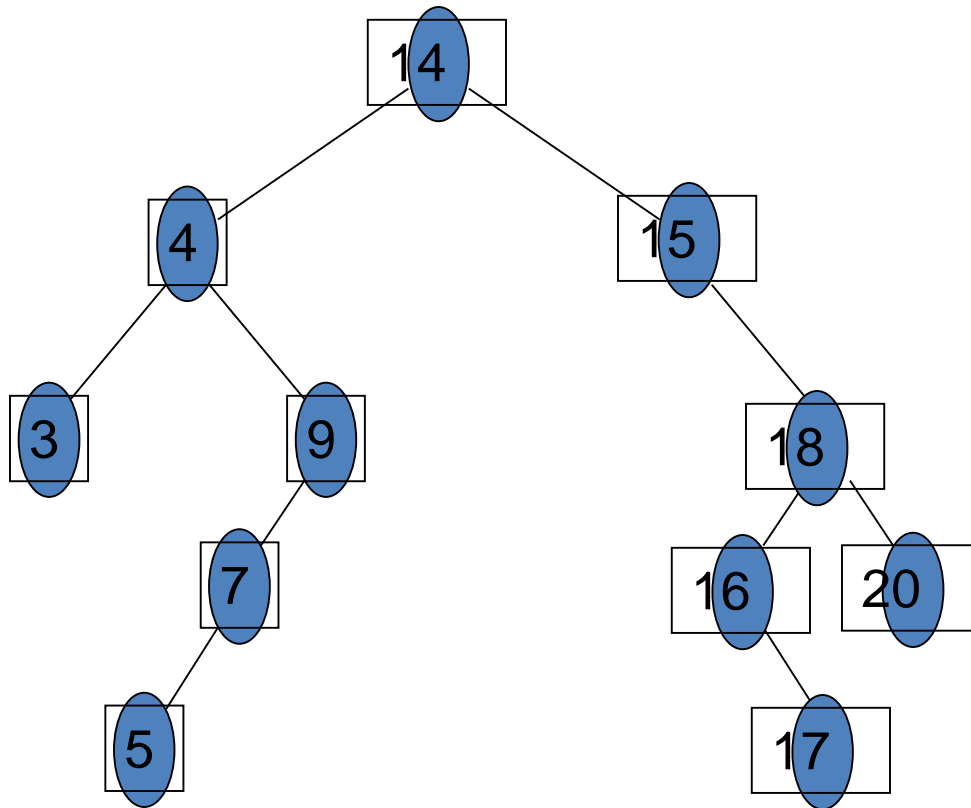


```
..inorder(15)
....inorder(null)
15
....inorder(18)
.....inorder(16)
.....inorder(null)
16
.....inorder(17)
.....inorder(null)
17
.....inorder(null)
18
.....inorder(20)
.....inorder(null)
20
.....inorder(null)
```

# Non Recursive Traversal

- We can implement non-recursive versions of the preorder, inorder and postorder traversal by using an explicit stack.
- The stack will be used to store the tree nodes in the appropriate order.
- Here, for example, is the routine for inorder traversal that uses a stack.

# Nonrecursive Inorder

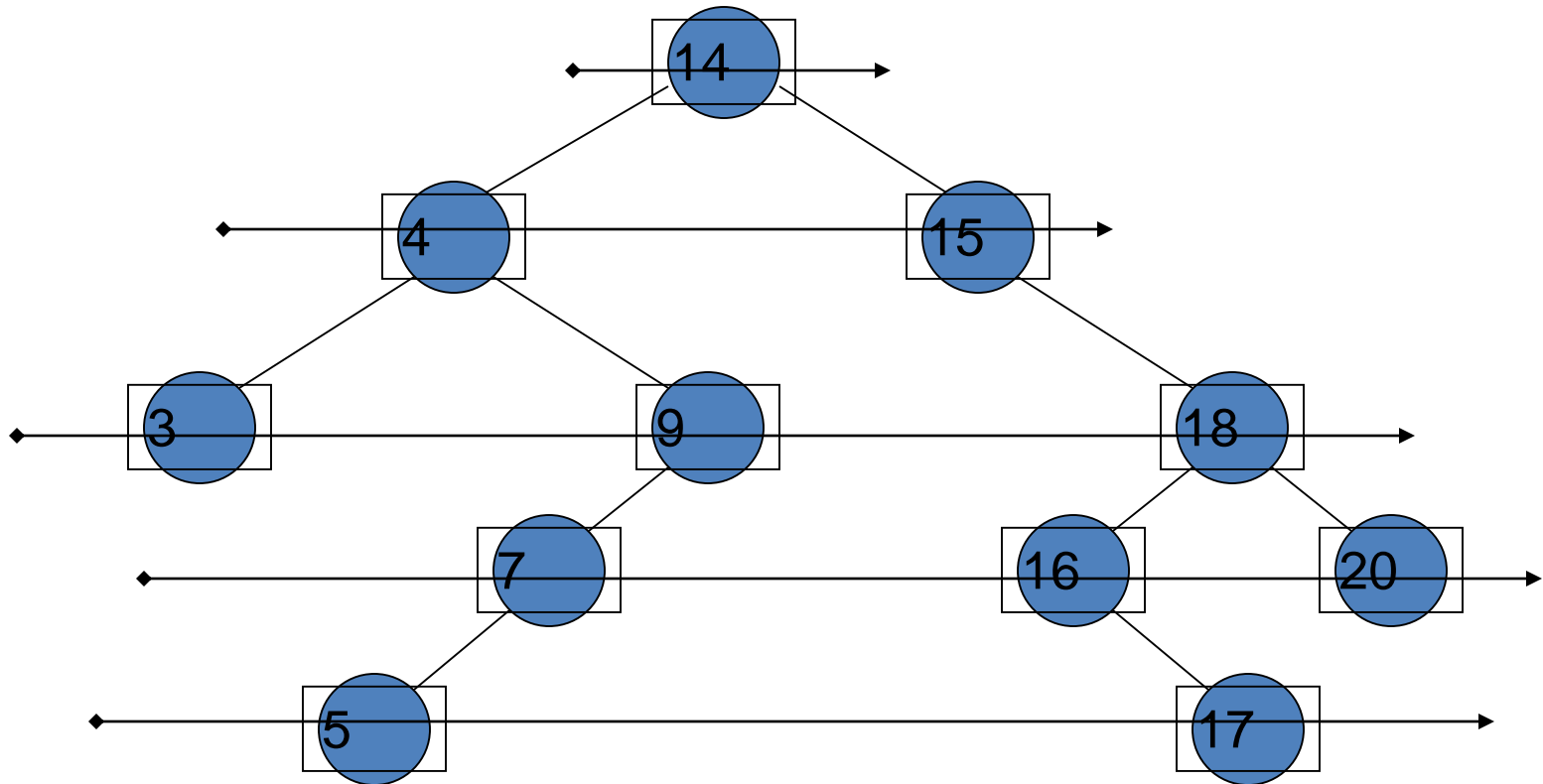


```
push(14)
..push(4)
....push(3)
3
4
..push(9)
....push(7)
.....push(5)
5
7
9
14
push(15)
15
push(18)
..push(16)
16
..push(17)
17
18
push(20)
20
```

# Level-order Traversal

- There is yet another way of traversing a binary tree that is not related to recursive traversal procedures discussed previously.
- In level-order traversal, we visit the nodes at each level before proceeding to the next level.
- At each level, we visit the nodes in a left-to-right order.

# Level-order Traversal



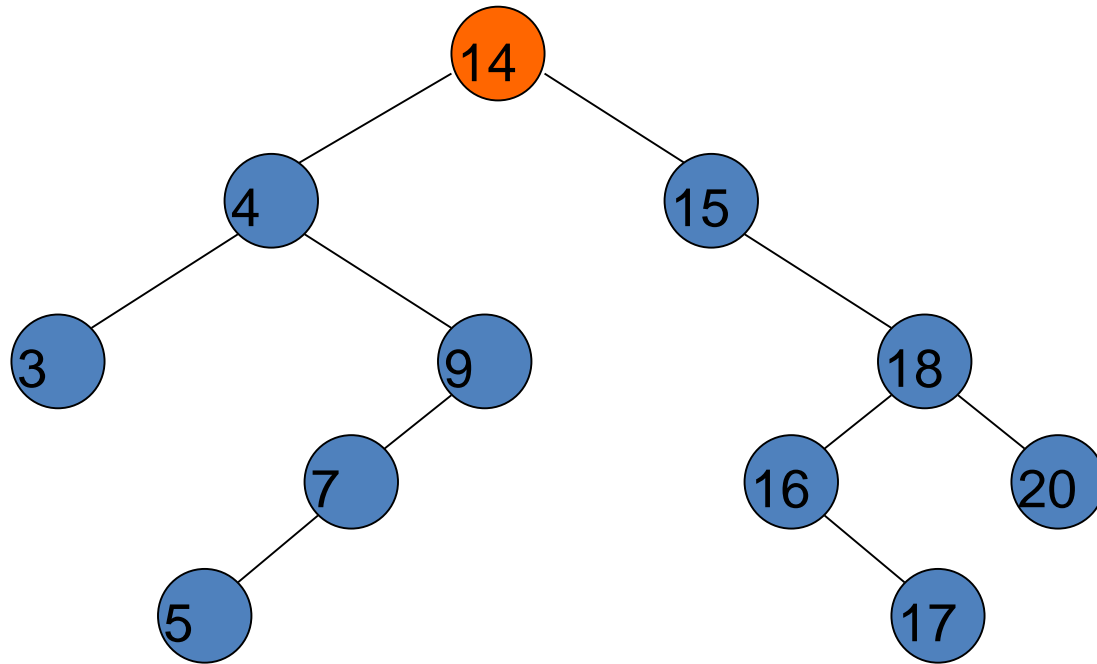
Level-order: 14 4 15 3 9 18 7 16 20 5 17



# Level-order Traversal

- How do we do level-order traversal?
- Surprisingly, if we use a queue instead of a stack, we can visit the nodes in level-order.
- Here is the code for level-order traversal:

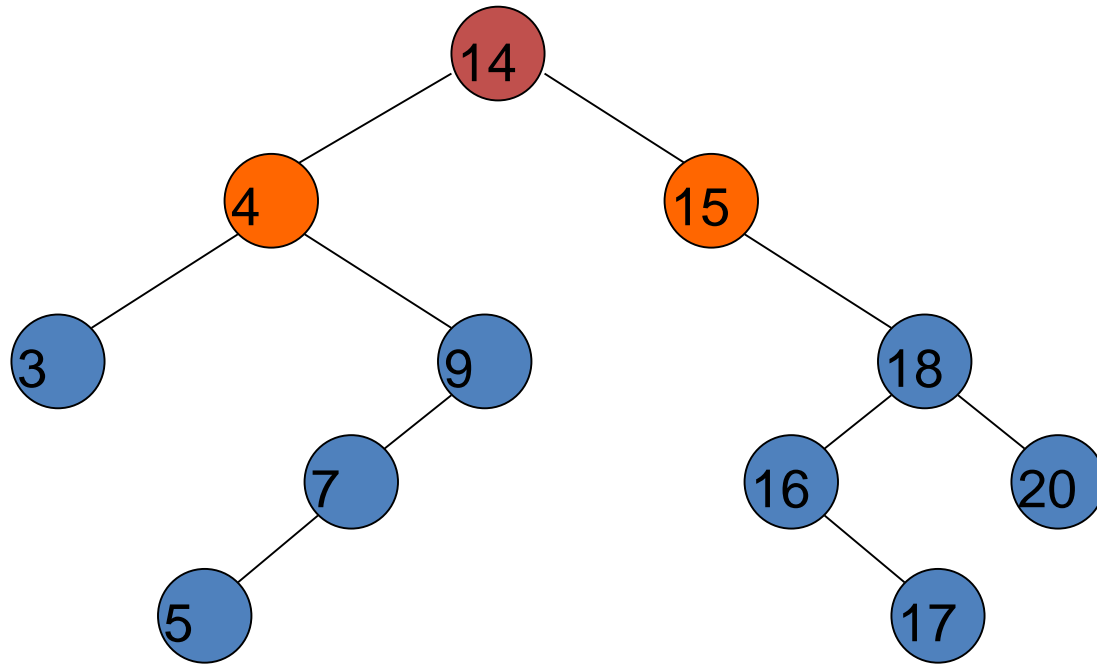
# Level-order Traversal



Queue: 14

Output:

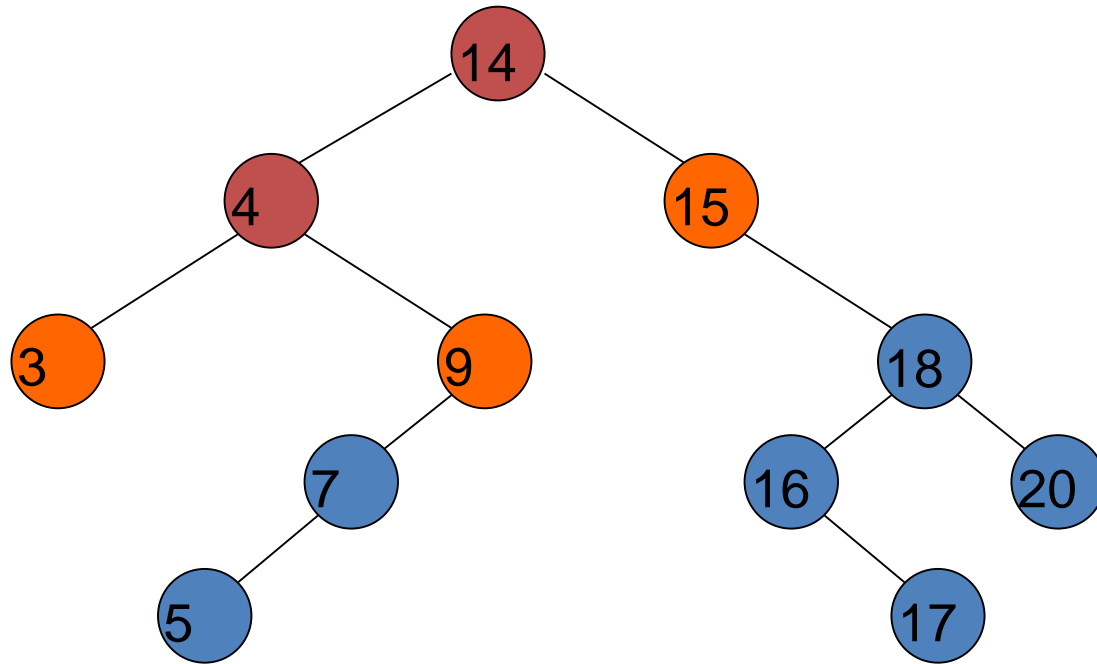
# Level-order Traversal



Queue: 4 15

Output: 14

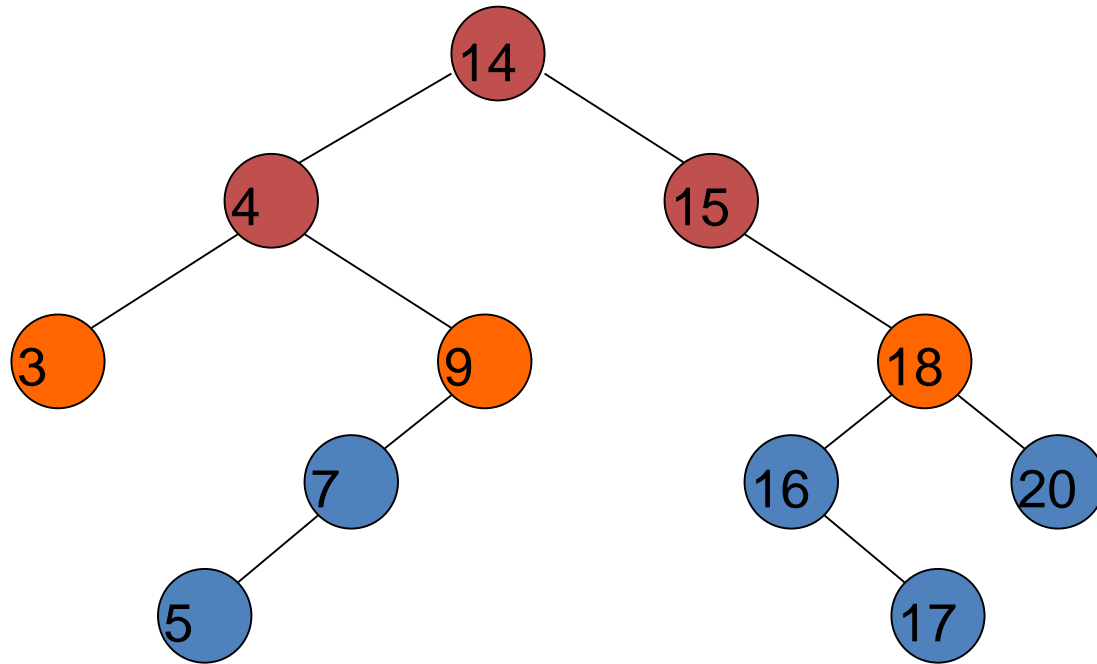
# Level-order Traversal



Queue: 15 3 9

Output: 14 4

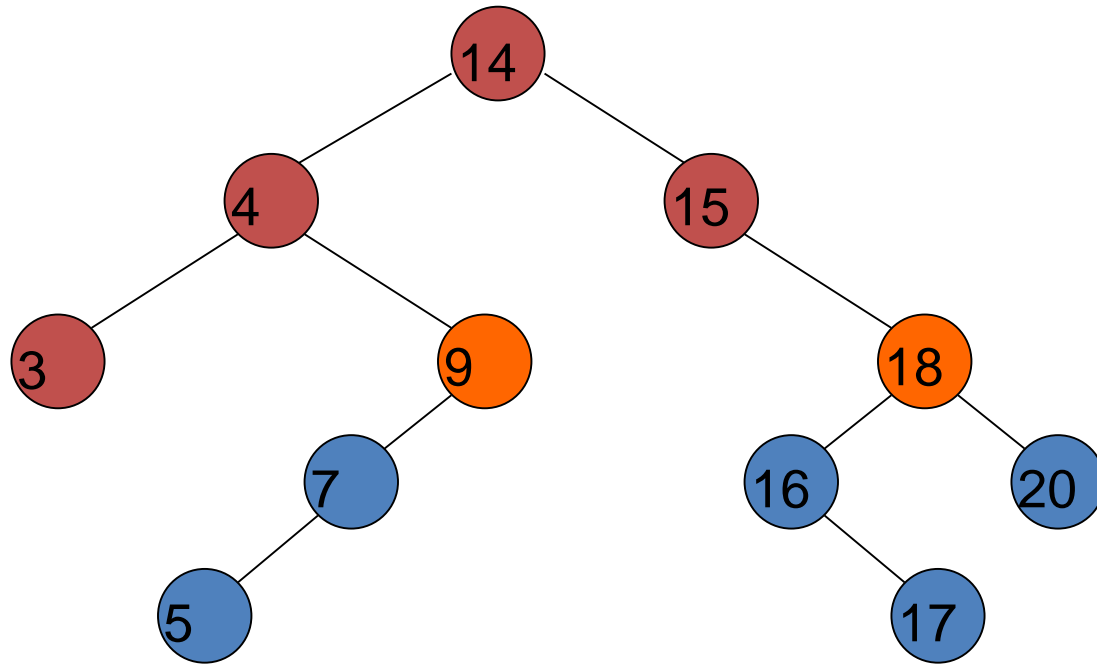
# Level-order Traversal



Queue: 3 9 18

Output: 14 4 15

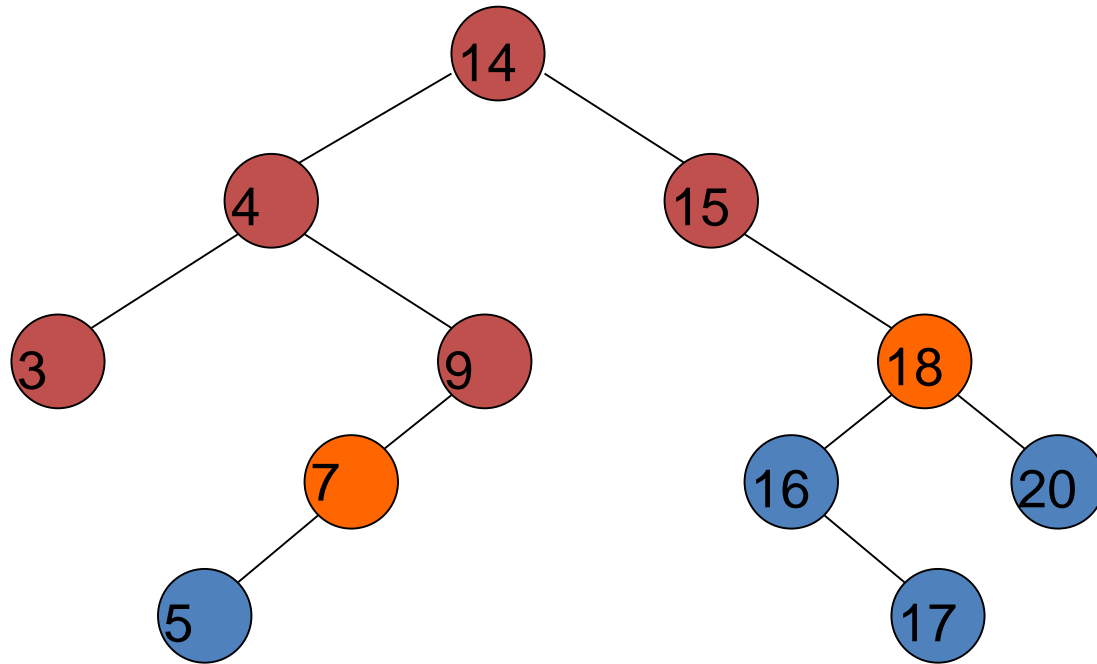
# Level-order Traversal



Queue: 9 18

Output: 14 4 15 3

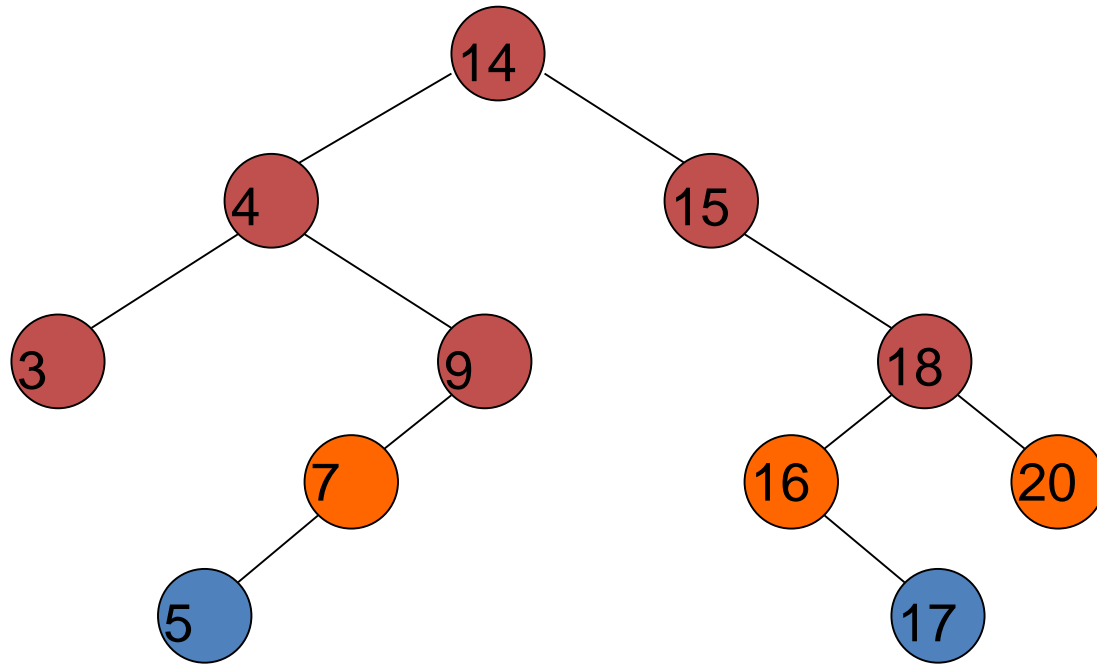
# Level-order Traversal



Queue: 18 7

Output: 14 4 15 3 9

# Level-order Traversal

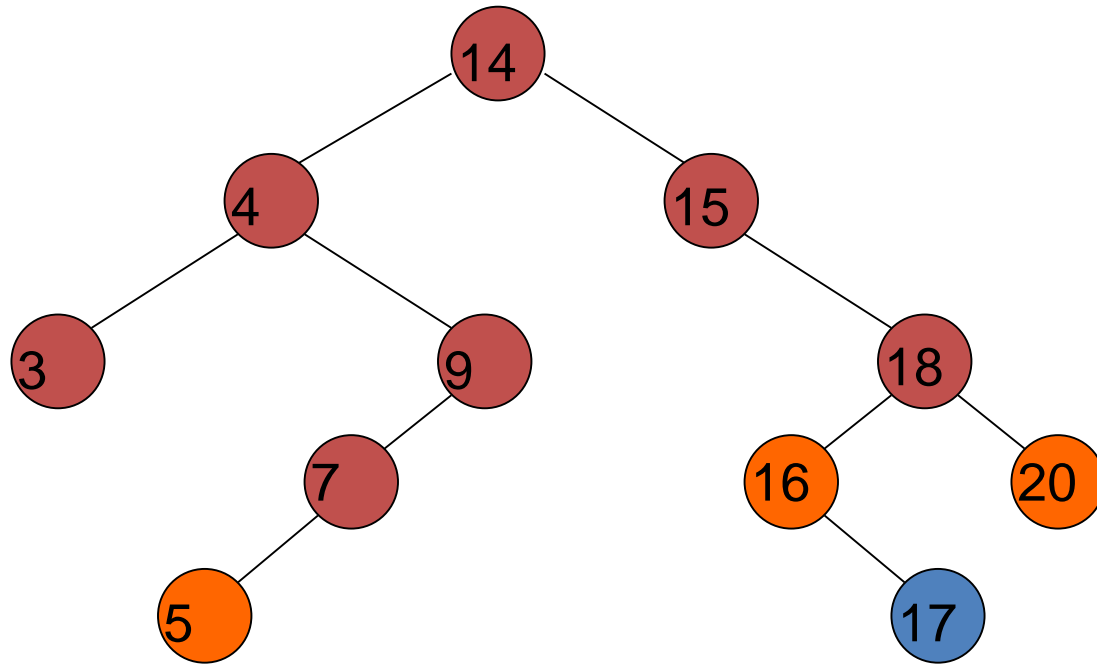


Queue: 7 16 20

Output: 14 4 15 3 9 18



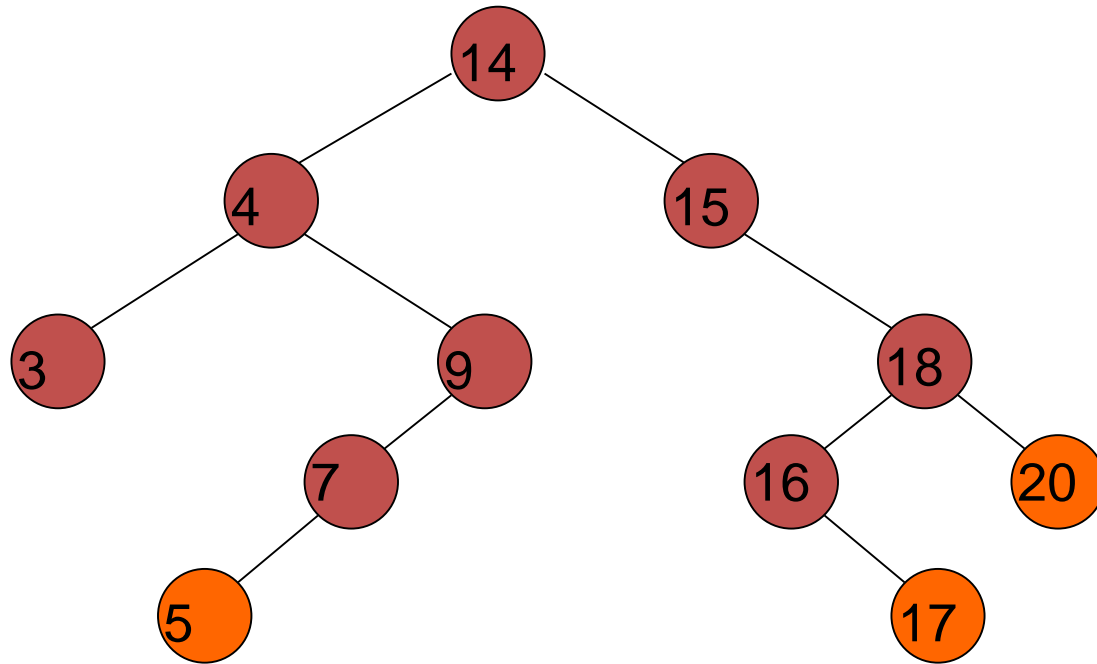
# Level-order Traversal



Queue: 16 20 5

Output: 14 4 15 3 9 18 7

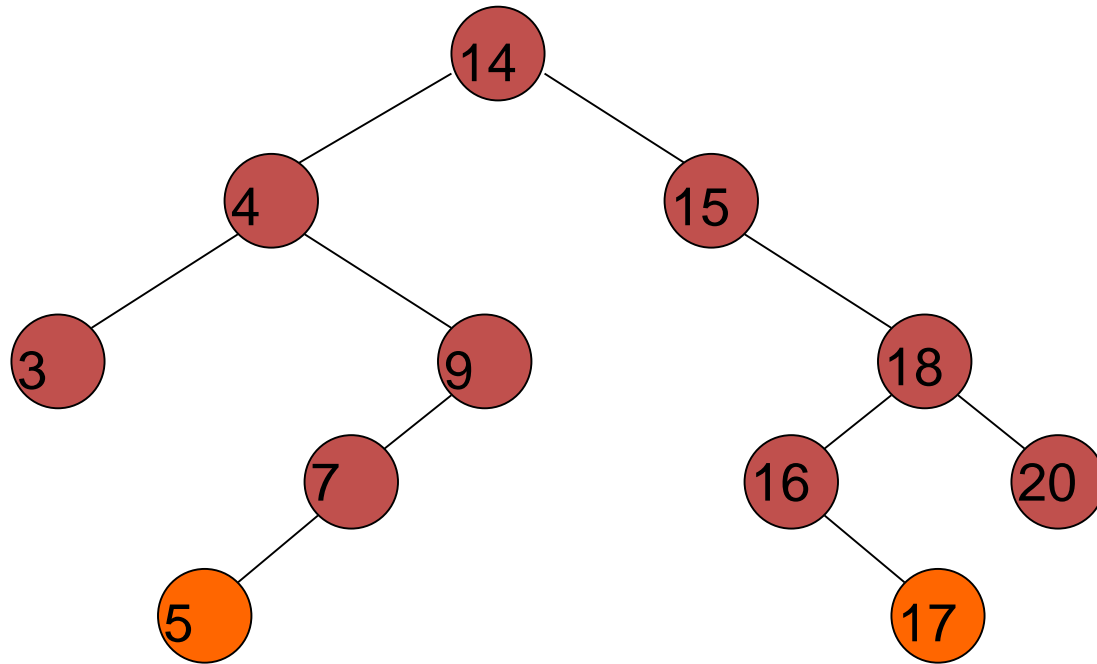
# Level-order Traversal



Queue: 20 5 17

Output: 14 4 15 3 9 18 7 16

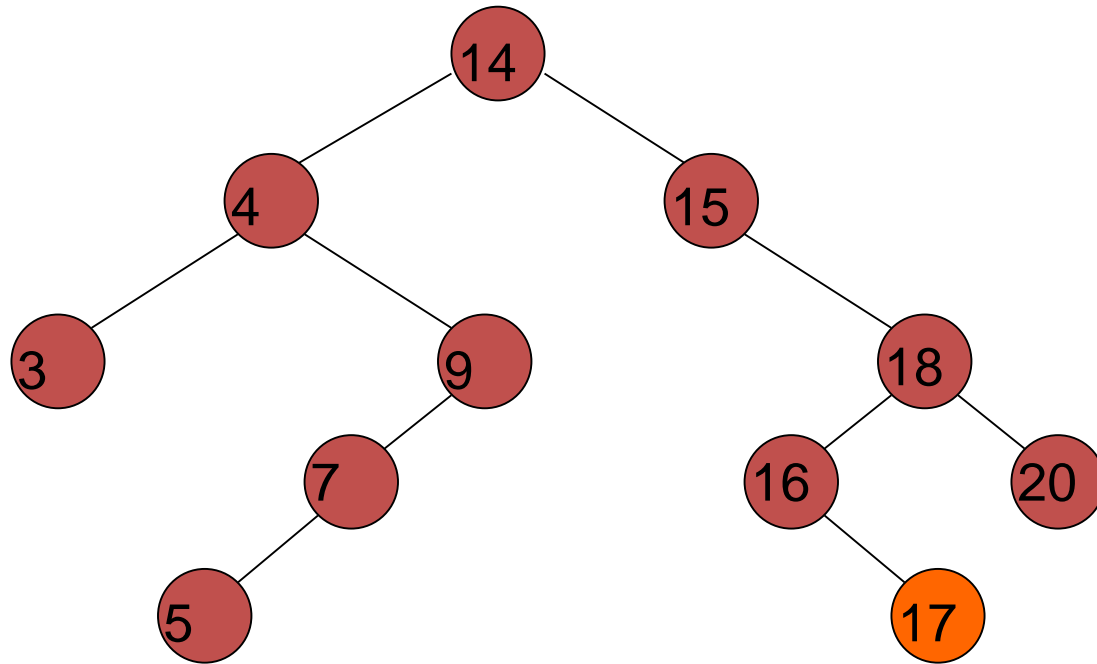
# Level-order Traversal



Queue: 5 17

Output: 14 4 15 3 9 18 7 16 20

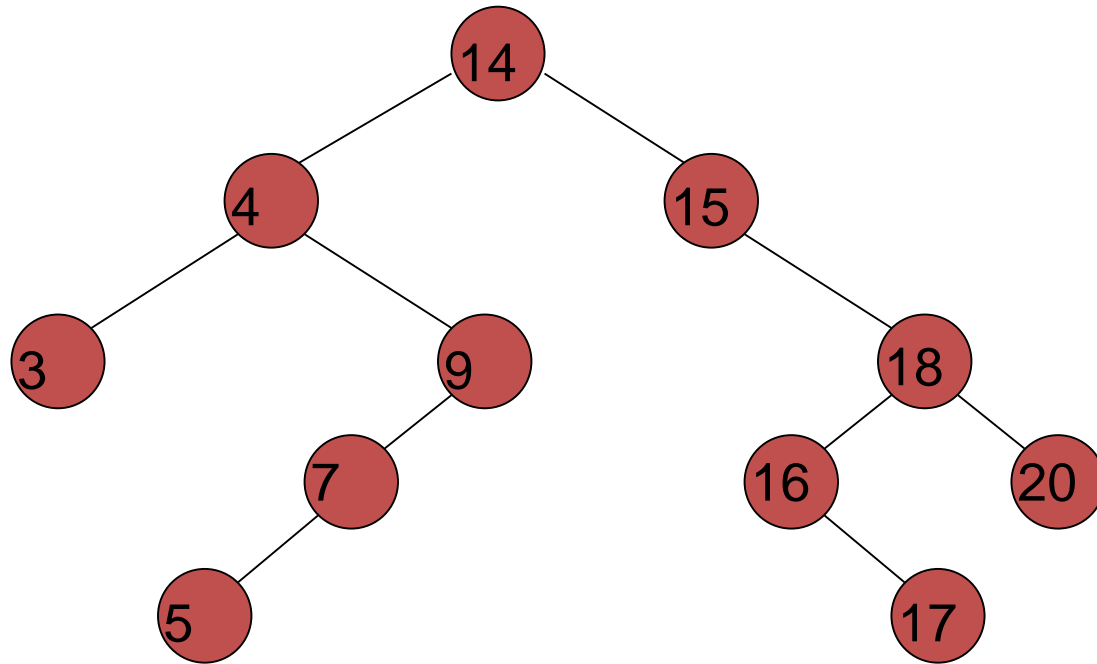
# Level-order Traversal



Queue: 17

Output: 14 4 15 3 9 18 7 16 20 5

# Level-order Traversal



Queue:

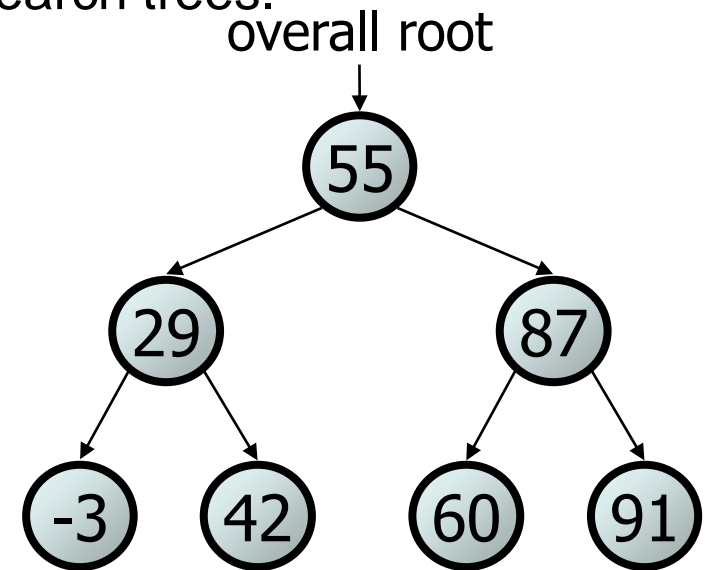
Output: 14 4 15 3 9 18 7 16 20 5 17

# Binary search trees

- **binary search tree** ("BST"): a binary tree where each non-empty node R has the following properties:

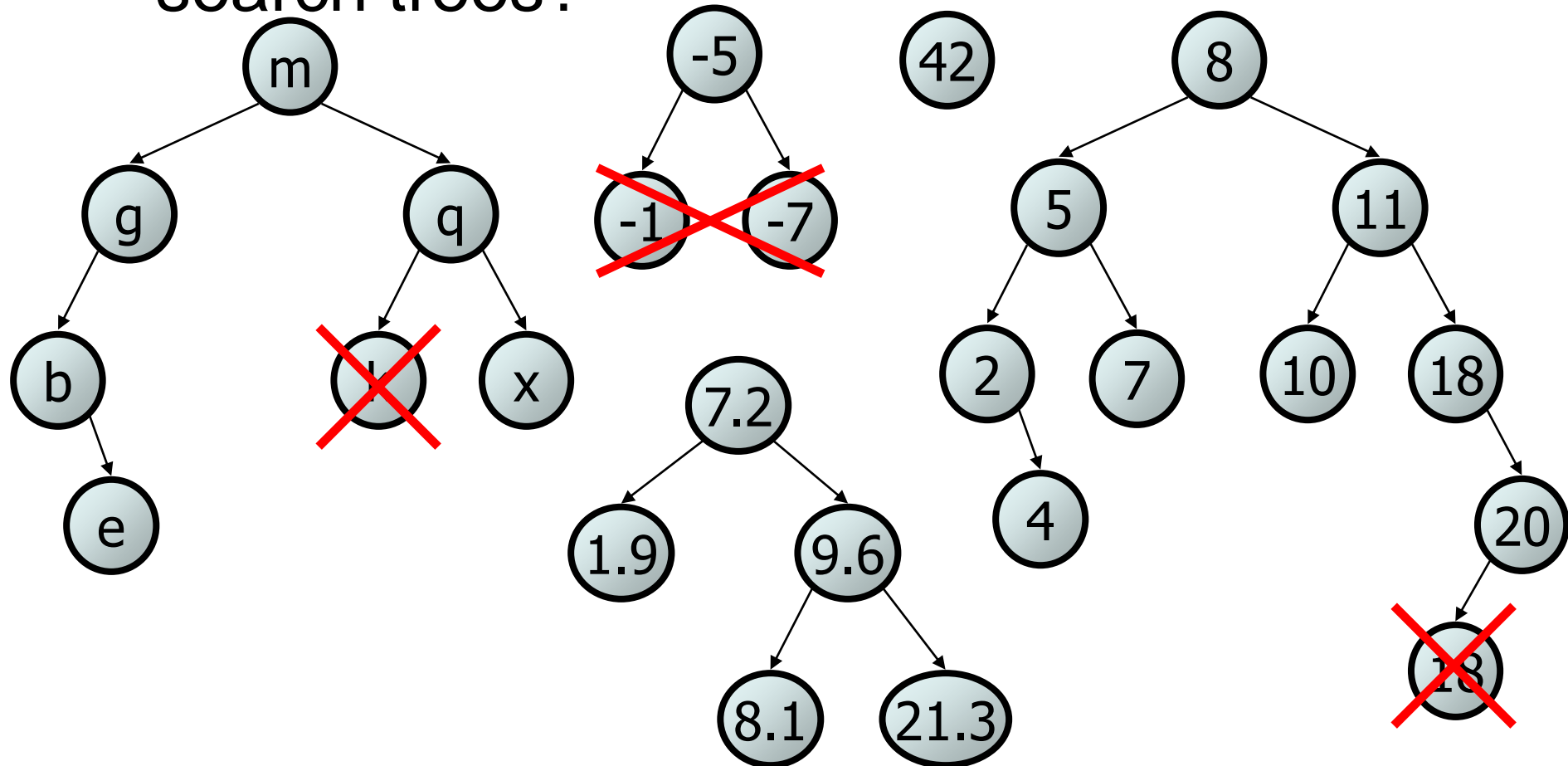
- every element of R's left subtree contains data "less than" R's data,
- every element of R's right subtree contains data "greater than" R's,
- R's left and right subtrees are also binary search trees.

- BSTs store their elements in sorted order, which is helpful for searching/sorting tasks.

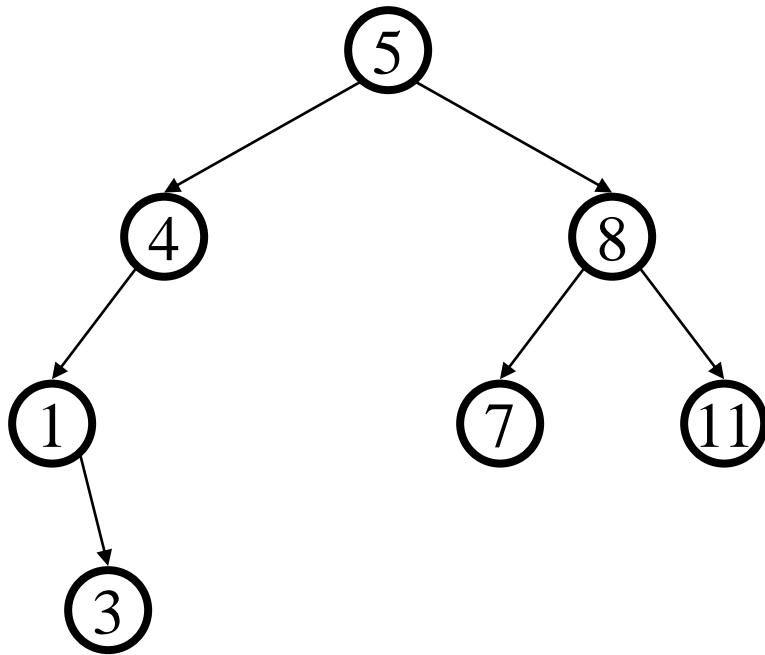


# BST examples

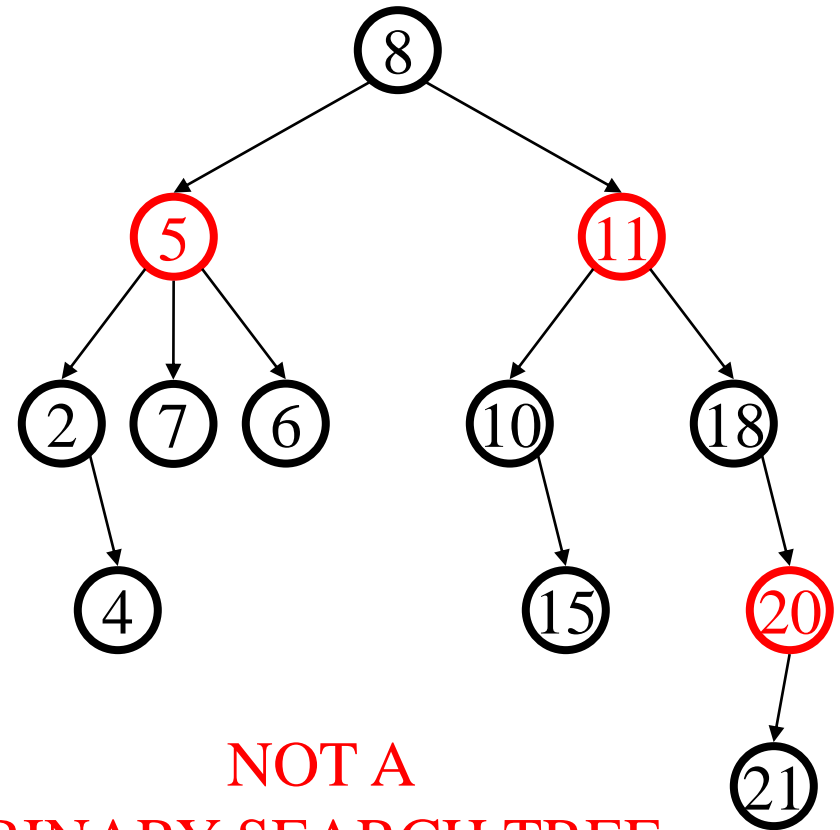
- Which of the trees shown are legal binary search trees?



# Example and Counter-Example



BINARY SEARCH TREE

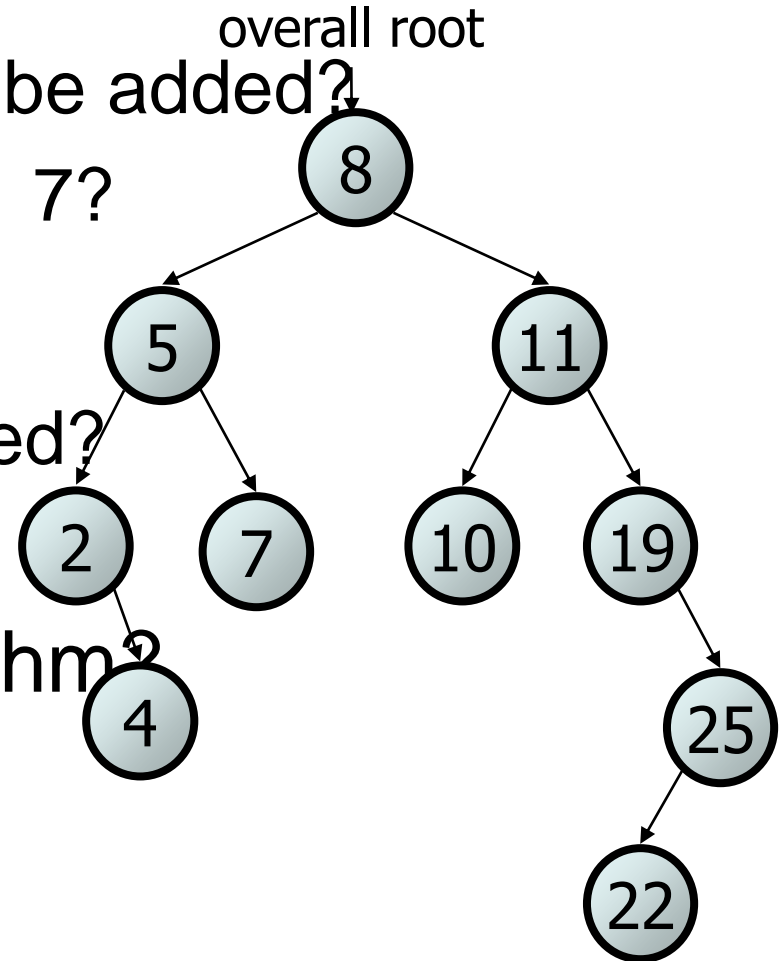


NOT A  
BINARY SEARCH TREE



# Adding to a BST

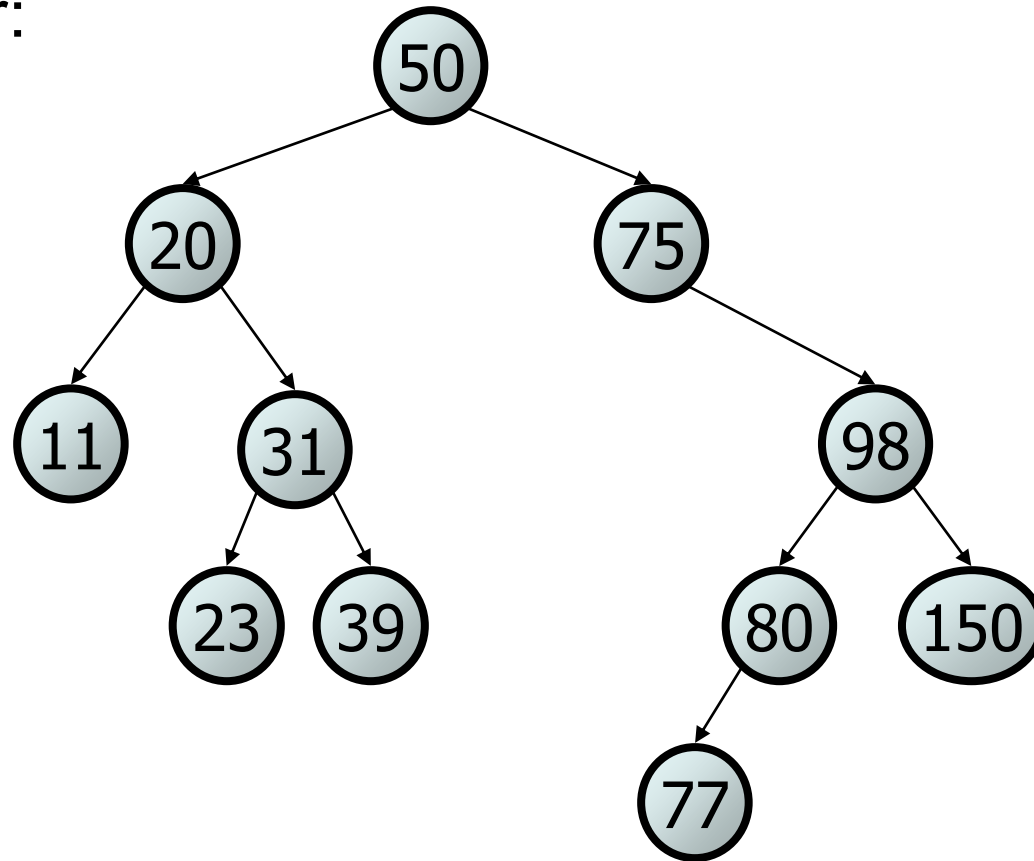
- Suppose we want to add new values to the BST below.
  - Where should the value 14 be added?
  - Where should 3 be added? 7?
  - If the tree is empty, where should a new value be added?
- What is the general algorithm?



# Adding exercise

- Draw what a binary search tree would look like if the following values were added to an initially empty tree in this order:

50  
20  
75  
98  
80  
31  
150  
39  
23  
11  
77

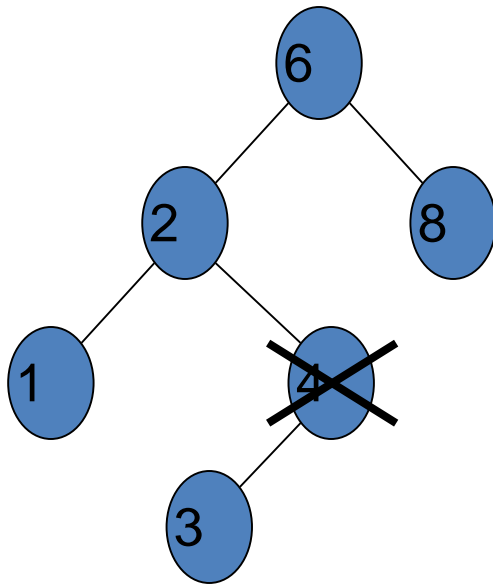


# Deleting a node in BST

- As is common with many data structures, the hardest operation is deletion.
- Once we have found the node to be deleted, we need to consider several possibilities.
- If the node is a *leaf*, it can be deleted immediately.

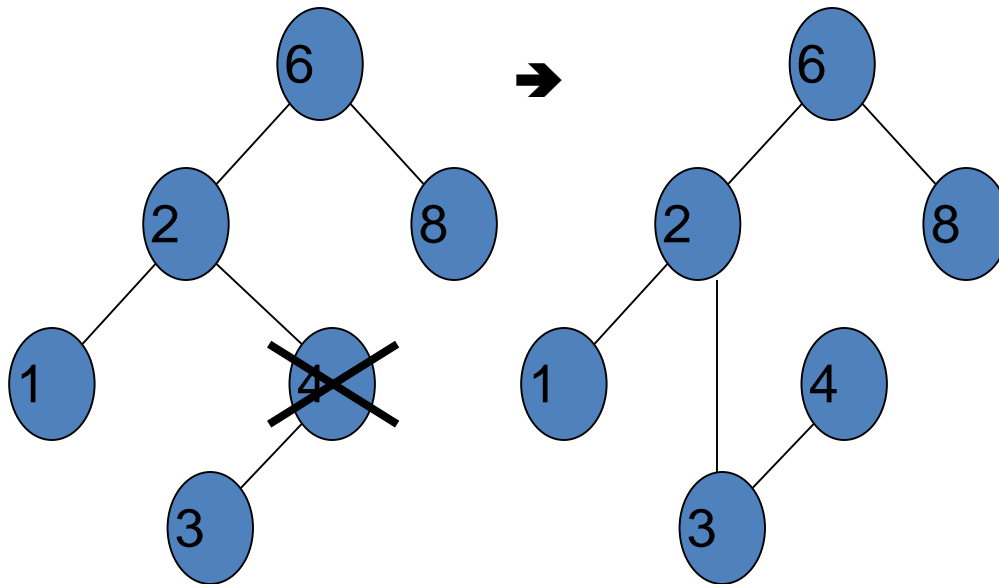
# Deleting a node in BST

- If the node has one child, the node can be deleted after its parent adjusts a pointer to bypass the node and connect to inorder successor.



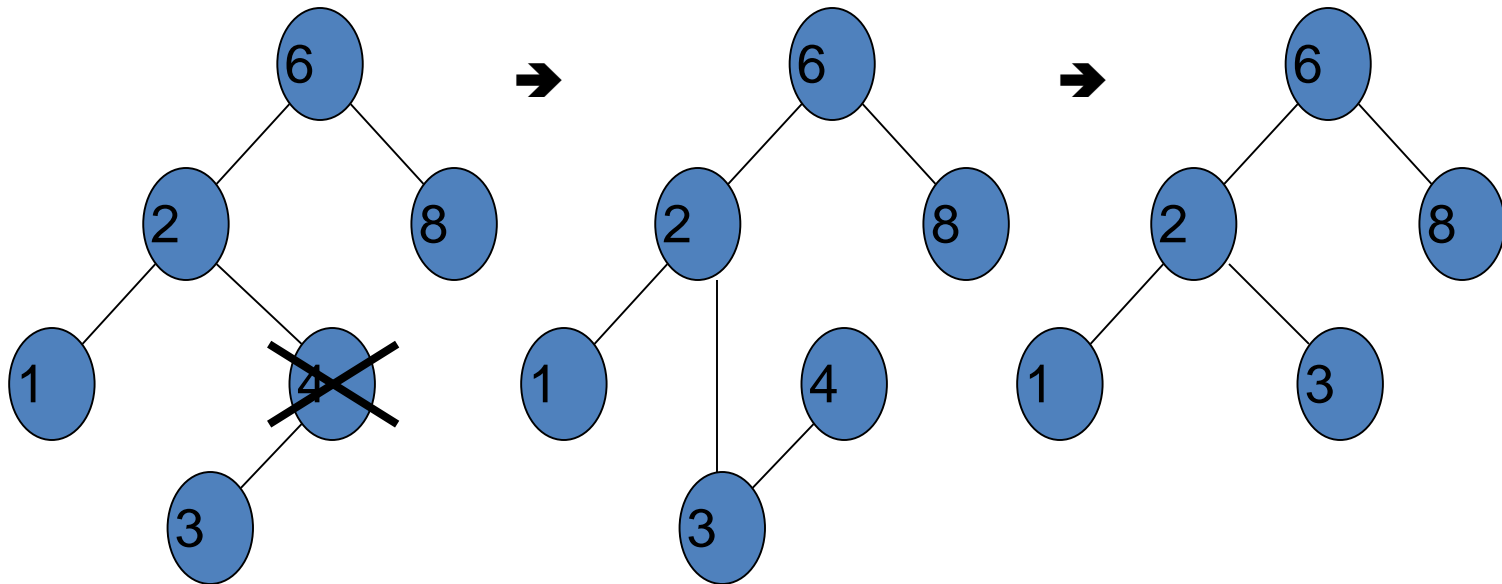
# Deleting a node in BST

- The inorder traversal order has to be maintained after the delete.



# Deleting a node in BST

- The inorder traversal order has to be maintained after the delete.

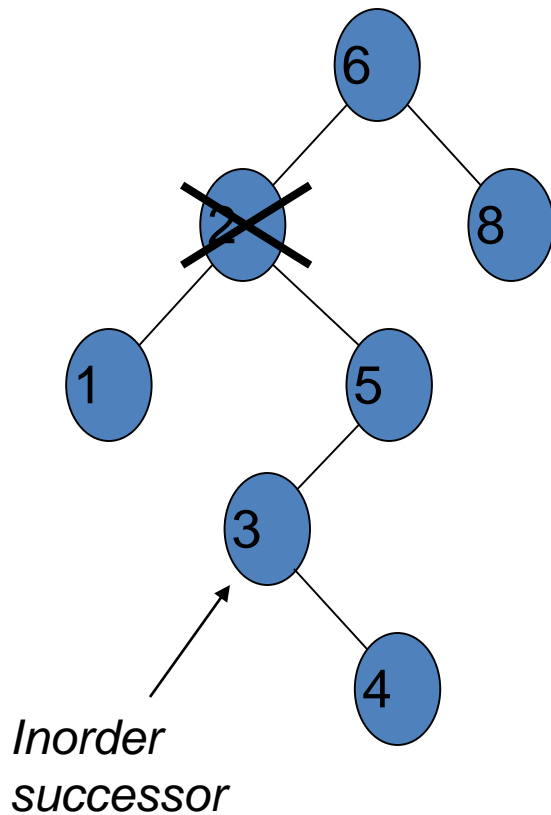


# Deleting a node in BST

- The complicated case is when the node to be deleted has both left and right subtrees.
- The strategy is to replace the data of this node with the smallest data of the right subtree and recursively delete that node.

# Deleting a node in BST

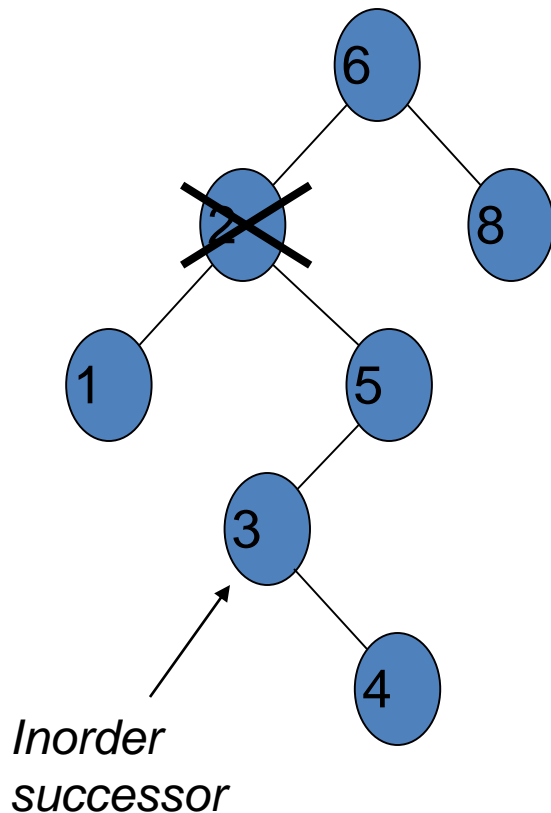
Delete(2): locate inorder successor





# Deleting a node in BST

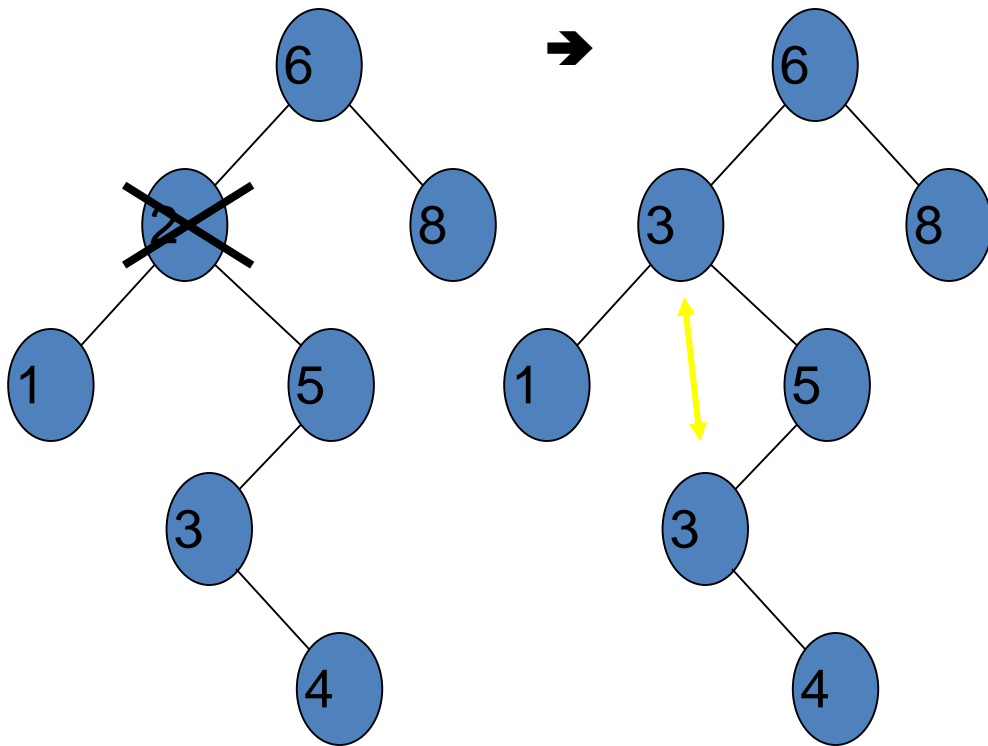
Delete(2): locate inorder successor



- Inorder successor will be the left-most node in the right subtree of 2.
- The inorder successor will not have a left child because if it did, that child would be the left-most node.

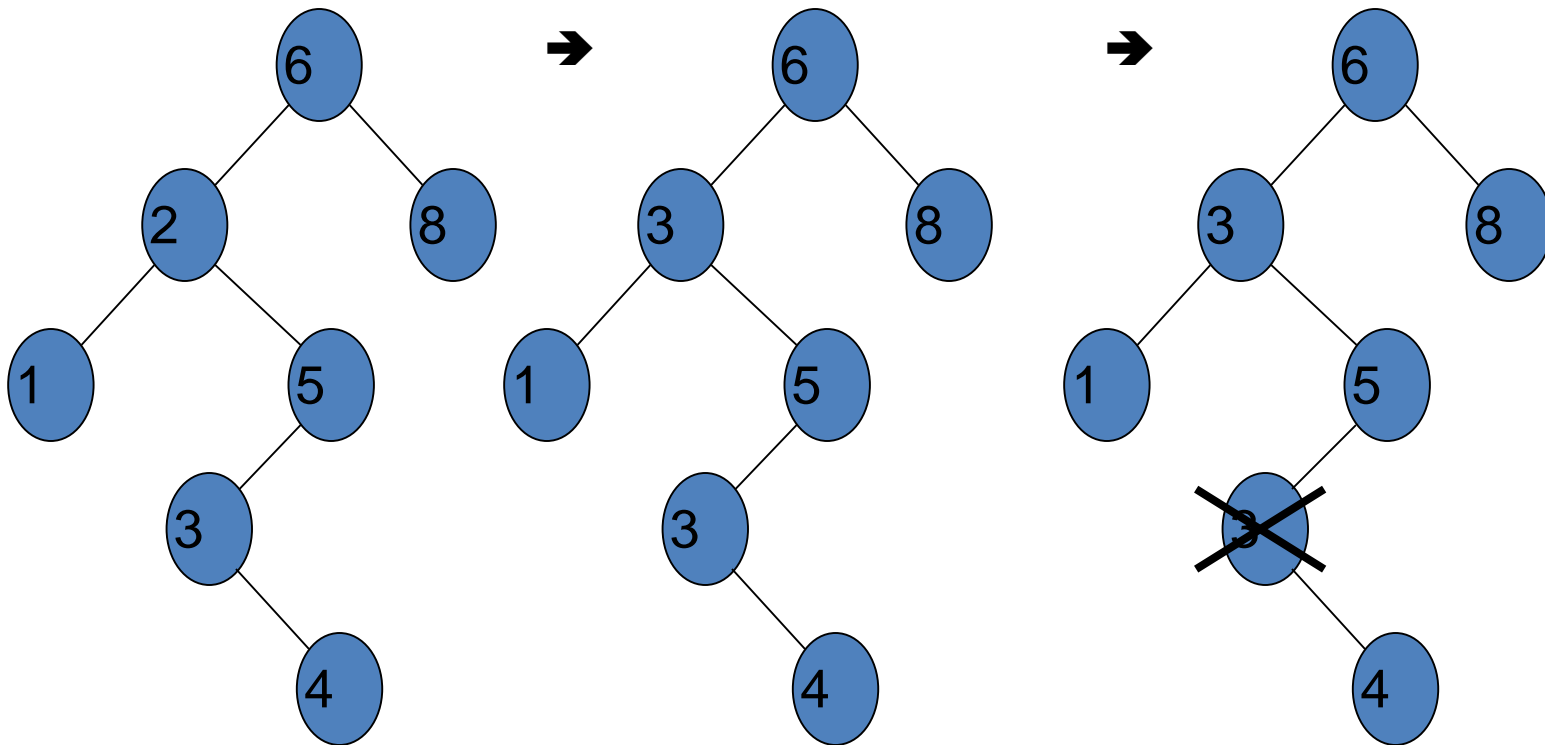
# Deleting a node in BST

Delete(2): copy data from inorder successor



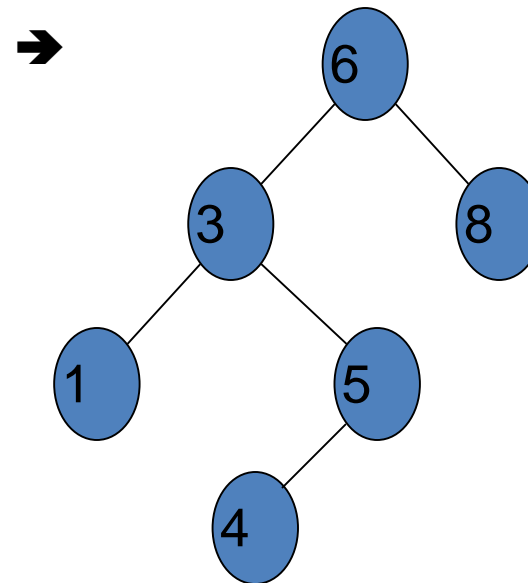
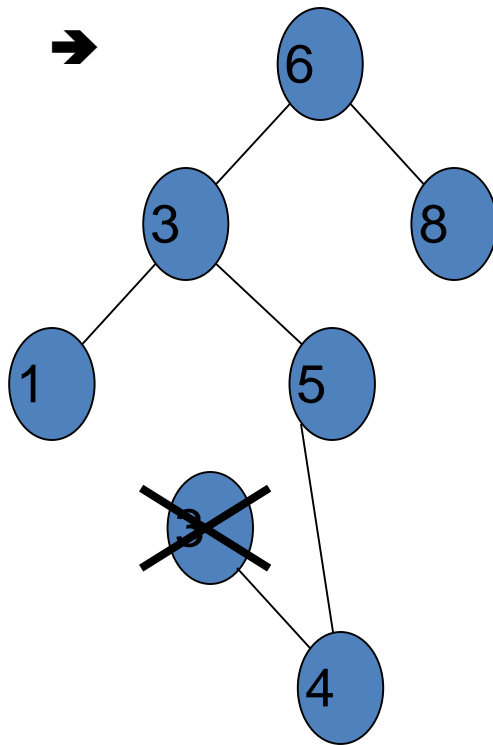
# Deleting a node in BST

Delete(2): remove the inorder successor



# Deleting a node in BST

Delete(2)



# delete

- When we delete a node, we need to consider how we take care of the children of the deleted node.
  - This has to be done such that the property of the **search tree** is maintained.

# delete

Three cases:

(1) the node is a leaf

- Delete it immediately

(2) the node has one child

- Adjust a pointer from the parent to bypass that node

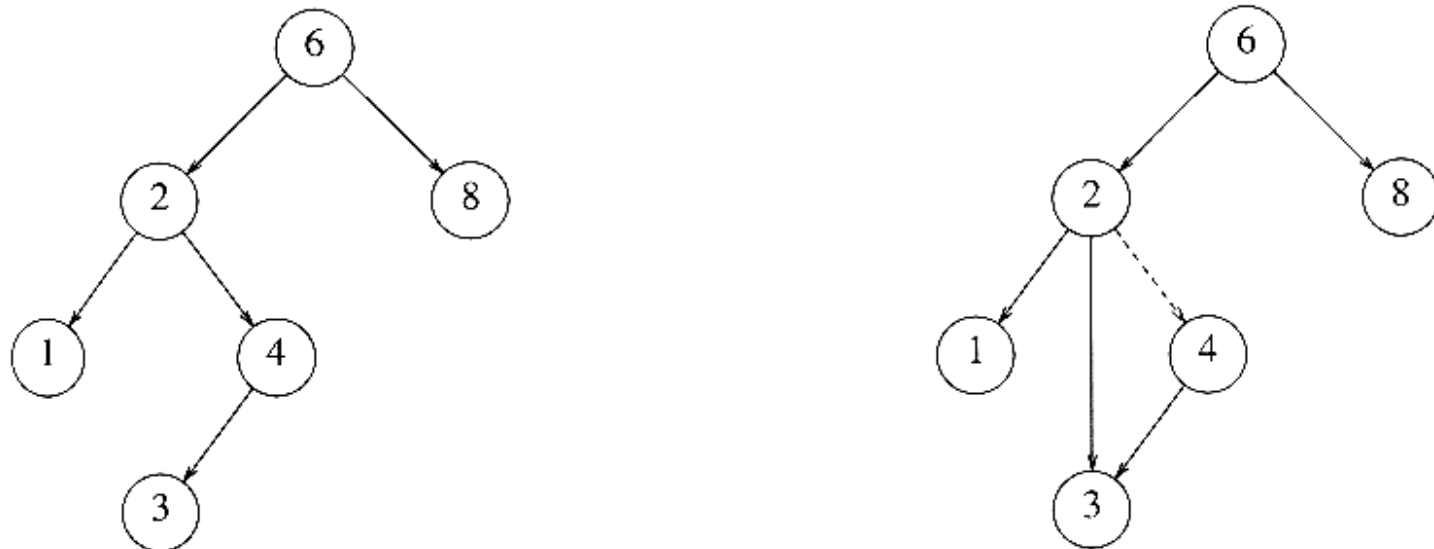
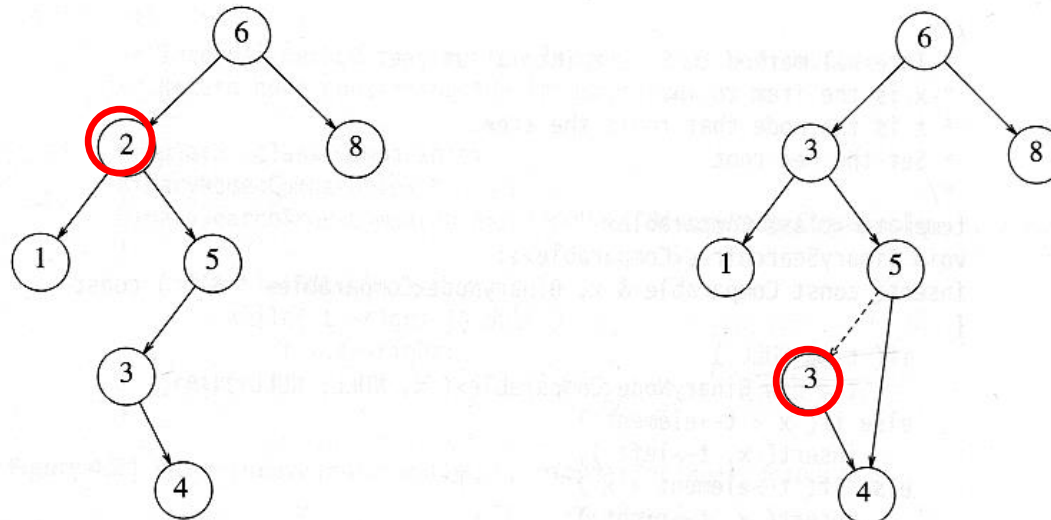


Figure 4.24 Deletion of a node (4) with one child, before and after

# delete

## (3) the node has 2 children

- replace the key of that node with the minimum element at the right subtree
- delete the minimum element
  - Has either no child or only right child because if it has a left child, that left child would be smaller and would have been chosen. So invoke case 1 or 2.

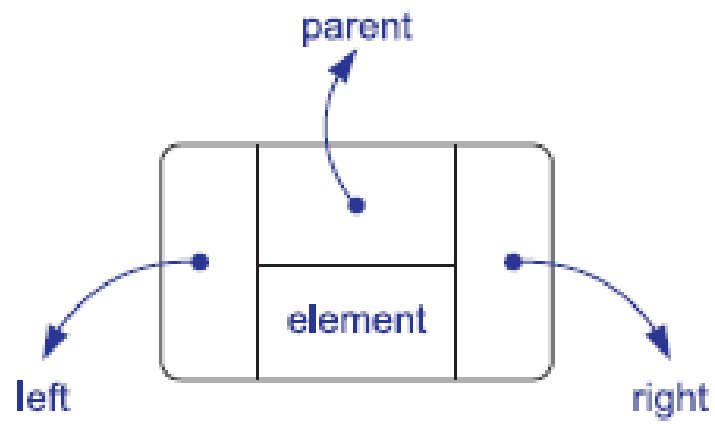


- Time com

Figure 4.25 Deletion of a node (2) with two children, before and after

**IMPLEMENTATION**





# Binary Search Tree Implementation

- We use two classes:
  - **BSTNode**
  - **BST**
- Declare **BSTNode** class for the nodes
  - info: **int**-type data in this example

```
class BSTNode {  
private:  
    int                info;           // data  
    BSTNode*          left;           // pointer to Left  
    BSTNode*          right;          // pointer to Right  
    BSTNode*          parent;         // pointer to Parent  
  
};
```

# BSTNode Implementation

```
class BSTNode{
public:
    //constructors
    BSTNode() ;
    BSTNode( int info);
    int getInfo();
    void setInfo(int info);
    BSTNode* getLeft();
    void setLeft(BSTNode *left);
    BSTNode *getRight();
    void setRight(BSTNode *right);
    int isLeaf( );
    BSTNode* getParent();
    void setParent(BSTNode *parent);

private:
    int    info;           // data
    BSTNode* left;         // pointer to Left
    BSTNode* right;        // pointer to Right
    BSTNode* parent;       // pointer to Parent

};
```

# BSTNode Constructors

```
BSTNode::BSTNode()  
{  
    this->info= -1;  
    this->left = this->right = NULL;  
    this->parent = NULL;  
  
}  
BSTNode:: BSTNode( int info)  
{  
    this->info= info;  
    this->left = this->right = NULL;  
    this->parent = NULL;  
  
}
```

# BSTNode Methods

```
int BSTNode::getInfo()  
{  
    return this->info;  
}  
void BSTNode::setInfo(int info)  
{  
    this->info = info;  
}  
BSTNode* BSTNode::getLeft()  
{  
    return this->left;  
}  
void BSTNode::setLeft (BSTNode *left)  
{  
    this->left = left;  
}
```

# BSTNode Methods

```
BSTNode* BSTNode::getRight()
{
    return this->right;
}

void BSTNode::setRight(BSTNode *right)
{
    this->right = right;
}

int BSTNode::isLeaf( )
{
    if( this->left == NULL && this->right == NULL )
        return 1;
    return 0;
}

BSTNode* BSTNode::getParent()
{
    return this->parent;
}

void BSTNode::setParent(BSTNode *parent)
{
    this->parent = parent;
}

}
```

# BST

```
class BST{
public:
    void insert(int info);
    bool search(int key);
    BSTNode* remove(BSTNode* tree,int key);
    BSTNode* findMin(BSTNode* Node);
    BSTNode* findMax(BSTNode* Node);
    int height();
    void preorder(BSTNode* root);
    void postorder(BSTNode* root);
    void inorder(BSTNode* root);

public:
    BSTNode*          root;

};
```

# Insert

```
void BST::insert(int value)
{
    BSTNode *p = root, *prev = 0;
    // find a place for inserting new node;
    while (p != 0)
    {
        prev = p;
        if (value < p->getInfo())
            p = p->left;
        else
            p = p->right;
    }
    if (root == 0) // tree is empty;
        root = new BSTNode(value);
    else if (value < prev->getInfo()){
        BSTNode temp= new BSTNode(value);
        temp->setParent (prev);
        prev->setLeft (temp);
    }else{
        BSTNode temp = new BSTNode(value);
        temp->setParent ( prev);
        prev->setright (temp);
    }
}
```



# C++ code for delete

```
BSTNode* BST::remove(BSTNode* tree, int info)
{
    BSTNode* t;
    int cmp = info - tree->getInfo();
    if( cmp < 0 ){
        t = remove(tree->getLeft(), info);
        tree->setLeft( t );
    }
    else if( cmp > 0 ){
        t = remove(tree->getRight(), info);
        tree->setRight( t );
    }
}
```

# C++ code for delete

```
BSTNode* BST::remove(BSTNode* tree, int info)
{
```

```
    BSTNode* t;
```

```
    int cmp = info - tree->getInfo();
```

```
    if( cmp < 0 ){
```

```
        t = remove(tree->getLeft(), info);
```

```
        tree->setLeft( t );
```

```
    }
```

```
    else if( cmp > 0 ){
```

```
        t = remove(tree->getRight(), info);
```

```
        tree->setRight( t );
```

```
    }
```

# C++ code for delete

```
BSTNode* BST::remove(BSTNode* tree, int info)
{
    BSTNode* t;
    int cmp = info - tree->getInfo();
    if( cmp < 0 ){
        t = remove(tree->getLeft(), info);
        tree->setLeft( t );
    }
    else if( cmp > 0 ){
        t = remove(tree->getRight(), info);
        tree->setRight( t );
    }
}
```

# C++ code for delete

```
BSTNode* BST::remove (BSTNode* tree, int info)
{
    BSTNode* t;

    int cmp = info - tree->getInfo();
    if( cmp < 0 ){
        t = remove(tree->getLeft(), info);
        tree->setLeft( t );
    }
    else if( cmp > 0 ){
        t = remove(tree->getRight(), info);
        tree->setRight( t );
    }
}
```

# C++ code for delete

```
➡ //two children, replace with inorder successor
else if(tree->getLeft() != NULL
        && tree->getRight() != NULL ){
    BSTNode* minNode;
    minNode = findMin(tree->getRight());
    tree->setInfo( minNode->getInfo() );
    t = remove(tree->getRight(),
               minNode->getInfo());
    tree->setRight( t );
}
```

# C++ code for delete

```
//two children, replace with inorder successor  
else if(tree->getLeft() != NULL
```

```
&& tree->getRight() != NULL ){
```

```
    BSTNode* minNode;
```

```
    minNode = findMin(tree->getRight());
```

```
    tree->setInfo( minNode->getInfo() );
```

```
    t = remove(tree->getRight(),
```

```
                minNode->getInfo());
```

```
    tree->setRight( t );
```

```
}
```

# C++ code for delete

```
//two children, replace with inorder successor
else if(tree->getLeft() != NULL
        && tree->getRight() != NULL ){
    BSTNode* minNode;
    minNode = findMin(tree->getRight());
    tree->setInfo( minNode->getInfo() );
    t = remove(tree->getRight(),
               minNode->getInfo());
    tree->setRight( t );
}
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# C++ code for delete

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               minNode->getInfo());
    tree->setRight( t );
}
```



# C++ code for delete

```
//two children, replace with inorder successor
else if(tree->getLeft() != NULL
        && tree->getRight() != NULL ){
    BSTNode* minNode;
    minNode = findMin(tree->getRight());
    tree->setInfo( minNode->getInfo() );
    t = remove(tree->getRight(),
                minNode->getInfo());
    tree->setRight( t );
}
```

# C++ code for delete

```
else { // case 1
    BSTNode* nodeToDelete = tree;
    if( tree->getLeft() == NULL ) //will handle 0 children
        tree = tree->getRight();
    else if( tree->getRight() == NULL )
        tree = tree->getLeft();
    else tree = NULL;

    delete nodeToDelete;
}
return tree;
}
```

# C++ code for delete

```
else { // case 1
    BSTNode<int>* nodeToDelete = tree;
    if( tree->getLeft() == NULL ) //will handle 0 children
        tree = tree->getRight();
    else if( tree->getRight() == NULL )
        tree = tree->getLeft();
    else tree = NULL;

    delete nodeToDelete;
}
return tree;
}
```

# C++ code for delete

```
else { // case 1
    BSTNode<int>* nodeToDelete = tree;
    if( tree->getLeft() == NULL ) //will handle 0 children
        tree = tree->getRight();
    else if( tree->getRight() == NULL )
        tree = tree->getLeft();
    else tree = NULL;
```



```
    delete nodeToDelete;
```

```
}
```

```
return tree;
```

```
}
```

# FindMin

```
➤ BSTNode* findMin(BSTNode* Node)
{
    if( Node== NULL )
        return NULL;
    if( Node->getLeft() == NULL )
        return Node; // this is it.
    return findMin( Node->getLeft() );
}
```

# FindMin

```
BSTNode* findMin(BSTNode* Node)
{
    if( Node== NULL )
        return NULL;
    if( Node->getLeft() == NULL )
        return Node; // this is it.
    return findMin( Node->getLeft() );
}
```

# FindMin

```
BSTNode* findMin(BSTNode* Node)
{
    if( Node== NULL )
        return NULL;
    if( Node->getLeft() == NULL )
        return Node; // this is it.
    ➡ return findMin( Node->getLeft() );
}
```

# FindMax

```
BSTNode* findMax(BSTNode* Node)
{
    if( Node== NULL )
        return NULL;
    if( Node->getRight() == NULL )
        return Node; // this is it.
    return findMax( Node->getRight() );
}
```



# Search

```
bool search(int value)
{
    BSTNode* p = root;
    while (p != 0)
    {
        if (value == p->getInfo())
            return true;
        else if (value < p->getInfo())
            p = p->getLeft();
        else
            p = p->getRight();
    }
    return 0;
}
```

Lecture content adapted from Michael T. Goodrich textbook,  
chapters 7 and 10.