Lecture 7 Artificial Intelligence

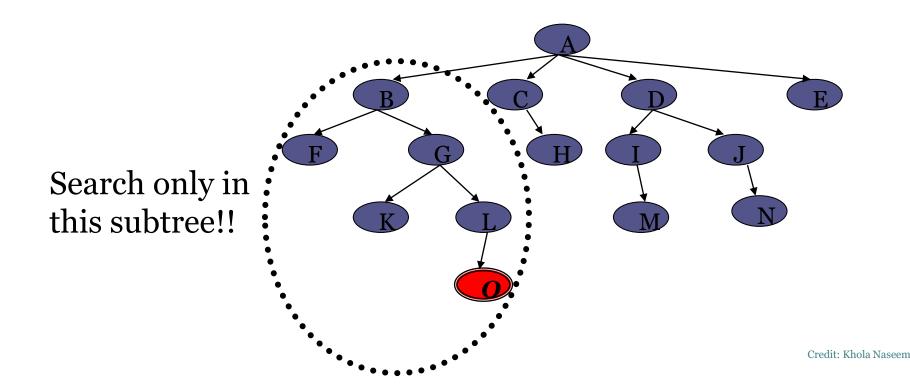
Khola Naseem khola.naseem@uet.edu.pk

Search strategies

- > Informed search, or heuristic search
 - ➤a cleverer strategy that searches toward the goal,
 - based on the information from the current state so far
 - ➤ E.g. A*, Heuristic DFS, Best first search

Informed search, or heuristic search

- ➤ With knowledge, one can search the state space as if he was given "hints" when exploring a maze.
- > Heuristic information in search = Hints
- Leads to dramatic speed up in efficiency.



Informed search, or heuristic search

- > Previous: at most b choices at each node and a depth of d at the goal node
 - An uninformed search algorithm would have to, in the worst case, search around O(b^d) nodes before finding a solution (Exponential Time Complexity).
- ➤ Heuristics improve the efficiency of search algorithms by reducing the effective branching factor from b to (ideally) a lower constant b* such that

$$>1 \le b^* << b$$

Evaluation Function

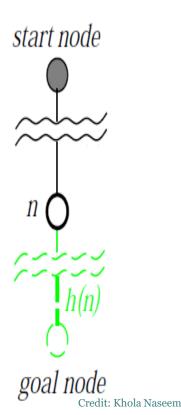
- An evaluation function f(n) gives an estimation on the "cost" of node n in a tree/graph
- So that the node with the least cost among all possible choices can be selected for expansion first

One Approach to calculate heuristic function > Evaluation function f measures the estimated cost of getting to the goal

from the current state:

$$\succ f(n) = h(n)$$

- > where
 - $\triangleright h(n)$ = an estimate of the cost to get from state n to a goal state

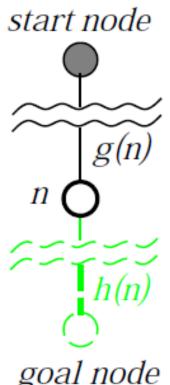


Another way to find the cost function

Evaluation function *f* measures the estimated cost of getting to the goal state from the current state and the cost of the existing path to it:

$$f(n) = g(n) + h(n)$$

where
 $g(n) = \text{an } exact \text{ cost to get to } n \text{ (from initial state)}$
 $h(n) = \text{an estimate of the cost to get}$
from state n to a goal state



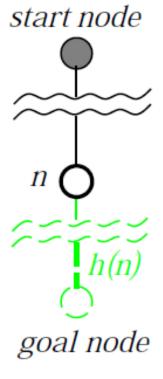
Example 8-Puzzle tree

- ▶h1: The number of misplaced tiles (squares with number).
- ➤ h2: The sum of the distances of the tiles from their goal positions.
 - ➤ h1: The number of misplaced tiles (not including the blank)

	100000000	Carle at	
C	1	2	3
Current State	4	5	6
	7		8
			MINE.
			_
a 1	1	2	3
Goal State	4	5	6



- Notation: f(n) = h(n)
- h(current state) = 1
- Because this state has one Y



Example 8-Puzzle tree

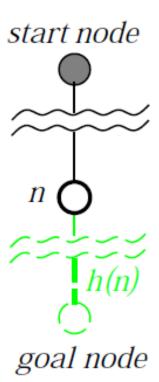
➤ h1: The number of misplaced tiles (not including the blank)

		femilie at	
Comment	1	2	3
Current State	4	5	6
	7		8
a 1	1	2	3
Goal State	4	5	6
	7	8	

N	N	N
N	N	N
N	Y	

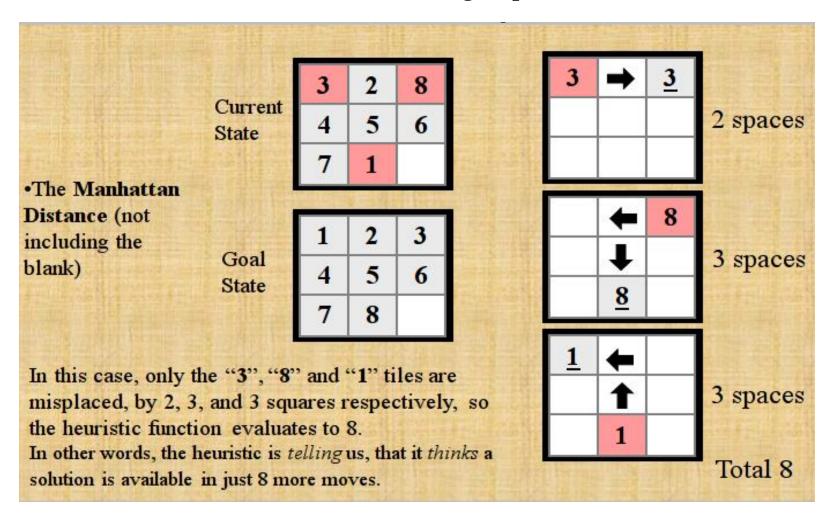
- Notation: f(n) = h(n)
- *h(current state) = 1*
- Because this state has one Y

- Only "**8**" is misplaced
- So the heuristic function evaluates to 1.
- The heuristic is *telling* us, that it *thinks* a solution might be available in just 1 more move.



Example 8-Puzzle tree

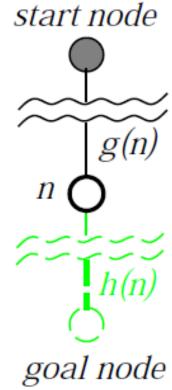
➤ h2: The sum of the distances of tiles from goal positions)



Notation: f(n) = h(n), h(current state) = 8

Another way to find the cost function

Evaluation function *f* measures the estimated cost of getting to the goal state from the current state and the cost of the existing path to it:



Summary of Evaluation / Heuristic Function

- \triangleright Evaluation function f(n) = g(n) + h(n)
 - \triangleright g(n) = exact cost so far to reach n
 - > h(n) = estimated cost to goal from n
 - \succ f(n) = estimated total cost of cheapest path through n to goal

Special cases:

- \triangleright Greedy (best-first) Search: f(n) = h(n)
- \triangleright A* Search: f(n) = g(n) + h(n)

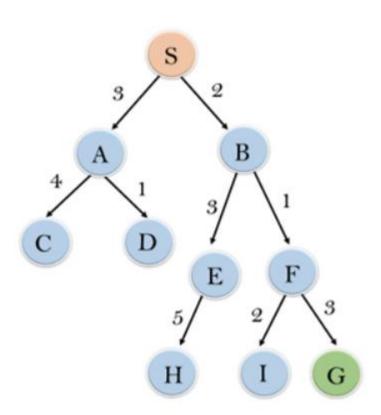
 $\widetilde{g(n)}$ $n \bigcirc$

start node

- **▶ Greedy Best First Search :** eval-fn: f(n) = h(n)
- ➤ Greedy best-first search tries to expand the node that is closest to the goal, on the grounds that this is likely to lead to a solution quickly.
- > Greedy best-first search algorithm always selects the path which appears best at that moment.
- ➤ It is the combination of depth-first search and breadth-first search algorithms. It uses the heuristic function and search.
- ➤ In the best first search algorithm, we expand the node which is closest to the goal node and the closest cost is estimated by heuristic function

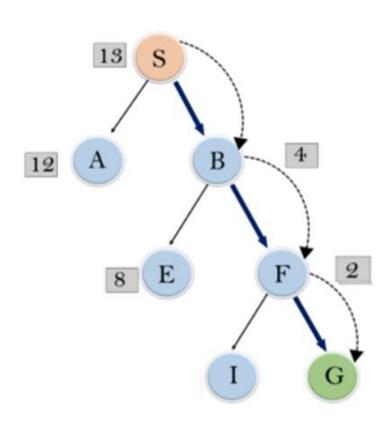
- **> Greedy Best First Search :** eval-fn: f(n) = h(n)
- ➤ Best first search can switch between BFS and DFS by gaining the advantages of both the algorithms.
- This algorithm is more efficient than BFS and DFS algorithms

- **Greedy Best First Search :** eval-fn: f(n) = h(n)
- > Example:



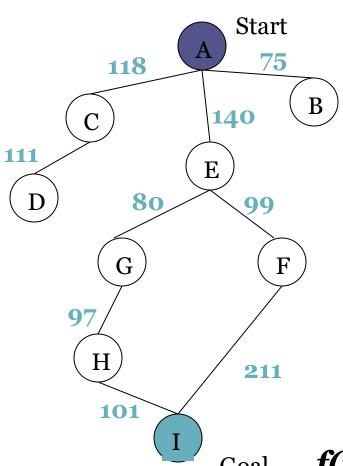
node	H (n)
A	12
В	4
C	7
D	3
E	8
F	2
Н	4
I	9
S	13
G	0

- **Greedy Best First Search :** eval-fn: f(n) = h(n)
- **Example:**

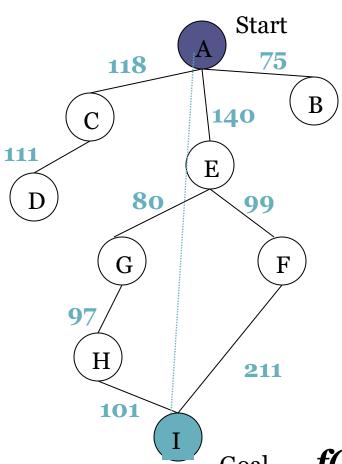


node	H (n)
A	12
В	4
C	7
D	3
E	8
F	2
Н	4
I	9
S	13
G	0

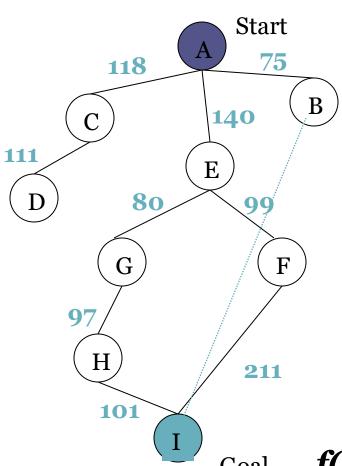
Greedy Search Example 2:



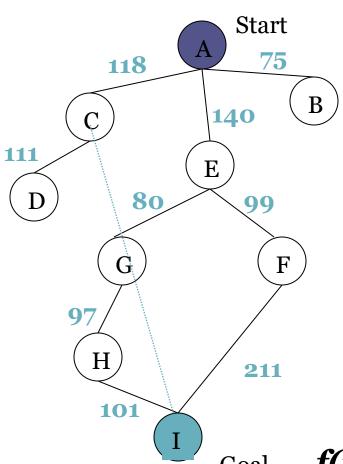
State	Heuristic: h(n)
Α	366
В	374
С	329
D	244
E	253
F	178
G	193
Н	98
I	0



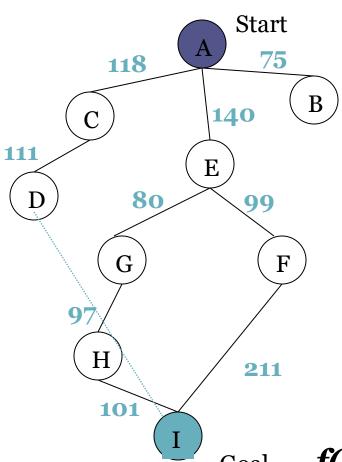
State	Heuristic: h(n)
A	366
В	374
С	329
D	244
E	253
F	178
G	193
Н	98
I	0



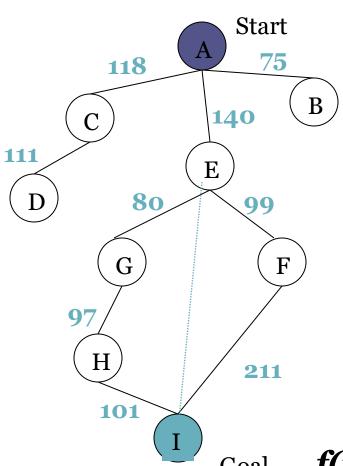
State	Heuristic: h(n)
А	366
В	374
С	329
D	244
Е	253
F	178
G	193
Н	98
I	0



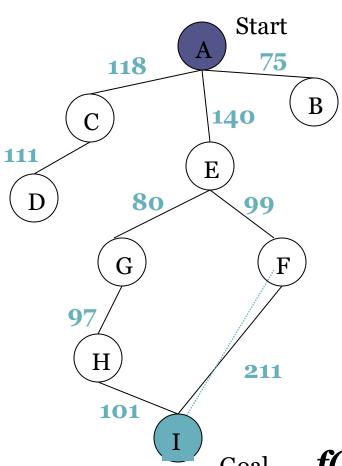
State	Heuristic: h(n)
Α	366
В	374
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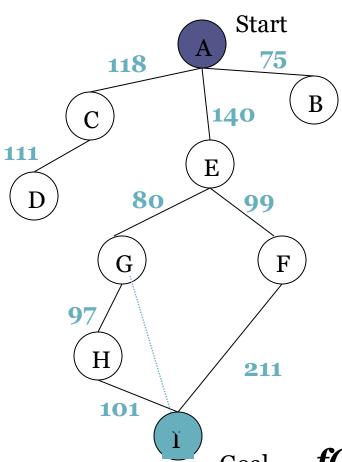
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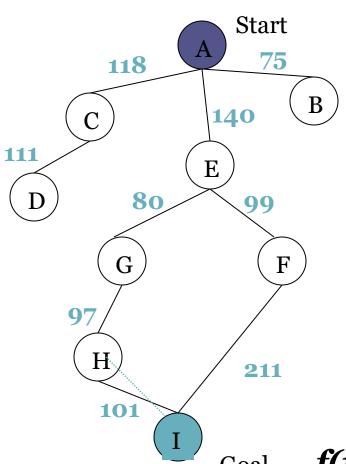
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E	253
F	178
G	193
Н	98
I	0



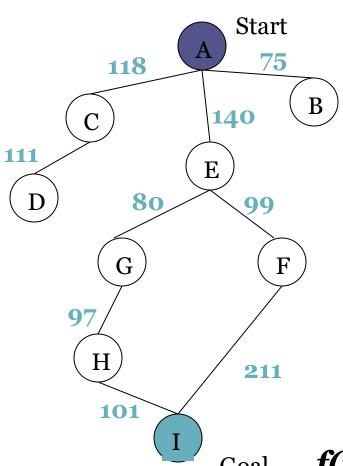
State	Heuristic: h(n)
Α	366
В	374
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State	Heuristic: h(n)
Α	366
В	374
С	329
D	244
E	253
F	178
G	193
Н	98
I	0

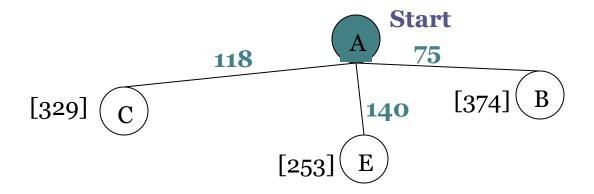


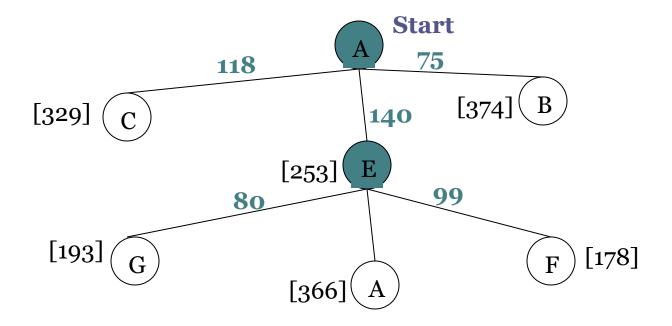
State	Heuristic: h(n)
Α	366
В	374
С	329
D	244
Е	253
F	178
G	193
Н	98
I	0

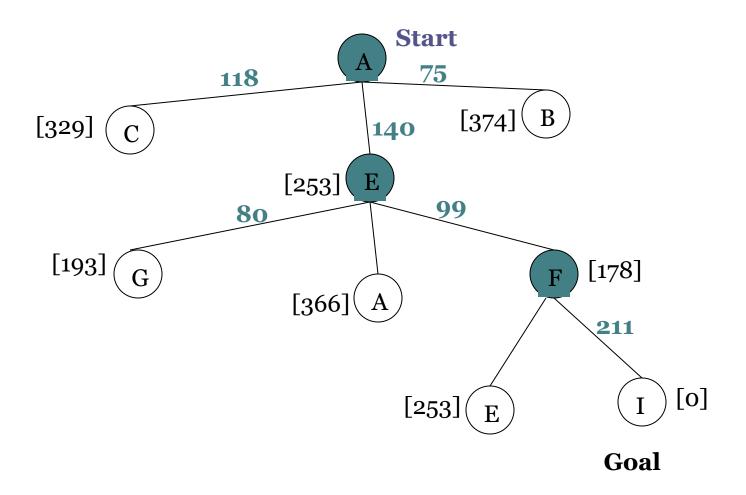


State	Heuristic: h(n)
Α	366
В	374
С	329
D	244
E	253
F	178
G	193
Н	98
I	0

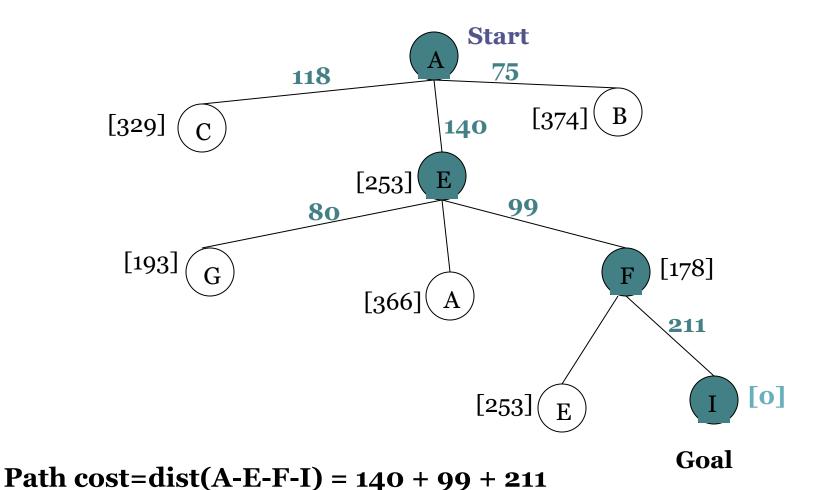




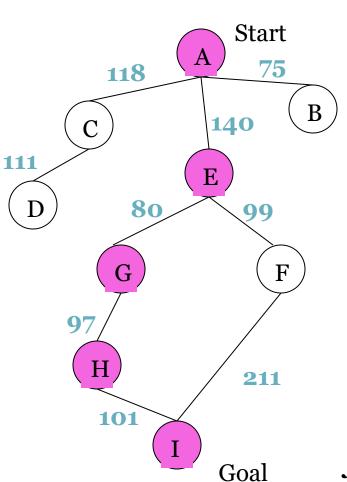




= 450



Greedy Search: Optimal?

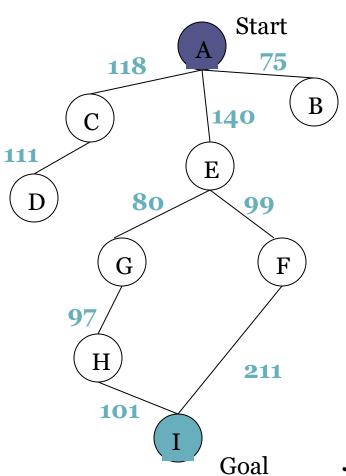


State	Heuristic: h(n)
Α	366
В	374
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D	244
E	253
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I	0

Not optimal!

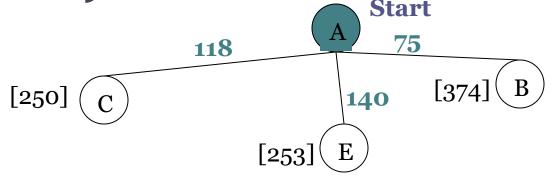
$$f(n) = h(n)$$
 = straight-line distance heuristic dist(A-E-G-H-I) = 140+80+97+101=418

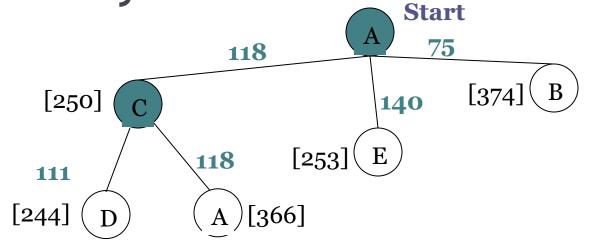
Greedy Search: Complete?



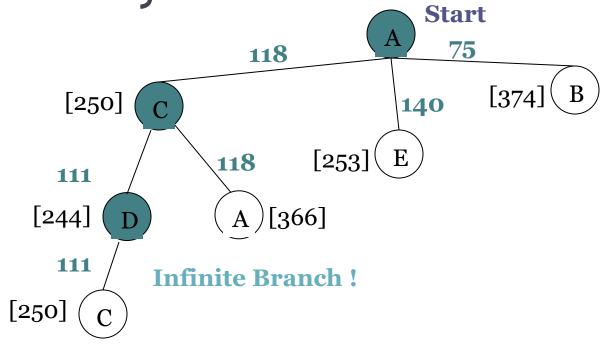
State	Heuristic: h(n)
Α	366
В	374
** C	250
D	244
E	253
F	178
G	193
Н	98
I	0



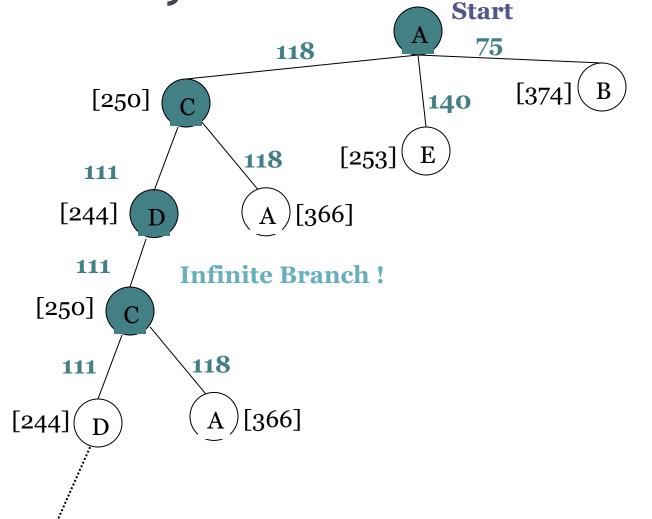




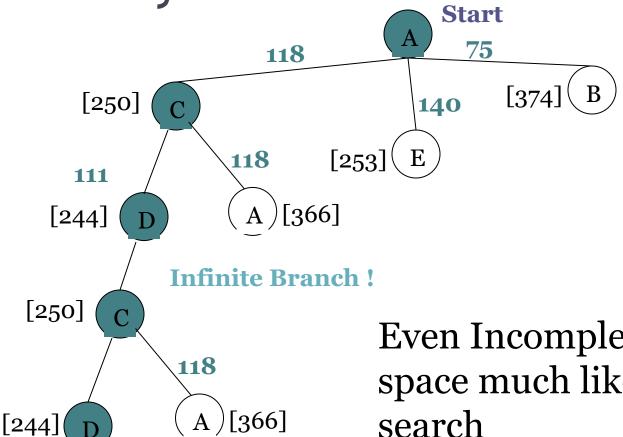
Greedy Search: Tree Search



Greedy Search: Tree Search



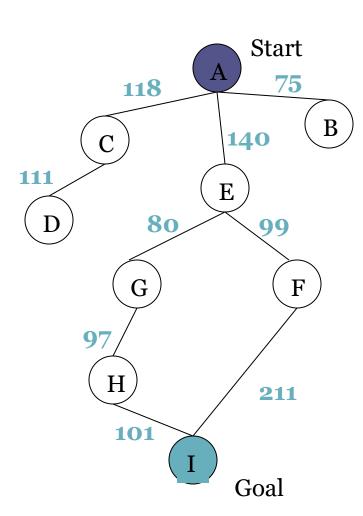
Greedy Search: Tree Search



Even Incomplete in finite state space much like depth first search

(Complete in finite space with repeated-state checking)

Greedy Search: Time and Space Complexity?

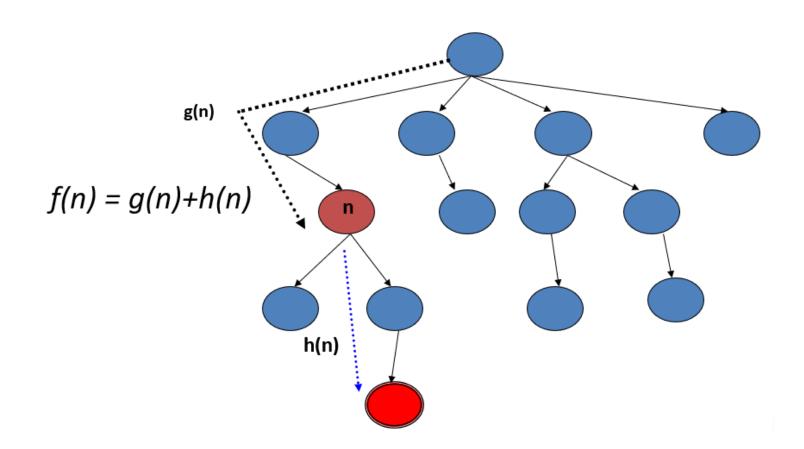


- Greedy search is <u>not optimal</u>.
- Greedy search is <u>incomplete</u>
- In the worst case, the Time and Space Complexity of Greedy Search are both O(b^m)

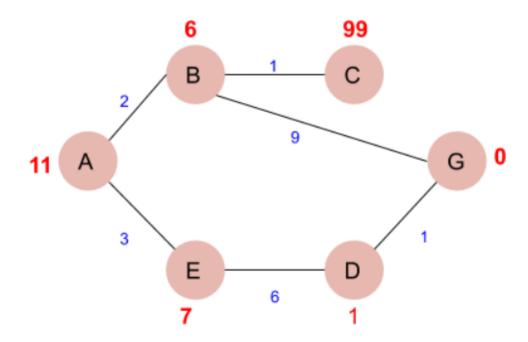
Where b is the branching factor and m the maximum path length

- A^* Search eval-fn: f(n)=g(n)+h(n)
- ➤ Although Greedy Search can considerably cut the search time (efficient), it is neither optimal nor complete.
- A* uses a priority function which combines g(n) and h(n): f(n) = g(n) + h(n)
- > g(n) is the exact cost to reach node n from the initial state. Cost so far up to node n.
- \triangleright h(n) is an estimation of the remaining cost to reach the goal.

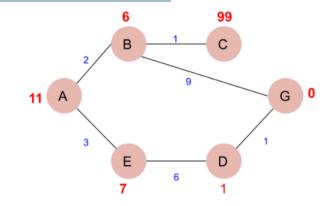
 A^* Search eval-fn: f(n)=g(n)+h(n)



 A^* Search eval-fn: f(n)=g(n)+h(n)

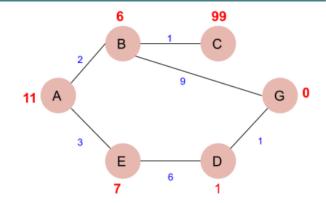


- \triangleright A* Search eval-fn: f(n)=g(n)+h(n)
- > Example:
- \triangleright g(x) + h(x) = f(x)
- \rightarrow 0+ 11 =11 Thus for A, we can write A=11



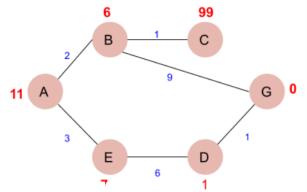
- Now from A, we can go to point B or point E, so we compute f(x) for each of them
- \rightarrow A \rightarrow B = 2 + 6 = 8
- \rightarrow A \rightarrow E = 3 + 7 = 10 (hold)
- \triangleright Since the cost for $A \rightarrow B$ is less, we move forward with this path and compute the f(x) for the children nodes of B
- ➤ Since there is no path between C and G, the heuristic cost is set to infinity or a very high value
- $A \rightarrow B \rightarrow C = (2 + 1) + 99 = 102$
- \rightarrow A \rightarrow B \rightarrow G = (2 + 9) + 0 = 11

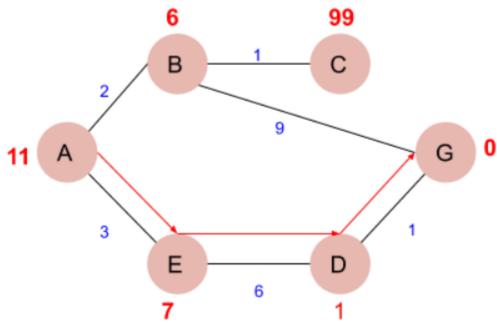
- \triangleright A* Search eval-fn: f(n)=g(n)+h(n)
- > Example:



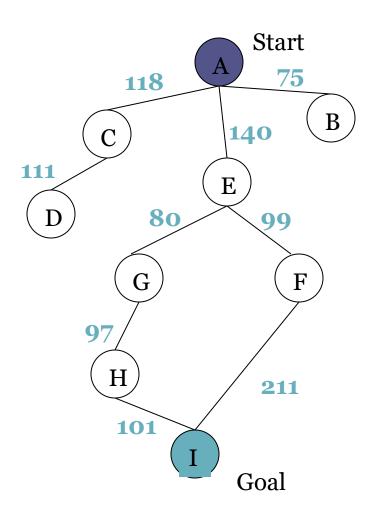
- ightharpoonup Here the path $A \to B \to G$ has the least cost but it is still more than the cost of $A \to E$, thus we explore this path further
- \rightarrow A \rightarrow E \rightarrow D = (3 + 6) + 1 = 10
- \triangleright Comparing the cost of A \rightarrow E \rightarrow D with all the paths we got so far and as this cost is least of all we move forward with this path.
- ➤ And compute the f(x) for the children of D
- \rightarrow A \rightarrow E \rightarrow D \rightarrow G = (3 + 6 + 1) +0 =10
- Now comparing all the paths that lead us to the goal, we conclude that A \rightarrow E \rightarrow D \rightarrow G is the most cost-effective path to get from A to G.

- A^* Search eval-fn: f(n)=g(n)+h(n)
- **Example:**





A* Search

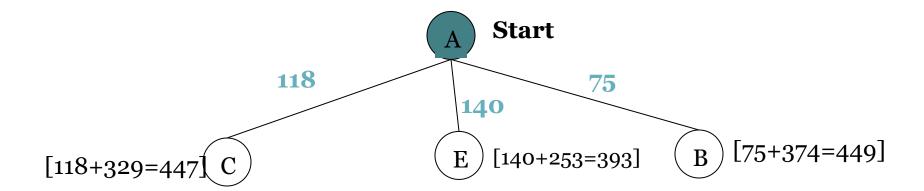


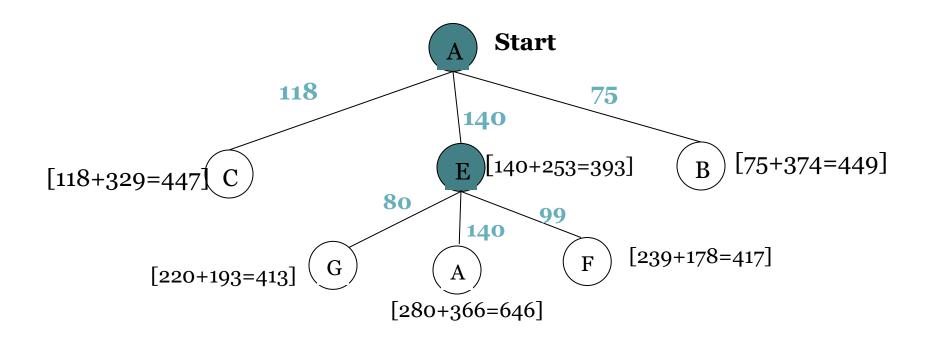
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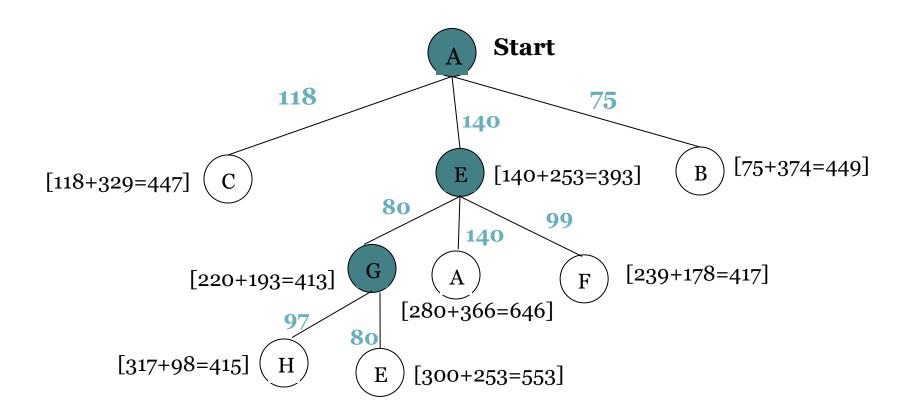
$$f(n) = g(n) + h(n)$$

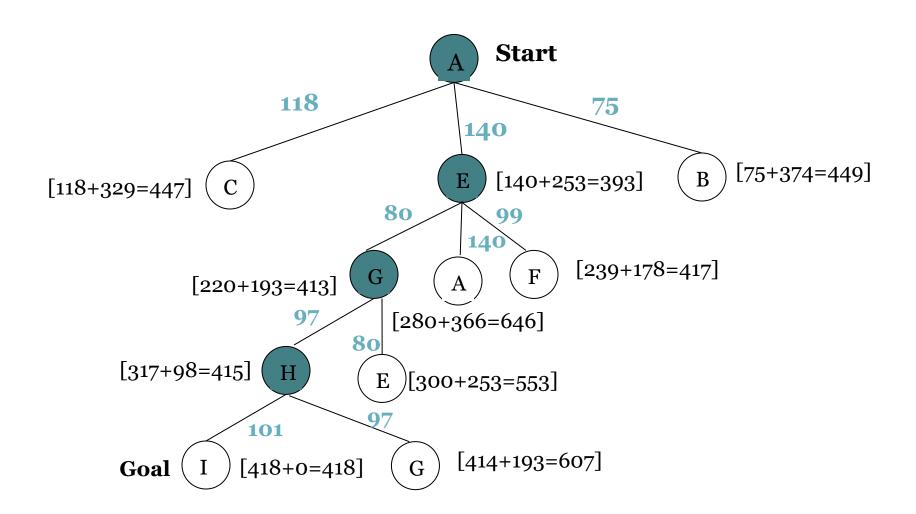
g(n): is the exact cost to reach node n from the initial state.

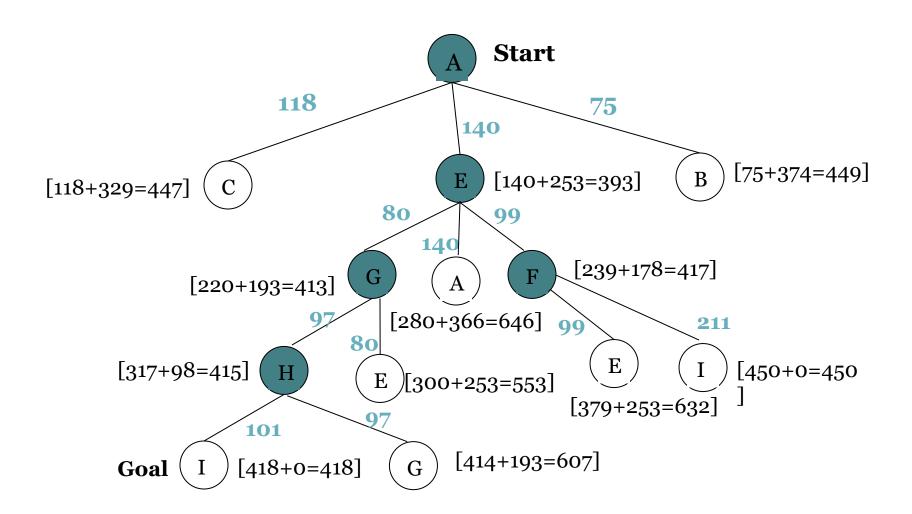


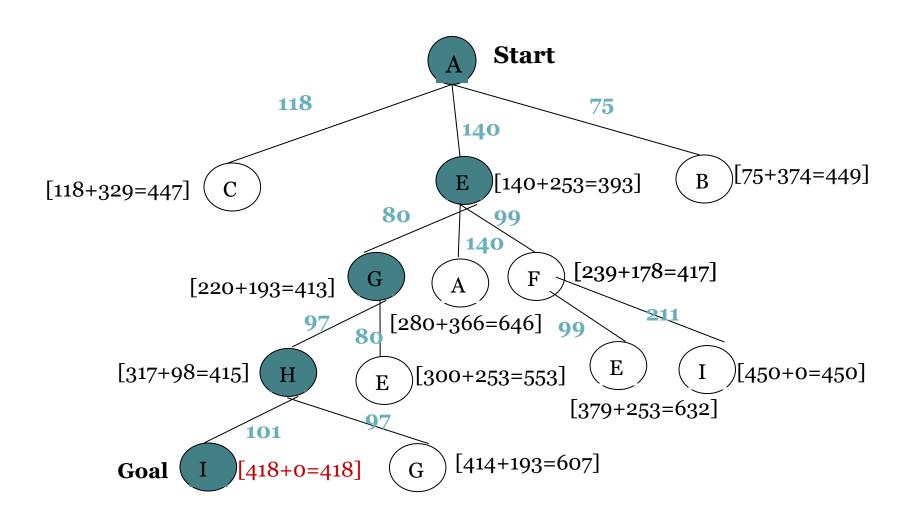


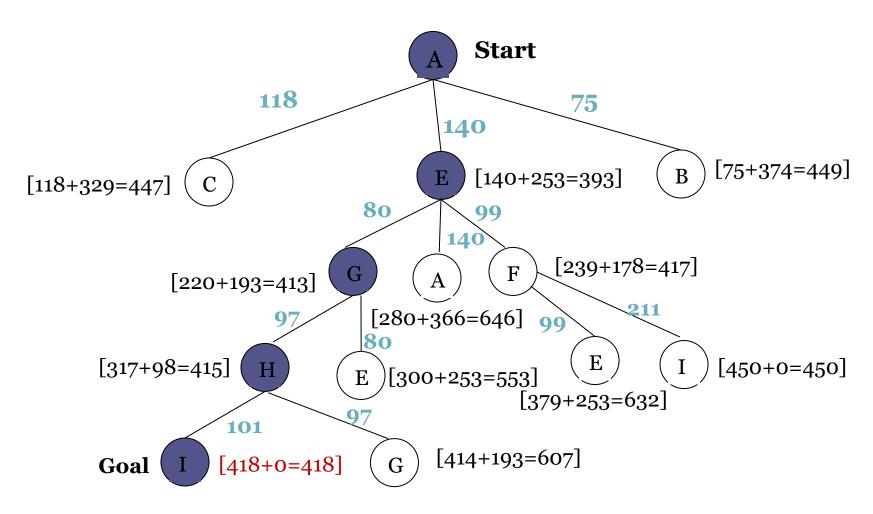






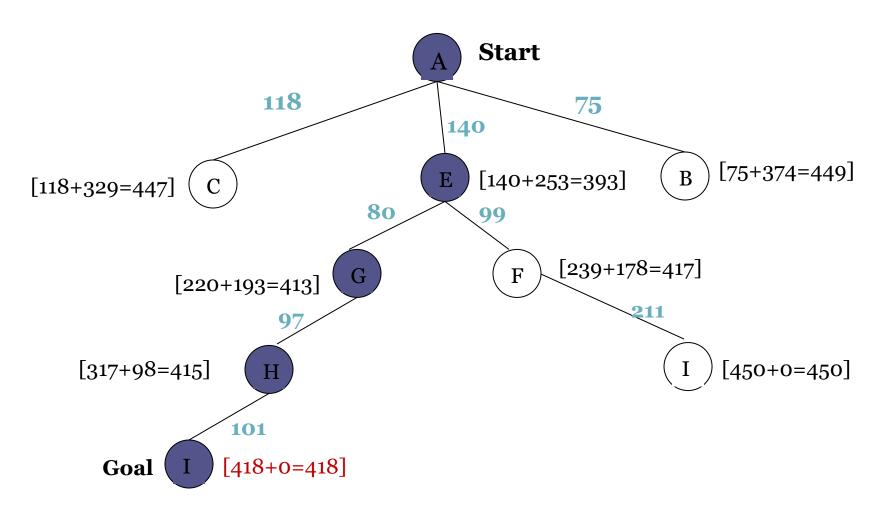






dist(A-E-G-H-I) = 140+80+97+101=418

A* Search: Graph Search



dist(A-E-G-H-I) = 140+80+97+101=418

A* Admissible

- An admissible heuristic is a heuristic that is guaranteed to find the shortest path from the current state to the goal state. In other words, it is an optimal heuristic.
- Admissible heuristics are often used in pathfinding algorithms such as A*.



Inadmissible (pessimistic) heuristics break optimality by trapping good plans on the fringe



Admissible (optimistic) heuristics slow down bad plans but never outweigh true costs

A* Admissible

- ➤ If h() overestimates the cost to reach the goal state
- ➤ Overestimate: not Admissible(If h() overestimates the cost to reach the goal state)
- > Underestimate: Admissible

```
h(n) \le h^*(n) :: Underestimation
h(n) \ge h^*(n) :: Overestimation
```

- \triangleright 1200 > 1000 i.e. h(n) ≥ h*(n) : Overestimation
- > 800 < 1000 i.e. $h(n) \le h^*(n)$: Underestimation

A* is Admissible if

- > we need to discover a solution to the problem, the estimated cost must be lower than or equal to the true cost of reaching the goal state.
- The heuristic function h(n) is called admissible if h(n) is never larger than h*(n), namely h(n) is always less or equal to true cheapest cost from n to the goal.

A* Admissible

> The evaluation function in A* looks like this:

```
f(n) = g(n) + h(n)

f(n) = Actual cost + Estimated cost

here,
```

n = current node.

f(n) = evaluation function.

g(n) = the cost from the initial node to the current node.

h(n) = estimated cost from the current node to the goal state.

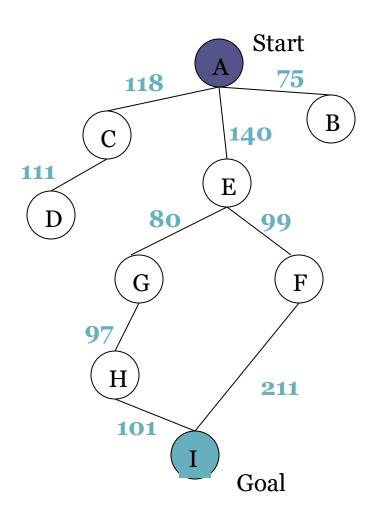
A* Admissible

➤ Now consider a fringe node n i.e. on an optimal solution path — example P in the example. If h(n) does not over estimate the cost of completing the solution path, then we know that

$$f(n) = g(n) + h(n) \le C^*$$

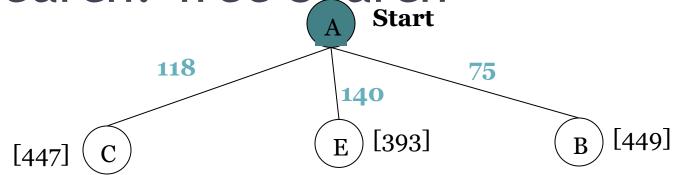
A* Search: if h not admissible

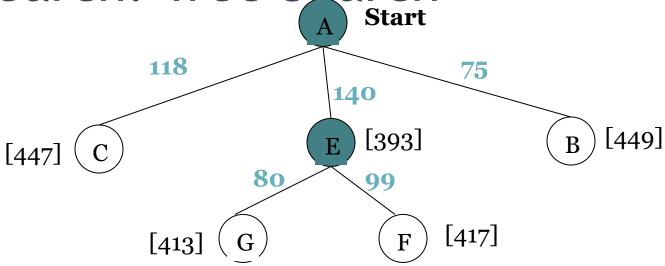
> Consistent Heuristic

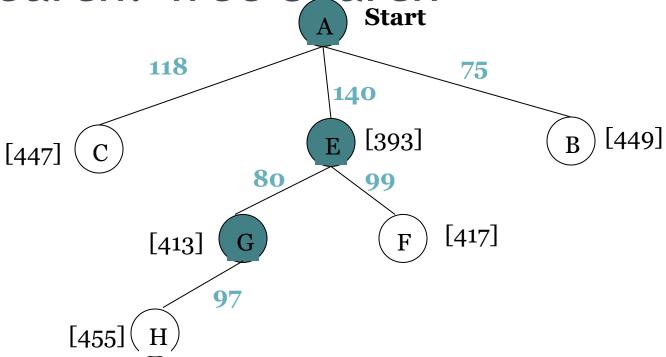


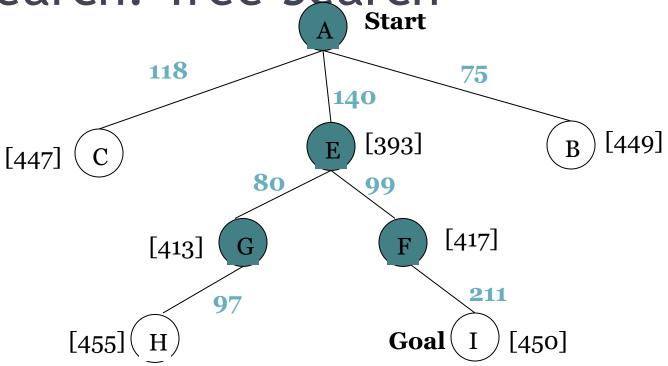
State	Heuristic: h(n)
Α	366
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I	0

A* Search: Tree Search A Start



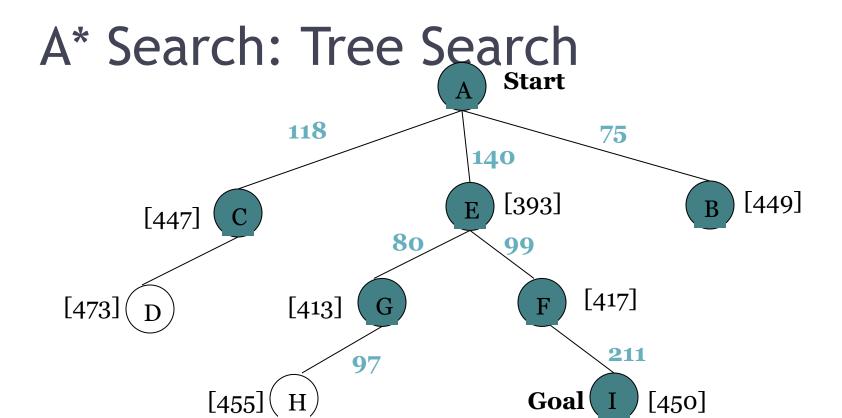






A* Search: Tree Search Start 118 **75 140** [449] [393] [447] 80 99 [417] [473] [413] **211** [455] Goal [450]

A* Search: Tree Search Start 118 **75 140** [449] [393] [447] 80 99 [417] [473] [413] **211** [455] Goal [450]



A* Search: Tree Search Start 118 **75 140** [449] [393] [447] 80 99 [417] [473] [413] 211 [455] [450] Goal

A* not optimal !!!

dist(A-E-F-I) =140+99+211=4

When heuristic h is not admissible