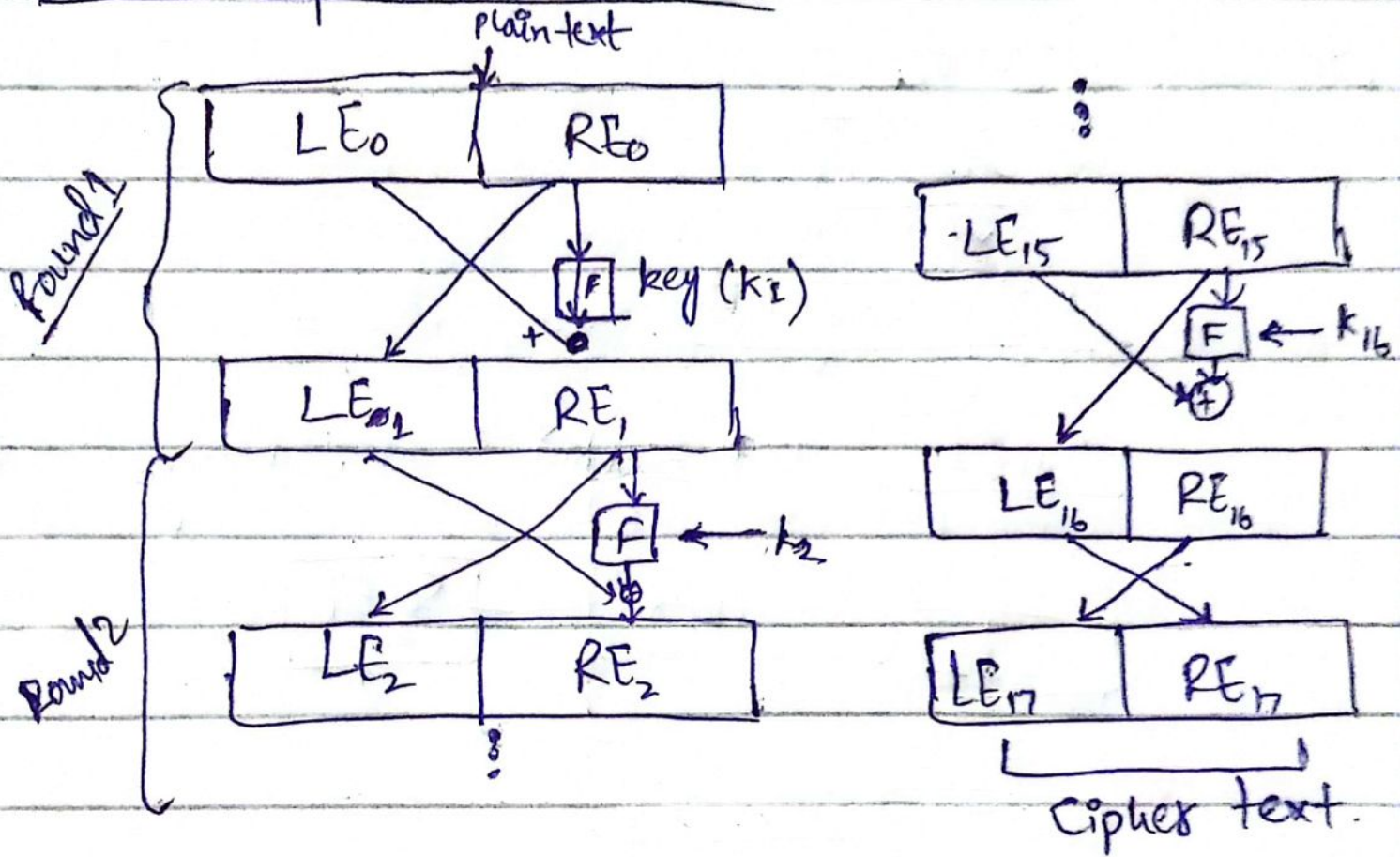


⇒ Feistel Cipher Structure :-



⇒ Block Cipher Design Principles

→ Block size 64 bits

→ Key size

→ No of Rounds

→ Subkey Count

→ Round Function

→ Plain divided into 2 halves

⇒ Data Encryption Standard (DES)

⇒ Block cipher ⇒ 64 bit Plaintext Block

⇒ Symmetric ⇒ 16 rounds (7 Feistel rounds)

⇒ 4 Steps

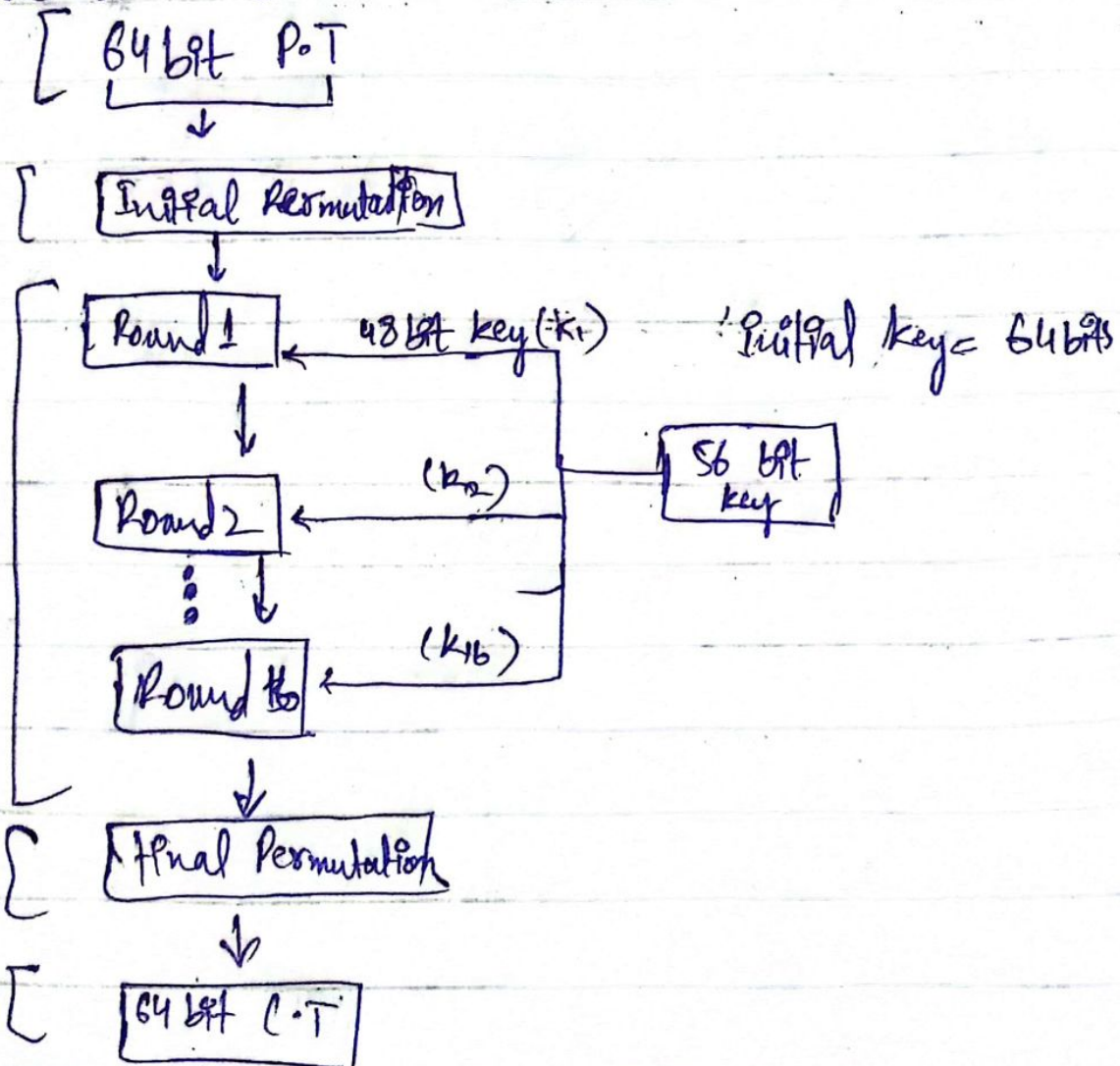
(i) Initial Permutation

(ii) 16 Feistel round

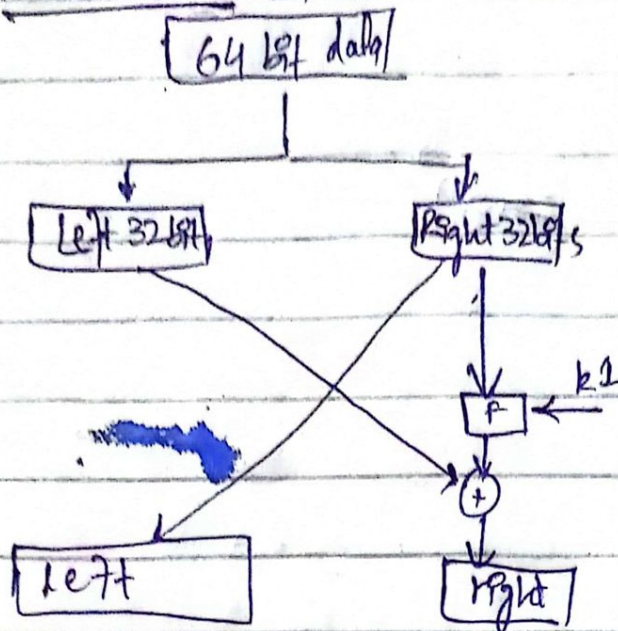
(iii) Swapping

(iv) Final Permutation

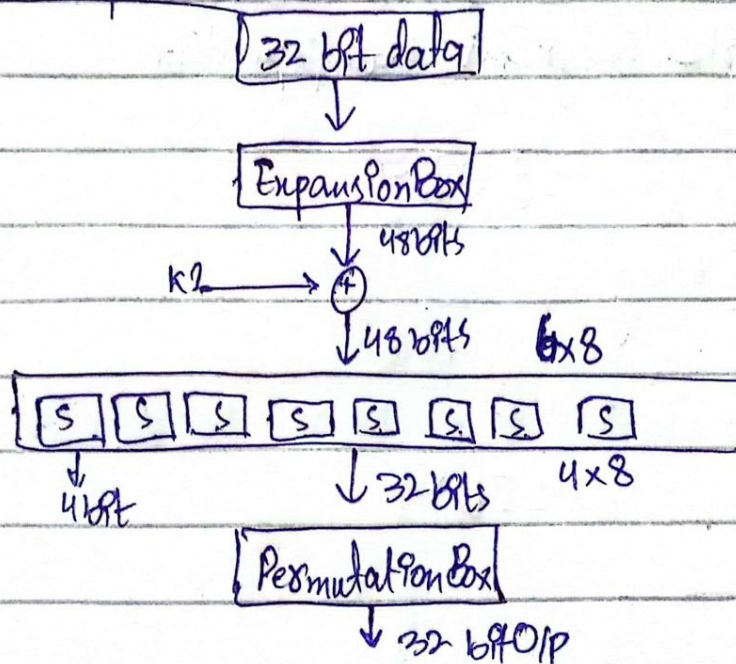
⇒ Basic Structure



⇒ Feistel Round



⇒ Round function



⇒ Expansion Box

data don't give the money to that person ever
 4 block size
 P DONT G I T GIVE T E THEM C

{IS}

⇒ Multiple DES

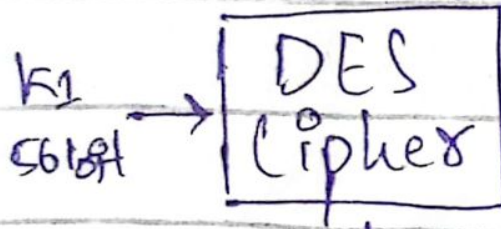
1) Double DES (2 DES)
↓ 64 Bit P.T

Encryption:

$$C.T = E(K_2, E(K_1, P.T))$$

Decryption:

$$P.T = D(K_1, D(K_2, C.T))$$



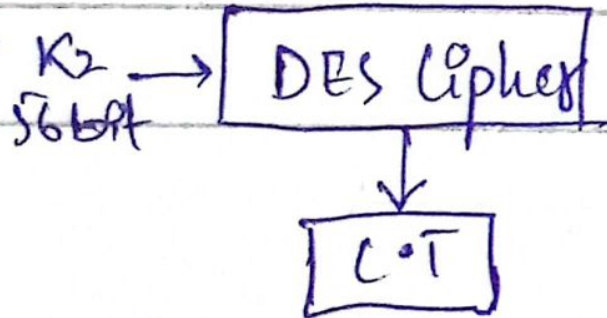
64 bit middle text

⇒ DES

key = 64 bit

$$64 - 8 = 56$$

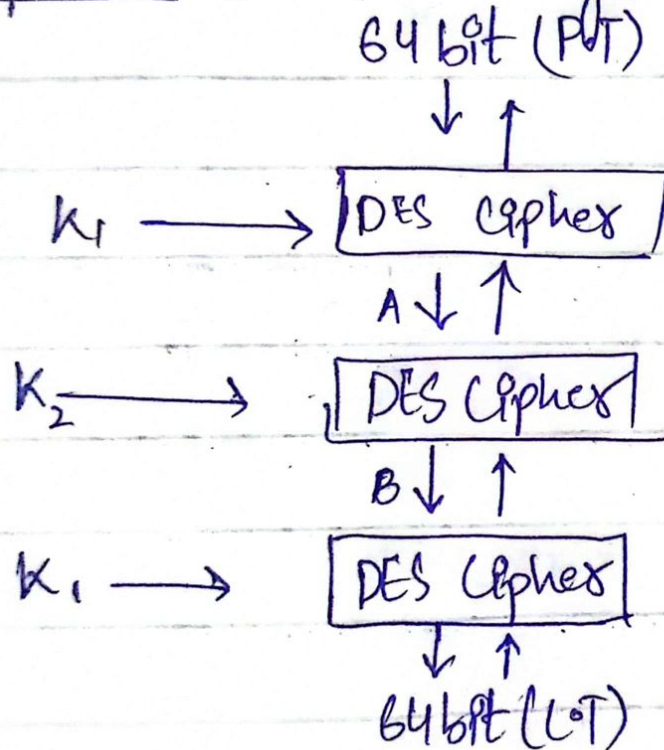
$$2^{56} \cdot 2 = 2^{57}$$



⇒ Meet-in-the-middle Attack :

If a person has both PT and CT he can generate possible outcomes of both keys using SB computation and the middle text.

⇒ Triple DES (3DES) with 2 keys



(3DES) using 3 keys

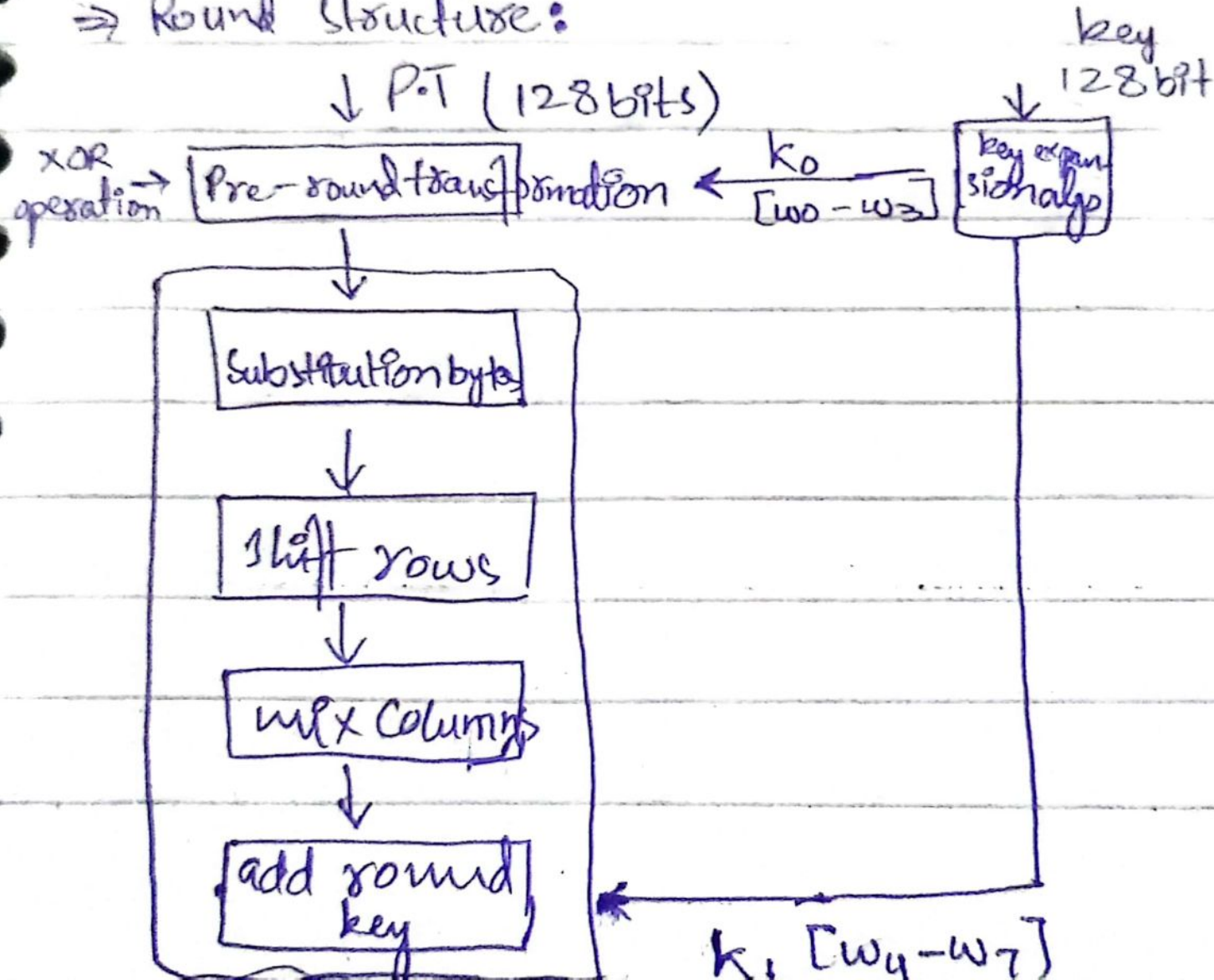
$$C = E(K_3, D(K_2, E(K_1, P.T)))$$

$$P.T = D(K_1, E(K_2, D(K_3, C)))$$

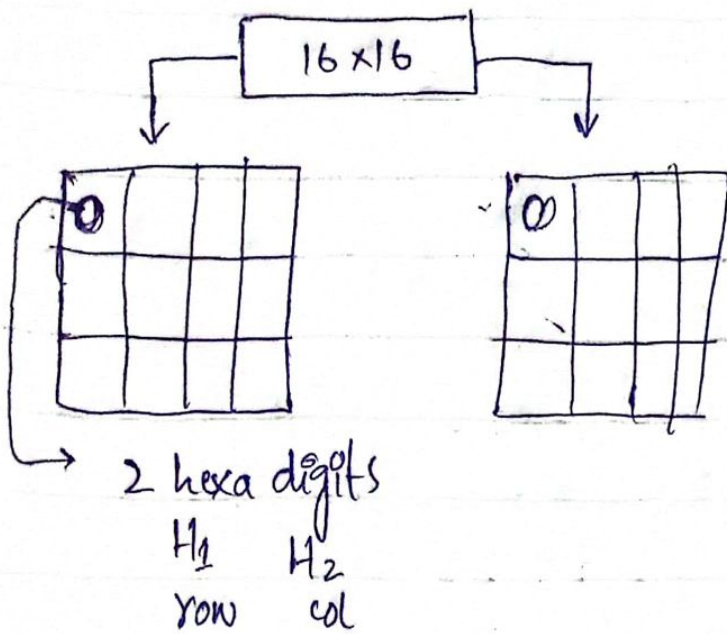
{IS}

Advanced Encryption Standard (AES)

⇒ Round Structure:



1) Substitution table (Bytes)



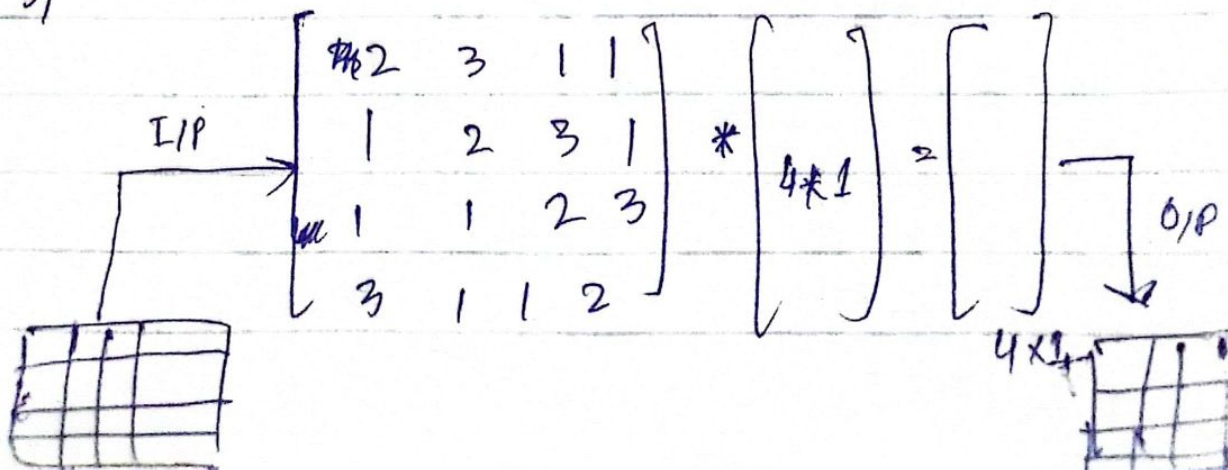
2)

R_0	63	C9	FE	30
R_1	F2	F2	63	26
R_2	C9	C3	7D	D4
R_3	BA	63	82	D4

↓ shift to left by Row #

63	C9	FE	30
F2	63	26	F2
7D	D4	C9	C3
D4	BA	63	82

3)



{IS}

Euler's Totient Function:

⇒ $\phi(n)$

⇒ no. of +ve integers less than 'n' that are co-prime to 'n' where $(n \geq 1)$

Ex $\phi(1) = 1$

$\phi(5) = ?$

No. < 5 are $\rightarrow 1, 2, 3, 4$

$\text{GCD}(1, 5) = 1$

$\text{GCD}(2, 5) = 1 \Rightarrow \phi(5) = 4$

$\text{GCD}(3, 5) = 1 \Rightarrow \phi(6) = ?$

$\text{GCD}(4, 5) = 1$

⇒ Theorem

if 'x' & 'n' are positive coprime integers

then $x^{\phi(n)} \equiv 1 \pmod{n}$

⇒ $x^{\phi(n)} \pmod{n} = 1 \pmod{n}$

Ex

$x = 11, n = 10$

$11^{\phi(10)} \equiv 1 \pmod{10}$

$\phi = \phi(10) = \phi(2) \times \phi(5)$

$= 1 \times 4 = 4$

$\Rightarrow 11^4 \equiv 1 \pmod{10}$

$1461 \equiv 1 \pmod{10}$

P

$x^{\phi(n)-a} \equiv 1 \pmod{n}$

Ex $11^{4+2} \equiv 1 \pmod{10}, 11^{40} \equiv 1 \pmod{10}$

Ex $4^{99} \mod 35$

$n=4, m=35 \rightarrow \text{co-prime}$

By Euler's theorem

$$4^{\phi(35)} = 1 \mod 35$$

$$\phi(35) = \phi(7) * \phi(5)$$

$$= 6 * 4 = 24$$

$$\Rightarrow 4^{24} = 1 \mod 35$$

$$\Rightarrow 4^{99} \rightarrow 4^{4*24 + 3}$$

$$= (4^{24})^4$$

$$+ 4^3 \mod 35$$

$$= 1 \cdot 4^3 \mod 35$$

$$= 64 \mod 35$$

$$= 29 \mod 35$$

~~Ex~~ $\mod 11$

Ex $4^{532} \mod 11$

$$x^{n-1} \equiv 1 \mod n$$

$$x^4, n=11$$

$$4^{11-1} \equiv 1 \mod 11$$

$$4^{10} \equiv 1 \mod 11$$

$$4^{532} = 4^{(10*53)+2}$$

$$= 4^{(10*53)} \cdot 4^2 \mod 11$$

$$= (1 \cdot \mod 11) \cdot (4^2 \mod 11)$$

$$= 16 \mod 11 = 5$$

Fermat's Theorem:

\Rightarrow specific case of Euler's theorem
where $m=n$

RSA
700m
7t